

# ECE 385 Lab Report 5

Experiment #5:

Simple Computer SLC-3.2 in SystemVerilog

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## Introduction

For Lab 5.1 and 5.2, we had to implement an SLC-3 Processor using SystemVerilog. This is a microprocessor which is a subset of LC-3 ISA, having 16-bit Program Counter (PC), 16-bit Instructions, and Registers. This processor was able to perform several operations using a certain number of opcodes, which were a subset of the LC-3 ISA. It followed the same Datapath and the State Machine diagram as an LC-3 Processor except without the logic and states for the TRAP instructions. For Lab 5.1, we had to implement the FETCH phase. For lab 5.2, we had to implement the DECODE and EXECUTE phases. After creating all the 3 phases, our SLC-3 microprocessor was complete.

## Written Description and Diagrams of SLC-3

Our SLC-3 microprocessor had 3 main phases that each instruction would have to go through. As mentioned earlier, these were FETCH, DECODE and EXECUTE. The FETCH and DECODE phase was common for all instructions. In the FETCH phase, the instruction is identified and taken from memory. The way the processor does this is by performing 4 steps:

1.  $MAR \leftarrow PC \Rightarrow$  The MAR stores the address that contains the instruction to be performed
2.  $MDR \leftarrow M(MAR) \Rightarrow$  MDR will now store the instruction to be performed
3.  $IR \leftarrow MDR \Rightarrow$  IR stores the instruction that needs to be decoded before performing
4.  $PC \leftarrow PC + 1 \Rightarrow$  PC is incremented to perform the next instruction

The next phase that each instruction must go through is the DECODE phase. In this phase, the instruction that is stored in IR (during the FETCH cycle) is put through the instruction decoder. In this phase, the control signals required for the instruction in the IR are set, for example, the select signals for the muxes or the load signals for the registers.

The final phase that every instruction must go through is the EXECUTE phase. In this phase, the operation based on the opcode in the instruction is performed depending on the control signals set in the DECODE phase. The result is then stored in memory or register, depending on the instruction. Our microprocessor could perform 8 instructions, shown below.

Instruction	Instruction(15 downto 0)							Operation
ADD	0001	DR	SR1	0	00	SR2		$R(DR) \leftarrow R(SR1) + R(SR2)$
ADDi	0001	DR	SR	1	imm5			$R(DR) \leftarrow R(SR) + \text{SEXT}(\text{imm5})$
AND	0101	DR	SR1	0	00	SR2		$R(DR) \leftarrow R(SR1) \text{ AND } R(SR2)$
ANDi	0101	DR	SR	1	imm5			$R(DR) \leftarrow R(SR) \text{ AND } \text{SEXT}(\text{imm5})$
NOT	1001	DR	SR		111111			$R(DR) \leftarrow \text{NOT } R(SR)$
BR	0000	n	z	p	PCOffset9			if ((nzp AND NZP) != 0) $PC \leftarrow PC + \text{SEXT}(\text{PCOffset9})$
JMP	1100		000	BaseR		000000		$PC \leftarrow R(\text{BaseR})$
JSR	0100	1	PCOffset11					$R(7) \leftarrow PC;$ $PC \leftarrow PC + \text{SEXT}(\text{PCOffset11})$
LDR	0110	DR	BaseR		offset6			$R(DR) \leftarrow M[R(\text{BaseR}) + \text{SEXT}(\text{offset6})]$
STR	0111	SR	BaseR		offset6			$M[R(\text{BaseR}) + \text{SEXT}(\text{offset6})] \leftarrow R(SR)$
PAUSE	1101	ledVect12						$LEDs \leftarrow \text{ledVect12}; \text{ Wait on Continue}$

Table 1: SLC-3 ISA

Once the EXECUTE phase is completed for one instruction, the processor loops back to the FETCH cycle for the next instruction stored in the updated Program Counter.

