

# Version 4.3

# C Cross Compiler User's Guide for ST Microelectronics STM8

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# Table of Contents

Pr	eface
Organization of this Manual	1
Cha	pter 1
Introdu	ction
Introduction	
Document Conventions.	
Typewriter font	
Italies	
[ Brackets ]	
Conventions	
Command Line	6
Flags	6
Compiler Architecture	
Predefined Symbol	
Linking	9
Programming Support Utilities	
Listings	10
Optimizations	10
Support for ROMable Code	11
Support for eeprom	12
Memory Models	12
Cha	pter 2
Tutorial Introdu	ction
Acia.c, Example file	16
Default Compiler Operation	
Compiling and Linking	19
Step 1: Compiling	19
Step 2: Assembler	20
Step 3: Linking	21
Step 4: Generating S-Records file	
Linking Your Application	24
Generating Automatic Data Initialization	
Specifying Command Line Options	28
Cha	pter 3
Programming Environn	
Introduction	
Modifying the Runtime Startup	
Description of Runtime Startup Code	

Initializing data in RAM	36
Memory Models for code smaller than 64K	39
Memory Models for code larger than 64K	39
Handling Large Code and Constants	41
Bit Variables	42
The const and volatile Type Qualifiers	43
Performing Input/Output in C	45
Referencing Absolute Addresses	46
Accessing Internal Registers	48
Placing Data Objects in The Bss Section	48
Placing Data Objects in Short Range Memory	49
Setting Short Range Size	49
Placing Data Objects in Long Range Memory	
Placing Data Objects in the EEPROM Space	51
Redefining Sections	53
Inserting Inline Assembly Instructions	55
Inlining with pragmas	55
Inlining with _asm	56
Inlining Labels	58
Writing Interrupt Handlers	
Placing Addresses in Interrupt Vectors	60
Inline Function	
Interfacing C to Assembly Language	
Register Usage	
Heap Management Control with the C Compiler	
Modifying The Heap Location	68
Data Representation	70
Chap	
Using The Comp	oiler
Invoking the Compiler	
Compiler Command Line Options	
File Naming Conventions	
Generating Listings.	
Generating an Error File	
Return Status	
Examples	81
C Library Support	
How C Library Functions are Packaged	
Inserting Assembler Code Directly	
Linking Libraries with Your Program	
Integer Library Functions	
Common Input/Output Functions	83

Functions Implemented as Macros	83
Including Header Files	
Descriptions of C Library Functions	
Generate inline assembly code	
Abort program execution	
Find absolute value	
Arccosine	
Arcsine	90
Arctangent	91
Arctangent of y/x	
Convert buffer to double	93
Convert buffer to integer	94
Convert buffer to long	
Test or get the carry bit	96
Round to next higher integer	97
Verify the recorded checksum	98
Verify the recorded checksum	99
Verify the recorded checksum	100
Verify the recorded checksum	101
Cosine	102
Hyperbolic cosine	103
Divide with quotient and remainder	104
Copy a buffer to an eeprom buffer	105
Erase the eeprom space	
Propagate fill character throughout eeprom buffer	107
Exit program execution	108
Exponential	109
Find double absolute value	110
Copy a moveable code segment in RAM	
Round to next lower integer	
Find double modulus	
Extract fraction from exponent part	
Get character from input stream	115
Get a text line from input stream	
Test the interrupt mask bit	117
Test the interrupt line level	
Test for alphabetic or numeric character	119
Test for alphabetic character	
Test for control character	121
Test for digit	
Test for graphic character	
Test for lowercase character	124
Test for printing character	125

Test for punctuation character	. 126
Test for whitespace character	
Test for uppercase character	
Test for hexadecimal digit	. 129
Find long absolute value	. 130
Scale double exponent	
Long divide with quotient and remainder	. 132
Natural logarithm	
Common logarithm	. 134
Test for maximum	
Scan buffer for character	. 136
Compare two buffers for lexical order	. 137
Copy one buffer to another	. 138
Copy one buffer to another	
Propagate fill character throughout buffer	. 140
Test for minimum	
Extract fraction and integer from double	
Raise x to the y power	
Output formatted arguments to stdout	
Put a character to output stream	
Put a text line to output stream	
Generate pseudo-random number	
Sin	
Hyperbolic sine	
Output arguments formatted to buffer	
Real square root	
Seed pseudo-random number generator	
Concatenate strings.	
Scan string for first occurrence of character	
Compare two strings for lexical order	
Copy one string to another	
Find the end of a span of characters in a set	
Find length of a string	
Concatenate strings of length n	
Compare two n length strings for lexical order	. 164
Copy n length string	
Find occurrence in string of character in set	
Scan string for last occurrence of character	
Find the end of a span of characters not in set	
Scan string for first occurrence of string	
Convert buffer to double	
Convert buffer to long	
Convert buffer to unsigned long	

Tangent	173
Hyperbolic tangent	174
Convert character to lowercase if necessary	
Convert character to uppercase if necessary	
11	
Chaj	pter 5
Using The Assen	nbler
Invoking castm8	
Object File	
Listings	181
Assembly Language Syntax	
Instructions	
Labels	183
Temporary Labels	184
Constants	
Expressions	
Macro Instructions	
Conditional Directives	190
Sections	191
Bit Handling	
Includes	
Branch Optimization	194
Old Syntax	
C Style Directives	
Assembler Directives	
Align the next instruction on a given boundary	
Define the default base for numerical constants	
Switch to the predefined .bsct section.	
Turn listing of conditionally excluded code on or off	
Allocate constant(s)	
Allocate constant block	
Turn listing of debug directives on or off	
Allocate variable(s)	
Conditional assembly	
Conditional assembly	
Stop the assembly	
End conditional assembly	
End conditional assembly	
End macro definition	
End repeat section	
Give a permanent value to a symbol	
Assemble next byte at the next even address relative to	
ž	

start of a section	212
Generate error message.	213
Conditional assembly	214
Conditional assembly	215
Conditional assembly	216
Conditional assembly	217
Conditional assembly	218
Conditional assembly	219
Conditional assembly	220
Conditional assembly	221
Conditional assembly	222
Conditional assembly	223
Conditional assembly	
Include text from another text file	225
Turn on listing during assembly	226
Give a text equivalent to a symbol	227
Create a new local block	228
Define a macro	229
Send a message out to STDOUT	231
Terminate a macro definition	232
Turn on or off listing of macro expansion	233
Turn off listing.	234
Disable pagination in the listing file	235
Creates absolute symbols	236
Sets the location counter to an offset from the beginning	of a
section	237
Start a new page in the listing file	238
Specify the number of lines per pages in the listing file	239
Repeat a list of lines a number of times	240
Repeat a list of lines a number of times	241
Restore saved section	243
Terminate a repeat definition	244
Save section	245
Turn on or off section crossing	246
Define a new section	
Give a resetable value to a symbol	249
Insert a number of blank lines before the next statement is	n the
listing file	250
Place code into a section.	251
Specify the number of spaces for a tab character in the lis	sting
file	252
Define default header	253
Declare hit symbol as being defined elsewhere	254

Declare a variable to be visible	255
Declare symbol as being defined elsewhere	
	Chapter 6
Using Th	e Linker
Introduction	
Overview	
Linker Command File Processing	
Inserting comments in Linker commands	
Linker Options	
Global Command Line Options	
Segment Control Options	
Segment Grouping	
Linking Files on the Command line	
Example	
Include Option	
Example	
Private Region Options	
Symbol Definition Option	
Reserve Space Option	
Handle Dependencies	
Section Relocation	
Address Specification	
Overlapping Control	
Setting Bias and Offset	
Setting the Bias	
Setting the Offset	
Using Default Placement	
Bit Segment Handling	
Linking Objects	
Linking Library Objects	
Library Order	
Libraries Setup Search Paths	
Automatic Data Initialization	
Descriptor Format	
Moveable Code	
Checksum Computation	
DEFs and REFs.	
Special Topics	
Private Name Regions	
Renaming Symbols	
Absolute Symbol Tables	
1.2001400 07111001 140100	

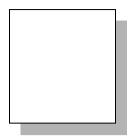
Description of The Map File	293
Return Value	
Linker Command Line Examples	295
•	
	Chapter 7
Debugging	<b>Support</b>
Generating Debugging Information	
Generating Line Number Information	
Generating Data Object Information	298
The cprd Utility	300
Command Line Options	
Examples	
The clst utility	302
Command Line Options	
	Chapter 8
Programming	g Support
The chex Utility	306
Command Line Options	306
Return Status	308
Examples	308
The clabs Utility	
Command Line Options	
Return Status	311
Examples	311
The clib Utility	313
Command Line Options	
Return Status	314
Examples	314
The cobj Utility	316
Command Line Options	316
Return Status	317
Examples	317
The Ctat Utility	318
Command Line Options	
Return Status	319
Examples	319
The cvdwarf Utility	320
Command Line Options	
Return Status	
Examples	322

Chapter A	L
Compiler Error Messages	
Parser (cpstm8) Error Messages	
Code Generator (cgstm8) Error Messages	
Assembler (castm8) Error Messages	
Linker (clnk) Error Messages	
Chapter E	3
Modifying Compiler Operation	
The Configuration File	
Changing the Default Options	
Creating Your Own Options	
Example	
Chapter C	7
STM8 Machine Library	
Update an int bitfield in near memory354	
Eeprom char bit field update	
Write a char in eeprom	,
Write a long int in eeprom357	
Write a short int in eeprom	
Move a structure in eeprom	,
Add float to float360	ļ
Compare floats	
Divide float by float	
Add float to float in memory	
Multiply float by float in memory	
Subtract float from float in memory	
Multiply float by float	
Negate a float	
Subtract float from float	
Convert float to integer	
Convert float into long integer	
Compare a float in memory to zero	
Get a long word from memory	
Get a long word from memory	
Get a word from far memory	
Get a word from far memory	
Get a word from far memory	
Get a word from far memory	
Get a word from far memory	

Get a word from far memory	379
Quotient of integer division	380
Integer multiplication	381
Convert integer into float	382
Convert integer into long	383
Convert integer into long	384
Convert integer into long	385
Perform C switch statement on char	386
Perform C switch statement on long	387
Perform C switch statement on integer	388
Long integer addition	
Long integer addition	390
Bitwise AND for long integers	391
Long integer compare	392
Quotient of long integer division	
Long addition	
Long addition	395
Long bitwise AND	
Long shift left	397
Long multiplication in memory	
Negate a long integer in memory	399
Long bitwise OR	400
Signed long shift right	
Long subtraction	
Long subtraction	403
Unsigned long shift right	404
Long bitwise exclusive OR	405
Long integer shift left	
Remainder of long integer division	407
Multiply long integer by long integer	408
Negate a long integer	
Bitwise OR with long integers	410
Long integer right shift	
Long test against zero	412
Long integer subtraction	413
Long integer subtraction	
Long integer compare with overflow	415
Convert long integer into float	416
Load memory into long register	
Quotient of unsigned long integer division	
Remainder of unsigned long integer division	
Unsigned long integer shift right	

Bitwise exclusive OR with long integers	s 421
Compare a long integer to zero	422
Far pointer addition	423
Far pointer addition	424
Put a long integer in memory	425
Put a long integer in memory	426
Put a word in far memory	427
Get a far pointer from far memory	428
Get a far pointer from far memory	429
Get a far pointer from far memory	430
Get a far pointer from far memory	431
Store long register in far memory	432
Store long register in memory	433
Quotient of signed char division	434
Quotient of signed char division	
Remainder of signed char division	
Remainder of signed char division	437
Multiply long integer by unsigned byte.	438
Convert unsigned integer into float	439
Convert unsigned integer into long	440
Convert unsigned integer into long	441
Convert unsigned integer into long	442
Convert unsigned long integer into float	
Multiply unsigned integers with long res	
Multiply signed integers with long resul	t445
Store a far pointer into far memory	446
Copy a structure into another	447
Copy a structure in far memory	448
Copy a large structure in far memory	449
Copy a structure into another	
Copy a large structure into another	451
Store a far pointer into far memory	452
Copy a structure into another	453
Copy a structure in far memory	454
Copy a large structure in far memory	
Copy a structure into another	456
Copy a large structure into another	457
	CI + D
	Chapter D
	Compiler Passes
The cpstm8 Parser	
Command Line Options	460

Extra verifications	465
Return Status	466
Example	466
The cgstm8 Code Generator	467
Command Line Options	467
Return Status	469
Example	469
The costm8 Assembly Language Optimizer	
Command Line Options	470
Disabling Optimization	471
Return Status	471
Example	471



# **Preface**

The Cross Compiler User's Guide for STM8 is a reference guide for programmers writing C programs for STM8 microcontroller environments. It provides an overview of how the cross compiler works, and explains how to compile, assemble, link and debug programs. It also describes the programming support utilities included with the cross compiler and provides tutorial and reference information to help you configure executable images to meet specific requirements. This manual assumes that you are familiar with your host operating system and with your specific target environment.

# Organization of this Manual

This manual is divided into eight chapters and four appendixes.

**Chapter 1,** "<u>Introduction</u>", describes the basic organization of the C compiler and programming support utilities.

**Chapter 2,** "*Tutorial Introduction*", is a series of examples that demonstrates how to compile, assemble and link a simple C program.

Chapter 3, "<u>Programming Environments</u>", explains how to use the features of C for STM8 to meet the requirements of your particular application. It explains how to create a runtime startup for your application, and how to write C routines that perform special tasks such as: serial I/O, direct references to hardware addresses, interrupt handling, and assembly language calls.

**Chapter 4,** "*Using The Compiler*", describes the compiler options. This chapter also describes the functions in the C runtime library.

**Chapter 5,** "<u>Using The Assembler</u>", describes the STM8 assembler and its options. It explains the rules that your assembly language source must follow, and it documents all the directives supported by the assembler.

**Chapter 6,** "*Using The Linker*", describes the linker and its options. This chapter describes in detail all the features of the linker and their use.

**Chapter 7,** "<u>Debugging Support</u>", describes the support available for COSMIC's C source level cross debugger and for other debuggers or incircuit emulators

**Chapter 8,** "*Programming Support*", describes the programming support utilities. Examples of how to use these utilities are also included.

**Appendix A**, "Compiler Error Messages", is a list of compile time error messages that the C compiler may generate.

**Appendix B**, "*Modifying Compiler Operation*", describes the "configuration file" that serves as default behaviour to the C compiler.

**Appendix C**, "<u>STM8 Machine Library</u>", describes the assembly language routines that provide support for the C runtime library.

**Appendix D**, "<u>Compiler Passes</u>", describes the specifics of the parser, code generator and assembly language optimizer and the command line options that each accepts.

This manual also contains an Index.

# CHAPTER 1

# Introduction

This chapter explains how the compiler operates. It also provides a basic understanding of the compiler architecture. This chapter includes the following sections:

- Introduction
- Document Conventions
- Compiler Architecture
- · Predefined Symbol
- Linking
- Programming Support Utilities
- Listings
- Optimizations
- Support for ROMable Code
- Support for eeprom
- Memory Models

#### Introduction

The C cross compiler targeting the STM8 microcontroller reads C source files, assembly language source files, and object code files, and produces an executable file. You can request listings that show your C source interspersed with the assembly language code and object code that the compiler generates. You can also request that the compiler generate an object module that contains debugging information that can be used by COSMIC's C source level cross debugger or by other debuggers or in-circuit emulators.

You begin compilation by invoking the **cxstm8** compiler driver with the specific options you need and the files to be compiled.

#### **Document Conventions**

In this documentation set, we use a number of styles and typefaces to demonstrate the syntax of various commands and to show sample text you might type at a terminal or observe in a file. The following is a list of these conventions.

#### Typewriter font

Used for user input/screen output. Typewriter (or courier) font is used in the text and in examples to represent what you might type at a terminal: command names, directives, switches, literal filenames, or any other text which must be typed exactly as shown. It is also used in other examples to represent what you might see on a screen or in a printed listing and to denote executables.

To distinguish it from other examples or listings, input from the user will appear in a shaded box throughout the text. Output to the terminal or to a file will appear in a line box.

For example, if you were instructed to type the compiler command that generates debugging information, it would appears as:

#### cxstm8 +debug acia.c

Typewriter font enclosed in a shaded box indicates that this line is entered by the user at the terminal.

If, however, the text included a partial listing of the file *acia.c* 'an example of text from a file or from output to the terminal' then type-writer font would still be used, but would be enclosed in a line box:

```
/* defines the ACIA as a structure */
struct acia {
   char status;
   char data;
} acia @0x6000;
```

#### NOTE

Due to the page width limitations of this manual, a single invocation line may be represented as two or more lines. You should, however, type the invocation as one line unless otherwise directed.

#### **Italics**

Used for value substitution. *Italic* type indicates categories of items for which you must substitute appropriate values, such as arguments or hypothetical filenames. For example, if the text was demonstrating a hypothetical command line to compile and generate debugging information for any file, it might appear as:

```
cxstm8 +debug file.c
```

In this example, cxstm8 +debug file.c is shown in typewriter font because it must be typed exactly as shown. Because the filename must be specified by the user, however, *file* is shown in italics.

#### [Brackets]

Items enclosed in brackets are optional. For example, the line:

```
[ options ]
```

means that zero or more options may be specified because options appears in brackets. Conversely, the line:

```
options
```

means that one or more options must be specified because options is not enclosed by brackets.

As another example, the line:

```
file1.[o|stm8]
```

means that one file with the extension .o or .sm8 may be specified, and the line:

```
file1 [ file2 . . . ]
```

means that additional files may be specified.

#### **Conventions**

All the compiler utilities share the same optional arguments syntax. They are invoked by typing a command line.

#### Command Line

A command line is generally composed of three major parts:

```
program_name [<flags>] <files>
```

where cprogram\_name is the name of the program to run, flags an optional series of flags, and files a series of files. Each element of a command line is usually a string separated by whitespace from all the others.

#### Flags

Flags are used to select options or specify parameters. Options are recognized by their first character, which is always a '-' or a '+', followed by the name of the flag (usually a single letter). Some flags are simply **yes** or **no** indicators, but some must be followed by a value or some additional information. The value, if required, may be a character string, a single character, or an integer. The flags may be given in any order, and two or more may be combined in the same argument, so long as the second flag can't be mistaken for a value that goes with the previous one.

It is possible for each utility to display a list of accepted options by specifying the **-help** option. Each option will be displayed alphabetically on a separate line with its name and a brief description. If an option requires additional information, then the type of information is

indicated by one of the following code, displayed immediately after the option name:

Code	Type of information
*	character string
#	short integer
##	long integer
?	single character

If the code is immediately followed by the character '>', the option may be specified more than once with different values. In that case, the option name must be repeated for every specification.

For example, the options of the **chex** utility are:

_		
chex	[option:	s] file
	-a##	absolute file start address
	-b##	address bias
	-e##	entry point address
	-f?	output format
	-h	suppress header
	+h*	specify header string
	-m#	maximum data bytes per line
	-n*>	output only named segments
	-0*	output file name
	-p	use paged address format
	-pa	use paged address for data
	-p1##	page numbers for linear mapping
	-pn	use paged address in bank only
	-pp	use paged address with mapping
	-s	output increasing addresses
	-w	output word addresses
	-x*	exclude named segment

chex accepts the following distinct flags:

Flag	Function
-a	accept a long integer value
-b	accept a long integer value
-е	accept a long integer value
-f	accept a single character
-h	simply a flag indicator
+h	accept a character string
-m	accept a short integer value
-n	accept a character string and may be repeated
-0	accept a character string
-p	simply a flag indicator
-pl	accept a long integer value
-pn	simply a flag indicator
-рр	simply a flag indicator
-S	simply a flag indicator
-W	simply a flag indicator
-X	accept a character string and may be repeated

# **Compiler Architecture**

The C compiler consists of several programs that work together to translate your C source files to executable files and listings. **cxstm8** controls the operation of these programs automatically, using the options you specify, and runs the programs described below in the order listed:

**cpstm8** - the C preprocessor and language parser. *cpstm8* expands directives in your C source and parses the resulting text.

**cgstm8** - the code generator. *cgstm8* accepts the output of *cpstm8* and generates assembly language statements.

**costm8** - the assembly language optimizer. *costm8* optimizes the assembly language code that *cgstm8* generates.

**castm8** - the assembler. *castm8* converts the assembly language output of *costm8* to a relocatable object module.

# **Predefined Symbol**

The COSMIC compiler defines the \_\_CSMC\_\_ preprocessor symbol. It expands to a numerical value whose each bit indicates if a specific option has been activated:

bit 2	set if unsigned char option specified (-pu)
bit 4	set if reverse bitfield option specified (+rev)
bit 5	set if no enum optimization specified (-pne)

### Linking

**clnk** combines all the object modules that make up your program with the appropriate modules from the C library. You can also build your own libraries and have the linker select files from them as well. The linker generates an executable file which, after further processing with the *chex* utility, can be downloaded and run on your target system. If you specify debugging options when you invoke **cxstm8**, the compiler will generate a file that contains debugging information. You can then use the COSMIC's debugger to debug your code.

# **Programming Support Utilities**

Once object files are produced, you run **clnk** (the linker) to produce an executable image for your target system; you can use the programming support utilities to inspect the executable.

**chex** - absolute hex file generator. *chex* translates executable images produced by the linker into hexadecimal interchange formats, for use

with in-circuit emulators and PROM programmers. *chex* produces the following formats:

- Motorola S-record format
- standard Intel hex format

**clabs** - absolute listing utility. *clabs* translates relocatable listings produced by the assembler by replacing all relocatable information by absolute information. This utility must to be used only after the linker.

**clib** - build and maintain object module libraries. *clib* allows you to collect related files into a single named library file for convenient storage. You use it to build and maintain object module libraries in standard library format.

**cobj** - object module inspector. *cobj* allows you to examine standard format executable and relocatable object files for symbol table information and to determine their size and configuration.

**cvdwarf** - ELF/DWARF format converter. *cvdwarf* allows you to convert a file produced by the linker into an ELF/DWARF format file.

## Listings

Several options for listings are available. If you request no listings, then error messages from the compiler are directed to your terminal, but no additional information is provided. Each error is labelled with the C source file name and line number where the error was detected.

If you request an assembly language and object code listing with interspersed C source, the compiler merges the C source as comments among the assembly language statements and lines of object code that it generates. Unless you specify otherwise, the error messages are still written to your terminal. Your listing is the listing output from the assembler.

# **Optimizations**

The C cross compiler performs a number of compile time and optimizations that help make your application smaller and faster:

- The compiler will perform arithmetic operations in 8-bit precision if the operands are 8-bit.
- The compiler eliminates unreachable code.
- Branch shortening logic chooses the smallest possible jump/ branch instructions. Jumps to jumps and jumps over jumps are eliminated as well.
- Integer and float constant expressions are folded at compile time.
- Redundant load and store operations are removed.
- **enum** is large enough to represent all of its declared values, each of which is given a name. The names of **enum** values occupy the same space as type definitions, functions and object names. The compiler provides the ability to declare an **enum** using the smallest type char, int or long:
- The compiler performs multiplication by powers of two as faster shift instructions
- An optimized switch statement produces combinations of tests and branches, jump tables for closely spaced case labels, a scan table for a small group of loosely spaced case labels, or a sorted table for an efficient search.
- The functions in the C library are packaged in three separate libraries; one of them is built without floating point support. If your application does not perform floating point calculations, you can decrease its size and increase its runtime efficiency by linking with the non-floating-point version of the modules needed.

# Support for ROMable Code

The compiler provides the following features to support ROMable code production. See Chapter 3 for more information.

- · Referencing of absolute hardware addresses;
- Control of the STM8 interrupt system;

- Automatic data initialization;
- User configurable runtime startup file;
- Support for mixing C and assembly language code; and
- User configurable executable images suitable for direct input to a PROM programmer or for direct downloading to a target system.

# Support for eeprom

The compiler provides the following features to support **eeprom** handling:

- @eeprom type qualifier to describe a variable as an eeprom location. The compiler generates special sequences when the variable is modified
- Library functions for erasure, initialization and copy of eeprom locations.

#### NOTE

The basic routine to program an eeprom byte is located in the library file **eeprom.s** and has been written using the default input/output address. This file must be modified if using a different base address.

These basic routines are not updating any watchdog, so applications enabling a watchdog must modify these routines to add watchdog updates in the wait loops.

# **Memory Models**

The STM8 compiler supports several memory models allowing you to choose the best configuration for your processor and your application. You can choose to use 2-bytes addressing (applications smaller than 64k bytes) or 3-bytes addressing. You can also specify where variables are located in the memory space by default: inside or outside the Short Range (short addressing). For more information, please refer to Chapter 3.

For information on using the compiler, see <u>Chapter 4</u>.
For information on using the assembler, see <u>Chapter 5</u>.
For information on using the linker, see <u>Chapter 6</u>.
For information on debugging support, see <u>Chapter 7</u>.
For information on using the programming utilities, see <u>Chapter 8</u>.
For information on the compiler passes, see <u>Appendix D</u>.

CHAPTER 2

# **Tutorial Introduction**

This chapter will demonstrate, step by step, how to compile, assemble and link the example program **acia.c**, which is included on your distribution media. Although this tutorial cannot show all the topics relevant to the COSMIC tools, it will demonstrate the basics of using the compiler for the most common applications.

In this tutorial you will find information on the following topics:

- Default Compiler Operation
- · Compiling and Linking
- Linking Your Application
- Generating Automatic Data Initialization
- Specifying Command Line Options

# Acia.c, Example file

The following is a listing of *acia.c.* This C source file is copied during the installation of the compiler:

```
STM8 EXAMPLE PROGRAM WITH INTERRUPT HANDLING
    Copyright (c) 2008 by COSMIC Software
* /
#include <iostm8.h>
#define SIZE64
#define TRDE0x80
                         /* buffer size */
                         /* transmit ready bit */
   Authorize interrupts.
#define rim() asm("rim")
/* Some variables
* /
/* write pointer */
char *ptecr;
/* Character reception.
   Loops until a character is received.
* /
char getch (void)
     char c;
                         /* character to be returned */
     while (ptlec == ptecr) /* equal pointers => loop */
     c = *ptlec++; /* get the received char */
     if (ptlec >= &buffer[SIZE])/* put in in buffer */
         ptlec = buffer;
     return (c);
   Send a char to the SCI.
void outch(char c)
     while (!(USART SR & TRDE))/* wait for READY */
      ;
     USART DR = c; /* send it */
```

```
/* Character reception routine.
      This routine is called on interrupt.
      It puts the received char in the buffer.
 * /
@interrupt void recept(void)
                                      /* clear interrupt */
      USART SR;
      *ptecr++ = USART DR;
                                      /* get the char */
      if (ptecr >= &buffer[SIZE]) /* put it in buffer */
            ptecr = buffer;
/* Main program.
      Sets up the SCI and starts an infinite loop
      of receive transmit.
 * /
void main(void)
      ptecr = ptlec = buffer; /* initialize pointers */
      USART_BRR1 = 0xc9; /* parameter for baud rate */
USART_CR1 = 0x00; /* param. for word length */
USART_CR2 = 0x2c; /* parameters for interrupt */
                               /* authorize interrupts */
      rim();
                                /* loop */
      for (;;)
            outch(getch()); /* get and put a char */
```

#### **Default Compiler Operation**

By default, the compiler compiles and assembles your program. You may then link object files using **clnk** to create an executable program.

The compiler supports several memory models, for applications smaller or larger than 64K, defining how the stack is used and where variables are allocated. A model option should always be specified on the command line; if nothing is specified, the compiler assumes the +modsl option (physical stack and globals in long range). See "Memory Models for code smaller than 64K" in Chapter 3 and "Memory Models for <u>code larger than 64K</u>" in Chapter 3 for more information.

As it processes the command line, cxstm8 echoes the name of each input file to the standard output file (your terminal screen by default). You can change the amount of information the compiler sends to your terminal screen using command line options, as described later.

According to the options you will use, the following files, recognized by the COSMIC naming conventions, will be generated:

file.s Assembler source module file.o Relocatable object module

**file.sm8** input (e.g. libraries) or output (e.g. absolute executable)

file for the linker

# **Compiling and Linking**

To compile and assemble *acia.c* using the short stack model, type:

```
cxstm8 +mods acia.c
```

The compiler writes the name of the input file it processes:

```
acia.c:
```

The result of the compilation process is an object module named *acia.o* produced by the assembler. We will, now, show you how to use the different components.

#### **Step 1: Compiling**

The first step consists in compiling the C source file and producing an assembly language file named **acia.s**.

```
cxstm8 +mods -s acia.c
```

The -s option directs cxstm8 to stop after having produced the assembly file *acia.s.* You can then edit this file with your favourite editor. You can also visualize it with the appropriate system command (*type*, *cat*, *more*,...). For example under MS/DOS you would type:

```
type acia.s
```

If you wish to get an interspersed C and assembly language file, you should type:

```
cxstm8 +mods -1 acia.c
```

The -I option directs the compiler to produce an assembly language file with C source line interspersed in it. Please note that the C source lines are commented in the assembly language file: they start with ';'.

As you use the C compiler, you may find it useful to see the various actions taken by the compiler and to verify the options you selected.

The **-v** option, known as verbose mode, instructs the C compiler to display all of its actions. For example if you type:

```
cxstm8 +mods -v -s acia.c
```

the display will look like something similar to the following:

```
acia.c:
    cpstm8 -o \2.cx1 -i\cxstm8\hstm8 -u -hmods.h acia.c
    cgstm8 -o \2.cx2 \2.cx1
    costm8 -o acia.s \2.cx2
```

The compiler runs each pass:

cpstm8	the C parser
cgstm8	the assembly code generator
costm8	the optimizer

#### Step 2: Assembler

The second step of the compilation is to assemble the code previously produced. The relocatable object file produced is *acia.o*.

```
cxstm8 acia.s
```

or

```
castm8 -i\cxstm8\hstm8 acia.s
```

if you want to use directly the macro cross assembler.

The cross assembler can provide, when necessary, listings, symbol table, cross reference and more. The following command will generate a listing file named *acia.ls* that will also contain a cross reference:

```
castm8 -c -l acia.s
```

For more information, see Chapter 5, "Using The Assembler".

#### Step 3: Linking

This step consists in linking relocatable files, also referred to as object modules, produced by the compiler or by the assembler (**<files>.0**) into an absolute executable file: **acia.sm8** in our example. Code and data sections will be located at absolute memory addresses. The linker is used with a command file (*acia.lkf* in this example).

An application that uses one or more object module(s) may require several sections (code, data, interrupt vectors, etc.,...) located at different addresses. Each object module contains several sections. The compiler creates the following sections:

Туре	Description
.text	code (or program) section (e.g. ROM)
.fconst	large constant and literal data (e.g. ROM, see @far)
.const	constant and literal data (e.g. ROM)
.data	initialized data in <b>long addressing range</b> memory (see @near in chapter 3) (e.g. RAM)
.bss	all non initialized data in <b>long range</b> memory (e.g. <b>RAM</b> )
.bsct	initialized data in the first 256 bytes (see <b>@tiny</b> in chapter 3), also called <b>short range</b> or <b>short addressing range</b> .
.ubsct	non initialized data in the short range
.fdata	large variables (@far)
.eeprom	any variable in eeprom (@eeprom)
.bit	bit variables

In our example, and in the test file provided with the compiler, the *acia.lkf* file contains the following information:

```
line 1 # LINK COMMAND FILE FOR TEST PROGRAM
line 2 # Copyright (c) 2002 by COSMIC Software
line 3 #
line 4 +seg .text -b0xf000 -n.text# program start address
line 5 +seg .const -a .text # constants follow code
```

```
line 6 +seg .bsct -b0x0 -m0x100# short range start address
line 7 +seg .ubsct  # uninitialized short range
line 8 crts.o  # startup routine
line 9 acia.o  # application program
line 10 \cx32\lib\libis.sm8  # C library (if needed)
line 11 \cx32\lib\libm.sm8  # machine library
line 12 +seg .vector -b0x8000 # vectors start address
line 13 vector.o  # interrupt vectors
```

You can create your own link command file by modifying the one provided with the compiler.

Here is the explanation of the lines in *acia.lkf*:

**lines 1 to 3:** These are comment lines. Each line can include comments. They must be prefixed by the "#" character.

line 4: +seg .text -b0xf000 creates a text (code) segment located at f000 (hex address) named .text.

line 5: +seg .const -n.text creates a constant segment following the text segment.

line 6: +seg .bsct -b0x0 -m0x100 creates a short range segment located at 0 (hex address) with a maximum size of 256 bytes.

**line 7:** +seg .ubsct creates an uninitialized short range segment located by default after the .bsct segment.

**line 8:** crts.o runtime startup code. It will be located at 0xf000 (code segment)

**line 9: acia.o**, the file that constitutes your application. It follows the startup routine for code and data

line 10: libis.sm8 the integer library to resolve references

line 11: libm. sm8 the machine library to resolve references

**line 12:** +seg .vector -b0x8000 creates a segment vector (const) segment located at 8000 (hex address)

line 13: vectors .o interrupt vectors file

By default and in our example, the .bss segment follows the .data segment

The *crts.o* file contains the runtime startup that performs the following operations:

- initialize the *bss*, if any
- initialize the stack pointer
- call *main()* or any other chosen entry point.

For more information, see "<u>Modifying the Runtime Startup</u>" in **Chapter 3**.

After you have modified the linker command file, you can link by typing:

```
clnk -o acia.sm8 acia.lkf
```

#### **Step 4: Generating S-Records file**

Although *acia.sm8* is an executable image, it may not be in the correct format to be loaded on your target. Use the **chex** utility to translate the format produced by the linker into standard formats. To translate *acia.sm8* to *Motorola standard S-record* format:

```
chex acia.sm8 > acia.hex
```

or

```
chex -o acia.hex acia.sm8
```

**acia.hex** is now an executable image in *Motorola S-record* format and is ready to be loaded in your target system.

For more information, see "The chex Utility" in Chapter 8.

# **Linking Your Application**

You can create as many *text*, *data* and *bss* segments as your application requires. For example, assume we have one *bss*, two *data* and two *text* segments. Our link command file will look like:

```
+seq .text -b 0xf000 -n .text # program start address
+seg .const -a .text # constant follow
+seq .data -b 0x100
                           # data start address
+seg .bss -b 0x200
                           # bss start address
                          # zpage start address
+seg .bsct -b0x0 -m0x100
+seq .ubsct -n iram
                           # uninitialized short range
                            # startup routine
crts.o
acia.o
                            # main program
module1.o
                            # application program
+seq .text -b 0x2000
                            # start new text section
module2.o
                            # application program
module3.o
                            # application program
\cx\lib\libis.sm8
                           # C library (if needed)
                           # machine library
\cx\lib\libm.sm8
+seq .vector -b0x8000
                           # vectors start address
vector.o
                            # interrupt vectors
# define these symbols if crtsi is used
#+def endzp=@.ubsct
                           # end of uninitialized zpage
#+def memory=@.bss
                            # symbol used by startup
```

In this example the linker will locate and merge crts.o, acia.o and module1.o in a text segment at  $0 \times 100$ , a text segment at  $0 \times 100$ , a text segment at  $0 \times 100$ , a text segment, if needed at  $0 \times 200$ . text segment are text segment at text segment.

For more information about the linker, see **Chapter 6**, "<u>Using The Linker</u>".

# **Generating Automatic Data Initialization**

Usually, in embedded applications, your program must reside in ROM.

This is not an issue when your application contains code and read-only data (such as string or const variables). All you have to do is burn a PROM with the correct values and plug it into your application board.

The problem comes up when your application uses initial data values that you have defined with initialized static data. These static data values must reside in RAM.

There are two types of static data initializations:

1) data that is explicitly initialized to a non-zero value:

```
char var1 = 25;
```

which is generated into the .bsct section and

2) data that is explicitly initialized to zero or left uninitialized:

```
char var2;
```

which is generated into the .ubsct section.

There is one exception to the above rules when you declare data that will be located in the **external memory**, using the **@near** type qualifier. In this case, the data is generated into the **.data** section if it is initialized or in the **.bss** section otherwise.

The first method to ensure that these values are correct consists in adding code in your application that reinitializes them from a copy that you have created and located in ROM, at each restart of the application.

The second method is to use the **crtsi.sm8** start-up file:

- that defines a symbol that will force the linker to create a copy of the initialized RAM in ROM
- 2) and that will do the copy from ROM to RAM

The following link file demonstrates how to achieve automatic data initialization.

```
# demo.lkf: link command with automatic init
+seg .text -b 0xf000 -n .text # program start address
+seg .const -a .text # constant follow
+seg .bsct -b 0x0 -m 0x100 \# zpage start address
+seq .data -b0x100
                          # data start address
\cx32\lib\crtsi.sm8
                          # startup with auto-init
acia.o
                          # main program
module1.o
                          # module program
\cx32\lib\libis.sm8
                         # C library (if needed)
\cx32\lib\libm.sm8
                          # machine library
+seg .vector -b 0x8000  # vectors start address vector.o
# define these symbols if crtsi is used
+def __endzp=@.ubsct # end of uninitialized zpage
+def memory=@.bss
                           # end of bss segment
```

In the above example, the *text* segment is located at address  $0 \times f000$ , the *data* segment is located at address  $0 \times 100$ , immediately followed by the *bss* segment that contains uninitialized data. The initialized data in ROM will follow the descriptor created by the linker after the code segment.

In case of multiple code and data segments, a link command file could be:

```
# demoinit.lkf: link command with automatic init
+seg .text -b 0xf000 -n .text # program start address
+seg .data -b0x100
                         # data start address
\cx32\lib\crtsi.sm8
                         # startup with auto-init
acia.o
                          # main program
module1.o
                         # module program
+seq .text -b0xf800
                         # new code segment
module2.o
                          # module program
module3.o
                         # module program
                         # C library (if needed)
\cx32\lib\libis.sm8
\cx32\lib\libm.sm8
                         # machine library
+seg .vector -b 0x8000
                       # vectors start address
vector.o
                          # interrupt vectors
# define these symbols if crtsi is used
+def endzp=@.ubsct
                         # end of uninitialized zpage
+def memory=@.bss
                        # end of bss segment
```

or

```
# demoinit.lkf: link command with automatic init
+seq .text -b 0xf000 -n .text # program start address
+seg .const -a .text  # constant follow
+seg .bsct -b 0x0 -m 0x100  # zpage start address
+seg .data -b0x100  # data start address
\cx32\lib\crtsi.sm8
                             # startup with auto-init
acia.o
                              # main program
module1.o
                              # module program
+seg .text -b0xf800 -it
                             # sets the section attribute
module2.o
                              # module program
module3.o
                               # module program
\cx32\lib\libis.sm8
                              # C library (if needed)
\cx32\lib\libm.sm8
                              # machine library
+seg .vector -b 0x8000 # vectors start address
vector.o
                               # interrupt vectors
# define these symbols if crtsi is used
+def endzp=@.ubsct # end of uninitialized zpage
+def memory=@.bss
                               # end of bss segment
```

In the first case, the initialized data will be located after the first code segment. In the second case, the -it option instructs the linker to locate the initialized data after the segment marked with this flag. The initialized data will be located after the second code segment located at address 0xf800

For more information, see "Initializing data in RAM" in Chapter 3 and "Automatic Data Initialization" in Chapter 6.

# **Specifying Command Line Options**

You specify command line options to **cxstm8** to control the compilation process.

To compile and get a relocatable file, type:

```
cxstm8 +mods acia.c
```

The file produced is *acia.o.* 

The -v option instructs the compiler driver to echo the name and options of each program it calls. The -l option instructs the compiler driver to create a mixed listing of C code and assembly language code in the file acia.ls.

To perform the operations described above, enter the command:

```
cxstm8 +mods -v -l acia.c
```

When the compiler exits, the following files are left in your current directory:

- the C source file acia.c
- the C and assembly language listing acia.ls
- the object module acia.o

It is possible to locate listings and object files in specified directories if they are different from the current one, by using respectively the -cl and -co options:

```
cxstm8 +mods -cl\mylist -co\myobj -l acia.c
```

This command will compile the *acia.c* file, create a listing named *acia.ls* in the *\mylist* directory and an object file named *acia.o* in the *\myobj* directory.

**cxstm8** allows you to compile more than one file. The input files can be C source files or assembly source files. You can also mix all of these files.

If your application is composed with the following files: two C source files and one assembly source file, you would type:

```
cxstm8 +mods -v start.s acia.c getchar.c
```

This command will assemble the *start.s* file, and compile the two C source files.

See "<u>Compiler Command Line Options</u>" in **Chapter 4** for information on these and other command line options.

CHAPTER 3

# **Programming Environments**

This chapter explains how to use the COSMIC program development system to perform special tasks required by various STM8 applications.

#### Introduction

This chapter provides details about:

- Modifying the Runtime Startup
- Initializing data in RAM
- Memory Models for code smaller than 64K
- Memory Models for code larger than 64K
- Handling Large Code and Constants
- Bit Variables
- The const and volatile Type Qualifiers
- Performing Input/Output in C
- Referencing Absolute Addresses
- Accessing Internal Registers
- Placing Data Objects in The Bss Section
- Placing Data Objects in Short Range Memory
- Placing Data Objects in Long Range Memory
- Placing Data Objects in the EEPROM Space
- Redefining Sections
- Inserting Inline Assembly Instructions
- Writing Interrupt Handlers
- Placing Addresses in Interrupt Vectors
- Inline Function

- Interfacing C to Assembly Language
- Register Usage
- Heap Management Control with the C Compiler
- Data Representation

# **Modifying the Runtime Startup**

The runtime startup module performs many important functions to establish a runtime environment for C. The runtime startup file included with the standard distribution provides the following:

- Initialization of the bss section if any,
- ROM into RAM copy if required,
- Initialization of the stack pointer,
- f main or other program entry point call, and
- An exit sequence to return from the C environment. Most users must modify the exit sequence provided to meet the needs of their specific execution environment.

The following is a listing of the standard runtime startup file **crts.sm8** included on your distribution media. It does not perform automatic data initialization. The runtime startup file can be placed anywhere in memory. Usually, the startup will be "linked" with the **RESET** interrupt, and the startup file may be at any convenient location.

```
; C STARTUP FOR STM8
  ; WITHOUT ANY DATA INITIALIZATION
3
  ; Copyright (c) 2006 by COSMIC Software
4
5
   xref f_main, __stack
    xdef f exit, stext, f stext
7
 ; startup routine from reset vector
10
    switch .text
11 stext:
12 f stext:
13 ;
14; initialize stack pointer
15;
18 ;
19; execute main() function
```

```
20; may be called by a 'jpf' instruction if no return expected
21;
22 callf f_main ; execute main
23 f_exit:
24 jra f_exit ; and stay here
25;
26 end
```

## **Description of Runtime Startup Code**

**f** main is the entry point into the user C program.

Line 15 resets the stack pointer.

Line 22 calls *main()* in the user's C program.

Lines 23 to 24 trap a return from *main()*. If your application must return to a monitor, for example, you must modify this line.

# Initializing data in RAM

If you have initialized static variables, which are located in **RAM**, you need to perform their initialization before you start your C program. The **clnk** linker will take care of that: it moves the initialized data segments after the **first** text segment, or the one you have selected with the **-it** option, and creates a descriptor giving the starting address, destination and size of each segment.

The table thus created and the copy of the **RAM** are located in **ROM** by the linker, and used to do the initialization. An example of how to do this is provided in the **crtsi.s** file, located in the headers sub-directory.

```
C STARTUP FOR STM8
     WITH AUTOMATIC DATA INITIALISATION
     Copyright (c) 2006 by COSMIC Software
     xref f main, memory, idesc , stack
     xref.b c_x, c_y, __endzp
     xdef f exit, stext, f stext
    start address of zpage
:
     switch .ubsct
 suzp:
     start address of bss
     switch .bss
 sbss:
    startup routine from reset vector
     switch.text
 stext:
f stext:
     initialize stack pointer
;
     ldw x,# stack ; stack pointer
     ldw sp,x ; in place
     setup initialized data
;
     ldw y, _idesc__ ; data start address
```

```
ldw
           x,# idesc +2 ; descriptor address
ibcl:
     ld
           a,(x)
                     ; test flag byte
     jreg zero
                     ; no more segment
           a,#$60
                     ; test for moveable code segment
     bcp
     jreq qbcl
                     ; yes, skip it
     ldw
          C X,X
                     ; save pointer
     1dw x, (3,x)
                     ; move end address
     ldw cy,x
                     ; in memory
     ldw x,c x
                     ; reload pointer
     ldw x, (1,x)
                     ; start address
dbc1:
     1d
           a,(y)
                     ; transfer
     14
                     ; byte
          (x),a
     incw x
                     ; increment
     incw y
                     ; pointers
     cpw
           у,с_у
                     ; last byte ?
     jrne dbcl
                     ; no, loop again
     ldw
          x,c_x
                     ; reload pointer
qbcl:
     addw x,#5
                      ; next descriptor
     jra
           ibcl
                      ; and loop
;
     reset uninitialized data in zero page
zero:
     ldw
         x,# suzp ; start of uninitialized zpage
     ira loop
                  ; test segment end first
zbcl:
          (x),a
                     ; clear byte
     incw x
                     ; next byte
loop:
           x, # endzp ; end of zpage
     сри
     jrne zbcl
                    ; no, continue
;
     reset uninitialized data in bss
;
           x, # sbss ; start address
     ldw
     jra
           ok
                     ; test segment end first
bbcl:
     ld
          (x),a
                     ; clear byte
     incw x
                     ; next byte
ok:
     cpw x,# memory; compare end
     jrne bbcl
                 ; not equal, continue
;
```

```
; execute main() function
; may be called by a 'jpf' instruction if no return
expected
;
    callf f_main    ; execute main
f_exit:
    jra f_exit    ; and stay here
;
    end
```

*crtsi.s* performs the same function as described with the *crts.s*, but with one additional step. Lines (marked in bold) in *crtsi.s* include code to copy the contents of initialized static data, which has been placed in the text section by the linker, to the desired location in RAM.

The compiler is provided with several startup files implementing the automatic data initialization depending on the range of variables to be initialized and the range of the initialization table:

Startup	Initialize	From Table in
crtsi(0).s	@near	@near
crtsx(0).s	@near and @far	@near
crtsif(0).s	@near	@far
crtsxf(0).s	@near and @far	@far

#### - NOTE ·

When using a model for application smaller than 64K, you must use the specific startup (name ending with '0').

For more information, see "<u>Generating Automatic Data Initialization</u>" in **Chapter 2** and "<u>Automatic Data Initialization</u>" in **Chapter 6**.

## Memory Models for code smaller than 64K

The STM8 compiler supports two memory models for application smaller than 64K, allowing you to choose the most efficient behavior depending on your processor configuration and your application. All these models handle code smaller than 64K and then function pointers are defaulted to @near pointers (2 bytes). Data pointers are defaulted to @near pointers (2 bytes). The supported models are:

- Stack Short (mods0) global variables are defaulted to *short range*. Any global object in long range will have to be accessed explicitly with the @near modifier unless accessed through a pointer.
- Stack Long (modsl0) global variables are defaulted to long range.
   Any object in short range will have to be accessed explicitly with the atiny modifier.

