

# OS Structures

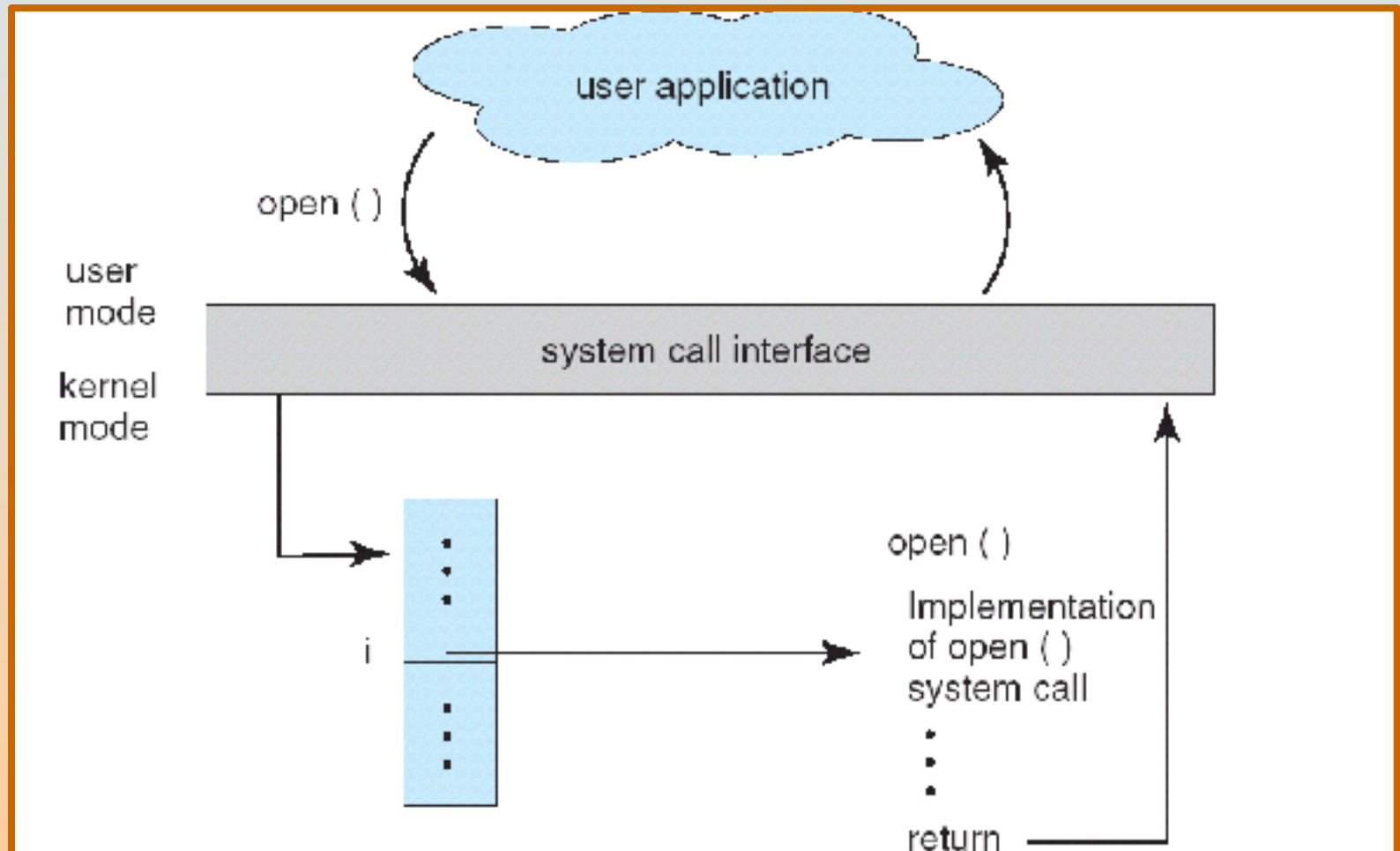
1

- Simple
  - Only one or two levels of code
- Layered
  - Lower levels independent of upper levels
- Microkernel
  - OS built from many user-level processes
- Modular
  - Core kernel with Dynamically loadable modules

# System Calls

2

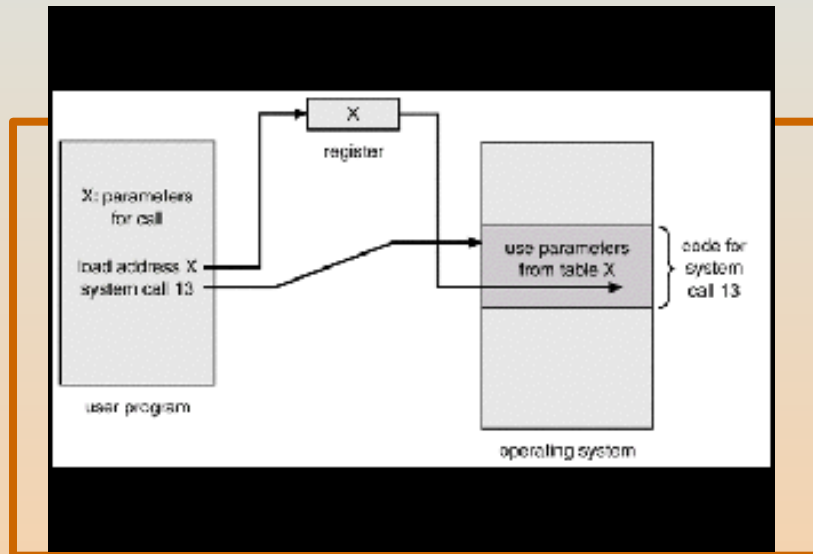
*System calls* provide the interface between a running program and the operating system



# System Calls

3

- Three general methods are used to *pass parameters* between a running program and the operating system
  - Pass parameters in *registers*
  - Store the parameters in a table in memory, and the table address is passed as a parameter in a register



- *Push* (store) the parameters onto the *stack* by the program, and *pop* off the stack by operating system

# Types of System Calls

4

- Process control
- File management
- Device management
- Information maintenance 维持
- Communications

# Abstraction from the Hardware

5

- From transistors to 0/1 bits  
晶体管
- Logic gates abstract away the details of CMOS.
- Machine language abstracts away the details of logic gates.
- Assembly language abstracts away the details of machine languages.
- Programming language abstracts away the details of assembly languages

