

- Formulate problems as LPs.
- Equivalence theorem
- Minimum cost trading to achieve target Sector exposures.

Midterm 25 points x 4.

- Least squares
- Multi-objective / constrained LS.
- Linear optimization modeling
- Linear optimization geometry
(equivalence theorem)

no coding / no simplex

You are allowed to use all course materials (lecture slides, precept, homework, books)

No internet No communication

NO precept next week!

(Q1) Formulate problems as LPs

(a) $A \in \mathbb{R}^{M \times n}$ $b \in \mathbb{R}^M$ find the vector $x \in \mathbb{R}^n$ that minimizes.

$$\sum_{i=1}^M \max\{0, a_i^T x + b_i\}. \quad (1)$$

$$c_i = \max\{0, a_i^T x + b_i\}.$$

$c_i \geq 0$

$$c_i \geq a_i^T x + b_i$$

$$\underbrace{\qquad}_{\text{minimize}} \quad \underbrace{1^T \cdot c.}_{\text{---}}$$

$$c = \begin{pmatrix} c_1 \\ \vdots \\ c_M \end{pmatrix}$$

$$\text{minimize}_{(C, X)} \quad (\mathbb{I}^T, 0) \begin{pmatrix} C \\ X \end{pmatrix} \quad \mathbb{I}^T \cdot C$$

s.t.

$$\left\{ \begin{array}{l} c_i \geq 0 \\ c_i - a_i^T x \geq b_i \end{array} \right.$$

(2)

Let X_1 minimize (1)
 (C^*, X_2) minimize (2)

By definition of (1). $\forall X$

$$\sum \max\{0, a_i^T x_1 + b_i\} \leq \sum \max\{0, a_i^T x_2 + b_i\}$$

$$\leq \sum \max\{0, a_i^T x + b_i\}.$$

for any C satisfies $\left\{ \begin{array}{l} c_i \geq 0 \\ c_i - a_i^T x \geq b_i \end{array} \right.$

(C, X_1) is a feasible solution of (2)

(C^*, X_2)

$$\Rightarrow \mathbb{I}^T C \geq \mathbb{I}^T C^*$$

$$(\mathbb{I}^T, 0) \begin{pmatrix} C \\ X_1 \end{pmatrix} \geq (\mathbb{I}^T, 0) \begin{pmatrix} C^* \\ X_2 \end{pmatrix}$$

$$\underset{(c, x)}{\text{minimize}} \quad (\mathbb{1}^T, 0) \begin{pmatrix} c \\ x \end{pmatrix}$$

s.t.

$$\left\{ \begin{array}{l} c_i \geq 0 \\ c_i - a_i^T x \geq b_i \end{array} \right.$$

(2)

$$c_i \geq \max \{0, a_i^T x + b_i\}.$$

$$\Rightarrow \mathbb{1}^T c = \sum \max \{0, a_i^T x_1 + b\}.$$

$$\mathbb{1}^T c^* = \sum \max \{0, a_i^T x_2 + b\}.$$

$$\sum \max \{0, a_i^T x_1 + b\} \geq \sum \max \{0, a_i^T x_2 + b\}$$

$$\Rightarrow \sum \max \{0, a_i^T x_1 + b\}$$

$$= \sum \max \{0, a_i^T x_2 + b\}.$$

CQ2 equivalence theorem.

$$P = \{x \mid Ax \leq b\}.$$

(1) Show that every vertex is an extreme point.

vertex: if $\exists c$ such that x is the unique optimum of
minimize $c^T y$
s.t. $y \in P$.

extreme point: x not on a straight line between any other points in the set.

$$\underline{x = \lambda y + (1-\lambda)z}$$

vertex \Rightarrow extreme point.

(proof by contradiction)

Assume x is vertex, but not extreme point.

$$\exists y, z \neq x \quad y, z \in P \quad x = \lambda y + (1-\lambda)z$$

vertex $\exists c$. such that $\min_{x \in P} c^T x$
 $c^T y < c^T z \quad \forall w \in P$ st. $w \in P$

$$\underline{c^T y, c^T z > c^T x}$$

$$\Rightarrow \underline{c^T x = \lambda c^T y + (1-\lambda) c^T z > c^T x}$$

Contradiction! $P = \{x | Ax \leq b\}$

x

x basic

② extreme point \Rightarrow feasible solution
 $\{a_1, a_2, a_3, a_4\}$.