The choice of the appropriate model for a given application should be driven by the amount of globals compared with the available space in memory. Short Range variables are more efficient than Long Range variables but the Short Range size is limited to 256 bytes.

#### - NOTE

When using a model for application smaller than 64K, you must link with the specific set of libraries and startup (names ending with '0').

# Memory Models for code larger than 64K

The STM8 compiler supports two memory models for application larger than 64K, allowing you to choose the most efficient behavior depending on your processor configuration and your application. All these models allow the code to be larger than 64K and then function pointers are defaulted to <code>@far</code> pointers (3 bytes). Data pointers are defaulted to <code>@near</code> pointers (2 bytes) unless explicitly declared with the <code>@far</code> modifier. The supported models are:

• Stack Short (mods) global variables are defaulted to *short range*. Any global object in long range will have to be accessed explicitly with the @near modifier unless accessed through a pointer.

• Stack Long (modsl) global variables are also defaulted to *long range*. Any object in short range will have to be accessed explicitly with the @tiny modifier.

The choice of the appropriate model for a given application should be driven by the amount of globals compared with the available space in memory. Short Range variables are more efficient than Long Range variables but the Short Range size is limited to 256 bytes.

# **Handling Large Code and Constants**

The STM8 addresses more than 64K of code, dividing the total space in 64K large sections. The compiler allows by default functions to be as large as needed and to be allocated anywhere in the total space, regardless of any sections boundary crossing. The assembly name of functions allocated anywhere, implicitly declared as @far functions, is prefixed with the sequence "f\_", the function being called by a callf instruction and ended with a retf instruction, while the name of functions declared explicitly as @near functions is prefixed with a single "\_", the function being called by a call instruction and ended with a ret instruction. Such @near functions must be completely located inside the same section and called only by functions located in the same section. Unless trying to optimize the code space, the @near modifier should not be used. Such a difference in the function names allows the linker to check that @far and @near functions are called properly, any mixing being forbidden as the stack display would not be the same.

In order to allow large functions to cross section boundaries, the compiler has to use the far jump instruction **jpf** for any unconditional jump inside the function when the range is too large for a **jra** instruction. This makes the code larger even if such a function is not actually crossing a section boundary. It is possible to force the compiler using the **jp** instruction on any function by using the **-gnc** option. Such functions may still be **@far** or **@near** but cannot be larger than 64K and cannot be allocated across section. Such functions and **@near** functions are checked by the linker in order to verify this constraint.

Constants and literals are normally produced in the .const section which must be located in the first section. Large constants may be declared with an @far modifier. In such a case, they are produced in a .fconst section which may be located anywhere. Although the far space is used for code and constants, the compiler allows variables to be declared with the @far modifier. Such variables are allocated in a .fdata section which may be located anywhere.

## **Bit Variables**

The C compiler implements bit variables using the \_Bool type name as defined by the new ANSI/ISO standard C99 (also referenced as C9X). A \_Bool variable is a boolean object which only contains the values true and false. They are implemented on single bits with value 1 for true and 0 for false. When assigning an expression to a \_Bool variable, the compiler compares the expression against zero and assigns the boolean result to the boolean variable. So, any integer, real or pointer expression can be assigned to a boolean variable. It is not possible to declare arrays of booleans nor pointers to booleans, but booleans can be used as structure or union fields. Consecutive \_Bool fields will be packed in bytes.

The compiler packs global \_Bool variables in bytes using a bit section named .bit. Local \_Bool variables are also packed in bytes regardless of the memory model. \_Bool arguments are passed widened to a single byte.

# The const and volatile Type Qualifiers

You can add the type qualifiers **const** and **volatile** to any base type or pointer type attribute.

Volatile types are useful for declaring data objects that appear to be in conventional storage but are actually represented in machine registers with special properties. You use the type qualifier volatile to declare memory mapped input/output control registers, shared data objects, and data objects accessed by signal handlers. The compiler will not optimize references to volatile data.

An expression that stores a value in a data object of *volatile* type stores the value immediately. An expression that accesses a value in a data object of *volatile* type obtains the stored value for each access. Your program will not reuse the value accessed earlier from a data object of *volatile* type.

#### NOTE

The volatile keyword must be used for any data object (variables) that can be modified outside of the normal flow of the function. Without the volatile keyword, all data objects are subject to normal redundant code removal optimizations. Volatile MUST be used for the following conditions:

- 1) All data objects or variables associated with a memory mapped hardware register e.g. volatile char DDRB @0x05;
- 2) All global variable that can be modified (written to) by an interrupt service routine either directly or indirectly. e.g. a global variable used as a counter in an interrupt service routine.

You use *const* to declare data objects whose stored values you do not intend to alter during execution of your program. You can therefore place data objects of *const* type in ROM or in write protected program segments. The cross compiler generates an error message if it encounters an expression that alters the value stored in a *const* data object.

If you declare a static data object of *const* type at either file level or at block level, you may specify its stored value by writing a data initializer. The compiler determines its stored value from its data initializer before program startup, and the stored value continues to exist unchanged until program termination. If you specify no data initializer, the stored value is zero. If you declare a data object of *const* type at argument level, you tell the compiler that your program will not alter the value stored in that argument in the related function. If you declare a data object of *const* type and dynamic lifetime at block level, you must specify its stored value by writing a data initializer. If you specify no data initializer, the stored value is undefined.

The *const* keyword implies the **@near** memory space to allow such a variable to be located in the code space. If a memory space modifier is explicitly given on a declaration using the *const* keyword, the compiler uses the given space instead of the default one, meaning that the object may not be located in the code space depending on the memory space given. In such a case, the *const* keyword still enforces the assignment checking.

You may specify *const* and *volatile* together, in either order. A *const volatile* data object could be a Read-only status register, or a status variable whose value may be set by another program.

Examples of data objects declared with type qualifiers are:

```
char * const x;  /* const pointer to char */
int * volatile y;  /* volatile pointer to int */
const float pi = 355.0 / 113.0; /* pi is never changed */
```

# **Performing Input/Output in C**

You perform input and output in C by using the C library functions *getchar, gets, printf, putchar, puts* and *sprintf*. They are described in chapter 4.

The C source code for these and all other C library functions is included with the distribution, so that you can modify them to meet your specific needs. Note that all input/output performed by C library functions is supported by underlying calls to *getchar* and *putchar*. These two functions provide access to all input/output library functions. The library is built in such a way so that you need only modify *getchar* and *putchar*, the rest of the library is independent of the runtime environment.

Function definitions for *getchar* and *putchar* are:

```
char getchar(void);
char putchar(char c);
```

# **Referencing Absolute Addresses**

This C compiler allows you to read from and write to absolute addresses, and to assign an absolute address to a function entry point or to a data object. You can give a memory location a symbolic name and associated type, and use it as you would do with any C identifier. This feature is usefull for accessing memory mapped I/O ports or for calling functions at known addresses in ROM.

References to absolute addresses have the general form **@<address>**, where <address> is a valid memory location in your environment. For example, to associate an I/O port at address **0x20** with the identifier name *ttystat*, write a definition of the form:

```
char MISCR1 @0x20;
```

where @0x20 indicates an absolute address specification and not a data initializer. Since input/output on the STM8 architecture is memory mapped, performing I/O in this way is equivalent to writing in any given location in memory.

Such a declaration does not reserve any space in memory. The compiler still creates a label, using an *equate* definition, in order to reference the C object symbolically. This symbol is made *public* to allow external usage from any other file.

Individual bits can also be declared as **\_Bool** objects by adding a bit number to the address using the syntax **@**<address>:<bitnum>, where <address> is a byte memory location and <bitnum> a bit number expressed by a constant value (or expression) between 0 and 7. For example, to define bit 3 of memory byte at **0x5001** as *PB3*:

```
_Bool PB3 @0x5001:3;
```

To use the I/O port in your application, write:

```
char c;
c = MISCR1; /* to read from input port */
MISCR1 = c; /* to write to output port */
```

Another solutions is to use a **#define** directive with a cast to the type of the object being accessed, such as:

```
#define MISCR1 *(char *)0x20
```

which is both inelegant and confusing. The COSMIC implementation is more efficient and easier to use, at the cost of a slight loss in portability.

Note that COSMIC C does support the pointer and #define methods of implementing I/O access.

Another example of how to reference a direct memory address, defines a structure at absolute address **0x6000**:

```
struct acia
{
    char status;
    char data;
} acia @0x6000
```

Using this declaration, references to acia.status will refer to memory location 0x6000 and acia.data will refer to memory location 0x6001. This is very useful if you are building your own custom I/O hardware that must reside at some location in the STM8 memory map.

# **Accessing Internal Registers**

All registers are declared in **io\*.h** files provided with the compiler, each one dedicated to a specific processor. One of these files should be included in each file using the input-output registers, for example by a:

#### #include <iostm8.h>

Note that the compiler will access to these registers as standard variables. In some case of reading or writing some "*int*" registers, you should declare an **union** (with two char and one int) instead of using directly the I/O register.

# **Placing Data Objects in The Bss Section**

The compiler automatically reserves space for uninitialized data object. All such data are placed in the .bss section. All initialized static data are placed in the .data section. The bss section is located, by default, after the data section by the linker.

The run-time startup files, *crtsi.s* and *crtsi.s*, contain code which initializes the bss section space to zero.

The compiler provides a special option, **+nobss**, which forces uninitialized data to be explicitly located in the **.data** section. In such a case, these variables are considered as being explicitly initialized to zero.

# **Placing Data Objects in Short Range Memory**

The **Stack Short** model allocates global variables by default in *short* range memory. Such variables are located into the section **.bsct** if they are initialized, or in the section **.ubsct** otherwise. An external object name is published via a **xref.b** declaration at the assembly language level. A variable can be explicitly allocated in *short* range by using the **@tiny** modifier:

#### @tiny char c;

To place data objects into short range memory on a file basis, you can use the **#pragma** directive of the compiler:

#### #pragma space extern [] @tiny

instructs the compiler to place all data objects of storage class **extern** or **static** into short range memory for the current unit of compilation (usually a file). The section must end with a **#pragma space extern** [] **@near** to revert to the default compiler behaviour.

The compiler provides the *ctat* utility, which helps optimizing short range allocation by automatically extracting the list of all variables which can be efficiently moved to the short range area, by producing a header file which must be used to rebuild the application. For more information, see "*The Ctat Utility*" in **Chapter 8**.

#### NOTE

The code generator does not check for short range overflow.

#### **Setting Short Range Size**

You can set the size of the *short range* section of your object image at link time by specifying the following options on the linker command line:

#### +seg .bsct -m##

where ## represents the size of the *short range* section in bytes. Note that the size of the *short range* section can never exceed **256** bytes.

# **Placing Data Objects in Long Range Memory**

The **Stack Long** model allocates global variables by default in *long range* memory. Such variables will be located into the **.data** section if they are initialized, or in the **.bss** section otherwise. An external object name is published via a **xref** declaration at the assembly language level. A variable can be explicitly allocated in *Long Range* memory by using the **@near** modifier. The following declaration:

#### @near char ext;

instructs the compiler to locate the variable *ext* in the long range memory.

To place data objects into long range memory on a file basis, if not yet selected by the memory model, you can use the **#pragma** directive of the compiler:

#### #pragma space extern [] @near

instructs the compiler to place all data objects of storage class **extern** or **static** into *Long Range memory* for the current unit of compilation (usually a file).

The section must end with a **#pragma space extern** [] **@tiny** to revert to the default compiler behaviour.

## Placing Data Objects in the EEPROM Space

The compiler allows the use to define a variable as an **eeprom** location, using the type qualifier **@eeprom**. This causes the compiler to produce special code when such a variable is modified. When the compiler detects a write to an *eeprom* location, it calls a machine library function which performs the actual write. An example of such a definition is:

```
@eeprom char var;
```

To place all data objects from a file into *eeprom*, you can use the **#pragma** directive of the compiler:

```
#pragma space extern [] @eeprom @near
```

instructs the compiler to treat all *extern* and *static* data in the current file as *eeprom* locations. The **@near** modifier is necessary because the eeprom is located outside the short range.

The section must end with a **#pragma space extern** [] @near.

The compiler allocates @eeprom variables in a separate section named .eeprom, which will be located at link time.

The linker directive:

```
seg .eeprom -b0x4000 -m2048
var_eeprom.o
```

will create a segment located at address **0x4000**, with a maximum size of 2048 bytes.

The library routines are optimized to take advantage of the four byte word programming feature wherever possible. The library routines are also using a few external symbols to locate the command registers in the I/O register space. When a standard header file provided with the compiler is included in at least one source file of the application, these symbols are automatically found. Otherwise, it is necessary to provide these symbols either by extra C declarations:

```
volatile char FLASH_IAPSR @0x505f;
volatile char FLASH_CR2 @0x505b;
volatile char FLASH_NCR2 @0x505c;
```

or by +def directives in the linker command file:

```
+def _FLASH_IAPSR=0x505f
+def _FLASH_CR2=0x505b
+def _FLASH_NCR2=0x505c
```

The addresses provided must match the proper register addresses of the target device.

#### NOTE -

The code generator cannot check if data address will be eeprom addresses after linkage.

# **Redefining Sections**

The compiler uses by default predefined sections to output the various component of a C program. The default sections are:

Section	Description
.text	executable code
.const	text string and constants
.fconst	large constant variables (@far)
.data	initialized variables (@near)
.bss	uninitialized variables (@near)
.bsct	initialized variables in short range (@tiny by default)
.ubsct	uninitialized variables in short range (@tiny by default)
.fdata	large variables (@far)
.eeprom	any variable in eeprom (@eeprom)
.bit	bit variables

It is possible to redirect any of these components to any user defined section by using the following pragma definition:

#### #pragma section <attribute> <qualified\_name>

where < attribute > is either **empty** or one of the following sequences:

const
\_Bool
@tiny
@near
@far
@eeprom

and <qualified name > is a section name enclosed as follows:

(name) - parenthesis indicating a code section
 [name] - square brackets indicating uninitialized data
 {name} - curly braces indicating initialized data

A section name is a plain C identifier which *does not* begin with a dot character, and which is no longer than **13** characters. The compiler will prefix automatically the section name with a dot character when passing this information to the assembler. It is possible to switch back to the default sections by omitting the section name in the *qualified\_name* sequence.

Each pragma directive starts redirecting the selected component from the next declarations. Redefining the *bss* section forces the compiler to produce the memory definitions for all the previous *bss* declarations before to switch to the new section.

The following directives:

```
#pragma section (code)
#pragma section const {string}
#pragma section @near [udata]
#pragma section @near {idata}
#pragma section @tiny [uzpage]
#pragma section @tiny {izpage}
#pragma section @eeprom @near {e2prom}
#pragma section _Bool {bdata}
```

redefine the default sections (or the previous one) as following:

- executable code is redirected to section .code
- strings and constants are redirected to section .string
- uninitialized variables are redirected to section .udata
- initialized data are redirected to section .idata
- uninitialized zpage variables are redirected to section .uzpage
- initialized zpage variables are redirected to section .izpage
- eeprom variables are redirected to section .e2prom
- bit variables are redirected to section .bdata

Note that {name} and [name] are equivalent for constant and eeprom sections as they are all considered as initialized.

The following directive:

```
#pragma section ()
```

switches back the code section to the default section .text.

# **Inserting Inline Assembly Instructions**

The compiler features two ways to insert assembly instructions in a C file. The first method uses **#pragma** directives to enclose assembly instructions. The second method uses a special function call to insert assembly instructions. The first one is more convenient for large sequences but does not provide any connection with C object. The second one is more convenient to interface with C objects but is more limited regarding the code length.

## **Inlining with pragmas**

The compiler accepts the following pragma sequences to start and finish assembly instruction blocks:

Directive	Description
#pragma asm	start assembler block
#pragma endasm	end assembler block

The compiler also accepts shorter sequences with the same meaning:

Directive	Description
#asm	start assembler block
#endasm	end assembler block

Such an assembler block may be located anywhere, inside or outside a function. Outside a function, it behaves syntactically as a declaration. This means that such an assembler block cannot split a C declaration somewhere in the middle. Inside a function, it behaves syntactically as one C instruction. This means that there is no trailing semicolon at the end, and no need for enclosing braces. It also means that such an assembler block cannot split a C instruction or expression somewhere in the middle.

The following example shows a correct syntax:

## Inlining with asm

The **\_asm()** function inserts inline assembly code in your C program. The syntax is:

```
_asm("string constant", arguments...);
```

The "string constant" argument is the assembly code you want embedded in your C program. "arguments" follow the standard C rules for passing arguments. The string you specify follows standard C rules. For example, carriage returns can be denoted by the '\n' character.

For example, to produce the following assembly sequence:

```
ldw x,sp
callf f_main
```

you would write

```
_asm("ldw x,sp\n callf f_main\n");
```

The '\n' character is used to separate the instructions when writing multiple instructions in the same line.

#### NOTE

The argument string **must** be shorter than 255 characters. If you wish to insert longer assembly code strings you will have to split your input among consecutive calls to asm().

\_asm() does not perform any checks on its argument string. Only the assembler can detect errors in code passed as argument to an \_asm() call.

\_asm() can be used in expressions, if the code produced by \_asm complies with the rules for function returns. For example:

allows to rotate the variable *var* passed as argument in the a register, and store the result in the same variable. The variable *var* is supposed to be declared as a *char*, and is loaded in the a register because it is considered as a first argument. The result is expected in the a register in order to comply with the return register convention, as described below.

#### NOTE -

With both methods, the assembler source is added as is to the code during the compilation. The optimizer does not modify the specified instructions, unless the -a option is specified on the code generator. The assembler input can use lowercase or uppercase mnemonics, and may include assembler comments.

By default, \_asm() is returning an **int** as any undeclared function. To avoid the need of several definitions (usually conflictuous) when \_asm() is used with different return types, the compiler implements a special behaviour when a cast is applied to \_asm(). In such a case, the cast is considered to define the return type of \_asm() instead of asking for a type conversion. There is no need for any prototype for the \_asm() function as the parser verifies that the first argument is a string constant.

#### **Inlining Labels**

When labels are necessary in the inlined assembly code, the compiler provides a special syntax allowing local labels to be created and handled without interaction with other labels and the optimizer. The sequence \$N in the assembly source is replaced by a new label name while the sequence \$L is replaced by the label name created by the last \$N. Using this syntax, a simple wait loop may be entered as follow:

```
#asm
    ld a,#7
$N:
    dec a
    jrne $L ; loop on the previous label
#endasm
```

## **Writing Interrupt Handlers**

A function declared with the type qualifier **@interrupt** is suitable for direct connection to an interrupt (hardware or software). **@interrupt** functions may not return a value. **@interrupt** functions are allowed to have arguments, although hardware generated interrupts are not likely to supply anything meaningful.

When you define an @interrupt function, the compiler uses the "iret" instruction for the return sequence, and saves, if necessary, the memory bytes used by the compiler for its internal usage. Those areas are c\_x (3 bytes), c\_y (3 bytes) and c\_lreg (4 bytes). Those bytes will be saved and restored if the interrupt function uses them directly. If the interrupt function does not uses these areas directly, but calls another C function, the c\_x and c\_y areas will be automatically saved and restored, unless using the type qualifier @nosvf on the interrupt function definition. This qualifier can be used when the called functions are known not using those areas, but the compiler does not perform any verification. The c\_lreg area is not saved implicitly in such a case, in order to keep the interrupt function as efficient as possible. If any function called by the interrupt function uses longs or floats, the c\_lreg area can be saved by using the type qualifier @svlreg on the interrupt function definition. Whatever the model used is, these copies are made directly on the stack.

You define an @interrupt function by using the type qualifier @interrupt to qualify the type returned by the function you declare. An example of such a definition is:

```
@interrupt void it_handler(void)
{
    ...
}
```

#### NOTE

The @interrupt modifier is an extension to the ANSI standard.

## **Placing Addresses in Interrupt Vectors**

The compiler implements the special modifier **@vector** in order to allow the interrupt table to be declared directly in C. Each entry must be declared as a function pointer and initialized with a function name. The compiler will produce instead of an address constant, a 4 byte entry containing the special opcode **0x82** followed by the 24 bit address of the interrupt function to be reached.

Refer to the *vector.c* file provided with the compiler example for a more accurate implementation description.

A small C construct would be:

```
extern void handler1(), handler2(), handler3();
void (* const @vector vectab[])() =
     {
     handler1,
     handler2,
     handler3,
     };
```

where *handler1* and so forth are interrupt handlers, which can be located anywhere in the code space, meaning **@near** or **@far** functions. Then, in the linker command file, include the following options on the directive line:

```
+seg .const -b0x8000 vector.o
```

where *vector.o* is the file which contains the vector table. This file is provided in the compiler package.

## **Inline Function**

The compiler is able to inline a function body instead of producing a function call. This feature allows the program to run faster but produces a larger code. A function to be inlined has to be defined with the **@inline** modifier. Such a function is kept by the compiler and does not produced any code yet. Each time this function is called in the same source file, the call is replaced by the full body of the inlined function. Because inlined functions are in fact local to a source file, they should be defined in a header file if they have to be used by several source files. To allow the arguments to be passed properly, inlined functions must be defined with prototypes.

#### - NOTE

Inline functions cannot declare static local variables and cannot call themselves either directly or indirectly.

The compiler allows access to specific instructions or features of the STM8 processor, using **@inline** functions. Such functions shall be declared as external functions with the **@inline** modifier. The compiler recognizes three predefined functions when explicitly declared as follows:

```
@inline char carry(void);
@inline char irq(void);
@inline char imask(void);
```

carry	the <i>carry</i> function is used to test or get the carry bit from the condition register. If the <i>carry</i> function is used in a test, the compiler produces a <i>jnrc</i> or <i>jrc</i> instruction. If the <i>carry</i> function is used in any other expression, the compiler produces a code sequence setting the a register to 0 or 1 depending on the carry bit value.
irq	the <i>irq</i> function is used to test the interrupt line level using the <b>jrih</b> or <b>jril</b> instruction. The <i>irq</i> function can be used only in a test
imask	the <i>imask</i> function is used to test the interrupt mask bit in the condition register using the <b>jrm</b> or <b>jrnm</b> instruction. The <i>imask</i> function can be used only in a test.

These functions are predeclared in the *processor.h* header file. A full description with examples is provided in Chapter 4.

Any other function declared as an *@inline* will be translated into a call to a user provided macro. The macro name is obtained by prefixing the *@inline* function name with the '\_' character. Arguments are allowed but should be restricted to variable references. Each reference is translated into the proper assembler expression (same translation as applied by the compiler) and then passed to the macro as a quoted text string.

*@inline* functions may use the registers **a** and/or **x**, but the compiler can not check their use and will not save them. To save the registers before they are used by *@inline* functions, you must add the *@usea* and/or the *@usex* modifiers.

For example:

```
@inline @usea lsub();
```

tells the compiler that *lsub()* uses the register **a**, so that the compiler will save it. If both registers are used, you must specify both modifiers.

## Interfacing C to Assembly Language

The C cross compiler translates C programs into assembly language according to the specifications described in this section.

You may write external identifiers in both uppercase and lowercase. The compiler prepends an underscore '\_' character to each identifier. If the identifier is the name of an **@far** function, the compiler prepends a 'f' character to the extra underscore.

The compiler places function code in the .text section. Function code is not to be altered or read as data. External function names are published via xdef declarations.

Compiler literals, containing text strings, float and long constants, and switch tables, are generated by default into the .const section. An option on the code generator allows such constants to be produced in the .text section. In any case, these literals must be linked in the @near space.

The compiler generates initialized data declared with the **@near** modifier into the **.data** section. Such external data names are published via **xref** declarations. Data you declare to be of "const" type by adding the type qualifier *const* to its base type is generated by default into the **.const** section. Const variables declared with the **@far** modifier are published in the **.fconst** section. Initialized data declared with the **@tiny** space modifier will be generated into the **.bsct** section. Such external data names are published via **xref.b** declarations. Uninitialized data are normally generated into the **.bss** section for **@near** variables or the **.ubsct** section for **@tiny** variables, unless forced to the **.data** or **.bsct** section by the compiler option **+nobss**. Variables declared with the **@far** modifier are published in the **.fdata** section. **\_Bool** data is generated in the **.bit** section and external names are published via **xbit.b** declarations.

Section	Declaration	Reference
.bsct	@tiny char i =2;	xdef
.ubsct	@tiny char i;	xdef

Section	Declaration	Reference
.data	int init = 1	xdef
.bss	int uninit	xdef
.text	char putchar(c);	xdef
.const	const char c = 1;	xdef
.fconst	@far const char c = 1;	xdef
.fdata	@far char i;	xdef
.bit	_Bool Pb3;	xdef
Any of above	extern int out;	xref, xbit

Function calls are performed according to the following:

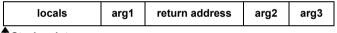
- 1) Arguments are evaluated from right to left. The first argument is stored in the a register if it is a char or a @tiny pointer, or in the x register if its type is short, int or @near pointer, or the x register and the memory location c\_x if its type is @far pointer, and if the function does not return a structure larger than 2 bytes. Other arguments are pushed to the stack. If the two first arguments are of type char or unsigned char, they are both stored in registers, the first in the xh register, and the second in the xl register as if they were considered as an int.
- 2) The function is called via a *callf f func* instruction.
- 3) Stacked arguments are cleaned out.

## **Register Usage**

Except for the return value, the registers  $\mathbf{a}$ ,  $\mathbf{x}$ ,  $\mathbf{y}$  and the condition codes are undefined on return from a function call. The return value is in  $\mathbf{a}$  if it is of type char,  $\mathbf{@tiny}$  pointer or a one byte structure,  $\mathbf{x}$  if it is of type short, integer,  $\mathbf{@near}$  pointer or a two byte structure. The return value is in the memory located at symbol  $\mathbf{c}$ \_lreg if it is of type long or float. The return value is in the memory location  $\mathbf{c}$ \_ $\mathbf{x}$  (upper byte) and the  $\mathbf{x}$  register (lower word) if it is of type  $\mathbf{@far}$  pointer.

The first argument may be hold in register, and will be stored at the function entry. Such a function declaration:

will create the following memory area:



TStack pointer

## **Heap Management Control with the C Compiler**

The name **heap** designates a memory area in which are allocated and deallocated memory blocks for temporary usage. A memory block is allocated with the *malloc()* function, and is released with the *free()* function. The *malloc()* function returns a pointer to the allocated area which can be used until it is released by the *free()* function. Note that the *free()* function has to be called with the pointer returned by *malloc*. The *heap* allocation differs from a local variable allocation because its life is not limited to the life of the function performing the allocation.

In an embedded application, the *malloc-free* mechanism is available and automatically set up by the compiler environment and the library. But it is possible to control externally the *heap* size and location. The default compiler behaviour is to create a data area containing application variables and *heap* in the following way:

initialized variables (data segment)	uninitialized variables (bss segment)	heap growing upward until upper limit
	<i>heap</i> starts here	heap ends here

The heap start is the *bss* end, and is equal to the \_\_**startmem** symbol defined by the linker with an appropriate +**def** directive. The heap end is equal to the \_\_**endmem** symbol defined by the linker with an appropriate +**def** directive. A typical heap reservation may be implemented in the linker command file by using the +**spc** directive:

```
+def __startmem=@.bss # heap start symbol
+spc .bss=0x400 # reserve 1K space
+def endmem=@.bss # heap end
```

The *heap* management functions maintain a global pointer named *heap pointer*, or simply **HP**, pointing to the *heap top*, and a linked list of memory blocks, free or allocated, in the area between the *heap start* and the *heap top*. In order to be able to easily modify the heap implementation, the heap management functions use a dedicated function to move the heap pointer whenever necessary. The heap pointer is initialized to the heap start: the heap is initially empty. When *malloc* needs some memory and no space is available in the free list, it calls this dedicated function named \_sbreak to move the heap pointer upwards if possible.

\_sbreak will return a NULL pointer if this move is not possible (usually this is because the heap would overlap the stack). Therefore it is possible to change the heap default location by rewriting the \_sbreak function.

The default *sbreak* function provided by the library is as follows:

If the new top is outside the authorized limits defined by the address of the two objects **\_startmem** and **\_endmem**, the function returns a NULL pointer, otherwise, it returns the start of the new allocated area. Note that the top variable **\_brk** is a static variable initialized to zero (NULL pointer). It is set to the heap start on the first call. It is also possible to initialize it directly within the declaration, but in this case, we create an initialized variable in the **data** segment which needs to be initialized by the startup. The current code avoids such a requirement by initializing the variable to zero (in the **bss** segment), which is simply done by the standard startup sequence.

## **Modifying The Heap Location**

It is easy to modify the *\_sbreak* function in order to handle the heap in a standard C array, which will be part of the application variables.

The heap area is declared as an array of **char** simply named *heap*. The algorithm is mainly the same, and once the new top is computed, it is compared with the array limits. Note that the array is declared as a static local variable. It is possible to have it declared as a static global variable. If you want it to be global, be careful on the selected name. You should start it with a '\_' character to avoid any conflict with the application variables.

The modified *sbreak* function using an array is as follows:

If you need to place the heap array at a specific location, you need to locate this module at a specific address using the linker options. In the above example, the heap array will be located in the .bss segment, thus, complicating the startup code which would need to zero two bss sections instead of one. Compiling this function, with the +nobss option, will force allocation of the heap, in the data segment and you can locate it easily with linker directives as:

```
+seg .data -b 0x2000  # heap start
sbreak.o  # sbreak function
```

It is also possible to handle the heap area outside of any C object, just by redefining the heap start and end values using the linker **+def** directives:

```
+def __startmem=0x2000 # heap start
+def __endmem=0x3000 # heap end
```

#### - NOTE -

Since the initial content of the area may be undefined, the **-ib** option should be specified to exclude the segment in the automatic RAM initialization.

Note that it is possible to use this \_sbreak function as a malloc equivalent function with some restrictions. The malloc function should be used when the allocated memory has to be released, or if the application has no idea about the total amount of space needed. If memory can be allocated and never released, the free mechanism is not necessary, nor the linked list of memory blocks built by malloc. In that case, simply rename the \_sbreak function as malloc, regardless of its implementation, and you will get a very efficient and compact malloc mechanism. You may do the renaming in the function itself, which needs to be recompiled, or by using a #define at C level, or by renaming the function at link time with a +def directive such as:

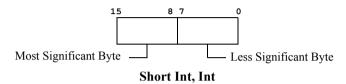
```
+pri  # enter a private region
+def _malloc=__sbreak  # defines malloc as _sbreak
+new  # close region and forget malloc
libi.sm8  # load library containing sbreak
```

This sequence has to be placed just before loading libraries, or before placing the module containing the *\_sbreak* function. The private region is used to forget the *\_malloc* reference once it has been aliased to *sbreak*.

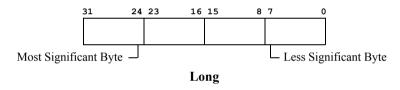
## **Data Representation**

Data objects of type *char* are stored as one byte. A plain *char* is defaulted to type *unsigned char*.

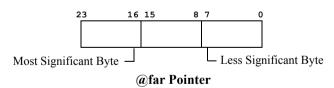
Data objects of type *short int* and *int* are stored as two bytes, more significant byte first.



Data objects of type *long int* are stored as four bytes, in descending order of significance.

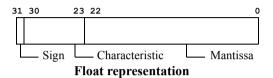


*@tiny* pointers (short range) are stored as one byte. *@near* pointers (long range) are stored as two bytes. *@far* pointers are stored as three bytes, in descending order of significance.



Data objects of type *float* are represented as for the proposed IEEE Floating Point Standard; four bytes stored in descending order of significance. The IEEE representation is: most significant bit is one for negative numbers, and zero otherwise; the next eight bits are the characteristic, biased such that the binary exponent of the number is the characteristic minus 126; the remaining bits are the fraction, starting with the weighted bit. If the characteristic is zero, the entire number is

taken as zero, and should be all zeros to avoid confusing some routines that do not process the entire number. Otherwise there is an assumed 0.5 (assertion of the weighted bit) added to all fractions to put them in the interval [0.5, 1.0). The value of the number is the fraction, multiplied by -1 if the sign bit is set, multiplied by 2 raised to the exponent.



**CHAPTER** 

4

# **Using The Compiler**

This chapter explains how to use the C cross compiler to compile programs on your host system. It explains how to invoke the compiler, and describes its options. It also describes the functions which constitute the C library. This chapter includes the following sections:

- Invoking the Compiler
- File Naming Conventions
- Generating Listings
- Generating an Error File
- · C Library Support
- Descriptions of C Library Functions

## **Invoking the Compiler**

To invoke the cross compiler, type the command **cxstm8**, followed by the compiler options and the name(s) of the file(s) you want to compile. All the valid compiler options are described in this chapter. Commands to compile source files have the form:

#### cxstm8 [options] <files>.[c|s]

**cxstm8** is the name of the *compiler*. The option list is optional. You must include the name of at least one input file <*file*>. <*file*> can be a C source file with the suffix '.c', or an assembly language source file with the suffix '.s'. You may specify multiple input files with any combination of these suffixes in any order.

If you do not specify any command line options, **cxstm8** will compile your *<files>* with the default options. It will also write the name of each file as it is processed. It writes any error messages to STDERR.

The following command line:

#### cxstm8 +mods acia.c

compiles and assembles the *acia.c* C program, using the *Stack Short* model, creating the relocatable program **acia.o**.

If the compiler finds an error in your program, it halts compilation. When an error occurs, the compiler sends an error message to your terminal screen unless the option -e has been specified on the command line. In this case, all error messages are written to a file whose name is obtained by replacing the suffix .c of the source file by the suffix .err. An error message is still output on the terminal screen to indicate that errors have been found. Appendix A, "Compiler Error Messages" lists the error messages the compiler generates. If one or more command line arguments are invalid, cxstm8 processes the next file name on the command line and begins the compilation process again.

The example command above does not specify any compiler options. In this case, the compiler will use only default options to compile and assemble your program. You can change the operation of the compiler by specifying the options you want when you run the compiler.

To specify options to the compiler, type the appropriate option or options on the command line as shown in the first example above. Options should be separated with spaces. You must include the '-' or '+' that is part of the option name.

## **Compiler Command Line Options**

The **cxstm8** compiler accepts the following command line options, each of which is described in detail below:

```
cxstm8 [options] <files>
     -a*> assembler options
     -ce* path for errors
     -cl* path for listings
     -co* path for objects
     -d*> define symbol
         create error file
     -e
     -ec all C files
     -es all assembler files
     -ex* prefix executables
     -f* configuration file
     -g*> code generator options
     -i*> path for include
     -1
         create listing
     -no do not use optimizer
     -o*> optimizer options
     -p*> parser options
     -s create only assembler file
     -sm create only dependencies
     -sp create only preprocessor file
          path for temporary files
     -t*
         verbose
     -v
          do not execute
     -x
     +*> select compiler options
```

## **Cxstm8 Option Usage**

Option	Description	
-a*>	specify assembler options. Up to 60 options can be specified on the same command line. See <b>Chapter 5</b> , " <i>Using The Assembler</i> ", for the list of all accepted options.	
-ce*	specify a path for the error files. By default, errors are created in the same directory than the source files.	
-cl*	specify a path for the listing files. By default, listings are created in the same directory than the source files.	
-co*	specify a path for the object files. By default, objects are created in the same directory than the source files.	
-d*^	specify * as the name of a user-defined preprocessor symbol (#define). The form of the definition is -dsymbol[=value]; the symbol is set to 1 if value is omitted. You can specify up to 60 such definitions.	
-е	log errors from parser in a file instead of displaying them on the terminal screen. The error file name is defaulted to <file>.err, and is created only if there are errors.</file>	
-ec	treat all files as C source files.	
-es	treat all files as assembler source files.	
-ex	use the compiler driver's path as prefix to quickly locate the executable passes. Default is to use the path variable environment. This method is faster than the default behavior but reduces the command line length.	
-f*	specify * as the name of a configuration file. This file contains a list of options, which will be automatically used by the compiler. If no file name is specified, then the compiler looks for a default configuration file named <code>cxstm8.cxf</code> in the compiler directory as specified in the installation process. For more information, see <code>Appendix B</code> , "The Configuration File".	
-g*>	specify code generation options. Up to 60 options can be specified. See " <i>The cgstm8 Code Generator</i> " in <b>Appendix D</b> , for the list of all accepted options.	

## **Cxstm8 Option Usage (cont.)**

Option	Description
-i*>	define include path. You can define up to 128 different paths. Each path is a directory name, <b>not</b> terminated by any directory separator character, or a file containing a list of directory names.
-1	merge C source listing with assembly language code; listing output defaults to <i><file>.ls</file></i> .
-no	do not use the optimizer.
-0*>	specify optimizer options. Up to 60 options can be specified. See "The costm8 Assembly Language Optimizer" in Appendix D, for the list of all accepted options.
-p*>	specify parser options. Up to 60 options can be specified. See " <i>The cpstm8 Parser</i> " in <b>Appendix D</b> , for the list of all accepted options.
-S	create only assembler files and stop. Do not assemble the files produced.
-sm	create only a list of 'make' compatible dependencies consisting for each source file in the object name followed by a list of header files needed to compile that file.
-sp	create only preprocessed files and stop. Do not compile files produced. Preprocessed output defaults to <file>.p. The produced files can be compiled as C source files.</file>
-t*	specify path for temporary files. The path is a directory name, <b>not</b> terminated by any directory separator character.
-v	be "verbose". Before executing a command, print the command, along with its arguments, to STDOUT. The default is to output only the names of each file processed. Each name is followed by a colon and newline.
-X	do not execute the passes, instead write to STDOUT the commands which otherwise would have been performed.

## **Cxstm8 Option Usage (cont.)**

Option	Description
+*>	select a predefined compiler option. These options are pre- defined in the configuration file. You can specify up to 60 compiler options on the command line. The following docu- ments the available options as provided by the default con- figuration file
+compact	produce a smaller code but slower than the default behaviour. Smaller code is produced by enabling the optimizer factorization feature with a default depth of seven instructions.
+debug	produce debug information to be used by the debug utilities provided with the compiler and by any external debugger.
+fast	produce faster code by inlining machine library calls for long integers handling and integer switches. The code produced will be larger than without this option. For more information, see " <i>Inline Function</i> " in <b>Chapter 3</b> .
+mods	select the <b>Stack Short</b> model. Variables are in <i>short range</i> memory but pointers are pointing to <i>long range</i> memory. See " <u>Memory Models for code larger than 64K</u> " in <b>Chapter 3</b> .
+modsl	select the <b>Stack Long</b> mode. Variables and pointers are in and pointing to <i>long range</i> memory. See " <u>Memory Models</u> <u>for code larger than 64K</u> " in <b>Chapter 3</b> .
+mods0	select the <b>Stack Short</b> model for application smaller than 64K. Variables are in <i>short range</i> memory but pointers are pointing to <i>long range</i> memory. See " <u>Memory Models for code smaller than 64K</u> " in <b>Chapter 3</b> .
+modsl0	select the <b>Stack Long</b> model for application smaller than 64K. Variables and pointers are in and pointing to <i>long range</i> memory. See " <u>Memory Models for code smaller than 64K"</u> in <b>Chapter 3</b> .
+nobss	do not use the .bss section for variables allocated in external memory. By default, such uninitialized variables are defined into the .bss section. This option is useful to force all variables to be grouped into a single section.

#### **Cxstm8 Option Usage (cont.)**

Option	Description
+nocst	output literals and constants in the code section .text instead of the specific section .const.
+proto	enforce prototype declaration for functions. An error message is issued if a function is used and no prototype declaration is found for it. By default, the compiler accepts both syntaxes without any error.
+rev	reverse the bitfield filling order. By default, bitfields are filled from the Less Significant Bit (LSB) towards the Most Significant Bit (MSB) of a memory cell. If the <b>+rev</b> option is specified, bitfields are filled from the <i>msb</i> to the <i>lsb</i> .
+split	create a separate sub-section per function, up to a maximum number of 256 sections, thus allowing the linker to suppress unused functions if the <b>-k</b> option has been specified on at least one segment in the linker command file. For objects with more than 256 functions, the functions will be grouped together to a minimum number of functions per sub-section to not exceed the maximum number of 256 sub-sections. See "Segment Control Options" in Chapter 6.
+strict	direct the compiler to enforce stronger type checking. For more information, see " <i>Extra verifications</i> " in <b>Appendix D</b> .
+warn	enable warnings.

Note that some compiler options are specified to a default value in the configuration file, <code>cxstm8.cxf</code>, and cannot be modified without modiying the configuration file itself (see **Appendix B**, "<u>Modifying Compiler Operation</u>"). These options are -pu (char is unsigned by default) and -ppb (pack local bit variables). Take special care if you remove the -pu option: as libraries expect unsigned chars as argument, it is now the user responsibility to manually declare the relevant variables as unsigned char. If you forget this, run time errors may occur.

## **File Naming Conventions**

The programs making up the C cross compiler generate the following output file names, by default. See the documentation on a specific program for information about how to change the default file names accepted as input or generated as output.

Program	Input File Name	Output File Name
cpstm8	<file>.c</file>	<file>.1</file>
cgstm8	<file>.1</file>	<file>.2</file>
costm8	<file>.2</file>	<file>.s</file>
error listing	<file>.c</file>	<file>.err</file>
assembler listing	<file>.[c s]</file>	<file>.ls</file>
C header files	<file>.h</file>	

castm8	<file>.s</file>	<file>.o</file>
source listing	<file>.s</file>	<file>.ls</file>

cink	<file>.o</file>	name required
------	-----------------	---------------

chex	<file></file>	STDOUT
clabs	<file.sm8></file.sm8>	<files>.la</files>
clib	<file></file>	name required
cobj	<file></file>	STDOUT
cvdwarf	<file.sm8></file.sm8>	<file>.elf</file>

## **Generating Listings**

You can generate listings of the output of any (or all) the compiler passes by specifying the -l option to **cxstm8**. You can locate the listing file in a different directory by using the -cl option.

The example program provided in the package shows the listing produced by compiling the C source file *acia.c* with the **-l** option:

```
cxstm8 +mods -1 acia.c
```

## **Generating an Error File**

You can generate a file containing all the error messages output by the parser by specifying the **-e** option to **cxstm8**. You can locate the error file in a different directory by using the **-ce** option. For example, you would type:

```
cxstm8 +mods -e prog.c
```

The error file name is obtained from the source filename by replacing the .c suffix by the .err suffix.

## **Return Status**

**cxstm8** returns success if it can process all files successfully. It prints a message to STDERR and returns failure if there are errors in at least one processed file.

## **Examples**

To echo the names of each program that the compiler runs:

```
cxstm8 +mods -v file.c
```

To save the intermediate files created by the code generator and halt before the assembler:

```
cxstm8 +mods -s file.c
```

## **C Library Support**

This section describes the facilities provided by the C library. The C cross compiler for STM8 includes all useful functions for programmers writing applications for ROM-based systems.

## **How C Library Functions are Packaged**

The functions in the C library are packaged in three separate sub-libraries; one for machine-dependent routines (the *machine* library), one that does not support floating point (the *integer* library) and one that provides full floating point support (the *floating point* library). If your application does not perform floating point calculations, you can decrease its size and increase its runtime efficiency by including only the integer library.

## **Inserting Assembler Code Directly**

Assembler instructions can be quoted directly into C source files, and entered unchanged into the output assembly stream, by use of the \_asm() function. This function is not part of any library as it is recognized by the compiler itself. See "<u>Inserting Inline Assembly Instructions</u>" in **Chapter 3**.

## **Linking Libraries with Your Program**

If your application requires floating point support, you must specify the floating point library **before** the integer library in the linker command file. Modules common to both libraries will therefore be loaded from the floating point library, followed by the appropriate modules from the floating point and integer libraries, in that order.

#### NOTE

When using a model for application smaller than 64K, you must link with the specific set of libraries (names ending with '0').

## **Integer Library Functions**

The following table lists the C library functions in the integer library.

_asm	irq	malloc	strcat
abs	isalnum	memchr	strchr
atoi	isalpha	memcmp	strcmp
atol	iscntrl	memcpy	strcpy

calloc	isdigit	memmove	strcspn
checksum	isgraph	memset	strlen
checksum16	islower	printf	strncat
checksum16x	isprint	putchar	strncmp
checksumx	ispunct	puts	strncpy
div	isqrt	rand	strpbrk
еерсру	isspace	realloc	strrchr
eepset	isupper	sbreak	strspn
free	isxdigit	scanf	strstr
getchar	labs	sprintf	strtol
gets	ldiv	srand	tolower
imask	lsqrt	sscanf	toupper

## **Floating Point Library Functions**

The following table lists the C library functions in the float library.

acos	cosh	log	sprintf
asin	exp	log10	sqrt
atan	fabs	modf	strtod
atan2	floor	pow	tan
atof	fmod	printf	tanh
ceil	frexp	sin	
cos	ldexp	sinh	

## **Common Input/Output Functions**

Two of the functions that perform stream output are included in both the integer and floating point libraries. The functionalities of the versions in the integer library are a subset of the functionalities of their floating point counterparts. The versions in the integer library cannot print or manipulate floating point numbers. These functions are: *printf*, *sprintf*.

## **Functions Implemented as Macros**

Two of the functions in the C library are actually implemented as "macros". Unlike other functions, which (if they do not return *int*) are declared in header files and defined in a separate object module that is linked in with your program later, functions implemented as macros are defined using **#define** preprocessor directives in the header file that declares them. Macros can therefore be used independently of any library by including the header file that defines and declares them with your program, as explained below. The functions in the C library that are implemented as macros are: *max* and *min*.

## **Including Header Files**

If your application calls a C library function, you must include the header file that declares the function at compile time, in order to use the proper return type and the proper function prototyping, so that all the expected arguments are properly evaluated. You do this by writing a preprocessor directive of the form:

#### #include <header\_name>

in your program, where < header\_name > is the name of the appropriate header file enclosed in angle brackets. The required header file should be included before you refer to any function that it declares.

The names of the header files packaged with the C library and the functions declared in each header are listed below.

<assert.h> - Header file for the assertion macro: assert.

<ctype.h> - Header file for the character functions: *isalnum*, *isalpha*, *iscntrl*, *isgraph*, *isprint*, *ispunct*, *isspace*, *isxdigit*, *isdigit*, *isupper*, *islower*, *tolower* and *toupper*.

<float.h> - Header file for limit constants for floating point values.

<io\*.h> - Header files for input-output registers. Each register has an upper-case name which matches the data sheet definition. The compiler provides a large set of header files for most derivative processors.

Header file for limit constants of the compiler.