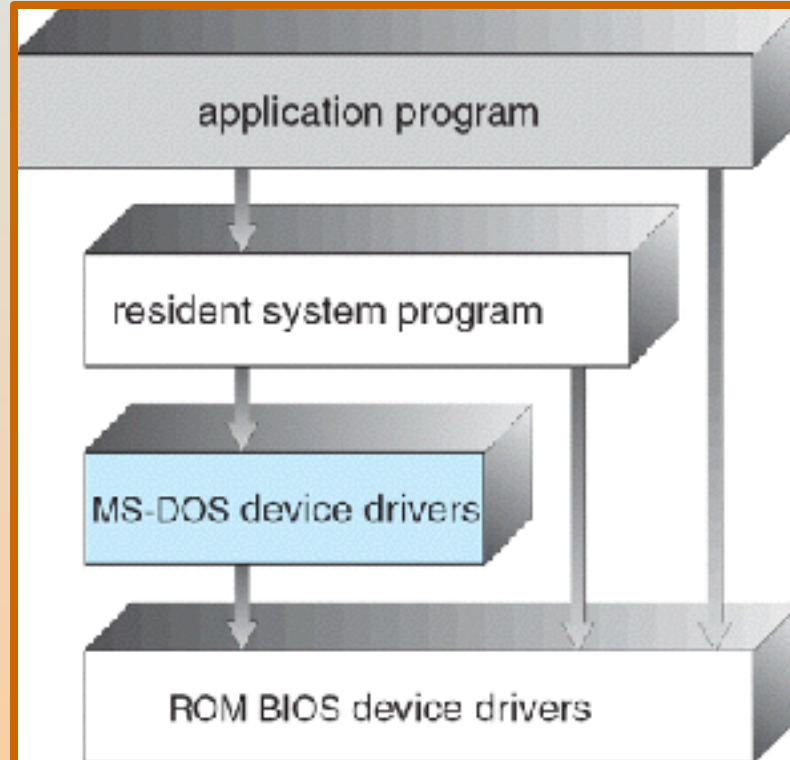
# The History of OS

- 1940s and 1950s
  - IOCS, Storage, Batch processing
- 1960s
  - Time sharing, Multiprogramming
  - IBM OS/360, Multics
- 1960s-1970s
  - UNIX
  - “the genius of the UNIX system is its framework, which enables programmers to stand on the work of others”
- 1980s
  - PC
  - Apple, IBM, CP/M, Bill Gates
  - Macintosh, Windows
- 1990s
  - Windows, Unix, Linux

# Simple Structure

7

- MS-DOS – written to provide the most functionality in the least space
  - Not divided into modules
  - Interfaces and levels of functionality not well separated



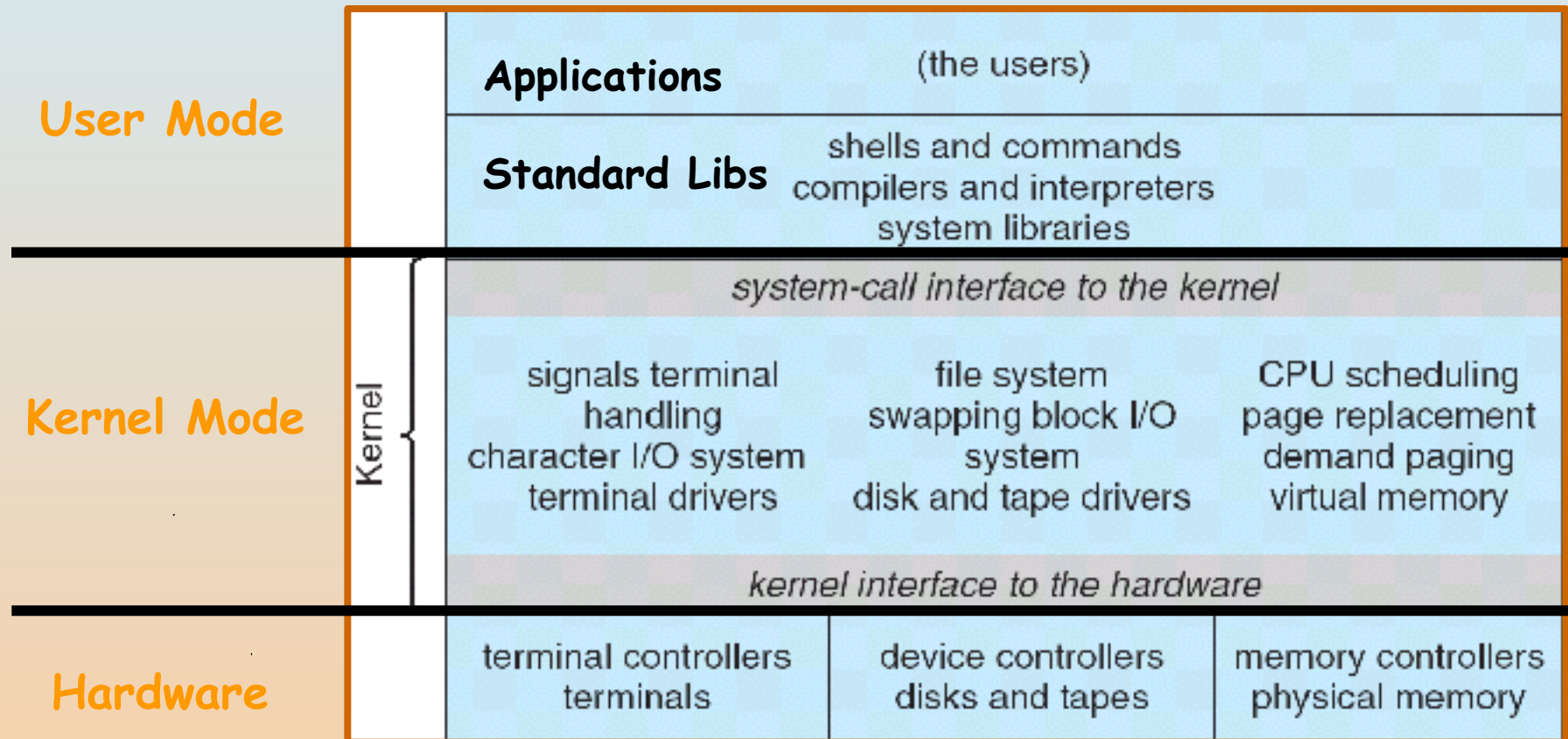
# UNIX: Also “Simple” Structure

- UNIX – limited by hardware functionality
- Original UNIX operating system consists of two separable parts:
  - Systems programs
  - The kernel
    - Consists of everything below the system-call interface and above the physical hardware
    - Provides the file system, CPU scheduling, memory management, and other operating-system functions;
    - Many interacting functions for one level



