Heuristic Search

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CAP4630 – Artificial Intelligence

Today

- Informed Search
- Search Heuristics
- Greedy Search
- A* Search
- Admissible Heuristics
- Semi-Lattice of Heuristics
- Consistency of Heuristics
- Effect of Heuristics on Performance
- Optimality

Informed Search

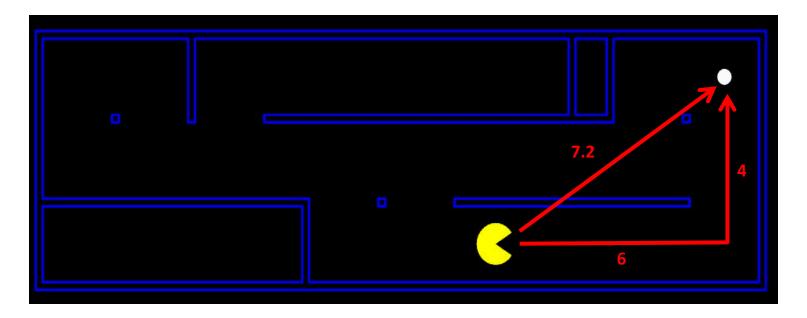
- Basic idea: Use some information to know if we are making progress towards goal
 - Are we getting "hotter"
 - Are we getting "colder"



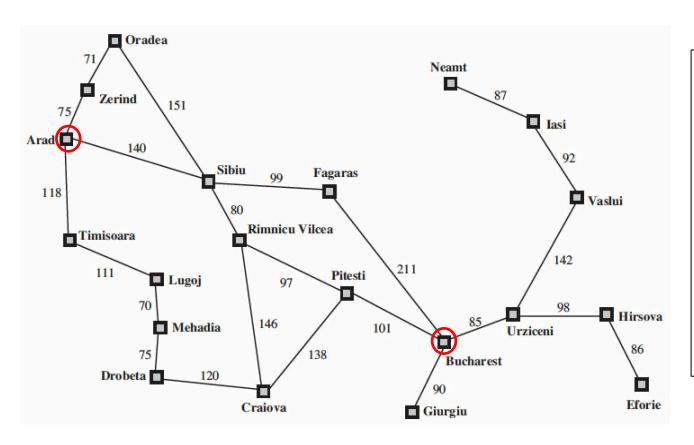


Search Heuristics

- Heuristic defined:
 - A function f:{States} → R
 - that estimates "how close" a state is to a goal
 - heuristics are problem-specific
 - Examples for Pac-Man: Manhattan distance, Euclidean distance



Example: Heuristic Function

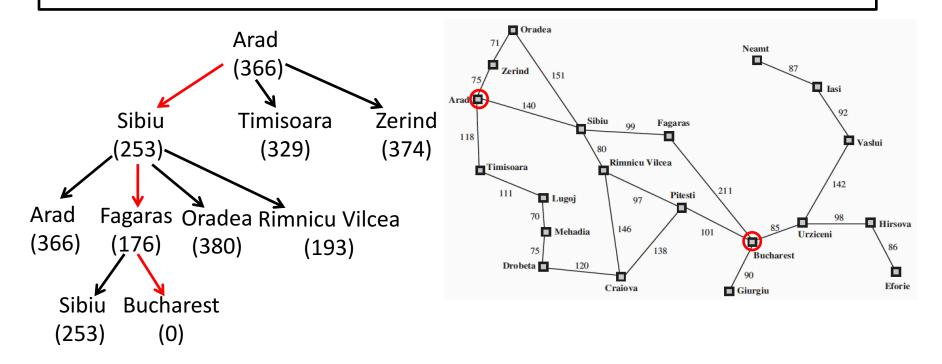


Heuristic function h(x)

Straight-line distance to Bucharest		
Arad	366	
Bucharest	0	
Craiova	160	
Fagaras	176	
Sibiu	253	
 Zerind	374	

Greedy Search

Strategy: Expand the node that is closest to goal according to the heuristic function

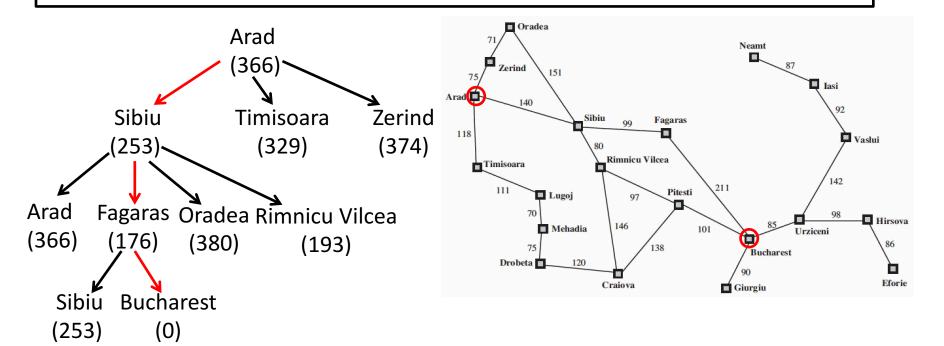


Total distance traveled: 450



Greedy Search

Strategy: Expand the node that is closest to goal according to the heuristic function



Total distance traveled: 450

Shortest path distance: 418 Q: What went wrong, here?

