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34.1 E8 OSU ST 6540 HW#4

Let $\{x_n\}$ be successive generations wealth

$$P_{ij} = \begin{array}{|ccc|} \hline & L & M & U \\ \hline .7 & .2 & .1 & L \\ .2 & .6 & .2 & M \\ .1 & .4 & .5 & U \\ \hline \end{array} \quad S = \{L, M, U\}$$

The fraction of time spent in state U

is given by π_U .

We note that $P > 0 \Rightarrow P^n > 0$, also

$P > 0 \Rightarrow$ our 2 conditions for regularity;

$\forall i, \exists k_1, \dots, k_r \Rightarrow P_{ik_1}, \dots, P_{ik_r} > 0$ and

$\exists i \ni P_{ii} > 0$. So we know π exists

and don't have to check $(P^{(N^2)})$

Since P is regular, then $\forall i \Rightarrow \pi_i^+ = \pi_i^+ P$

$$\Rightarrow \pi_L P_{LL} + \pi_M P_{ML} + \pi_U P_{UL} = \pi_L$$

$$\pi_L P_{LM} + \pi_M P_{MM} + \pi_U P_{UM} = \pi_M$$

$$\pi_L P_{LU} + \pi_M P_{MU} + \pi_U P_{UU} = \pi_U$$

$$\left\{ \begin{array}{l} \pi_L \frac{1}{10} + \pi_M \frac{2}{10} + \pi_U \frac{1}{10} = \pi_L \\ \pi_L \frac{2}{10} + \pi_M \frac{6}{10} + \pi_U \frac{4}{10} = \pi_M \end{array} \right.$$

$$\left. \begin{array}{l} \pi_L \frac{1}{10} + \pi_M \frac{2}{10} + \pi_U \frac{5}{10} = \pi_U \\ \pi_L + \pi_M + \pi_U = 1 \end{array} \right.$$

$$\pi_L + \pi_M + \pi_U = 1$$

$$\textcircled{1} \quad -3 \pi_L + 2 \pi_M + \pi_U = 0 \Rightarrow \pi_U = 3\pi_L - 2\pi_M$$

$$\textcircled{2} \quad 2 \pi_L - 4 \pi_M + 4 \pi_U = 0 \Rightarrow \pi_L = 6/7 \pi_M$$

$$\textcircled{3} \quad \pi_L + 2\pi_M - 5\pi_U = 0 \Rightarrow \frac{6}{7}\pi_M + 2\pi_M - 5\left(3 \cdot \frac{6}{7}\pi_M - 2\pi_M\right) = 0$$

can't use eq \textcircled{3} have to use constraint $\sum \pi_i = 1$

$$\frac{6}{7}\pi_M + \pi_M + \left(3 \cdot \frac{6}{7}\pi_M - 2\pi_M\right) = 1$$

$$\frac{13}{7}\pi_M + \frac{4}{7}\pi_M = 1 \Rightarrow \pi_M = 7/17$$

$$\pi_L = \frac{6}{17}$$

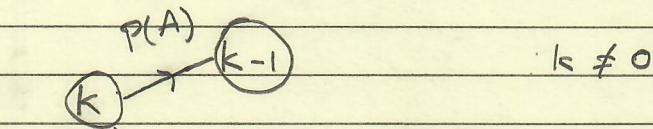
$$\pi_U = 4/17$$

Same as
Ma lab!
②

§ 4.1 p1

Let $\{X_n\}$ be the MC with $S = \{0, 1, \dots, 5\}$

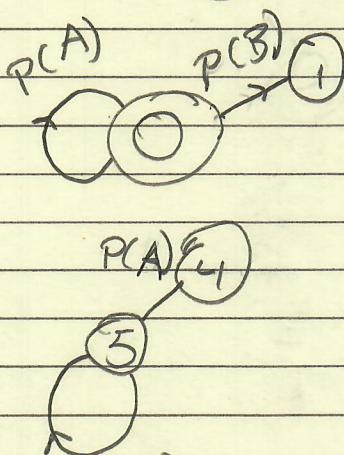
$X_n = \# \text{ balls in urn } A.$



$$k \neq 0$$

$$P(A) = \text{Prob}(A \text{ is selected}) =: p$$

$$P(B) = 1 - P(A) =: q$$



Then

$$P = \begin{array}{|cccccc|c} \hline & 0 & 1 & 2 & 3 & 4 & 5 & \\ \hline 0 & p & q & 0 & 0 & 0 & 0 & 0 \\ p & 0 & q & 0 & 0 & 0 & 0 & 1 \\ 0 & p & 0 & q & 0 & 0 & 0 & 2 \\ 0 & 0 & p & 0 & q & 0 & 0 & 3 \\ 0 & 0 & 0 & p & 0 & q & 0 & 4 \\ 0 & 0 & 0 & 0 & p & 0 & q & 5 \\ \hline \end{array}$$

