


Agenda

- Branch & Bound Overview
- Disjunctive constraints problem
- Production/Distribution problem

Logistics

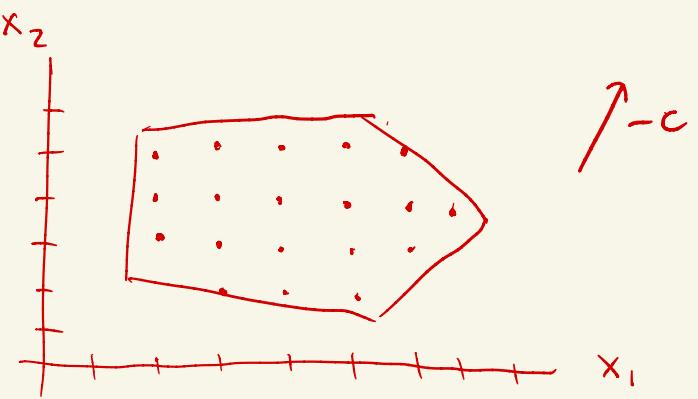
- HW 8 out due Friday April 30 9pm
- This is the last precept
- Final project
 - more details to be announced
 - coding project
 - 24 hours to complete
 - Window open May 5
 - Window closed May 7
- Midterm 2 - 9 point bonus given

Graduate School

- Reach out to me if you're interested/ have questions

Mixed integer program

$$\begin{aligned} \min_{\mathbf{x}} \quad & \mathbf{c}^T \mathbf{x} \\ \text{s.t.} \quad & \mathbf{A}\mathbf{x} \leq \mathbf{b} \\ & \mathbf{x}_i \in \mathbb{Z}, i \in \mathcal{I} \end{aligned}$$



Relaxation

Branch & Bound

Branch & Bound Algorithm

$$\begin{array}{ll} \min_x c^T x & P_{IP} = \{x \mid Ax \leq b, x_i \in \mathbb{Z} \text{ } i \in I\} \\ \text{s.t. } x \in P_{IP} & \end{array}$$

Iterations

1. Branch: create/refine the partition of P_{IP} and get S^j

2. Bound:

- Compute lower and upper bounds

$$L_j = \Phi_{LB}(S^j) \quad U_j = \Phi_{UB}(S^j) \quad \forall j$$

- Update global bounds on $c^T x^*$

$$L = \min_j L_j \quad U = \max_j U_j$$

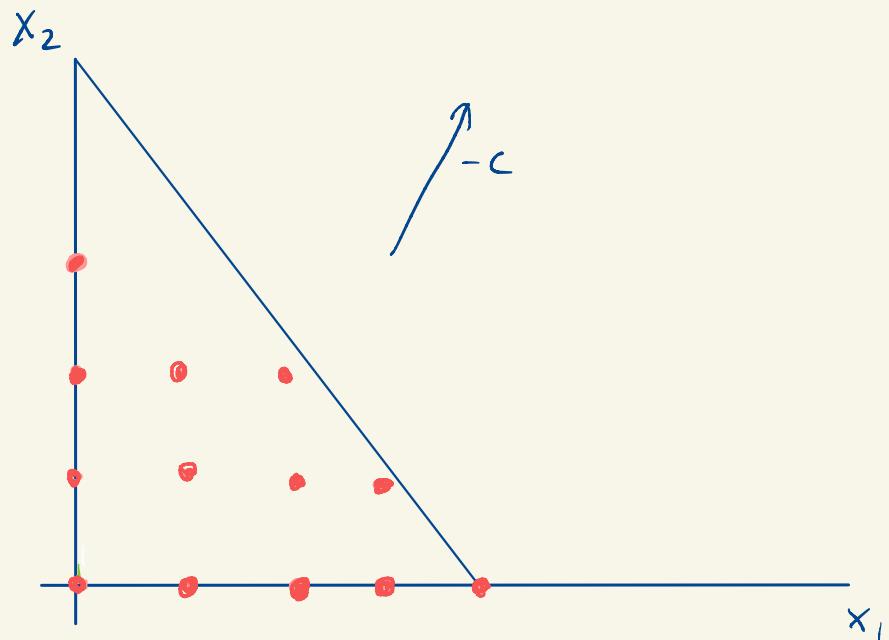
3. If $U - L \leq \epsilon$ break

- Assume U is nonincreasing, L nondecreasing
- Pruning: S^j is active if $L_j \leq \min_j U_j$
else: it is inactive and we can prune it

Branch & Bound example

$$\begin{aligned} \min_x \quad & -2x_1 - 3x_2 \\ \text{s.t.} \quad & \frac{2}{9}x_1 + \frac{1}{4}x_2 \leq 1 \\ & \frac{1}{7}x_1 + \frac{1}{3}x_2 \leq 1 \\ & x \geq 0 \\ & x \in \mathbb{Z} \end{aligned}$$

$$x^{\star \text{ IP}} = (2, 2)$$



Disjunctive constraints

Suppose we are given m constraints $a_i^T x \geq b_i; i=1, \dots, m$
Model the requirement that at least k of them are satisfied

Production/ Distribution problem

A company produces a set of K products at I plants and ships these products to J market zones ($k=1, \dots, K$) ($i=1, \dots, I$) ($j=1, \dots, J$)

v_{ik} : cost of producing 1 unit of product k at plant i

c_{ijk} : cost of shipping 1 unit of product k from plant i to zone j

f_{ik} : fixed cost of producing product k at plant i

M_{ik} : maximal quantity of product k produced at plant i

m_{ik} : minimal quantity of product k that can be produced at plant i if plant i introduces a nonzero quantity

ϱ_{ik} : capacity of plant i to produce 1 unit of product k

Q_i : capacity of plant i

d_{jk} : demand for product k at market zone j

(a) formulate minimizing the total cost as an integer program

- (b) No plant can produce more than k_1 products
- (c) Every product can be produced in at most I_1 plants
- (d) For a particular product k_0 , plant 3 must produce it if neither plant 1 or plant 2 produces it
- (e) Each market zone must be sourced by exactly one plant for all products