<math.h> - Header file for mathematical functions: acos, asin, atan, atan2, ceil, cos, cosh, exp, fabs, floor, fmod, frexp, ldexp, log, log10, modf, pow, sin, sinh, sqrt, tan and tanh.

cessor.h> - Header file for inline functions: carry, irq, imask.

<stdbool.h> - Header file for type bool and values true, false.

<stddef.h> - Header file for types: size\_t, wchar\_t and ptrdiff\_t.

<stdio.h> - Header file for stream input/output: getchar, gets, printf, putchar, puts and sprintf.

<stdlib.h> - Header file for general utilities: abs, abort, atof, atoi, atol, div, exit, labs, ldiv, rand, srand, strtod, strtol and strtoul.

<string.h> - Header file for string functions: memchr, memcmp, memcpy, memmove, memset, strcat, strchr, strcmp, strcpy, strcspn, strlen, strncat, strncmp, strncpy, strpbrk, strrchr, strspn and strstr.

**Functions returning int** - C library functions that return *int* and can therefore be called without any header file are: *isalnum, isalpha, iscntrl, isgraph, isprint, ispunct, isspace, isxdigit, isdigit, isupper, islower, sbreak, tolower* and *toupper*.

## **Descriptions of C Library Functions**

The following pages describe each of the functions in the C library in quick reference format. The descriptions are in alphabetical order by function name.

The *syntax* field describes the function prototype with the return type and the expected arguments, and if any, the header file name where this function has been declared

## asm

## **Description**

Generate inline assembly code

## **Syntax**

```
_asm("string constant", arguments...)
```

#### **Function**

\_asm() generates inline assembly code by copying <string constant> and quoting it into the output assembly code stream. <arguments> are first evaluated following the usual rules for passing arguments. The first argument is kept in the a register or the x:a register pair whenever possible, and all other arguments are pushed onto the stack. After the <string constant> code is output, arguments pushed to the stack are removed before to continue

#### Return Value

Nothing, unless \_asm() is used in an expression. In that case, standard return conventions must be followed. See "<u>Register Usage</u>" in **Chapter 3**.

## **Example**

The sequence inc x; call main, may be generated by the following call:

```
_{asm("inc x n call _main");}
```

Note that the string-quoting syntax matches the familiar *printf()* function.

#### Notes

\_asm() is not packaged in any library. It is recognized by the compiler itself.

For more information, see "<u>Inserting Inline Assembly Instructions</u>" in **Chapter 3**.

## abort

## **Description**

Abort program execution

## **Syntax**

```
#include <stdlib.h>
void abort(void)
```

#### Function

**abort** stops the program execution by calling the *exit* function which is placed by the startup module just after the call to the main function.

#### Return Value

abort never returns.

## **Example**

To abort in case of error:

```
if (fatal_error)
    abort();
```

### See Also

exit

#### Notes

abort is a macro equivalent to the function name exit.

## abs

## **Description**Find absolute value

## **Syntax**

```
#include <stdlib.h>
int abs(int i)
```

#### **Function**

**abs** obtains the absolute value of i. No check is made to see that the result can be properly represented.

#### Return Value

abs returns the absolute value of i, expressed as an int.

## **Example**

To print out a debit or credit balance:

```
printf ("balance ds \n", abs(bal), (bal < 0)? "CR": "");
```

#### See Also

labs, fabs

#### **Notes**

abs is packaged in the integer library, and may be implemented as a macro.

## acos

## **Description**

Arccosine

## **Syntax**

```
#include <math.h>
double acos(double x)
```

#### Function

**acos** computes the angle in radians the cosine of which is  $\mathbf{x}$ , to full double precision.

#### Return Value

acos returns the closest internal representation to acos(x), expressed as a double floating value in the range [0, pi]. If x is outside the range [-1, 1], acos returns zero.

## Example

To find the arccosine of x:

```
theta = acos(x);
```

#### See Also

asin, atan, atan2

#### Notes

acos is packaged in the floating point library.

## asin

## **Description**

Arcsine

## **Syntax**

```
#include <math.h>
double asin(double x)
```

#### **Function**

**asin** computes the angle in radians the sine of which is  $\mathbf{x}$ , to full double precision.

#### Return Value

asin returns the nearest internal representation to asin(x), expressed as a double floating value in the range [-pi/2, pi/2]. If **x** is outside the range [-1, 1], asin returns zero.

## **Example**

To compute the arcsine of y:

```
theta = asin(y);
```

#### See Also

acos, atan, atan2

#### **Notes**

asin is packaged in the floating point library.

## atan

## **Description**

Arctangent

## **Syntax**

```
#include <math.h>
double atan(double x)
```

#### Function

atan computes the angle in radians; the tangent of which is  $\mathbf{x}$ , atan computes the angle in radians; the tangent of which is  $\mathbf{x}$ , to full double precision.

#### **Return Value**

atan returns the nearest internal representation to atan(x), expressed as a double floating value in the range [-pi/2, pi/2].

## Example

To find the phase angle of a vector in degrees:

```
theta = atan(y/x) * 180.0 / pi;
```

#### See Also

acos, asin, atan2

#### **Notes**

atan is packaged in the floating point library.

## atan2

## **Description**

Arctangent of y/x

## **Syntax**

```
#include <math.h>
double atan2(double y, double x)
```

#### **Function**

atan2 computes the angle in radians the tangent of which is y/x to full double precision. If y is negative, the result is negative. If x is negative, the magnitude of the result is greater than pi/2.

#### Return Value

atan2 returns the closest internal representation to atan(y/x), expressed as a double floating value in the range [-pi, pi]. If both input arguments are zero, atan2 returns zero.

## **Example**

To find the phase angle of a vector in degrees:

```
theta = atan2(y/x) * 180.0/pi;
```

#### See Also

acos, asin, atan

#### **Notes**

atan2 is packaged in the floating point library.

## atof

### **Description**

Convert buffer to double

## **Syntax**

```
#include <stdlib.h>
double atof(char *nptr)
```

#### **Function**

**atof** converts the string at *nptr* into a double. The string is taken as the text representation of a decimal number, with an optional fraction and exponent. Leading whitespace is skipped and an optional sign is permitted; conversion stops on the first unrecognizable character. Acceptable inputs match the pattern:

where **d** is any decimal digit and **e** is the character '**e**' or '**E**'. No checks are made against overflow, underflow, or invalid character strings.

#### **Return Value**

*atof* returns the converted double value. If the string has no recognizable characters, it returns zero.

## **Example**

To read a string from STDIN and convert it to a double at d:

```
gets(buf);
d = atof(buf);
```

### See Also

atoi, atol, strtol, strtod

#### Notes

atof is packaged in the floating point library.

## atoi

## **Description**

Convert buffer to integer

## **Syntax**

```
#include <stdlib.h>
int atoi(char *nptr)
```

#### **Function**

**atoi** converts the string at *nptr* into an integer. The string is taken as the text representation of a decimal number. Leading whitespace is skipped and an optional sign is permitted; conversion stops on the first unrecognizable character. Acceptable characters are the decimal digits. If the stop character is **l** or **L**, it is skipped over.

No checks are made against overflow or invalid character strings.

### Return Value

*atoi* returns the converted integer value. If the string has no recognizable characters, zero is returned.

## **Example**

To read a string from STDIN and convert it to an *int* at **i**:

```
gets(buf);
i = atoi(buf);
```

#### See Also

atof, atol, strtol, strtod

#### Notes

atoi is packaged in the integer library.

# atol

### **Description**

Convert buffer to long

### **Syntax**

```
#include <stdlib.h>
long atol(char *nptr)
```

#### **Function**

**atol** converts the string at *nptr* into a long integer. The string is taken as the text representation of a decimal number. Leading whitespace is skipped and an optional sign is permitted; conversion stops on the first unrecognizable character. Acceptable characters are the decimal digits. If the stop character is **l** or **L** it is skipped over.

No checks are made against overflow or invalid character strings.

#### **Return Value**

*atol* returns the converted long integer. If the string has no recognizable characters, zero is returned.

# Example

To read a string from STDIN and convert it to a long 1:

```
gets(buf);
l = atol(buf);
```

#### See Also

atof, atoi, strtol, strtod

#### **Notes**

atol is packaged in the integer library.

# carry

# **Description**

Test or get the carry bit

# **Syntax**

#### **Function**

**carry** is an *inline* function allowing to test or get the value of the *carry* bit. When used in an *if* construct, this function expands directly to a **bcc** or **bcs** instruction. When used in an expression, it expands in order to build in the **a** register the value **0** or **1** depending on the *carry* bit value.

#### **Return Value**

carry returns 0 or 1 in the a register if such a value is needed.

# **Example**

```
low <<= 1;
                  produces
                                     sll
                                          low
if (carry())
                                     jruge L1
      ++high;
                                    inc
                                          high
                              T.1:
                  produces
low <<= 1;
                                     sll
                                           low
high = carry()
                                    clr
                                           а
                                     rlc
                                     ld
                                          high,a
```

#### Notes

carry is an *inline* function and then is not defined in any library. It is therefore not possible to take its address. For more information, see "*Inline Function*" in **Chapter 3**.

# ceil

# **Description**

Round to next higher integer

# **Syntax**

```
#include <math.h>
double ceil(double x)
```

### **Function**

**ceil** computes the smallest integer greater than or equal to  $\mathbf{x}$ .

#### Return Value

ceil returns the smallest integer greater than or equal to x, expressed as a double floating value.

# **Example**

X	ceil(x)
5.1	6.0
5.0	5.0
0.0	0.0
-5.0	-5.0
-5.1	-5.0

### See Also

floor

# Notes

ceil is packaged in the floating point library.

# checksum

# **Description**

Verify the recorded checksum

### **Syntax**

```
int _checksum()
```

#### **Function**

**\_checksum** scans the descriptor built by the linker and controls that the computed checksum is equal to the one expected. For more infomation, see "*Checksum Computation*" in **Chapter 6**.

#### Return Value

\_checksum returns 0 if the checksum is correct, or a value different of 0 otherwise.

# **Example**

```
if (_checksum())
          abort();
```

#### Notes

The descriptor is built by the linker only if the *\_checksum* function is called by the application, even if there are segments marked with the **-ck** option.

\_checksum is packaged in the integer library.

```
checksumx, checksum16, checksum16x
```

# \_checksumx

# **Description**

Verify the recorded checksum

# **Syntax**

```
int _checksumx()
```

#### Function

**\_checksumx** scans the descriptor built by the linker and controls at the end that the computed 8 bit checksum is equal to the one expected. For more infomation, see "*Checksum Computation*" in **Chapter 6**.

#### Return Value

\_checksumx returns 0 if the checksum is correct, or a value different of 0 otherwise.

### Example

```
if (_checksumx())
     abort();
```

#### Notes

The descriptor is built by the linker only if the *\_checksumx* function is called by the application, even if there are segments marked with the **-ck** option.

\_checksumx is packaged in the integer library.

```
checksum, checksum16, checksum16x
```

# checksum16

# **Description**

Verify the recorded checksum

# **Syntax**

```
int checksum16()
```

#### Function

checksum16 scans the descriptor built by the linker and controls at the end that the computed 16 bit checksum is equal to the one expected. For more infomation, see "Checksum Computation" in Chapter 6.

#### Return Value

checksum16 returns 0 if the checksum is correct, or a value different of 0 otherwise.

# Example

```
if (checksum16())
      abort();
```

#### Notes

The descriptor is built by the linker only if the *checksum16* function is called by the application, even if there are segments marked with the -ck option.

*checksum16* is packaged in the integer library.

```
checksum, checksumx, checksum16x
```

# checksum16x

# **Description**

Verify the recorded checksum

# **Syntax**

```
int checksum16x()
```

#### Function

checksum16x scans the descriptor built by the linker and controls at the end that the computed 16 bit checksum is equal to the one expected. For more infomation, see "Checksum Computation" in Chapter 6.

#### Return Value

checksum 16x returns 0 if the checksum is correct, or a value different of 0 otherwise.

### Example

```
if (checksum16x())
     abort();
```

#### Notes

The descriptor is built by the linker only if the *checksum16x* function is called by the application, even if there are segments marked with the -ck option.

*checksum16x* is packaged in the integer library.

```
checksum, checksumx, checksum16
```

# COS

# **Description**

Cosine

# **Syntax**

```
#include <math.h>
double cos(double x)
```

#### **Function**

**cos** computes the cosine of  $\mathbf{x}$ , expressed in radians, to full double precision. If the magnitude of  $\mathbf{x}$  is too large to contain a fractional quadrant part, the value of  $\cos$  is 1.

#### Return Value

cos returns the nearest internal representation to cos(x) in the range [0, pi], expressed as a double floating value. A large argument may return a meaningless value.

# Example

To rotate a vector through the angle **theta**:

```
xnew = xold * cos(theta) - yold * sin(theta);
ynew = xold * sin(theta) + yold * cos(theta);
```

#### See Also

sin, tan

#### **Notes**

cos is packaged in the floating point library.

# cosh

# **Description**

Hyperbolic cosine

# **Syntax**

```
#include <math.h>
double cosh(double x)
```

#### **Function**

 $\cosh$  computes the hyperbolic cosine of x to full double precision.

#### **Return Value**

cosh returns the nearest internal representation to cosh(x) expressed as a double floating value. If the result is too large to be properly represented, cosh returns zero.

# Example

To use the Moivre's theorem to compute (cosh x + sinh x) to the nth power:

```
demoivre = cosh(n * x) + sinh(n * x);
```

### See Also

exp, sinh, tanh

#### Notes

cosh is packaged in the floating point library.

# div

# **Description**

Divide with quotient and remainder

# **Syntax**

```
#include <stdlib.h>
div_t div(int numer, int denom)
```

#### Function

**div** divides the integer *numer* by the integer *denom* and returns the quotient and the remainder in a structure of type *div\_t*. The field *quot* contains the quotient and the field *rem* contains the remainder.

#### Return Value

div returns a structure of type div\_t containing both quotient and remainder.

# **Example**

To get minutes and seconds from a delay in seconds:

```
div_t result;
    result = div(time, 60);
    min = result.quo;
    sec = result.rem;
```

#### See Also

ldiv

#### Notes

div is packaged in the integer library.

# eepcpy

# **Description**

Copy a buffer to an **eeprom** buffer

# **Syntax**

```
#include <string.h>
void *eepcpy(void *s1, void *s2, unsigned int n)
```

### **Function**

**eepcpy** copies the first *n* characters starting at location *s2* into the *eep*rom buffer beginning at s1.

#### Return Value

Nothing.

# **Example**

To place "first string, second string" in eepbuf[]:

```
eepcpy(eepbuf, "first string", 12);
eepcpy(eepbuf + 13, ", second string", 15);
```

#### See Also

eepset, eepera

#### Notes

eepcpy is packaged in the integer library.

# eepera

# **Description**

Erase the *eeprom* space

# **Syntax**

```
void eepera(void *s, unsigned int n)
```

### **Function**

**eepera** erases the first n bytes starting at location s.

#### Return Value

Nothing.

# **Example**

To erase the *eeprom* space:

```
eepera((void *)0x4000, 1024);
```

#### See Also

eepcpy, eepset

#### Notes

eepera is a macro declared in the <string.h> header file, calling the eepset function with argument 0.

# eepset

# **Description**

Propagate fill character throughout *eeprom* buffer

# **Syntax**

```
#include <string.h>
void eepset(void *s, char c, unsigned int n)
```

### **Function**

**eepset** floods the n character buffer starting at *eeprom* location s with fill character c. The function waits for all bytes to be programmed.

#### Return Value

Nothing.

# **Example**

To flood a 512 byte *eeprom* buffer with NULs:

```
eepset(eepbuf, '\0', BUFSIZ);
```

#### See Also

eepcpy, eepera

#### Notes

eepset is packaged in the integer library.

# exit

# **Description**

Exit program execution

# **Syntax**

```
#include <stdlib.h>
void exit(int status)
```

# **Function**

**exit** stops the execution of a program by switching to the startup module just after the call to the *main* function. The status argument is not used by the current implementation.

#### Return Value

exit never returns.

# **Example**

To *exit* in case of error:

```
if (fatal_error)
     exit();
```

#### See Also

abort

#### **Notes**

exit is in the startup module.



# **Description**

Exponential

# **Syntax**

```
#include <math.h>
double exp(double x)
```

#### Function

exp computes the exponential of x to full double precision.

#### Return Value

exp returns the nearest internal representation to exp  $\mathbf{x}$ , expressed as a double floating value. If the result is too large to be properly represented, exp returns zero.

# Example

To compute the hyperbolic sine of  $\mathbf{x}$ :

```
sinh = (exp(x) - exp(-x)) / 2.0;
```

# See Also

log

#### Notes

exp is packaged in the floating point library.

# fabs

**Description**Find double absolute value

# **Syntax**

```
#include <math.h>
double fabs(double x)
```

#### **Function**

**fabs** obtains the absolute value of x.

#### **Return Value**

fabs returns the absolute value of x, expressed as a double floating value.

# **Example**

Х	fabs(x)
5.0	5.0
0.0	0.0
-3.7	3.7

# See Also

abs. labs

#### **Notes**

fabs is packaged in the floating point library.

# fctcpy

# **Description**

Copy a moveable code segment in RAM

### **Syntax**

```
int fctcpy(char name);
```

#### **Function**

fctcpy copies a moveable code segment in RAM from its storage location in ROM. fctcpy scans the descriptor built by the linker and looks for a moveable segment whose flag byte matches the given argument. If such a segment is found, it is entirely copied in RAM. Any function defined in that segment may then be called directly. For more information, see "Moveable Code" in Chapter 6.

#### **Return Value**

fctcpy returns a non zero value if a segment has been found and copied. It returns 0 otherwise.

### Example

```
if (fctcpy('b'))
     flash();
```

#### **Notes**

fctcpy is packaged in the machine library.

# floor

# **Description**

Round to next lower integer

# **Syntax**

```
#include <math.h>
double floor(double x)
```

### **Function**

**floor** computes the largest integer less than or equal to  $\mathbf{x}$ .

#### **Return Value**

*floor* returns the largest integer less than or equal to  $\mathbf{x}$ , expressed as a double floating value.

# **Example**

Х	floor(x)
5.1	5.0
5.0	5.0
0.0	0.0
-5.0	-5.0
-5.1	-6.0

# See Also

ceil

#### **Notes**

floor is packaged in the floating point library.

# fmod

# **Description**

Find double modulus

# **Syntax**

```
#include <math.h>
double fmod(double x, double y)
```

### **Function**

**fmod** computes the floating point remainder of  $\mathbf{x} / \mathbf{y}$ , to full double precision. The return value of **f** is determined using the formula:

$$f = x - i * y$$

where i is some integer, f is the same sign as x, and the absolute value of  $\mathbf{f}$  is less than the absolute value of  $\mathbf{y}$ .

### **Return Value**

*fmod* returns the value of **f** expressed as a double floating value. If **y** is zero, fmod returns zero.

# **Example**

X	У	fmod(x,	λ)
5.5	5.0	0.5	
5.0	5.0	0.0	
0.0	0.0	0.0	
-5.5	5.0	-0.5	

### **Notes**

*fmod* is packaged in the floating point library.

# frexp

# **Description**

Extract fraction from exponent part

# **Syntax**

```
#include <math.h>
double frexp(double val, int *exp)
```

#### **Function**

frexp partitions the double at val, which should be non-zero, into a fraction in the interval [1/2, 1) times two raised to an integer power. It then delivers the integer power to \*exp, and returns the fractional portion as the value of the function. The exponent is generally meaningless if val is zero.

#### Return Value

*frexp* returns the power of two fraction of the double at *val* as the return value of the function, and writes the exponent at \**exp*.

# **Example**

To implement the sqrt(x) function:

```
double sqrt(double x)
{
    extern double newton(double);
    int n;

    x = frexp(x, &n);
    x = newton(x);
    if (n & 1)
        x *= SQRT2;
    return (ldexp(x, n / 2));
}
```

#### See Also

ldexp

#### **Notes**

*frexp* is packaged in the floating point library.

# getchar

### **Description**

Get character from input stream

# **Syntax**

```
#include <stdio.h>
int getchar(void)
```

#### **Function**

getchar obtains the next input character, if any, from the user supplied input stream. This user must rewrite this function in C or in assembly language to provide an interface to the input mechanism of the C library.

#### **Return Value**

getchar returns the next character from the input stream. If end of file (break) is encountered, or a read error occurs, getchar returns EOF.

# **Example**

To copy characters from the input stream to the output stream:

```
while ((c = getchar()) != EOF)
      putchar(c);
```

#### See Also

putchar

#### Notes

getchar is packaged in the integer library, and is by default using the first serial port SCI 1.

# gets

# **Description**

Get a text line from input stream

### **Syntax**

```
#include <stdio.h>
char *gets(char *s)
```

#### **Function**

**gets** copies characters from the input stream to the buffer starting at **s**. Characters are copied until a newline is reached or end of file is reached. If a newline is reached, it is discarded and a NUL is written immediately following the last character read into **s**.

gets uses getchar to read each character.

#### Return Value

gets returns **s** if successful. If end of file is reached, gets returns NULL. If a read error occurs, the array contents are indeterminate and gets returns NULL.

# **Example**

To copy input to output, line by line:

```
while (puts(gets(buf)))
;
```

#### See Also

puts

#### **Notes**

There is no assured limit on the size of the line read by *gets*.

gets is packaged in the integer library.

# imask

# **Description**

Test the interrupt mask bit

# **Syntax**

```
#include cessor.h>
@inline char imask(void)
```

#### **Function**

imask is an inline function allowing to test the interrupt mask bit. The imask function can only be used in an if construct. This function expands directly to a bms or bmc instruction.

#### Return Value

None

# **Example**

```
produces
if (imask())
                                   jrnm L3
     ++high;
                                   inc
                                        high
                             L3:
if (!imask())
                 produces
                                   bms
                                        L1
     ++high
                                   inc high
                             L1:
```

#### **Notes**

imask is an inline function and then is not defined in any library. It is therefore not possible to take its address. For more information, see "Inline Function" in Chapter 3.

# irq

# **Description**

Test the interrupt line level

### **Syntax**

```
#include  processor.h>
@inline char irq(void)
```

#### **Function**

**irq** is an *inline* function allowing to test the interrupt line level. The *irq* function can only be used in an *if* construct. This function expands directly to a **bih** or **bil** instruction.

#### Return Value

None.

# **Example**

#### Notes

*irq* is an *inline* function and then is not defined in any library. It is therefore not possible to take its address. For more information, see "*Inline Function*" in **Chapter 3**.

# isalnum

# **Description**

Test for alphabetic or numeric character

# **Syntax**

```
#include <ctype.h>
int isalnum(char c)
```

#### Function

isalnum tests whether c is an alphabetic character (either upper or lower case), or a decimal digit.

#### Return Value

isalnum returns nonzero if the argument is an alphabetic or numeric character; otherwise the value returned is zero.

# **Example**

To test for a valid C identifier:

```
if (isalpha(*s) || *s == ' ')
      for (++s; isalnum(*s) || *s == ' '; ++s)
```

#### See Also

isalpha, isdigit, islower, isupper, isxdigit, tolower, toupper

#### **Notes**

If the argument is outside the range [-1, 255], the result is undefined.

*isalnum* is packaged in the integer library.

# isalpha

# **Description**

Test for alphabetic character

# **Syntax**

```
#include <ctype.h>
int isalpha(char c)
```

#### **Function**

isalpha tests whether c is an alphabetic character, either upper or lower case.

#### Return Value

isalpha returns nonzero if the argument is an alphabetic character. Otherwise the value returned is zero.

# **Example**

To find the end points of an alphabetic string:

#### See Also

isalnum, isdigit, islower, isupper, isxdigit, tolower, toupper

#### Notes

If the argument is outside the range [-1, 255], the result is undefined.

isalpha is packaged in the integer library.

# iscntrl

# **Description**

Test for control character

# **Syntax**

```
#include <ctype.h>
int iscntrl(char c)
```

#### Function

iscntrl tests whether c is a delete character (0177 in ASCII), or an ordinary control character (less than 040 in ASCII).

#### Return Value

iscntrl returns nonzero if c is a control character; otherwise the value is zero.

# Example

To map control characters to percent signs:

```
for (; *s; ++s)
      if (iscntrl(*s))
             *s = '%';
```

#### See Also

isgraph, isprint, ispunct, isspace

#### **Notes**

If the argument is outside the range [-1, 255], the result is undefined.

iscntrl is packaged in the integer library.

# isdigit

# **Description**

Test for digit

# **Syntax**

```
#include <ctype.h>
int isdigit(char c)
```

#### Function

isdigit tests whether c is a decimal digit.

#### Return Value

isdigit returns nonzero if c is a decimal digit; otherwise the value returned is zero.

# **Example**

To convert a decimal digit string to a number:

```
for (sum = 0; isdigit(*s); ++s)
sum = sum * 10 + *s - '0';
```

#### See Also

isalnum, isalpha, islower, isupper, isxdigit, tolower, toupper

#### **Notes**

If the argument is outside the range [-1, 255], the result is undefined.

isdigit is packaged in the integer library.

# isgraph

# **Description**

Test for graphic character

# **Syntax**

```
#include <ctype.h>
int isgraph(char c)
```

#### Function

isgraph tests whether c is a graphic character; i.e. any printing character except a space (040 in ASCII).

#### Return Value

isgraph returns nonzero if c is a graphic character. Otherwise the value returned is zero.

# Example

To output only graphic characters:

```
for (; *s; ++s)
      if (isgraph(*s))
            putchar(*s);
```

#### See Also

iscntrl, isprint, ispunct, isspace

#### **Notes**

If the argument is outside the range [-1, 255], the result is undefined.

isgraph is packaged in the integer library.

# islower

**Description**Test for lowercase character

# **Syntax**

```
#include <ctype.h>
int islower(char c)
```

#### **Function**

**islower** tests whether **c** is a lowercase alphabetic character.

#### Return Value

islower returns nonzero if c is a lowercase character; otherwise the value returned is zero.

# **Example**

To convert to uppercase:

```
if (islower(c))
      c += 'A' - 'a';  /* also see toupper() */
```

#### See Also

isalnum, isalpha, isdigit, isupper, isxdigit, tolower, toupper

#### Notes

If the argument is outside the range [-1, 255], the result is undefined.

*islower* is packaged in the integer library.

# isprint

# **Description**

Test for printing character

# **Syntax**

```
#include <ctype.h>
int isprint(char c)
```

#### Function

**isprint** tests whether c is any printing character. Printing characters are all characters between a space (040 in ASCII) and a tilde '~' character (0176 in ASCII).

#### Return Value

isprint returns nonzero if c is a printing character; otherwise the value returned is zero.

# **Example**

To output only printable characters:

```
for (; *s; ++s)
      if (isprint(*s))
            putchar(*s);
```

#### See Also

iscntrl, isgraph, ispunct, isspace

### **Notes**

If the argument is outside the range [-1, 255], the result is undefined.

isprint is packaged in the integer library.

# ispunct

# **Description**

Test for punctuation character

# **Syntax**

```
#include <ctype.h>
int ispunct(char c)
```

#### **Function**

**ispunct** tests whether **c** is a punctuation character. Punctuation characters include any printing character except space, a digit, or a letter.

#### Return Value

*ispunct* returns nonzero if  $\mathbf{c}$  is a punctuation character; otherwise the value returned is zero.

# **Example**

To collect all punctuation characters in a string into a buffer:

```
for (i = 0; *s; ++s)
    if (ispunct(*s))
        buf[i++] = *s;
```

#### See Also

iscntrl, isgraph, isprint, isspace

#### Notes

If the argument is outside the range [-1, 255], the result is undefined.

ispunct is packaged in the integer library.

# isspace

# **Description**

Test for whitespace character

# **Syntax**

```
#include <ctype.h>
int isspace(char c)
```

#### **Function**

**isspace** tests whether c is a whitespace character. Whitespace characters are horizontal tab ('\t'), newline ('\n'), vertical tab ('\v'), form feed ('\f'), carriage return ('\r'), and space ('').

#### **Return Value**

isspace returns nonzero if c is a whitespace character; otherwise the value returned is zero.

# Example

To skip leading whitespace:

```
while (isspace(*s))
      ++s;
```

#### See Also

iscntrl, isgraph, isprint, ispunct

### **Notes**

If the argument is outside the range [-1, 255], the result is undefined.

isspace is packaged in the integer library.

# isupper

# **Description**

Test for uppercase character

# **Syntax**

```
int isupper(char c)
```

#### **Function**

**isupper** tests whether c is an uppercase alphabetic character.

#### Return Value

isupper returns nonzero if c is an uppercase character; otherwise the value returned is zero.

# **Example**

To convert to lowercase:

#### See Also

isalnum, isalpha, isdigit, islower, isxdigit, tolower, toupper

#### Notes

If the argument is outside the range [-1, 255], the result is undefined.

isupper is packaged in the integer library.

# isxdigit

### **Description**

Test for hexadecimal digit

# **Syntax**

```
#include <ctype.h>
int isxdigit(char c)
```

#### **Function**

isxdigit tests whether c is a hexadecimal digit, i.e. in the set [0123456789abcdefABCDEF].

#### **Return Value**

isxdigit returns nonzero if c is a hexadecimal digit; otherwise the value returned is zero.

# Example

To accumulate a hexadecimal digit:

```
for (sum = 0; isxdigit(*s); ++s)
     if (isdigit(*s)
            sum = sum * 10 + *s - '0';
     else
            sum = sum * 10 + tolower(*s) + (10 - 'a');
```

### See Also

isalnum, isalpha, isdigit, islower, isupper, tolower, toupper

#### Notes

If the argument is outside the range [-1, 255], the result is undefined.

isxdigit is packaged in the integer library.

# labs

# **Description**

Find long absolute value

# **Syntax**

```
#include <stdlib.h>
long labs(long 1)
```

#### **Function**

**labs** obtains the absolute value of **l**. No check is made to see that the result can be properly represented.

#### Return Value

labs returns the absolute value of I, expressed as an long int.

# **Example**

To print out a debit or credit balance:

```
printf("balance %ld%s\n",labs(bal),(bal < 0) ? "CR" : "");</pre>
```

#### See Also

abs, fabs

#### **Notes**

labs is packaged in the integer library.

# ldexp

## **Description**

Scale double exponent

## **Syntax**

```
#include <math.h>
double ldexp(double x, int exp)
```

### Function

**ldexp** multiplies the double **x** by two raised to the integer power **exp**.

#### Return Value

*ldexp* returns the double result  $x * (1 \le exp)$  expressed as a double floating value. If a range error occurs, *ldexp* returns **HUGE VAL**.

## **Example**

Х	exp	ldexp(x,	exp)
1.0	1	2.0	
1.0	0	1.0	
1.0	-1	0.5	
0.0	0	0.0	

#### See Also

frexp, modf

### **Notes**

*ldexp* is packaged in the floating point library.

## ldiv

## **Description**

Long divide with quotient and remainder

## **Syntax**

```
#include <stdlib.h>
ldiv_t ldiv(long numer, long denom)
```

#### Function

**Idiv** divides the long integer *numer* by the long integer *denom* and returns the quotient and the remainder in a structure of type *ldiv\_t*. The field *quot* contains the quotient and the field *rem* contains the remainder.

#### Return Value

*ldiv* returns a structure of type *ldiv\_t* containing both quotient and remainder.

## **Example**

To get minutes and seconds from a delay in seconds:

```
ldiv_t result;
    result = ldiv(time, 60L);
    min = result.quo;
    sec = result.rem;
```

#### See Also

div

#### **Notes**

ldiv is packaged in the integer library.

# log

## **Description**

Natural logarithm

## **Syntax**

```
#include <math.h>
double log(double x)
```

#### Function

 $\log$  computes the natural logarithm of x to full double precision.

#### **Return Value**

log returns the closest internal representation to log(x), expressed as a double floating value. If the input argument is less than zero, or is too large to be represented, log returns zero.

## Example

To compute the hyperbolic arccosine of  $\mathbf{x}$ :

```
arccosh = log(x + sqrt(x * x - 1));
```

## See Also

exp

#### Notes

log is packaged in the floating point library.

# log10

## **Description**

Common logarithm

## **Syntax**

```
#include <math.h>
double log10(double x)
```

#### **Function**

log10 computes the common log of x to full double precision by computing the natural log of x divided by the natural log of 10. If the input argument is less than zero, a domain error will occur. If the input argument is zero, a range error will occur.

#### Return Value

log10 returns the nearest internal representation to log10 **x**, expressed as a double floating value. If the input argument is less than or equal to zero, log10 returns zero.

## **Example**

To determine the number of digits in  $\mathbf{x}$ , where  $\mathbf{x}$  is a positive integer expressed as a double:

```
ndig = log10(x) + 1;
```

## See Also

log

### **Notes**

log10 is packaged in the floating point library.

## max

## **Description**

Test for maximum

## **Syntax**

```
#include <stdlib.h>
max(a,b)
```

#### **Function**

max obtains the maximum of its two arguments, a and b. Since max is implemented as a C preprocessor macro, its arguments can be any numerical type, and type coercion occurs automatically.

#### **Return Value**

max is a numerical rvalue of the form ((a < b) ? b : a), suitably parenthesized.

## **Example**

To set a new maximum level:

```
hiwater = max(hiwater, level);
```

## See Also

min

#### Notes

max is an extension to the proposed ANSI C standard.

max is a macro declared in the <stdlib.h> header file. You can use it by including <stdlib.h> with your program. Because it is a macro, max cannot be called from non-C programs, nor can its address be taken. Arguments with side effects may be evaluated other than once.

## memchr

## **Description**

Scan buffer for character

## **Syntax**

```
#include <string.h>
void *memchr(void *s, char c, unsigned char n)
```

#### **Function**

**memchr** looks for the first occurrence of a specific character c in an n character buffer starting at s.

#### Return Value

*memchr* returns a pointer to the first character that matches c, or NULL if no character matches.

## **Example**

To map *keybuf[]* characters into *subst[]* characters:

```
if ((t = memchr(keybuf, *s, KEYSIZ)) != NULL)
     *s = subst[t - keybuf];
```

### See Also

strchr, strcspn, strpbrk, strrchr, strspn

#### Notes

*memchr* is packaged in the integer library.

## memcmp

## **Description**

Compare two buffers for lexical order

## **Syntax**

```
#include <string.h>
int memcmp(void *s1, void *s2, unsigned char n)
```

#### **Function**

memcmp compares two text buffers, character by character, for lexical order in the character collating sequence. The first buffer starts at s1, the second at s2; both buffers are n characters long.

#### **Return Value**

memcmp returns a short integer greater than, equal to, or less than zero, according to whether s1 is lexicographically greater than, equal to, or less than s2.

## Example

To look for the string "include" in name:

```
if (memcmp(name, "include", 7) == 0)
      doinclude();
```

#### See Also

strcmp, strncmp

#### Notes

*memcmp* is packaged in the integer library.

## memcpy

## **Description**

Copy one buffer to another

## **Syntax**

```
#include <string.h>
void *memcpy(void *s1, void *s2, unsigned char n)
```

#### **Function**

**memcpy** copies the first n characters starting at location s2 into the buffer beginning at s1.

#### Return Value

memcpy returns s1.

## **Example**

To place "first string, second string" in buf[]:

```
memcpy(buf, "first string", 12);
memcpy(buf + 13, ", second string", 15);
```

### See Also

strcpy, strncpy

#### Notes

*memcpy* is packaged in the integer library and may be implemented as an *inline* function.

## memmove

## **Description**

Copy one buffer to another

## **Syntax**

```
#include <string.h>
void *memmove(void *s1, void *s2, unsigned char n)
```

#### **Function**

memmove copies the first n characters starting at location s2 into the buffer beginning at s1. If the two buffers overlap, the function performs the copy in the appropriate sequence, so the copy is not corrupted.

#### Return Value

memmove returns \$1.

## Example

To shift an array of characters:

```
memmove(buf, &buf[5], 10);
```

#### See Also

тетсру

#### Notes

*memmove* is packaged in the integer library.

## memset

## **Description**

Propagate fill character throughout buffer

## **Syntax**

```
#include <string.h>
void *memset(void *s, char c, unsigned char n)
```

#### **Function**

memset floods the n character buffer starting at s with fill character c.

#### Return Value

memset returns s.

## **Example**

To flood a 512-byte buffer with NULs:

```
memset(buf,'\0', BUFSIZ);
```

#### Notes

memset is packaged in the integer library and may be implemented as an inline function



## **Description**

Test for minimum

## **Syntax**

```
#include <stdlib.h>
min(a, b)
```

#### **Function**

min obtains the minimum of its two arguments, a and b. Since min is implemented as a C preprocessor macro, its arguments can be any numerical type, and type coercion occurs automatically.

#### **Return Value**

min is a numerical rvalue of the form ((a < b) ? a : b), suitably parenthesized.

## **Example**

To set a new minimum level:

```
nmove = min(space, size);
```

#### See Also

max

#### Notes

min is an extension to the ANSI C standard.

min is a macro declared in the <stdlib.h> header file. You can use it by including <stdlib.h> with your program. Because it is a macro, min cannot be called from non-C programs, nor can its address be taken. Arguments with side effects may be evaluated more than once.

## modf

## **Description**

Extract fraction and integer from double

## **Syntax**

```
#include <math.h>
double modf(double val, double *pd)
```

#### **Function**

modf partitions the double val into an integer portion, which is delivered to \*pd, and a fractional portion, which is returned as the value of the function. If the integer portion cannot be represented properly in an int, the result is truncated on the left without complaint.

#### Return Value

modf returns the signed fractional portion of val as a double floating value, and writes the integer portion at \*pd.

## **Example**

val	*pd	modf(val,	*pd)
5.1	5	0.1	
5.0	5	0.0	
4.9	4	0.9	
0.0	0	0.0	
-1.4	-1	-0.4	

#### See Also

frexp, ldexp

#### Notes

*modf* is packaged in the floating point library.

## **Description**

Raise x to the y power

## **Syntax**

```
#include <math.h>
double pow(double x, double y)
```

## **Function**

**pow** computes the value of x raised to the power of y.

#### Return Value

pow returns the value of x raised to the power of y, expressed as a double floating value. If x is zero and y is less than or equal to zero, or if xis negative and y is not an integer, pow returns zero.

## **Example**

Х	У	pow(x,	λ)
2.0	2.0	4.0	
2.0	1.0	2.0	
2.0	0.0	1.0	
1.0	any	1.0	
0.0	-2.0	0	
-1.0	2.0	1.0	
-1.0	2.1	0	

### See Also

exp

#### Notes

pow is packaged in the floating point library.

# printf

### Description

Output formatted arguments to stdout

## **Syntax**

```
#include <stdio.h>
int printf(@near char *fmt, ...)
```

#### **Function**

**printf** writes formatted output to the output stream using the format string at *fmt* and the arguments specified by ..., as described below.

printf uses putchar to output each character.

## **Format Specifiers**

The format string at *fint* consists of literal text to be output, interspersed with conversion specifications that determine how the arguments are to be interpreted and how they are to be converted for output. If there are insufficient arguments for the format, the results are undefined. If the format is exhausted while arguments remain, the excess arguments are evaluated but otherwise ignored. *printf* returns when the end of the format string is encountered.

Each < conversion specification > is started by the character '%'. After the '%', the following appear in sequence:

<flags> - zero or more which modify the meaning of the conversion
specification.

**field width>** - a decimal number which optionally specifies a minimum field width. If the converted value has fewer characters than the field width, it is padded on the left (or right, if the left adjustment flag has been given) to the field width. The padding is with spaces unless the field width digit string starts with zero, in which case the padding is with zeros.

cision> - a decimal number which specifies the minimum number of digits to appear for d, i, o, u, x, and X conversions, the number of digits to appear after the decimal point for e, E, and f conversions, the maximum number of significant digits for the g and G conversions, or the maximum number of characters to be printed from a string in an s conversion. The precision takes the form of a period followed by a decimal digit string. A null digit string is treated as zero.

**h** - optionally specifies that the following **d**, **i**, **o**, **u**, **x**, or **X** conversion character applies to a short int or unsigned short int argument (the argument will have been widened according to the integral widening conversions, and its value must be cast to short or unsigned short before printing). It specifies a short pointer argument if associated with the p conversion character. If an h appears with any other conversion character, it is ignored.

**l** - optionally specifies that the **d**, **i**, **o**, **u**, **x**, and **X** conversion character applies to a long int or unsigned long int argument. It specifies a long or far pointer argument if used with the p conversion character. If the l appears with any other conversion character, it is ignored.

L - optionally specifies that the following e, E, f, g, and G conversion character applies to a long double argument. If the L appears with any other conversion character, it is ignored.

<conversion character> - character that indicates the type of conversion to be applied.

A field width or precision, or both, may be indicated by an asterisk '\*' instead of a digit string. In this case, an int argument supplies the field width or precision. The arguments supplying field width must appear before the optional argument to be converted. A negative field width argument is taken as a - flag followed by a positive field width. A negative precision argument is taken as if it were missing.

The *<flags>* field is zero or more of the following:

**space** - a space will be prepended if the first character of a signed conversion is not a sign. This flag will be ignored if space and + flags are both specified.

# - result is to be converted to an "alternate form". For c, d, i, s, and u conversions, the flag has no effect. For o conversion, it increases the precision to force the first digit of the result to be zero. For p, x and X conversion, a non-zero result will have Ox or OX prepended to it. For e, E, f, g, and G conversions, the result will contain a decimal point, even if no digits follow the point. For g and G conversions, trailing zeros will not be removed from the result, as they normally are. For p conversion, it designates hexadecimal output.

- + result of signed conversion will begin with a plus or minus sign.
- - result of conversion will be left justified within the field.

The *<conversion character>* is one of the following:

- % a '%' is printed. No argument is converted.
- **c** the char argument is converted to a character and printed.
- d, i, o, u, x, X the int argument is converted to signed decimal (d or i), unsigned octal (o), unsigned decimal (u), or unsigned hexadecimal notation (x or X); the letters **abcdef** are used for x conversion and the letters **ABCDEF** are used for x conversion. The precision specifies the minimum number of digits to appear; if the value being converted can be represented in fewer digits, it will be expanded with leading zeros. The default precision is x. The result of converting a zero value with precision of zero is no characters.
- **e, E** the double argument is converted in the style [-]d.ddde+dd, where there is one digit before the decimal point and the number of digits after it is equal to the precision. If the precision is missing, six digits are produced; if the precision is zero, no decimal point appears. The **E** format code will produce a number with **E** instead of **e** introducing the exponent. The exponent always contains at least two digits. However, if the magnitude to be printed is greater than or equal to 1E+100, additional exponent digits will be printed as necessary.
- **f** the double argument is converted to decimal notation in the style **[-]ddd.ddd**, where the number of digits following the decimal point is equal to the precision specification. If the precision is missing, it is

taken as 6. If the precision is explicitly zero, no decimal point appears. If a decimal point appears, at least one digit appears before it.

- **g.** G the double argument is printed in style f or e (or in style E in the case of a G format code), with the precision specifying the number of significant digits. The style used depends on the value converted; style e will be used only if the exponent resulting from the conversion is less than -4 or greater than the precision. Trailing zeros are removed from the result; a decimal point appears only if it is followed by a digit.
- **n** the argument is taken to be an int \* pointer to an integer into which is written the number of characters written to the output stream so far by this call to *printf*. No argument is converted.
- **p** the argument is taken to be a *void* \* pointer to an object. The value of the pointer is converted to a sequence of printable characters, and printed as a hexadecimal number with the number of digits printed being determined by the field width.
- **s**, **S** the argument is taken to be a *char* \* pointer to a string. When the Compact model is used, the s format will use a @tiny pointer and the S format will use a @near pointer. Characters from the string are written up to, but not including, the terminating NUL, or until the number of characters indicated by the precision are written. If the precision is missing, it is taken to be arbitrarily large, so all characters before the first NUL are printed.

If the character after '%' is not a valid conversion character, the behavior is undefined

If any argument is or points to an aggregate (except for an array of characters using %s conversion or any pointer using %p conversion), unpredictable results will occur.

A nonexistent or small field width does not cause truncation of a field; if the result is wider than the field width, the field is expanded to contain the conversion result.

#### Return Value

printf returns the number of characters transmitted, or a negative number if a write error occurs

#### Notes

A call with more conversion specifiers than argument variables will cause unpredictable results.

## **Example**

To print arg, which is a double with the value 5100.53:

```
printf("%8.2f\n", arg);
printf("%*.*f\n", 8, 2, arg);
```

both forms will output: 05100.53

#### See Also

sprintf

#### Notes

printf is packaged in both the integer library and the floating point library. The functionality of the integer only version of *printf* is a subset of the functionality of the floating point version. The integer only version cannot print or manipulate floating point numbers. If your programs call the integer only version of printf, the following conversion specifiers are invalid: e, E, f, g and G. The L modifier is also invalid.

If printf encounters an invalid conversion specifier, the invalid specifier is ignored and no special message is generated.

# putchar

## **Description**

Put a character to output stream

## **Syntax**

```
#include <stdio.h>
int putchar(char c)
```

#### **Function**

putchar copies c to the user specified output stream.

You must rewrite *putchar* in either C or assembly language to provide an interface to the output mechanism to the C library.

#### Return Value

putchar returns c. If a write error occurs, putchar returns EOF.

## Example

To copy input to output:

```
while ((c = getchar()) != EOF)
      putchar(c);
```

#### See Also

getchar

#### **Notes**

putchar is packaged in the integer library, and is by default using the first serial port SCI 1.

# puts

## **Description**

Put a text line to output stream

## **Syntax**

```
#include <stdio.h>
int puts(char *s)
```

#### **Function**

**puts** copies characters from the buffer starting at **s** to the output stream and appends a newline character to the output stream.

*puts* uses *putchar* to output each character. The terminating NUL is not copied.

#### Return Value

puts returns zero if successful, or else nonzero if a write error occurs.

## **Example**

To copy input to output, line by line:

```
while (puts(gets(buf)))
;
```

#### See Also

gets

#### **Notes**

puts is packaged in the integer library.

## rand

## **Description**

Generate pseudo-random number

## **Syntax**

```
#include <stdlib.h>
int rand(void)
```

### Function

rand computes successive pseudo-random integers in the range [0, 32767], using a linear multiplicative algorithm which has a period of 2 raised to the power of 32.

## **Example**

```
int dice()
      return (rand() % 6 + 1);
```

#### **Return Value**

rand returns a pseudo-random integer.

### See Also

srand

#### **Notes**

rand is packaged in the integer library.

## sin

## **Description**

Sin

## **Syntax**

```
#include <math.h>
double sin(double x)
```

#### **Function**

sin computes the sine of  $\mathbf{x}$ , expressed in radians, to full double precision. If the magnitude of  $\mathbf{x}$  is too large to contain a fractional quadrant part, the value of sin is 0.

#### Return Value

sin returns the closest internal representation to sin(x) in the range [-pi/2, pi/2], expressed as a double floating value. A large argument may return a meaningless result.