### Block Diagram of slc3.sv

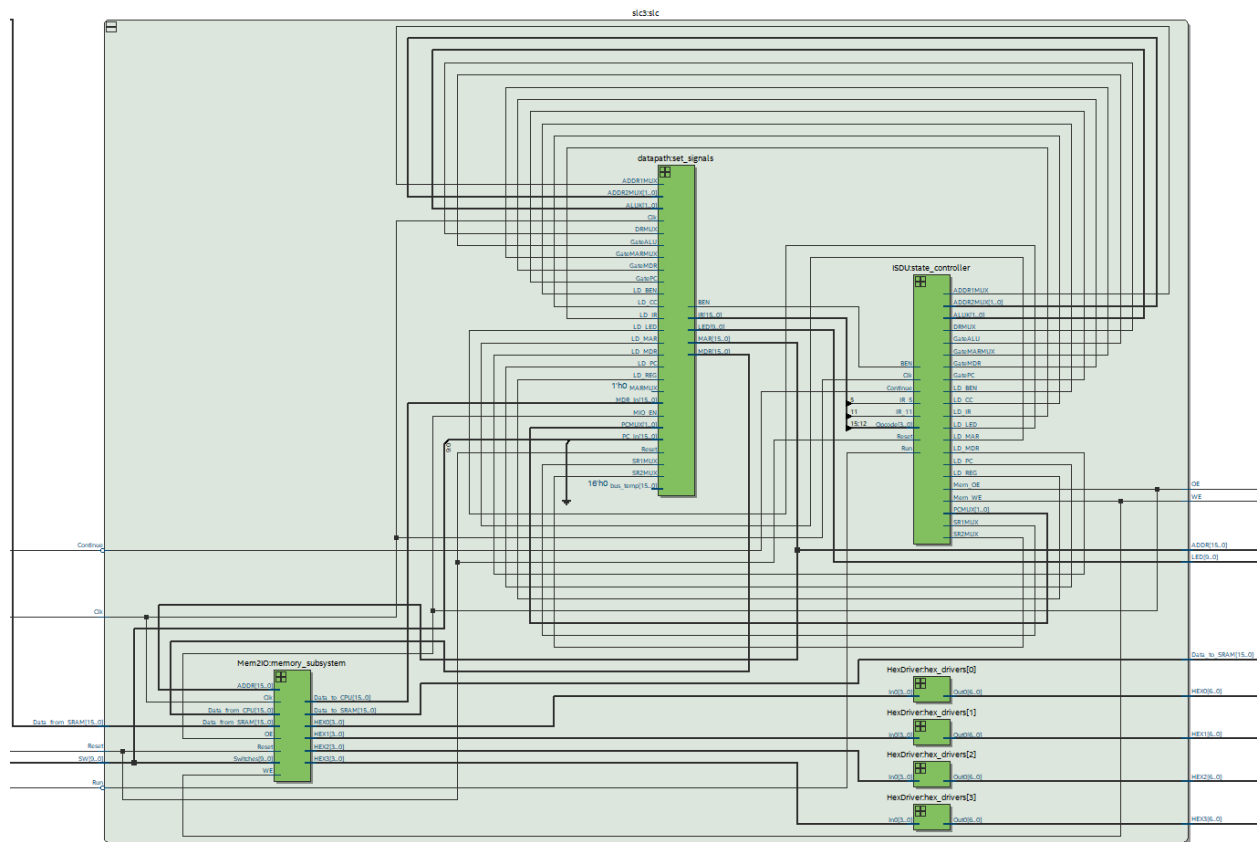


Figure 1: Block diagram of slc3.sv module

## State Diagram of ISDU

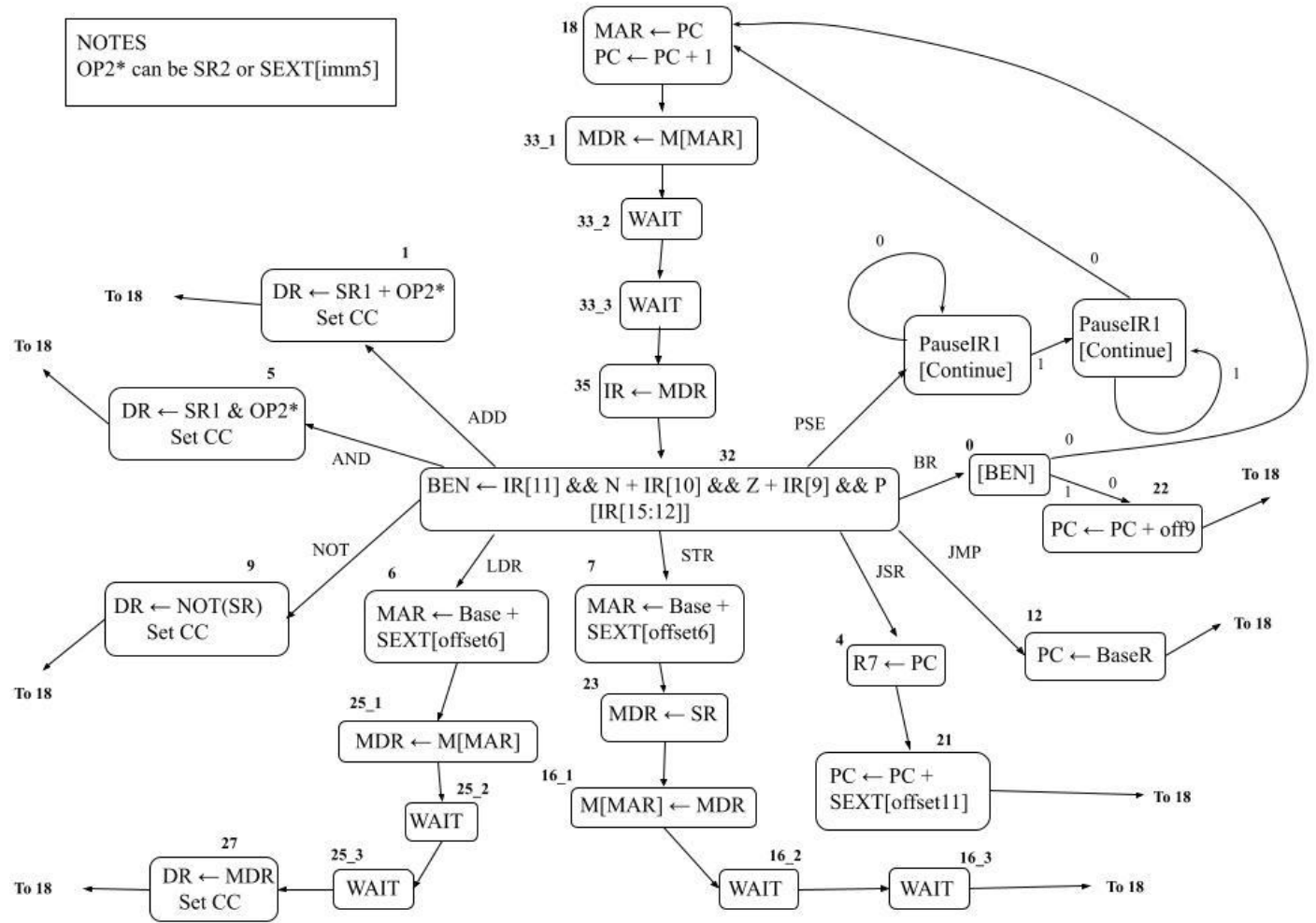


Figure 2: State diagram of ISDU

## Simulations of SLC-3 Instructions

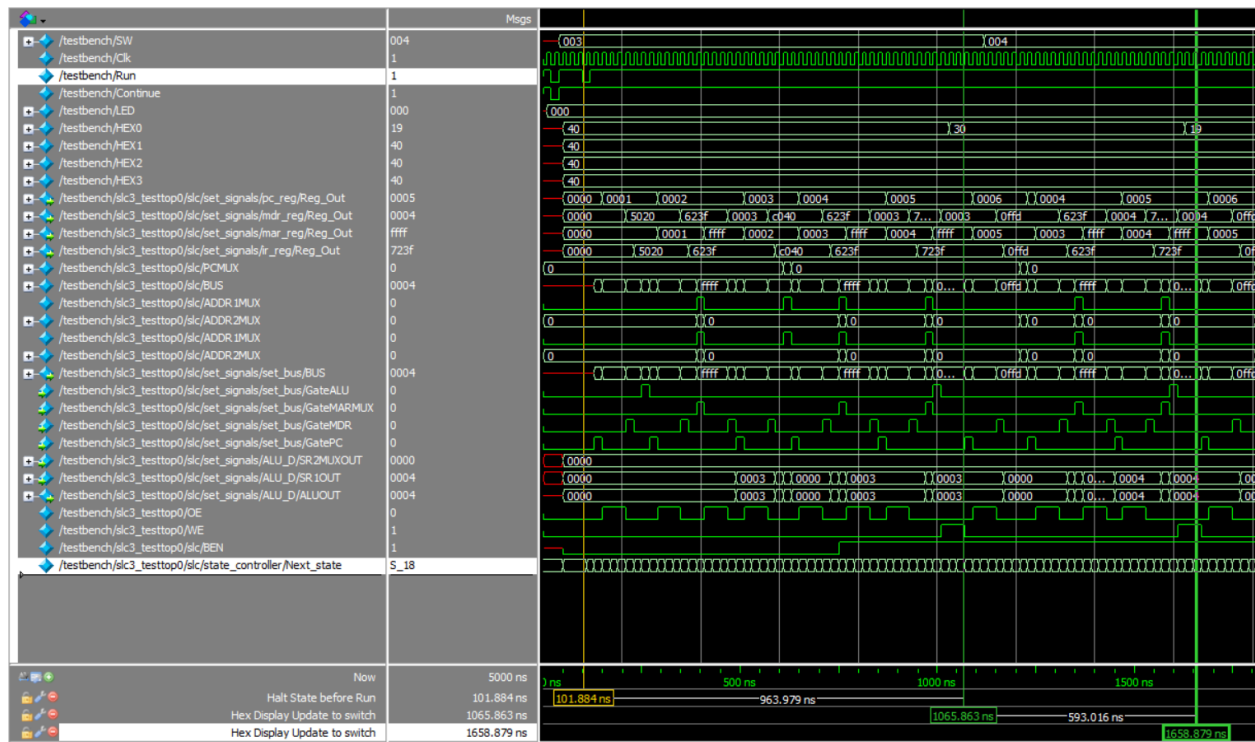


Figure 3: IO Test 1

- Test 1 Simulation
  - FPGA is reset
  - OPCode is set to switches and run is pressed
  - As switches are changed the value on the Hex Display is updated to be the value of the switches

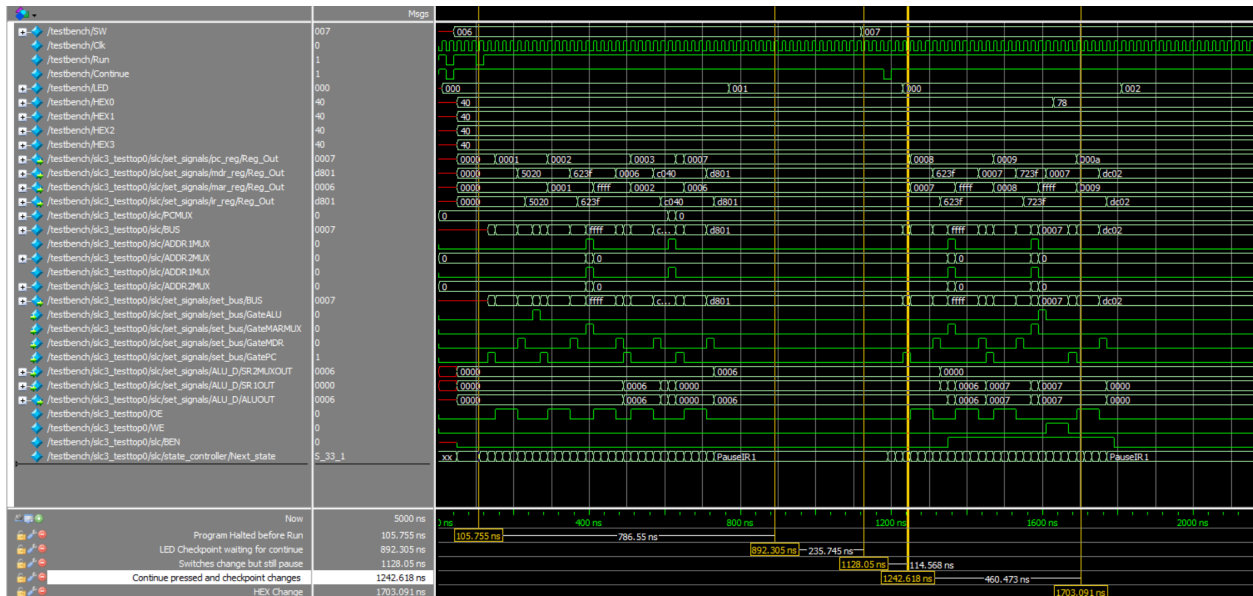


Figure 4: IO Test 2

- Test 2 Simulation
  - FPGA is reset
  - OPCode is set to switches and run is pressed
  - As switches are changed the value on the Hex Display does not update
  - Only when continue is also pressed does the Hex Display update

