

CS 325 Project 3: Linear Programming

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Problem 1

Problem 1 Part A

Determine the number of refrigerators to be shipped from plants to warehouses, and then warehouses to retailers to minimize the cost.

Solution

- In all 1000 units will travel through the network at a minimum cost of \$17100.
- Ship 150 units from plant #1 to warehouse #1 at a cost of \$1500.
- Ship 200 units from plant #2 to warehouse #1 at a cost of \$2200.
- Ship 250 units from plant #2 to warehouse #2 at a cost of \$2000.
- Ship 150 units from plant #3 to warehouse #2 at a cost of \$1200.
- Ship 100 units from plant #3 to warehouse #3 at a cost of \$ 900.
- Ship 150 units from plant #4 to warehouse #3 at a cost of \$1200.
- Ship 100 units from warehouse #1 to retailer #1 at a cost of \$ 500.
- Ship 150 units from warehouse #1 to retailer #2 at a cost of \$ 900.
- Ship 100 units from warehouse #1 to retailer #3 at a cost of \$ 700.
- Ship 200 units from warehouse #2 to retailer #4 at a cost of \$1600.
- Ship 200 units from warehouse #2 to retailer #5 at a cost of \$2000.
- Ship 150 units from warehouse #3 to retailer #6 at a cost of \$1800.
- Ship 100 units from warehouse #3 to retailer #7 at a cost of \$ 600.
- 150 total units will leave plant #1 (capacity is 150).
- 450 total units will leave plant #2 (capacity is 450).
- 250 total units will leave plant #3 (capacity is 250).
- 150 total units will leave plant #4 (capacity is 150).
- 350 total units will enter warehouse #1, 350 units will leave.

- 400 total units will enter warehouse #2, 400 units will leave.
- 250 total units will enter warehouse #3, 250 units will leave.
- 100 total units will enter retailer #1 (demand is 100).
- 150 total units will enter retailer #2 (demand is 150).
- 100 total units will enter retailer #3 (demand is 100).
- 200 total units will enter retailer #4 (demand is 200).
- 200 total units will enter retailer #5 (demand is 200).
- 150 total units will enter retailer #6 (demand is 150).
- 100 total units will enter retailer #7 (demand is 100).

Linear Program Formulation

1. Overall idea of problem

- Refrigerators moving from $n = 4$ plants to $q = 3$ warehouses to $m = 7$ retailers.
- Not all plants deliver to all warehouses.
- Not all warehouses deliver to all retailers.
- Costs of shipping from plants to warehouses vary by pair.
- Costs of shipping from warehouses to retailers vary by pair.
- Each plant has a capacity in terms of number of refrigerators it can supply.
- Each retailer has a capacity in terms of number of refrigerators it demands.

2. What is the goal? What are you trying to achieve?

- Determine optimal shipping routes (n to q and q to m).
- Determine number of refrigerators moving along each route (n to q and q to m).
- Satisfy the demand of the retailers.
- Minimize the cost.

3. Identify variables

- cp_{ij} = cost of moving a refrigerator between plant i and warehouse j
 - ex. $cp_{32} = 8$ = cost of moving from plant 3 to warehouse 2
 - 9 variables
- cw_{jk} = cost of moving a refrigerator between warehouse j and retailer k
 - ex. $cp_{14} = 10$ = cost of moving from warehouse 1 to retailer 4

- 12 variables
- s_i = capacity (supply) of each plant
 - ex. $s_2 = 450$ = number of refrigerators that plant 2 can supply
 - 4 variables
- d_k = capacity (demand) of each retailer
 - ex. $d_6 = 150$ = number of refrigerators that plant 6 demands
 - 7 variables
- np_{ij} = number of refrigerators shipped from plant i to warehouse j
 - 9 variables
- nw_{jk} = number of refrigerators shipped from warehouse j to retailer k
 - 12 variables

4. Identify constraints

- $s_1 \leq 150$
- $s_2 \leq 450$
- $s_3 \leq 250$
- $s_4 \leq 150$
- $d_1 \geq 100$
- $d_2 \geq 150$
- $d_3 \geq 100$
- $d_4 \geq 200$
- $d_5 \geq 200$
- $d_6 \geq 150$
- $d_7 \geq 100$
- $np_{11} + np_{21} + np_{31} = nw_{11} + nw_{12} + nw_{13} + nw_{14}$
- $np_{12} + np_{22} + np_{32} + np_{42} = nw_{23} + nw_{24} + nw_{25} + nw_{26}$
- $np_{33} + np_{43} = nw_{34} + nw_{35} + nw_{36} + nw_{37}$
- $s_1 = np_{11} + np_{12}$
- $s_2 = np_{21} + np_{22}$
- $s_3 = np_{31} + np_{32} + np_{33}$
- $s_4 = np_{42} + np_{43}$
- $d_1 = nw_{11}$
- $d_2 = nw_{12}$
- $d_3 = nw_{13} + nw_{23}$
- $d_4 = nw_{14} + nw_{24} + nw_{34}$

- $d_5 = nw_{25} + nw_{35}$
- $d_6 = nw_{26} + nw_{36}$
- $d_7 = nw_{37}$
- $np_{11} \geq 0$
- $np_{12} \geq 0$
- $np_{21} \geq 0$
- $np_{22} \geq 0$
- $np_{31} \geq 0$
- $np_{32} \geq 0$
- $np_{33} \geq 0$
- $np_{42} \geq 0$
- $np_{43} \geq 0$
- $nw_{11} \geq 0$
- $nw_{12} \geq 0$
- $nw_{13} \geq 0$
- $nw_{14} \geq 0$
- $nw_{23} \geq 0$
- $nw_{24} \geq 0$
- $nw_{25} \geq 0$
- $nw_{26} \geq 0$
- $nw_{34} \geq 0$
- $nw_{35} \geq 0$
- $nw_{36} \geq 0$
- $nw_{37} \geq 0$

5. Identify inputs and outputs that you can control

- np_{ij}
- nw_{jk}
- $cost$

6. Specify all quantities mathematically

- Many have been defined above already. A few more will be added here.
- $cost = \sum_{i=1}^n \sum_{j=1}^q np_{ij} * cp_{ij} + \sum_{j=1}^q \sum_{k=1}^m nw_{jk} * cw_{jk}$
- $cp_{11} = 10$
- $cp_{12} = 15$

- $cp_{21} = 11$
- $cp_{22} = 8$
- $cp_{31} = 13$
- $cp_{32} = 8$
- $cp_{33} = 9$
- $cp_{42} = 14$
- $cp_{43} = 8$
- $cw_{11} = 5$
- $cw_{12} = 6$
- $cw_{13} = 7$
- $cw_{14} = 10$
- $cw_{23} = 12$
- $cw_{24} = 8$
- $cw_{25} = 10$
- $cw_{26} = 14$
- $cw_{34} = 14$
- $cw_{35} = 12$
- $cw_{36} = 12$
- $cw_{37} = 6$

7. Check the model for completeness and correctness

- All variables are positive.

Matlab Code

```

1 % -----
2 % reference: array index to variable mapping
3 % -----
4 % (1) np11
5 % (2) np12
6 % (3) np21
7 % (4) np22
8 % (5) np31
9 % (6) np32
10 % (7) np33
11 % (8) np42
12 % (9) np43
13 % (10) nw11
14 % (11) nw12
15 % (12) nw13

```

```

16 % (13) nw14
17 % (14) nw23
18 % (15) nw24
19 % (16) nw25
20 % (17) nw26
21 % (18) nw34
22 % (19) nw35
23 % (20) nw36
24 % (21) nw37
25 % (22) s1
26 % (23) s2
27 % (24) s3
28 % (25) s4
29 % (26) d1
30 % (27) d2
31 % (28) d3
32 % (29) d4
33 % (30) d5
34 % (31) d6
35 % (32) d7
36
37 % -----
38 % lower bounds vector
39 %   note matlab arrays/vectors start at index 1 (not 0)
40 % -----
41 lb = zeros(32,1);
42 lb(26) = 100;    % d1
43 lb(27) = 150;    % d2
44 lb(28) = 100;    % d3
45 lb(29) = 200;    % d4
46 lb(30) = 200;    % d5
47 lb(31) = 150;    % d6
48 lb(32) = 100;    % d7
49
50 % -----
51 % upper bounds vector
52 %   note matlab arrays/vectors start at index 1 (not 0)
53 % -----
54 ub = Inf(32,1);
55 ub(1)  = 150;    % np11
56 ub(2)  = 150;    % np12
57 ub(3)  = 450;    % np21
58 ub(4)  = 450;    % np22
59 ub(5)  = 250;    % np31
60 ub(6)  = 250;    % np32
61 ub(7)  = 250;    % np33
62 ub(8)  = 150;    % np42
63 ub(9)  = 150;    % np43
64 ub(22) = 150;    % s1
65 ub(23) = 450;    % s2
66 ub(24) = 250;    % s3
67 ub(25) = 150;    % s4
68
69 % -----

```

```

70 % linear inequality matrix and vector
71 %   note matlab arrays/vectors start at index 1 (not 0)
72 % -----
73 A = [];
74 b = [];
75
76 % -----
77 % linear equality matrix and vector
78 %   note matlab arrays/vectors start at index 1 (not 0)
79 %   14 equations in 32 variables
80 % -----
81 Aeq = zeros(14, 32);
82 beq = zeros(14, 1);
83 %np11 + np21 + np31          = nw11 + nw12 + nw13 + nw14
84 %np11 + np21 + np31 - nw11 - nw12 - nw13 - nw14 = 0
85 Aeq(1, [1, 3, 5, 10, 11, 12, 13]) = [1, 1, 1, -1, -1, -1, -1];
86 %np12 + np22 + np32 + np42 =          nw23 + nw24 + nw25 + nw26
87 %np12 + np22 + np32 + np42 - nw23 - nw24 - nw25 - nw26 = 0
88 Aeq(2, [2, 4, 6, 8, 14, 15, 16, 17]) = [1, 1, 1, 1, -1, -1, -1, -1];
89 %          np33 + np43 =          nw34 + nw35 + nw36 + nw37
90 %np33 + np43 - nw34 - nw35 - nw36 - nw37 = 0
91 Aeq(3, [7, 9, 18, 19, 20, 21]) = [1, 1, -1, -1, -1, -1];
92 %s1 = np11 + np12
93 %s1 - np11 - np12 = 0
94 Aeq(4, [22, 1, 2]) = [1, -1, -1];
95 %s2 = np21 + np22
96 %s2 - np21 - np22 = 0
97 Aeq(5, [23, 3, 4]) = [1, -1, -1];
98 %s3 = np31 + np32 + np33
99 %s3 - np31 - np32 - np33 = 0
100 Aeq(6, [24, 5, 6, 7]) = [1, -1, -1, -1];
101 %s4 =          np42 + np43
102 %s4 - np42 - np43 = 0
103 Aeq(7, [25, 8, 9]) = [1, -1, -1];
104 %d1 = nw11
105 %d1 - nw11 = 0
106 Aeq(8, [26, 10]) = [1, -1];
107 %d2 = nw12
108 %d2 - nw12 = 0
109 Aeq(9, [27, 11]) = [1, -1];
110 %d3 = nw13 + nw23
111 %d3 - nw13 - nw23 = 0
112 Aeq(10, [28, 12, 14]) = [1, -1, -1];
113 %d4 = nw14 + nw24 + nw34
114 %d4 - nw14 - nw24 - nw34 = 0
115 Aeq(11, [29, 13, 15, 18]) = [1, -1, -1, -1];
116 %d5 =          nw25 + nw35
117 %d5 - nw25 - nw35 = 0
118 Aeq(12, [30, 16, 19]) = [1, -1, -1];
119 %d6 =          nw26 + nw36
120 %d6 - nw26 - nw36 = 0
121 Aeq(13, [31, 17, 20]) = [1, -1, -1];
122 %d7 =          nw37
123 %d7 - nw37 = 0