# UNIX System Structure

9



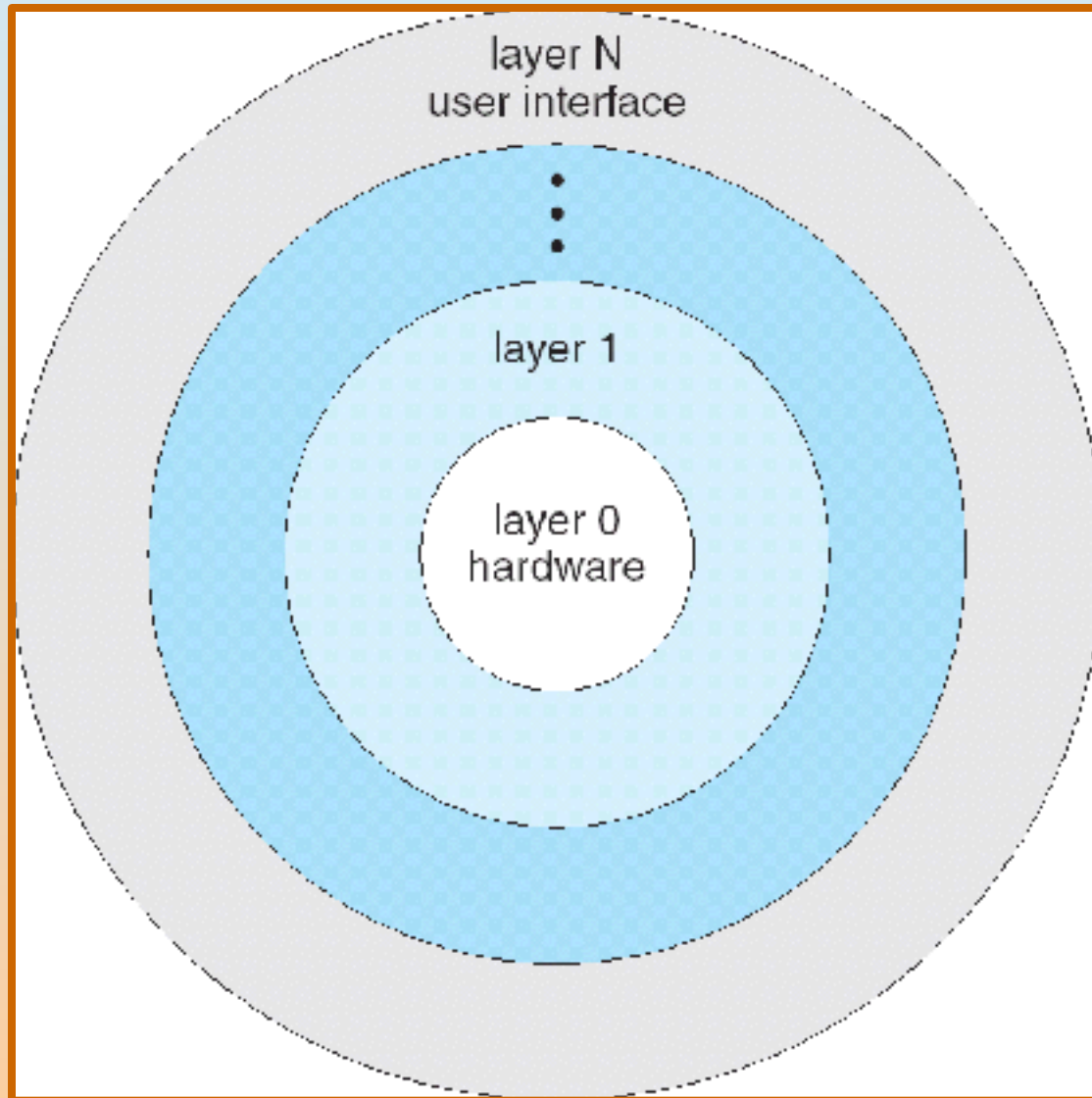
# Layered Structure

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- Operating system is divided many layers (levels)
  - Each built on top of lower layers
  - Bottom layer (layer 0) is hardware
  - Highest layer (layer N) is the user interface

# Layered Operating System

11



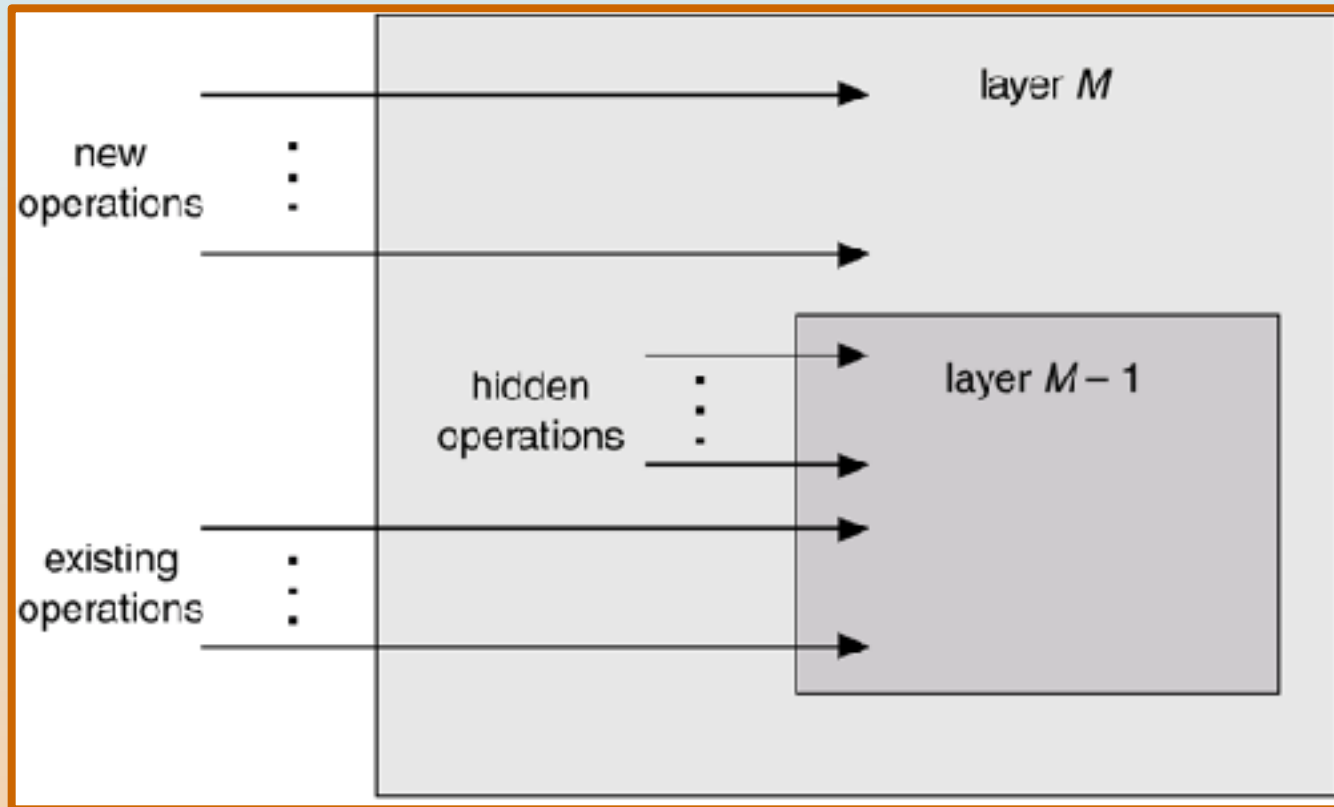
# Layered Structure

12

- Each layer uses functions (operations) and services of only lower-level layers
  - Advantage: modularity  $\Rightarrow$  Easier debugging/Maintenance
  - Not always possible: Does process scheduler lie above or below virtual memory layer?
    - Need to reschedule processor while waiting for paging
    - May need to page in information about tasks
- Machine-dependent vs independent layers
  - Easier migration between platforms
  - Easier evolution of hardware platform