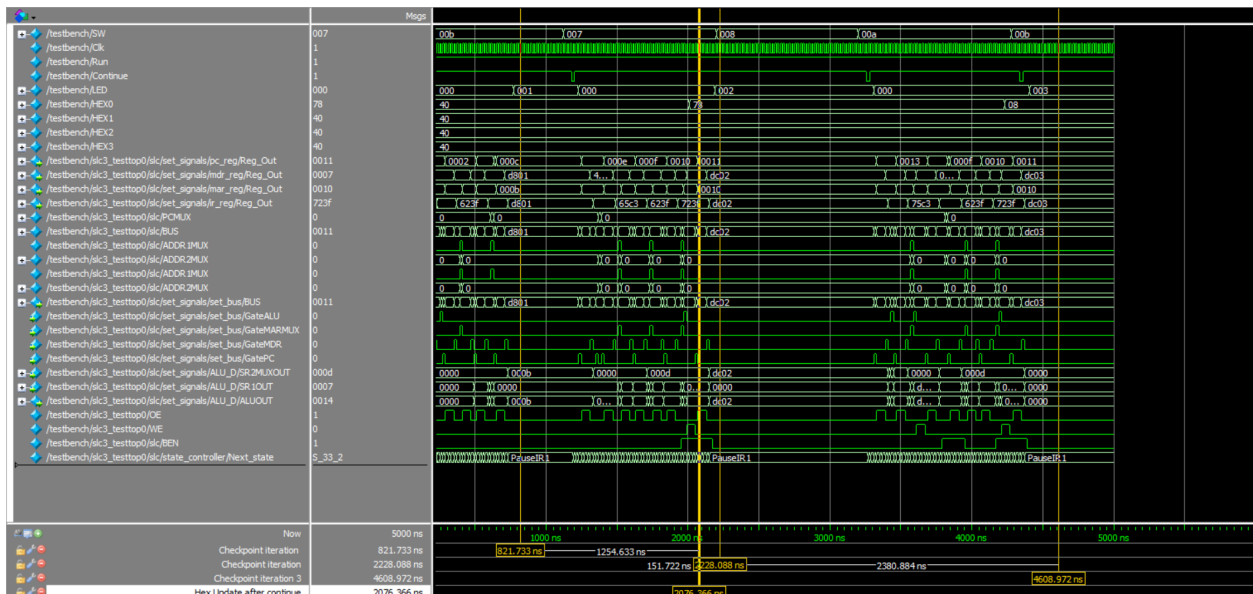


Figure 5: Self Modifying Code Test

- Test 3 Simulation
  - FPGA is reset

- [illegible]

- Test 4 Simulation

- FPGA is reset
- OPCode is set to switches and run is pressed
- To perform XOR, the switches are set to some value and continue is pressed (0101 in this case)
- After Pause is reached switches are set to a different value and continue is pressed once more (0011 in this case)
- The Checkpoints on the LED's update at each pause and once the final value is outputted on the HEX Displays
- The Checkpoint is reset and the program is ready to be run again



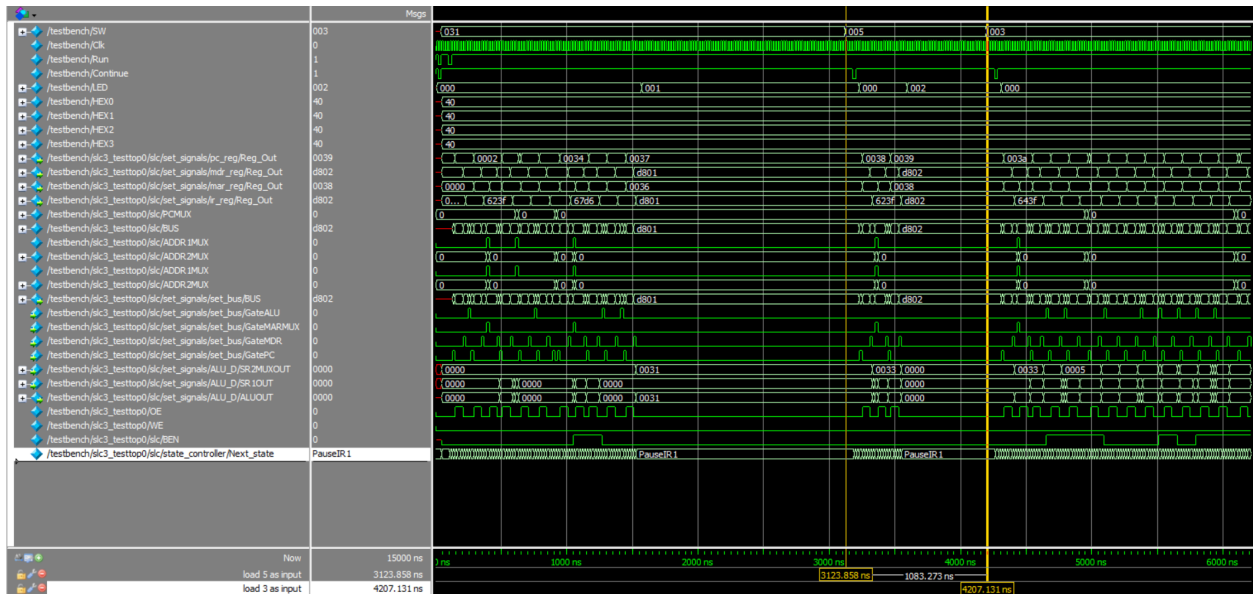


Figure 7.1: Multiplier Test

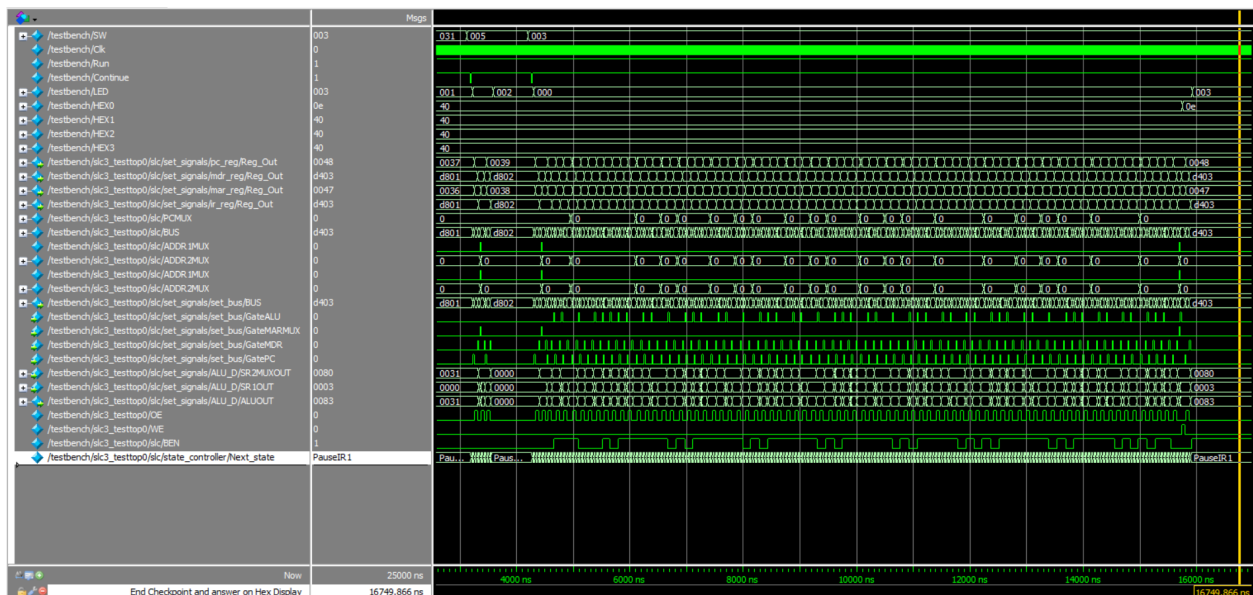


Figure 7.2: Multiplier Test

- Test 5 Simulation
  - FPGA is reset
  - OPCode is set to switches and run is pressed
  - To perform multiplication, the switches are set to some value and continue is pressed (0101 in this case)
  - After Pause is reached switches are set to a different value and continue is pressed once more (0011 in this case)

- The Checkpoints on the LED's update at each pause and once the final value is outputted on the HEX Displays
- The Checkpoint is reset and the program is ready to be run again