A* Search

- Pronounced: "A-star" search
- Basic Idea: Combine UCS and Greedy
 - UCS path cost is "backward cost" g(n)
 - Greedy value is "forward cost" h(n)
- A* Search orders by the minimum value of the sum of these two numbers:

$$f(n) = g(n) + h(n)$$

Q: If h(n) = 0, what kind of search is this?

Q: If g(n) = 0, what kind of search is this?

A* Search

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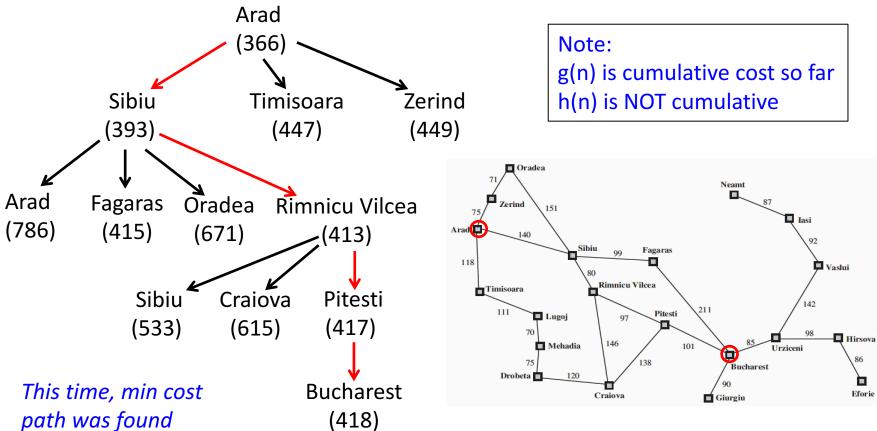
$$f(n) = g(n) + h(n)$$

Q: If h(n) = 0, what kind of search is this? UCS

Q: If g(n) = 0, what kind of search is this? Greedy

A* Search Example

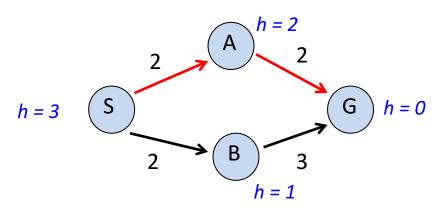
Strategy: Order by minimum value of f(n) = g(n) + h(n)

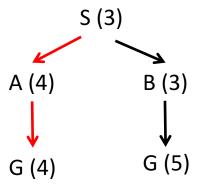


demo: astar

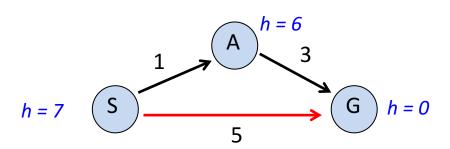
A* Implementation Notes

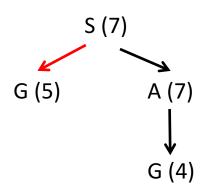
• Stop when dequeue a goal, not when enqueue it





A* can fail if the heuristic is "bad"





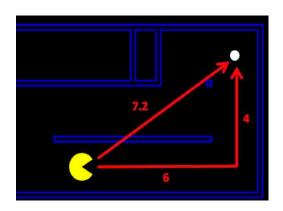
Admissible Heuristics

- Basic idea: The heuristic value should never be greater than the actual cost
- A heuristic h is admissible (optimistic) if:

$$0 \le h(n) \le h^*(n)$$

where $h^*(n)$ is the **true** cost to the nearest goal

- Example:
 - Manhattan and Euclidean distance for Pac-Man
- In practice:
 - Finding an admissible heuristic can be challenging
 - Often done by considering a "relaxed" model



Semi-Lattice of Heuristics

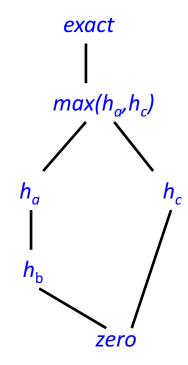
• Dominance: $h_a \ge h_b$ if

$$\forall n: h_a(n) \geq h_b(n)$$

Max of admissible heuristics is admissible

$$h_{new}(n) = max[h_a(n), h_c(n)]$$

- → heuristics form a *semi-lattice*
- Trivial heuristics
 - Top of lattice is the exact heuristic
 - Bottom of lattice is the zero heuristic (what search method is this?)