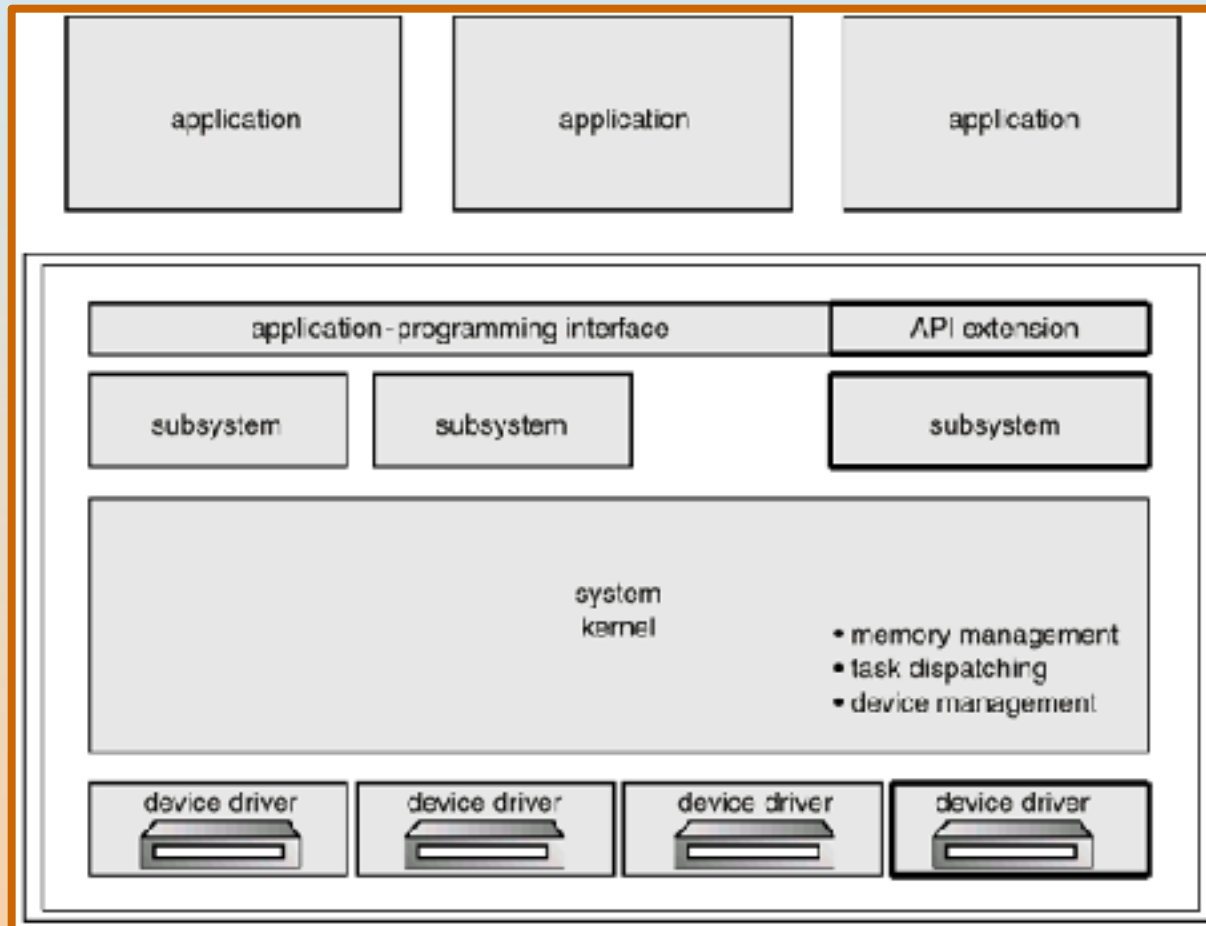
# An Operating System Layer

13



# OS/2 Layer Structure

14



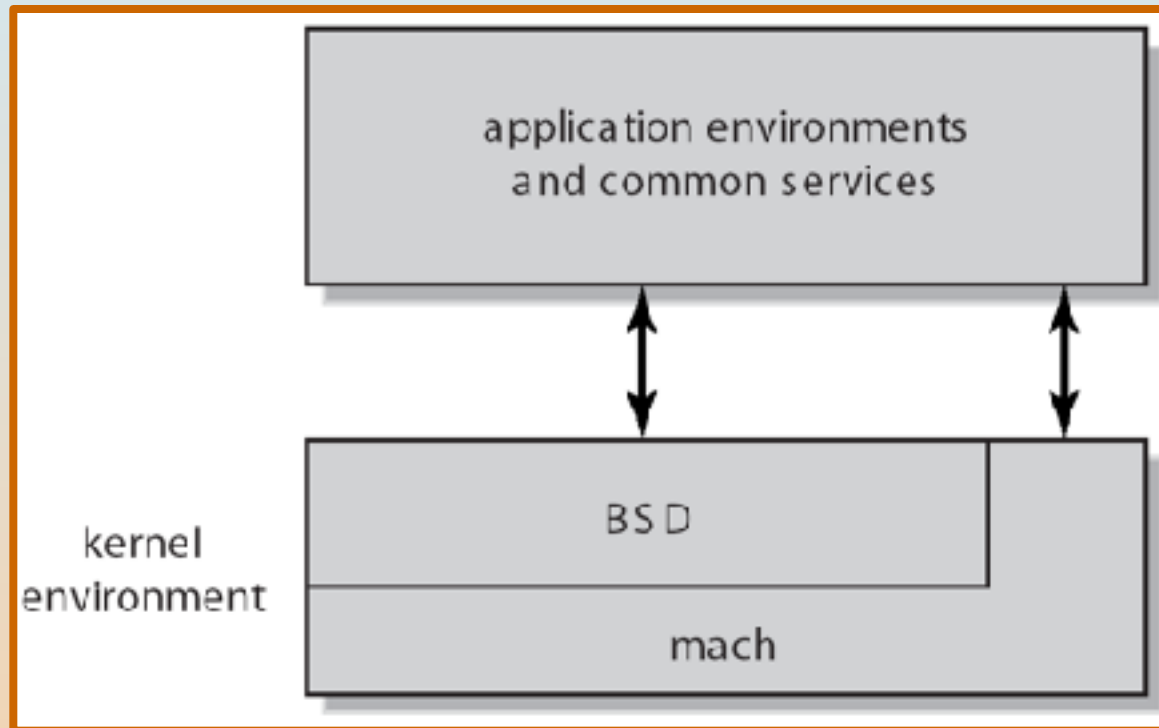
# Microkernel System Structure

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- Moves as much from the kernel into “*user*” space as possible
- Communication takes place between user modules using message passing
- Benefits:
  - Easier to extend a microkernel
  - Easier to port the operating system to new architectures
  - More reliable (less code is running in kernel mode)
  - More secure
- Detriments:
  - Performance overhead of user space to kernel space communication