## **Example**

To rotate a vector through the angle *theta*:

```
xnew = xold * cos(theta) - yold * sin(theta);
ynew = xold * sin(theta) + yold * cos(theta);
```

#### See Also

cos. tan

#### **Notes**

*sin* is packaged in the floating point library.

## sinh

## **Description**

Hyperbolic sine

## **Syntax**

```
#include <math.h>
double sinh(double x)
```

#### **Function**

sinh computes the hyperbolic sine of x to full double precision.

#### **Return Value**

sinh returns the closest internal representation to sinh(x), expressed as a double floating value. If the result is too large to be properly represented, sinh returns zero.

## Example

To obtain the hyperbolic sine of complex **z**:

```
typedef struct
      double x, iy;
      }complex;
complex z;
      z.x = sinh(z.x) * cos(z.iy);
      z.iy = cosh(z.x) * sin(z.iy);
```

#### See Also

cosh, exp, tanh

#### Notes

*sinh* is packaged in the floating point library.

# sprintf

## **Description**

Output arguments formatted to buffer

## **Syntax**

```
#include <stdio.h>
int sprintf(char *s, @near char fmt, ...)
```

#### **Function**

sprintf writes formatted to the buffer pointed at by s using the format string at fmt and the arguments specified by ..., in exactly the same way as printf. See the description of the printf function for information on the format conversion specifiers. A NUL character is written after the last character in the buffer

#### Return Value

sprintf returns the numbers of characters written, not including the terminating NUL character.

## Example

To format a double at **d** into *buf*:

```
sprintf(buf, "%10f\n", d);
```

#### See Also

printf

#### Notes

sprintf is packaged in both the integer library and the floating point library. The functionality of the integer only version of sprintf is a subset of the functionality of the floating point version. The integer only version cannot print or manipulate floating point numbers. If your programs call the integer only version of sprintf, the following conversion specifiers are invalid: e, E, f, g and G. The L flag is also invalid.

# sqrt

## **Description**

Real square root

## **Syntax**

```
#include <math.h>
double sqrt(double x)
```

#### Function

**sqrt** computes the square root of  $\mathbf{x}$  to full double precision.

#### **Return Value**

sqrt returns the nearest internal representation to sqrt(x), expressed as a double floating value. If x is negative, sqrt returns zero.

## **Example**

To use sqrt to check whether n > 2 is a prime number:

```
if (!(n & 01))
      return (NOTPRIME);
sq = sqrt((double)n);
for (div = 3; div \le sq; div += 2)
      if (!(n % div))
             return (NOTPRIME);
return (PRIME);
```

### **Notes**

*sqrt* is packaged in the floating point library.

# srand

## **Description**

Seed pseudo-random number generator

## **Syntax**

```
#include <stdlib.h>
void srand(unsigned char nseed)
```

#### **Function**

**srand** uses *nseed* as a seed for a new sequence of pseudo-random numbers to be returned by subsequent calls to *rand*. If *srand* is called with the same seed value, the sequence of pseudo-random numbers will be repeated. The initial seed value used by *rand* and *srand* is 1.

#### Return Value

Nothing.

## **Example**

To set up a new sequence of random numbers:

srand(103);

#### See Also

rand

#### Notes

*srand* is packaged in the integer library.

## streat

## **Description**

Concatenate strings

## **Syntax**

```
#include <string.h>
char *strcat(char *s1, char *s2)
```

#### **Function**

streat appends a copy of the NUL terminated string at s2 to the end of the NUL terminated string at s1. The first character of s2 overlaps the NUL at the end of s1. A terminating NUL is always appended to s1.

#### **Return Value**

strcat returns s1.

## **Example**

To place the strings "first string, second string" in buf[]:

```
buf[0] = ' \setminus 0';
strcpy(buf, "first string");
strcat(buf, ", second string");
```

#### See Also

strncat

#### **Notes**

There is no way to specify the size of the destination area to prevent storage overwrites.

strcat is packaged in the integer library.

## strchr

## **Description**

Scan string for first occurrence of character

## **Syntax**

```
#include <string.h>
char *strchr(char *s1, char c)
```

#### **Function**

strchr looks for the first occurrence of a specific character c in a NUL terminated target string s.

#### Return Value

strchr returns a pointer to the first character that matches c, or NULL if none does.

## **Example**

To map *keystr[]* characters into *subst[]* characters:

```
if (t = strchr(keystr, *s))
      *s = subst[t - keystr];
```

### See Also

memchr, strcspn, strpbrk, strrchr, strspn

#### **Notes**

strchr is packaged in the integer library.

# stremp

## **Description**

Compare two strings for lexical order

## **Syntax**

```
#include <string.h>
int strcmp(char *s1, char *s2)
```

#### **Function**

strcmp compares two text strings, character by character, for lexical order in the character collating sequence. The first string starts at s1, the second at s2. The strings must match, including their terminating NUL characters, in order for them to be equal.

#### Return Value

strcmp returns an integer greater than, equal to, or less than zero, according to whether s1 is lexicographically greater than, equal to, or less than s2.

## **Example**

To look for the string "include":

```
if (strcmp(buf, "include") == 0)
      doinclude();
```

#### See Also

memcmp, strncmp

#### Notes

*strcmp* is packaged in the integer library.

# strcpy

## **Description**

Copy one string to another

## **Syntax**

```
#include <string.h>
char *strcpy(char *s1, char *s2)
```

#### **Function**

**strcpy** copies the NUL terminated string at **s2** to the buffer pointed at by **s1**. The terminating NUL is also copied.

#### Return Value

strcpy returns s1.

## **Example**

To make a copy of the string s2 in dest:

```
strcpy(dest, s2);
```

#### See Also

memcpy, strncpy

#### **Notes**

There is no way to specify the size of the destination area, to prevent storage overwrites.

*strcpy* is packaged in the integer library, and may be implemented as an *inline* function.

# strcspn

## **Description**

Find the end of a span of characters in a set

## **Syntax**

```
#include <string.h>
unsigned int strcspn(char *s1, char *s2)
```

#### Function

strcspn scans the string starting at s1 for the first occurrence of a character in the string starting at s2. It computes a subscript i such that:

- **s1[i]** is a character in the string starting at **s1**
- s1[i] compares equal to some character in the string starting at s2, which may be its terminating null character.

#### Return Value

strcspn returns the lowest possible value of i. s1[i] designates the terminating null character if none of the characters in s1 are in s2.

## Example

To find the start of a decimal constant in a text string:

```
if (!str[i = strcspn(str, "0123456789+-")])
      printf("can't find number\n");
```

#### See Also

memchr, strchr, strpbrk, strrchr, strspn

#### Notes

strcspn is packaged in the integer library.

## strlen

## **Description**

Find length of a string

## **Syntax**

```
#include <string.h>
unsigned int strlen(char *s)
```

#### **Function**

strlen scans the text string starting at s to determine the number of characters before the terminating NUL.

#### Return Value

The value returned is the number of characters in the string before the NUL.

#### **Notes**

strlen is packaged in the integer library and may be implemented as an inline function.

## strncat

## **Description**

Concatenate strings of length n

## **Syntax**

```
#include <string.h>
char *strncat(char *s1, char *s2, unsigned char n)
```

#### **Function**

strncat appends a copy of the NUL terminated string at s2 to the end of the NUL terminated string at s1. The first character of s2 overlaps the NUL at the end of s1. n specifies the maximum number of characters to be copied, unless the terminating NUL in s2 is encountered first. A terminating NUL is always appended to s1.

#### **Return Value**

strncat returns s1.

## Example

To concatenate the strings "day" and "light":

```
strcpy(s, "day");
strncat(s + 3, "light", 5);
```

#### See Also

strcat

### **Notes**

*strncat* is packaged in the integer library.

# strncmp

## **Description**

Compare two n length strings for lexical order

## **Syntax**

```
#include <string.h>
int strncmp(char *s1, char *s2, unsigned char n)
```

#### **Function**

**strncmp** compares two text strings, character by character, for lexical order in the character collating sequence. The first string starts at **s1**, the second at **s2**. **n** specifies the maximum number of characters to be compared, unless the terminating NUL in **s1** or **s2** is encountered first. The strings must match, including their terminating NUL character, in order for them to be equal.

#### Return Value

strncmp returns an integer greater than, equal to, or less than zero, according to whether s1 is lexicographically greater than, equal to, or less than s2.

## **Example**

To check for a particular error message:

#### See Also

memcmp, strcmp

#### Notes

*strncmp* is packaged in the integer library.

# strncpy

### **Description**

Copy n length string

## **Syntax**

```
#include <string.h>
char *strncpy(char *s1, char *s2, unsigned char n)
```

#### **Function**

strncpy copies the first n characters starting at location s2 into the buffer beginning at s1. n specifies the maximum number of characters to be copied, unless the terminating NUL in s2 is encountered first. In that case, additional NUL padding is appended to s2 to copy a total of n characters.

#### **Return Value**

strncpy returns s1.

## Example

To make a copy of the string **s2** in *dest*:

```
strncpy(dest, s2, n);
```

#### See Also

memcpy, strcpy

#### Notes

If the string s2 points at is longer than n characters, the result may not be NUL-terminated.

*strncpy* is packaged in the integer library.

# strpbrk

## **Description**

Find occurrence in string of character in set

## **Syntax**

```
#include <string.h>
char *strpbrk(char *s1, char *s2)
```

#### **Function**

**strpbrk** scans the NUL terminated string starting at **s1** for the first occurrence of a character in the NUL terminated set **s2**.

#### Return Value

strpbrk returns a pointer to the first character in s1 that is also contained in the set s2, or a NULL if none does.

## **Example**

To replace unprintable characters (as for a 64 character terminal):

```
while (string = strpbrk(string, "`{|}~"))
     *string = '@';
```

#### See Also

memchr, strchr, strcspn, strrchr, strspn

#### Notes

strpbrk is packaged in the integer library.

# strrchr

# **Description**

Scan string for last occurrence of character

# **Syntax**

```
#include <string.h>
char *strrchr(char *s, char c)
```

#### **Function**

strrchr looks for the last occurrence of a specific character c in a NUL terminated string starting at s.

#### Return Value

strrchr returns a pointer to the last character that matches c, or NULL if none does.

# Example

To find a filename within a directory pathname:

```
if (s = strrchr("/usr/lib/libc.user", '/')
      ++s;
```

## See Also

memchr, strchr, strpbrk, strcspn, strspn

#### Notes

strrchr is packaged in the integer library.

# strspn

# **Description**

Find the end of a span of characters not in set

# **Syntax**

```
#include <string.h>
unsigned int strspn(char *s1, char *s2)
```

## **Function**

**strspn** scans the string starting at **s1** for the first occurrence of a character not in the string starting at **s2**. It computes a subscript **i** such that

- s1[i] is a character in the string starting at s1
- s1[i] compares equal to no character in the string starting at s2, except possibly its terminating null character.

## Return Value

strspn returns the lowest possible value of i. s1[i] designates the terminating null character if all of the characters in s1 are in s2.

# **Example**

To check a string for characters other than decimal digits:

```
if (str[strspn(str, "0123456789")])
    printf("invalid number\n");
```

#### See Also

memchr, strcspn, strchr, strpbrk, strrchr

#### Notes

strspn is packaged in the integer library.

# strstr

# **Description**

Scan string for first occurrence of string

# **Syntax**

```
#include <string.h>
char *strstr(char *s1, char *s2)
```

#### Function

strstr looks for the first occurrence of a specific string s2 not including its terminating NUL, in a NUL terminated target string s1.

#### Return Value

strstr returns a pointer to the first character that matches c, or NULL if none does.

# Example

To look for a keyword in a string:

```
if (t = strstr(buf, "LIST"))
      do list(t);
```

## See Also

memchr, strcspn, strpbrk, strrchr, strspn

#### Notes

strstr is packaged in the integer library.

# strtod

# **Description**

Convert buffer to double

# **Syntax**

```
#include <stdlib.h>
double strtod(char *nptr, char **endptr)
```

#### **Function**

**strtod** converts the string at *nptr* into a double. The string is taken as the text representation of a decimal number, with an optional fraction and exponent. Leading whitespace is skipped and an optional sign is permitted; conversion stops on the first unrecognizable character. Acceptable inputs match the pattern:

where **d** is any decimal digit and **e** is the character '**e**' or '**E**'. If *endptr* is not a null pointer, \**endptr* is set to the address of the first unconverted character remaining in the string nptr. No checks are made against overflow, underflow, or invalid character strings.

#### Return Value

*strtod* returns the converted double value. If the string has no recognizable characters, it returns zero.

# Example

To read a string from STDIN and convert it to a double at d:

```
gets(buf);
d = strtod(buf, NULL);
```

#### See Also

atoi, atol, strtol, strtoul

#### Notes

strtod is packaged in the floating point library.

# strtol

## **Description**

Convert buffer to long

# **Syntax**

```
#include <stdlib.h>
long strtol(char *nptr, char **endptr, char base)
```

#### **Function**

**strtol** converts the string at *nptr* into a long integer. Leading whitespace is skipped and an optional sign is permitted; conversion stops on the first unrecognizable character. If base is not zero, characters a-z or A-Z represents digits in range 10-36. If base is zero, a leading "0x" or "0x" in the string indicates hexadecimal, a leading "0" indicates octal, otherwise the string is take as a decimal representation. If base is 16 and a leading "0x" or "0x" is present, it is skipped before to convert. If endptr is not a null pointer, \*endptr is set to the address of the first unconverted character in the string *nptr*.

No checks are made against overflow or invalid character strings.

#### Return Value

strtol returns the converted long integer. If the string has no recognizable characters, zero is returned.

# Example

To read a string from STDIN and convert it to a long I:

```
gets (buf);
l = strtol(buf, NULL, 0);
```

#### See Also

atof, atoi, strtoul, strtod

#### **Notes**

*strtol* is packaged in the integer library.

# strtoul

# **Description**

Convert buffer to unsigned long

# **Syntax**

#### Function

**strtoul** converts the string at *nptr* into a long integer. Leading whitespace is skipped and an optional sign is permitted; conversion stops on the first unrecognizable character. If base is not zero, characters **a-z** or **A-Z** represents digits in range 10-36. If base is zero, a leading "**0x**" or "**0X**" in the string indicates hexadecimal, a leading "**0**" indicates octal, otherwise the string is take as a decimal representation. If base is 16 and a leading "**0x**" or "**0X**" is present, it is skipped before to convert. If *endptr* is not a null pointer, \**endptr* is set to the address of the first unconverted character in the string *nptr*.

No checks are made against overflow or invalid character strings.

#### **Return Value**

*strtoul* returns the converted long integer. If the string has no recognizable characters, zero is returned.

# **Example**

To read a string from STDIN and convert it to a long 1:

```
gets(buf);
l = strtoul(buf, NULL, 0);
```

#### See Also

atof, atoi, strtol, strtod

## Notes

strtoul is a macro redefined to strtol.

# **Description**

Tangent

# **Syntax**

```
#include <math.h>
double tan(double x)
```

#### Function

tan computes the tangent of x, expressed in radians, to full double precision.

#### Return Value

tan returns the nearest internal representation to tan(x), in the range [-pi/2, pi/2], expressed as a double floating value. If the number in x is too large to be represented, tan returns zero. An argument with a large size may return a meaningless value, i.e. when x/(2 \* pi) has no fraction bits.

## Example

To compute the tangent of *theta*:

```
y = tan(theta);
```

#### See Also

cos. sin

#### Notes

tan is packaged in the floating point library.

# tanh

# **Description**

Hyperbolic tangent

## **Syntax**

```
#include <math.h>
double tanh(double x)
```

#### Function

tanh computes the value of the hyperbolic tangent of x to double precision.

#### Return Value

tanh returns the nearest internal representation to tanh(x), expressed as a double floating value. If the result is too large to be properly represented, tanh returns zero.

# **Example**

To compute the hyperbolic tangent of  $\mathbf{x}$ :

```
y = tanh(x);
```

## See Also

cosh, exp, sinh

#### **Notes**

tanh is packaged in the floating point library.

# tolower

# **Description**

Convert character to lowercase if necessary

# **Syntax**

```
#include <ctype.h>
int tolower(char c)
```

#### Function

tolower converts an uppercase letter to its lowercase equivalent, leaving all other characters unmodified.

#### Return Value

tolower returns the corresponding lowercase letter, or the unchanged character.

# Example

To accumulate a hexadecimal digit:

```
for (sum = 0; isxdigit(*s); ++s)
     if (isdigit(*s)
            sum = sum * 16 + *s - '0';
     else
            sum = sum * 16 + tolower(*s) + (10 - 'a');
```

## See Also

toupper

#### Notes

tolower is packaged in the integer library.

# toupper

# **Description**

Convert character to uppercase if necessary

# **Syntax**

```
#include <ctype.h>
int toupper(char c)
```

#### **Function**

**toupper** converts a lowercase letter to its uppercase equivalent, leaving all other characters unmodified.

#### Return Value

toupper returns the corresponding uppercase letter, or the unchanged character.

# **Example**

To convert a character string to uppercase letters:

```
for (i = 0; i < size; ++i)
    buf[i] = toupper(buf[i]);</pre>
```

#### See Also

tolower

#### Notes

toupper is packaged in the integer library.

# **CHAPTER**

# **Using The Assembler**

The castm8 cross assembler translates your assembly language source files into relocatable object files. The C cross compiler calls *castm8* to assemble your code automatically, unless specified otherwise. castm8 generates also listings if requested. This chapter includes the following sections:

- Invoking castm8
- Object File
- Listings
- Assembly Language Syntax
- **Branch Optimization**
- **Old Syntax**
- C Style Directives
- **Assembler Directives**

# **Invoking castm8**

castm8 accepts the following command line options, each of which is described in detail below:

castm8 [op	tions] <files></files>
-a	absolute assembler
-b	do not optimizes branches
-c	output cross reference
-d*>	define symbol=value
+e*	error file name
-ff	use formfeed in listing
-ft	force title in listing
-f#	fill byte value
-h*	include header
-i*>	include path
-1	output listing
+1*	listing file name
-m	accept old syntax
-md	make dependencies
-mi	accept label syntax
-0*	output file name
-pe	all equates public
-p	all symbols public
-pl	keep local symbol
-si	suppress .info. section
-u	undefined in listing
-v	be verbose
-x	include line debug info
-хр	no path in debug info
-xx	include full debug info

# **Castm8 Option Usage**

Option	Description
-a	map all sections to absolute, including the predefined ones.
-b	do not optimize branch instructions. By default, the assembler replaces long branches by short branches wherever a shorter instruction can be used, and short branches by long branches wherever the displacement is too large. This optimization also applies to jump and jump to subroutines instructions.

# Castm8 Option Usage (cont.)

Option	Description
-с	produce cross-reference information. The cross-reference information will be added at the end of the listing file. This option enforces the -I option.
-d*>	where * has the form name=value, defines name to have the value specified by value. This option is equivalent to using an equ directive in each of the source files.
+e*	log errors from assembler in the text file * instead of displaying the messages on the terminal screen.
-ff	use <i>formfeed</i> character to skip pages in listing instead of using blank lines.
-ft	output a title in listing (date, file name, page). By default, no title is output.
-f#	define the value of the filling byte used to fill any gap created by the assembler directives. Default is ${\bf 0}$ .
-h*	include the file specified by * before starting assembly. It is equivalent to an <b>include</b> directive in each source file.
-i*>	define a path to be used by the <b>include</b> directive. Up to 20 paths can be defined. A path is a directory name and is <b>not</b> ended by any directory separator character.
-1	create a listing file. The name of the listing file is derived from the input file name by replacing the suffix by the './s' extension, unless the +I option has been specified.
+ *	create a listing file in the text file *. If both -I and +I are specified, the listing file name is given by the +I option.
-m	accept the old syntax.
-md	create only a list of 'make' compatible dependencies consisting for each source file in the object name followed by a list of included files needed to assemble that file.
-mi	accept label that is not ended with a ':' character.

## Castm8 Option Usage (cont.)

Option	Description
-0*	write object code to the file *. If no file name is specified, the output file name is derived from the input file name, by replacing the rightmost extension in the input file name with the character 'o'. For example, if the input file name is <i>prog.s</i> , the default output file name is <i>prog.o</i> .
-pe	mark all symbols defined by an <b>equ</b> directive as <b>public</b> . This option has the same effect than adding a <b>xdef</b> directive for each of those symbols.
-pl	put locals in the symbol table. They are not published as externals and will be only displayed in the linker map file.
-p	mark all defined symbols as <b>public</b> . This option has the same effect than adding a <b>xdef</b> directive for each label.
-si	suppress the .info. section produced automatically and containing the object name, date and assembler options.
-u	produce an error message in the listing file for all occurrence of an undefined symbol. This option enforces the -l option.
-v	display the name of each file which is processed.
-x	add line debug information to the object file.
-хр	do not prefix filenames in the debug information with any absolute path name. Debuggers will have to be informed about the actual files location.
-xx	add debug information in the object file for any label defining code or data. This option disables the -p option as only public or used labels are selected.

Each source file specified by <files> will be assembled separately, and will produce separate object and listing files. For each source file, if no errors are detected, castm8 generates an object file. If requested by the l or -c options, *castm8* generates a listing file even if errors are detected. Such lines are followed by an error message in the listing.

# **Object File**

The object file produced by the assembler is a relocatable object in a format suitable for the linker clnk. This will normally consist of machine code, initialized data and relocation information. The object file also contains information about the sections used, a symbol table, and a debug symbol table.

The object file contains by default a special section named .info. containing text strings. The assembler produces the object name followed by the current date and time, and the assembler name followed by the specified options. The other compiler passes will complete this information to add their own name and options. This section can be suppressed for compatibility reason by using the -si option.

# Listings

The listing stream contains the source code used as input to the assembler, together with the hexadecimal representation of the corresponding object code and the address for which it was generated. The contents of the listing stream depends on the occurrence of the list, nolist, clist, dlist and mlist directives in the source. The format of the output is as follows:

```
<address> <generated code> <source line>
```

where <address> is the hexadecimal relocatable address where the <source line> has been assembled, <generated code> is the hexadecimal representation of the object code generated by the assembler and <source line> is the original source line input to the assembler. If expansion of data, macros and included files is not enabled, the <generated code> print will not contain a complete listing of all generated code.

Addresses in the listing output are the offsets from the start of the current section. After the linker has been executed, the listing files may be updated to contain absolute information by the clabs utility. Addresses and code will be updated to reflect the actual values as built by the linker

Several directives are available to modify the listing display, such as title for the page header, plen for the page length, page for starting a new page, tabs for the tabulation characters expansion. By default, the listing file is not paginated. Pagination is enabled by using at least one title directive in the source file, or by specifying the -ft option on the command line. Otherwise, the plen and page directives are simply ignored. Some other directives such as clist, mlist or dlist control the amount of information produced in the listing.

A **cross-reference** table will be appended to the listing if the **-c** option has been specified. This table gives for each symbol its value, its attributes, the line number of the line where it has been defined, and the list of line numbers where it is referenced.

# **Assembly Language Syntax**

The assembler *castm8* conforms to the STM8 syntax as described in the document Assembly Language Input Standard. The assembly language consists of lines of text in the form:

```
[label:] [command [operands]] [; comment]
or
      : comment
```

where ':' indicates the end of a label and ';' defines the start of a comment. The end of a line terminates a comment. The *command* field may be an instruction, a directive or a macro call.

Instruction mnemonics and assembler directives may be written in upper or lower case. The C compiler generates lowercase assembly language.

A source file must end with the **end** directive. All the following lines will be ignored by the assembler. If an end directive is found in an included file, it stops only the process for the included file.

#### Instructions

castm8 recognizes the following instructions:

adc	ccf	incw	jrnh	ldf	rim	sra
add	clr	iret	jrnm	ldw	rlc	sraw
addw	clrw	gţ	jrnv	mov	rlcw	srl
and	ср	jpf	jrpl	mul	rlwa	srlw
bccm	cpl	jra	jrsge	neg	rrc	sub
bcp	cplw	jrc	jrsgt	negw	rrcw	subw
bcpl	cpw	jreq	jrsle	nop	rrwa	swap
bkpt	dec	jrf	jrslt	or	rvf	swapw
bres	decw	jrh	jrt	pop	sbc	tnz
bset	div	jrih	jruge	popw	scf	tnzw
btjf	divw	jril	jrugt	push	sim	trap
btjt	exg	jrm	jrule	pushw	sla	wfe
call	exgw	jrmi	jrult	rcf	slaw	wfi
callf	halt	jrnc	jrv	ret	sll	xor
callr	inc	jrne	ld	retf	sllw	

The operand field of an instruction uses an addressing mode to describe the instruction argument. The following examples demonstrate the accepted syntax:

```
rcf
                     ; implicit
push y ld a,#1
                     ; register
                     ; immediate
      a,var
and
                     ; short, long
      a,(2,x)
a,[var]
                     ; indexed
ld
      a,[var] ; indirect
a,([var.w],y) ; indirect indexed long
ld
jrne
      loop
2,#1
                     ; relative
                     ; bit number
bset
btjt
      2,#1,loop ; bit test and branch
```

The assembler chooses the smallest addressing mode where several solutions are possible. Short addressing mode is selected when using a label defined in the .bsct section.

For an exact description of the above instructions, refer to the ST Microelectronics's STM8 Family Programming Manual.

#### Labels

A source line may begin with a label. Some directives require a label on the same line, otherwise this field is optional. A label must begin with an alphabetic character, the underscore character '\_' or the period character '.'. It is continued by alphabetic (A-Z or a-z) or numeric (0-9) characters, underscores, dollar signs (\$) or periods. Labels are case sensitive. The processor register names 'a' and 'x' are reserved and cannot be used as labels

```
data1: dc.b $56
c x: ds.b 1
```

When a label is used within a macro, it may be expanded more than once and in that case, the assembler will fail with a *multiply defined symbol* error. In order to avoid that problem, the special sequence '\@' may be used as a label prefix. This sequence will be replaced by a unique sequence for each macro expansion. This prefix is only allowed inside a macro definition.

```
wait: macro
\@loop:btjf PORTA,#1,\@loop
    endm
```

## **Temporary Labels**

The assembler allows temporary labels to be defined when there is no need to give them an explicit name. Such a label is composed by a decimal number immediately followed by a '\$' character. Such a label is valid until the next standard label or the *local* directive. Then the same temporary label may be redefined without getting a multiply defined error message.

```
1$: dec
    jrne 1$
2$: dec
    jrne 2$
```

Temporary labels do not appear in the symbol table or the cross reference list

For example, to define 3 different local blocks and create and use 3 different local labels named 10\$:

```
function1:
10$: ld a,var
    jreq 10$
    ld a,var2
```

```
local
10$:
     ld a, var2
            10$
     jreq
      1 d
            a,var
      ret.
function2:
10$: ld
            a,var2
      sub
             a,var
      jrne
            10$
      ret.
```

## Constants

The assembler accepts numeric constants and string constants. Numeric constants are expressed in different bases depending on a *prefix* character as follows:

Number	Base
10	decimal (no prefix)
%1010	binary
@12	octal
\$A	hexadecimal

The assembler also accepts numerics constants in different bases depending on a suffix character as follow:

Suffix	Base
D, d or none	decimal (no prefix)
B or b	binary
Q or q	octal
0AH or 0Ah	hexadecimal

The suffix letter can be entered uppercase or lowercase. Hexadecimal numbers still need to start with a digit.

String constants are a series of printable characters between single or double quote characters:

```
'This is a string'
"This is also a string"
```

Depending on the context, a string constant will be seen either as a series of bytes, for a data initialization, or as a numeric; in which case, the string constant should be reduced to only one character.

```
hexa: dc.b '0123456789ABCDEF'
start: cp x,#'A'; ASCII value of 'A'
```

## **Expressions**

An expression consists of a number of labels and constants connected together by operators. Expressions are evaluated to 32-bit precision. Note that operators have the same precedence than in the C language.

A special label written '\*' is used to represent the current location address. Note that when '\*' is used as the operand of an instruction, it has the value of the program counter **before** code generation for that instruction. The set of accepted operators is:

```
addition
        subtraction (negation)
        multiplication
        division
        remainder (modulus)
        bitwise and
ı
        bitwise or
        bitwise exclusive or
       bitwise complement
       left shift
<<
>>
       right shift
==
        equality
! =
        difference
<
        less than
<=
        less than or equal
>
        greater than
>=
        greater than or equal
        logical and
33
        logical or
11
        logical complement
```

These operators may be applied to constants without restrictions, but are restricted when applied to *relocatable* labels. For those labels, the **addition** and **substraction** operators only are accepted and only in the following cases:

```
label + constant
label - constant
label1 - label2
```

#### - NOTE

The difference of two relocatable labels is valid only if both symbols are not external symbols, and are defined in the same section.

An expression may also be constructed with a special operator. These expressions cannot be used with the previous operators and have to be specified alone.

```
high(expression)
                   upper byte
low(expression)
                   lower byte
page(expression)
                   page byte
```

These special operators evaluate an expression and extract the appropriate information from the result. The expression may be relocatable, and may use the set of operators if allowed.

high - extract the upper byte of the 16-bit expression

**low** - extract the lower byte of the 16-bit expression

page - extract the page value of the expression. It is computed by the linker according to the -bs option used. This is used to get the address extension when bank switching is used.

#### **Macro Instructions**

A macro instruction is a list of assembler commands collected under a unique name. This name becomes a new command for the following of the program. A macro begins with a macro directive and ends with a endm directive. All the lines between these two directives are recorded and associated with the macro name specified with the **macro** directive.

```
signex:macro
                   ; sign extension
          Х
     clr
                   ; prepare MSB
            a ; test sign
     jrpl \@pos ; if not positive
                   ; invert MSB
     cpl
\@pos:
      endm
                    ; end of macro
```

This macro is named signex and contains the code needed to perform a sign extension of **a** into **x**. Whenever needed, this macro can be expanded just by using its name in place of a standard instruction:

```
ld a,char+1; load LSB
signex ; expand macro
ld char,x ; store result
```

The resulting code will be the same as if the following code had been written:

```
ld a,char+1; load LSB
clr x ; prepare MSB
tnz a ; test sign
jrpl pos ; if not positive
cpl x ; invert MSB
pos:
ld char,x ; store result
```

A macro may have up to 35 parameters. A parameter is written \1, \2,...\9, \A,...,\Z inside the macro body and refers explicitly to the first, second,... ninth argument and \A to \Z to denote the tenth to 35th operand on the invocation line, which are placed after the macro name, and separated by commas. Each argument replaces each occurrence of its corresponding parameter. An argument may be expressed as a string constant if it contains a comma character.

A macro can also handle named arguments instead of numbered argument. In such a case, the macro directive is followed by a list of argument named, each prefixed by a \ character, and separated by commas. Inside the macro body, arguments will be specified using the same syntax or a sequence starting by a \ character followed by the argument named placed between parenthesis. This alternate syntax is useful to catenate the argument with a text string immediately starting with alphanumeric characters.

The special *parameter* \# is replaced by a numeric value corresponding to the number of *arguments* actually found on the invocation line.

In order to operate directly in memory, the previous macro may have been written using the **numbered** syntax:

```
signex:macro
                    ; sign extension
            Х
      clr
                    ; prepare MSB
      ld
             a,\1
                    ; load LSB
             \@pos ; if not positive
      jrpl
      cpl
                     : invert MSB
\@pos:
      1 d
              \1, x ; store MSB
      endm
                     ; end of macro
```

And called:

```
signex char ; sign extend char
```

This macro may also be written using the **named** syntax:

```
signex:macro
              \value
                             ; sign extension
      clr
                             ; prepare MSB
              Х
      ld
             a,\value
                            ; load LSB
      jrpl
             \@pos
                              ; if not positive
                             ; invert MSB
      cpl
\@pos:
      ld
              \(value),x
                            ; store MSB
      endm
                              ; end of macro
```

The form of a macro call is:

```
name>[.<ext>] [<arguments>]
```

The special parameter  $\setminus 0$  corresponds to an extension  $\langle ext \rangle$  which may follow the macro name, separated by the period character '.'. An extension is a single letter which may represent the size of the operands and the result. For example:

```
table: macro
      dc.\0
            1,2,3,4
      endm
```

When invoking the macro:

table.b

will generate a table of byte:

dc.b 1,2,3,4

When invoking the macro:

```
table.w
```

will generate a table of word:

```
1.2.3.4
dc.w
```

The special parameter \\* is replaced by a sequence containing the list of all the passed arguments separated by commas. This syntax is useful to pass all the macro arguments to another macro or a repeatl directive.

The directive **mexit** may be used at any time to stop the macro expansion. It is generally used in conjunction with a conditional directive.

A macro call may be used within another macro definition, all macros must then be defined before their first call. A macro definition cannot contain another macro definition

If a listing is produced, the macro expansion lines are printed if enabled by the **mlist** directive. If enabled, the invocation line is not printed, and all the expanded lines are printed with all the parameters replaced by their corresponding arguments. Otherwise, the invocation line only is printed.

#### **Conditional Directives**

A conditional directive allows parts of the program to be assembled or not depending on a specific condition expressed in an if directive. The condition is an expression following the if command. The expression cannot be relocatable, and shall evaluate to a numeric result. If the condition is false (expression evaluated to zero), the lines following the if directive are skipped until an endif or else directive. Otherwise, the lines are normally assembled. If an else directive is encountered, the condition status is reversed, and the conditional process continues until the next endif directive.

```
i f
         debug == 1
ld
         x, #message
         print
call
endif
```

If the symbol debug is equal to 1, the next two lines are assembled. Otherwise they are skipped.

```
offset != 1 ; if offset too large
addptr offset
                   ; call a macro
                   ; otherwise
else
inc x
                   ; increment X register
endif
```

If the symbol offset is not one, the macro addptr is expanded with offset as argument, otherwise the aix instruction is directly assembled.

Conditional directives may be nested. An else directive refers to the closest previous if directive, and an endif directive refers to the closest previous **if** or **else** directive.

If a listing is produced, the skipped lines are printed only if enabled by the **clist** directive. Otherwise, only the assembled lines are printed.

#### Sections

The assembler allows code and data to be splitted in **sections**. A section is a set of code or data referenced by a section name, and providing a contiguous block of relocatable information. A section is defined with a section directive, which creates a new section and redirects the following code and data thereto. The directive switch can be used to redirect the following code and data to another section.

```
data: section
                     ; defines data section
text: section
                    ; defines text section
start:
      1 d
            x, #value; fills text section
      jp print
      switch data ; use now data section
value:
      dc.b 1,2,3 ; fills data section
```

The assembler allows up to 255 different sections. A section name is limited to 15 characters. If a section name is too long, it is simply truncated without any error message.

The assembler predefines the following sections, meaning that a section directive is *not* needed before to use them:

Section	Description
.text	executable code
.data	initialized data
.bss	uninitialized data
.bsct	initialized data in short range
.ubsct	non initialized data in short rangeshort range

The sections .bsct and .ubsct are used for locating data in the short range of the processor. The short range is defined as the memory addresses between 0x00 and 0xFF inclusive, i.e. the memory directly addressable by a single byte. Several processors include special instructions and/or addressing modes that take advantage of this special address range. The Cosmic assembler will automatically use the most efficient addressing mode if the data objects are allocated in the .bsct, **.ubsct** or a section with the same attributes. If short range data objects are defined in another file then the directive xref.b must be used to externally reference the data object. This directive specifies that the address for these data object is only one byte and therefore the assembler may use 8 bit addressing modes.

```
switch
                .bsct
zvar2: ds.b
      switch .bss
var2: ds.b
       switch .text
       1 d
               a,var
       1 d
               a,zvar
       1 d
               a, var2
       ld
               a.var2
       end
```

## Bit Handling

The assembler allows symbols specifying bit addresses instead of byte addresses. A bit address is obtained from a byte address and a bit number by or'ing the bit number with the byte address 3-bit shifted to the left. Such symbol can be defined either by an equate definition or as

member of a bit section. Such a section can be defined by using the section directive with the bit attribute. In a bit section, any directive creating or reserving bytes can be used, but will create or reserve bits. Bit symbols can be directly used by the bit instructions with a shortened syntax, as a bit symbol is defining both a byte and a bit in this byte. Bit symbols can be declared as external by using the xbit directive. External definitions for bit symbols located in the short range will used the **xhit.h** directive

```
xbit.b b1
                     : external bit declaration
PA3: equ PORTA:3 ; bit 3 of byte PORTA
.bit: section zpage, bit; create bit section named ".bit"
h0·
     ds.b 1
                 ; allocates one bit
      switch .text
      btjf PA3, skip; use directly bit symbol
      bset
            b0 ; with bit instructions
skip:
      bres
             b1
```

Bit sections are located at link time either at specified bit addresses or attached to any short range section. The linker is computing the proper addresses when hooking bit sections to byte sections, or byte sections to bit sections

#### Includes

The **include** directive specifies a file to be included and assembled in place of the **include** directive. The file name is written between double quotes, and may be any character string describing a file on the host system. If the file cannot be found using the given name, it is searched from all the include paths defined by the -i options on the command line, and from the paths defined by the environment symbol **CXLIB**, if such a symbol has been defined before the assembler invocation. This symbol may contain several paths separated by the usual path separator of the host operating system (';' for MSDOS and ':' for UNIX).

The -h option can specify a file to be "included". The file specified will be included as if the program had an include directive at its very top. The specified file will be included before any source file specified on the command line.

# **Branch Optimization**

Branch instructions are by default automatically optimized to produce the shortest code possible. This behaviour may be disabled by the **-b** option. This optimization operates on conditional branches, on jumps and jumps to subroutine.

A *conditional* branch offset is limited to the range [-128,127]. If such an instruction cannot be encoded properly, the assembler will replace it by a sequence containing an inverted branch to the next location followed immediately by a jump to the original target address. The assembler keep track of the last replacement for each label, so if a long branch has already been expanded for the same label at a location close enough from the current instruction, the target address of the short branch will be changed only to branch on the already existing jump instruction to the specified label.

Note that a *jra* instruction will be replaced by a single *jp* instruction if it cannot be encoded as a relative branch.

A *jp* or *call* instruction will be replaced by a *jra* or *callr* instruction if the destination address is in the same section than the current one, and if the displacement is in the range allowed by a relative branch.

# **Old Syntax**

The -m option allows the assembler to accept old constructs which are now obsolete. The following features are added to the standard behaviour:

- a comment line may begin with a '\*' character;
- a label starting in the first column does not need to be ended with a ':' character;
- no error message is issued if an operand of the dc.b directive is too large;
- the **section** directive handles *numbered* sections;

The comment separator at the end of an instruction is still the ';' character because the '\*' character is interpreted as the multiply operator.

# C Style Directives

The assembler also supports C style directives matching the preprocessor directives of a C compiler. The following directives list shows the equivalence with the standard directives:

C Style	Assembler Style
#include "file"	include "file"
#define label expression	label: equ expression
#define label	label: equ 1
#if expression	if expression
#ifdef label	ifdef label
#ifndef label	ifndef label
#else	else
#endif	endif
#error "message"	fail "message"

#### NOTE

The #define directive does not implement all the text replacement features provided by a C compiler. It can be used only to define a symbol equal to a numerical value.

# **Assembler Directives**

This section consists of quick reference descriptions for each of the castm8 assembler directives.

# align

## **Description**

Align the next instruction on a given boundary

# **Syntax**

```
align <expression>,[<fill_value>]
```

#### **Function**

The **align** directive forces the next instruction to start on a specific boundary. The **align** directive is followed by a constant expression which must be positive. The next instruction will start at the next address which is a multiple of the specified value. If bytes are added in the section, they are set to the value of the filling byte defined by the **-f** option. If **<fill\_value>** is specified, it will be used locally as the filling byte, instead of the one specified by the **-f** option.

# **Example**

```
align 3; next address is multiple of 3 ds.b 1
```

#### See Also

even

# hase

# **Description**

Define the default base for numerical constants

# **Syntax**

```
base <expression>
```

# **Function**

The base directive sets the default base for numerical constants beginning with a digit. The base directive is followed by a constant expression which value must be one of 2, 8, 10 or 16. The decimal base is used by default. When another base is selected, it is no more possible to enter decimal constants.

# **Example**

```
base
                    ; select octal base
ld a,#377
                    ; load $FF
```

# bsct

# **Description**

Switch to the predefined **.bsct** section.

# **Syntax**

bsct

#### **Function**

The bsct directive switches input to a section named .bsct, also known as the short range section. The assembler will automatically select the short addressing mode when referencing an object defined in the .bsct section

# **Example**

#### **Notes**

The .bsct section is limited to 256 bytes, but the assembler does not check the .bsct section size. This will be done by the linker.

#### See Also

section, switch

# clist

# **Description**

Turn listing of conditionally excluded code on or off.

# **Syntax**

clist [on|off]

# **Function**

The **clist** directive controls the output in the listing file of conditionally excluded code. It is effective if and only if listings are requested; it is ignored otherwise.

The parts of the program to be listed are the program lines which are not assembled as a consequence of if, else and endif directives.

## See Also

if, else, endif

# dc

# **Description**

Allocate constant(s)

# **Syntax**

```
dc[.size] <expression>[,<expression>...]
```

#### Function

The **dc** directive allocates and initializes storage for constants. If <*expression>* is a string constant, one byte is allocated for each character of the string. Initialization can be specified for each item by giving a series of values separated by commas or by using a repeat count.

The **dc** and **dc.b** directives will allocate one byte per <*expression*>.

The **dc.w** directive will allocate one word per <*expression*>.

The **dc.l** directive will allocate one long word per *<expression>*.

## Example

```
digit: dc.b 10,'0123456789' dc.w digit
```

#### Note

For compatibility with previous assemblers, the directive **fcb** is alias to **dc.b**, and the directive **fdb** is alias to **dc.w**.

# dch

# **Description**

Allocate constant block

# **Syntax**

```
dcb.<size> <count>,<value>
```

#### Function

The dcb directive allocates a memory block and initializes storage for constants. The size area is the number of the specified value < count > of <size>. The memory area can be initialized with the <value> specified.

The **dcb** and **dcb.b** directives will allocate one **byte** per < count>.

The **dcb.w** directive will allocate one **word** per <*count*>.

The **dcb.l** directive will allocate one **long word** per <*count*>.

# **Example**

```
digit: dcb.b 10,5
                      ; allocate 10 bytes,
                      ; all initialized to 5
```

# dlist

# **Description**

Turn listing of debug directives on or off.

# **Syntax**

dlist [on|off]

## **Function**

The dlist directive controls the visibility of any debug directives in the listing. It is effective if and only if listings are requested; it is ignored otherwise.

Allocate variable(s)

## **Syntax**

```
ds[.size] <space>
```

#### **Function**

The **ds** directive allocates storage space for variables. <*space*> must be an absolute expression. Bytes created are set to the value of the filling byte defined by the -f option.

The **ds** and **ds.b** directives will allocate < space > bytes.

The **ds.w** directive will allocate <*space*> words.

The **ds.l** directive will allocate <*space*> long words.

## **Example**

```
ptlec: ds.b
ptecr: ds.b
chrbuf: ds.w 128
```

#### Note

For compatibility with previous assemblers, the directive **rmb** is alias to ds.b.

## else

## **Description**

Conditional assembly

## **Syntax**

```
if <expression>
instructions
else
instructions
endc
```

#### **Function**

The **else** directive follows an **if** directive to define an alternative conditional sequence. It reverts the condition status for the following instructions up to the next matching **endif** directive. An **else** directive applies to the closest previous **if** directive.

## **Example**

#### Note

The **else** and **elsec** directives are equivalent and may used without distinction. They are provided for compatibility with previous assemblers.

#### See Also

```
if, endif, clist
```

## elsec

### **Description**

Conditional assembly

## **Syntax**

```
if <expression>
instructions
elsec
instructions
```

#### Function

The elsec directive follows an if directive to define an alternative conditional sequence. It reverts the condition status for the following instructions up to the next matching endc directive. An elsec directive applies to the closest previous if directive.

## **Example**

```
ifge
      offset-127 ; if offset too large
addptr offset
                   ; call a macro
elsec
                   ; otherwise
inc x
                   ; increment X register
endc
```

#### Note

The elsec and else directives are equivalent and may used without distinction. They are provided for compatibility with previous assemblers.

#### See Also

```
if, endc, clist, else
```

## end

## **Description**

Stop the assembly

## **Syntax**

end

### **Function**

The **end** directive stops the assembly process. Any statements following it are ignored. If the **end** directive is encountered in an included file, it will stop the assembly process for the included file only.

## endc

## **Description**

End conditional assembly

## **Syntax**

```
if <cc> <expression>
instructions
endc
```

#### Function

The **endc** directive closes an **if**<**cc>** or **elsec** conditional directive. The conditional status reverts to the one existing before entering the if<ce> directives. The endc directive applies to the closest previous if <cc> or elsec directive.

### Example

```
ifae
      offset-127 ; if offset too large
addptr offset
                    ; call a macro
elsec
                   ; otherwise
inc
                   ; increment X register
endc
```

#### Note

The endc and endif directives are equivalent and may used without distinction. They are provided for compatibility with previous assemblers.

#### See Also

```
if < cc>, elsec, clist, end
```

## endif

## **Description**

End conditional assembly

## **Syntax**

```
if <expression>
instructions
endif
```

#### **Function**

The **endif** directive closes an **if**, or **else** conditional directive. The conditional status reverts to the one existing before entering the **if** directive. The **endif** directive applies to the closest previous **if** or **else** directive.

## **Example**

#### Note

The **endif** and **endc** directives are equivalent and may used without distinction. They are provided for compatibility with previous assemblers.

#### See Also

```
if, else, clist
```

## endm

## **Description**

End macro definition

## **Syntax**

```
label: macro
      <macro body>
      endm
```

### **Function**

The **endm** directive is used to terminate macro definitions.

## **Example**

```
; define a macro that places the length of
; a string in a byte prior to the string
ltext:
             macro
             \@2 - \@1
      ds.b
\@1:
      ds.b \1
\@2:
      endm
```

#### See Also

mexit, macro

## endr

## **Description**

End repeat section

## **Syntax**

```
repeat
<repeat_body>
endr
```

### **Function**

The endr directive is used to terminate repeat sections.

## **Example**

```
; shift a value n times
asln: macro
    repeat \1
    sla
    endr
    endm
; use of above macro
    asln 10    ;shift 10 times
```

### See Also

repeat, repeatl, rexit



Give a permanent value to a symbol

## **Syntax**

```
label: equ <expression>
```

#### **Function**

The **equ** directive is used to associate a permanent value to a symbol (label). Symbols declared with the equ directive may not subsequently have their value altered otherwise the **set** directive should be used. <expression> must be either a constant expression, or a relocatable expression involving a symbol declared in the same section as the current one.

The **equ** directive can also be used to define a **bit** symbol by suffixing the defining expression with an absolute bit number. The expression and the bit number are separated by a colon character ':'. The expression can be absolute or relocatable.