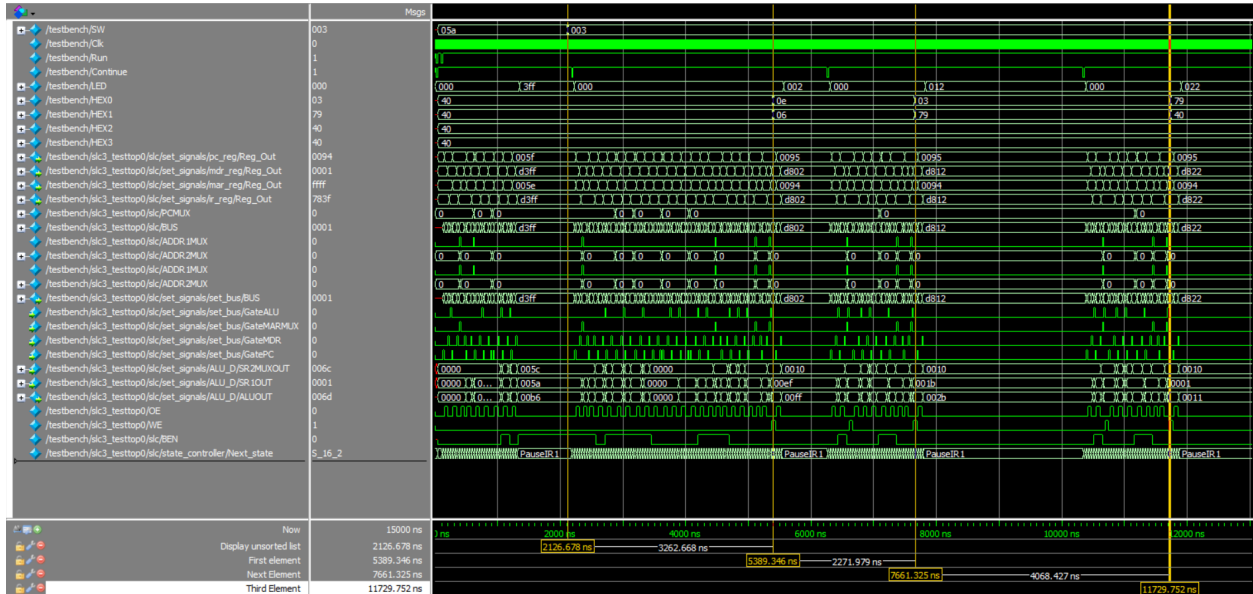


Figure 8.1: Sort Test - Unsorted Elements

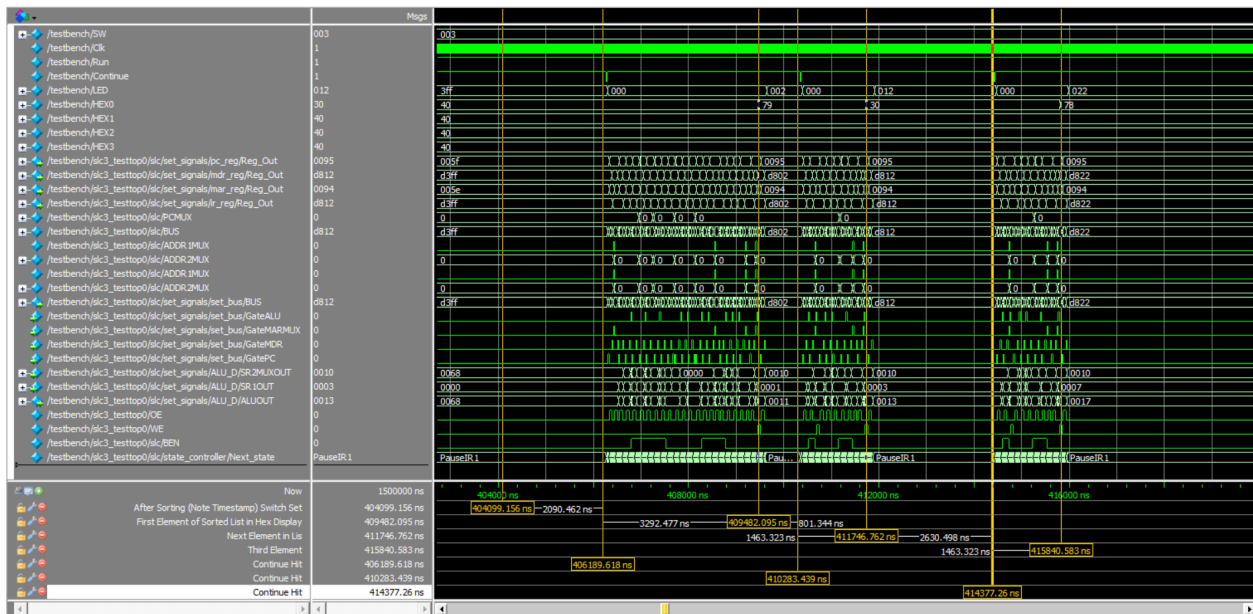


Figure 9: Sort Test - Sorted Elements

- Test 6 Simulation
  - FPGA is reset
  - OPCode is set to switches and run is pressed

- First switches are set to 3 and continue is pressed continuously to view unsorted list
- Once 3FFF is reached on the LED's, switches are set to 2, and continue is pressed once again, this sorts the list
- Finally switches are set to 3 and continue is pressed continuously to view the sorted list

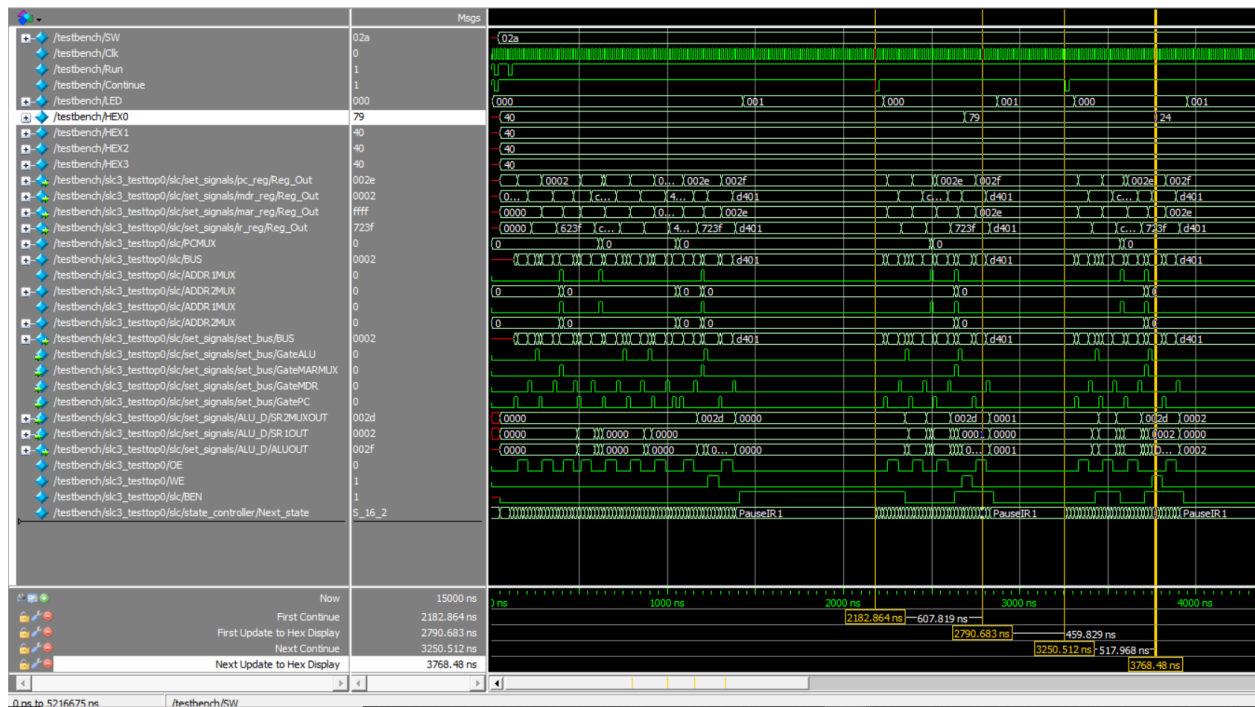


Figure 10: Act Once Test

### ● Test 7 Simulation

- FPGA is reset
- OPCode is set to switches and run is pressed
- To iterate through the counter continue is pressed continuously
- The checkpoint is continuously set to 1 as the Hex Display value increases by one each button press of continue

## Post-Lab Questions

1) Refer to the Design Resources and Statistics in IQT.17-19 and complete the following design statistics table.

LUT	866
DSP	0
Memory (BRAM)	18432
Flip-Flop	268
Frequency	71.98 MHz
Static Power	89.94 mW
Dynamic Power	0.00 mW
Total Power	98.63 mW

2) What is MEM2IO used for, i.e. what is its main function?

Answer: MEM2IO is the interface between the FPGA's memory and the switches, it treats the address of xFFFF as a read from the switches and treats the instruction xFFFF as a write to the Hex seven segment display to output values. This is important for the LDR and STR functions, if the address or instruction is not xFFFF the FPGA will read and write data to memory.

3) What is the difference between BR and JMP instructions?

Answer: Both of these instructions directly affect the PC. However, for the JMP instruction, the PC is changed to a PC of a subroutine. The PC changes back to the PC where the JMP was executed once the subroutine is fully executed and the RET (return) instruction is executed towards the end. For the BR instruction, the change in PC only occurs if the NZP signal in the BR instruction matches the condition codes set in the previous instruction. The PC can be changed by a small, limited number and it does not return to the PC where the BR instruction was executed. For JMP, the PC can be changed to any 16-bit value stored in the base register, and for BR, the PC can only be changed by a 9-bit offset.

4) What is the purpose of the R signal in Patt and Patel? How do we compensate for the lack of the signal in our design? What implications does this have for synchronization?

Answer: The R signal is known as the Ready signal. This signal, once asserted, indicates that the memory access is completed, and the data is ready for computation. Since we lack this signal in our SLC-3 microprocessor, we compensate by having 3 wait or pause states for all instructions that require the system to wait for the R signal. Since the states don't need to wait for an R signal to be raised, all the states that require the system to wait, wait for the same number of cycles, meaning our processor is asynchronous.

## **Conclusion**

Given that this is considered one of the hardest labs in this class, we were extremely happy to have a fully functioning SLC-3 Microprocessor. The two most difficult and important parts of the lab were creating the Datapath and setting the right signals for all the phases and opcodes according to the State Diagram. The most difficult path in the lab was to understand what each signal did and what signals need to be on for each state in the state diagram. In our opinion, more than the experiment documentation, the additional resources on the Lab 5 page such as the resources from ECE 120 were more helpful in creating the microprocessor.