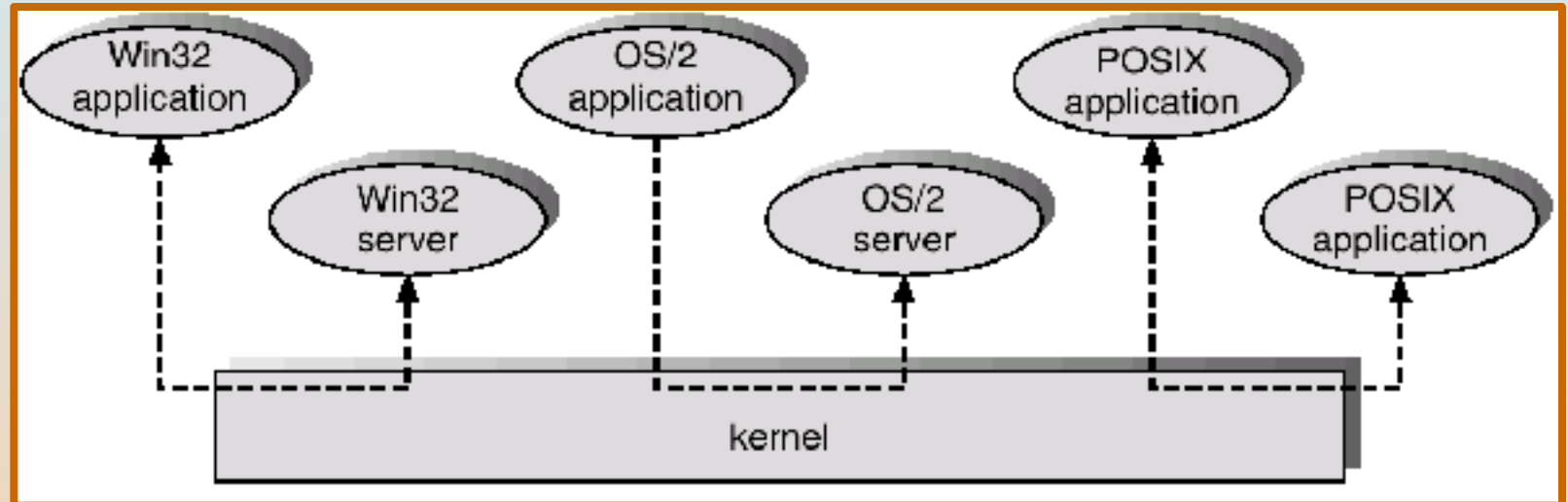
# Mac OS X Structure

16





# Windows NT Client-Server Structure 17



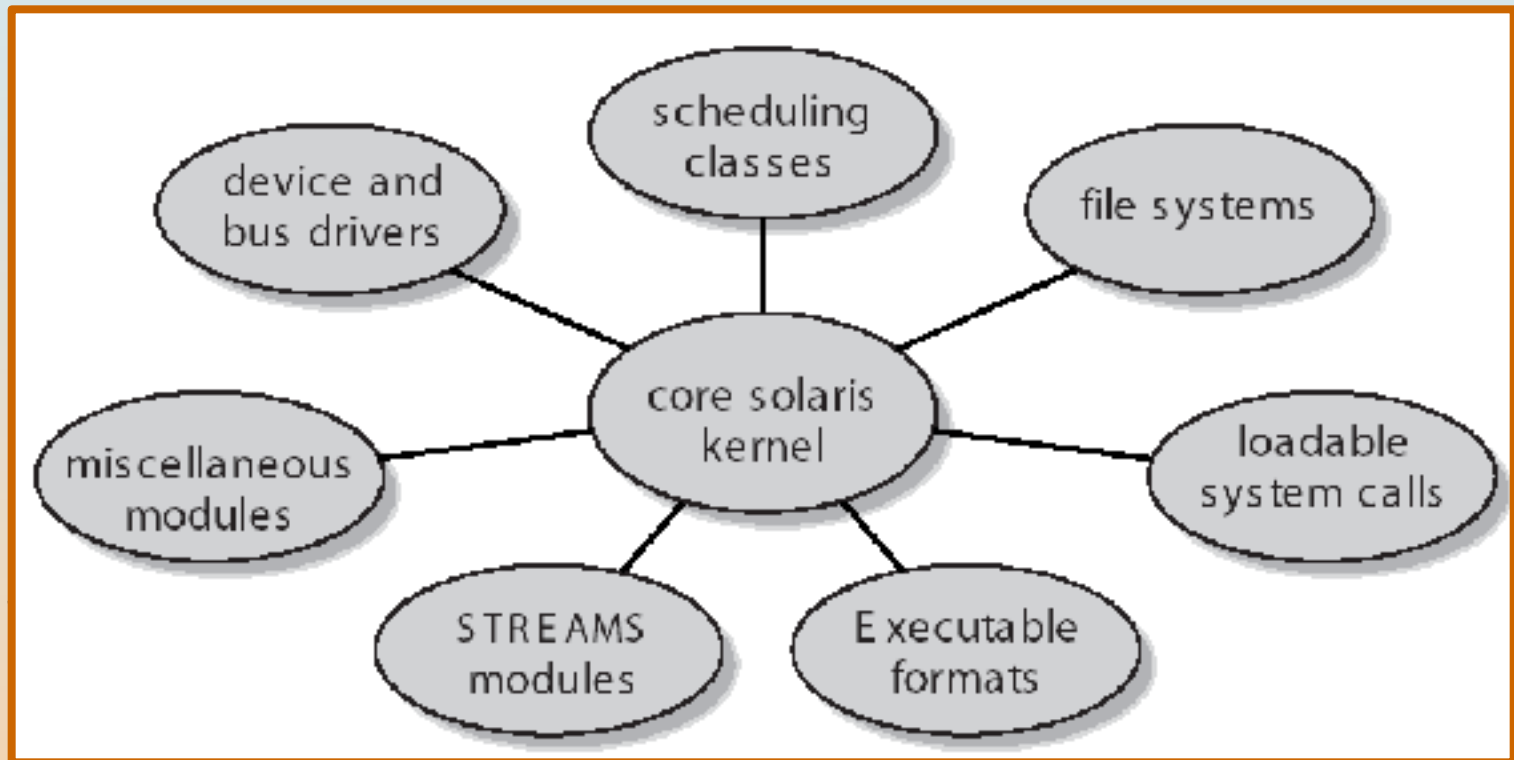
# Module based

18

- Most modern operating systems implement kernel *modules*
  - Uses object-oriented approach
  - Each core component is separate
  - Each talks to the others over known interfaces
  - Each is loadable as needed within the kernel
- Overall, similar to layers but with more flexible