## Example

```
false: equ
              0
                     ; initialize these values
true: equ
tablen: equ
            tabfin - tabsta ; compute table length
                     ; define strings for ascii characters
nul: equ
              $0
              $1
soh:
      equ
stx:
      equ
              $2
etx:
      equ
              $3
eot:
      equ
              $4
enq:
      equ
              $5
PORTB: equ
              $1
PB2:
      equ
             PORTB:2
```

#### See Also

lit. set

## even

## **Description**

Assemble next byte at the next even address relative to the start of a section

## **Syntax**

```
even [fill_<value>]
```

#### **Function**

The **even** directive forces the next assembled byte to the next even address. If a byte is added to the section, it is set to the value of the filling byte defined by the **-f** option. If **<fill\_value>** is specified, it will be used locally as the filling byte, instead of the one specified by the **-f** option.

## **Example**

```
vowtab: dc.b 'aeiou' even ; ensure aligned at even address tentab: dc.w 1, 10, 100, 1000
```

## fail

## **Description**

Generate error message.

## **Syntax**

```
fail "string"
```

### Function

The fail directive outputs "string" as an error message. No output file is produced as this directive creates an assembly error. fail is generally used with conditional directives.

## Example

Max:

```
equ
     512
ifge
     value - Max
fail
     "Value too large"
```

## if

### **Description**

Conditional assembly

## **Syntax**

if <expression></expression>	or	if <expression></expression>
instructions		instructions
endif		else
		instructions
		endif

#### **Function**

The **if**, **else** and **endif** directives allow conditional assembly. The **if** directive is followed by a constant expression. If the result of the expression is **not** zero, the following instructions are assembled up to the next matching **endif** or **else** directive; otherwise, the following instructions up to the next matching **endif** or **else** directive are skipped.

If the **if** statement ends with an **else** directive, the expression result is inverted and the same process applies to the following instructions up to the next matching **endif**. So, if the **if** expression was **not** zero, the instructions between **else** and **endif** are skipped; otherwise, the instructions between **else** and **endif** are assembled. An **else** directive applies to the closest previous **if** directive.

The **if** directives may be nested. The skipped lines may or may not be in the listing depending on the **clist** directive status.

## **Example**

#### See Also

else, endif, clist



Conditional assembly

### **Syntax**

```
ifc <string1>,<string2> orifc <string1>,<string2>
instructions
                           instructions
                           elsec
                           instructions
                           enda
```

#### **Function**

The ifc, else and endc directives allow conditional assembly. The ifc directive is followed by a constant expression. If <string1> and <string2> are equals, the following instructions are assembled up to the next matching endc or elsec directive; otherwise, the following instructions up to the next matching endc or elsec directive are skipped.

If the **ifc** statement ends with an **elsec** directive, the expression result is inverted and the same process applies to the following instructions up to the next matching endc. So, if the ifc expression was not zero, the instructions between **elsec** and **endc** are skipped; otherwise, the instructions between elsec and endc are assembled. An elsec directive applies to the closest previous if directive.

The **if** directives may be nested. The skipped lines may or may not be in the listing depending on the **clist** directive status.

## Example

```
"hello", \2
ifc
                        ; if "hello" equals argument
       a,#45
                         ; load 45
ld
elsec
                         ; otherwise ...
       a,#0
1 d
endc
```

### See Also

## ifdef

## **Description**

Conditional assembly

## **Syntax**

ifdef <label></label>	or	ifdef <label></label>	
instructions		instructions	
endc		elsec	
		instructions	
		endc	

#### **Function**

The **ifdef**, **elsec** and **endc** directives allow conditional assembly. The **ifdef** directive is followed by a label < *label*>. If < *label*> is defined, the following instructions are assembled up to the next matching **endc** or **elsec** directive; otherwise, the following instructions up to the next matching **endc** or **elsec** directive are skipped. < *label*> must be first defined. It cannot be a forward reference.

If the **ifdef** statement ends with an **elsec** directive, the expression result is inverted and the same process applies to the following instructions up to the next matching **endif**. So, if the **ifdef** expression was **not** zero, the instructions between **elsec** and **endc** are skipped; otherwise, the instructions between **elsec** and **endc** are assembled. An **elsec** directive applies to the closest previous **if** directive.

The **if** directives may be nested. The skipped lines may or may not be in the listing depending on the **clist** directive status.

## **Example**

```
ifdef offset1 ; if offset1 is defined
addptr offset1 ; call a macro
elsec ; otherwise
addptr offset2 ; call a macro
endif
```

#### See Also

ifndef, elsec, endc, clist



Conditional assembly

## **Syntax**

<pre>ifeq <expression> instructions endc</expression></pre>	or	<pre>ifeq <expression> instructions elsec</expression></pre>
		instructions
		endc

#### **Function**

The **ifeq**, **elsec** and **endc** directives allow conditional assembly. The ifeq directive is followed by a constant expression. If the result of the expression is equal to zero, the following instructions are assembled up to the next matching endc or elsec directive; otherwise, the following instructions up to the next matching **endc** or **elsec** directive are skipped.

If the **ifeq** statement ends with an **elsec** directive, the expression result is inverted and the same process applies to the following instructions up to the next matching endc. So, if the ifeq expression is equal to zero, the instructions between elsec and endc are skipped; otherwise, the instructions between elsec and endc are assembled. An elsec directive applies to the closest previous if directive.

The **if** directives may be nested. The skipped lines may or may not be in the listing depending on the **clist** directive status.

## Example

```
offset
ifea
                      ; if offset nul
tnz
                      ; just test it
elsec
                      ; otherwise
     a,#offset
add
                      ; add to accu
endc
```

#### See Also

# ifge

### **Description**

Conditional assembly

## **Syntax**

ifge <expression></expression>	or	ifge <expression></expression>
instructions		instructions
endc		elsec
		instructions
		endc

#### **Function**

The **ifge**, **elsec** and **endc** directives allow conditional assembly. The **ifge** directive is followed by a constant expression. If the result of the expression is **greater or equal** to zero, the following instructions are assembled up to the next matching **endc** or **elsec** directive; otherwise, the following instructions up to the next matching **endc** or **elsec** directive are skipped.

If the **ifge** statement ends with an **elsec** directive, the expression result is inverted and the same process applies to the following instructions up to the next matching **endc**. So, if the **ifge** expression is **greater or equal** to zero, the instructions between **elsec** and **endc** are skipped; otherwise, the instructions between **elsec** and **endc** are assembled. An **elsec** directive applies to the closest previous **if** directive.

The **if** directives may be nested. The skipped lines may or may not be in the listing depending on the **clist** directive status.

## Example

```
ifge offset-127 ; if offset too large
addptr offset ; call a macro
elsec ; otherwise
inc x ; increment X register
endc
```

#### See Also

elsec. endc. clist



Conditional assembly

## **Syntax**

<pre>ifgt <expression> instructions endc</expression></pre>	or	<pre>ifgt <expression> instructions elsec</expression></pre>
		instructions endc

#### Function

The **ifgt**, **elsec** and **endc** directives allow conditional assembly. The **ifgt** directive is followed by a constant expression. If the result of the expression is greater than zero, the following instructions are assembled up to the next matching endc or elsec directive; otherwise, the following instructions up to the next matching endc or elsec directive are skipped.

If the **ifgt** statement ends with an **elsec** directive, the expression result is inverted and the same process applies to the following instructions up to the next matching endc. So, if the ifgt expression was greater than zero, the instructions between **elsec** and **endc** are skipped; otherwise, the instructions between elsec and endc are assembled. An elsec directive applies to the closest previous **if** directive.

The **if** directives may be nested. The skipped lines may or may not be in the listing depending on the **clist** directive status.

### Example

```
ifgt
        offset-127
                      ; if offset too large
addptr offset
                      ; call a macro
elsec
                       ; otherwise
                       ; increment X register
inc
        Х
endc
```

#### See Also

## ifle

### **Description**

Conditional assembly

## **Syntax**

ifle <expression></expression>	or	ifle <expression></expression>
instructions		instructions
endc		elsec
		instructions
		endc

#### **Function**

The **ifle**, **elsec** and **endc** directives allow conditional assembly. The **ifle** directive is followed by a constant expression. If the result of the expression is **less or equal** to zero, the following instructions are assembled up to the next matching **endc** or **elsec** directive; otherwise, the following instructions up to the next matching **endc** or **elsec** directive are skipped.

If the **ifle** statement ends with an **elsec** directive, the expression result is inverted and the same process applies to the following instructions up to the next matching **endc.** So, if the **ifle** expression was **less or equal** to zero, the instructions between **elsec** and **endc** are skipped; otherwise, the instructions between **elsec** and **endc** are assembled. An **elsec** directive applies to the closest previous **if** directive.

The **if** directives may be nested. The skipped lines may or may not be in the listing depending on the **clist** directive status.

## **Example**

#### See Also



Conditional assembly

## **Syntax**

<pre>iflt <expression> instructions endc</expression></pre>	or	<pre>iflt <expression> instructions elsec instructions</expression></pre>
		instructions
		endc

#### Function

The **iflt**, **else** and **endc** directives allow conditional assembly. The **iflt** directive is followed by a constant expression. If the result of the expression is less than zero, the following instructions are assembled up to the next matching endc or elsec directive; otherwise, the following instructions up to the next matching endc or elsec directive are skipped.

If the **iflt** statement ends with an **elsec** directive, the expression result is inverted and the same process applies to the following instructions up to the next matching endc. So, if the iflt expression was less than zero, the instructions between **elsec** and **endc** are skipped; otherwise, the instructions between elsec and endc are assembled. An elsec directive applies to the closest previous if directive.

The **if** directives may be nested. The skipped lines may or may not be in the listing depending on the **clist** directive status.

## **Example**

```
i.flt
      offset-127
                     ; if offset small enough
                      ; increment X register
inc
elsec
                      ; otherwise
addptr offset
                      ; call a macro
endc
```

#### See Also

## ifndef

#### Description

Conditional assembly

### **Syntax**

ifndef <label></label>	or	ifndef <label></label>
instructions		instructions
endc		elsec
		instructions
		endc

#### **Function**

The **ifndef**, **else** and **endc** directives allow conditional assembly. The ifndef directive is followed by a label < label>. If < label> is not defined, the following instructions are assembled up to the next matching endc or elsec directive; otherwise, the following instructions up to the next matching **endc** or **elsec** directive are skipped. < label> must be first defined. It cannot be a forward reference

If the **ifndef** statement ends with an **elsec** directive, the expression result is inverted and the same process applies to the following instructions up to the next matching endif. So, if the ifndef expression was not zero, the instructions between **elsec** and **endc** are skipped; otherwise, the instructions between elsec and endc are assembled. An elsec directive applies to the closest previous if directive.

The **if** directives may be nested. The skipped lines may or may not be in the listing depending on the **clist** directive status.

## **Example**

```
; if offset1 is not defined
        offset1
ifndef
addptr offset2
                        ; call a macro
elsec
                        ; otherwise
addptr offset1
                        ; call a macro
endif
```

#### See Also

ifdef, elsec, endc, clist



Conditional assembly

## **Syntax**

ifne <expression></expression>	or	ifne <expression></expression>
instructions		instructions
endc		elsec
		instructions
		endc

#### Function

The **ifne**, **elsec** and **endc** directives allow conditional assembly. The ifne directive is followed by a constant expression. If the result of the expression is not equal to zero, the following instructions are assembled up to the next matching endc or elsec directive; otherwise, the following instructions up to the next matching endc or elsec directive are skipped.

If the **ifne** statement ends with an **elsec** directive, the expression result is inverted and the same process applies to the following instructions up to the next matching endc. So, if the ifne expression was not equal to zero, the instructions between **elsec** and **endc** are skipped; otherwise, the instructions between elsec and endc are assembled. An elsec directive applies to the closest previous if directive.

The **if** directives may be nested. The skipped lines may or may not be in the listing depending on the **clist** directive status.

## **Example**

```
ifne
       offset
                        ; if offset not nul
       a,#offset
add
                        ; add to accu
elsec
                        ; otherwise
                        ; just test it
tnz
endc
```

#### See Also

## ifnc

### Description

Conditional assembly

## **Syntax**

```
ifnc <string1>,string2> or ifnc <string1><string2>
instructions
                           instructions
endc
                           elsec
                           instructions
                           endc
```

#### Function

The **ifnc**, **elsec** and **endc** directives allow conditional assembly. The ifnc directive is followed by a constant expression. If <string1> and <string2> are different, the following instructions are assembled up to the next matching endc or elsec directive; otherwise, the following instructions up to the next matching **endc** or **elsec** directive are skipped.

If the **ifnc** statement ends with an **elsec** directive, the expression result is inverted and the same process applies to the following instructions up to the next matching endc. So, if the ifnc expression was not zero, the instructions between **elsec** and **endc** are skipped; otherwise, the instructions between **elsec** and **endc** are assembled. An **elsec** directive applies to the closest previous if directive.

The **if** directives may be nested. The skipped lines may or may not be in the listing depending on the **clist** directive status.

## **Example**

```
ifnc
         "hello", \2
addptr offset
                        ; call a macro
else
                        ; otherwise
inc
                        ; increment X register
endif
```

#### See Also

## include

## **Description**

Include text from another text file

### **Syntax**

```
include "filename"
```

#### **Function**

The **include** directive causes the assembler to switch its input to the specified *filename* until end of file is reached, at which point the assembler resumes input from the line following the **include** directive in the current file. The directive is followed by a string which gives the name of the file to be included. This string must match exactly the name and extension of the file to be included; the host system convention for uppercase/lowercase characters should be respected.

## **Example**

```
include "datstr" ; use data structure library
include "bldstd" ; use current build standard
include "matmac" ; use maths macros
include "ports82"; use ports definition
```

## list

## **Description**

Turn on listing during assembly.

## **Syntax**

```
list
```

### **Function**

The **list** directive controls the parts of the program which will be written to the listing file. It is effective if and only if listings are requested; it is ignored otherwise.

## **Example**

```
list ; expand source code until end or nolist dc.b 1,2,4,8,16 end
```

### See Also

nolist



Give a text equivalent to a symbol

## **Syntax**

```
label: lit "string"
```

## **Function**

The lit directive is used to associate a text string to a symbol (label). This symbol is replaced by the string content when parsed in any assembler instruction or directive.

### **Example**

```
"#5"
nbr:
     lit
     ld a, nbr
                    ; expand as 'ld a,#5'
```

### See Also

equ, set

## local

## **Description**

Create a new local block

## **Syntax**

```
local
```

### **Function**

The **local** directive is used to create a new local block. When the *local* directive is used, all temporary labels defined before the local directive will be undefined after the local label. New local labels can then be defined in the new local block. Local labels can only be referenced within their own local block. A local label block is the area between two standard labels or local directives or a combination of the two.

### **Example**

```
ds.b
                1
var2: ds.b
                1
function1:
10$:
       ld
                a, var
                10$
       jreq
       ld
                var2,a
local
10$:
       ld
                a.var2
       jreq
                10$
       ld
                var,a
       ret
```

## macro

### **Description**

Define a macro

## **Syntax**

```
label: macro
     <macro body>
     endm
```

#### Function

The **macro** directive is used to define a macro. The name may be any previously unused name, a name already used as a macro, or an instruction mnemonic for the microprocessor.

Macros are expanded when the name of a previously defined macro is encountered. Operands, where given, follow the name and are separated from each other by commas.

The *<argument list>* is optional and, if specified, is declaring each argument by name. Each argument name is prefixed by a \ character, and separated from any other name by a comma. An argument name is an identifier which may contain . and characters.

The *macro body* consists of a sequence of instructions not including the directives macro or endm. It may contain macro variables which will be replaced, when the macro is expanded, by the corresponding operands following the macro invocation. These macro variables take the form \1 to \9 to denote the first to ninth operand respectively and \A been defined without any < argument list>. Otherwise, macro variables are denoted by their name prefixed by a \ character. The macro variable name can also be enclosed by parenthesis to avoid unwanted concatenation with the remaining text. In addition, the macro variable \# contains the number of actual operands for a macro invocation.

The special parameter \\* is expanded to the full list of passed arguments separated by commas.

The special parameter  $\setminus 0$  corresponds to an extension  $\langle ext \rangle$  which may follow the macro name, separated by the period character '.'. For more information, see "<u>Macro Instructions</u>" on page 187.

A macro expansion may be terminated early by using the **mexit** directive which, when encountered, acts as if the end of the macro has been reached

The sequence '\@' may be inserted in a label in order to allow a unique name expansion. The sequence '\@' will be replaced by a unique number.

A macro can not be defined within another macro.

### **Example**

```
; define a macro that places the length of a string
; in a byte in front of the string using numbered syntax
ltext: macro
      dc.b \@2-\@1
\@1:
      dc.b \1; text given as first operand
\@2:
      endm
; define a macro that places the length of a string
; in a byte in front of the string using named syntax
ltext: macro
              \string
              \@2-\@1
      dc.b
\@1:
      dc.b \string; text given as first operand
\@2:
      endm
```

#### See Also

endm. mexit

## messg

## **Description**

Send a message out to STDOUT

## **Syntax**

```
messg "<text>"
       '<text>'
messq
```

### **Function**

The **messg** directive is used to send a message out to the host system's standard output (STDOUT).

## **Example**

```
messg "Test code for debug"
       ld a,_#2
ld _SCR,a
```

### See Also

title

## mexit

**Description**Terminate a macro definition

## **Syntax**

```
mexit
```

### **Function**

The mexit directive is used to exit from a macro definition before the endm directive is reached. mexit is usually placed after a conditional assembly directive.

## **Example**

```
ctrace:macro
       if tflag == 0
                mexit
       endif
       qŗ
                \1
       endm
```

## See Also

endm, macro

## mlist

## **Description**

Turn on or off listing of macro expansion.

## **Syntax**

mlist [on|off]

### **Function**

The **mlist** directive controls the parts of the program which will be written to the listing file produced by a macro expansion. It is effective if and only if listings are requested; it is ignored otherwise.

The parts of the program to be listed are the lines which are assembled in a macro expansion.

### See Also

macro

## nolist

## **Description**

Turn off listing.

## **Syntax**

nolist

## **Function**

The **nolist** directive controls the parts of the program which will be **not** written to the listing file until an **end** or a **list** directive is encountered. It is effective if and only if listings are requested; it is ignored otherwise.

### See Also

list

#### Note

For compatibility with previous assemblers, the directive **nol** is alias to **nolist**.

## nopage

## **Description**

Disable pagination in the listing file

## **Syntax**

```
nopage
```

### **Function**

The nopage directive stops the pagination mechanism in the listing output. It is ignored if no listing has been required.

## **Example**

```
xref mult, div
nopage
ds.b charin, charout
ds.w a, b, sum
```

#### See Also

plen, title

## offset

## **Description**

Creates absolute symbols

## **Syntax**

```
offset <expresion>
```

#### **Function**

The *offset* directive starts an absolute section which will only be used to define symbols, and not to produce any code or data. This section starts at the address specified by *<expression>*, and remains active while no directive or instructions producing code or data is entered. This absolute section is then destroyed and the current section is restored to the one which was active when the *offset* directive has been entered. All the labels defined is this section become absolute symbols.

<expression> must be a valid absolute expression. It must not contain any forward or external references.

## **Example**

```
offset 0
next:
    ds.b 2
buffer:
    ds.b 80
    switch .text
size:
    ld x,next ; ends the offset section
```

## org

## **Description**

Sets the location counter to an offset from the beginning of a section.

## **Syntax**

org <expresion>

### Function

<expression> must be a valid absolute expression. It must not contain any forward or external references.

For an absolute section, the first *org* before any code or data defines the starting address.

An org directive cannot define an address smaller than the location counter of the current section.

Any gap created by an org directive is filled with the byte defined by the **-f** option.

## page

## **Description**

Start a new page in the listing file

## **Syntax**

page

#### **Function**

The **page** directive causes a formfeed to be inserted in the listing output if pagination is enabled by either a **title** directive or the **-ft** option.

## **Example**

```
xref mult, div
page
ds.b charin, charout
ds.w a, b, sum
```

## See Also

plen, title

# plen

# **Description**

Specify the number of lines per pages in the listing file

# **Syntax**

```
plen <page length>
```

#### **Function**

The **plen** directive causes < page length > lines to be output per page in the listing output if pagination is enabled by either a title directive or the -ft option. If the number of lines already output on the current page is less than <page length>, then the new page length becomes effective with <page length>. If the number of lines already output on the current page is greater than or equal to <page length>, a new page will be started and the new page length is set to <page length>.

# **Example**

plen 58

#### See Also

page, title

# repeat

# **Description**

Repeat a list of lines a number of times

# **Syntax**

```
repeat <expression>
repeat_body
endr
```

#### **Function**

The repeat directive is used to cause the assembler to repeat the following list of source line up to the next endr directive. The number of times the source lines will be repeated is specified by the expression operand. The repeat directive is equivalent to a macro definition followed by the same number of calls on that macro.

# **Example**

```
; shift a value n times
asln: macro
    repeat \1
    sla (x)
    endr
    endm

; use of above macro
    asln 5
```

#### See Also

endr, repeatl, rexit

# repeatl

### **Description**

Repeat a list of lines a number of times

### **Syntax**

```
repeatl <arguments>
    repeat body
```

#### **Function**

The **repeatl** directive is used to cause the assembler to repeat the following list of source line up to the next endr directive. The number of times the source lines will be repeated is specified by the number of arguments, separated with commas (with a maximum of 36 arguments) and executed each time with the value of an argument. The repeatl directive is equivalent to a macro definition followed by the same number of calls on that macro with each time a different argument. The repeat argument is denoted \1 unless the argument list is starting by a name prefixed by a \ character. In such a case, the repeat argument is specified by its name prefixed by a \ character.

A repeatl directive may be terminated early by using the rexit directive which, when encountered, acts as if the end of the repeatl has been reached

### Example

```
; test a value using the numbered syntax
      repeatl 1,2,3
           add a, \# \ 1
                            ; add to accu
      endr
      end
or
      ; test a value using the named syntax
      repeatl \count,1,2,3
            add a, #\count ; add to accu
      endr
      end
```

#### will both produce:

```
; test a value
9 0000 ab01 add a,#1
                           ; add to accu
9 0002 ab02 add a,#2
                           ; add to accu
9 0004 ab03 add
                 a,#3
                           ; add to accu
10
            end
```

### See Also

endr, repeat, rexit

# restore

# **Description**

Restore saved section

# **Syntax**

```
restore
```

### **Function**

The **restore** directive is used to restore the last saved section. This is equivalent to a switch to the saved section.

# **Example**

```
switch
             .bss
      ds.b
              1
var:
var2: ds.b
              1
      save
      switch .text
function1:
10$:
      ld
              a,var
      jreq
             10$
      ld
              var2,a
function2:
10$:
              a,var2
      sub
              a,var
              10$
      jrne
      ret
      restore
var3: ds.b 1
var4: ds.b
             1
      switch .text
      ld
              a, var3
              var4,a
      ld
end
```

#### See Also

save, section

# rexit

# **Description**

Terminate a repeat definition

# **Syntax**

```
rexit
```

# **Function**

The **rexit** directive is used to exit from a **repeat** definition before the **endr** directive is reached. *rexit* is usually placed after a conditional assembly directive.

# **Example**

```
; shift a value n times
asln: macro
    repeat \1
    if \1 == 0
        rexit
    endif
    sla
    endr
    endm
; use of above macro
    asln 5
```

# See Also

endr, repeat, repeatl

# save

# **Description**

Save section

# **Syntax**

```
save
```

### **Function**

The save directive is used to save the current section so it may be restored later in the source file.

# **Example**

```
switch
               .bss
      ds.b
               1
var:
var2: ds.b
               1
       save
       switch .text
function1:
10$:
       ld
               a, var
       jreq
              10$
       ld
               var2,a
function2:
10$:
               a,var2
       sub
               a, var
               10$
       jrne
       ret
       restore
var3: ds.b
var4: ds.b
       switch .text
       ld
               a, var3
       ld
               var4,a
end
```

#### See Also

restore, section

# scross

# **Description**

Turn on or off section crossing

# **Syntax**

scross [on|off]

# **Function**

The scross directive controls the branch instructions optimization and forces the usage of **jpf** instruction if scross is set (on) or **jp** instruction otherwise. The assembler starts with scross on by default.

# section

### **Description**

Define a new section

# **Syntax**

```
<section name>: section [<attributes>]
```

#### **Function**

The **section** directive defines a new section, and indicates that the following program is to be assembled into a section named <section name>. The section directive cannot be used to redefine an already existing section. If no name and no attributes are specified to the section, the default is to defined the section as a text section with its same attributes. It is possible to associate *<attributes>* to the new section. An attribute is either the name of an existing section or an attribute keyword. Attributes may be added if prefixed by a '+' character or not prefixed, or deleted if prefixed by a '-' character. Several attributes may be specified separated by commas. Attribute keywords are:

abs	absolute section
bss	bss style section (no data)
hilo	values are stored in descending order of significance
even	enforce even starting address and size
zpage	enforce 8 bit relocation
long	enforce 32 bit relocation
bit	bit section

# Example

```
CODE: section .text ; section of text
lab1: ds.b 5
DATA: section .data ; section of data
lab2: ds.b 6
     switch CODE
lab3: ds.b
     switch DATA
lab4: ds.b
```

This will place **lab1** and then **lab3** into consecutive locations in section CODE and **lab2** and **lab4** in consecutive locations in section DATA

```
.frame: section .bsct.even
```

The *frame* section is declared with same attributes than the *.bsct* section and with the *even* attribute.

```
.bit: section +zpage, +even, -hilo
```

The .bit section is declared using 8 bit relocation, with an even alignment and storing data with an ascending order of significance.

When the -m option is used, the *section* directive also accepts a number as operand. In that case, a labelled directive is considered as a section definition, and an unlabelled directive is considered as a section opening (*switch*).

It is still possible to add attributes after the section number of a section definition line, separated by a comma.

#### See Also

switch, bsct

# set

# **Description**

Give a resetable value to a symbol

# **Syntax**

```
label: set <expression>
```

### Function

The set directive allows a value to be associated with a symbol. Symbols declared with set may be altered by a subsequent set. The equ directive should be used for symbols that will have a constant value. <expression> must be fully defined at the time the equ directive is assembled.

### **Example**

OFST: 10 set

#### See Also

equ, lit

# spc

# **Description**

Insert a number of blank lines before the next statement in the listing file

# **Syntax**

```
spc <num_lines>
```

### **Function**

The **spc** directive causes < *num\_lines*> blank lines to be inserted in the listing output before the next statement.

# **Example**

```
spc 5
title "new file"
```

If listing is requested, 5 blank lines will be inserted, then the title will be output.

### See Also

title

# **switch**

# **Description**

Place code into a section.

# **Syntax**

```
switch <section name>
```

### Function

The switch directive switches output to the section defined with the **section** directive. < section name > is the name of the target section, and has to be already defined. All code and data following the switch directive up to the next section, switch, bsct or end directive are placed in the section <section name>.

# **Example**

	switch	.bss
buffer:	ds.b	512
	xdef	buffer

This will place **buffer** into the .bss section.

#### See Also

section, bsct

# tabs

# **Description**

Specify the number of spaces for a tab character in the listing file

# **Syntax**

### **Function**

The **tabs** directive sets the number of spaces to be substituted to the tab character in the listing output. The minimum value of <*tab\_size*> is 0 and the maximum value is 128.

# **Example**

tabs 6



**Description**Define default header

# **Syntax**

```
title "name"
```

### **Function**

The title directive is used to enable the listing pagination and set the default page header used when a new page is written to the listing output.

# **Example**

title "My Application"

### See Also

page, plen

#### Note

For compatibility with previous assemblers, the directive ttl is alias to title.

# **xbit**

# **Description**

Declare bit symbol as being defined elsewhere

# **Syntax**

```
xbit[.b] identifier[,identifier...]
```

# **Function**

Visibility of bit symbols between modules is controlled by the xref and xbit directives. Symbols which are defined in other modules must be declared as xbit. A symbol may be declared both xdef and xbit in the same module, to allow for usage of common headers.

The directive xbit.b declares external symbols located in the .bsct section.

# **Example**

```
xbit
        otherprog
xbit.b
                        ; is in .bsct section
        zpage
```

#### See Also

xdef, xref

# xdef

# **Description**

Declare a variable to be visible

# **Syntax**

```
xdef identifier[,identifier...]
```

#### **Function**

Visibility of symbols between modules is controlled by the xdef and **xref** directives. A symbol may only be declared as *xdef* in one module. A symbol may be declared both xdef and xref in the same module, to allow for usage of common headers.

# Example

```
xdef
        sgrt
              ; allow sgrt to be called
                ; from another module
                ; routine to return a square root
                ; of a number >= zero
```

#### See Also

xbit, xref

sqrt:

# xref

# **Description**

Declare symbol as being defined elsewhere

# **Syntax**

```
xref[.b] identifier[,identifier...]
```

# **Function**

Visibility of symbols between modules is controlled by the **xref** and **xdef** directives. Symbols which are defined in other modules must be declared as *xref*. A symbol may be declared both *xdef* and *xref* in the same module, to allow for usage of common headers.

The directive *xref.b* declares external symbols located in the .bsct section.

# **Example**

```
xref otherprog
xref.b zpage ; is in .bsct section
```

#### See Also

xbit, xdef

# **CHAPTER**

6

# **Using The Linker**

This chapter discusses the **clnk** linker and details how it operates. It describes each linker option, and explains how to use the linker's many special features. It also provides example linker command lines that show you how to perform some useful operations. This chapter includes the following sections:

- Introduction
- Overview
- Linker Command File Processing
- · Linker Options
- Section Relocation
- · Setting Bias and Offset
- Linking Objects
- Linking Library Objects
- Automatic Data Initialization
- Moveable Code

- **Checksum Computation**
- DEFs and REFs
- Special Topics
- Description of The Map File
- Linker Command Line Examples

# Introduction

The linker combines relocatable object files, selectively loading from libraries of such files made with *clib*, to create an executable image for standalone execution or for input to other binary reformatters.

*clnk* will also allow the object image that it creates to have local symbol regions, so the same library can be loaded multiple times for different segments, and so that more control is provided over which symbols are exposed. On microcontroller architectures this feature is useful if your executable image must be loaded into several noncontiguous areas in memory.

#### NOTE

The terms "segment" and "section" refer to different entities and are carefully kept distinct throughout this chapter. A "section" is a contiguous subcomponent of an object module that the linker treats as indivisible.

The assembler creates several sections in each object module. The linker combines input sections in various ways, but will not break one up. The linker then maps these combined input sections into output segments in the executable image using the options you specify.

A "segment" is a logically unified block of memory in the executable image. An example is the code segment which contains the executable instructions

For most applications, the "sections" in an object module that the linker accepts as input are equivalent to the "segments" of the executable image that the linker generates as output.

### **Overview**

You use the linker to build your executable program from a variety of modules. These modules can be the output of the C cross compiler, or can be generated from handwritten assembly language code. Some modules can be linked unconditionally, while others can be selected only as needed from function libraries. All input to the linker, regardless of its source, must be reduced to object modules, which are then combined to produce the program file.

The linker can be used to build freestanding programs such as system bootstraps and embedded applications. It can also be used to make object modules that are loaded one place in memory but are designed to execute somewhere else. For example, a data segment in ROM to be copied into RAM at program startup can be linked to run at its actual target memory location. Pointers will be initialized and address references will be in place.

As a side effect of producing files that can be reprocessed, *clnk* retains information in the final program file that can be quite useful. The symbol table, or list of external identifiers, is handy when debugging programs, and the utility *cobj* can be made to produce a readable list of symbols from an object file. Finally, each object module has in its header useful information such as segment sizes.

In most cases, the final program file created by *clnk* is structurally identical to the object module input to *clnk*. The only difference is that the executable file is complete and contains everything that it needs to run. There are a variety of utilities which will take the executable file and convert it to a form required for execution in specific microcontroller environments. The linker itself can perform some conversions, if all that is required is for certain portions of the executable file to be stripped off and for segments to be relocated in a particular way. You can therefore create executable programs using the linker that can be passed directly to a PROM programmer.

#### The linker works as follows:

- Options applying to the linker configuration. These options are referred to in this chapter as "<u>Global Command Line Options</u>" on page 265.
- Command file options apply only to specific sections of the object being built. These options are referred to in this chapter as "<u>Seg-ment Control Options</u>" on page 267.
- Sections can be relocated to execute at arbitrary places in physical memory, or "stacked" on suitable storage boundaries one after the other.
- The final output of the linker is a header, followed by all the segments and the symbol table. There may also be an additional debug symbol table, which contains information used for debugging purposes.

# **Linker Command File Processing**

The command file of the linker is a small control language designed to give the user a great deal of power in directing the actions of the linker. The basic structure of the command file is a series of command items. A command item is either an explicit linker option or the name of an input file (which serves as an implicit directive to link in that file or, if it is a library, scan it and link in any required modules of the library).

An explicit linker option consists of an option keyword followed by any parameters that the option may require. The options fall into five groups:

#### Group 1

(+seg <section>) controls the creation of new segments and has parameters which are selected from the set of local flags.

(**+grp** <*section*>) controls the section grouping.

#### **Group 2**

(+inc\*) is used to include files

#### **Group 3**

(+new, +pub and +pri) controls name regions and takes no parameters.

#### Group 4

(+def <symbol>) is used to define symbols and aliases and takes one required parameter, a string of the form ident1=ident2, a string of the form ident1=constant, or a string of the form ident1=@segment.

### **Group 5**

(+spc <segment>) is used to reserve space in a particular <segment> and has a required parameter

A description of each of these command line options appears below.

The manner in which the linker relocates the various sections is controlled by the +seg option and its parameters. If the size of a current segment is zero when a command to start a new segment of the same name is encountered, it is discarded. Several different sections can be redirected directly to the same segment by using the +grp option.

*clnk* links the *<files>* you specify in order. If a file is a library, it is scanned as long as there are modules to load. Only those library modules that define public symbols for which there are currently outstanding unsatisfied references are included.

#### **Inserting comments in Linker commands**

Each input line may be ended by a comment, which must be prefixed by a # character. If you have to use the # as a significant character, you can escape it, using the syntax \#.

Here is an example for an indirect link file:

```
# Link for EPROM
+seq .text -b0x8100 -n .text # start eprom address
+seg .const -a .text  # constants follow program
+seg .bsct -b 0x0 -m 0x100  # short range start address
+seq .data -b 0x100 -n .data # uninitialized data
+seq .bss -a .data -n .bss # initialized data to 0
\cxstm8\lib\crts.sm8
                            # startup object file
mod1.o mod2.o
                             # input object files
\cxstm8\lib\libisl.sm8  # C library \cxstm8\lib\libm.sm8  # machine library
\cxstm8\lib\libm.sm8
+seg .vector -b0x8000 -0x7f # vectors eprom address
                            # reset and interrupt vectors
vector.o
```

# **Linker Options**

The linker accepts the following options, each of which is described in detail below.

```
clnk [options] <file.lkf> [<files>]
     -bs# bank size
          error file name
     -1*> library path
     -m* map file name
          output file name
     -0*
          phys addr in map
     -p
          symbol table only
     -s
     -sa sort symbol by address
     -si suppress .info. segment
     -sl output local symbols
     -u# display unused symbols
     -37
          verbose
```

The **output file name** and the **link command file must** be present on the command line. The options are described in terms of the two groups listed above; the global options that apply to the linker, and the segment control options that apply only to specific segments.

Global Command Line Options
The global command line options that the linker accepts are:

### **Global linker Options**

Option	Description
-bs#	set the window shift to #, which implies that the number of bytes in a window is <b>2**#</b> . The default value is . For more information, see the section " <u>Address Specification</u> " on page 277.
-e*	log errors in the text file * instead of displaying the messages on the terminal screen.
- *>	specify library path. You can specify up to 20 different paths. Each path is a directory name, <b>not</b> terminated by any directory separator character.
-m*	produce map information for the program being built to file $^{\star}.$
-0*	write output to the file *. This option is required and has no default value.
- <b>p</b>	display symbols with physical address instead of logical address in the map file.
-s	create an output file containing only an absolute symbol table, but still with an object file format. The resulting file can then be used in another link to provide the symbol table of an existing application.
-sa	display symbols sort by address instead of alphabetic order in the map file.
-si	suppress the .info. segment content for compatibility with tools not supporting this segment yet.
-sl	output local symbols in the executable file.

# **Global linker Options**

Option	Descrip	tion
-u#	display unused symbols. Valid options are:	
	-u1	display data symbols
	-u2	display code and constant symbols
	-u4	display absolute symbols (located variables)
	-u8	display symbols defined in the link file
	several care option hat the debug	lues can be combined (added or or'ed) to display ategories. This option has no effect when the debug is been set as symbols are at least referenced in a section. Symbols defined in removed sections are not displayed.
-v	be "verbo	se".

### **Segment Control Options**

This section describes the segment control options that control the structure of individual segments of the output module.

A group of options to control a specific segment must begin with a **+seg** option. Such an option must precede any group of options so that the linker can determine which segment the options that follow apply to. The linker allows up to **255** different segments.

**+seg <section> <options>** start a new segment loading assembler section type <*section>* and build it as directed by the <*options>* that follow:

#### **Segment Control Options Usage**

Option	Description
-a*	make the current segment follow the segment *, where * refers to a segment name given explicitly by a -n option. Options -b, -e and -o cannot be specified if -a has been specified.
-b*	set the physical start address of the segment to *. Option -e or -a cannot be specified if -b has been specified.
-с	do not output any code/data for the segment.
-ck	mark the segment you want to check. For more information, see " <i>Checksum Computation</i> " on page 285.
-ds#	set the bank size for paged addresses calculation. This option overwrites the global <b>-bs</b> option for that segment.
-e*	set the physical end address of the segment to *. Option -b or -a cannot be specified if -e has been specified.
-f#	fill the segment up to the value specified by the -m option with single bytes or two byte words whose value is #. This option has no effect if no -m option is specified for that segment.

# **Segment Control Options Usage (cont.)**

Option	Description	
-i?	define the initialization option. Valid options are:	
	-it use this segment to host the descriptor and images copies of initialized data used for automatic data initialization	
	-id initialize this segment	
	-ib do not initialize this segment	
	-ik mark this segment as checksum segment	
	-ic mark this segment as moveable segment	
-k	mark the segment as a root segment for the unused section suppression. This flags is usually applied on the reset and interrupt vectors section, and as soon as it is specified at least once in the linker command file, enables the section suppression mechanism. This option can be used on any other segment to force the linker to keep it even if it is not used.	
-m*	other segment to force the linker to keep it even if it is not	

# **Segment Control Options Usage (cont.)**

Option	Description
-n*	set the output name of the segment to *. Segment output names have at most 15 characters; longer names are truncated. If no name is given with a -n option, the segment inheritates a default name equal to its assembler section name.  For example, use this option when you want to generate the hex records for a particular PROM, such as:
	+seg .text -b0x2000 -n prom1 <object_files> +seg .text -b0x4000 -n prom2 <object_files></object_files></object_files>
	You can generate the hex records for prom1 by typing:
	chex -n prom1 file.sm8epd
	For more information, see " <i>The chex Utility</i> " in <b>Chapter 8</b> .
-0*	set the logical start address of the segment to * if -b option is specified or the logical end address if -e option is specified. The default is to set the logical address equal to the physical address. Options -b and -e cannot be specified both if -o has been specified.
-r*	round up the starting address of the segment and all the loaded sections. The expression defines the power of two of the alignment value. The option -r3 will align the start address to an 8 bytes boundary. This option has no effect if the start address is explicitly defined by a -b option.
-s*	define a space name for the segment. This segment will be verified for overlapping only against segments defined with the same space name. See "Overlapping Control" on page 277.
-v	do not verify overlapping for the segment.

#### Segment Control Options Usage (cont.)

Option	Description
-W*	set the window size for banked applications, and activate the automatic bank segment creation.
-х	expandable segment. Allow a segment to spill in the next segment of the same section type if its size exceeds the value given by the -m option. The next segment must be declared before the object causing the overflow. This option has no effect if no -m option is specified for the expendable segment. Option -e cannot be specified with option -x.

Options defining a numerical value (addresses and sizes) can be entered as constant, symbols, or simple expression combined them with '+' and '-' operators. Any symbol used has to be defined before to be used, either by a +def directive or loaded as an absolute symbol from a previously loaded object file. The operators are applied from left to right without any priority and parenthesis () are not allowed. Such expressions CANNOT contain any whitespace. For example:

```
+def START=0x1000
+def MAXSIZE=0x2000
+seg .text -bSTART+0x100 -mMAXSIZE-0x100
```

The first line defines the symbol START equals to the absolute value 1000 (hex value), the second line defines the symbol MAXSIZE equals to the absolute value 2000 (hex value). The last line opens a .text segment located at 1100 (hex value) with a maximum size of 1f00 (hex value). For more information, see the section "Symbol Definition Option" on page 274.

Unless -b\* is given to set the *bss* segment start address, the *bss* segment will be made to follow the last *data* segment in the output file. Unless -b\* is given to set the *data* segment start address, the *data* segment will be made to follow the last *bsct* segment in the output file. The *bsct* and *text* segments are set to start at zero unless you specify otherwise by using -b option. It is permissible for all segments to overlap, as far as *clnk* is concerned; the target machine may or may not make sense of this situation (as with separate instruction and data spaces).

#### NOTE

A new segment of the specified type will not actually be created if the last segment of the same name has a size of zero. However, the new options will be processed and will override the previous values.

### **Segment Grouping**

Different sections can be redirected directly to the same segment with the **+grp** directive:

+grp <section>=<section list> where <section> is the name of the target section, and <section list> a list of section names separated by commas. When loading an object file, each section listed in the right part of the declaration will be loaded as if it was named as defined in the left part of the declaration. The target section may be a new section name or the name of an existing section (including the predefined ones). When using a new name, this directive has to be preceded by a matching +seg definition.

#### NOTE

Whitespaces are **not** allowed aside the equal sign '=' and the commas.

# **Linking Files on the Command line**

The linker supports linking objects from the command line. The link command file has to be modified to indicate where the objects are to be loaded using the following @# syntax.

@1, @2,	include each individual object file at its positional location on the command line and insert them at the respective locations in the link file (@1 is the first object file, and so on).
@*	include all of the objects on the command line and insert

them at this location in the link file.

# **Example**

Linking objects from the command line:

```
clnk -o test.sm8 test.lkf file1.o file2.o

## Test.lkf:
+seg .text -b0x5000
+seg .data -b0x100
@1
+seg .text -b0x7000
@2

Is equivalent to

clnk -o test.sm8 test.lkf
## test.lkf
+seg .text -b0x5000
+seg .data -b0x100
file1.o
+seg .text -b0x7000
file2.o
```

# **Include Option**

Subparts of the link command file can be included from other files by using the following option:

+inc\*

include the file specified by \*. This is equivalent to expanding the text file into the link file directly at the location of the **+inc** line.

# Example

Include the file "seg2.txt" in the link file "test.lkf":

```
## Test.lkf:
+seg .text -b0x5000
+seg .data -b0x100
file1.o file2.o
+seg .text -b0x7000
+inc seg2.txt

## seg2.txt:
mod1.o mod2.o mod3.o

## Resultant link file
+seg .text -b0x5000
+seg .data -b0x100
file1.o file2.o
+seg .text -b0x7000
mod1.o mod2.o mod3.o
```

# **Private Region Options**

Options that control code regions are:

+new	start a new region. A "region" is a user definable group of input object modules which may have both public and private portions. The private portions of a region are local to that region and may not access or be accessed by anything outside the region. By default, a new region is given public access.	
+pub	make the following portion of a given region public.	
+pri	make the following portion of a given region private.	

# **Symbol Definition Option**

The option controlling symbol definition and aliases is:

+def\*

define new symbols to the linker. The string \* must be of the form:

ident=constant	where <i>ident</i> is a valid identifier and <i>constant</i> is a valid constant expressed with the standard C language syntax. This form is used to add <i>ident</i> to the symbol table as a defined absolute symbol with a value equal to <i>constant</i> .
Ident=constant:bitnum	where <i>ident</i> is a valid identifier, <i>constant</i> is a valid constant expressed with the standard C language syntax and <i>bitnum</i> a constant expression between 0 and 7. This form is used to add <i>ident</i> to the symbol table as a defined absolute bit symbol with a value equal to <i>constant</i> 3-bit left shifted and or'ed with <i>bitnum</i> .
ident1=ident2	where <i>ident1</i> and <i>ident2</i> are both valid identifiers. This form is used to define aliases. The symbol <i>ident1</i> is defined as the alias for the symbol <i>ident2</i> and goes in the symbol table as an external DEF (a DEF is an entity defined by a given module.) If <i>ident2</i> is not already in the symbol table, it is placed there as a REF (a REF is an entity referred to by a given module).
ident1=ident2:bitnum	where <i>ident1</i> and <i>ident2</i> are both valid identifiers, and <i>bitnum</i> a constant between 0 and 7. This form is used to define bit aliases. The symbol <i>ident1</i> is defined as the alias for the corresponding bit of symbol <i>ident2</i> which cannot be already a bit symbol itself, and goes in the symbol table as an external DEF (a DEF is an entity defined by a given module.) If <i>ident2</i> is not already in the symbol table, it is placed there as a REF (a REF is an entity referred to by a given module).
ident=@section	where <i>ident</i> is a valid identifier, and <i>section</i> is the name of a section specified as the first argument of a +seg directive. This form is used to add <i>ident</i> to the symbol table as a defined symbol whose value is the address of the next byte to be loaded in the specified section.

ident=start(segment)	where <i>segment</i> is the name given to a segment by the -n option. This form is used to add <i>ident</i> to the symbol table as a defined symbol whose value is the <i>logical</i> start address of the designated segment. This directive can be placed anywhere in the link command file, even before the segment is defined.
ident=end(segment)	where <i>segment</i> is the name given to a segment by the <b>-n</b> option. This form is used to add <i>ident</i> to the symbol table as a defined symbol whose value is the <i>logical</i> end address of the designated segment. This directive can be placed anywhere in the link command file, even before the segment is defined.
ident=pstart(segment)	where <i>segment</i> is the name given to a segment by the <b>-n</b> option. This form is used to add <i>ident</i> to the symbol table as a defined symbol whose value is the <i>physical</i> start address of the designated segment. This directive can be placed anywhere in the link command file, even before the segment is defined.
ident=pend(segment)	where <i>segment</i> is the name given to a segment by the <b>-n</b> option. This form is used to add <i>ident</i> to the symbol table as a defined symbol whose value is the <i>physical</i> end address of the designated segment. This directive can be placed anywhere in the link command file, even before the segment is defined.
ident=size(segment)	where <i>segment</i> is the name given to a segment by the <b>-n</b> option. This form is used to add <i>ident</i> to the symbol table as a defined symbol whose value is the <i>size</i> of the designated segment. This directive can be placed anywhere in the link command file, even before the segment is defined.