The lab was an amazing way for us to understand how a simple microprocessor such as the SLC-3 microprocessor works. We learned how the Datapath is created, the State diagram, and how the States control the microprocessor. Overall, our objective of creating an SLC-3 Microprocessor, a subset of an LC-3 processor was complete.

## Appendix A

Module: *hexdriver.sv*

Inputs: [3:0] In0,

Outputs: [6:0] Out0

Description: Takes a nibble or 4 bits and converts it to a 7 bit equivalent that is displayable on the HEX# output LEDs of the FPGA.

Purpose: To allow the seven segment displays to display what the input and output of our code are.

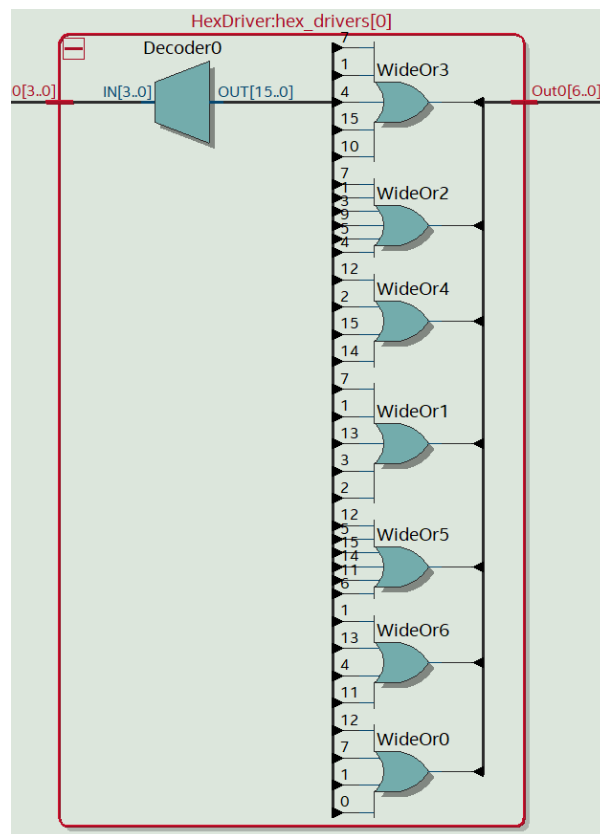


Figure 11: RTL Block diagram for `hexdriver.sv`

*Module: datapath.sv*

Inputs: Clk, Reset, LD\_MAR, LD\_MDR, LD\_IR, LD\_BEN, LD\_CC, LD\_REG, LD\_PC, LD\_LED, SR2MUX, ADDR1MUX, MARMUX, GatePC, GateMDR, GateALU, GateMARMUX, MIO\_EN, DRMUX, SR1MUX, PCMUX, ADDR2MUX, ALUK, MDR\_In, PC\_in, bus\_temp,

Outputs: LED, MAR, MDR, IR, PC, PC\_Now, BUS, BEN

Description: Construction of the LC3 datapath, a combination of MUXes and register signals connected to the bus via tristate gate,

Purpose: To implement the SLC3 the datapath handles directing all the signals in the correct way such that the outputs and inputs through MEM2IO are correct.

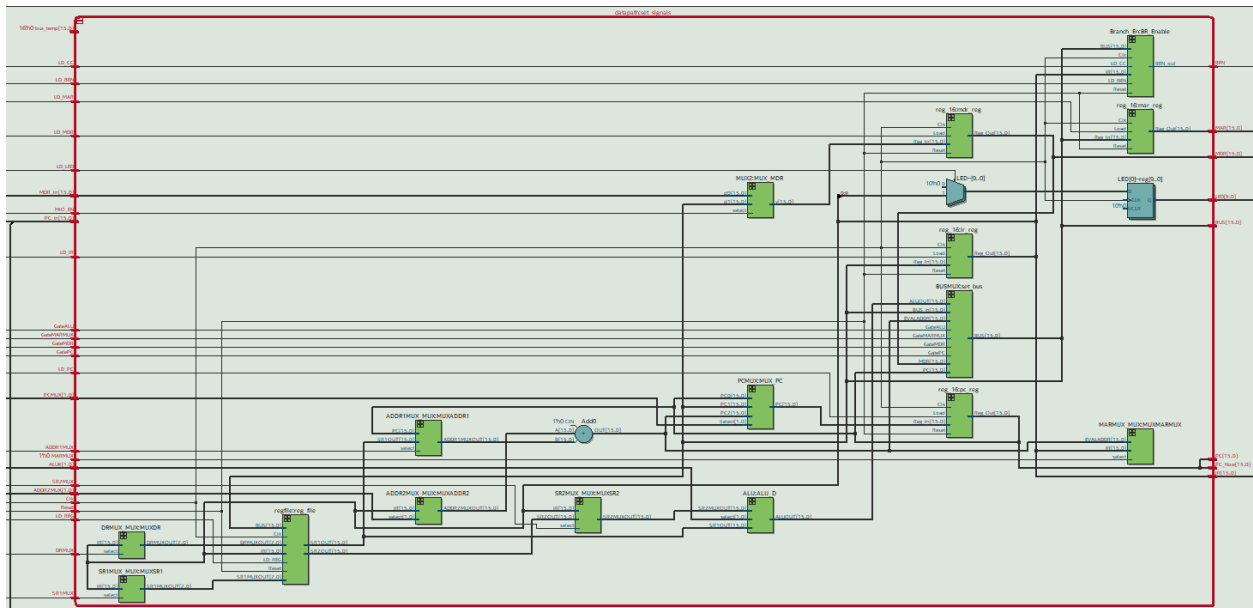


Figure 12: RTL Block diagram datapath.sv

*Module: slc3.sv*

Inputs: Clk, Reset, LD\_MAR, LD\_MDR, LD\_IR, LD\_BEN, LD\_CC, LD\_REG, LD\_PC, LD\_LED, SR2MUX, ADDR1MUX, MARMUX, GatePC, GateMDR, GateALU, GateMARMUX, MIO\_EN, DRMUX, SR1MUX, PCMUX, ADDR2MUX, ALUK, MDR\_In, PC\_in, bus\_temp,

Outputs: LED, MAR, MDR, IR, PC, PC\_Now, BUS, BEN

Description: instantiates datapath, MEM2IO, and the ISDU, as well as the WE register, and constructs the HexDriver variables that are used to condense the instructions.

Purpose: To start the LC3 so that it can begin to run instructions

*Module: SEXT.sv*

Inputs: signed in, parameter size

Outputs: signed out

Description: Takes the signed value of the inputted value of size (parameter - 1) and extends the output to be the 16-bit sign extension of the input

Purpose: To make sign extension easy with a simple call of a module rather than manually extending each signal.

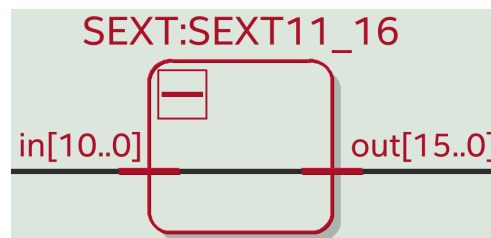


Figure 13: RTL Block diagram for SEXT.sv

*Module: MUX2.sv*

Inputs: d0, d1, select

Outputs: y

Description: parameterized 2 bit mux, used for MIO.En MUX to get MDR value.

Purpose: to implement MIO MUX in the datapath

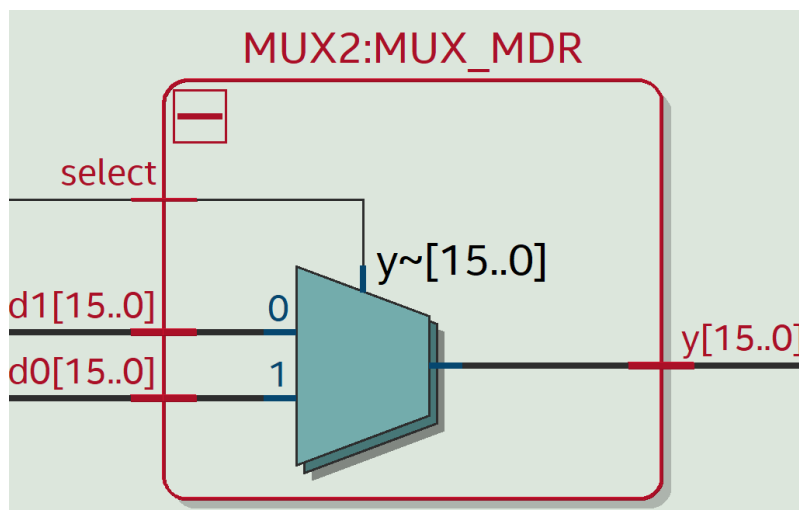


Figure 14: RTL Block diagram for MUX2.sv



Module: PCMUX.sv

Inputs: PC0, PC1, PC2,select

Outputs: PC

Description: MUX, used to choose the input for PC, and what will get passed on to the BUS if GatePC is High.