# Solaris Modular Approach

19



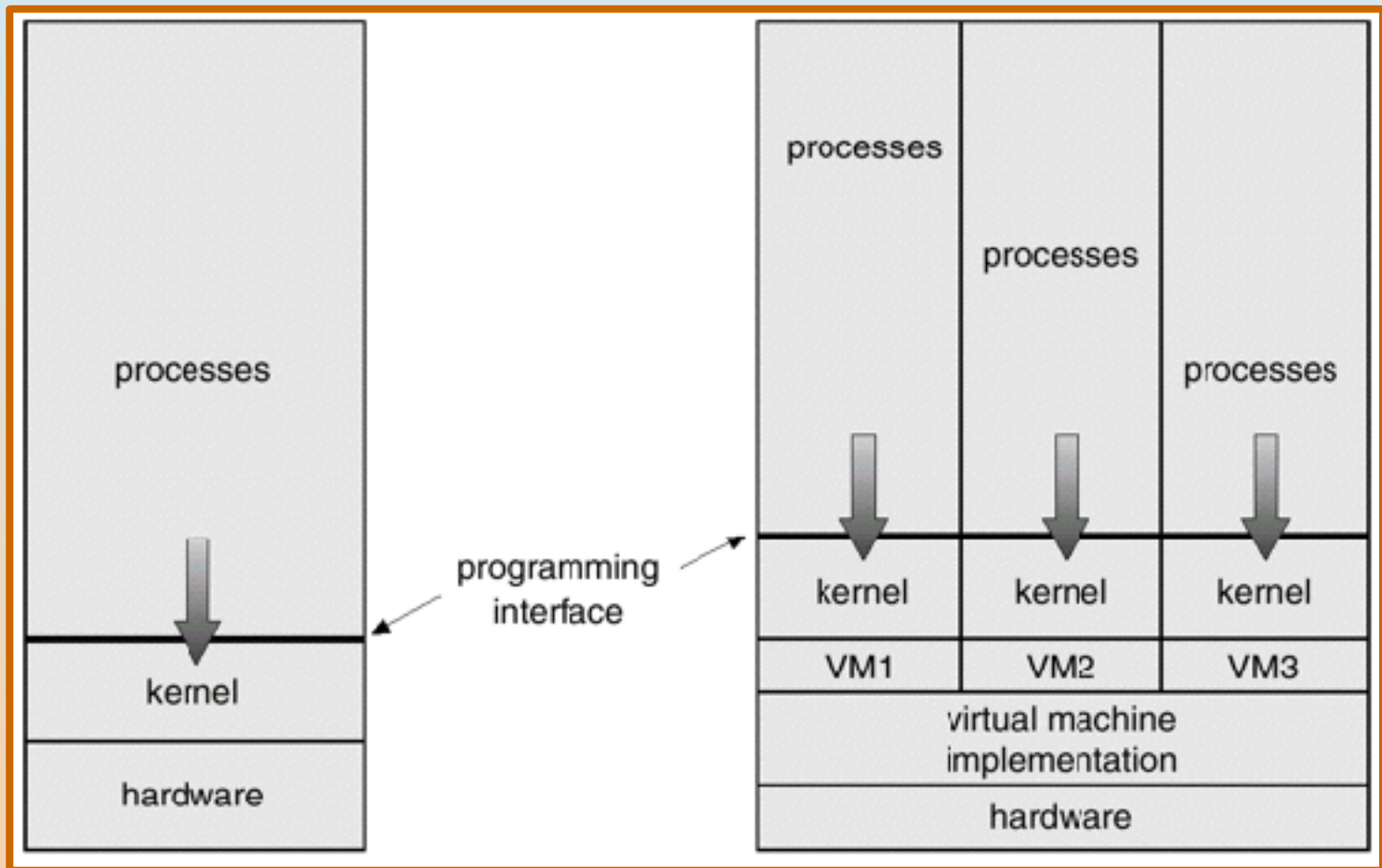
# Virtual Machines

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- A *virtual machine* takes the layered approach to its logical conclusion. It treats hardware and the operating system kernel as though they were all hardware
- A virtual machine provides an interface *identical* to the underlying bare hardware 相同
- The operating system creates the illusion of multiple processes, each executing on its own processor with its own (virtual) memory

# System Models

21



Non-virtual Machine

Virtual Machine

# Advantages/Disadvantages of Virtual Machines

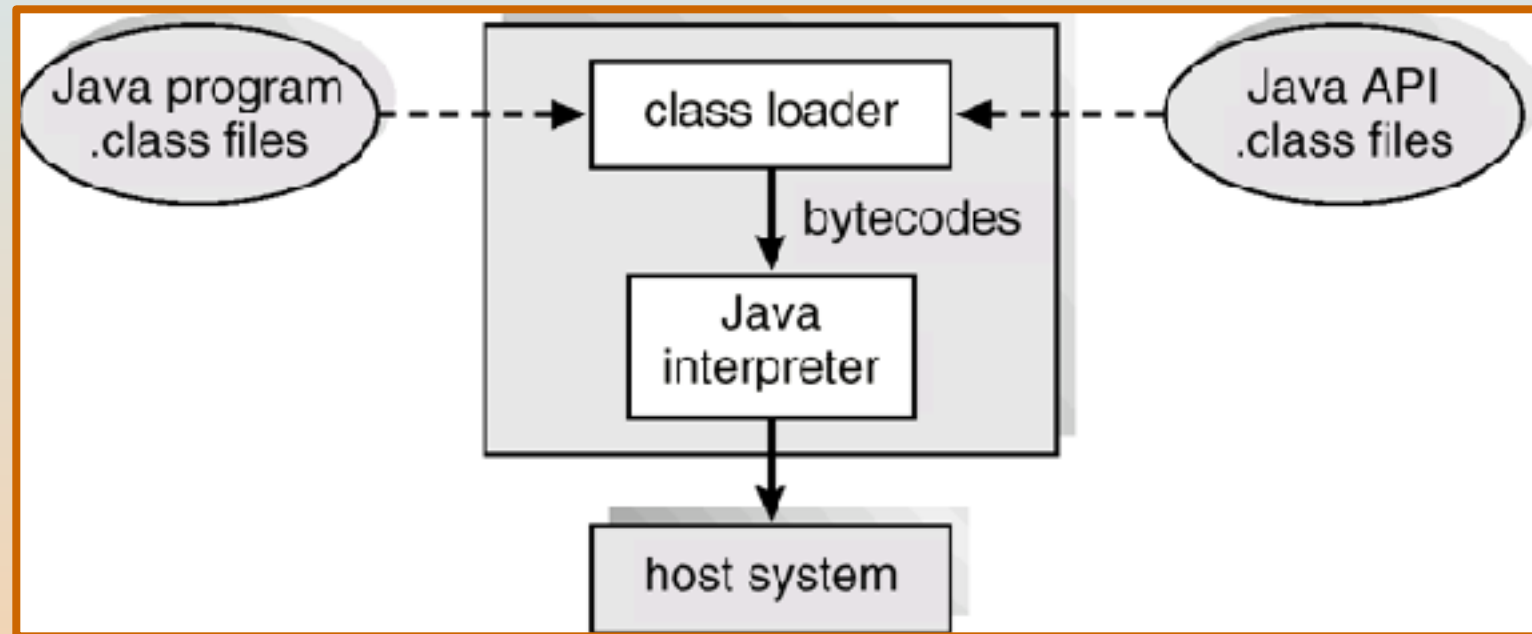
22

- Isolation from all other virtual machines.
- No disruption on normal system operation.
- Difficult to implement due to the effort required to provide an *exact* duplicate to the underlying machine