### - NOTE -

Whitespaces are **not** allowed aside the equal sign '='.

For more information about DEFs and REFs, refer to the section "DEFs and REFs" on page 287.

### **Reserve Space Option**

The following option is used to reserve space in a given segment:

+spc <segment>=<value></value></segment>	reserve <value> bytes of space at the current location in the segment named <segment>.</segment></value>
+spc <segment>=@section</segment>	reserve a space at the current location in the segment named <segment> equal to the current size of the opened segment where the given section is loaded. The size is evaluated at once, so if the reference segment grows after that directive, there is no further modification of the space reservation. If such a directive is used to duplicate an existing section, it has to be placed in the link command file after all the object files.</segment>

### - NOTE -

Whitespaces are **not** allowed aside the equal sign '='.

### **Handle Dependencies**

This directive allows creating or suppressing a dependency between two functions using their assembly level symbol:

+dep <func1>+<func2></func2></func1>	add a dependence marking <func1> as calling <func2>.</func2></func1>	
+dep <func1>-<func2></func2></func1>	suppress a dependence marking <func1> as not calling <func2>.</func2></func1>	

This directive is mostly used to help building complex applications using a static model.

### - NOTE ·

Whitespaces are **not** allowed aside the + and - signs.

### **Section Relocation**

The linker relocates the sections of the input files into the segments of the output file.

An absolute section, by definition, cannot and should not be relocated. The linker will detect any conflicts between the placement of this file and its absolute address given at compile/assemble time.

In the case of a bank switched system, it is still possible for an absolute section to specify a physical address different from the one and at compile/assembly time, the logical address MUST match the one specified at compile/assemble time.

### **Address Specification**

The two most important parameters describing a segment are its **bias** and its **offset**, respectively its physical and logical start addresses. In nonsegmented architectures there is no distinction between *bias* and *offset*. The *bias* is the address of the location in memory where the segment is relocated to run. The *offset* of a segment will be equal to the *bias*. In this case you must set only the *bias*. The linker sets the *offset* automatically.

### **Overlapping Control**

The linker is verifying that a segment does not overlap any other one, by checking the physical addresses (*bias*). This control can be locally disabled for one segment by using the -v option. For targets implementing separated address spaces (such as bank switching), the linker allows several segments to be isolated from the other ones, by giving them a *space* name with the -s option. In such a case, a segment in a named space is checked only against the other segments of the same space. The unnamed segments are checked together.

# **Setting Bias and Offset**

The bias and offset of a segment are controlled by the **-b\*** option and **-o\*** option. The rules for dealing with these options are described below

### **Setting the Bias**

If the -b\* option is specified, the bias is set to the value specified by \*. Otherwise, the bias is set to the end of the last segment of the same name. If the -e\* option is specified, the bias is set to value obtain by subtracting the segment size to the value specified by \*.

### **Setting the Offset**

If the **-0\*** option is specified, the offset is set to the value specified by \*. Otherwise, the offset is set equal to the bias.

### **Using Default Placement**

If none of -b, -e or -o options is specified, the segment may be placed after another one, by using the -a\* option, where \* is the name of another segment. Otherwise, the linker will try to use a default placement based on the segment name. The compiler produces specific sections for code (.text) and data (.data, .bss, .bsct and .ubsct). By default, .text and .bsct segments start at zero, .ubsct segment follows the latest .bsct segment, .data segment follows the latest .text segment, and .bss segment follows the latest .data segment. Note that there is no default placement for the constants segment .const and the bit segment .bit.

### **Bit Segment Handling**

Bit segments are allocated using bit addresses. A bit address is a value based on the byte address and the bit number in this byte. The bit address is equal to the byte address 3-bit left shifted or'ed with the bit number. The bias (or offset) value can be entered directly as a bit address or with a special syntax combining the byte address and the bit number. The following lines are identical:

```
+seg .bit -b 0x103
+seg .bit -b 0x20:3
```

When using the -a option, the linker automatically converts byte address to bit address when entering a bit segment from a byte segment, starting at bit 0, and converts a bit address to a byte address when leaving a bit segment to a byte segment, starting from next available byte.

Bit addresses are displayed in the map file using the combined syntax.

If the bit segment contains initialized variables, the code must be linked with an appropriate startup file and the bit segment must be declared with option -id in the linker command file. Otherwise, the bit segment must be declared with option -c to suppress it from the output file.

# **Linking Objects**

A new segment is built by concatenating the corresponding sections of the input object modules in the order the linker encounters them. As each input section is added to the output segment, it is adjusted to be relocated relative to the end portion of the output segment so far constructed. The first input object module encountered is relocated relative to a value that can be specified to the linker. The size of the output *bss* segment is the sum of the sizes of the input *bss* sections.

Unless the **-v** option has been specified on a segment definition, the linker checks that the segment physical address range does not overlap any other segment of the application. Logical addresses are not checked as bank switching creates several segments starting at the same logical address.

# **Linking Library Objects**

The linker will selectively include modules from a library when outstanding references to member functions are encountered. The library file must be place *after* all objects that may call it's modules to avoid unresolved references. The standard ANSI libraries are provided in two versions to provide the level of support that your application needs. This can save a significant amount of code space and execution time when full ANSI single precision floating point support is not needed. The first letter after "lib" in each library file denotes the library type (**f** for single precision, and **i** for integer). See below.

libf.sm8	Single Precision Library. This library is used for applications where only single precision floating point support is needed. Link this library <i>before</i> the other libraries when <i>only</i> single precision floats are used.
libi.sm8	Integer only Library. This library is designed for applications where <b>no</b> floating point is used. Floats can still be used for arithmetic but not with the standard library. Link this library <b>before</b> the other libraries when only integer libraries are needed.

Memory Model	Machine Library	Integer Only Library	Float Library
Stack Short	lib (0) 0	libis(0).sm8	libfs(0).sm8
Stack Long	libm( <mark>0</mark> ).sm8	libisl(0).sm8	libfsl(0).sm8

### NOTE

When using a model for application smaller than 64K, you must link with the specific set of libraries and startup (names ending with '0').

### **Library Order**

You should link your application with the libraries in the following orders:

Integer Only Application	Single Precision Float Application	
libi.sm8	libf.sm8	
libm.sm8	libi.sm8	
	libm.sm8	

For more information, see "Linker Command Line Examples" on page 295

### **Libraries Setup Search Paths**

The linker uses the environment variable **CXLIB** to search for objects and library files. If you don't specify the full path to the objects and/or libraries in the link command file AND they are not found in the local directory, the linker will then search all paths specified by the CXLIB environment variable. This allows you to specify just the names of the objects and libraries in your link command file. For example, setting the **CXLIB** environment variable to the C:\COSMIC\LIB directory is done as follow:

C>set CXLIB=C:\COSMIC\LIB

## **Automatic Data Initialization**

The linker is able to configure the executable for an automatic data initialization. This mechanism is initiated automatically when the linker finds the symbol <u>idesc</u> in the symbol table, as an *undefined* symbol. *clnk* first locates a segment behind which it will add an image of the data, so called the *host* segment. The default behaviour is to select the **first** *.text* segment in the executable file, but you can override this by marking one segment with the **-it** option.

Then, *clnk* looks in the executable file for initialized segments. All the segments *.data* and *.bsct* are selected by default, unless disabled explicitly by the **-ib** option. Otherwise, renamed segments may also be selected by using the **-id** option. The **-id** option cannot be specified on a bss segment, default or renamed. Once all the selected segments are located, *clnk* builds a descriptor containing the starting address and length of each such segment, and moves the descriptor and the selected segments to the end of the *host* segment, without relocating the content of the selected segments.

For more information, see "<u>Generating Automatic Data Initialization</u>" in **Chapter 2** and "<u>Initializing data in RAM</u>" in **Chapter 3**.

### **Descriptor Format**

The created descriptor has the following format:

The flag byte is used to detect the end of the descriptor, and also to specify a type for the data segment. The actual value is equal to the code of the first significant letter in the segment name.

If the RAM segment has been created using banked addresses (-b and -o values), the RAM start address is described using two words, the first

giving the page value for that segment and the second giving the matching value for the start address in that space. A segment description is displayed as:

```
dc.b flag
                                   ; segment type
dc.w paged_ram_address
dc.w start_ram_address
; paged value of ram start address
; start address of segment in ram
dc.w end prom address
                                  ; address of last data byte
```

The end address in PROM of one segment gives also the starting address in prom of the following segment, if any.

The address of the descriptor will be assigned to the symbol *idesc*, which is used by the *crtsi.s* startup routine. So all this mechanism will be activated just by linking the *crtsi.sm8* file with the application, or by referencing the symbol *idesc* in your own startup file.

If the *host* segment has been opened with a -m option giving a maximum size, clnk will check that there is enough space to move all the selected segments.

### Moveable Code

The linker allows a code segment to be stored in the ROM part, but linked at another address which is supposed to be located in RAM. This feature is specially designed to allow an application to run FLASH programming routines or bootloader from the RAM space. This feature is sharing the same global mechanism than initialized data, and the common descriptor built by the linker contains both record types. The flag byte is used to qualify each entry. In order to implement such a feature, the link command file should contain a dedicated code segment marked with the **-ic** option:

```
LINKER EXAMPLE FOR MOVEABLE CODE
# mark this segment with -ic and link it at RAM address
+seq .text -b 0x100 -n boot -ic
flash.o
+seg .text -b 0x8000 -n code# application code
file.o
. . .
```

The function contained in the object flash.o is now linked at the RAM address 0x100 but stored somewhere in the code space along with any other initialized data. It is not necessary to link the application with the startup routine *crtsi.s* if the application does not contain initialized data but the descriptor will be built as soon as a moveable function is used by the application, but if the *crts.s* startup is used, moveable code segments are **not** copied in RAM at the application start up.

In order to use such a function, it is necessary to first copy it from ROM to RAM. This is done by calling the library function **\_fctcpy()** with one character argument equal to the first significant letter of the moveable segment name. This argument allows an application to implement several different moveable segments for different kind of situations. In such a case, all the moveable segment names should have names with different first character. This function returns a boolean status equal to 0 if no moveable segment has been copied, or a value different of zero otherwise. Once the segment has been successfully copied, the RAM function can be called directly:

There is no possible name conflict between data segment names and moveable code segment names because the linker internally marks the flag byte differently.

# **Checksum Computation**

This feature is activated by the detection of the symbol **ckdesc** as an undefined symbol. This is practically done by calling one of the provided checksum functions, which uses that symbol and returns o if the checksum is correct. These functions are provided in the integer library and are the following:

_checksum()	check a 8 bit checksum stored once for all the selected segments.
_checksumx()	check a 8 bit checksum stored for every selected segments. This method allows a segment to be dynamically reloaded by updating the corresponding CRC byte.
_checksum16()	check a 16 bit checksum stored once for all the selected segments.
_checksum16x()	check a 16 bit checksum stored for every selected segments. This method allows a segment to be dynamically reloaded by updating the corresponding CRC word.

You then have to update the link command file in two ways:

- 1) Mark the segments (usually code segments) you want to check, by using the -ck option on the +seg line. Note that you need only to mark the first segment of a hooked list, meaning that if a segment is declared with -a option as following a segment which is marked with the -ck option, it will automatically inherit the -ck marker and will be also checked. Note also that if you are using the automatic initialization mechanism, and if the code segment hosting the init descriptor (-it) is also marked with -ck, the init segment and ALL the initialization copy segments will also be checked.
- 2) Create an empty segment which will contain the checksum descriptor. This has to be an empty segment, located wherever you want with a -b or -a option. This segment will NOT be checked, even if marked or hooked to a marked segment. The linker will fill this segment with a data descriptor allowing the checking function to scan all the requested segments and compute the final crc. This segment

has to be specially marked with the option -ik to allow the linker to recognize it as the checksum segment.

Here is an example of link command file showing how to use **-ck** and **-ik**:

```
# LINKER EXAMPLE FOR CHECKSUM IMPLEMENTATION

# mark the first segment of an attached list with -ck

# +seg .text -b 0x8000 -n code -ck# this segment is marked
+seg .const -a code -n const# this one is implicitly marked

# create an empty segment for checksum table marked with -ik

# +seg .cksum -a const -n cksum -ik # checksum segment

# remaining part should contain the verification code

# +seg .data -b 0x100
crtsix.sm8
test.o
libi.sm8
libm.sm8
+def __memory=@.bss
```

The descriptor built by the linker is a list of entries followed by the expected CRC value, only once if functions \_checksum() or \_checksum16() are called, or after each entry if functions \_checksumx() or \_checksum16x() are called. An entry contains a flag byte, a start address and an end address. The flag byte is non-zero, and is or'ed with 0x80 if the start address contains a bank value (two words, page first then start address), otherwise it is just one word with the start address. The end address is always one word. The last entry is always followed by a nul byte (seen as an ending flag), and immediately followed by the expected CRC if functions \_checksum() or \_checksum16() are called. The linker compresses the list of entries by creating only one entry for contiguous segments (as long as they are in the same space (-s\* option) and in the same bank/page).

The current linker implements only on algorithm. Starting with zero, the CRC byte/word is first rotated one bit left (a true bit rotation), then xor'ed with the code byte. The CRC values stored in the checksum descriptor are the one's complement value of the expected CRC.

### **DEFs and REFs**

The linker builds a new symbol table based on the symbol tables in the input object modules, but it is not a simple concatenation with adjustments. There are two basic type of symbols that the linker puts into its internal symbol table: REFs and DEFs. DEFs are symbols that are defined in the object module in which they occur. REFs are symbols that are referenced by the object module in which they occur, but are not defined there

The linker also builds a debug symbol table based on the debug symbol tables in any of the input object modules. It builds the debug symbol table by concatenating the debug symbol tables of each input object module in the order it encounters them. If debugging is not enabled for any of input object module, the debug symbol table will be of zero length.

An incoming REF is added to the symbol table as a REF if that symbol is not already entered in the symbol table; otherwise, it is ignored (that reference has already been satisfied by a DEF or the reference has already been noted). An incoming DEF is added to the symbol table as a DEF if that symbol is not already entered in the symbol table; its value is adjusted to reflect how the linker is relocating the input object module in which it occurred. If it is present as a REF, the entry is changed to a DEF and the symbol's adjusted value is entered in the symbol table entry. If it is present as a DEF, an error occurs (multiply defined symbol).

When the linker is processing a library, an object module in the library becomes an input object module to the linker only if it has at least one DEF which satisfies some outstanding REF in the linker's internal symbol table. Thus, the simplest use of clnk is to combine two files and check that no unused references remain

The executable file created by the linker must have no REFs in its symbol table. Otherwise, the linker emits the error message "undefined sym**bol**" and returns failure.

# **Special Topics**

This section explains some special linker capabilities that may have limited applicability for building most kinds of microcontroller applications.

### **Private Name Regions**

Private name regions are used when you wish to link together a group of files and expose only some to the symbol names that they define. This lets you link a larger program in groups without worrying about names intended only for local usage in one group colliding with identical names intended to be local to another group. Private name regions let you keep names truly local, so the problem of name space pollution is much more manageable.

An explicit use for private name regions in a STM8 environment is in building a paged program with duplication of the most used library functions in each page, in order to avoid extra page commutation. To avoid complaints when multiple copies of the same file redefine symbols, each such contribution is placed in a private name region accessible only to other files in the same page.

The basic sequence of commands for each island looks like:

#### +new <public files> +pri <private libraries>

Any symbols defined in <public files> are known outside this private name region. Any symbols defined in <pri>private libraries> are known only within this region; hence they may safely be redefined as private to other regions as well.

### NOTE

All symbols defined in a private region are local symbols and will not appear in the symbol table of the output file.

### **Renaming Symbols**

At times it may be desirable to provide a symbol with an alias and to hide the original name (*i.e.*, to prevent its definition from being used by the linker as a DEF which satisfies REFs to that symbol name). As an

example, suppose that the function *func* in the C library provided with the compiler does not do everything that is desired of it for some special application. There are three methods of handling this situation (we will ignore the alternative of trying to live with the existing function's deficiencies).

The first method is to write a new version of the function that performs as required and link it into the program being built before linking in the libraries. This will cause the new definition of func to satisfy any references to that function, so the linker does not include the version from the library because it is not needed. This method has two major drawbacks: first, a new function must be written and debugged to provide something which basically already exists; second, the details of exactly what the function must do and how it must do it may not be available. thus preventing a proper implementation of the function.

The second approach is to write a new function, say my func, which does the extra processing required and then calls the standard function func. This approach will generally work, unless the original function func is called by other functions in the libraries. In that case, the extra function behavior cannot occur when func is called from library functions, since it is actually my func that performs it.

The third approach is to use the aliasing capabilities of the linker. Like the second method, a new function will be written which performs the new behavior and then calls the old function. The twist is to give the old function a new name and hide its old name. Then the new function is given the old function's name and, when it calls the old function, it uses the new name, or alias, for that function. The following linker script provides a specific example of this technique for the function *func*:

```
line 1 +seg .text -b 0x1000
line 2 +seg .data -b0
line 3 +new
line 4 Crts.xx
line 5 +def oldfunc= func
line 6 +pri func.o
line 7 +new
line 8 prog.o newfunc.o
line 9 <libraries>
```

### NOTE

The function name func as referenced here is the name as seen by the C programmer. The name which is used in the linker for purposes of aliasing is the name as seen at the object module level. For more information on this transformation, see the section "Interfacing C to Assembly Language" in Chapter 3.

The main thing to note here is that *func.o* and *new\_func.o* both define a (different) function named *func*. The second function *func* defined in *newfunc.o* calls the old *func* function by its alias *oldfunc*.

Name regions provide limited scope control for symbol names. The **+new** command starts a new name region, which will be in effect until the next **+new** command. Within a region there are public and private name spaces. These are entered by the **+pub** and **+pri** commands; by default, **+new** starts in the public name space.

**Lines 1,2** are the basic linker commands for setting up a separate I/D program. Note that there may be other options required here, either by the system itself or by the user.

Line 3 starts a new region, initially in the public name space.

Line 4 specifies the startup code for the system being used.

**Line 5** establishes the symbol \_oldfunc as an alias for the symbol \_func. The symbol \_oldfunc is entered in the symbol table as a public definition. The symbol \_func is entered as a private reference in the current region.

**Line 6** switches to the private name space in the current region. Then *func.o* is linked and provides a definition (private, of course) which satisfies the reference to *func*.

**Line 7** starts a new name region, which is in the public name space by default. Now no reference to the symbol *\_func* can reach the definition created on **Line 6**. That definition can only be reached now by using the symbol *\_oldfunc*, which is publicly defined as an alias for it.

**Line 8** links the user program and the module *newfunc.o*, which provides a new (and public) definition of *\_func*. In this module the old version is accessed by its alias. This new version will satisfy all references to *func* made in *prog.o* and the libraries.

### **Line 9** links in the required libraries.

The rules governing which name space a symbol belongs to are as follows:

- Any symbol definition in the public space is public and satisfies all outstanding and future references to that symbol.
- Any symbol definition in the private space of the current region is private and will satisfy any private reference in the current region.
- All private definitions of a symbol must occur before a public definition of that symbol. After a public definition of a symbol, any other definition of that symbol will cause a "multiply defined symbol" error.
- Any number of private definitions are allowed, but each must be in a separate region to prevent a multiply defined symbol error.
- Any new reference is associated with the region in which the reference is made. It can be satisfied by a private definition in that region, or by a public definition. A previous definition of that symbol will satisfy the reference if that definition is public, or if the definition is private and the reference is made in the same region as the definition.
- If a new reference to a symbol occurs, and that symbol still has an
  outstanding unsatisfied reference made in another region, then
  that symbol is marked as requiring a public definition to satisfy it.
- Any definition of a symbol must satisfy all outstanding references to that symbol; therefore, a private definition of a symbol which requires a public definition causes a blocked symbol reference error.

 No symbol reference can "reach" any definition made earlier than the most recent definition.

### **Absolute Symbol Tables**

Absolute Symbol tables are used to export symbols from one application to another, to share common functions for instance, or to use functions already built in a ROM, from an application downloaded into RAM. The linker option -s will modify the output file in order to contain only a symbol table, without any code, but still with an object file format, by using the same command file used to build the application itself. All symbols are flagged as absolute symbols. This file can be used in another link, and will then transmit its symbol table, allowing another application to use those symbols as externals. Note that the linker does not produce any map even if requested, when used with the -s option.

The basic sequence of commands looks like:

```
clnk -o appli.sm8 -m appli.map appli.lkf
clnk -o appli.sym -s appli.lkf
```

The first link builds the application itself using the *appli.lkf* command file. The second link uses the same command file and creates an object file containing only an absolute symbol table. This file can then be used as an input object file in any other link command file.

# **Description of The Map File**

The linker can output a map file by using the -m option. The map file contains 4 sections: the Segment section, the Modules section, the Stack Usage section and the Symbols section.

**Segment** Describe the different segments which compose the application, specifying for each of them: the start address (in hexa), the end address (in hexa), the length (in decimal), and the name of the segment. Note that the end value is the address of the byte following the last one of the segment, meaning that an empty segment will have the same start and end addresses. If a segment is initialized, it is displayed twice, the first time with its final address, the second time with the address of the image copy.

**Modules** List all the modules which compose the application, giving for each the description of all the defined sections with the same format as in the Segment section. If an object has been assembled with the -pl option, local symbols are displayed just after the module description.

**Stack Usage** Describe the amount of memory needed for the stack. Each function of the application is listed by its name, followed by a '>' character indicating that this function is not called by any other one (the main function, interrupt functions, task entries...). The first number is the total size of the stack used by the function including all the internal calls. The second number between braces shows the stack need for that function alone. The entry may be flagged by the keyword "Recursive" meaning that this function is itself recursive or is calling directly or indirectly a recursive function, and that the total stack space displayed is not accurate. The linker may detect potential but not actual recursive functions when such functions are called by pointer. The linker displays at the end of the list a total stack size assuming interrupt functions cannot be themselves interrupted. Interrupt frames and machine library calls are properly counted.

Call Tree List all the functions sorted alphabetically followed by all the functions called inside. The display goes on recursively unless a function has already been listed. In such a case, the name is followed by the line number where the function is expanded. If a line becomes too long, the process is suspended and the line ends with a ... sequence indicating that this function is listed later. Functions called by pointer are listed between parenthesis, or between square brackets if called from an array of pointers.

Symbols List all the symbols defined in the application specifying for each its name, its value, the section where it is defined, and the modules where it is used. If the target processor supports bank switching, addresses are displayed as logical addresses by default. Physical addresses can be displayed by specifying the -p option on the linker command line. Addresses of bit symbols are displayed with the byte address followed by a colon character and the bit number.

### **Return Value**

*clnk* returns success if no error messages are printed to STDOUT; that is, if no undefined symbols remain and if all reads and writes succeed. Otherwise it returns failure.

# **Linker Command Line Examples**

This section shows you how to use the linker to perform some basic operations.

A linker command file consists of linker options, input and output file, and libraries. The options and files are read from a command file by the linker. For example, to create an STM8 file from *file.o* you can type at the system prompt:

```
clnk -o myapp.sm8 myapp.lkf
```

where *myapp.lkf* contains:

```
+seg .text -b0xf000 -n .text # start eprom address
+seg .const -a .text # constants follow program
+seg .bsct -b0x0 -niram -m 0x100# initialized short range
+seg .data -b0x100  # start data address
\cxstm8\lib\crts.sm8  # startup object file
file1.o file2.o  # input object files
                                   # startup object file
\cxstm8\lib\libis.sm8  # C library \cxstm8\lib\libm.sm8  # machine library
+def memory=@.bss
                                   # symbol used by startup
```

The following link command file is an example for an application that does **not** use floating point data types and does **not** require automatic initialization

```
# demo.lkf: link command WITHOUT automatic init.
+seg .text -b 0xf000 -n.text # program start address
+seg .const -a .text # constants follow program
+seg .bsct -b0x0 -niram -m 0x100# initialized short range
+seg .data -b0x100 \# start data address
\cxstm8\lib\crts.sm8
                          # startup with NO-INIT
acia.o
                             # main program
" module program
\cxstm8\lib\libis.sm8  # C library
\cxstm8\lib\libm.sm8  # machine library
+seg .const -b0x8000  # vectors eprom address
vector.o  # defire...
module1.o
                             # module program
                              # reset & interrupt vectors
# define these symbols if crtsi is used
```

The following link command file is an example for an application that uses single precision floating point data types and utilizes automatic data initialization.

```
# demo.lkf: link command WITH automatic init
+seq .text -bf000 0x -n.text # program start address
+seg .const -a .text # constants follow program
+seg .bsct -b0x80 -niram -m 0x80# initialized short range
+seg .ubsct -n iram # uninitialized short range
+seg .data -b0x100 # start data address
\cxstm8\lib\crtsi.sm8 # startup with auto-init
acia.o
                                 # main program
module1.o
                                  # module program
\cxstm8\lib\libfs.sm8
                                # single precision library
\cxstm8\lib\libis.sm8
                                # integer library
                              # machine library
# vectors eprom address
\cxstm8\lib\libm.sm8
+seg .const -b0x8000
vector.o
                                 # reset & interrupt vectors
# define these symbols if crtsi is used
+def __endzp=@.ubsct # end of uninitialized zpage
+def memory=@.bss # end of bss segment
+def memory=@.bss
```

**CHAPTER** 

# **Debugging Support**

This chapter describes the debugging support available with the cross compiler targeting the STM8. There are two levels of debugging support available, so you can use either the COSMIC's Zap C source level cross debugger or your own debugger or in-circuit emulator to debug your application. This chapter includes the following sections:

- Generating Debugging Information
- Generating Line Number Information
- Generating Data Object Information
- The cprd Utility
- The clst utility

# **Generating Debugging Information**

The compiler generates debugging information in response to command line options you pass to the compiler as described below. The compiler can generate the following debugging information:

- 1 line number information that allows COSMIC's C source level debugger or another debugger or emulator to locate the address of the code that a particular C source line (or set of lines) generates. You may put line number information into the object module in either of the two formats, or you can generate both line number information and information about program data and function arguments, as described below
- 2 information about the name, type, storage class and address (absolute or relative to a stack offset) of program static data objects, function arguments, and automatic data objects that functions declare. Information about what source files produced which relocatable or executable files. This information may be localized by address (where the output file resides in memory). It may be written to a file, sorted by address or alphabetical order, or it may be output to a printer in paginated or unpaginated format.

### **Generating Line Number Information**

The compiler puts line number information into a special debug symbol table. The debug symbol table is part of the relocatable object file produced by a compilation. It is also part of the output of the *clnk* linker. You can therefore obtain line number information about a single file, or about all the files making up an executable program. However, the compiler can produce line number information only for files that are fewer than 65,535 lines in length.

### **Generating Data Object Information**

The +debug option directs the compiler to generate information about data objects and function arguments and return types. The debugging information the compiler generates is the information used by the COSMIC's C source level cross debugger or another debugger or emulator. The information produced about data objects includes their name, scope, type and address. The address can be either absolute or relative to a stack offset.

As with line number information alone, you can generate debugging information about a single file or about all the files making up an executable program.

*cprd* may be used to extract the debugging information from files compiled with the **+debug** option, as described below.

# The cprd Utility

**cprd** extracts information about functions and data objects from an object module or executable image that has been compiled with the **+debug** option. *cprd* extracts and prints information on the name, type, storage class and address (absolute or offset) of program static data objects, function arguments, and automatic data objects that functions declare. For automatic data, the address provided is an offset from the frame pointer. For function arguments, the address provided is an offset from the stack pointer.

### **Command Line Options**

cprd accepts the following command line options, each of which is described in detail below:

```
cprd [options] file
-fc* select function name
-fl* select file name
-o* output file name
-r recurse structure fields
-s display object size
-u display unused object
```

where *<file>* is an object file compiled from C source with the compiler command line option **+debug** set.

### **Cprd Option Usage**

Option	Description
-fc*	print debugging information only about the function *. By default, <i>cprd</i> prints debugging information on all functions in <i><file></file></i> . Note that information about global data objects is always displayed when available.
-fl*	print debugging information only about the file *. By default, cprd prints debugging information on all C source files.
-O*	print debugging information to file *. Debugging information is written to your terminal screen by default.
-r	Display structure fields with their offset.

### **Cprd Option Usage (cont.)**

Option	Description
-S	Display object size in bytes.
-u	display only unused global variables.

By default, cprd prints debugging information about all functions and global data objects in < file>.

### **Examples**

The following example show sample output generated by running the cprd utility on an object file created by compiling the program acia.c with the compiler option +debug set.

### cprd acia.sm8

```
Information extracted from acia.sm8
source file acia.c:
unsigned char buffer[64] at 0x0104
unsigned char *ptlec at 0x0102
unsigned char *ptecr at 0x0100
unsigned char getch() lines 26 to 36 at 0x810a-0x8135
    auto unsigned char c at -1 from frame pointer
void outch() lines 40 to 45 at 0x8136-0x8144
    argument unsigned char c at 0 from frame pointer
void recept() lines 51 to 57 at 0x8145-0x8167
    (no locals)
void main() lines 63 to 72 at 0x8168-0x818a
    (no locals)
source file vector.c:
void (* vectab[16])() at 0x8000
```

# The clst utility

The **clst** utility takes relocatable or executable files as arguments, and creates listings showing the C source files that were compiled or linked to obtain those relocatable or executable files. It is a convenient utility for finding where the source statements are implemented.

To use *clst* efficiently, its argument files must have been compiled with the **+debug** option.

*clst* can be instructed to limit its display to files occupying memory in a particular range of addresses, facilitating debugging by excluding extraneous data. *clst* will display the entire content of any files located between the endpoints of its specified address range.

### **Command Line Options**

clst accepts the following command line options, each of which is described in detail below:

```
clst [options> file
    -a list file alphabetically
    -b display physical address
    -f*> process selected file
    -i*> source file
    -l# page length
    -o* output file name
    -p suppress pagination
    -r* specify a line range #:#
```

### **Clst Option Usage**

Option	Description
-a	when set, cause <i>clst</i> to list files in alphabetical order. The default is that they are listed by increasing addresses.
-b	display physical address instead of logical address in the listing file.
-f*>	specify * as the file to be processed. Default is to process all the files of the application. Up to 10 files can be specified.

### Clst Option Usage (cont.)

Option	Description			
-i*>	read string * to locate the source file in a specific directory. Source files will first be searched for in the current directory, then in the specified directories in the order they were given to clst. You can specify up to 20 different paths Each path is a directory name, not terminated by any directory separator character.			
-1#	when paginating output, make the listings # lines long. By default, listings are paginated at 66 lines per page.			
-o*	redirect output from <i>clst</i> to file *. You can achieve a similar effect by redirecting output in the command line.			
	clst -o acia.lst acia.sm8			
	is equivalent to:			
	clst acia.sm8 >acia.lst			
-p	suppress pagination. No page breaks will be output.			
-r#:#	where #:# is a range specification. It must be of the form <number>:<number>. When this flag is specified, only those source files occupying memory in the specified range will be listed. If part of a file occupies memory in the specified range, that file will be listed in its entirety. The following is a valid use of -r:</number></number>			
	-r 0xe000:0xe200			

# **CHAPTER**

# **Programming Support**

This chapter describes each of the programming support utilities packaged with the C cross compiler targeting the STM8. The following utilities are available:

Utility	Description
chex	translate object module format
clabs	generate absolute listings
clib	build and maintains libraries
cobj	examine objects modules
ctat	generate short range variables header file
cvdwarf	generate ELF/DWARF format

The assembler is described in Chapter 5, "Using The Assembler". The linker is described in Chapter 6, "Using The Linker". Support for debugging is described in **Chapter 7**, "Debugging Support".

The description of each utility tells you what tasks it can perform, the command line options it accepts, and how you use it to perform some commonly required operations. At the end of the chapter are a series of examples that show you how to combine the programming support utilities to perform more complex operations.

# The chex Utility

You use the **chex** utility to translate executable images produced by *clnk* to one of several hexadecimal interchange formats. These formats are: *Motorola S-record* format, and *Intel standard hex* format. You can also use *chex* to override text and data biases in an executable image or to output only a portion of the executable.

The executable image is read from the input file <*file*>.

### **Command Line Options**

chex accepts the following command line options, each of which is described in detail below:

chex	[option	ns] file
	-a##	absolute file start address
	-b##	address bias
	-e##	entry point address
	-f?	output format
	-h	suppress header
	+h*	specify header string
	-m#	maximum data bytes per line
	-n*>	output only named segments
	-0*	output file name
	-p	use paged address format
	-pa	use paged address for data
	-pl##	page number for linear mapping
	-pn	use paged address in bank only
	-pp	use paged address with mapping
	-s	output increasing addresses
	-w	output word addresses
	-x*>	exclude named segments

### **Chex Option Usage**

Option	Description
-a##	the argument file is a considered as a pure binary file and ## is the output address of the first byte.
-b##	substract ## to any address before output.

## **Chex Option Usage (cont.)**

Option	Description	
-e##	define ## as the entry point address encoded in the dedicated record of the output format, if available.	
-f?	define output file format. Valid options are:	
	i Intel Hex Format	
	m Motorola S19 format	
	2 Motorola S2 format	
	3 Motorola S3 format	
	Default is to produced Motorola S-Records (-fm). Any other letter will select the default format	
-h	do not output the header sequence if such a sequence exists for the selected format.	
+h*	insert * in the header sequence if such a sequence exists for the selected format.	
-m#	output # maximum data bytes per line. Default is to output 32 bytes per line.	
-n*>	output only segments whose name is equal to the string *. Up to twenty different names may be specified on the command line. If there are several segments with the same name, they will all be produced. This option is used in combination with the -n option of the linker.	
<b>-0</b> *	write output module to file *. The default is STDOUT.	
-р	output addresses of banked segments using a paged for- mat <page_number><logical_address>, instead of the default format <physical>.</physical></logical_address></page_number>	
-ра	output addresses of banked data segments using a paged format <page_number><logical_address>, instead of the default format <physical>.</physical></logical_address></page_number>	

### **Chex Option Usage (cont.)**

Option	Description
-pl##	specify the page value of the segment localized between 0x8000 and 0xc000 when using a linear non-banked application. This option enforces a paged format for this segment.
-pn	behaves as <b>-p</b> but only when logical address is inside the banked area. This option has to be selected when producing an hex file for the Noral debugger.
-рр	behaves as -p but uses paged addresses for all banked segments, mapped or unmapped. This option has to be selected when producing an hex file for Promic tools.
-s	sort the output addresses in increasing order.
-w	output word addresses. Addresses must be aligned on even addresses. This option is useful for word processor type.
-x*>	do not output segments whose name is equal to the string *. Up to twenty different names may be specified on the command line. If there are several segments with the same name, they will not all be output.

### **Return Status**

*chex* returns success if no error messages are printed; that is, if all records are valid and all reads and writes succeed. Otherwise it returns failure

### **Examples**

The file *hello.c*, consisting of:

<pre>char *p = {"hello world"};</pre>	
---------------------------------------	--

when compiled produces the following the following *Motorola* S-record format:

chex hello.o		

S00A000068656C6C6F2E6F44 S1110000020068656C6C6F20776F726C640090 S9030000FC

### and the following Intel standard hex format:

### chex -fi hello.o

:0E000000020068656C6C6F20776F726C640094 :0000001FF

# The clabs Utility

**clabs** processes assembler listing files with the associated executable file to produce listing with updated code and address values.

clabs decodes an executable file to retrieve the list of all the files which have been used to create the executable. For each of these files, clabs looks for a matching listing file produced by the compiler (".ls" file). If such a file exists, clabs creates a new listing file (".la" file) with absolute addresses and code, extracted from the executable file.

To be able to produce any results, the compiler **must** have been used with the '-l' option.

### **Command Line Options**

clabs accepts the following command line options, each of which is described in detail below.

clabs [opti	ons] file
-a	process also library files
-c1*	listings files
-1	restrict to local directory
-p	use paged address format
-pn	use paged address in bank only
-pp	use paged address with mapping
-r*	relocatable listing suffix
-s*	absolute listing suffix
-v	echo processed file names

### **Clabs Option Usage**

Option	Description
-a	process also files located in libraries. Default is to process only all the files of the application.
-cl*	specify a path for the listing files. By default, listings are created in the same directory than the source files.
-1	process files in the current directory only. Default is to process all the files of the application.

#### Clabs Option Usage (cont.)

Option	Description
-р	output addresses of banked segments using a paged for- mat <page_number><logical_address>, instead of the default format <physical>.</physical></logical_address></page_number>
-pn	behaves as -p but only when logical address is inside the banked area.
-рр	behaves as <b>-p</b> but uses paged addresses for all banked segments, mapped or unmapped.
-r*	specify the input suffix, including or not the dot '.' character. Default is ".ls"
-s*	specify the output suffix, including or not the dot '.' character. Default is ".la"
-v	be verbose. The name of each module of the application is output to STDOUT.

<sup>&</sup>lt;file> specifies one file, which must be in executable format.

# **Return Status**

clabs returns success if no error messages are printed; that is, if all reads and writes succeed. An error message is output if no relocatable listing files are found. Otherwise it returns failure.

# **Examples**

The following command line:

```
clabs -v acia.sm8
```

#### will output:

crts.ls acia.ls vector.ls

## and creates the following files:

crts.la acia.la vector.la The following command line:

```
clabs -r.lst acia.sm8
```

will look for files with the suffix ".lst":

The following command line:

```
clabs -s.lx acia.sm8
```

will generate:

crts.lx
acia.lx
vector.lx

# The clib Utility

clib builds and maintains object module libraries. clib can also be used to collect arbitrary files in one place. < library> is the name of an existing library file or, in the case of replace or create operations, the name of the library to be constructed.

# **Command Line Options**

clib accepts the following command line options, each of which is described in detail below:

clib [optio	ons] <library> <files></files></library>
- <b>-</b>	
-a	accept absolute symbols
-c	create a new library
-d	delete modules from library
-е	accept empty module
-i*	object list filename
-1	load all library at link
-r	replace modules in library
-s	list symbols in library
-t	list files in library
-v	be verbose
-x	extract modules from library

#### **Clib Option Usage**

Option	Description
-a	include absolute symbols in the library symbol table.
-с	create a library containing <i><files></files></i> . Any existing <i><li>library&gt;</li></i> of the same name is removed before the new one is created.
-d	delete from the library the zero or more files in <files>.</files>
-е	accept module with no symbol.
-i*	take object files from a list *. You can put several files per line or put one file per line. Each lines can include comments. They must be prefixed by the '#' character. If the command line contains <files>, then <files> will be also added to the library.</files></files>

#### Clib Option Usage (cont.)

Option	Description
-1	when a library is built with this flag set, all the modules of the library will be loaded at link time. By default, the linker only loads modules necessary for the application.
-r	in an existing library, replace the zero or more files in <files>. If no library <li>library&gt; exists, create a library containing <files>. The files in <files> not present in the library are added to it.</files></files></li></files>
-s	list the symbols defined in the library with the module name to which they belong.
-t	list the files in the library.
-v	be verbose
-X	extract the files in <files> that are present in the library into discrete files with the same names. If no <files> are specified, all files in the library are extracted.</files></files>

At most one of the options -[c r t x] may be specified at the same time. If none of these is specified, the -t option is assumed.

#### **Return Status**

*clib* returns success if no problems are encountered. Otherwise it returns failure. After most failures, an error message is printed to STDERR and the library file is not modified. Output from the -t, -s options, and verbose remarks, are written to STDOUT.

# **Examples**

To build a library and check its contents:

```
clib -c libc one.o two.o three.o clib -t libc
```

#### will output:

one.o two.o three.o To build a library from a list file:

```
clib -ci list libc six.o seven.o
```

where *list* contains:

```
# files for the libc library
one.o
two.o
three.o
four.o
five.o
```

# The cobj Utility

You use **cobj** to inspect relocatable object files or executable. Such files may have been output by the assembler or by the linker. *cobj* can be used to check the size and configuration of relocatable object files or to output information from their symbol tables.

## **Command Line Options**

cobj accepts the following options, described in detail below.

cobj	[options] file		
	-d	output data flows	
	-h	output header	
	-i	display info section	
	-n	output sections	
	-0*	output file name	
	-r	output relocation flows	
	-s	output symbol table	
	-v	display file addresses	
	-x	output debug symbols	

<file> specifies a file, which must be in relocatable format or executable format.

### **Cobj Option Usage**

Option	Description
-d	output in hexadecimal the data part of each section.
-h	display all the fields of the object file header.
-i	display the content of the .info. section in a readable format.
-n	display the name, size and attribute of each section.
-o*	write output module to file *. The default is STDOUT.
-r	output in symbolic form the relocation part of each section.
-S	display the symbol table.
-v	display seek addresses inside the object file.
- <b>x</b>	display the debug symbol table.

If none of these options is specified, the default is **-hns**.

#### Return Status

cobj returns success if no diagnostics are produced (i.e. if all reads are successful and all file formats are valid).

## **Examples**

For example, to get the symbol table:

```
cobj -s acia.o
symbols:
main:
               0000003e section .text defined public
outch:
               0000001b section .text defined public
buffer:
              00000000 section .bss defined public
__ptecr: 00000000 section .bsct defined public zpage getch: 00000000 section .text defined public __ptlec: 00000002 section .bsct defined public zpage
               00000028 section .text defined public
recept:
```

The information for each symbol is: name, address, section to which it belongs and attribute.

# The Ctat Utility

You use the **ctat** utility to produce a file which can be preincluded to the next compilation. The output file contains variables which can be efficiently allocated in the short range section.

## **Command Line Options**

ctat accepts the following command line options, each of which is described in detail below:

```
ctat [options] file
     -c1*
            listings directory
     -f*
            input header file
     -1
            restrict to local directory
     -m#
            maximum short range size
     -0*
            output file name
            verbose
     -v
```

#### **Ctat Option Usage**

Option	Description
-cl*	specify a path for the listing files. By default, listings are in the same directory than the source files.
-f*	specify an input header file.
-1	specify only listings in the local directory are used.
-m#	specify the maximum short range size.
-o*	write output to the file *. The default is STDOUT.
-v	be "verbose".

<file> is an executable whose objects are compiled with command line options -l and +debug, in order to produce listings and debug information.

This utility scans the debug information to build a list of global variables not already defined in a dedicated space and the listing files to create a count usage of these variables. If a header file is entered with the option -f containing a list of variables described using the same syntax than the output file, the specified variables will be ignored for the allocation

The most used variables are then allocated in the short range section by producing a specific header file containing a list of #pragma attribute directives. The maximum available space in the short range section is defaulted to 256 bytes minus the space already used by the provided application. The maximum space can be modified by using the -m option.

The created header file can then be included by recompiling the application sources with the option **-ph** followed by the header file name.

#### **Return Status**

ctat returns success if no problems are encountered. Otherwise it returns failure

## **Examples**

The following command:

```
ctat -o zpage.h test.sm8
```

will generate **zpage.h** as the result file from the application *test.sm8*. The application must be recompiled by adding to the already specified options:

```
cxstm8 -phzpage.h ...
```

# The cvdwarf Utility

**cvdwarf** is the utility used to convert a file produced by the linker into an ELF/DWARF format file.

# **Command Line Options**

cvdwarf accepts the following options, each of which is described in detail below

```
cvdwarf [options] file

-bp## bank start address

-bs# bank shift

+dup accept duplicate headers

-loc complex location description

-o* output file name

+page# define pagination (HC12/HCS08 only)

-rb reverse bitfield (L to R)

-so add stack offset

-v be verbose
```

#### **Cvdwarf Option usage**

Option	Description
-bp#	start address of the banking page.
-bs#	set the window shift to #, which implies that the number of bytes in a window is 2**#.
	THESE FLAGS ARE CURRENTLY ONLY MEANINGFULL FOR THE HC11K4.
+dup	handle duplicate header files individually. By default, the converter assumes that all header files sharing the same name do have the same content or with conditional behaviours.
-loc	location lists are used in place of location expressions whenever the object whose location is being described can change location during its lifetime. THIS POSSIBILITY IS NOT SUPPORTED BY ALL DEBUGGERS.

<sup>&</sup>lt;file> specifies a file, which must be in executable format.

# **Cvdwarf Option usage (cont.)**

Option	Description			
-o*	where * is a filename. * is used to specify the output file for <i>cvdwarf</i> . By default, if -o is not specified, <i>cvdwarf</i> send its output to the file whose name is obtained from the input file by replacing the filename extension with ".elf".			
+page#	output addresses in paged mode where # specifies the page type:			
	# Valid usage for Paging Window			
	1 for banked code Code Code Code Code Code Code Code C			
	2 for banked when Data Paging used RAM 0x7000 to 0x7fff			
	3 both (code and data)  Only for HC12A4 when Data and Code Paging used  Only for HC12A4 px8000 to 0xbfff RAM 0x7000 to 0x7fff			
	By default, the banked mode is disable.  THIS FLAG IS CURRENTLY ONLY MEANINGFULL FOR THE HC12/HCS12 and HCS08.			
	THIS FLAG IS NOT TO BE USED ON ANY S12X PAGING, BASED ON THE EXISTING GLOBAL ADDRESS MODE.			
-rb	reverse bitfield from left to right.			
- <b>s</b> 0	add stack offset. This option has to be selected when using debuggers using the SP value directly.			
	THIS FLAG IS CURRENTLY ONLY MEANINGFULL FOR THE HC08/HCS08.			

#### Cvdwarf Option usage (cont.)

Option	Description
-v	select verbose mode. <i>cvdwarf</i> will display information about its activity.

#### **Return Status**

cvdwarf returns success if no problems are encountered. Otherwise it returns failure.

## **Examples**

Under MS/DOS, the command could be:

```
cvdwarfC:\test\acia.sm8
```

and will produce: C:\test\acia.elf

and the following command:

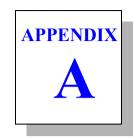
```
cvdwarf -o file C:\test\acia.sm8
```

will produce: file

Under UNIX, the command could be:

```
cvdwarf /test/acia.sm8
```

and will produce: test/acia.elf



# Compiler Error Messages

This appendix lists the error messages that the compiler may generate in response to errors in your program, or in response to problems in your host system environment, such as inadequate space for temporary intermediate files that the compiler creates.

The first pass of the compiler generally produces all user diagnostics. This pass deals with # control lines and lexical analysis, and then with everything else having to do with semantics. Only machine-dependent extensions are diagnosed in the code generator pass. If a pass produces diagnostics, later passes will not be run.