```



```

124 Aeq(14,[32,21]) = [1,-1];
125
126 % -----
127 % objective function vector
128 %   note matlab arrays/vectors start at index 1 (not 0)
129 % -----
130 f = zeros(32,1);
131 f(1) = 10; % np11(value in f is cp11)
132 f(2) = 15; % np12(value in f is cp12)
133 f(3) = 11; % np21(value in f is cp21)
134 f(4) = 8; % np22(value in f is cp22)
135 f(5) = 13; % np31(value in f is cp31)
136 f(6) = 8; % np32(value in f is cp32)
137 f(7) = 9; % np33(value in f is cp33)
138 f(8) = 14; % np42(value in f is cp42)
139 f(9) = 8; % np43(value in f is cp43)
140 f(10) = 5; % nw11(value in f is cw11)
141 f(11) = 6; % nw12(value in f is cw12)
142 f(12) = 7; % nw13(value in f is cw13)
143 f(13) = 10; % nw14(value in f is cw14)
144 f(14) = 12; % nw23(value in f is cw23)
145 f(15) = 8; % nw24(value in f is cw24)
146 f(16) = 10; % nw25(value in f is cw25)
147 f(17) = 14; % nw26(value in f is cw26)
148 f(18) = 14; % nw34(value in f is cw34)
149 f(19) = 12; % nw35(value in f is cw35)
150 f(20) = 12; % nw36(value in f is cw36)
151 f(21) = 6; % nw37(value in f is cw37)
152
153 % -----
154 % call solver and obtain solution
155 % -----
156 [x fval] = linprog(f,A,b,Aeq,beq,lb,ub);
157
158 % -----
159 % print the optimum shipping routes and min cost
160 % -----
161 fileID = fopen('partA.out','w');
162 fprintf(fileID, '-----\n');
163 fprintf(fileID, 'Project 3 Problem 1 Part A Solution\n');
164 fprintf(fileID, '-----\n');
165 fprintf(fileID, '\n');
166 fprintf(fileID, ['Ship %3.0f units from plant #%d to warehouse #%d ', ...
167     'at a cost of $%4.0f.\n'], x(1), 1, 1, x(1) * f(1));
168 fprintf(fileID, ['Ship %3.0f units from plant #%d to warehouse #%d ', ...
169     'at a cost of $%4.0f.\n'], x(2), 1, 2, x(2) * f(2));
170 fprintf(fileID, ['Ship %3.0f units from plant #%d to warehouse #%d ', ...
171     'at a cost of $%4.0f.\n'], x(3), 2, 1, x(3) * f(3));
172 fprintf(fileID, ['Ship %3.0f units from plant #%d to warehouse #%d ', ...
173     'at a cost of $%4.0f.\n'], x(4), 2, 2, x(4) * f(4));
174 fprintf(fileID, ['Ship %3.0f units from plant #%d to warehouse #%d ', ...
175     'at a cost of $%4.0f.\n'], x(5), 3, 1, x(5) * f(5));
176 fprintf(fileID, ['Ship %3.0f units from plant #%d to warehouse #%d ', ...
177     'at a cost of $%4.0f.\n'], x(6), 3, 2, x(6) * f(6));

```

```

178 fprintf(fileID, ['Ship %3.0f units from plant #%d to warehouse #%d ', ...
179     'at a cost of $%4.0f.\n'], x(7), 3, 3, x(7) * f(7));
180 fprintf(fileID, ['Ship %3.0f units from plant #%d to warehouse #%d ', ...
181     'at a cost of $%4.0f.\n'], x(8), 4, 2, x(8) * f(8));
182 fprintf(fileID, ['Ship %3.0f units from plant #%d to warehouse #%d ', ...
183     'at a cost of $%4.0f.\n'], x(9), 4, 3, x(9) * f(9));
184 fprintf(fileID, '\n');
185 fprintf(fileID, ['Ship %3.0f units from warehouse #%d to retailer ', ...
186     '#%d at a cost of $%4.0f.\n'], x(10), 1, 1, x(10) * f(10));
187 fprintf(fileID, ['Ship %3.0f units from warehouse #%d to retailer ', ...
188     '#%d at a cost of $%4.0f.\n'], x(11), 1, 2, x(11) * f(11));
189 fprintf(fileID, ['Ship %3.0f units from warehouse #%d to retailer ', ...
190     '#%d at a cost of $%4.0f.\n'], x(12), 1, 3, x(12) * f(12));
191 fprintf(fileID, ['Ship %3.0f units from warehouse #%d to retailer ', ...
192     '#%d at a cost of $%4.0f.\n'], x(13), 1, 4, x(13) * f(13));
193 fprintf(fileID, ['Ship %3.0f units from warehouse #%d to retailer ', ...
194     '#%d at a cost of $%4.0f.\n'], x(14), 2, 3, x(14) * f(14));
195 fprintf(fileID, ['Ship %3.0f units from warehouse #%d to retailer ', ...
196     '#%d at a cost of $%4.0f.\n'], x(15), 2, 4, x(15) * f(15));
197 fprintf(fileID, ['Ship %3.0f units from warehouse #%d to retailer ', ...
198     '#%d at a cost of $%4.0f.\n'], x(16), 2, 5, x(16) * f(16));
199 fprintf(fileID, ['Ship %3.0f units from warehouse #%d to retailer ', ...
200     '#%d at a cost of $%4.0f.\n'], x(17), 2, 6, x(17) * f(17));
201 fprintf(fileID, ['Ship %3.0f units from warehouse #%d to retailer ', ...
202     '#%d at a cost of $%4.0f.\n'], x(18), 3, 4, x(18) * f(18));
203 fprintf(fileID, ['Ship %3.0f units from warehouse #%d to retailer ', ...
204     '#%d at a cost of $%4.0f.\n'], x(19), 3, 5, x(19) * f(19));
205 fprintf(fileID, ['Ship %3.0f units from warehouse #%d to retailer ', ...
206     '#%d at a cost of $%4.0f.\n'], x(20), 3, 6, x(20) * f(20));
207 fprintf(fileID, ['Ship %3.0f units from warehouse #%d to retailer ', ...
208     '#%d at a cost of $%4.0f.\n'], x(21), 3, 7, x(21) * f(21));
209 fprintf(fileID, '\n');
210 fprintf(fileID, ['%3.0f total units will leave plant #%d (capacity ', ...
211     'is %3.0f).\n'], x(1) + x(2), 1, x(22));
212 fprintf(fileID, ['%3.0f total units will leave plant #%d (capacity ', ...
213     'is %3.0f).\n'], x(3) + x(4), 2, x(23));
214 fprintf(fileID, ['%3.0f total units will leave plant #%d (capacity ', ...
215     'is %3.0f).\n'], x(5) + x(6) + x(7), 3, x(24));
216 fprintf(fileID, ['%3.0f total units will leave plant #%d (capacity ', ...
217     'is %3.0f).\n'], x(8) + x(9), 4, x(25));
218 fprintf(fileID, '\n');
219 temp = x(10) + x(11) + x(12) + x(13);
220 fprintf(fileID, ['%3.0f total units will enter warehouse #%d, %3.0f ', ...
221     'units will leave.\n'], x(1) + x(3) + x(5), 1, temp);
222 temp = x(14) + x(15) + x(16) + x(17);
223 fprintf(fileID, ['%3.0f total units will enter warehouse #%d, %3.0f ', ...
224     'units will leave.\n'], x(2) + x(4) + x(6) + x(8), 2, temp);
225 temp = x(18) + x(19) + x(20) + x(21);
226 fprintf(fileID, ['%3.0f total units will enter warehouse #%d, %3.0f ', ...
227     'units will leave.\n'], x(7) + x(9), 3, temp);
228 fprintf(fileID, '\n');
229 fprintf(fileID, ['%3.0f total units will enter retailer #%d (demand ', ...
230     'is %3.0f).\n'], x(10), 1, x(26));
231 fprintf(fileID, ['%3.0f total units will enter retailer #%d (demand ', ...

```

```

232     'is %3.0f).\n'], x(11), 2, x(27));
233 fprintf(fileID, ['%3.0f total units will enter retailer #%d (demand ', ...
234     'is %3.0f).\n'], x(12) + x(14), 3, x(28));
235 fprintf(fileID, ['%3.0f total units will enter retailer #%d (demand ', ...
236     'is %3.0f).\n'], x(13) + x(15) + x(18), 4, x(29));
237 fprintf(fileID, ['%3.0f total units will enter retailer #%d (demand ', ...
238     'is %3.0f).\n'], x(16) + x(19), 5, x(30));
239 fprintf(fileID, ['%3.0f total units will enter retailer #%d (demand ', ...
240     'is %3.0f).\n'], x(17) + x(20), 6, x(31));
241 fprintf(fileID, ['%3.0f total units will enter retailer #%d (demand ', ...
242     'is %3.0f).\n'], x(21), 7, x(32));
243 fprintf(fileID, '\n');
244 total = x(22) + x(23) + x(24) + x(25);
245 fprintf(fileID, ['In all %3.0f units will travel through the network ', ...
246     'at a minimum cost of $%5.0f.\n'], total, fval);
247 fclose(fileID);

```

Problem 1 Part B

Determine the number of refrigerators to be shipped from plants to warehouses, and then warehouses to retailers to minimize the cost. For part B warehouse 2 is closed along with all associated routes. Changes to the problem statement or solution relative to Problem 1 Part A are highlighted in red.

Solution

There is no solution when warehouse 2 is closed. The following is the error message that is returned by Matlab function `linprog()`:

Exiting: One or more of the residuals, duality gap, or total relative error has grown 100000 times greater than its minimum value so far:

the primal appears to be infeasible (and the dual unbounded).
(The dual residual < TolFun=1.00e-08.)