- $\{a_i | i \in I(x)\}$ has n linear independent vectors.

$$I(x) = \{i \in \{1, \dots, m\} | a_i^T x = b_i\}$$

$$I(x) = \{1, 2, 3, 4\}$$

(proof by contradiction)

Assume x is extreme point but not basic feasible solution.

Then $\{a_i : i \in I(x)\}$ does not span

$$\mathbb{R}^n \Rightarrow \exists d \in \mathbb{R}^n \quad a_i^T d = 0 \quad i \in I(x)$$
$$x = \lambda y + (1-\lambda) z$$

Let $y = x + \varepsilon d$ $z = x - \varepsilon d$.

for any i

$$\begin{array}{l} a_i^T y \leq b_i \\ a_i^T z \leq b_i \end{array} \quad \begin{array}{l} \text{if } i \in I(x) \quad a_i^T x = b_i \Rightarrow a_i^T y = b_i \\ a_i^T y = a_i^T (x + \varepsilon d) = b_i \\ a_i^T z = b_i \end{array}$$

$$\forall i \notin I(x) \quad a_i^T x < b_i. \quad \underbrace{\varepsilon \text{ small enough}}$$

$$\Rightarrow \underbrace{a_i^T y < b_i}_{\text{and}} \quad \underbrace{a_i^T z < b_i}_{\text{and}}$$

$$x = 0.5y + 0.5z \quad \text{and } y, z \in P.$$

$\Rightarrow x$ not extreme point

contradiction!

$$\underline{a_i^T y} = \underline{a_i^T x} + (\varepsilon) \underline{a_i^T d}$$

Dm.

$$\underbrace{\quad}_{\text{bi}} \quad \underbrace{\quad}_{\text{I}} \quad \underline{a_i^T x}$$

Q3. h^{curr} → n-vector (portfolio)
with entries giving the dollar
value invested in the n assets.

$$\text{total value} = \mathbb{1}^T h^{\text{curr}}$$

we seek a new portfolio. h .

$$\text{s.t. } \mathbb{1}^T h = \mathbb{1}^T h^{\text{curr}}$$

Same total ~~vector~~ value.

m - industry sectors.
(Tech / bio).

s m-vector denotes the
sector exposures (dollar value)
to the m sectors. $s = Sh$.

S is $m \times n$ matrix

$$\begin{cases} S_{ij} = 1 & \text{asset } j \text{ in sector } i \\ S_{ij} = 0 & \text{asset } j \text{ not in} \\ & \text{sector } i. \end{cases}$$

The new portfolio must
have a given sector exposure
sdes.

$$s_{des} = Sh$$

Among all portfolios that have the same value as our current portfolio and achieve the desire exposures

We wish to minimize the trading cost.

$$\sum_{i=1}^n k_i (h_i - h_i^{\text{curr}})^2$$

$$k_i \geq 0$$

Explain how to find h .
give KKT equations.

$$h^{\text{curr}} = (v_1, v_2, \dots, v_n) \cdot n$$

$$\underline{h} = (h_1, h_2, \dots, h_n) \cdot n$$

$$\text{constraints: } \underline{Sh} = \underline{s^{\text{des}}}$$

$$\underline{U}^T h = \underline{I}^T h^{\text{curr}}$$

$$\text{minimize } \sum_{i=1}^n k_i (\underline{h}_i - \underline{h}_i^{\text{curr}})^2$$

$$\text{Let } K = \begin{pmatrix} k_1 & & \\ & \ddots & \\ & & k_n \end{pmatrix}$$

$$\underline{U}^T h^{\text{curr}} = \underline{V}$$

$$\min_{\underline{h}} \|\underline{K}\underline{h} - \underline{K}\underline{h}^{\text{curr}}\|_2$$

$$\begin{bmatrix} S \\ I^T \end{bmatrix} \underline{h} = \begin{bmatrix} S^{\text{des}} \\ V \end{bmatrix}$$

$$\min \|Ax-b\|^2$$

[\leq KT equation:

$$A = \underline{K}, \quad b = \underline{K}\underline{h}^{\text{curr}}$$

$$C = \begin{bmatrix} S \\ I^T \end{bmatrix}, \quad d = \begin{bmatrix} S^{\text{des}} \\ V \end{bmatrix}$$

$$\begin{bmatrix} 2A^TA & C^T \\ C & 0 \end{bmatrix} \begin{bmatrix} X^* \\ z \end{bmatrix} = \begin{bmatrix} 2A^Tb \\ d \end{bmatrix}$$