Purpose: to implement the PCMUX in datapath

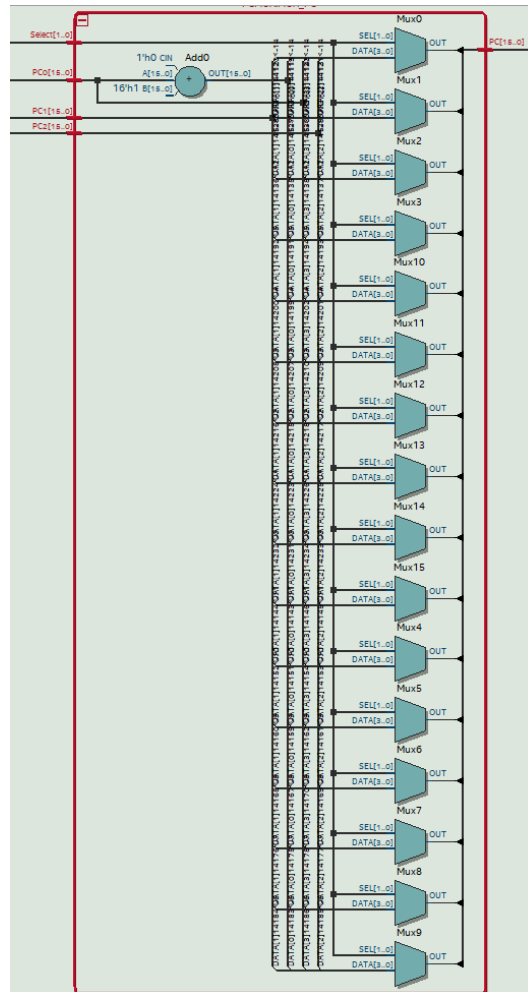


Figure 15: RTL Block diagram for PCMUX.sv

Module: *BUSMUX.sv*

Inputs: GateALU, GateMARMUX, GateMDR, GatePC, MDR, PC, BUS\_in, EVALADDR, ALUOUT,

Outputs: BUS

Description: Tristate buffer, used to choose which signal from MAR, MDR, PC, or ALU gets to be on the BUS

Purpose: to implement the tristate buffer in a way that is executable on the FPGA

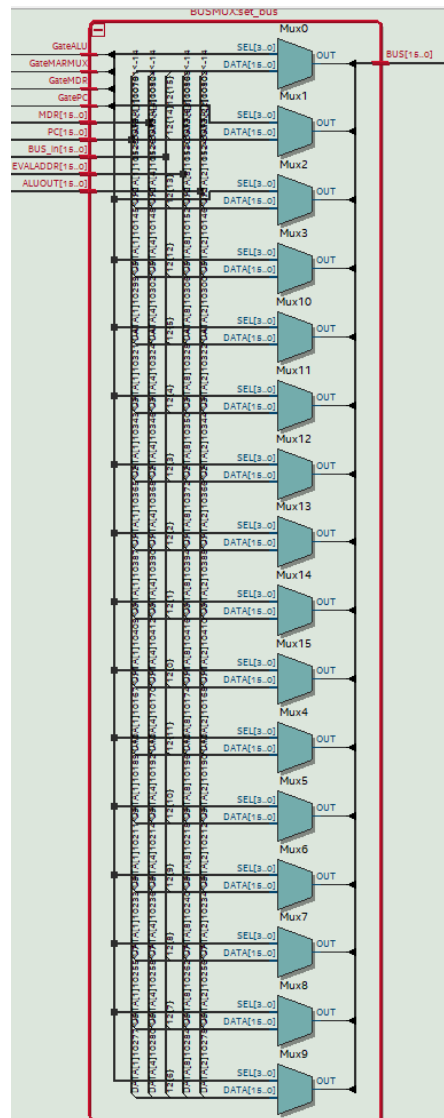


Figure 16: RTL Block diagram for *BUSMUX.sv*

Module: SR1MUX\_MUX.sv

Inputs: IR, select

Outputs: SR1MUXOUT

Description: MUX that chooses the value of the SR1out that goes into the regfile

Purpose: to implement the SR1MUX in datapath

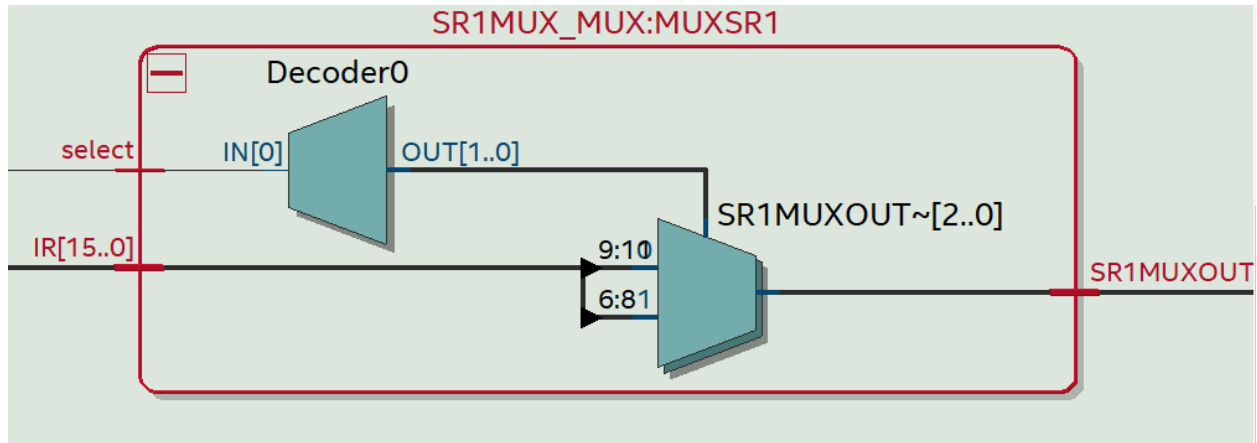


Figure 17: RTL Block diagram for SR1MUX\_MUX.sv

Module: DRMUX\_MUX.sv

Inputs: IR, select

Outputs: DRMUXOUT

Description: MUX to choose the value of the DR between 0-7 to be used in the regfile later

Purpose: to implement DRMUX in datapath

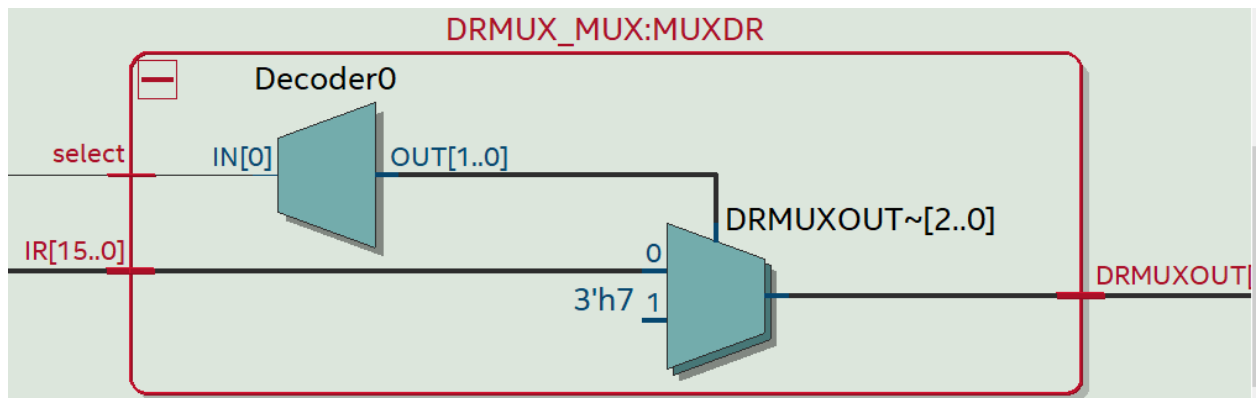


Figure 18: RTL Block diagram for DRMUX\_MUX.sv

Module: *SR2MUX\_MUX.sv*

Inputs: IR, SR2OUT, select

Outputs: SR2MUXOUT

Description: 2 bit MUX to choose the value of the B input for the ALU to be used in the ALU module later

Purpose: to implement SR2MUX in datapath

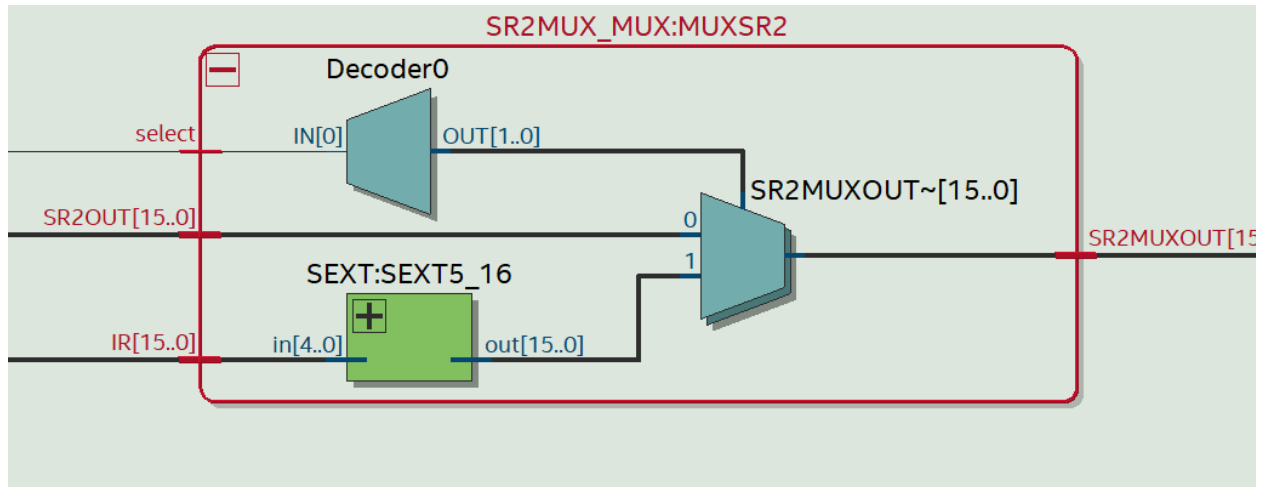


Figure 19: RTL Block diagram for *SR2MUX\_MUX.sv*

Module: *ADDR2MUX\_MUX.sv*

Inputs: IR, select

Outputs: ADDR2MUXOUT

Description: 4 bit MUX that chooses either 0 or some sign extended value of the IR, to be used in the Address addition