So Matlab tells us that there is no feasible solution. Why is that? If you look back at the network diagram and the supply and demand tables, you'll note that with warehouse 2 out of commission, retailers 5, 6, and 7 can only receive shipments from warehouse 3 and their total demand is 450 units. At the same time, warehouse 3 can only receive shipments from plant 3 and 4. The total supply capacity of those plants is only 400 units. Therefore, warehouse 3 gets 50 less units from plants 3 and 4 than are demanded from retailers 5, 6, and 7.

Linear Program Formulation

1. Overall idea of problem

- Refrigerators moving from $n = 4$ plants to $q = 2$ warehouses to $m = 7$ retailers.
- Not all plants deliver to all warehouses.
- Not all warehouses deliver to all retailers.
- Costs of shipping from plants to warehouses vary by pair.
- Costs of shipping from warehouses to retailers vary by pair.
- Each plant has a capacity in terms of number of refrigerators it can supply.
- Each retailer has a capacity in terms of number of refrigerators it demands.
- Warehouse 2 has closed and all associate routes have been eliminated.

2. What is the goal? What are you trying to achieve?

- Unchanged from part A.

3. Identify variables

- Unchanged from part A.

4. Identify constraints

- All constraints from part A remain in effect with the addition of two new constraints:
- $np_{12} + np_{22} + np_{32} + np_{42} = 0$
- $nw_{23} + nw_{24} + nw_{25} + nw_{26} = 0$

5. Identify inputs and outputs that you can control

- Unchanged from part A.

6. Specify all quantities mathematically

- Unchanged from part A.

7. Check the model for completeness and correctness

- All variables are positive.

Matlab Code

Code minimally changed from part A. Only changes are 2 additional constraints (16 total equations) in the linear equality matrix and vector. Identical code from part A is not shown below (to save space).

```
1 % -----
2 % linear equality matrix and vector
3 %   note matlab arrays/vectors start at index 1 (not 0)
4 %   16 equations in 32 variables
5 % -----
6 Aeq = zeros(16, 32);
7 beq = zeros(16, 1);
8 %np11 + np21 + np31          = nw11 + nw12 + nw13 + nw14
9 %np11 + np21 + np31 - nw11 - nw12 - nw13 - nw14 = 0
10 Aeq(1, [1, 3, 5, 10, 11, 12, 13]) = [1, 1, 1, -1, -1, -1, -1];
11 %np12 + np22 + np32 + np42 =          nw23 + nw24 + nw25 + nw26
12 %np12 + np22 + np32 + np42 - nw23 - nw24 - nw25 - nw26 = 0
13 Aeq(2, [2, 4, 6, 8, 14, 15, 16, 17]) = [1, 1, 1, 1, -1, -1, -1, -1];
14 %          np33 + np43 =          nw34 + nw35 + nw36 + nw37
15 %np33 + np43 - nw34 - nw35 - nw36 - nw37 = 0
16 Aeq(3, [7, 9, 18, 19, 20, 21]) = [1, 1, -1, -1, -1, -1];
17 %s1 = np11 + np12
18 %s1 - np11 - np12 = 0
19 Aeq(4, [22, 1, 2]) = [1, -1, -1];
20 %s2 = np21 + np22
21 %s2 - np21 - np22 = 0
22 Aeq(5, [23, 3, 4]) = [1, -1, -1];
23 %s3 = np31 + np32 + np33
24 %s3 - np31 - np32 - np33 = 0
25 Aeq(6, [24, 5, 6, 7]) = [1, -1, -1, -1];
26 %s4 =          np42 + np43
27 %s4 - np42 - np43 = 0
```

```

28 Aeq(7,[25,8,9]) = [1,-1,-1];
29 %d1 = nw11
30 %d1 - nw11 = 0
31 Aeq(8,[26,10]) = [1,-1];
32 %d2 = nw12
33 %d2 - nw12 = 0
34 Aeq(9,[27,11]) = [1,-1];
35 %d3 = nw13 + nw23
36 %d3 - nw13 - nw23 = 0
37 Aeq(10,[28,12,14]) = [1,-1,-1];
38 %d4 = nw14 + nw24 + nw34
39 %d4 - nw14 - nw24 - nw34 = 0
40 Aeq(11,[29,13,15,18]) = [1,-1,-1,-1];
41 %d5 =          nw25 + nw35
42 %d5 - nw25 - nw35 = 0
43 Aeq(12,[30,16,19]) = [1,-1,-1];
44 %d6 =          nw26 + nw36
45 %d6 - nw26 - nw36 = 0
46 Aeq(13,[31,17,20]) = [1,-1,-1];
47 %d7 =          nw37
48 %d7 - nw37 = 0
49 Aeq(14,[32,21]) = [1,-1];
50 %np12 + np22 + np32 + np42 = 0
51 Aeq(15,[2,4,6,8]) = [1,1,1,1];
52 %nw23 + nw24 + nw25 + nw26 = 0
53 Aeq(16,[14,15,16,17]) = [1,1,1,1];

```

Problem 1 Part C

Determine the number of refrigerators to be shipped from plants to warehouses, and then warehouses to retailers to minimize the cost. For part C warehouse 2 is limited to just 100 units per week. Changes to the problem statement or solution relative to Problem 1 Part A are highlighted in blue.

Solution

- In all 1000 units will travel through the network at a minimum cost of \$18300.
- Ship 150 units from plant #1 to warehouse #1 at a cost of \$1500.
- Ship 350 units from plant #2 to warehouse #1 at a cost of \$3850.
- Ship 100 units from plant #2 to warehouse #2 at a cost of \$800.
- Ship 0 units from plant #3 to warehouse #2 at a cost of \$ 0.
- Ship 250 units from plant #3 to warehouse #3 at a cost of \$2250.
- Ship 150 units from plant #4 to warehouse #3 at a cost of \$1200.
- Ship 100 units from warehouse #1 to retailer #1 at a cost of \$ 500.
- Ship 150 units from warehouse #1 to retailer #2 at a cost of \$ 900.
- Ship 100 units from warehouse #1 to retailer #3 at a cost of \$ 700.
- Ship 150 units from warehouse #1 to retailer #4 at a cost of \$1500.
- Ship 50 units from warehouse #2 to retailer #4 at a cost of \$ 400.
- Ship 50 units from warehouse #2 to retailer #5 at a cost of \$ 500.
- Ship 150 units from warehouse #3 to retailer #5 at a cost of \$1800.
- Ship 150 units from warehouse #3 to retailer #6 at a cost of \$1800.
- Ship 100 units from warehouse #3 to retailer #7 at a cost of \$ 600.
- 150 total units will leave plant #1 (capacity is 150).
- 450 total units will leave plant #2 (capacity is 450).
- 250 total units will leave plant #3 (capacity is 250).
- 150 total units will leave plant #4 (capacity is 150).

- 500 total units will enter warehouse #1, 500 units will leave.
- 100 total units will enter warehouse #2, 100 units will leave.
- 400 total units will enter warehouse #3, 400 units will leave.
- 100 total units will enter retailer #1 (demand is 100).
- 150 total units will enter retailer #2 (demand is 150).
- 100 total units will enter retailer #3 (demand is 100).
- 200 total units will enter retailer #4 (demand is 200).
- 200 total units will enter retailer #5 (demand is 200).
- 150 total units will enter retailer #6 (demand is 150).
- 100 total units will enter retailer #7 (demand is 100).

Linear Program Formulation

1. Overall idea of problem

- Refrigerators moving from $n = 4$ plants to $q = 3$ warehouses to $m = 7$ retailers.
- Not all plants deliver to all warehouses.
- Not all warehouses deliver to all retailers.
- Costs of shipping from plants to warehouses vary by pair.
- Costs of shipping from warehouses to retailers vary by pair.
- Each plant has a capacity in terms of number of refrigerators it can supply.
- Each retailer has a capacity in terms of number of refrigerators it demands.
- Warehouse 2 is limited to just 100 units in and out per week.

2. What is the goal? What are you trying to achieve?

- Unchanged from part A.

3. Identify variables

- Unchanged from part A.

4. Identify constraints

- All constraints from part A remain in effect with the addition of two new constraints:
- $np_{12} + np_{22} + np_{32} + np_{42} = 100$
- $nw_{23} + nw_{24} + nw_{25} + nw_{26} = 100$

5. Identify inputs and outputs that you can control
 - Unchanged from part A.
6. Specify all quantities mathematically
 - Unchanged from part A.
7. Check the model for completeness and correctness
 - All variables are positive.

Matlab Code

Code minimally changed from part A. Only changes are 2 additional constraints (16 total equations) in the linear equality matrix and vector. Identical code from part A is not shown below (to save space).

```

1 % -----
2 % linear equality matrix and vector
3 %   note matlab arrays/vectors start at index 1 (not 0)
4 %   16 equations in 32 variables
5 % -----
6 Aeq = zeros(16, 32);
7 beq = zeros(16, 1);
8 %np11 + np21 + np31          = nw11 + nw12 + nw13 + nw14
9 %np11 + np21 + np31 - nw11 - nw12 - nw13 - nw14 = 0
10 Aeq(1, [1, 3, 5, 10, 11, 12, 13]) = [1, 1, 1, -1, -1, -1, -1];
11 %np12 + np22 + np32 + np42 =          nw23 + nw24 + nw25 + nw26
12 %np12 + np22 + np32 + np42 - nw23 - nw24 - nw25 - nw26 = 0
13 Aeq(2, [2, 4, 6, 8, 14, 15, 16, 17]) = [1, 1, 1, 1, -1, -1, -1, -1];
14 %          np33 + np43 =          nw34 + nw35 + nw36 + nw37
15 %np33 + np43 - nw34 - nw35 - nw36 - nw37 = 0
16 Aeq(3, [7, 9, 18, 19, 20, 21]) = [1, 1, -1, -1, -1, -1];
17 %s1 = np11 + np12
18 %s1 - np11 - np12 = 0
19 Aeq(4, [22, 1, 2]) = [1, -1, -1];
20 %s2 = np21 + np22
21 %s2 - np21 - np22 = 0
22 Aeq(5, [23, 3, 4]) = [1, -1, -1];
23 %s3 = np31 + np32 + np33
24 %s3 - np31 - np32 - np33 = 0
25 Aeq(6, [24, 5, 6, 7]) = [1, -1, -1, -1];
26 %s4 =          np42 + np43
27 %s4 - np42 - np43 = 0
28 Aeq(7, [25, 8, 9]) = [1, -1, -1];
29 %d1 = nw11
30 %d1 - nw11 = 0
31 Aeq(8, [26, 10]) = [1, -1];
32 %d2 = nw12
33 %d2 - nw12 = 0
34 Aeq(9, [27, 11]) = [1, -1];

```

```

35 %d3 = nw13 + nw23
36 %d3 - nw13 - nw23 = 0
37 Aeq(10,[28,12,14]) = [1,-1,-1];
38 %d4 = nw14 + nw24 + nw34
39 %d4 - nw14 - nw24 - nw34 = 0
40 Aeq(11,[29,13,15,18]) = [1,-1,-1,-1];
41 %d5 =          nw25 + nw35
42 %d5 - nw25 - nw35 = 0
43 Aeq(12,[30,16,19]) = [1,-1,-1];
44 %d6 =          nw26 + nw36
45 %d6 - nw26 - nw36 = 0
46 Aeq(13,[31,17,20]) = [1,-1,-1];
47 %d7 =          nw37
48 %d7 - nw37 = 0
49 Aeq(14,[32,21]) = [1,-1];
50 %np12 + np22 + np32 + np42 = 100
51 Aeq(15,[2,4,6,8]) = [1,1,1,1];
52 beq(15) = 100;
53 %nw23 + nw24 + nw25 + nw26 = 100
54 Aeq(16,[14,15,16,17]) = [1,1,1,1];
55 beq(16) = 100;

```

Problem 1 Part D

Formulate a generalized linear programming model for the transshipment problem. Give the objective function and constraints as mathematical formulas.