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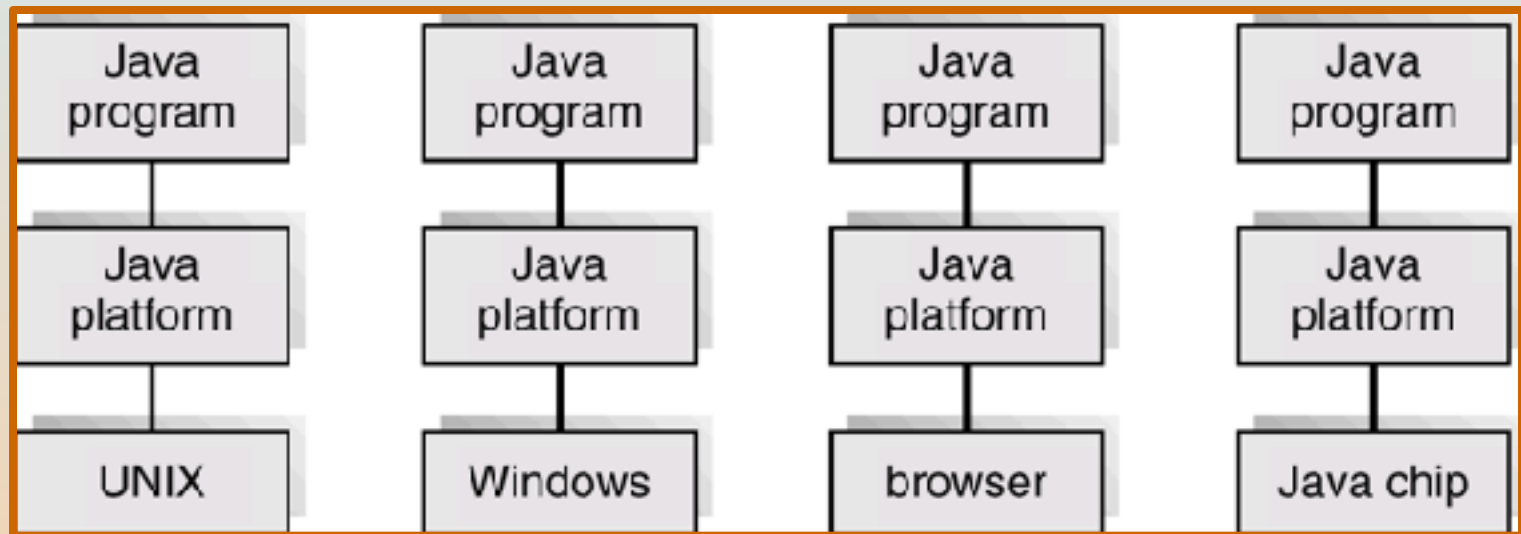
# The Java Virtual Machine

23



# The Java Platform

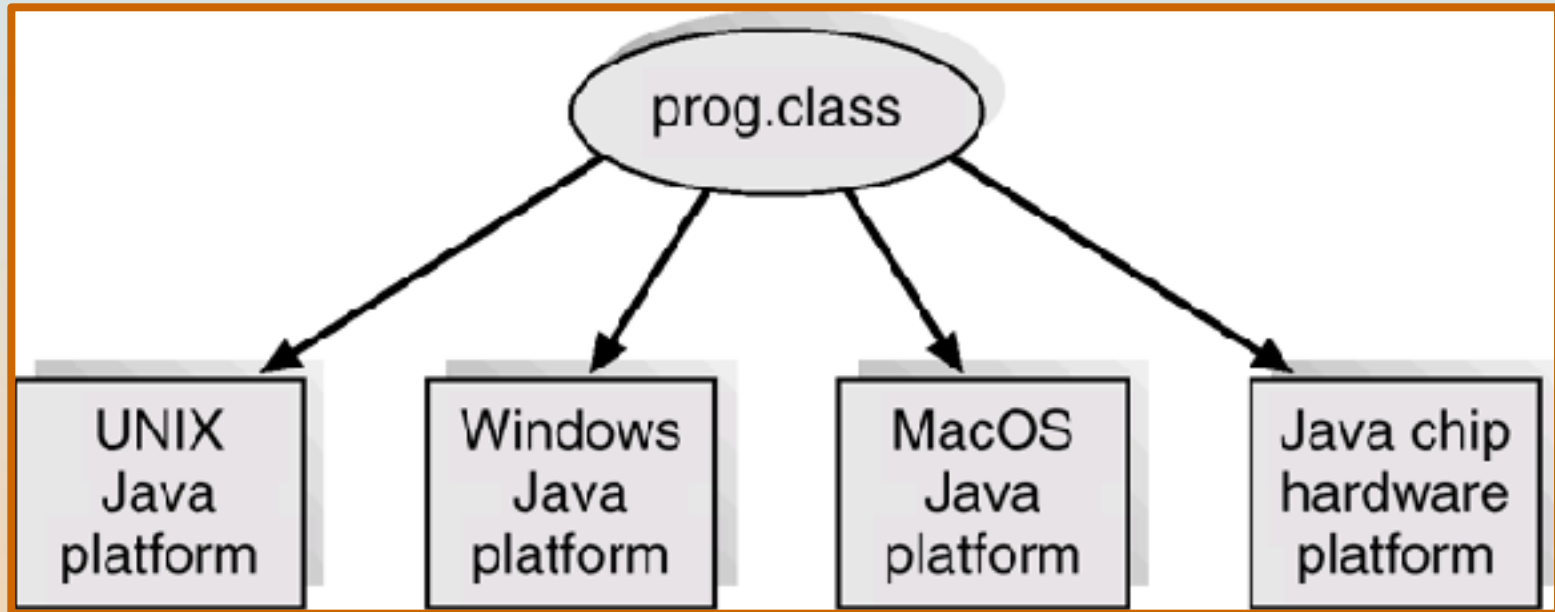
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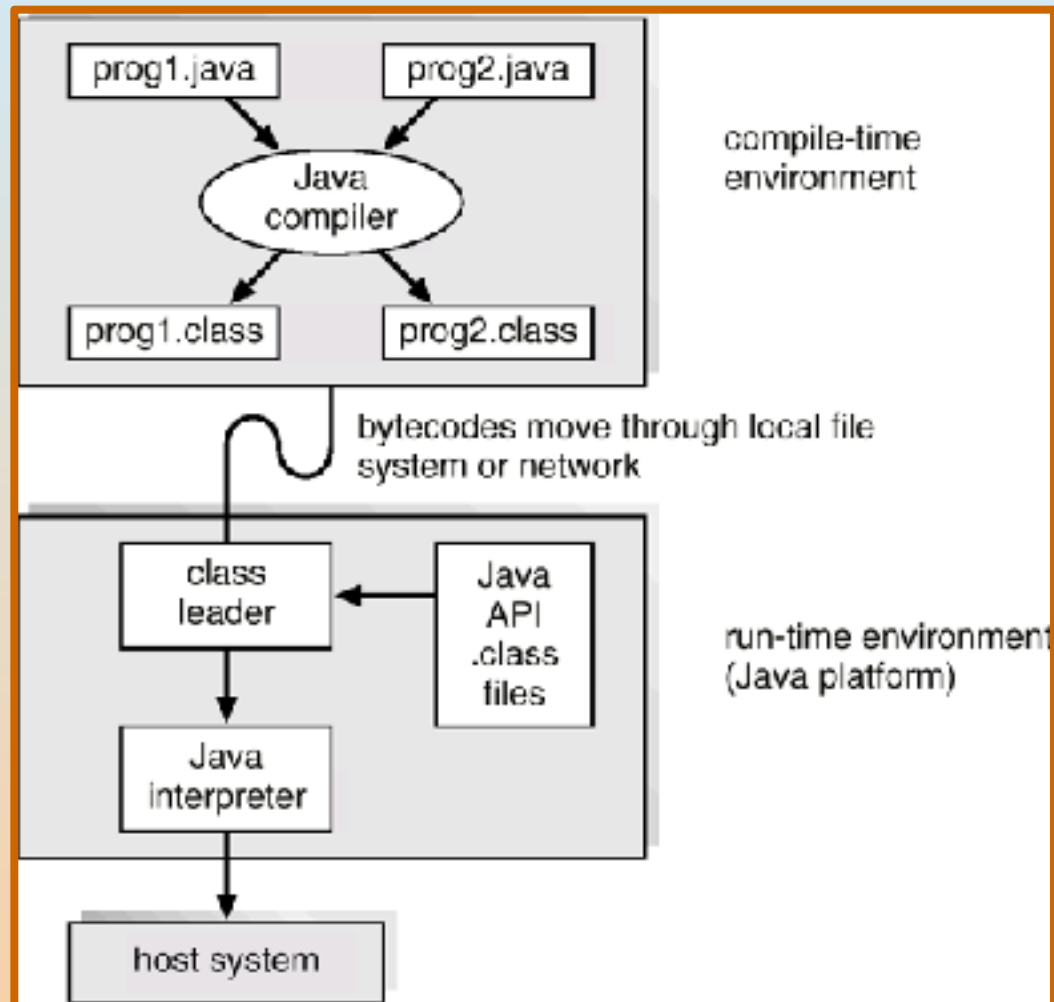
# Java .class File on Cross Platforms