Purpose: to implement ADDR2 in datapath

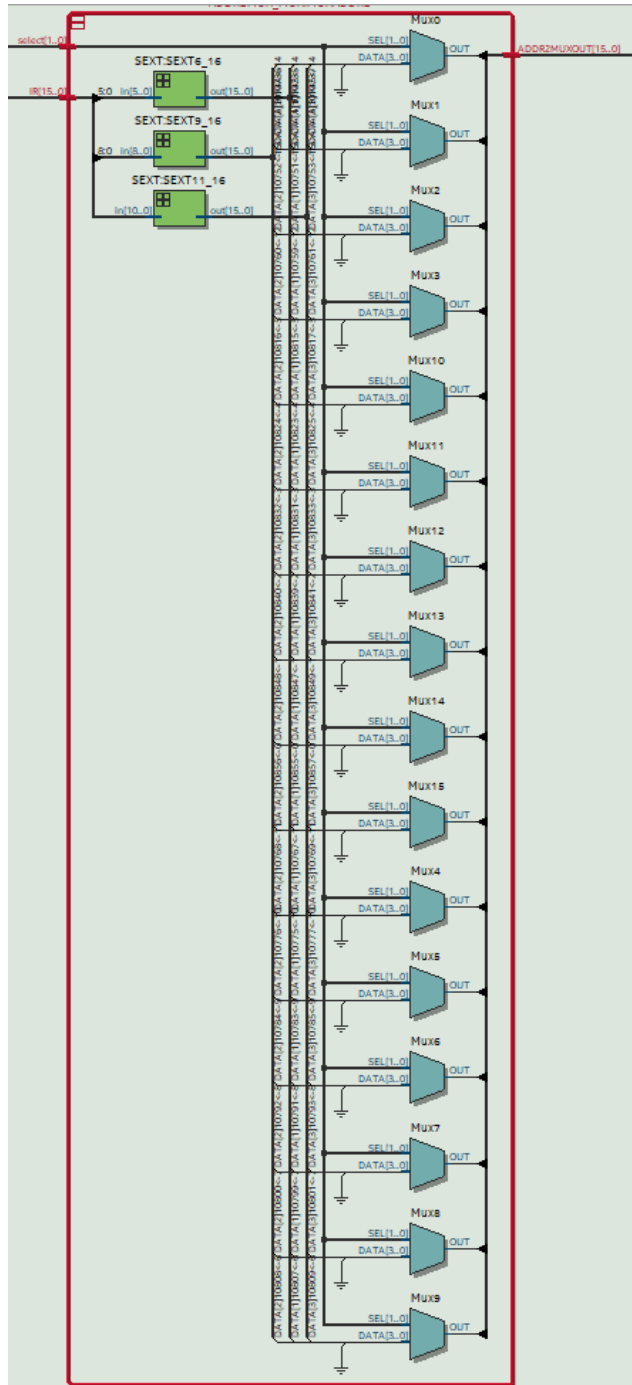


Figure 20: RTL Block diagram for ADDR2MUX\_MUX.sv

Module: ADDR1MUX\_MUX.sv

Inputs: SR1OUT, PC, , select

Outputs: ADDR1MUXOUT

Description: 2 bit MUX that chooses either the PC or the SR1OUT value from the reg file for the address addition.

Purpose: to implement ADDR1 in datapath

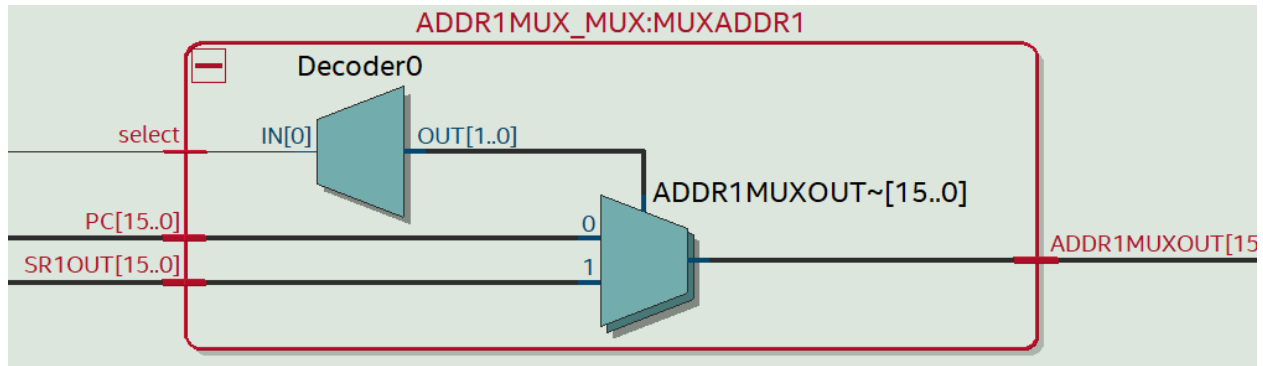


Figure 21: RTL Block diagram for ADDR1MUX\_MUX.sv

Module: MARMUX\_MUX.sv

Inputs: EVALADDR, IR, select

Outputs: MARMUXOUT

Description: MUX that chooses the value of MAR from either memory or the BUS.

Purpose: to implement the MUX feeding into MAR

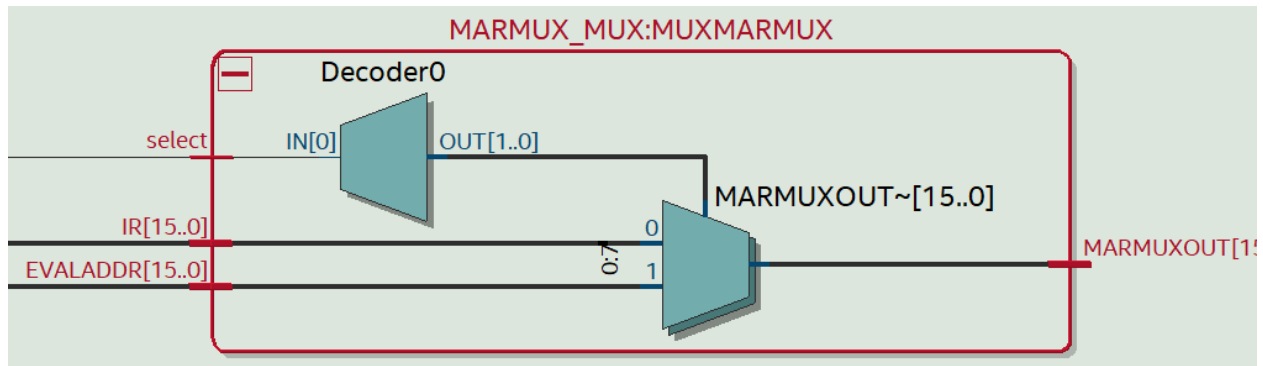


Figure 22: RTL Block diagram for MARMUX\_MUX.sv

Module: regfile.sv

Inputs: LD\_REG, Clk, Reset, IR, BUS, DRMUXOUT, SR1MUXOUT,

Outputs: SR1OUT, SR2OUT, R0, R1, R2, R3, R4, R5, R6, R7

Description: creates 8 registers 0-7 that are reset on reset, and are loaded with the value of the BUS if LD\_REG is high, otherwise the outputs from the regfile are defined as the register corresponding to the output of the SR MUXes.

Purpose: to create the temp registers that the LC3 needs to run its functions.