Note

I can't help but feel that I've already done this in Parts A through C but I'll repeat (and try to expand) here in a more general form.

Objective Function

Whereas...

n = number of plants ≥ 1

q = number of warehouses ≥ 1

m = number of retailers ≥ 1

np_{ij} = number of refrigerators shipped from plant p_i to warehouse $w_j \geq 0$

cp_{ij} = cost of moving a refrigerator between plant p_i and warehouse $w_j > 0$

nw_{jk} = number of refrigerators shipped from warehouse w_j to retailer $r_k \geq 0$

cw_{jk} = cost of moving a refrigerator between warehouse w_j and retailer $r_k > 0$

The objective function is to...

$$\text{minimize } cost = \sum_{i=1}^n \sum_{j=1}^q np_{ij} * cp_{ij} + \sum_{j=1}^q \sum_{k=1}^m nw_{jk} * cw_{jk}$$

Constraints

The objective function (designed to minimize $cost$) is subject to the following constraints...

$0 \leq \text{capacity of each plant} = s_i \leq \text{some maximum}$

$0 \leq \text{throughput capability of each warehouse} = q_j \leq \text{some maximum}$

$0 \leq \text{demand of each retailer} = m_k \leq \text{some maximum}$

and also...

$$\sum_{i=1}^n np_{ij} = \text{capacity of a given plant } p_i$$

$$\sum_{i=1}^n np_{ij} = \sum_{k=1}^m nw_{jk} = \text{throughput capability for a given warehouse } w_j$$

$$\sum_{k=1}^m nw_{jk} = \text{demand of a given retailer } r_k$$

Problem 2

Problem 2 Part A

Determine the combination of ingredients that minimizes calories but meets all nutritional requirements.

- i. Formulate the problem as a linear program with an objective function and all constraints.

We can formulate this problem as the following linear program:

$$\begin{aligned} \text{minimize} \quad & 21 \cdot w_{\text{tomato}} + 16 \cdot w_{\text{lettuce}} + 40 \cdot w_{\text{spinach}} + 41 \cdot w_{\text{carrot}} \\ & + 585 \cdot w_{\text{sunflower seed}} + 120 \cdot w_{\text{smoked tofu}} + 164 \cdot w_{\text{chickpea}} + 884 \cdot w_{\text{oil}} \\ \\ \text{subject to} \quad & -0.85 \cdot w_{\text{tomato}} - 1.62 \cdot w_{\text{lettuce}} - 2.86 \cdot w_{\text{spinach}} - 0.93 \cdot w_{\text{carrot}} \\ & - 23.4 \cdot w_{\text{sunflower seed}} - 16 \cdot w_{\text{smoked tofu}} - 9 \cdot w_{\text{chickpea}} - 0 \cdot w_{\text{oil}} \leq -15, \\ \\ & -0.33 \cdot w_{\text{tomato}} - 0.2 \cdot w_{\text{lettuce}} - 0.39 \cdot w_{\text{spinach}} - 0.24 \cdot w_{\text{carrot}} \\ & - 48.7 \cdot w_{\text{sunflower seed}} - 5 \cdot w_{\text{smoked tofu}} - 2.6 \cdot w_{\text{chickpea}} - 100 \cdot w_{\text{oil}} \leq -2, \\ \\ & 0.33 \cdot w_{\text{tomato}} + 0.2 \cdot w_{\text{lettuce}} + 0.39 \cdot w_{\text{spinach}} + 0.24 \cdot w_{\text{carrot}} \\ & + 48.7 \cdot w_{\text{sunflower seed}} + 5 \cdot w_{\text{smoked tofu}} + 2.6 \cdot w_{\text{chickpea}} + 100 \cdot w_{\text{oil}} \leq 8, \\ \\ & -4.64 \cdot w_{\text{tomato}} - 2.37 \cdot w_{\text{lettuce}} - 3.63 \cdot w_{\text{spinach}} - 9.58 \cdot w_{\text{carrot}} \\ & - 15 \cdot w_{\text{sunflower seed}} - 3 \cdot w_{\text{smoked tofu}} - 27 \cdot w_{\text{chickpea}} - 0 \cdot w_{\text{oil}} \leq -4, \\ \\ & 9 \cdot w_{\text{tomato}} + 28 \cdot w_{\text{lettuce}} + 65 \cdot w_{\text{spinach}} + 69 \cdot w_{\text{carrot}} \\ & + 3.8 \cdot w_{\text{sunflower seed}} + 120 \cdot w_{\text{smoked tofu}} + 78 \cdot w_{\text{chickpea}} + 0 \cdot w_{\text{oil}} \leq 200, \\ \\ & 0.4 \cdot w_{\text{tomato}} - 0.6 \cdot w_{\text{lettuce}} - 0.6 \cdot w_{\text{spinach}} + 0.4 \cdot w_{\text{carrot}} \\ & + 0.4 \cdot w_{\text{sunflower seed}} + 0.4 \cdot w_{\text{smoked tofu}} + 0.4 \cdot w_{\text{chickpea}} + 0.4 \cdot w_{\text{oil}} \leq 0, \\ \\ & -w_{\text{tomato}} \leq 0 \\ & -w_{\text{lettuce}} \leq 0 \\ & -w_{\text{spinach}} \leq 0 \\ & -w_{\text{carrot}} \leq 0 \\ & -w_{\text{sunflower seed}} \leq 0 \\ & -w_{\text{smoked tofu}} \leq 0 \\ & -w_{\text{chickpea}} \leq 0 \\ & -w_{\text{oil}} \leq 0 \end{aligned}$$

where the w parameters are the weight of each ingredient, in 100's of grams. The equations, from top to bottom, are:

- 1) Minimize total calories
 - 2) Subject to total protein ≥ 15 grams
 - 3) Subject to total fat ≥ 2 grams
 - 4) Subject to total fat ≤ 8 grams
 - 5) Subject to total carbohydrates ≥ 4 grams
 - 6) Subject to total sodium ≤ 200 milligrams
 - 7) Subject to total leafy green mass $\geq 40\%$ of total mass
 - 8) Subject to individual ingredient weights ≥ 0
- ii. Determine the optimal solution for the linear program using any software you want. Include a copy of the code/file in the report.

The following MATLAB code was used to generate the solution:

```
1 f = [21, 16, 40, 41, 585, 120, 164, 884];
2
3 A = [-0.85, -1.62, -2.86, -0.93, -23.4, -16, -9, 0;
4      -0.33, -0.2, -0.39, -0.24, -48.7, -5, -2.6, -100;
5      0.33, 0.2, 0.39, 0.24, 48.7, 5, 2.6, 100;
6      -4.64, -2.37, -3.63, -9.58, -15, -3, -27, 0;
7      9, 28, 65, 69, 3.8, 120, 78, 0;
8      0.4, -0.6, -0.6, 0.4, 0.4, 0.4, 0.4, 0.4;
9      -eye(8)];
10
11 b = [-15; -2; 8; -4; 200; 0; zeros(8,1)];
12
13 [X,FVAL,EXITFLAG] = linprog(f,A,b);
```

where **X** stores the resulting weights of ingredients (in 100's of grams), **FVAL** stores the minimized number of calories, and **EXITFLAG** stores the status of the `linprog` optimization.

- iii. What is the cost of the low calorie salad?

The optimal low calorie salad contains the following weights of ingredients (in 100's of grams):

$$\begin{aligned}
w_{\text{tomato}} &\approx 0 \\
w_{\text{lettuce}} &\approx 0.5855 \\
w_{\text{spinach}} &\approx 0 \\
w_{\text{carrot}} &\approx 0 \\
w_{\text{sunflower seed}} &\approx 0 \\
w_{\text{smoked tofu}} &\approx 0.8782 \\
w_{\text{chickpea}} &\approx 0 \\
w_{\text{oil}} &\approx 0
\end{aligned}$$

The optimal low calorie salad costs approximately \$2.33. It contains approximately 114.75 kcal

Problem 2 Part B

Veronica realizes that it is also important to minimize the cost associated with the new salad. Unfortunately some of the ingredients can be expensive. Determine the combination of ingredients that minimizes cost.

- i. Formulate the problem as a linear program with an objective function and all constraints.

We can formulate this problem as the following linear program:

$$\begin{aligned}
 &\text{minimize} && 1 \cdot w_{\text{tomato}} + 0.75 \cdot w_{\text{lettuce}} + 0.5 \cdot w_{\text{spinach}} + 0.5 \cdot w_{\text{carrot}} \\
 &&& + 0.45 \cdot w_{\text{sunflower seed}} + 2.15 \cdot w_{\text{smoked tofu}} + 0.95 \cdot w_{\text{chickpea}} + 2.00 \cdot w_{\text{oil}} \\
 \\
 &\text{subject to} && -0.85 \cdot w_{\text{tomato}} - 1.62 \cdot w_{\text{lettuce}} - 2.86 \cdot w_{\text{spinach}} - 0.93 \cdot w_{\text{carrot}} \\
 &&& - 23.4 \cdot w_{\text{sunflower seed}} - 16 \cdot w_{\text{smoked tofu}} - 9 \cdot w_{\text{chickpea}} - 0 \cdot w_{\text{oil}} \leq -15, \\
 \\
 &&& -0.33 \cdot w_{\text{tomato}} - 0.2 \cdot w_{\text{lettuce}} - 0.39 \cdot w_{\text{spinach}} - 0.24 \cdot w_{\text{carrot}} \\
 &&& - 48.7 \cdot w_{\text{sunflower seed}} - 5 \cdot w_{\text{smoked tofu}} - 2.6 \cdot w_{\text{chickpea}} - 100 \cdot w_{\text{oil}} \leq -2, \\
 \\
 &&& 0.33 \cdot w_{\text{tomato}} + 0.2 \cdot w_{\text{lettuce}} + 0.39 \cdot w_{\text{spinach}} + 0.24 \cdot w_{\text{carrot}} \\
 &&& + 48.7 \cdot w_{\text{sunflower seed}} + 5 \cdot w_{\text{smoked tofu}} + 2.6 \cdot w_{\text{chickpea}} + 100 \cdot w_{\text{oil}} \leq 8, \\
 \\
 &&& -4.64 \cdot w_{\text{tomato}} - 2.37 \cdot w_{\text{lettuce}} - 3.63 \cdot w_{\text{spinach}} - 9.58 \cdot w_{\text{carrot}} \\
 &&& - 15 \cdot w_{\text{sunflower seed}} - 3 \cdot w_{\text{smoked tofu}} - 27 \cdot w_{\text{chickpea}} - 0 \cdot w_{\text{oil}} \leq -4, \\
 \\
 &&& 9 \cdot w_{\text{tomato}} + 28 \cdot w_{\text{lettuce}} + 65 \cdot w_{\text{spinach}} + 69 \cdot w_{\text{carrot}} \\
 &&& + 3.8 \cdot w_{\text{sunflower seed}} + 120 \cdot w_{\text{smoked tofu}} + 78 \cdot w_{\text{chickpea}} + 0 \cdot w_{\text{oil}} \leq 200, \\
 \\
 &&& 0.4 \cdot w_{\text{tomato}} - 0.6 \cdot w_{\text{lettuce}} - 0.6 \cdot w_{\text{spinach}} + 0.4 \cdot w_{\text{carrot}} \\
 &&& + 0.4 \cdot w_{\text{sunflower seed}} + 0.4 \cdot w_{\text{smoked tofu}} + 0.4 \cdot w_{\text{chickpea}} + 0.4 \cdot w_{\text{oil}} \leq 0, \\
 \\
 &&& -w_{\text{tomato}} \leq 0 \\
 &&& -w_{\text{lettuce}} \leq 0 \\
 &&& -w_{\text{spinach}} \leq 0 \\
 &&& -w_{\text{carrot}} \leq 0 \\
 &&& -w_{\text{sunflower seed}} \leq 0 \\
 &&& -w_{\text{smoked tofu}} \leq 0 \\
 &&& -w_{\text{chickpea}} \leq 0 \\
 &&& -w_{\text{oil}} \leq 0
 \end{aligned}$$

where the w parameters are the weight of each ingredient, in 100's of grams. The equations, from top to bottom, are:

- 1) Minimize total cost
 - 2) Subject to total protein ≥ 15 grams
 - 3) Subject to total fat ≥ 2 grams
 - 4) Subject to total fat ≤ 8 grams
 - 5) Subject to total carbohydrates ≥ 4 grams
 - 6) Subject to total sodium ≤ 200 milligrams
 - 7) Subject to total leafy green mass $\geq 40\%$ of total mass
 - 8) Subject to individual ingredient weights ≥ 0
- ii. Determine the optimal solution for the linear program using any software you want. Include a copy of the code/file in the report.

The following MATLAB code was used to generate the solution:

```

1 f = [1, 0.75, 0.5, 0.5, 0.45, 2.15, 0.95, 2.0];
2
3 A = [-0.85, -1.62, -2.86, -0.93, -23.4, -16, -9, 0;
4      -0.33, -0.2, -0.39, -0.24, -48.7, -5, -2.6, -100;
5      0.33, 0.2, 0.39, 0.24, 48.7, 5, 2.6, 100;
6      -4.64, -2.37, -3.63, -9.58, -15, -3, -27, 0;
7      9, 28, 65, 69, 3.8, 120, 78, 0;
8      0.4, -0.6, -0.6, 0.4, 0.4, 0.4, 0.4, 0.4;
9      -eye(8)];
10
11 b = [-15; -2; 8; -4; 200; 0; zeros(8,1)];
12
13 [X,FVAL,EXITFLAG] = linprog(f,A,b);

```

where **X** stores the resulting weights of ingredients (in 100's of grams), **FVAL** stores the minimized number of calories, and **EXITFLAG** stores the status of the `linprog` optimization.

- iii. How many calories are in the low cost salad?

The optimal low cost salad contains the following weights of ingredients (in 100's of grams):

$$\begin{aligned}
w_{\text{tomato}} &\approx 0 \\
w_{\text{lettuce}} &\approx 0 \\
w_{\text{spinach}} &\approx 0.8323 \\
w_{\text{carrot}} &\approx 0 \\
w_{\text{sunflower seed}} &\approx 0.0961 \\
w_{\text{smoked tofu}} &\approx 0 \\
w_{\text{chickpea}} &\approx 1.1524 \\
w_{\text{oil}} &\approx 0
\end{aligned}$$

The optimal low cost salad costs approximately \$1.55. It contains approximately 278.49 kcal

Problem 2 Part C

Compare the results from part A and B. Veronicas goal is to create a Very Veggie Salad that is both low calorie and low cost. She would like to sell the salad for \$5.00 and still have a profit of at least \$3.00. However if she can advertise that the salad has under 250 calories then she may be able to sell more.

- i. Suggest some possible ways that she select a combination of ingredients that is near optimal for both objectives. This is a type of multi-objective optimization.

What Veronica should do is attempt to find a “pareto optimal” solution that accounts for both calories and cost. In order to do this, she should introduce a new parameter λ to her optimization formulation. In essence, she would like to:

$$\begin{array}{ll}\text{minimize} & (1 - \lambda) \cdot \text{CALORIES} + \lambda \cdot \text{PRICE} \\ \text{subject to} & \text{CONSTRAINTS}\end{array}$$

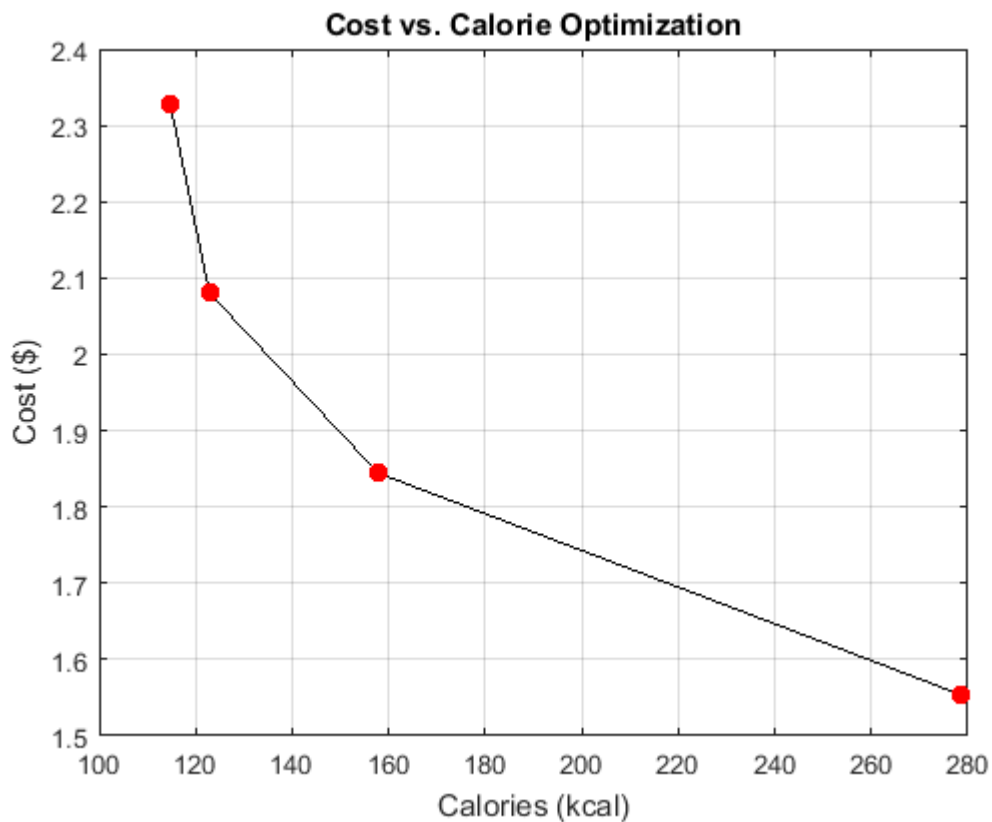
for values of λ between 0 and 1. When $\lambda = 0$, she will be finding the minimum calorie combination of ingredients. When $\lambda = 1$, she will be finding the minimum price combination of ingredients. For all values of λ inbetween, she will be finding the “pareto optimal” combination of ingredients. So, she should solve the optimization problem for values of λ between 0 and 1. She can then examine these pareto optimal combinations of ingredients to determine which would best meet her goals of \$3.00 profit and under 250 calories.

We can express this as the following linear program:

$$\begin{aligned}
\text{minimize} \quad & (1 - \lambda) \cdot (21 \cdot w_{\text{tomato}} + 16 \cdot w_{\text{lettuce}} + 40 \cdot w_{\text{spinach}} + 41 \cdot w_{\text{carrot}} \\
& + 585 \cdot w_{\text{sunflower seed}} + 120 \cdot w_{\text{smoked tofu}} + 164 \cdot w_{\text{chickpea}} + 884 \cdot w_{\text{oil}}) \\
& + \lambda \cdot (1 \cdot w_{\text{tomato}} + 0.75 \cdot w_{\text{lettuce}} + 0.5 \cdot w_{\text{spinach}} + 0.5 \cdot w_{\text{carrot}} \\
& + 0.45 \cdot w_{\text{sunflower seed}} + 2.15 \cdot w_{\text{smoked tofu}} + 0.95 \cdot w_{\text{chickpea}} + 2.00 \cdot w_{\text{oil}}) \\
\\
\text{subject to} \quad & -0.85 \cdot w_{\text{tomato}} - 1.62 \cdot w_{\text{lettuce}} - 2.86 \cdot w_{\text{spinach}} - 0.93 \cdot w_{\text{carrot}} \\
& - 23.4 \cdot w_{\text{sunflower seed}} - 16 \cdot w_{\text{smoked tofu}} - 9 \cdot w_{\text{chickpea}} - 0 \cdot w_{\text{oil}} \leq -15, \\
\\
& -0.33 \cdot w_{\text{tomato}} - 0.2 \cdot w_{\text{lettuce}} - 0.39 \cdot w_{\text{spinach}} - 0.24 \cdot w_{\text{carrot}} \\
& - 48.7 \cdot w_{\text{sunflower seed}} - 5 \cdot w_{\text{smoked tofu}} - 2.6 \cdot w_{\text{chickpea}} - 100 \cdot w_{\text{oil}} \leq -2, \\
\\
& 0.33 \cdot w_{\text{tomato}} + 0.2 \cdot w_{\text{lettuce}} + 0.39 \cdot w_{\text{spinach}} + 0.24 \cdot w_{\text{carrot}} \\
& + 48.7 \cdot w_{\text{sunflower seed}} + 5 \cdot w_{\text{smoked tofu}} + 2.6 \cdot w_{\text{chickpea}} + 100 \cdot w_{\text{oil}} \leq 8, \\
\\
& -4.64 \cdot w_{\text{tomato}} - 2.37 \cdot w_{\text{lettuce}} - 3.63 \cdot w_{\text{spinach}} - 9.58 \cdot w_{\text{carrot}} \\
& - 15 \cdot w_{\text{sunflower seed}} - 3 \cdot w_{\text{smoked tofu}} - 27 \cdot w_{\text{chickpea}} - 0 \cdot w_{\text{oil}} \leq -4, \\
\\
& 9 \cdot w_{\text{tomato}} + 28 \cdot w_{\text{lettuce}} + 65 \cdot w_{\text{spinach}} + 69 \cdot w_{\text{carrot}} \\
& + 3.8 \cdot w_{\text{sunflower seed}} + 120 \cdot w_{\text{smoked tofu}} + 78 \cdot w_{\text{chickpea}} + 0 \cdot w_{\text{oil}} \leq 200, \\
\\
& 0.4 \cdot w_{\text{tomato}} - 0.6 \cdot w_{\text{lettuce}} - 0.6 \cdot w_{\text{spinach}} + 0.4 \cdot w_{\text{carrot}} \\
& + 0.4 \cdot w_{\text{sunflower seed}} + 0.4 \cdot w_{\text{smoked tofu}} + 0.4 \cdot w_{\text{chickpea}} + 0.4 \cdot w_{\text{oil}} \leq 0, \\
\\
& -w_{\text{tomato}} \leq 0 \\
& -w_{\text{lettuce}} \leq 0 \\
& -w_{\text{spinach}} \leq 0 \\
& -w_{\text{carrot}} \leq 0 \\
& -w_{\text{sunflower seed}} \leq 0 \\
& -w_{\text{smoked tofu}} \leq 0 \\
& -w_{\text{chickpea}} \leq 0 \\
& -w_{\text{oil}} \leq 0
\end{aligned}$$