Any compiler message containing an exclamation mark! or the word 'PANIC' indicates that the compiler has detected an inconsistent internal state. Such occurrences are uncommon and should be reported to the maintainers

- Parser (cpstm8) Error Messages
- Code Generator (cgstm8) Error Messages
- Assembler (castm8) Error Messages
- Linker (clnk) Error Messages

# Parser (cpstm8) Error Messages

<name> not a member - field name not recognized for this struct/ union

<name> not an argument - a declaration has been specified for an argument not specified as a function parameter

<name> undefined - a function or a variable is never defined

FlexLM <message>- an error is detected by the license manager

asm string too long - the string constant passed to asm is larger than 255 characters

ambiguous space modifier - a space modifier attempts to redefine an already specified modifier

**array size unknown** - the *sizeof* operator has been applied to an array of unknown size

bad # argument in macro <name> - the argument of a # operator in a #define macro is not a parameter

bad # directive: <name> - an unknown #directive has been specified

**bad** # **syntax** - # is not followed by an identifier

bad ## argument in macro <name> - an argument of a ## operator in a #define macro is missing

bad #asm directive - a #asm directive is not entered at a valid declaration or instruction boundary

**bad #define syntax -** a #define is not followed by an identifier

**bad #elif expression -** a *#elif* is not followed by a constant expression

**bad #else -** a *#else* occurs without a previous *#if*, *#ifdef*, *#ifndef* or *#elif* 

**bad #endasm directive -** a *#endasm* directive is not closing a previous #asm directive

**bad #endif -** a #endif occurs without a previous #if, #ifdef, #ifndef, #elif or #else

**bad #if expression -** the expression part of a #if is not a constant expression

bad #ifdef syntax - extra characters are found after the symbol name

bad #ifndef syntax - extra characters are found after the symbol name

bad #include syntax - extra characters are found after the file name

bad #pragma attribute directive - syntax for the #pragma attribute directive is incorrect

bad #pragma section directive - syntax for the #pragma section directive is incorrect

bad #pragma space directive - syntax for the #pragma space directive is incorrect

bad #pragma unroll directive - syntax for the #pragma unroll directive is incorrect

**bad #undef syntax -** #undef is not followed by an identifier

**bad asm() argument type -** the first argument passed to asm is missing or is not a character string

**bad alias expression** - alias definition is not a valid expression

bad alias value - alias definition is not a constant expression

**bad bit number -** a bit number is not a constant between 0 and 7

**bad character < character> - < character>** is not part of a legal token

bad defined syntax - the defined operator must be followed by an identifier, or by an identifier enclosed in parenthesis

bad function declaration - function declaration has not been terminated by a right parenthesis

bad integer constant - an invalid integer constant has been specified

bad invocation of macro <name> - a #define macro defined without arguments has been invoked with arguments

**bad macro argument** - a parameter in a #define macro is not an identifier

bad macro argument syntax - parameters in a #define macro are not separated by commas

bad proto argument type - function prototype argument is declared without an explicit type

bad real constant - an invalid real constant has been specified

**bad space modifier -** a modifier beginning with a @ character is not followed by an identifier

**bad structure for return -** the structure for return is not compatible with that of the function

bad struct/union operand - a structure or an union has been used as operand for an arithmetic operator

bad symbol definition - the syntax of a symbol defined by the -d option on the command line is not valid

bad void argument - the type void has not been used alone in a prototyped function declaration

can't create <name> - file <name> cannot be created for writing

**can't open <name> -** file <*name>* cannot be opened for reading

can't redefine macro <name> - macro <name> has been already defined

can't undef macro <name> - a #undef has been attempted on a predefined macro

compare out of range - a comparison is detected as beeing always true or always false (+strict)

const assignment - a const object is specified as left operand of an assignment operator

constant assignment in a test - an assignment operator has been used in the test expression of an if, while, do, for statements or a conditional expression (+strict)

duplicate #pragma attibute name < name > - two objects have been declared with the same < name > in #pragma attribute directives

**duplicate case** - two *case* labels have been defined with the same value in the same *switch* statement

**duplicate default** - a *default* label has been specified more than once in a *switch* statement

embedded usage of tag name <name> - a structure/union definition contains a reference to itself

enum size unknown - the range of an enumeration is not available to choose the smallest integer type

**exponent overflow in real -** the exponent specified in a real constant is too large for the target encoding

file too large for label information - the source file is producing too many labels in the code and debug parts for the coding restrictions

float value too large for integer cast - a float constant is too large to be casted in an integer (+strict)

hexadecimal constant too large - an hexadecimal constant is too large to be represented on an integer

**illegal storage class -** storage class is not legal in this context

**illegal type specification -** type specification is not recognizable

illegal void operation - an object of type void is used as operand of an arithmetic operator

illegal void usage - an object of type void is used as operand of an assignment operator

implicit int type in argument declaration - an argument has been declared without any type (+strict)

implicit int type in global declaration - a global variable has been declared without any type (+strict)

implicit int type in local declaration - a local variable has been declared without any type (+strict)

implicit int type in struct/union declaration - a structure or union field has been declared without any type (+strict)

incompatible argument type - the actual argument type does not match the corresponding type in the prototype

incompatible compare type - operands of comparison operators must be of scalar type

incompatible operand types - the operands of an arithmetic operator are not compatible

incompatible pointer assignment - assigned pointers must have the same type, or one of them must be a pointer to void

incompatible pointer operand - a scalar type is expected when operators += and -= are used on pointers

incompatible pointer operation - pointers are not allowed for that kind of operation

incompatible pointer types - the pointers of the assignment operator must be of equal or coercible type

**incompatible return type -** the return expression is not compatible with the declared function return type

incompatible struct/union assignment - a structure or an union has been used as operand for an assignement operator and the other operand is not a structure or an union

incompatible struct/union operation - a structure or an union has been used as operand of an arithmetic operator

incompatible types in struct/union assignment - structure or union types must be identical for assignment

incomplete #elif expression - a #elif is followed by an incomplete expression

**incomplete** #if expression - a #if is followed by an incomplete expression

**incomplete type -** structure type is not followed by a tag or definition

incomplete type for debug information - a structure or union is not completely defined in a file compiled with the debug option set

integer constant too large - a decimal constant is too large to be represented on an integer

invalid #pragma attribute syntax - a syntax error has been detected in a #pragma attribute directive

invalid? test expression - the first expression of a ternary operator (?:) is not a testable expression

invalid address expression - the "address of" operator has been applied to a rvalue expression

invalid address operand - the "address of" operator has been applied to a register variable

invalid address type - the "address of" operator has been applied to a bitfield

invalid alias - an alias has been applied to an extern object

invalid arithmetic operand - the operands of an arithmetic operator are not of the same or coercible types

invalid array dimension - an array has been declared with a dimension which is not a constant expression

invalid binary number - the syntax for a binary constant is not valid

invalid bit assignment - the expression assigned to a bit variable must be scalar

**invalid bit initializer -** the expression initializing a bit variable must be scalar

invalid bitfield size - a bitfield has been declared with a size larger than its type size

invalid bitfield type - a type other than int, unsigned int, char, unsigned char has been used in a bitfield.

invalid break - a break may be used only in while, for, do, or switch statements

invalid case - a case label has been specified outside of a switch statement

invalid case operand - a case label has to be followed by a constant expression

invalid cast operand - the operand of a cast operator in not an expression

invalid cast type - a cast has been applied to an object that cannot be coerced to a specific type

invalid conditional operand - the operands of a conditional operator are not compatible

invalid constant expression - a constant expression is missing or is not reduced to a constant value

**invalid continue** - a continue statement may be used only in while, for, or do statements

invalid default - a default label has been specified outside of a switch statement

invalid do test type - the expression of a do ... while() instruction is not a testable expression

invalid expression - an incomplete or ill-formed expression has been detected

invalid external initialization - an external object has been initialized

invalid floating point operation - an invalid operator has been applied to floating point operands

invalid for test type - the second expression of a for(::) instruction is not a testable expression

invalid function member - a function has been declared within a structure or an union

invalid function type - the function call operator () has been applied to an object which is not a function or a pointer to a function

invalid if test type - the expression of an if () instruction is not a testable expression

invalid indirection operand - the operand of unary \* is not a pointer

invalid line number - the first parameter of a #line directive is not an integer

invalid local initialization - the initialization of a local object is incomplete or ill-formed

invalid lvalue - the left operand of an assignment operator is not a variable or a pointer reference

invalid narrow pointer cast - a cast operator is attempting to reduce the size of a pointer

invalid operand type - the operand of a unary operator has an incompatible type

invalid pointer cast operand - a cast to a function pointer has been applied to a pointer that is not a function pointer

invalid pointer initializer - initializer must be a pointer expression or the constant expression 0

invalid pointer operand - an expression which is not of integer type has been added to a pointer

invalid pointer operation - an illegal operator has been applied to a pointer operand

invalid pointer types - two incompatible pointers have been substracted

invalid shift count type - the right expression of a shift operator is not an integer

invalid sizeof operand type - the sizeof operator has been applied to a function

invalid space for argument < name > - an argument has been declared with a space modifier incompatible with the stack allocation

**invalid space for function -** a function has been declared with a space modifier incompatible with the function allocation

invalid space for local <name> - a local variable has been declared with a space modifier incompatible with the stack allocation

invalid storage class - storage class is not legal in this context

invalid struct/union operation - a structure or an union has been used as operand of an arithmetic operator

invalid switch test type - the expression of a switch () instruction must be of integer type

invalid typedef usage - a typedef identifier is used in an expression

invalid void pointer - a void pointer has been used as operand of an addition or a substraction

invalid while test type - the expression of a while () instruction is not a testable expression

misplaced #pragma section directive - a #pragma section directive has been placed inside the body of a C function

misplaced #pragma atrtibute name - a #pragma attribute directive is not declaring any object

missing ## argument in macro <name> - an argument of a ## operator in a #define macro is missing

missing '>' in #include - a file name of a #include directive begins with '<' and does not end with '>'

missing) in defined expansion - a '(' does not have a balancing ')' in a defined operator

missing; in argument declaration - the declaration of a function argument does not end with ':'

missing; in local declaration - the declaration of a local variable does not end with ':'

missing; in member declaration - the declaration of a structure or union member does not end with ':'

missing? test expression - the test expression is missing in a ternary operator (?:)

missing asm() argument - the asm function needs at least one argument

missing argument - the number of arguments in the actual function call is less than that of its prototype declaration

missing argument for macro <name> - a macro invocation has fewer arguments than its corresponding declaration

missing argument name - the name of an argument is missing in a prototyped function declaration

missing array subscript - an array element has been referenced with an empty subscript

missing do test expression - a do ... while () instruction has been specified with an empty while expression

missing enumeration member - a member of an enumeration is not an identifier

missing explicit return - a return statement is not ending a non-void function (+strict)

missing exponent in real - a floating point constant has an empty exponent after the 'e' or 'E' character

missing expression - an expression is needed, but none is present

missing file name in #include - a #include directive is used, but no file name is present

missing goto label - an identifier is needed after a goto instruction

**missing if test expression -** an *if ()* instruction has been used with an empty test expression

missing initialization expression - a local variable has been declared with an ending '=' character not followed by an expression

missing initializer - a simple object has been declared with an ending '=' character not followed by an expression

missing line number - a line number is missing in a #line directive

missing local name - a local variable has been declared without a name

missing member declaration - a structure or union has been declared without any member

missing member name - a structure or union member has been declared without a name

missing name in declaration - a variable has been declared without a name

missing prototype - a function has been used without a fully prototyped declaration (+strict)

missing prototype for inline function - an inline function has been declared without a fully prototyped syntax

missing return expression - a simple return statement is used in a nonvoid function (+strict)

missing switch test expression - an expression in a switch instruction is needed, but is not present

missing while - a 'while' is expected and not found

missing while test expression - an expression in a while instruction is needed, but none is present

missing: - a ':' is expected and not found

missing; - a ';' is expected and not found. The parser reports such an error on the previous element as most of the time the; is missing at the end of the declaration. When this error occurs on top of a file or just after a file include, the line number reported may not match the exact location where the problem is detected.

missing ( - a '(' is expected and not found

missing) - a')' is expected and not found

missing | - a '/' is expected and not found

missing { - a 'f' is expected and not found

missing \} - a '\}' is expected and not found

missing } in enum definition - an enumeration list does not end with a '?' character

missing } in struct/union definition - a structure or union member list does not end with a '}' character

redeclared #pragma attribute name < name > - a #pragma attribute object is already declared by another #pragma attribute directive

redeclared argument <name> - a function argument has conflicting declarations

**redeclared enum member <name> -** an *enum* element is already declared in the same scope

redeclared external <name> - an external object or function has conflicting declarations

**redeclared local <name> -** a *local* is already declared in the same scope

redeclared proto argument <name> - an identifier is used more than once in a prototype function declaration

**redeclared typedef <name> -** a typedef is already declared in the same scope

**redefined alias <name> -** an alias has been applied to an already declared object

**redefined label <name> -** a label is defined more than once in a function

redefined member <name> - an identifier is used more than once in structure member declaration

**redefined tag <name> -** a tag is specified more than once in a given scope

**repeated type specification** - the same type modifier occurs more than once in a type specification

scalar type required - type must be integer, floating, or pointer

shift count out of range - a constant shift count is larger than the shifted object size (+strict)

size unknown - an attempt to compute the size of an unknown object has occurred

**space attribute conflict** - a space modifier attempts to redefine an already specified modifier

space conflict with #pragma attribute - a space modifier declared with a #pragma attribute mismatches the space modifier specified in the object declaration

stack attibute conflict on cast - a cast is attempting to change the (a) stack/(a) nostack attribute of an object (+strict)

string too long - a string is used to initialize an array of characters shorter than the string length

struct/union size unknown - an attempt to compute a structure or union size has occurred on an undefined structure or union

syntax error - an unexpected identifier has been read

token overflow - an expression is too complex to be parsed

too many argument - the number of actual arguments in a function declaration does not match that of the previous prototype declaration

too many arguments for macro <name> - a macro invocation has more arguments than its corresponding macro declaration

too many initializers - initialization is completed for a given object before initializer list is exhausted

too many spaces modifiers - too many different names for '@' modifiers are used

truncating assignment - the right operand of an assignment is larger than the left operand (+strict)

truncating constant cast - a cast is attempting to narrow down the value of a constant (+strict)

unbalanced '- a character constant does not end with a simple quote

unbalanced " - a string constant does not end with a double quote

<name> undefined - an undeclared identifier appears in an expression

undefined label <name> - a label is never defined

undefined struct/union - a structure or union is used and is never defined

**unexpected end of file** - last declaration is incomplete

**unexpected return expression -** a return with an expression has been used within a void function

**unknown enum definition** - an enumeration has been declared with no member

**unknown structure** - an attempt to initialize an undefined structure has been done

unknown union - an attempt to initialize an undefined union has been done

**unreachable code -** a code sequence cannot be accessed (+strict)

value out of range - a constant is assigned to a variable too small to represent its value (+strict)

variable arguments in nostack mode - a function has been declared with the ... syntax and the@nostack modifier (+strict)

**zero divide -** a divide by zero was detected

zero modulus - a modulus by zero was detected

# Code Generator (cgstm8) Error Messages

**bad builtin** - the @builtin type modifier can be used only on functions

**bad** @interrupt usage - the @interrupt type modifier can only be used on functions.

invalid indirect call - a function has been called through a pointer with more than one *char* or *int* argument, or is returning a structure.

redefined space - the version of cpstm8 you used to compile your program is incompatible with cgstm8.

**unknown space** - you have specified an invalid space modifier @xxx

unknown space modifier - you have specified an invalid space modifier @xxx

PANIC! bad input file - cannot read input file

PANIC! bad output file - cannot create output file

PANIC! can't write - cannot write output file

All other PANIC! messages should never happen. If you get such a message, please report it with the corresponding source program to COSMIC

# Assembler (castm8) Error Messages

The following error messages may be generated by the assembler. Note that the assembler's input is machine-generated code from the compiler. Hence, it is usually impossible to fix things 'on the fly'. The problem must be corrected in the source, and the offending program(s) recompiled.

**bad** .source directive - a .source directive is not followed by a string giving a file name and line numbers

bad addressing mode - an invalid addressing mode have been constructed

**bad argument number-** a parameter sequence n uses a value negative or greater than 9

bad character constant - a character constant is too long for an expression

bad comment delimiter- an unexpected field is not a comment

**bad constant -** a constant uses illegal characters

**bad else** - an *else* directive has been found without a previous *if* directive

**bad endif** - an *endif* directive has been found without a previous *if* or else directive

**bad file name** - the *include* directive operand is not a character string

bad index register - an invalid register has been used in an indexed addressing mode

**bad register** - an invalid register has been specified as operand of an instruction

**bad relocatable expression -** an external label has been used in either a constant expression, or with illegal operators

bad string constant - a character constant does not end with a single or double quote

**bad symbol name: <name> -** an expected symbol is not an identifier can't create <name> - the file <name> cannot be opened for writing **can't open <name> -** the file <*name>* cannot be opened for reading can't open source <name> - the file <name> cannot be included cannot include from a macro - the directive *include* cannot be specified within a macro definition

cannot move back current pc - an org directive has a negative offset

**illegal size** - the size of a ds directive is negative or zero

missing label - a label must be specified for this directive

missing operand - operand is expected for this instruction

**missing register -** a register is expected for this instruction

missing string - a character string is expected for this directive

relocatable expression not allowed - a constant is needed

section name <name> too long - a section name has more than 15 characters

string constant too long - a string constant is longer than 255 characters

symbol <name> already defined - attempt to redefine an existing symbol

**symbol <name> not defined -** a symbol has been used but not declared **syntax error** - an unexpected identifier or operator has been found

too many arguments - a macro has been invoked with more than 9 arguments

too many back tokens - an expression is too complex to be evaluated **unclosed if -** an *if* directive is not ended by an *else* or *endif* directive unknown instruction <name> - an instruction not recognized by the processor has been specified

value too large - an operand is too large for the instruction type zero divide - a divide by zero has been detected

# Linker (clnk) Error Messages

-a not allowed with -b or -o - the after option cannot be specified if any start address is specified.

+def symbol <symbol > multiply defined - the symbol defined by a +def directive is already defined.

bad address (<value>) for zero page symbol <name> - a symbol declared in the zero page is allocated to an address larger than 8 bits.

**bad file format** - an input file has not an object file format.

**bad number in +def** - the number provided in a **+def** directive does not follow the standard C syntax.

**bad number in +spc <segment> -** the number provided in a +spc directive does not follow the standard C syntax.

bad processor type - an object file has not the same configuration information than the others

bad reloc code - an object file contains unexpected relocation information

bad section name in +def - the name specified after the '@' in a +def directive is not the name of a segment.

can't create map file <file> - map file cannot be created.

**can't create <file> -** output file cannot be created.

can't locate .text segment for initialization - initialized data segments have been found but no host segment has been specified.

can't locate shared segment - shared datas have been found but no host segment has been specified.

**can't open file <file> -** input file cannot be found.

file already linked - an input file has already been processed by the linker

**function <function> is recursive -** a **nostack** function has been detected as recursive and cannot be allocated.

function <function > is reentrant - a function has been detected as reentrant. The function is both called in an interrupt function and in the main code

**incomplete** +def directive - the +def directive syntax is not correct.

**incomplete** +seg directive - the +seg directive syntax is not correct.

incomplete +spc directive - the +spc directive syntax is not correct.

init segment cannot be initialized - the host segment for initialization cannot be itself initialized

**invalid** (a) argument - the syntax of an optional input file is not correct.

invalid -i option - the -i directive is followed by an unexpected character.

missing command file - a link command file must be specified on the command line

missing output file - the -o option must be specified.

missing '=' in +def - the +def directive syntax is not correct.

missing '=' in +spc <segment> - the +spc directive syntax is not correct

named segment < segment > not defined - a segment name does not match already existing segments.

**no default placement for segment <segment> -** a segment is missing -a or -b option.

**prefixed symbol <name> in conflict -** a symbol beginning by 'f' (for a banked function) also exists without the 'f' prefix.

read error - an input object file is corrupted

segment <segment> and <segment> overlap - a segment is overlapping an other segment.

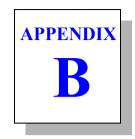
**segment <segment> size overflow -** the size of a segment is larger than the maximum value allowed by the -m option.

shared segment not empty - the host segment for shared data is not empty and cannot be used for allocation.

symbol <symbol> multiply defined - an object file attempts to redefine a symbol.

symbol <symbol> not defined - a symbol has been referenced but never defined

unknown directive - a directive name has not been recognized as a linker directive.



# **Modifying Compiler Operation**

This chapter tells you how to modify compiler operation by making changes to the standard configuration file. It also explains how to create your own programmable options" which you can use to modify compiler operation from the cxstm8.cxf.

### The Configuration File

The configuration file is designed to define the default options and behaviour of the compiler passes. It will also allow the definition of programmable options thus simplifying the compiler configuration. A configuration file contains a list of options similar to the ones accepted for the compiler driver utility **cxstm8**.

These options are described in Chapter 4, "Using The Compiler". There are two differences: the option -f cannot be specified in a configuration file, and the extra -m option has been added to allow the definition of a programmable compiler option, as described in the next paragraph.

The contents of the configuration file cxstm8.cxf as provided by the default installation appears below:

```
# CONFIGURATION FILE FOR STM8 COMPILER
# Copyright (c) 2008 by COSMIC Software
#
                       # unsigned char
-pu
                       # pack local bit variables
-ppb
-i c:\cx32\hstm8
                       # include path
-m debug:x
                       # debug: produce debug info
-m fast:,i,,,compact # fast: inline long transferts
-m compact:,,f7
                       # compact: do not factorize code
                      # nobss: do not use bss
-m nobss:,bss
-m nocst:,ct
                       # nocst: constant in text section
-m nocross:,nc
                       # functions do not cross boundaries
                       # proto: enable prototype checking
-m proto:p
                      # rev: reverse bit field order
-m rev:rb
-m strict:ck
                      # strict: enforce type checking
-m split:,sf
                      # functions in different sections
-m mods:hmods.h,fl # stack model
-m modsl:hmodsl.h,fl
                     # stack long model
-m mods0:hmods0.h
                      # stack model 64K
-m modsl0:hmodsl0.h
                      # stack long model 64K
-m warn:w1
                       # warn: enable warnings
```

The following command line:

```
cxstm8 hello.c
```

in combination with the above configuration file directs the cxstm8 compiler to execute the following commands:

```
cpstm8 -o \2.cx1 -u -i\cosmic\hstm8 hello.c
cgstm8 -o \2.cx2 \2.cx1
costm8 -o \2.cx1 \2.cx2
castm8 -o hello.o -i\cosmic\hstm8 \2.cx1
```

### **Changing the Default Options**

To change the combination of options that the compiler will use, edit the configuration file and add your specific options using the -p (for the parser), -g (for the code generator), -o (for the optimizer) and -a (for the assembler) options. If you specify an invalid option or combination of options, compilation will not proceed beyond the step where the error occurred. You may define up to 60 such options.

### **Creating Your Own Options**

To create a programmable option, edit the configuration file and define the parametrable option with the -m\* option. The string \* has the following format:

```
name:popt,gopt,oopt,aopt,exclude...
```

The first field defines the option *name* and must be ended by a colon character ':'. The four next fields describe the effect of this option on the four passes of the compiler, respectively the *parser*, the *generator*, the *optimizer* and the *assembler*. These fields are separated by a comma character ','. If no specific option is needed on a pass, the field has to be specified empty. The remaining fields, if specified, describe a exclusive relationship with other defined options. If two *exclusive* options are specified on the command line, the compiler will stop with an error message. You may define up to 20 programmable options. At least one field has to be specified. Empty fields need to be specified only if a useful field has to be entered after.

In the following example:

```
-m dl1:1,dl1,,,dl2# dl1: line option 1
-m dl2:1,dl2,,,dl1# dl1: line option 2
```

the two options dl1 and dl2 are defined. If the option +dl1 is specified on the compiler command line, the specific option -l will be used for the parser and the specific option -dl1 will be used for the code generator. No specific option will be used for the optimizer and for the assembler. The option dl1 is also declared to be exclusive with the option dl2, meaning that dl1 and dl2 will not be allowed together on the compiler command line. The option dl2 is defined in the same way.

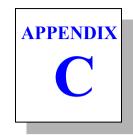
### **Example**

The following command line

```
cxstm8 +nobss hello.c
```

in combination with the previous configuration file directs the cxstm8 compiler to execute the following commands:

```
cpstm8 -o \2.cx1 -u -i\cosmic\hstm8 hello.c
cgstm8 -o \2.cx2 -bss \2.cx1
costm8 -o \2.cx1 \2.cx2
castm8 -o hello.o -i\cosmic\hstm8 \2.cx1
```



# **STM8 Machine Library**

This appendix describes each of the functions in the Machine Library (libm, using d prefix and libm0 for application smaller than 64K, using c prefix). These functions provide the interface between the STM8 microcontroller hardware and the functions required by the code generator. They are described in reference form, and listed alphabetically.

Note that machine library functions handle values as follows:

- integer in the registers x or y.
- longs and floats in the four byte memory location c lreg, ("float register" or "long register" depending on context).
- pointer to long or float in registers x or y for @near pointers and in  $\mathbf{x}$  and the memory location  $\mathbf{c}$   $\mathbf{x}$  or  $\mathbf{v}$  and the memory location  $\mathbf{c}$   $\mathbf{v}$ for **a**far pointers.

The library functions using a pointer to far memory have a name beginning with the 'f' letter, and the pointer is located in the pair composed by the x register for the lower word, and the memory location  $\mathbf{c} \times \mathbf{x}$  for the upper byte. The following describes only the function handling data in near memory. Their equivalent functions have the same description except for the pointer location and size.

### c\_bitfw

### **Description**

Update an int bitfield in near memory

### **Syntax**

```
; bitfield address in y
; mask in c_x and c_x+1
; value in a and x1
    call c_bitfw
```

### **Function**

 $c\_bitfw$  is used to update a 16 bits bitfield located in extended memory by a new value located in the **a** and **xl** register pair. The value loaded from extended memory is first and ed with the mask located in the **c\_x** memory location. It is then or ed with the value in the **a** and **xl** register pair and stored back in memory.

### Return Value

None

### c eewbf

### **Description**

Eeprom char bit field update

### **Syntax**

```
; address in c x and extension
; mask in a
     call c eewbfb
```

#### Function

c eewbf updates a char bit field (8 bits sized) located in eeprom with a new value. The new value is in register x and is right justified. The byte address in eeprom is in c x and c x+1, and the mask, giving the bit field size and location, is in register a. The function waits for the time necessary to program the new value.

### See Also

c eewstr

### c eewrc

### **Description**

Write a char in eeprom

### **Syntax**

```
; value in a
; address in x
     call c eewrc
```

#### Function

c eewrc writes a byte in eeprom. The new byte value is in the a register and its address in eeprom is in x. The function tests if the erasure is necessary, and do it only in that case. Then if the new value is different from one in eeprom, the new byte is programmed. The function waits for the time necessary to program correctly the byte. The function does not test if the byte address is in the address range corresponding to the existing eeprom.

#### See Also

c eewrl, c eewrw

### c eewrl

### **Description**

Write a long int in *eeprom* 

### **Syntax**

```
; value in c lreg
; address in x
     call c eewrl
```

### **Function**

c eewrl writes a long int in eeprom. The new value is in the long register, and its address in eeprom is in x. Each byte is programmed independently by the *c* eewrc function.

### See Also

c eewrc, c eewrw

### c\_eewrw

### **Description**

Write a short int in *eeprom* 

### **Syntax**

```
; value on the stack
; address in x
    call c_eewrw
```

### **Function**

 $c\_eewrw$  writes a short int in *eeprom*. The new value is on the stack, and its address in *eeprom* is in  $\mathbf{x}$ . Each byte is programmed independently by the  $c\_eewrc$  function.

#### See Also

c eewrc, c eewrl

### c eewstr

### **Description**

Move a structure in *eeprom* 

### **Syntax**

```
; source address in y
; destination address in x
; size in a
     call c eewstr
```

### Function

c eewstr moves a structure into an eeprom memory location. Pointer to source is in  $\mathbf{y}$ , and pointer to destination is in  $\mathbf{x}$ . The structure size is in register a. Each byte is programmed independently by the c eewrc function

#### See Also

c eewbfb, c eewrc

### c fadd

# **Description**Add float to float

### **Syntax**

```
; left in float register
; pointer to right in x register
   call c fadd
; result in float register
```

### **Function**

c fadd adds the float in float register to the float indicated by the  $\mathbf{x}$  register. No check is made for overflow.

#### **Return Value**

The resulting value is in *float register*. Flags have no meaningful value upon return.

### See Also

c fsub

### c fcmp

### **Description**

Compare floats

### **Syntax**

```
; left in float register
; pointer to right in x register
    call c fcmp
; result in flags
```

#### Function

c fcmp compares the float in float register with the float pointered at by the x register.

### Return Value

The N and Z flags are set to reflect the value (left-right).

### c fdiv

### **Description**

Divide float by float

### **Syntax**

```
; left in float register
; pointer to right in x register
   call c fdiv
; result in float register
```

### **Function**

c fdiv divides the float in float register by the float pointered at by the  $\mathbf{x}$ register.

### **Return Value**

The resulting value is in *float register*. Flags have no meaningful value upon return.

# c fgadd

### **Description**

Add float to float in memory

### **Syntax**

```
; pointer to left in x register
; right in float register
     call c fgadd
; result in memory
```

### Function

c fgadd adds the float in the float pointered at by the  $\mathbf{x}$  register to the float register.

### Return Value

The resulting value is stored at the location pointered at by the  $\mathbf{x}$  register, meaning that the left operand is updated. Flags have no meaningful value upon return.

#### See Also

c fgsub

# c fgmul

### **Description**

Multiply float by float in memory

### **Syntax**

```
; pointer to left in x register
; right in float register
     call c fgmul
; result in memory
```

### **Function**

c fgmul multiplies the float in float register by the float pointered at by the x register.

#### **Return Value**

The resulting value is stored at the location pointered at by x. Flags have no meaningful value upon return.

# c fgsub

### **Description**

Subtract float from float in memory

### **Syntax**

```
; pointer to left in x register
; right in float register
    call c fgsub
; result in memory
```

#### Function

c fgsub subtracts the float pointered at by the  $\mathbf{x}$  register from the float in float register. No check is made for overflow.

### Return Value

The resulting value is stored at the location pointered at by  $\mathbf{x}$ . Flags have no meaningful value upon return.

### See Also

c fgadd

### c fmul

### **Description**

Multiply float by float

### **Syntax**

```
; left in float register
; pointer to right in x register
   call c fmul
; result in float register
```

#### **Function**

c fmul multiplies the float in float register by the float pointered at by the x register.

### **Return Value**

The resulting value is in *float register*. Flags have no meaningful value upon return

# c fneg

# **Description**Negate a float

### **Syntax**

```
; operand in float register
     call c fneg
; result in operand
```

### **Function**

c fneg negates the float pointered at by the float register.

### Return Value

The resulting value is in *float register*. Flags have no meaningful value upon return.

### c fsub

**Description**Subtract float from float

### **Syntax**

```
; left in float register
; pointer to right in x register
   call c fsub
; result in float register
```

### **Function**

c fsub subtracts the float pointed to by the  $\mathbf{x}$  register from the float in float register. No check is made for overflow.

### **Return Value**

The resulting value is in *float register*. Flags have no meaningful value upon return.

### See Also

c fadd

### c ftoi

### **Description**

Convert float to integer

### **Syntax**

```
; float in float register
     call c ftoi
; result in register x
```

### **Function**

c ftoi converts the float in float register to a two byte integer in the  $\mathbf{x}$ register. No check is made for overflow.

#### **Return Value**

The resulting value is in x. Flags have no meaningful value upon return.

```
c ftol, c itof, c itol, c ltof
```

### c\_ftol

### **Description**

Convert float into long integer

### **Syntax**

```
; float in float register
    call c_ftol
;result in long register
```

### **Function**

*c\_ftol* converts the float in *float register* to a four byte integer in *long register*. No check is made for overflow.

### **Return Value**

The resulting value is in *long register*. Flags have no meaningful value upon return.

```
c ftoi, c itof, c itol, c ltof
```

### c fzmp

### **Description**

Compare a float in memory to zero

### **Syntax**

```
; pointer to operand in x register
     call c fzmp
; result in flags
```

### **Function**

c fzmp compares the float pointered by the **x** register against zero.

### Return Value

The N and Z flags are set to reflect the operand value.

### c\_getlx

### **Description**

Get a long word from memory

### **Syntax**

```
; long address in x
    call c_getlx
; result in long register
```

### **Function**

 $c\_getlx$  gets a long integer from memory using a pointer loaded in **x**. The result is left in the *long register*.

### **Return Value**

The byte is loaded in the *long register*. Flags have no meaningful value upon return.

```
c getly, c getwfx, c getwfy
```

### c getly

### **Description**

Get a long word from memory

### **Syntax**

```
; long address in y
     call c gety
; result in long register
```

### **Function**

c getly gets a long integer from memory using a pointer loaded in y. The result is left in the *long register*.

### Return Value

The byte is loaded in the *long register*. Flags have no meaningful value upon return.

```
c getlx, c getwfx, c getwfy
```

# c\_getwfx

### **Description**

Get a word from far memory

### **Syntax**

```
; word address in x and c_x
     call c_getwfx
; result in a and xl
```

### **Function**

 $c\_getwfx$  gets a word from far memory using a pointer loaded in **x** and  $c\_x$ . The result is left in the **a** and **xl** registers.

#### **Return Value**

The word is loaded in the **a** and **xl** registers. Flags have no meaningful value upon return.

```
c getlx, c getly, c getwfy
```

# c getwfy

### **Description**

Get a word from far memory

### **Syntax**

```
; word address in y and c y
    call c getwfy
; result in a and xl
```

### Function

c\_getwfy gets a word from far memory using a pointer loaded in y and c y. The result is left in the a and xl registers.

### Return Value

The word is loaded in the a and xl registers. Flags have no meaningful value upon return.

```
c getlx, c getly, c getwfx
```

### c\_getxfx

### **Description**

Get a word from far memory

### **Syntax**

```
; word address in x and c_x
     call c_getxfx
; result in x
```

### **Function**

 $c\_getxfx$  gets a word from far memory using a pointer loaded in **x** and  $c\_x$ . The result is left in the **x** register.

### **Return Value**

The word is loaded in the  $\mathbf{x}$  register. Flags have no meaningful value upon return.

```
c getlx, c getly, c getwfy
```

# c getxfy

### **Description**

Get a word from far memory

### **Syntax**

```
; word address in y and c y
    call c getxfy
; result in x
```

### Function

c getxfy gets a word from far memory using a pointer loaded in y and c v. The result is left in the x register.

#### Return Value

The word is loaded in the x register. Flags have no meaningful value upon return.

```
c getlx, c getly, c getwfy
```

# c\_getyfx

### **Description**

Get a word from far memory

### **Syntax**

```
; word address in x and c_x
     call c_getyfx
; result in y
```

### **Function**

 $c\_getyfx$  gets a word from far memory using a pointer loaded in **x** and **c x**. The result is left in the **y** register.

#### **Return Value**

The word is loaded in the y register. Flags have no meaningful value upon return.

```
c getlx, c getly, c getwfy
```

# c\_getyfy

### **Description**

Get a word from far memory

### **Syntax**

```
; word address in y and c y
    call c getyfy
; result in y
```

### Function

c getyfy gets a word from far memory using a pointer loaded in y and **c v**. The result is left in the **v** register.

### Return Value

The word is loaded in the y register. Flags have no meaningful value upon return.

```
c getlx, c getly, c getwfy
```

### c idiv

### **Description**

Quotient of integer division

### **Syntax**

```
; dividend in x
; divisor in y
   call c idiv
; quotient in x
```

### **Function**

c idiv divides the two byte integer in the x register, by the two byte integer in the y register. Values are assumed to be signed. If division by zero is attempted, the result is the unchanged dividend.

#### Return Value

The quotient is placed in x. Flags have no meaningful value upon return.

#### See Also

c udiv

### c imul

### **Description**

Integer multiplication

### **Syntax**

```
; left in x
; right in y
   call c_imul
; result in x
```

### **Function**

c imul multiplies the two byte integer in the  $\mathbf{x}$  register, by the two byte integer in the y register. No check is made for overflow.

### **Return Value**

The resulting value is in x. Flags have no meaningful value upon return

### c itof

### **Description**

Convert integer into float

### **Syntax**

```
; operand in x
     call c itof
; result in float register
```

### **Function**

c it of converts the two byte integer in the x register, to a float stored in float register.

### **Return Value**

The resulting value is in *float register*. Flags have no meaningful value upon return.

```
c ltof, c ultof, c xtof, c uitof, c uxtof
```

# c itol

## **Description**

Convert integer into long

# **Syntax**

```
; operand in a and xl
     call c itol
; result in long register
```

#### **Function**

c itol converts the two byte integer in the a and xl register pair, to a long integer stored in long register.

# **Return Value**

The resulting value is in long register. Flags have no meaningful value upon return.

#### See Also

c xtol, c uitol, c uxtol

# c\_itolx

## **Description**

Convert integer into long

# **Syntax**

```
; operand in x
     call c_itolx
; result in long register
```

#### **Function**

 $c\_itolx$  converts the two byte integer in the **x** register, to a long integer stored in *long register*.

#### **Return Value**

The resulting value is in *long register*. Flags have no meaningful value upon return.

#### See Also

c xtol, c uitol, c uxtol

# c\_itoly

## **Description**

Convert integer into long

# **Syntax**

```
; operand in y
    call c itoly
; result in long register
```

#### **Function**

c itoly converts the two byte integer in the y register, to a long integer stored in long register.

#### **Return Value**

The resulting value is in long register. Flags have no meaningful value upon return.

#### See Also

```
c xtol, c uitol, c uxtol
```

# c\_jctab

# **Description**

Perform C switch statement on char

# **Syntax**

```
; value in a
; table address after the call
  call c_jctab
```

#### **Function**

 $c\_jctab$  is called to switch to the proper code segment, depending on a value in the **a** register and an address table found just after the call instruction, and consisting in a list of two bytes signed offsets.

## **Return Value**

c jctab jumps to the proper code. It never returns.

#### See Also

c jstab, c jltab

# c iltab

## **Description**

Perform C switch statement on long

# **Syntax**

```
; value in long register
; table address in x
    jp c jltab
```

#### Function

c jltab is called to switch to the proper code segment, depending on a value and an address table. The top of the table is found in the x register, and consists of a count followed by a list of pairs. A pair consists of a value followed by an address. The pair list is ended by the default address. All values are four byte integers. All addresses and the count are two byte integers.

#### **Return Value**

c jltab jumps to the proper code. It never returns.

#### See Also

c jctab, c jstab

# c\_jstab

## **Description**

Perform C switch statement on integer

# **Syntax**

```
; value in x
; table address after the call
  call c_jstab
```

#### **Function**

 $c\_jstab$  is called to switch to the proper code segment, depending on a value in the **x** register and an address table found just after the call instruction, and consisting in a list of two bytes signed offsets.

## **Return Value**

c jstab jumps to the proper code. It never returns.

#### See Also

c jctab, c jltab

# c\_ladc

# **Description**

Long integer addition

# **Syntax**

```
; left in long register
; right in a register
 call c ladc
; result in long register
```

#### Function

c ladc adds the four byte integer in long register and the unsigned char in the a register.

#### Return Value

The result is in *long register*. Flags are not meaningful upon return.

#### See Also

c ladd

# c\_ladd

## **Description**

Long integer addition

# **Syntax**

```
; left in long register
; pointer to right in x register
     call c_ladd
; result in long register
```

#### **Function**

 $c\_ladd$  adds the four byte integer in *long register* and the four byte integer pointered at by the x register.

#### **Return Value**

The result is in *long register*. Flags are not meaningful upon return.

#### See Also

c lcmp, c lsub

# c land

## **Description**

Bitwise AND for long integers

# **Syntax**

```
; left in long register
; pointer to right in x register
 call c land
; result in long register
```

#### Function

c land operates a bitwise AND between the operands. Each operand is taken to be a four byte integer.

#### Return Value

The result is in *long register*. Flags are not meaningful upon return.

#### See Also

c lor, c lxor

# c\_lcmp

## **Description**

Long integer compare

# **Syntax**

```
; left in long register
; pointer to right in x register
     call c lcmp
; result in long flags
```

#### **Function**

c lcmp compares the four byte integer in long register to the four byte integer pointered by the  $\mathbf{x}$  register.

#### **Return Value**

Flags are set accordingly.

#### See Also

c ladd, c lsub, c lsmp



## **Description**

Quotient of long integer division

# **Syntax**

```
; dividend in long register
; pointer to divisor in x register
    call c ldiv
;quotient in long register
```

#### Function

c ldiv divides the four byte integer in long register by the four byte integer pointered by the x register. Values are assumed to be signed. If division by zero is attempted, the result is the unchanged dividend.

#### Return Value

The quotient is in *long register* and the flags are not meaningful upon return

#### See Also

c ludv, c lmod, c lumd

# c\_lgadc

# **Description**

Long addition

# **Syntax**

```
; pointer to left in x register
; right in a register
    call c_lgadc
; result in left
```

#### **Function**

 $c\_lgadc$  performs the long addition of the unsigned char value in the a register and the long value pointed at by the x register.

#### **Return Value**

The result is stored at the location in the  $\mathbf{x}$  register. Flags are not meaningful upon return.

#### See Also

c lgadd, c pgadc

# c lgadd

## **Description**

Long addition

# **Syntax**

```
; pointer to left in x register
; right in long register
     call c lgadd
; result in left
```

#### Function

c lgadd performs the long addition of the value pointered by the x register and the value in long register.

## **Return Value**

The result is stored at the location pointered by the  $\mathbf{x}$  register. Flags are not meaningful upon return.

#### See Also

c lgadc

# c lgand

## **Description**

Long bitwise AND

# **Syntax**

```
; left in long register
; pointer to right in x register
   call c lgand
; result in left
```

## **Function**

c lgand performs the long bitwise AND of the value in long register and the value pointered by the  $\mathbf{x}$  register.

## **Return Value**

The result is stored in long register. Flags are not meaningful upon return.

# c lglsh

## **Description**

Long shift left

## **Syntax**

```
; pointer to long in x register
; shift count in a register
    call c lglsh
; result in memory
```

### Function

c lglsh performs the long left shift of the value pointered by the x register by the bit count in the a register. No check is done against silly counts.

## **Return Value**

The result is stored in the location pointered by x. Flags are not meaningful upon return.

# c\_lgmul

## **Description**

Long multiplication in memory

# **Syntax**

```
; pointer to left in x register
; right in long register
    call c_lgmul
; result in left
```

## **Function**

 $c\_lgmul$  performs the long multiplication of the value pointered by the **x** register, by the value in *long register*.

## **Return Value**

The result is stored in the location pointered by  $\mathbf{x}$ . Flags are not meaningful upon return.

#### See Also

c lmul

# c lgneg

# **Description**

Negate a long integer in memory

# **Syntax**

```
; pointer to operand in x register
    call c lneg
; result in memory
```

#### **Function**

c lgneg negates the four byte integer pointered by the  $\mathbf{x}$  register.

#### Return Value

The result is in the location pointered by x. The flags are not meaningful upon return.

#### See Also

c lneg

# c\_lgor

# **Description**

Long bitwise OR

# **Syntax**

```
; pointer to left in x register
; right in long register
     call c lgor
; result in left
```

#### **Function**

c lgor performs the long bitwise OR of the value in long register and the value pointered by the  $\mathbf{x}$  register.

## **Return Value**

The result is stored in long register. Flags are not meaningful upon return.

# c lgrsh

# **Description**

Signed long shift right

# **Syntax**

```
; pointer to long in x register
; shift count in a register
    call c lgrsh
; result in memory
```

#### Function

c lgrsh performs the signed long right shift of the value pointered by the x register and the value in *long register*. No check is done against silly counts. Because the value is signed, arithmetic shift instructions are used

#### Return Value

The result is stored in the location pointered by x. Flags are not meaningful upon return.

# c lgsbc

# **Description**

Long subtraction

# **Syntax**

```
; pointer to left in x register
; right in a register
     call c lgsbc
; result in left
```

#### **Function**

c lgsbc evaluates the (long) difference between the value pointered by the x register and the unsigned char value in the a register.

#### **Return Value**

The result is stored in the location pointered by x. Flags are not meaningful upon return.

#### See Also

c lgsub

# c lgsub

## **Description**

Long subtraction

# **Syntax**

```
; pointer to left in x register
; right in long register
     call c lgsub
; result in left
```

#### Function

c lgsub evaluates the (long) difference between the value pointered by the x register and the value in *long register*.

#### Return Value

The result is stored in the location pointered by x. Flags are not meaningful upon return.

#### See Also

c lgsbc

# c lgursh

## **Description**

Unsigned long shift right

# **Syntax**

```
; pointer to long in x register
; shift count in a register
     call c lgursh
; result in memory
```

#### **Function**

c lgursh performs the unsigned long right shift of the value pointered by the x register and the value in *long register*. No check is done against silly counts. Because the value is unsigned, logical shift instructions are used

#### Return Value

The result is stored in the location pointered by x. Flags are not meaningful upon return.

# c lgxor

## **Description**

Long bitwise exclusive OR

# **Syntax**

```
; pointer to right in x register
; left in long register
    call c lgxor
; result in left
```

#### Function

c lgxor performs the long bitwise Exclusive OR of the value in long register and the value pointered by the x register.

#### Return Value

The result is stored in long register. Flags are not meaningful upon return.

# c\_llsh

## **Description**

Long integer shift left

# **Syntax**

```
; operand in long register
; shift count in a register
     call c_llsh
; result in long register
```

## **Function**

 $c\_llsh$  shifts left four byte integer in *long register* by the number of places specified by the **a** register. A zero count leaves the *long register* unchanged. No check is made for invalid counts.

#### Return Value

The resulting value is in *long register*. Flags are not meaningful upon return.

#### See Also

c lrsh, c lursh

# c lmod

# **Description**

Remainder of long integer division

# **Syntax**

```
; left in long register
; pointer to right in x register
     call c lmod
; remainder in long register
```

#### Function

c lmod divides the four byte integer in long register by the four byte integer pointered by the x register. Values are assumed to be signed. The dividend is returned if a division by zero is attempted.

#### Return Value

The remainder is stored in *long register*. Flags are not meaningful upon return

#### See Also

```
c lumd, c ldiv, c udiv
```

# c\_lmul

## **Description**

Multiply long integer by long integer

# **Syntax**

```
; left in long register
; pointer to right in x register
        call c_lmul
; result in long register
```

## **Function**

 $c\_lmul$  multiplies the four byte integer in *long register* by the four byte integer pointered by the x register. No check is made for overflow.

#### **Return Value**

The resulting value is in *long register*. Flags are not meaningful upon return.

#### See Also

c lgmul

# c lneg

# **Description**

Negate a long integer

## **Syntax**

```
; operand in long register
    call c lneg
; result in long register
```

#### **Function**

c lneg negates the four byte integer in long register.

#### Return Value

The result is in *long register*. The flags are not meaningful upon return.