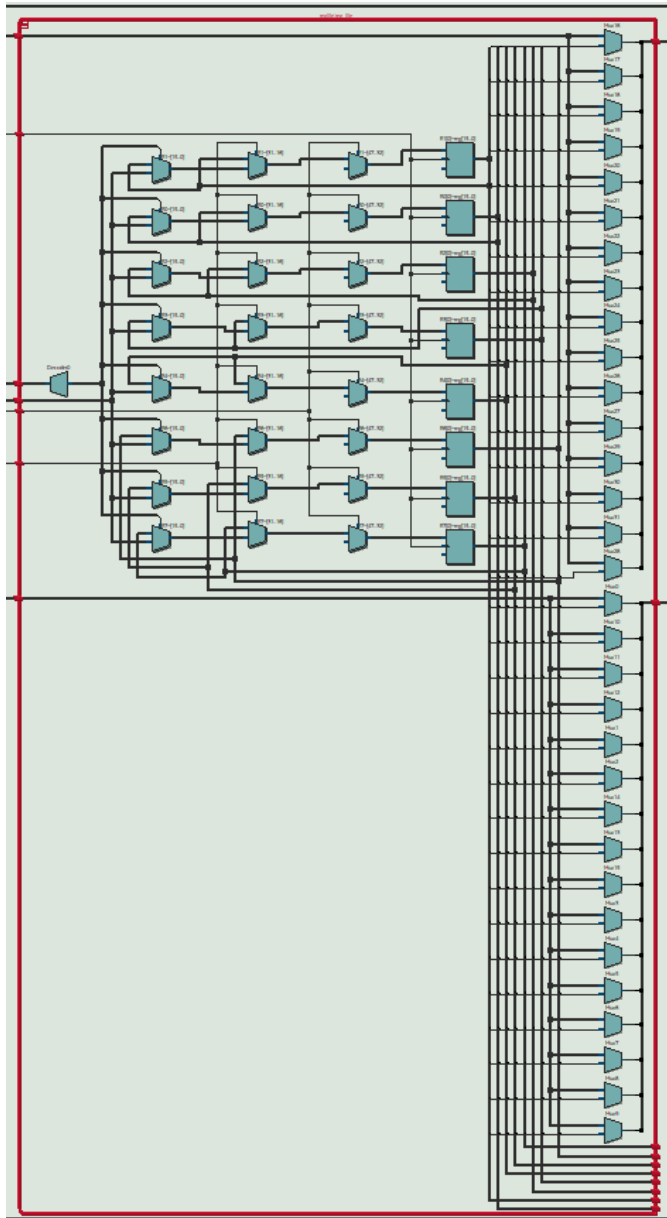


Figure 23: RTL Block diagram for regfile.sv

Module: *reg\_16.sv*

Inputs: Clk, Reset, Load, RegIn

Outputs: RegOut

Description: 16 bit registers that construct MAR, PC, MDR, and IR, the registers reset to 0 on reset and are filled with the input value taken off the bus.

Purpose: To properly store the values of PC, MDR, MAR and IR

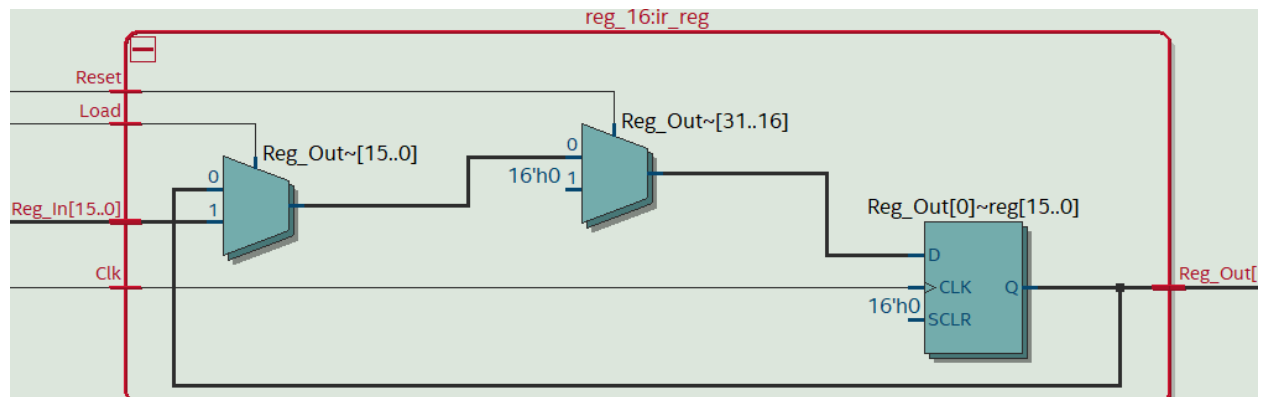


Figure 24: RTL Block diagram for *reg\_16.sv*



Module: ISDU.sv

Inputs: Clk, Reset, Continue, Run, Opcode, IR\_5, IR\_11, BEN

Outputs: LD\_MAR, LD\_MDR, LD\_BEN, LD\_IR, LD\_CC, LD\_PC, LD\_LED, GatePC, GateALU, GateMDR, GateMARMUX, PCMUX, DRMUX, SR1MUX, SR2MUX, ADDR1MUX, ADDR2MUX, ALUK, MEM\_OE, MEM\_WE.

Description: State Machine for the LC3, goes through fetch decode and execute states

Purpose: To allow proper function of the LC3 opcodes.

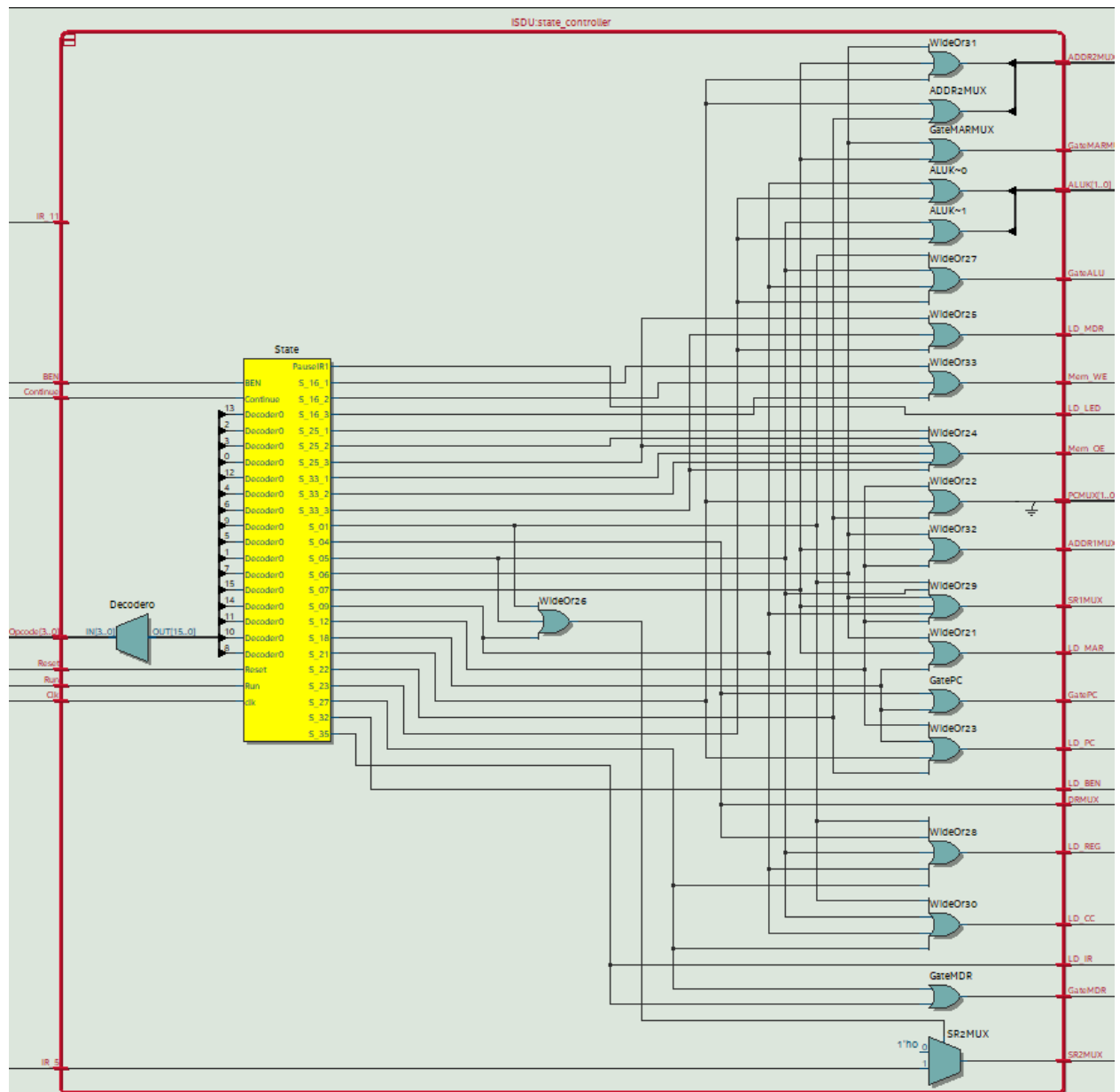
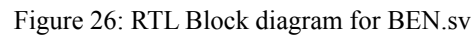


Figure 25: RTL Block diagram for ISDU.sv

Description: Branch logic to go to the control unit (ISDU). Given LD\_CC and LD\_BEN. Logic from IR [11:9] is loaded into NZP if LD\_CC is on, the resulting logic is a SOP of all NZP and Purpose: so that when a BR is hit, the next state branches accordingly



Purpose: to put register file outputs on the bus with the ALU operation applied if GateALU is high.