- ii. What combination of ingredient would you suggest and what is the associated cost and calorie.

Varying λ between 0 and 1, and solving the resulting optimization problems, we obtain the following pareto optimal combinations of total calories and total price:



Based on this result, I would suggest the following combination of ingredients:

$$\begin{aligned}
 w_{\text{tomato}} &\approx 0 \\
 w_{\text{lettuce}} &\approx 0 \\
 w_{\text{spinach}} &\approx 0.5346 \\
 w_{\text{carrot}} &\approx 0 \\
 w_{\text{sunflower seed}} &\approx 0.0865 \\
 w_{\text{smoked tofu}} &\approx 0.7154 \\
 w_{\text{chickpea}} &\approx 0 \\
 w_{\text{oil}} &\approx 0
 \end{aligned}$$

This pareto optimal salad costs approximately \$1.84. It contains approximately 157.86 kcal. This meets both of Veronicas goals.

iii. Note: There is not one right answer. Discuss how you derived your solution.

The following MATLAB code was used to derive the answer, based on the discussion above.

```

1 lambda = 0:0.0005:1;
2

```

```

3  for j = 1:numel(lambda)
4      L = lambda(j);
5      f1 = [21, 16, 40, 41, 585, 120, 164, 884];
6      f2 = [1, 0.75, 0.5, 0.5, 0.45, 2.15, 0.95, 2.0];
7
8      f = (1 - L)*f1 + L*f2;
9
10     A = [-0.85, -1.62, -2.86, -0.93, -23.4, -16, -9, 0;
11          -0.33, -0.2, -0.39, -0.24, -48.7, -5, -2.6, -100;
12          0.33, 0.2, 0.39, 0.24, 48.7, 5, 2.6, 100;
13          -4.64, -2.37, -3.63, -9.58, -15, -3, -27, 0;
14          9, 28, 65, 69, 3.8, 120, 78, 0;
15          0.4, -0.6, -0.6, 0.4, 0.4, 0.4, 0.4, 0.4;
16          -eye(8)];
17
18     b = [-15; -2; 8; -4; 200; 0; zeros(8,1)];
19
20     [X,FVAL,EXITFLAG] = linprog(f,A,b);
21
22     calories_per_ingredient = [21, 16, 40, 41, 585, 120, 164, 884];
23     cost_per_ingredient = [1, 0.75, 0.5, 0.5, 0.45, 2.15, 0.95, 2];
24
25     if (EXITFLAG == 1)
26         cost(j) = cost_per_ingredient * X;
27         calories(j) = calories_per_ingredient * X;
28     else
29         cost(j) = NaN;
30         calories(j) = NaN;
31     end
32 end
33
34 figure(1)
35 plot(calories, cost, 'k-')
36 hold on
37 plot(calories, cost, 'r.', 'markersize', 25)
38 grid on
39 xlabel('Calories (kcal)')
40 ylabel('Cost ($)')
41 title('Cost vs. Calorie Optimization')

```

Problem 3

a) What are the lengths of the shortest paths from vertex a to all other vertices?

The shortest path problem can be solved using a linear programming model of the following form:

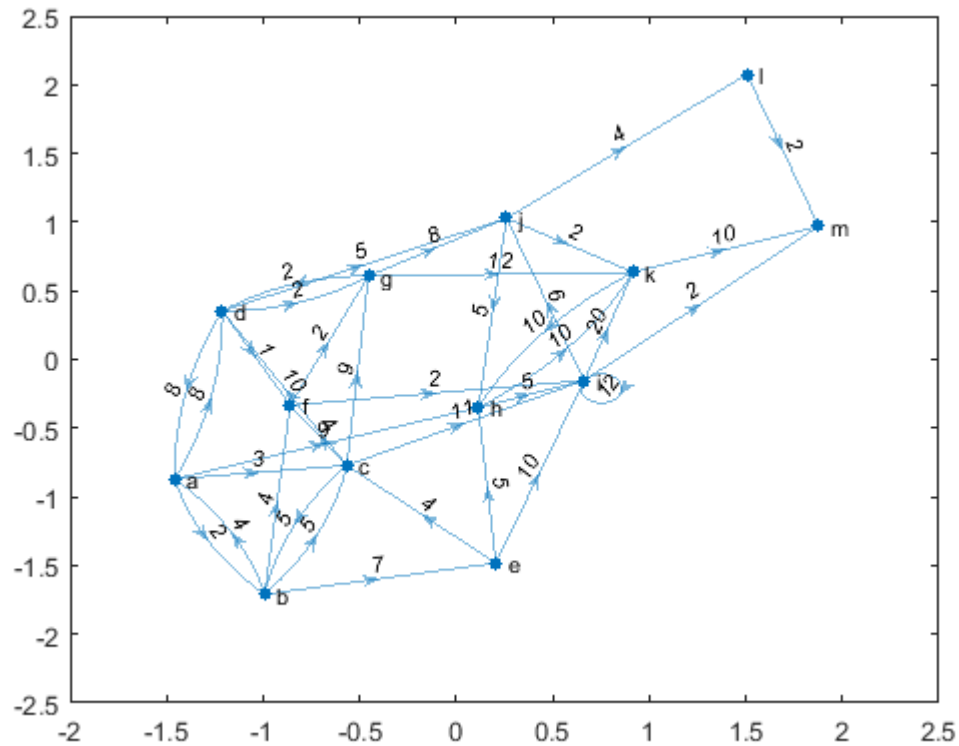
maximize d_t subject to $d_s = 0$ $d_v - d_u \leq l_{u \rightarrow v}$ for every edge $u \rightarrow v$

Based on the paths and weights denoted from the Project3Problem3-1.txt file the expressions for the linear programming solution to find shortest paths will be:

- $d_b - d_a \leq 2$
- $d_c - d_a \leq 3$
- $d_d - d_a \leq 8$
- $d_h - d_a \leq 9$
- $d_a - d_b \leq 4$
- $d_c - d_b \leq 5$
- $d_e - d_b \leq 7$
- $d_f - d_b \leq 4$
- $d_d - d_c \leq 10$
- $d_b - d_c \leq 5$
- $d_g - d_c \leq 9$
- $d_i - d_c \leq 11$
- $d_f - d_c \leq 4$
- $d_a - d_d \leq 8$
- $d_g - d_d \leq 2$
- $d_j - d_d \leq 5$
- $d_f - d_d \leq 1$
- $d_h - d_e \leq 5$
- $d_c - d_e \leq 4$
- $d_i - d_e \leq 10$
- $d_i - d_f \leq 2$
- $d_g - d_f \leq 2$
- $d_d - d_g \leq 2$
- $d_j - d_g \leq 8$
- $d_k - d_g \leq 12$
- $d_i - d_h \leq 5$
- $d_k - d_h \leq 10$

- $d_a - d_i \leq 20$
- $d_k - d_i \leq 6$
- $d_j - d_i \leq 2$
- $d_m - d_i \leq 12$
- $d_i - d_j \leq 2$
- $d_k - d_j \leq 4$
- $d_l - d_j \leq 5$
- $d_h - d_k \leq 10$
- $d_m - d_k \leq 10$
- $d_m - d_l \leq 2$

An image of the graph is below:



Code used to find the shortest paths:

```

1 clear variables
2 close all
3 clc
4
5 %% PARSE INPUT DATA
6 % Open file

```

```

7 fid = fopen('p3.input.txt');
8
9 % Read line while there is data remaining
10 tline = fgets(fid);
11 rowIdx = 0;
12 while ischar(tline)
13     % Increment row counter to store data
14     rowIdx = rowIdx + 1;
15
16     % fgetl reads line in as char array -- split on spaces
17     C = strsplit(tline, ' ');
18
19     % Convert node letter to index using ASCII codes (a = 1, b = 2, ...)
20     edgeStart(rowIdx) = double(C{1}) - double('a') + 1;
21     edgeEnd(rowIdx) = double(C{2}) - double('a') + 1;
22     edgeWeight(rowIdx) = str2num(C{3});
23     % Grab the next line
24     tline = fgetl(fid);
25 end
26
27 fclose(fid);
28
29 %% SHORTEST PATH FROM A TO OTHER VERTICES
30
31 % Number of nodes is highest numbered node in our data
32 numberOfNodes = max([edgeStart, edgeEnd]);
33
34 % Build A and b matrices from edgeEnd and edgeStart
35 % Size of A is number of inequalities by number of nodes
36 A = zeros(numel(edgeWeight), numberOfNodes);
37
38 for j = 1:numel(edgeWeight)
39     A(j, edgeStart(j)) = -1;
40     A(j, edgeEnd(j)) = 1;
41 end
42
43 b = edgeWeight';
44
45 % Additional constraint that all distances must be greater than 0
46 A = [A; -eye(numberOfNodes)];
47
48 b = [b; zeros(numberOfNodes, 1)];
49
50 % Single equality constraint -- distance to a must be zero, since we start at a
51 Aeq = zeros(1, numberOfNodes);
52 Aeq(1, 1) = 1;
53 beq = 0;
54
55 % Minimization constraint is to maximize negative sum of distances
56 f = -ones(numberOfNodes, 1);
57
58 [x, fval, exitflag] = linprog(f,A,b,Aeq,beq);
59
60 fid = fopen('P3A.solution.txt','w');