(dominance hierarchy)

Consistency of Heuristics

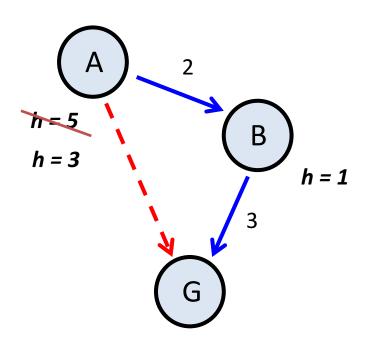
- Consistency is a stronger property than admissibility
- Admissibility:

$$h(A) \le actual cost from A to G$$

Consistency:

$$h(A) - h(B) \le actual cost(A, B)$$

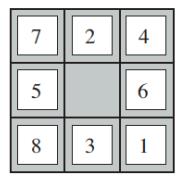
- Impact of heuristic being consistent:
 - The f-value along a path never decreases

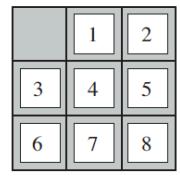


Effect of Heuristic on Performance

• 8-puzzle:

- h1 = number of misplaced tiles
- h2 = sum of Manhattan distances of misplaced tiles





Start State

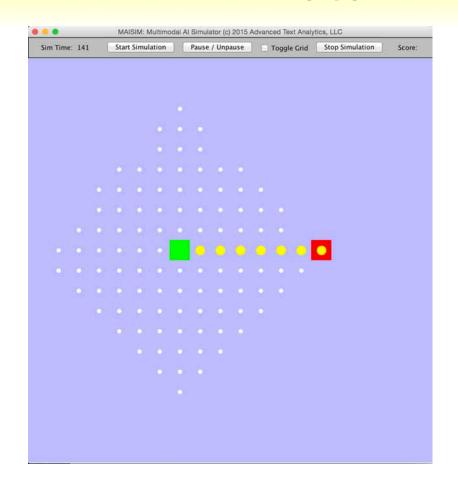
Goal State

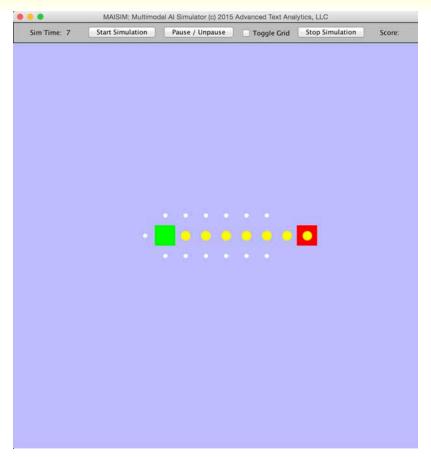
Number of nodes examined:

# moves in solution	h1	h2
6	20	18
12	277	73
18	3056	363
24	39135	1641



UCS v. A* Contours





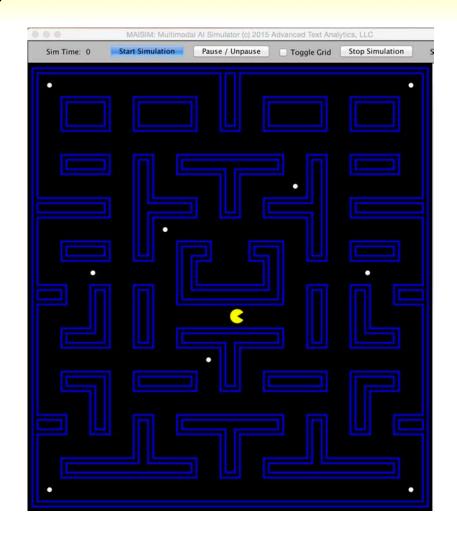
UCS expands equally in all directions

A* expands toward goal, but also some in other directions to ensure optimality



A* Applications

- Pathing/routing applications
- Resource planning/logistics
- Robot motion planning
- Video games/simulations
- Language understanding
- Machine translation
- Speech recognition
- ... and more
- ... including our Pac-Man tour



A* for Pac-Man Tour

• A* heuristic: Manhattan distance to nearest remaining food, plus maximum manhattan distance from any remaining food to any other remaining food

9 food pellets:

run time nodes examined solution path length

UCS	A*
1.230 sec	0.565 sec
162,926	70,160
153	153

• 14 food pellets:

run time nodes examined solution path length

UCS	A*
63.902	28.800 sec
5,303,471	2,683,625
185	185

A* and Optimality

- General tree search (states can appear multiple times)
 - A* is optimal if heuristic is admissible
 - UCS is optimal (special case, h=0, so also admissible)
- Graph search (no state appears more than once)
 - A* is optimal if heuristic is consistent
 - UCS is optimal (special case, h=0, so also consistent)
- Notes:
 - Consistency implies admissibility
 - Most admissible heuristics from relaxed problems are also consistent