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# Java Development Environment

26

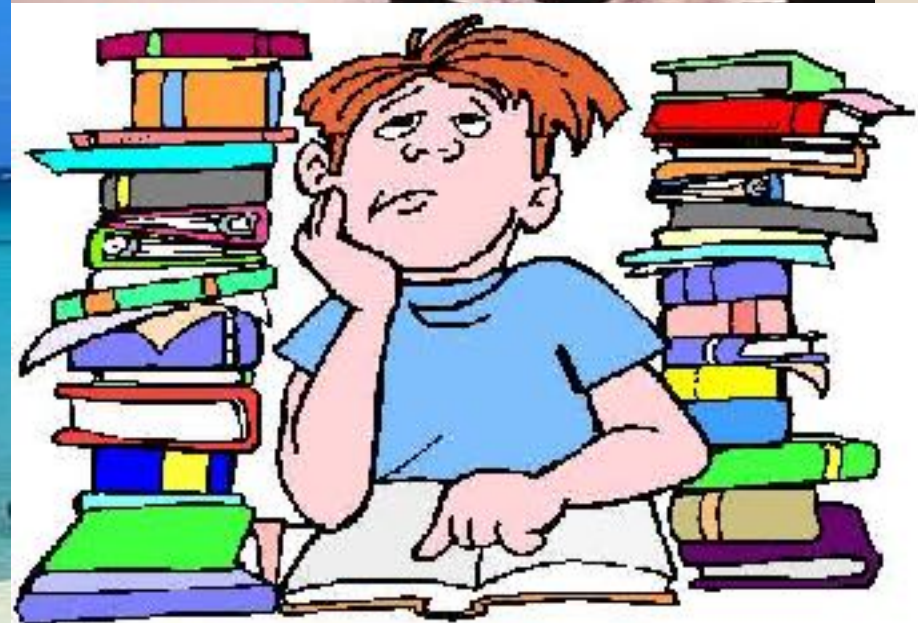


# Implementation Issues

- Policy vs. Mechanism
  - Policy: **What** will be done?
  - Mechanism: **How** to do it?
  - Should be separated, since both change
- High-level language?
- Backward compatibility issues
  - Very important for Windows 2000/XP

# Implementation Issues

- Algorithms used
  - Linear, Tree-based, Log Structured, etc...
- Event models used
  - threads vs event loops
- System generation/configuration
  - How to make generic OS fit on specific hardware
- Rapid Change in Hardware Leads to changing OS
  - Batch  $\Rightarrow$  Multiprogramming  $\Rightarrow$  Timeshare  $\Rightarrow$  Graphical UI  $\Rightarrow$  Ubiquitous Devices  $\Rightarrow$  Cyberspace/Metaverse/??



# 作业调度

30

- 作业集
- 作业发布时间, 耗费时间, 截止期, 容易程度

# 作业集

31

作业名称	发布时间	耗费时间	距截止期剩余时间	容易程度
AI, paper	4 days ago	2.5 days	25 days left	* * * *
Game design, homework	today	4 day	20 days left	* * * * *
Stat Learning, project	2 days ago	6 days	20 days left	* *
Adv. Prog, midterm	yesterday	3 days	10 days left	*
Senior soft. & Engi., project	today	5 days	50 days left	* * *

# 调度策略

32

- First come first serve
- Easiest first
- Shortest time-cost first
- Earliest deadline first



# First come first serve

33

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# Easiest first

34

作业名称	发布时间	耗费时间	距截止期剩余时间	容易程度
AI, paper			25 days left	* * * *
Game design, homework			20 days left	* * * * *
Stat Learning, project			20 days left	* *
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Senior soft. & Engi., project			50 days left	* * *

# Shortest time-cost first

35

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# Earliest deadline first

36

作业名称	发布时间	耗费时间	距截止期剩余时间	容易程度
AI, paper			25 days left	
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Stat Learning, project			20 days left	
Adv. Prog, midterm			10 days left	
Senior soft. & Engi., project			50 days left	

# 最优策略?

37

- First come first serve?
- Easiest first?
- Shortest time-cost first?
- Earliest deadline first?

# 如何判断策略是否可行?

38

- 可调度性
  - The tasks are **schedulable** if there exists a scheduling solution such that all the tasks can be scheduled to meet their deadlines

- 可调度性测试

- $U =$

$$\sum_{i=1}^n \frac{T_{\text{cost}}(i)}{T_{\text{remains}}(i)}$$

The tasks are **schedulable** if  **$U \leq 1$**  !

# 可调度性

39

作业名称	发布时间	耗费时间	距截止期剩余时间	容易程度
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# Earliest deadline first (EDF)

40

- Order jobs by deadline
- EDF is optimal
  - EDF can always produce a feasible schedule for a set of tasks if they are **schedulable** ( $U \leq 1$ ).



# Question

41

你的作业是可调度的么？

- Schedulability test result ?

# 延伸思考

42

- What about the overload case ( $U > 1$ )?
- What if the objective is to minimize the sum of the lateness?
  - EDF does not seem to work