```



```

61
62 fprintf(fid, '----- P3.A SOLUTION -----\n');
63 for j = 1:numel(x)
64     fprintf(fid, 'Distance from a to %c = %2.0f \n', char('a'+ j - 1), x(j));
65 end
66
67 fclose(fid);

```

The lengths of the shortest paths from a to the other vertices in the graph are listed below:

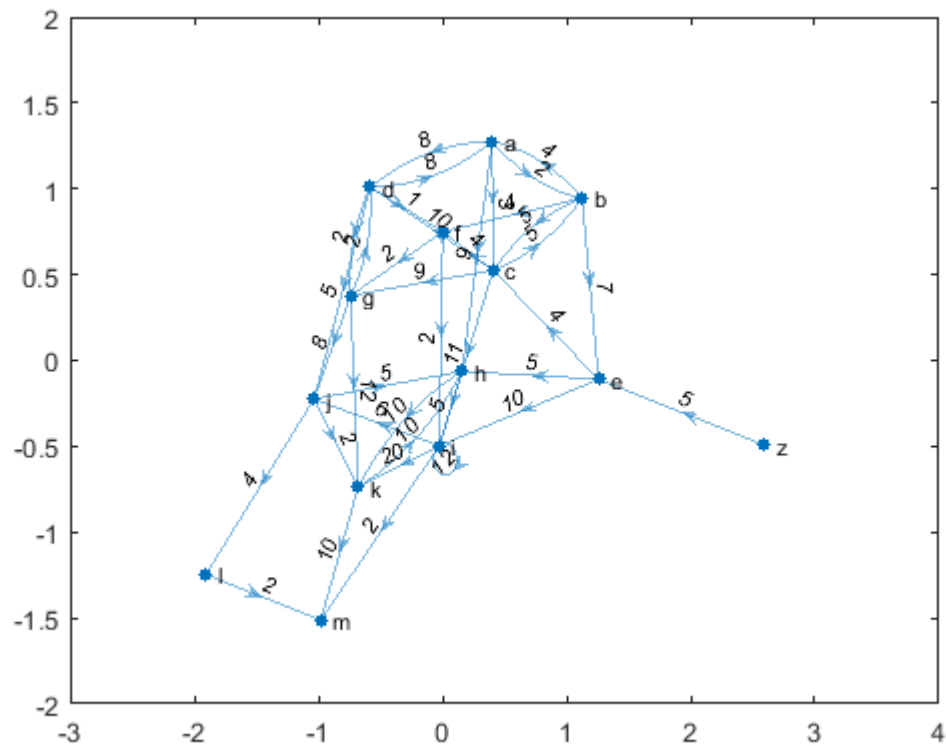
```

1 ----- P3.A SOLUTION -----
2 Distance from a to a = -0
3 Distance from a to b = 2
4 Distance from a to c = 3
5 Distance from a to d = 8
6 Distance from a to e = 9
7 Distance from a to f = 6
8 Distance from a to g = 8
9 Distance from a to h = 9
10 Distance from a to i = 8
11 Distance from a to j = 10
12 Distance from a to k = 14
13 Distance from a to l = 15
14 Distance from a to m = 17

```

- b) If a vertex z is added to the graph for which there is no path from vertex a to vertex z , what will be the result when you attempt to find the lengths of shortest paths as in part a).

Graphic with the new vertex z added:



```

1 clear variables
2 close all
3 clc
4
5 %% PARSE INPUT DATA
6 % Open file
7 fid = fopen('p3.input.unreachable.z.txt');
8
9 % Read line while there is data remaining
10 tline = fgets(fid);
11 rowIdx = 0;
12 while ischar(tline)
13     % Increment row counter to store data
14     rowIdx = rowIdx + 1;
15
16     % fgetl reads line in as char array -- split on spaces
17     C = strsplit(tline, ' ');
18
19     % Convert node letter to index using ASCII codes (a = 1, b = 2, ...)
20     edgeStart(rowIdx) = double(C{1}) - double('a') + 1;
21     edgeEnd(rowIdx) = double(C{2}) - double('a') + 1;
22     edgeWeight(rowIdx) = str2num(C{3});
23     % Grab the next line
24     tline = fgetl(fid);
25 end
26

```

```

27 fclose(fid);
28
29 %% SHORTEST PATH FROM A TO OTHER VERTICES
30
31 % Number of nodes is highest numbered node in our data
32 numberOfNodes = max([edgeStart, edgeEnd]);
33
34 % Build A and b matrices from edgeEnd and edgeStart
35 % Size of A is number of inequalities by number of nodes
36 A = zeros(numel(edgeWeight), numberOfNodes);
37
38 for j = 1:numel(edgeWeight)
39     A(j, edgeStart(j)) = -1;
40     A(j, edgeEnd(j)) = 1;
41 end
42
43 b = edgeWeight';
44
45 % Additional constraint that all distances must be greater than 0
46 A = [A; -eye(numberOfNodes)];
47
48 b = [b; zeros(numberOfNodes, 1)];
49
50 % Single equality constraint -- distance to a must be zero, since we start at a
51 Aeq = zeros(1, numberOfNodes);
52 Aeq(1, 1) = 1;
53 beq = 0;
54
55 % Minimization constraint is to maximize negative sum of distances
56 f = -ones(numberOfNodes, 1);
57
58 [x, fval, exitflag] = linprog(f,A,b,Aeq,beq);

```

When using the script above with the new vertex z the matlab code fails with this message:
 "Exiting: One or more of the residuals, duality gap, or total relative error has stalled:
 the dual appears to be infeasible and the primal unbounded since the primal objective is
 $-1e+10$ and the dual objective is $1e+6$."

This error is consistent with what is expected because we are trying to find the shortest path from a to all other vertices however that is not feasible with the new vertex z because we made it inaccessible.

However a workaround for this result would be to look at the preliminary optimization results to see which nodes are causing problems. We can specify those problem nodes to fixed dummy distances that are very large relative to the longest feasible path(s). We then rerun the script to obtain the shortest paths for the remaining reachable nodes.

- c) What are the lengths of the shortest paths from each vertex to vertex m ? How can you solve this problem with just one linear program?

For this part we just swapped the directionality on all the edges from part A and set the origin to be vertex m . This yields the distance from m to every other vertex at once which the distances would be equivalent regardless of direction so it works out.

Code used to find the shortest paths from all other vertices using a single linear program:

```
1 clear variables
2 close all
3 clc
4
5 %% PARSE INPUT DATA
6 % Open file
7 fid = fopen('p3.input.txt');
8
9 % Read line while there is data remaining
10 tline = fgets(fid);
11 rowIdx = 0;
12 while ischar(tline)
13     % Increment row counter to store data
14     rowIdx = rowIdx + 1;
15
16     % fgetl reads line in as char array -- split on spaces
17     C = strsplit(tline, ' ');
18
19     % Convert node letter to index using ASCII codes (a = 1, b = 2, ...)
20     edgeStart(rowIdx) = double(C{1}) - double('a') + 1;
21     edgeEnd(rowIdx) = double(C{2}) - double('a') + 1;
22     edgeWeight(rowIdx) = str2num(C{3});
23     % Grab the next line
24     tline = fgetl(fid);
25 end
26
27 fclose(fid);
28
29 %% SHORTEST PATH FROM OTHER VERTICES TO M
30 % Going from other vertices to M is equivalent to reversing the edge
31 % directions and going from M to other vertcies.
32
33 % Number of nodes is highest numbered node in our data
34 numberOfNodes = max([edgeStart, edgeEnd]);
35
36 % Build A and b matrices from edgeEnd and edgeStart
37 % Size of A is number of inequalities by number of nodes
38 A = zeros(numel(edgeWeight), numberOfNodes);
39
40 % Since we're reversing the direction of the edges, we'll swap the +1 and
41 % -1 values to capture that.
42 for j = 1:numel(edgeWeight)
43     A(j, edgeStart(j)) = 1;
44     A(j, edgeEnd(j)) = -1;
45 end
46
47 b = edgeWeight';
48
49 % Additional constriant that all distances must be greater than 0
50 A = [A; -eye(numberOfNodes)];
51
```

```

52 b = [b; zeros(numberOfNodes, 1)];
53
54 % Single equality constraint -- distance to m must be zero, since we start
55 % at m
56 Aeq = zeros(1, numberOfNodes);
57 startNode = double('m') - double('a') + 1;
58 Aeq(1, startNode) = 1;
59 beq = 0;
60
61 % Minimization constraint is to maximize negative sum of distances
62 f = -ones(numberOfNodes, 1);
63
64 [x, fval, exitflag] = linprog(f,A,b,Aeq,beq);
65
66 fid = fopen('P3C.solution.txt', 'w');
67
68 fprintf(fid, '----- P3.C SOLUTION ----- \n');
69 for j = 1:numel(x)
70     fprintf(fid, 'Distance from %c to m = %2.0f \n', char('a'+ j - 1), x(j));
71 end
72
73 fclose(fid);

```

Output of the paths after running the above script:

```

1 ----- P3.C SOLUTION -----
2 Distance from a to m = 17
3 Distance from b to m = 15
4 Distance from c to m = 15
5 Distance from d to m = 12
6 Distance from e to m = 19
7 Distance from f to m = 11
8 Distance from g to m = 14
9 Distance from h to m = 14
10 Distance from i to m = 9
11 Distance from j to m = 7
12 Distance from k to m = 10
13 Distance from l to m = 2
14 Distance from m to m = 0

```

- d) Suppose that all paths must pass through vertex i . How can you calculate the length of the shortest path from any vertex x to vertex y that pass through vertex i (for all $x, y \in V$)? Calculate the lengths of these paths for the given graph. (Note: for some vertices x and y , it may be impossible to pass through vertex i).

