



Diffie Hellman Key exchange

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Code (Python 3)

```
n = 10001
```

```
g = 7919
```

```
def a_b_for_all(x): #finds A for sender and B for receiver
```

```
    a = (g**x)%n ## (g^x)modn
```

```
    return a
```

```
def key_of_sender(b): #Prepares the key for sender
```

```
    k1 = (b**x)%n ## (b^x)modn
```

```
    print(k1)
```

```
def key_of_receiver(a): #Prepares the key for sender
```

```
    k2 = (a**y)%n ## (a^y) modn
```

```
    print(k2)
```

```
x = 850
```

```
y = 215
```

```
A = a_b_for_all(x)
```

```
print(A)
```

```
B = a_b_for_all(y)
```

```
print(B)
```

```
print(key_of_sender(B))
```

```
print(key_of_receiver(A))
```

Screenshots

```
class crypto project Logout  
File Edit View Insert Cell Kernel Widgets Help  
In [8]:  
n = 10001  
g = 7919  
In [9]:  
def a_b_for_all(x): #finds A for sender and B for receiver  
    a = (g**x)%n ## (g^x)modn  
    return a  
def key_of_sender(b): #Prepares the key for sender  
    k1 = (b**x)%n ## (b^x)modn  
    print(k1)  
def key_of_receiver(a): #Prepares the key for sender  
    k2 = (a**y)%n ## (a^y) modn  
    print(k2)  
In [10]:  
x = 850  
y = 215
```

```
class crypto project Logout  
File Edit View Insert Cell Kernel Widgets Help  
In [11]:  
A = a_b_for_all(x)  
print(A)  
8942  
In [12]:  
B = a_b_for_all(y)  
print(B)  
6107  
In [14]:  
print(key_of_sender(B))  
8320  
None  
In [15]:  
print(key_of_receiver(A))  
8320  
None
```

Explanation:

The process of the code is explained below in steps

- In Diffie Helmann key exchange algorithm , n and g are chosen as large prime numbers. For sake of coding, we chose n and g as medium sized prime numbers, where $n = 10001$ and $g = 7919$
- x and y are secret numbers of sender and receiver respectively. For this code we chose $x = 850$ and $y = 215$
- According to Diffie Helmann key exchange $A = (g^x) \bmod n$ (for sender) and $B = (g^y) \bmod n$ (for receiver). In our code we wrote a function named `a_b_for_all` which returned A to receiver and B to sender
- We defined function called `key_of_sender`, which returns sender its key, using the formula $\text{Key} = (B^x) \bmod n$
- And we defined function called `key_of_receiver`, which returns receiver its key, using the formula $\text{Key} = (A^y) \bmod n$
- After computation, both the keys must be equal, and according to our code, the value of our key is 8320
- The screenshots of page 3 can be referred for the details of our code along with its output