#### See Also

c lgneg

# c lor

# **Description**

Bitwise OR with long integers

# **Syntax**

```
; left in long register
; pointer to right in x register
   call c lor
; result in long register
```

# **Function**

c lor operates a bitwise OR between the contents of long register and the long pointered by the x register. Each operand is taken to be a four byte integer.

#### **Return Value**

The result is in *long register*. The flags are not meaningful upon return.

#### See Also

c land, c lxor

# c lrsh

# **Description**

Long integer right shift

# **Syntax**

```
; operand in long register
; shift count in a register
     call c lrsh
; result in long register
```

#### Function

c lrsh right shifts the four byte integer in long register by the number of bits specified by the a register. A zero count leaves the long register unchanged. No check is made for invalid counts. The value is assumed to be signed, so a negative value will stay negative as by an arithmetic shift

#### Return Value

The resulting value stays in long register. Flags are not meaningful upon return.

#### See Also

c llsh, c lursh

# c\_lrzmp

# **Description**

Long test against zero

# **Syntax**

```
; operand in long register
     call c lrzmp
; result in the flags
```

#### **Function**

c lrzmp tests the value in the long register and updates the sign and zero flags.

#### **Return Value**

Nothing, but the flags.

#### See Also

c lzmp

# c lsbc

# **Description**

Long integer subtraction

# **Syntax**

```
; left in long register
; right in a register
  call c lsbc
; result in long register
```

#### Function

c lsbc subtracts the unsigned char in the a register from the four byte integer in long register.

#### **Return Value**

The result is in *long register*. Flags are not meaningful upon return.

#### See Also

c lsub

# c Isub

## **Description**

Long integer subtraction

# **Syntax**

```
; left in long register
; pointer to right in x register
   call c lsub
; result in long register
```

#### **Function**

c lsub subtracts the four byte integer pointered by the  $\mathbf{x}$  register from the four byte integer in long register.

## **Return Value**

The result is in *long register*. Flags are not meaningful upon return.

#### See Also

c ladd, c lcmp



## **Description**

Long integer compare with overflow

# **Syntax**

```
; left in long register
; pointer to right in x register
    call c lsmp
; result in long flags
```

#### Function

c lsmp compares the four byte integer in long register to the four byte integer pointered by the x register and updates the resulting flags in case of overflow.

#### Return Value

Flags are set accordingly.

#### See Also

c ladd, c lsub, c lcmp

# c ltof

## **Description**

Convert long integer into float

# **Syntax**

```
; operand in float integer
     call c ltof
; result in float register
```

#### **Function**

c ltof converts the four byte integer in float register to a float.

#### **Return Value**

The resulting value is in *float register*. Flags have no meaningful value upon return.

#### See Also

```
c ftoi, c ftol, c itof, c itol
```

# c ltor

# **Description**

Load memory into long register

# **Syntax**

```
; pointer to operand in x register
    call c ltor
; result in float register
```

#### **Function**

c ltor loads the four byte integer pointered by the x register into the long register.

#### Return Value

The resulting value is in *long register*. Flags have no meaningful value upon return.

#### See Also

c rtol

# c ludv

## **Description**

Quotient of unsigned long integer division

# **Syntax**

```
; left in long register
; pointer to right in x register
   call c ludv
; quotient in long register
```

#### **Function**

c ludv divides the four byte integer in long register by the four byte integer pointered by the x register. Values are assumed to be unsigned. The dividend is returned if a division by zero is attempted.

#### Return Value

The quotient is in *long register*. Flags are not meaningful upon return.

#### See Also

c ldiv, c lmod, c lumd

# c lumd

# **Description**

Remainder of unsigned long integer division

## **Syntax**

```
; left in long register
; pointer to right in x register
    call c lumd
; remainder in long register
```

#### Function

c lumd divides the four byte integer in long register by the four byte integer pointered by the x register. Values are assumed to be unsigned. The dividend is returned if a division by zero is attempted.

#### Return Value

The remainder is in *long register*. Flags are not meaningful upon return.

#### See Also

c lmod, c ldiv, c ludv

# c lursh

## **Description**

Unsigned long integer shift right

## **Syntax**

```
; operand in long register
; shift count in a register
    call c lursh
; result in long register
```

## **Function**

c lursh right shifts the four byte integer in long register by the number of bits specified by the a register. A zero count leaves the long register unchanged. No check is made for invalid counts. The value is assumed to be unsigned. The shift instruction used is therefore a logical shift.

#### Return Value

The resulting value is in *long register*. Flags are not meaningful upon return

#### See Also

c llsh, c lrsh

# c lxor

## **Description**

Bitwise exclusive OR with long integers

# **Syntax**

```
; left in long integer
; pointer to right in x register
    call c lxor
; result in long register
```

#### Function

c lxor operates a bitwise Exclusive OR between the contents of long register and the long pointered by the x register. Each operand is taken to be a four byte integer.

#### Return Value

The result is in *long register*. The flags are not meaningful upon return.

#### See Also

c land, c lor

# c\_lzmp

# **Description**

Compare a long integer to zero

# **Syntax**

```
; pointer to operand in x register
     call c lzmp
; result in the flags
```

#### **Function**

c lzmp compares the four byte integer pointered by the x register to

#### **Return Value**

Nothing, but the flags.

#### See Also

c lrzmp

# c\_pgadc

## **Description**

Far pointer addition

## **Syntax**

```
; pointer to left in x register
; right in a register
     call c pgadc
; result in left
```

#### Function

c pgadc performs the long addition of the unsigned char value in the a register and the far pointer pointed at by the x register.

#### Return Value

The result is stored at the location in the x register. Flags are not meaningful upon return.

#### See Also

c lgadd, c lgadc

# c\_pgadd

## **Description**

Far pointer addition

## **Syntax**

```
; pointer to left in x register
; right in long register
     call c_pgadd
; result in left
```

## **Function**

 $c\_pgadd$  performs the long addition of the value in *long register* and the far pointer pointed at by the **x** register.

### **Return Value**

The result is stored at the location in the  $\mathbf{x}$  register. Flags are not meaningful upon return.

#### See Also

c lgadd, c lgadc

# c putlx

# **Description**

Put a long integer in memory

## **Syntax**

```
; long address in x
; value in long register
     call c putlx
```

#### **Function**

c putlx puts the value in long register into memory using a pointer loaded in x.

#### **Return Value**

None

#### See Also

c getlx, c getwx, c putly, c putw

# c\_putly

# **Description**

Put a long integer in memory

# **Syntax**

```
; long address in y
; value in long register
     call c putly
```

#### **Function**

c putly puts the value in long register into memory using a pointer loaded in y.

#### **Return Value**

None

#### See Also

c getly, c getwy, c putlx, c putw

# c\_putwf

# **Description**

Put a word in far memory

## **Syntax**

```
; word address in y and c y
; value in a and xl
     call c putwf
```

#### **Function**

c putwf puts the value in a and x1 registers into far memory using a pointer loaded in y and c y.

#### **Return Value**

None.

#### See Also

c getlx, c getly, c getw, c putlx, c putly

# c pxtox

## **Description**

Get a far pointer from far memory

# **Syntax**

```
; far pointer address in x and c x
     call c pxtox
; result in x and c x
```

#### **Function**

c pxtox gets a far pointer from far memory using a pointer loaded in x and  $\mathbf{c} \times \mathbf{x}$ . The result is left in the  $\mathbf{x}$  and  $\mathbf{c} \times \mathbf{x}$ .

#### **Return Value**

The far pointer is loaded in the x register and the  $c_x$  memory location. Flags have no meaningful value upon return.

```
c pxtoy, c pytox, c pytoy
```

# c pxtoy

## **Description**

Get a far pointer from far memory

## **Syntax**

```
; far pointer address in x and c x
     call c pxtoy
; result in y and c y
```

#### **Function**

c pxtoy gets a far pointer from far memory using a pointer loaded in x and c x. The result is left in the y and c y.

#### Return Value

The far pointer is loaded in the y register and the c y memory location. Flags have no meaningful value upon return.

```
c pxtox, c pytox, c pytoy
```

# c\_pytox

## **Description**

Get a far pointer from far memory

## **Syntax**

```
; far pointer address in y and c_y
    call c_pytox
; result in x and c_x
```

#### **Function**

 $c\_pytox$  gets a far pointer from far memory using a pointer loaded in **y** and  $c\_y$ . The result is left in the **x** and  $c\_x$ .

#### Return Value

The far pointer is loaded in the  $\mathbf{x}$  register and the  $\mathbf{c}_{\mathbf{x}}$  memory location. Flags have no meaningful value upon return.

```
c pxtox, c pxtoy, c pytoy
```

# c pytoy

## **Description**

Get a far pointer from far memory

## **Syntax**

```
; far pointer address in y and c y
    call c pytoy
; result in y and c y
```

#### Function

c pytoy gets a far pointer from far memory using a pointer loaded in y and c v. The result is left in the v and c v.

#### Return Value

The far pointer is loaded in the y register and the c y memory location. Flags have no meaningful value upon return.

```
c pxtox, c pxtoy, c pytox
```

# c rtofl

## **Description**

Store long register in far memory

# **Syntax**

```
; pointer to destination in x register and c x
; operand in long integer
     call c rtofl
```

#### **Function**

c rtofl store the four byte integer in long register into the memory location pointered by the x register and c x.

#### **Return Value**

The resulting value is in the memory location pointered by  $\mathbf{x}$  and  $\mathbf{c}_{\mathbf{x}}$ . Flags have no meaningful value upon return.

#### See Also

c fltor

# c rtol

## **Description**

Store long register in memory

# **Syntax**

```
; pointer to destination in x register
; operand in long integer
     call c rtol
```

#### **Function**

c rtol store the four byte integer in long register into the memory location pointered by the x register.

#### Return Value

The resulting value is in the memory location pointered by x. Flags have no meaningful value upon return.

#### See Also

c ltor

# c\_sdivx

## **Description**

Quotient of signed char division

# **Syntax**

```
; dividend in x register
; divisor in a register
    call c_sdivx
; quotient in x
```

## **Function**

 $c\_sdivx$  divides the signed integer in **x** by the signed byte in the **a** register. Values are assumed to be signed. If division by zero is attempted, the result is the unchanged dividend.

#### Return Value

The quotient is in x. Flags are not meaningful upon return.

```
c cdivx, c cdivy, c sdivy
```

# c sdivy

## **Description**

Quotient of signed char division

# **Syntax**

```
; dividend in y register
; divisor in a register
  call c sdivy
; quotient in y
```

#### Function

c sdivy divides the signed integer in y by the signed byte in the a register. Values are assumed to be signed. If division by zero is attempted, the result is the unchanged dividend.

#### Return Value

The quotient is in y. Flags are not meaningful upon return.

```
c cdivx, c cdivy, c sdivy
```

# c\_smodx

## **Description**

Remainder of signed char division

## **Syntax**

```
; dividend in x register
; divisor in a register
     call c_smodx
; remainder in x
```

## **Function**

 $c\_smodx$  divides the signed integer in **x** by the signed byte in the **a** register. Values are assumed to be signed. If division by zero is attempted, the result is the unchanged dividend.

#### Return Value

The remainder is in x. Flags are not meaningful upon return.

#### See Also

c smody

# c smody

## **Description**

Remainder of signed char division

# **Syntax**

```
; dividend in y register
; divisor in a register
    call c smody
; remainder in y
```

#### Function

c smody divides the signed integer in y by the signed byte in the a register. Values are assumed to be signed. If division by zero is attempted, the result is the unchanged dividend.

#### Return Value

The remainder is in y. Flags are not meaningful upon return.

#### See Also

c smodx

# c\_smul

## **Description**

Multiply long integer by unsigned byte

# **Syntax**

```
; left in long register
; right byte in a register
     call c_smul
; result in long register
```

#### **Function**

 $c\_smul$  multiplies the four byte integer in *long register* by the unsigned byte in the a register. No check is made for overflow.

#### **Return Value**

The resulting value is in *long register*. Flags are not meaningful upon return.

#### See Also

c lgmul

# c uitof

## **Description**

Convert unsigned integer into float

# **Syntax**

```
; operand in x
     call c uitof
; result in float register
```

#### **Function**

c uitof converts the two byte unsigned integer in the  $\mathbf{x}$  register to a float stored in float register.

#### **Return Value**

The resulting value is in *float register*. Flags have no meaningful value upon return.

```
c itof, c ltof, c ultof, c xtof, c uxtof
```

# c\_uitol

## **Description**

Convert unsigned integer into long

## **Syntax**

```
; operand in a and xl
     call c_uitol
; result in long register
```

#### **Function**

*c\_uitol* converts the two byte unsigned integer in the **a** and **xl** register pair, to a long integer stored in *long register*.

#### **Return Value**

The resulting value is in *long register*. Flags have no meaningful value upon return.

#### See Also

c itol, c xtol, c uxtol

# c\_uitolx

## **Description**

Convert unsigned integer into long

## **Syntax**

```
; operand in x
    call c uitolx
; result in long register
```

#### **Function**

c uitolx converts the two byte unsigned integer in the x register, to a long integer stored in *long register*.

#### Return Value

The resulting value is in long register. Flags have no meaningful value upon return.

#### See Also

c itol, c xtol, c uxtol

# c uitoly

## **Description**

Convert unsigned integer into long

## **Syntax**

```
; operand in y
    call c uitoly
; result in long register
```

#### **Function**

c uitoly converts the two byte unsigned integer in the y register, to a long integer stored in long register.

# **Return Value**

The resulting value is in long register. Flags have no meaningful value upon return.

#### See Also

c itol, c xtol, c uxtol

# c ultof

## **Description**

Convert unsigned long integer into float

## **Syntax**

```
; long in long register
    call c ultof
; result in float register
```

#### **Function**

c ultof converts the four unsigned byte integer in long register to a float.

#### Return Value

The resulting value is in *float register*. Flags have no meaningful value upon return.

```
c itof, c ltof, c xtof, c uitof, c uxtof
```

# c umul

# **Description**

Multiply unsigned integers with long result

## **Syntax**

```
; left in x register
; right in y register
   call c umul
; result in long register
```

#### **Function**

c umul multiplies the two byte unsigned integer in the x register by the two byte unsigned integer in the y register.

#### **Return Value**

The resulting value is in long register. Flags are not meaningful upon return.

#### See Also

c vmul

# c vmul

## **Description**

Multiply signed integers with long result

## **Syntax**

```
; left in x register
; right in y register
 call c vmul
; result in long register
```

#### Function

c vmul multiplies the two byte signed integer in the  $\mathbf{x}$  register by the two byte signed integer in the y register.

## **Return Value**

The resulting value is in long register. Flags are not meaningful upon return.

#### See Also

c umul

# c\_xtopy

## **Description**

Store a far pointer into far memory

# **Syntax**

```
; value in x and c_x
; far pointer address in y and c_y
      call c_xtopy
```

#### **Function**

 $c\_xtopy$  stores the far pointer in **x** and  $c\_x$  into far memory using a pointer loaded in **y** and  $c\_y$ .

#### **Return Value**

Flags have no meaningful value upon return.

#### See Also

c ytopx

# c\_xymov

## **Description**

Copy a structure into another

# **Syntax**

```
; pointer to source in y
; pointer to destination in x
; size in a register
     call c xymov
```

#### Function

c xymov copy the source structure pointed by the memory location y into the structure pointed by the x register. The structure size is located in the a register.

#### Return Value

None.

#### See Also

c xymov, c yxmov, c xymvx, c yxmvx

# c\_xymvf

## **Description**

Copy a structure in far memory

# **Syntax**

```
; pointer to source in y and c_y
; pointer to destination in x and c x
; size in a
     call c xymvf
```

## **Function**

c xymvf copy the source structure pointed by the y register and the memory location pointed by  $\mathbf{c}$  y into the structure pointed by the x register and the memory location pointed by c x. The structure size is in the a register.

#### **Return Value**

None

#### See Also

c xymov, c yxmov, c xymvy

# c\_xymvfl

## **Description**

Copy a large structure in far memory

# **Syntax**

```
; pointer to source in y and c_y
; pointer to destination in c_x
; size in x
     call c xymvfl
```

#### Function

c xymvfl copy the source structure pointed by the y register and the memory location pointed by c y into the structure pointed by the memory location pointed by  $\mathbf{c}$   $\mathbf{x}$ . The structure size is in the  $\mathbf{x}$  register.

#### Return Value

None.

#### See Also

c xymov, c yxmov, c xymvy

# c\_xymvx

# **Description**

Copy a structure into another

# **Syntax**

```
; pointer to source in y
; pointer to destination in x
; size in a register
    call c_xymvx
```

#### **Function**

 $c\_xymvx$  copy the source structure pointed by the memory location **y** into the structure pointed by the **x** register. The structure size is located in the **a** register.

#### Return Value

None.

#### See Also

c xymov, c yxmov, c xymvx, c yxmvx

# c xymvxl

## **Description**

Copy a large structure into another

# **Syntax**

```
; pointer to source in y
; pointer to destination in c x
; size in x register
     call c xymvxl
```

#### Function

c xymvxl copy the source structure pointed by the memory location y into the structure pointed by the  $\mathbf{c} \times \mathbf{x}$  memory location. The structure size is located in the x register.

#### Return Value

None.

#### See Also

c xymov, c yxmov, c xymvx, c yxmvx

# c ytopx

## **Description**

Store a far pointer into far memory

# **Syntax**

```
; value in y and c y
; far pointer address in x and c x
     call c_ytopx
```

#### **Function**

c ytopx stores the far pointer in y and c y into far memory using a pointer loaded in  $\mathbf{x}$  and  $\mathbf{c}$   $\mathbf{x}$ .

#### **Return Value**

Flags have no meaningful value upon return.

#### See Also

c xtopy

# c\_yxmov

## **Description**

Copy a structure into another

## **Syntax**

```
; pointer to source in x
; pointer to destination in y
; size in a
     call c yxmov
```

#### **Function**

c yxmov copy the source structure pointed by the  $\mathbf{x}$  register into the structure pointed by the y register. The structure size is in the a register.

#### Return Value

None

#### See Also

c xymov, c xymvx, c yxmvx

# c yxmvf

## **Description**

Copy a structure in far memory

# **Syntax**

```
; pointer to source in x and c x
; pointer to destination in y and c y
; size in a
     call c yxmvf
```

## **Function**

c yxmvf copy the source structure pointed by the x register and the memory location pointed by  $\mathbf{c} \times \mathbf{x}$  into the structure pointed by the  $\mathbf{y}$  register and the memory location pointed by c y. The structure size is in the a register.

#### **Return Value**

None

#### See Also

c xymov, c yxmov, c xymvy



# **Description**

Copy a large structure in far memory

# **Syntax**

```
; pointer to source c x
; pointer to destination in y and c y
; size in x
     call c yxmvfl
```

# Function

c yxmvfl copy the source structure pointed by the memory location pointed by  $\mathbf{c} \times \mathbf{x}$  into the structure pointed by the  $\mathbf{y}$  register and the memory location pointed by  $\mathbf{c}$   $\mathbf{y}$ . The structure size is in the  $\mathbf{x}$  register.

#### Return Value

None.

# See Also

c xymov, c yxmov, c xymvy

# c\_yxmvx

# **Description**

Copy a structure into another

# **Syntax**

```
; pointer to source in x
; pointer to destination in y
; size in a
     call c_yxmvx
```

# **Function**

 $c\_yxmvx$  copy the source structure pointed by the x register into the structure pointed by the y register. The structure size is in the a register.

# **Return Value**

None

# See Also

c xymov, c xymvx, c yxmvx

# c\_yxmvxl

# **Description**

Copy a large structure into another

# **Syntax**

```
; pointer to source in c x
; pointer to destination in y
; size in x
     call c yxmvxl
```

# Function

c yxmvxl copy the source structure pointed by the c x memory location into the structure pointed by the y register. The structure size is in the xregister.

#### Return Value

None.

# See Also

c xymov, c xymvx, c yxmvx



# **Compiler Passes**

The information contained in this appendix is of interest to those users who want to modify the default operation of the cross compiler by changing the configuration file that the **cxstm8** compiler uses to control the compilation process.

This appendix describes each of the passes of the compiler:

cpstm8	the parser
cgstm8	the code generator
costm8	the assembly language optimizer

# The cpstm8 Parser

# **The cpstm8 Parser**

**cpstm8** is the parser used by the C compiler to expand #defines, #includes, and other directives signalled by a #, parse the resulting text, and outputs a sequential file of flow graphs and parse trees suitable for input to the code generator **cgstm8**.

# **Command Line Options**

*cpstm8* accepts the following options, each of which is described in detail below:

cpstm8 [option	onsl file
-ad	expand defines in assembly
-c99	c99 type behaviour
-cc	do not cast const expressions
-ck	extra type checkings
-cp	no constant propagation
-csb	check signed bitfields
-d*>	define symbol=value
-e	run preprocessor only
+e*	error file name
-h*>	include header
-i*>	include path
-1	output line information
-md	make dependencies
-m#	model configuration
-nb	no bitfield packing
-nc	no const replacement
-ne	no enum optimization
-np	allow pointer narrowing
-ns	do not share locals
-0*	output file name
-p	need prototypes
-pb	pack bit variables
-rb	reverse bitfield order
-s	do not reorder locals
-sa	strict ANSI conformance
-u	plain char is unsigned
-w#	enable warnings
-xd	debug info for data
-xp	no path in debug info
-xu	no debug info if unused
-xx	extended debug info
-x	output debug info

# **Parser Option Usage**

Option	Description
-ad	enable #define expansion inside inline assembly code between #asm and #endasm directives. By default, #define symbols are expanded only in the C code.

Option	Description
-c99	authorize the repetition of the const and volatile modifiers in the declaration either directly or indirectly in the typedef.
-cc	do not apply standard type casting to the result of a constant expression. This option allows compatibility with parsers previous to version V4.5p. These previous parsers were behaving as if all constants were considered of type long instead of the default type int. Such expressions were allowing intermediate results to become larger that an int without any truncation.
-ck	enable extra type checking. For more information, see " <i>Extra verifications</i> " below.
-ср	disable the constant propagation optimization. By default, when a variable is assigned with a constant, any subsequent access to that variable is replaced by the constant itself until the variable is modified or a flow break is encountered (function call, loop, label).
-csb	produce an error message if a bitfield is declared explicitly with the <b>signed</b> keyword. By default, the compiler silently ignores the signed feature and handles all bitfields as unsigned values.
-d*^	specify * as the name of a user-defined preprocessor symbol (#define). The form of the definition is -dsymbol[=value]; the symbol is set to 1 if value is omitted. You can specify up to 60 such definitions.
-е	run preprocessor only. cpstm8 only outputs lines of text.
+e*	log errors in the text file * instead of displaying the messages on the terminal screen.
-h*>	include files before to start the compiler process. You can specify up to 60 files.
-i*>	specify include path. You can specify up to 128 different paths. Each path is a directory name, <b>not</b> terminated by any directory separator character, or a file containing a list of directory names.
-1	output line number information for listing or debug.

Option	Description	
-md	create only a list of 'make' compatible dependencies consisting for each source file in the object name followed by a list of header files needed to compile that file.	
-m#	the value # is used to configure the parser behaviour. It is a two bytes value, the upper byte specifies the default space for variables, and the lower byte specifies the default space for functions. A space byte is the or'ed value between a size specifier and several optional other specifiers. The allowed size specifiers are:	
	0x10 @tiny	
	0x20 @near	
	0x30 @far	
	Allowed optional specifiers are:	
	0x02 @pack	
	0x04 @nostack	
	Note that all the combinations are not significant for all the target processors.	
-nb	do not pack bitfields. By default, trailing unused bits in the last bitfield of a structure are removed if this saves at least one byte.	
-nc	do not replace an access to an initialized const object by its value. By default, the usage of a const object whose value is known is replaced by its constant value.	

Option	Description
-ne	do not optimize size of <i>enum</i> variables. By default, the compiler selects the smallest integer type by checking the range of the declared <i>enum</i> members. This mechanism does not allow incomplete <i>enum</i> declaration. When the <b>-ne</b> option is selected, all <i>enum</i> variables are allocated as <i>int</i> variables, thus allowing incomplete declarations, as the knowledge of all the members is no more necessary to choose the proper integer type.
-np	allow pointer narrowing. By default, the compiler refuses to cast the pointer into any smaller object. This option should be used carefully as such conversions are truncating addresses.
-ns	do not share independent local variables. By default, the compiler tries to overlay variables in the same memory location or register if they are not used concurrently.
-O*	write the output to the file *. Default is STDOUT for output if -e is specified. Otherwise, an output file name is required.
-р	enforce prototype declaration for functions. An error message is issued if a function is used and no prototype declaration is found for it. By default, the compiler accepts both syntaxes without any error.
-pb	pack <b>_Bool</b> local variables. By default, <b>_Bool</b> local variables are allocated on one byte each.
-rb	reverse the bitfield fill order. By default, bitfields are filled from <b>less</b> significant bit (LSB) to <b>most</b> significant bit (MSB). If this option is specified, filling works from most significant bit to less significant bit.
-S	do not reorder local variables. By default, the compiler sorts the local variables of a function in order to allocate the most used variables as close as possible to the frame pointer. This allows to use the shortest addressing modes for the most used variables.

Option	Description
-sa	enforce a strict ANSI checking by rejecting any syntax or semantic extension. This option also disables the enum size optimization (-ne).
-u	take a plain char to be of type <b>unsigned char</b> , not signed char. This also affects in the same way strings constants.
-w#	enable warnings if # is greater or equal to 0. By default, warnings are disabled.
-х	generate debugging information for use by the cross debugger or some other debugger or in-circuit emulator. The default is to generate no debugging information.
-xd	add debug information in the object file only for data objects, hiding any function.
-хр	do not prefix filenames in the debug information with any absolute path name. Debuggers will have to be informed about the actual files location.
-xu	do not produce debug information for localized variables if they are not used. By default, the compiler produces a com- plete debug information regardless the variable is accessed or not.
-XX	add debug information in the object file for any label defining code or data.

#### Extra verifications

This paragraph describes the checkings done by the -ck parser option (+strict compiler option) according to the error message produced.

implicit int type in struct/union declaration implicit int type in global declaration implicit int type in local declaration

implicit int type in argument declaration - an object is declared without an explicit type and is defaulted to int according to the ANSI standard.

float value too large for integer cast - a float constant is cast to an integer or a long but is larger than the maximum value of the cast type.

compare out of range - a comparison is made with a constant larger (or smaller) than the possible values for the type of the compared expression

**shift count out of range** - a shift count is larger than the bit size of the shifted expression.

constant assignment in a test - a constant is assigned to a variable in a test expression.

**unreachable code** - a code sequence cannot be reached due to previous optimizations.

missing return expression - a return statement without expression is specified in a function with a non void return type.

missing explicit return - a function is not ending with a return statement.

value out of range - a constant is assigned to a variable and is larger (or smaller) than the possible set of values for that type.

**truncating assignment** - an expression is assigned to a variable and has a type larger than the variable one.

The -ck option also enables internally the prototype checking.

#### Return Status

*cpstm8* returns success if it produces no error diagnostics.

# Example

*cpstm8* is usually invoked before *cgstm8* the code generator, as in:

```
cpstm8 -o \2.cx1 -u -i \cosmic\hstm8 file.c
cgstm8 -o \2.cx2 \2.cx1
```

# The cgstm8 Code Generator

cgstm8 is the code generating pass of the C compiler. It accepts a sequential file of flow graphs and parse trees from cpstm8 and outputs a sequential file of assembly language statements.

As much as possible, the compiler generates freestanding code, but, for those operations which cannot be done compactly, it generates inline calls to a set of machine-dependent runtime library routines.

# **Command Line Options**

cgstm8 accepts the following options, each of which is described in detail below:

cgstm8 [opti	ons] file
-a	optimize _asm code
-bss	do not use bss
-ck	check stack frame
-ct	constants in code
-dl#	output line information
+e*	error file name
- <b>f</b>	full source display
-fl	far library calls
-i	inline long machine calls
-1	output listing
-na	do not xdef alias name
-nc	functions do not cross section
-no	do not use optimizer
-0*	output file name
-sf	split function sections
-v	verbose

# **Code generator Option Usage**

Option	Description
-a	optimize _asm code. By default, the assembly code inserted by a _asm call is left unchanged by the optimizer.
-bss	inhibit generating code into the bss section.
-ck	enable stack overflow checking.

# **Code generator Option Usage (cont.)**

Option	Description
-ct	output constant in the .text section. By default, the compiler outputs literals and constants in the .const section.
-dl#	produce line number information. # must be either '1' or '2'. Line number information can be produced in two ways: 1) function name and line number is obtained by specifying -dl1; 2) file name and line number is obtained by specifying -dl2. All information is coded in symbols that are in the debug symbol table.
+e*	log errors in the text file * instead of displaying the messages on the terminal screen.
-f	merge all C source lines of functions producing code into the C and Assembly listing. By default, only C lines actually producing assembly code are shown in the listing.
-fl	use <i>callf</i> instruction for machine library calls, used for models allowing large applications. By default, machine library functions are called with a <i>call</i> instruction allowing only 64K application. This option is configured by the memory model selection.
-i	produce faster code by inlining machine library calls for long integers handling. The code produced will be larger than without this option.
-1	merge C source listing with assembly language code; listing output defaults to <file>.ls.</file>
-na	do not produce an <i>xdef</i> directive for the <i>equate</i> names created for each C object declared with an absolute address.
-nc	do not allow functions to cross a section boundary.
-no	do not produce special directives for the post-optimizer.
-O*	write the output to the file * and write error messages to STDOUT. The default is STDOUT for output and STDERR for error messages.

# **Code generator Option Usage (cont.)**

Option	Description
-sf	produce each function in a different section, thus allowing the linker to suppress a function if it is not used by the application. By default, all the functions are packed in a single section.
-V	When this option is set, each function name is send to STDERR when <i>cgstm8</i> starts processing it.

# **Return Status**

cgstm8 returns success if it produces no diagnostics.

# **Example**

*cgstm8* usually follows *cpstm8* as follows:

```
cpstm8 -o \2.cx1 -u -i\cosmic\hstm8 file.c
cgstm8 -o \2.cx2 \2.cx1
```

# The costm8 Assembly Language Optimizer

costm8 is the code optimizing pass of the C compiler. It reads source files of STM8 assembly language source code, as generated by the cgstm8 code generator, and writes assembly language statements. costm8 is a peephole optimizer; it works by checking lines function by function for specific patterns. If the patterns are present, costm8 replaces the lines where the patterns occur with an optimized line or set of lines. It repeatedly checks replaced patterns for further optimizations until no more are possible. It deals with redundant load/store operations, constants, stack handling, and other operations.

# **Command Line Options**

costm8 accepts the following options, each of which is described in detail below:

costm8 [options] <file></file>		
-c	keep original lines as comments	
-d*	disable specific optimizations	
-f#	minimum code factorization	
-0*	output file name	
-v	print efficiency statistics	

## **Optimizer Option Usage**

	The state of the s			
Option	Description			
-c	leave removed instructions as comments in the output file.			
-d*	specify a list of codes allowing specific optimizations functions to be selectively disabled.			
-f#	define the minimum of bytes for activating the code factorization. Any value smaller than 4 disables the feature. The default value is 7.			
-0*	write the output to the file * and write error messages to STDOUT. The default is STDOUT for output and STDERR for error messages.			
-v	write a log of modifications to STDERR. This displays the number of removed instructions followed by the number of modified instructions.			

If *<file>* is present, it is used as the input file instead of the default STDIN

# **Disabling Optimization**

When using the optimizer with the -c option, lines which are changed or removed are kept in the assembly source as comment, followed by a code composed with a letter and a digit, identifying the internal function which performs the optimization. If an optimization appears to do something wrong, it is possible to disable selectively that function by specifying its code with the -d option. Several functions can be disabled by specifying a list of codes without any whitespaces. The code letter can be enter both lower or uppercase.

#### **Return Status**

costm8 returns success if it produces no diagnostics.

# Example

*costm8* is usually invoked after *cgstm8* as follows:

```
cpstm8 -o \2.cx1 -u -i\cosmic\hstm8 file.c
cgstm8 -o \2.cx2 \2.cx1
costm8 -o file.s \2.cx2
```

# Index

Symbols	modifier 61
#asm	@interrupt
directive 461	function 59
#asm directive 55	qualifier 59
#endasm	@near
directive 461	.data,.bss sections 53
#endasm directive 55	function 41
#pragma	modifier 50, 63
asm directive 55	modifier, mods 39
directive for inlining 55	modifier, mods 0 39
endasm directive 55	pointer, size 39
space directive 49	variable 63
+dep 1	@near type qualifier 25
linker dependency directive 276	@nosvf qualifier 59
+grp directive 271	@svlreg qualifier 59
+modsl memory model 17	@tiny
+seg option 267	.bsct,.ubsct sections 53
.bsct section 183	modifier, modsl 40
.const	modifier, modsl0 39
output section 468	space modifier 63
segment 278	variable 63
@ tiny 50	@vector modifier 60
@eeprom	ckdesc1 285
type qualifier 12, 51	idesc 282, 283
@far	_asm
.fconst section 41	argument string length 57
fdata section 41, 53	argument string size 57
function 41, 63	assembly sequence 56
modifer 39, 63	code optimization 467
pointer, size 39	in expression 57
@inline	lowercase mnemonics 57
functions 61	return type 57

uppercase mnemonics 57	acos function 89
_asm()	address
function 86	default format 307, 311
inserting assembler function 82	logical end 269
Bool	logical start segment 277
assign expression to 42	logical start set 269
consecutive fields 42	paged format 307, 311
data 63	physical 269
pack local variable 464	physical end 267
packed variables 42	physical start 267
referencing absolute address 46	physical start segment 277
type name 42	set logical 269
variable 42	align directive 196
checksum function 98	allocate memory block 201
_checksum16 function 100	allocate storage for constants 200
_checksum16x function 101	application
checksumx function 99	embedded 260
_fctcpy function 111	non-banked 308
sbreak function 66	system bootstrap 260
	Arccosine 89
Numerics	Arcsine 90
8-bit precision, operation 11	Arctangent 91
, , , , , , , , , , , , , , , , , , ,	Arctangent of y/x 92
A	argument
abort function 87	formatted output to buffer 154
abs function 88	formatted output to stdout 144
absolute	asin function 90
address 298	assembler
address in listing 310	branch shortening 194
hex file generator 9	C style directives 195
listing file 310	code inline 56
listing utility 10	conditional branch range 194
map section 178	conditional directive 190
path name 465	create listing file 179
reference address 46	debug information
section relocation 277	add line 180
symbol 270	label 180
symbol in library 313	endm directive 187
symbol table 265	expression 186
symbol tables 292	filling byte 179, 196
symbol,flagged 292	include directive 193
absolute section 237, 247	label 183
20001210 00011011 2019 211	

listing process 310	segment 278
listing stream 181	bit segment
macro	initialization 279
instruction 187	bitfield
macro argument 188	compiler reverse option 79
macro directive 187	default order 464
macro parameter 188	filling 464
old syntax 194	filling order 79
operator set 186	reverse order 464
section name 191	sign check 462
section predefinition 191	bootloader 283
sections 191	boundary
special parameter \# 188	round up 269
special parameter \* 190	bsct directive 198
special parameter \0 189	buffer
switch directive 191	convert to double 93, 170
xbit directive 254	convert to integer 94
assembly language	convert to long 95, 171
code optimizer 470	convert to unsigned long 172
atan function 91	copy from one to another 138, 139
atan2 function 92	copy to eeprom 105
atof function 93	1,5
atoi function 94	$\mathbf{C}$
atol function 95	C interface
	extra character for far function 63
B	underscore character prefix 63
bank	C interface to assembly language 63
automatic segment creation 270	C library
default mode 321	floating point functions 83
size setting 267	integer functions 82
switched system 277	macro functions 83
base directive 197	package 82
bias	C source
segment parameter 277	lines merging 468
setting 278	c_prefix 353
bit	c lreg
address 192	memory byte 59
address value 278	сх
allocated section 193	memory byte 59
attribute section 193	су
define aliases 274	memory byte 59
number 278	call

instruction 468	code optimizer 9
callf	combination of options 350
instruction 468	command line option 74
carry function 96	configuration file 348
casting 462	configuration file specification 76
ceil function 97	configuration file, predefined option
char	78
signed 465	create assembler file only 77
unsigned 465	debug information, produce 78
character	default behavior 74
fill throughout eeprom buffer 107	default configuration file 76
underscore, start 68	default file names 80
checksum	default operations 459
-ck option 285	default options 74, 348
crc 285	driver 4
functions 285	error files path specification 76
-ik option 286	error message 74
clabs utility 310	exclusive options 350
clib utillity 313	flags 6
clist directive 199, 214, 216, 217, 218,	generate error 323
219, 220, 221, 222, 223, 224	generate error file 81
clst utility 302	generate listing 81
cobj utility 316	header files 84
code	include path definition 77
factorization 78, 470	invoke 74
smaller 78	listing file 77
code generator	listing file path specification 76
compiler pass 467	log error file 76
error log file 468	name 74
code optimizer	object file path specification 76
compiler pass 470	optimizer option specification 77
code/data, no output 267	options 74
compare for lexical order 164	options request 74
compiler	parser option specification 77
ANSI checking 465	predefined option selection 78
assembler 9	preprocessed file only 77
assembler option specification 76	produce assembly file 19
C preprocessor and language parser 8	produce listing file 19
code generation option specification	programmable option 348, 350
76	specific options 4
code generator 9	specify options 75
code optimization 10	stack long model option 78

stack short model option 78	initialization 25
temporary files path 77	int representation 70
type checking 79, 462	long int representation 70
user-defined preprocessor symbol 76	short int representation 70
verbose mode 20, 77	data object
compute 173	automatic 300
const	scope 298
@near memory space 44	type 298
data 43	dc directive 200
qualifier 43	dcb directive 201
constant	debug information
in .text section 468	adding 465
numeric 185	debug symbol
prefix character 185	build table 287
string 185	in object file 181
string character 185	table 298
suffix character 185	debugging
convert	data 298
ELF/DWARF format 320	support tools 297
hex format 306	debugging information
copy 150	data object 298
cos function 102	extract 300
cosh function 103	generate 298, 465
cprd utility 300	line number 298
cross-reference	print file 300
information 179	print function 300
table in listing 182	default base for numerical constants 197
ctat utility 318	default placement
cvdwarf utility 320	.bit segment 278
	.bsct segment 278
D	.bss segment 278
d_prefix 353	.data segment 278
data	.text segment 278
@far modifier 21	definition 287
@far pointer representation 70	DEFs 287
@near modifier 21	dependency
@near pointer representation 70	between function 276
@tiny modifier 21	descriptor
atiny pointer representation 70	host to 268
automatic initialization 36	div function 104
char representation 70	dlist directive 202
float representation 70	ds directive 203
initalized 48	

E	F
eepcpy function 105	fabs function 110
eepera function 106	fail directive 213
eeprom	faster code
@near modifier 51	production 468
erase 106	file length restriction 298
location 12, 51	filling byte 203, 212, 237
eepset function 107	float
ELF/DWARF	single precision library 280
format converter 10	floating point library
else directive 204, 205, 208, 214, 216,	link 82
222	Floating Point Library Functions 83
end directive 206	floor function 112
end5 directive 210	fmod function 113
endc directive 216, 222	format
endif directive 204, 207, 208, 214	ELF/DWARF 320
endm directive 209, 229, 232, 244	frexp function 114
endr 240, 241	function
enum	@inline modifier 61
size optimization 464	arguments 300
environment symbol 193	enforce prototype declaration 79,
equ directive 211, 249	464
error	in separate section 79
assembler log file 179	prototype declaration 79, 464
file name 81	recursive 293
log file 265	returning int 85
message 10	suppress 469
message list 323	suppress unused 79
multiply defined symbol 184, 291	function arguments 300
undefined symbol 287	Functions Implemented as Macros 83
undefined symbol in listing 180	· ·
error message 213	G
even directive 212	generate
executable image 306	.bsct section 63
exit 108	hex record 269
exp function 109	in .bit section 63
expression	in .bss section 63
evaluation 187	in .const section 63
high 187	in .data section 63
low 187	in fconst section 63
page 187	in .fdata section 63
	in .text section 63

in .ubsct section 63	path specification 462
listing file 180	specify path 462
object file 180	include directive 225
getchar function 115	initialization
gets function 116	automatic 282
group	define option 268
option 262	descriptor 282
	descriptor address 283
H	descriptor format 282
heap	first segment 282
allocation 66	initialized segments 282
location 68	marker 268
name 66	startup routine 283
pointer 66	initialize storage for constants 200
size 66	inline
stack 66	@usea modifier 62
start 66	@usex modifier 62
top 66	assembly code 55, 56
-help option 6	block inside a function 55
1 1	block outside a function 55
I	carry function 61
IEEE	function 61
Floating Point Standard 70	header function 84
if directive 204, 208, 214	imask function 61
if directive 207	irq function 61
ifc directive 215	machine library calls 78, 468
ifdef directive 216	produce faster code 78
ifeq directive 217	user macro name 62
ifge directive 218	with _asm function 56, 57
ifgt directive 219	with pragma sequences 55
ifle directive 220	input and output 45
iflt directive 221	input to output 150
ifnc directive 224	input/output 46
ifndef directive 222	integer
ifne directive 223	library 280
imask function 117	interrupt
include	function 60
directory names list 77, 462	function in map 293
file 272	handler 59
file before 462	hardware 59
module 280	software 59
object file 271	vectors table 60

irq function 118	# character prefix,comment 263
isalnum function 119	build freestanding program 260
isalpha function 120	clnk 9
isentrl function 121	command file 262
isdigit function 122	command file example 295
isgraph function 123	command item 262
islower function 124	comment 263
isprint function 125	global command line options 265
ispunct function 126	output file 261
isspace function 127	physical memory 261
isupper function 128	list directive 226
isxdigit function 129	listing
	cross reference 20
L	file location 28
labs function 130	file path specification 310
ldexp function 131	interspersed C and assembly file 19
ldiv function 132	lit directive 227
library	literals
build and maintain 10	in @near space 63
building and maintaining 313	local
create 313	labels 58
delete file 313	local directive 184, 228
extract file 314	local variable
file 280	reorder 464
floating point 82	locate source file 303
integer 82, 280	log function 133
list file 314	log10 function 134
load all files 314	long multiplication 398
load modules 263	
machine 82	M
path specification 265	macro
replace file 314	exit 190
scanned 263	expansion in listing 190
single precision 280	internal label 184
Standard ANSI 280	named syntax 189
version 280	numbered syntax 188
line number	macro directive 229
information 468	main
link	function 293
command file 264	main() routine 35
user command file 22	map
linker	file description 293

modules section 293	module 260
produce information 265	module inspector 10
segment section 293	relocatable 316
stack usage section 293	relocatable file output 181
symbols section 293	relocatable file size 316
max function 135	size 316
memchr function 136	offset
memcmp function 137	segment parameter 277
memcpy function 138	setting 278
memmove function 139	offset directive 236
memory	optimization
location 46	disable selectively 471
long range 50	keep line 471
mapped I/O 46	specific code 470
memory models 12, 39	option
memset function 140	global 264
messg directive 231	org directive 237
mexit directive 230, 232	output
min function 141	default format 307
mlist directive 233, 246	file name 264
modf function 142	listing line number 462
Motorola	specify format 144
S-Records format 307	override
standard S-record, generating 23	data bias 306
moveable	text bias 306
code section 283	_
function used 284	P
moveable code segment 111	page
	address extension 187
N	value 187, 308
named syntax, example 230, 241	page directive 238
new	page header 253
segment control 262	paginating output 303
start region 273	parser
nolist directive 234	behaviour 463
nopage directive 235	compiler pass 460
numbered syntax, example 230, 241	error log file 462
	plen directive 239
O	pointer
object	narrow 464
file location 28	pow function 143
image 259	prefix

filename 465	.eeprom 21, 51
preprocessor	.fconst 21
#define 460	.fdata 21
#include 460	.info. 180, 181, 265, 316
run only 462	.text 21
printf function 144	.ubsct 21, 49
private name region	assembler directive 247
use 288	crossing boundary 41
program	curly braces, initialized data 53
stop execution of 108	definition 259
putchar function 149	name 54
puts function 150	parenthesis,code 53
	pragma definition 53
R	pragma directive 54
rand function 151	single 469
range specification 303	square brackets, uninitialized data 53
redirect output 303	unused 268
REFs 287	user defined 53
region	sections
name 262	default 53
private 273	predefined 53
public 273	relocation 277
use of private name 288	segment
register	bsct start address 270
input/output 48	bss start address 270
relative address 298	build new 280
repeat directive 240	control options 264, 267
repeatl directive 241	data start address 270
restore directive 243	definition 259
rexit directive 241, 244	fill 267
ROM 46	follow current 267
runtime startup	maximum size 268
modifying 34	name 269
	overlap checking 269, 277
S	overlapping 280
save directive 245	overlapping control 269
section	root 268
.bit 21, 42	round up address 269
.bsct 21, 49	section overlap 270
.bss 21, 50	space name 277
.const 21	start,new 267
.data 21, 50	text start address 270

zero size 263	streat function 157
separated address space 277	strehr function 158
set directive 249	stremp function 159
share	strepy function 160
local variable 464	strespn function 161
short addressing 12	strings 164
short range	strlen function 162
@tiny modifier 49	strncat function 163
section 198, 318, 319	strnemp function 164
size 39, 40, 49	strncpy function 165
space 319	strpbrk function 166
sin function 152	strrchr function 167
sinh function 153	strspn function 168
source files listing 302	strstr function 169
source listings 302	strtod function 170
space	strtol function 171
for function 463	strtoul function 172
for variable 463	suffix
space name	assembly file 74
definition 269	C file 74
spc directive 250	input 311
sprintf function 154	output 311
sqrt function 155	suffix letter 185
srand function 156	suppress pagination 303
stack	switch directive 251
amount of memory 293	symbol
check overflow 467	endmem 66
need 293	startmem 66
stack model	_endmem 67
long 40, 50	_startmem 67
long,64K 39	alias 288
short 39	define 262
stack pointer 35	define alias 274
standard ANSI libraries 280	define new 274
startup file	definition 274
crts.sm8 34	export 292
crtsi.s 38	logical end value, equal 275
static data 300	logical start value, equal 275
STM8	physical end value,equal 275
addressing mode 183	physical start value, equal 275
instruction set 183	size value,equal 275
ST-Microelectronics syntax 182	sort alphabetically 265

sort by address 265 user-defined 462 symbol table add 274 information 316 new 287	xdef directive 255, 256 xref directive 254, 255, 256
T tab character setting 252 tabs directive 252 tan function 173 tangent 173 tanh function 174 task entries 293 title directive 253 tolower function 175 toupper function 176 translate executable images 306	
U unreachable code eliminate 11	
variable reorder local 464 volatile data 43 qualifier 43 using keyword 43	
W warnings 79, 465 window set shift 265, 320 size 270	
X xbit assembler directive 193 xbit.b assembler directive 193	