Code used to find the paths from one vertex to another while passing through an intermediate vertex i :

```

1 clear variables
2 close all
3 clc

```

```

4
5 %% PARSE INPUT DATA
6 % Open file
7 fid = fopen('p3.input.txt');
8
9 % Read line while there is data remaining
10 tline = fgets(fid);
11 rowIdx = 0;
12 while ischar(tline)
13     % Increment row counter to store data
14     rowIdx = rowIdx + 1;
15
16     % fgetl reads line in as char array -- split on spaces
17     C = strsplit(tline, ' ');
18
19     % Convert node letter to index using ASCII codes (a = 1, b = 2, ...)
20     edgeStart(rowIdx) = double(C{1}) - double('a') + 1;
21     edgeEnd(rowIdx) = double(C{2}) - double('a') + 1;
22     edgeWeight(rowIdx) = str2num(C{3});
23     % Grab the next line
24     tline = fgetl(fid);
25 end
26
27 fclose(fid);
28
29 %% SHORTEST PATH FROM OTHER VERTICES TO I
30 % Going from other vertices to I is equivalent to reversing the edge
31 % directions and going from I to other vertcies.
32
33 % Number of nodes is highest numbered node in our data
34 numberOfNodes = max([edgeStart, edgeEnd]);
35
36 % Build A and b matrices from edgeEnd and edgeStart
37 % Size of A is number of inequalities by number of nodes
38 A = zeros(numel(edgeWeight), numberOfNodes);
39
40 % Since we're reversing the direction of the edges, we'll swap the +1 and
41 % -1 values to capture that.
42 for j = 1:numel(edgeWeight)
43     A(j, edgeStart(j)) = 1;
44     A(j, edgeEnd(j)) = -1;
45 end
46
47 b = edgeWeight';
48
49 % Additional constraint that all distances must be greater than 0
50 A = [A; -eye(numberOfNodes)];
51
52 b = [b; zeros(numberOfNodes, 1)];
53
54 % Single equality constraint -- distance to i must be zero, since we start
55 % at i
56 Aeq = zeros(3, numberOfNodes);
57 startNode = double('i') - double('a') + 1;

```

```

58 nodeL = double('l') - double('a') + 1;
59 nodeM = double('m') - double('a') + 1;
60 Aeq(1, startNode) = 1;
61 Aeq(2, nodeL) = 1;
62 Aeq(3, nodeM) = 1;
63 beq = [0; 99999; 99999];
64
65 % Minimization constraint is to maximize negative sum of distances
66 f = -ones(numberOfNodes, 1);
67
68 [x, fval, exitflag] = linprog(f,A,b,Aeq,beq);
69
70 fprintf('----- P3.C SOLUTION -----\n')
71 for j = 1:numel(x)
72     fprintf('Distance from %c to i = %2.0f \n', char('a'+ j - 1), x(j))
73 end
74
75 distanceToNodeI = x;
76
77 %% SHORTEST PATH FROM I TO OTHER VERTICES
78
79 % Number of nodes is highest numbered node in our data
80 numberOfNodes = max([edgeStart, edgeEnd]);
81
82 % Build A and b matrices from edgeEnd and edgeStart
83 % Size of A is number of inequalities by number of nodes
84 A = zeros(numel(edgeWeight), numberOfNodes);
85
86 for j = 1:numel(edgeWeight)
87     A(j, edgeStart(j)) = -1;
88     A(j, edgeEnd(j)) = 1;
89 end
90
91 b = edgeWeight';
92
93 % Additional constraint that all distances must be greater than 0
94 A = [A; -eye(numberOfNodes)];
95
96 b = [b; zeros(numberOfNodes, 1)];
97
98 % Single equality constraint -- distance to i must be zero, since we start
99 % at i
100 % Aeq = zeros(3, numberOfNodes);
101 Aeq = zeros(1, numberOfNodes);
102 startNode = double('i') - double('a') + 1;
103 % nodeL = double('l') - double('a') + 1;
104 % nodeM = double('m') - double('a') + 1;
105 Aeq(1, startNode) = 1;
106 % Aeq(2, nodeL) = 1;
107 % Aeq(3, nodeM) = 1;
108 % beq = [0; 99999; 99999];
109 beq = 0;
110
111 % Minimization constraint is to maximize negative sum of distances

```

```

112 f = -ones(numberOfNodes, 1);
113
114 [x, fval, exitflag] = linprog(f,A,b,Aeq,beq);
115
116 fprintf('----- P3.C SOLUTION -----\n')
117 for j = 1:numel(x)
118     fprintf('Distance from i to %c = %2.0f \n', char('a'+ j - 1), x(j))
119 end
120
121 distanceFromNodeI = x;
122
123 %% COMBINE THE RESULTS:
124 for i = 1:numberOfNodes
125     for j = 1:numberOfNodes
126         distFromTo(i,j) = distanceToNodeI(i) + distanceFromNodeI(j);
127         if distFromTo(i,j) > 999
128             distFromTo(i,j) = NaN;
129         end
130     end
131 end
132
133 fid = fopen('P3D.solution.txt','w');
134 fprintf(fid, '----- P3.D SOLUTION -----\n');
135 for i = 1:numberOfNodes
136     for j = 1:numberOfNodes
137         fprintf(fid, 'Distance from %c to %c through i = %2.0f \n', ...
138             char('a' + i - 1), char('a'+ j - 1), distFromTo(i,j));
139     end
140 end
141
142 fclose(fid);

```

The outputted paths after running the above script (if the destination node is not reachable the path distance will be displayed as NaN):

```

1 ----- P3.D SOLUTION -----
2 Distance from a to a through i = 28
3 Distance from a to b through i = 30
4 Distance from a to c through i = 31
5 Distance from a to d through i = 36
6 Distance from a to e through i = 37
7 Distance from a to f through i = 34
8 Distance from a to g through i = 36
9 Distance from a to h through i = 24
10 Distance from a to i through i = 8
11 Distance from a to j through i = 10
12 Distance from a to k through i = 14
13 Distance from a to l through i = 15
14 Distance from a to m through i = 17
15 Distance from b to a through i = 26
16 Distance from b to b through i = 28
17 Distance from b to c through i = 29
18 Distance from b to d through i = 34

```


19 Distance from b to e through i = 35
20 Distance from b to f through i = 32
21 Distance from b to g through i = 34
22 Distance from b to h through i = 22
23 Distance from b to i through i = 6
24 Distance from b to j through i = 8
25 Distance from b to k through i = 12
26 Distance from b to l through i = 13
27 Distance from b to m through i = 15
28 Distance from c to a through i = 26
29 Distance from c to b through i = 28
30 Distance from c to c through i = 29
31 Distance from c to d through i = 34
32 Distance from c to e through i = 35
33 Distance from c to f through i = 32
34 Distance from c to g through i = 34
35 Distance from c to h through i = 22
36 Distance from c to i through i = 6
37 Distance from c to j through i = 8
38 Distance from c to k through i = 12
39 Distance from c to l through i = 13
40 Distance from c to m through i = 15
41 Distance from d to a through i = 23
42 Distance from d to b through i = 25
43 Distance from d to c through i = 26
44 Distance from d to d through i = 31
45 Distance from d to e through i = 32
46 Distance from d to f through i = 29
47 Distance from d to g through i = 31
48 Distance from d to h through i = 19
49 Distance from d to i through i = 3
50 Distance from d to j through i = 5
51 Distance from d to k through i = 9
52 Distance from d to l through i = 10
53 Distance from d to m through i = 12
54 Distance from e to a through i = 30
55 Distance from e to b through i = 32
56 Distance from e to c through i = 33
57 Distance from e to d through i = 38
58 Distance from e to e through i = 39
59 Distance from e to f through i = 36
60 Distance from e to g through i = 38
61 Distance from e to h through i = 26
62 Distance from e to i through i = 10
63 Distance from e to j through i = 12
64 Distance from e to k through i = 16
65 Distance from e to l through i = 17
66 Distance from e to m through i = 19
67 Distance from f to a through i = 22
68 Distance from f to b through i = 24
69 Distance from f to c through i = 25
70 Distance from f to d through i = 30
71 Distance from f to e through i = 31
72 Distance from f to f through i = 28

73 Distance from f to g through i = 30
74 Distance from f to h through i = 18
75 Distance from f to i through i = 2
76 Distance from f to j through i = 4
77 Distance from f to k through i = 8
78 Distance from f to l through i = 9
79 Distance from f to m through i = 11
80 Distance from g to a through i = 25
81 Distance from g to b through i = 27
82 Distance from g to c through i = 28
83 Distance from g to d through i = 33
84 Distance from g to e through i = 34
85 Distance from g to f through i = 31
86 Distance from g to g through i = 33
87 Distance from g to h through i = 21
88 Distance from g to i through i = 5
89 Distance from g to j through i = 7
90 Distance from g to k through i = 11
91 Distance from g to l through i = 12
92 Distance from g to m through i = 14
93 Distance from h to a through i = 25
94 Distance from h to b through i = 27
95 Distance from h to c through i = 28
96 Distance from h to d through i = 33
97 Distance from h to e through i = 34
98 Distance from h to f through i = 31
99 Distance from h to g through i = 33
100 Distance from h to h through i = 21
101 Distance from h to i through i = 5
102 Distance from h to j through i = 7
103 Distance from h to k through i = 11
104 Distance from h to l through i = 12
105 Distance from h to m through i = 14
106 Distance from i to a through i = 20
107 Distance from i to b through i = 22
108 Distance from i to c through i = 23
109 Distance from i to d through i = 28
110 Distance from i to e through i = 29
111 Distance from i to f through i = 26
112 Distance from i to g through i = 28
113 Distance from i to h through i = 16
114 Distance from i to i through i = 0
115 Distance from i to j through i = 2
116 Distance from i to k through i = 6
117 Distance from i to l through i = 7
118 Distance from i to m through i = 9
119 Distance from j to a through i = 22
120 Distance from j to b through i = 24
121 Distance from j to c through i = 25
122 Distance from j to d through i = 30
123 Distance from j to e through i = 31
124 Distance from j to f through i = 28
125 Distance from j to g through i = 30
126 Distance from j to h through i = 18

```
127 Distance from j to i through i = 2
128 Distance from j to j through i = 4
129 Distance from j to k through i = 8
130 Distance from j to l through i = 9
131 Distance from j to m through i = 11
132 Distance from k to a through i = 35
133 Distance from k to b through i = 37
134 Distance from k to c through i = 38
135 Distance from k to d through i = 43
136 Distance from k to e through i = 44
137 Distance from k to f through i = 41
138 Distance from k to g through i = 43
139 Distance from k to h through i = 31
140 Distance from k to i through i = 15
141 Distance from k to j through i = 17
142 Distance from k to k through i = 21
143 Distance from k to l through i = 22
144 Distance from k to m through i = 24
145 Distance from l to a through i = NaN
146 Distance from l to b through i = NaN
147 Distance from l to c through i = NaN
148 Distance from l to d through i = NaN
149 Distance from l to e through i = NaN
150 Distance from l to f through i = NaN
151 Distance from l to g through i = NaN
152 Distance from l to h through i = NaN
153 Distance from l to i through i = NaN
154 Distance from l to j through i = NaN
155 Distance from l to k through i = NaN
156 Distance from l to l through i = NaN
157 Distance from l to m through i = NaN
158 Distance from m to a through i = NaN
159 Distance from m to b through i = NaN
160 Distance from m to c through i = NaN
161 Distance from m to d through i = NaN
162 Distance from m to e through i = NaN
163 Distance from m to f through i = NaN
164 Distance from m to g through i = NaN
165 Distance from m to h through i = NaN
166 Distance from m to i through i = NaN
167 Distance from m to j through i = NaN
168 Distance from m to k through i = NaN
169 Distance from m to l through i = NaN
170 Distance from m to m through i = NaN
```