We have to check P is regular

We see all states are communicating, and

that $\exists i \ni P_{ii} > 0$. Writing out $\pi^+ = \pi^+ P$

$$\pi_0 p + \pi_1 p = \pi_0 \Rightarrow \pi_0 q = \pi_1 p$$

$$\pi_0 q + \pi_2 p = \pi_1 \Rightarrow \pi_1 p + \pi_2 p = \pi_1 \Rightarrow \pi_1 q = \pi_2 p$$

$$\pi_1 q + \pi_3 p = \pi_2 \Rightarrow \pi_2 p + \pi_3 p = \pi_2 \Rightarrow \pi_2 q = \pi_3 p$$

$$\pi_2 q + \pi_4 p = \pi_3 \Rightarrow \pi_3 p + \pi_4 p = \pi_3 \Rightarrow \pi_3 q = \pi_4 p$$

$$\pi_3 q + \pi_5 p = \pi_4 \Rightarrow \pi_4 p + \pi_5 p = \pi_4 \Rightarrow \pi_4 q = \pi_5 p$$

$$\pi_4 q + \pi_5 q = \pi_5 \Rightarrow \pi_5 p + \pi_5 q = \pi_5 \Leftarrow \text{not useful}$$

$$\pi_1 + \pi_2 + \pi_3 + \pi_4 + \pi_5 = 1 \quad \text{will be last eq.}$$

S 4.1 P 1

Putting our expressions

$$\pi_{i+1} = \pi_i \frac{q}{p}$$

into $\sum \pi_i = 1$

$$\pi_0 + \pi_0 \frac{q}{p} + \pi_1 \left(\frac{q}{p} \right)^2 + \dots + \pi_1 \left(\frac{q}{p} \right)^5 = 1$$

$$\pi_0 = \sum_{i=0}^5 \left(\frac{q}{p} \right)^i, \quad \pi_i = \left(\frac{q}{p} \right)^i / \sum_{i=0}^5 \left(\frac{q}{p} \right)^i$$

Finally - the fraction of time that urn A
is empty in the long run is π_0

§ 4.1 p 4

$\{X_n\}$ mc $|S| < \infty$ has IP and a limiting distribution π . What fraction of transitions are from k to m in the long run?

We know the proportion of time in states

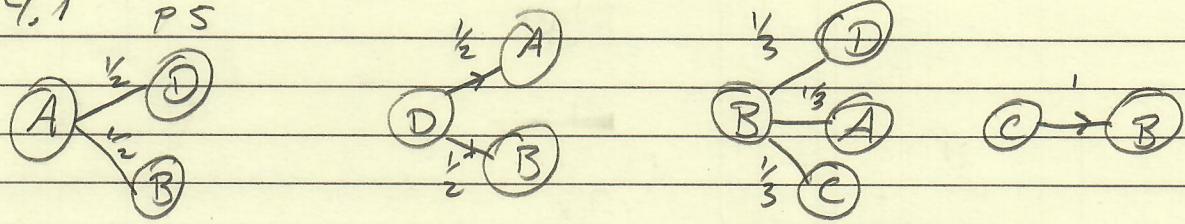
k and m are π_k, π_m

We expect the answer to be the proportion of time spent in k times the proportion of transitions from k to m , but this is just P_{km}

So

$\pi_k P_{km}$ is the long run fraction of transitions from $k \rightarrow m$.

S 4.1 PS



$$P = \begin{array}{c|ccccc} & A & B & C & D \\ \hline A & 0 & \frac{1}{2} & 0 & \frac{1}{2} \\ B & \frac{1}{3} & 0 & \frac{1}{3} & \frac{1}{3} \\ C & 0 & 1 & 0 & 0 \\ D & \frac{1}{2} & \frac{1}{2} & 0 & 0 \end{array}$$

We're asked to find π_D - the long run probability of finding state D. All states are communicating but, notice we don't have our regularity condition $\exists i \in \mathbb{N} : P_{ii} > 0$. We'll finish calculations in Matlab.

First checking $P^{(5^2)} > 0$ and then solving

$$\pi^+ = \pi^+ P$$

$$P = \begin{bmatrix} S^0_p & & \\ S^1_p & & \\ \vdots & & \\ S^{N-1}p & & \end{bmatrix} \quad \text{where } S \text{ is cyclic permutation}$$

$$S : (1, \dots, n) \rightarrow (n, 1, \dots, n-1)$$

$$p = (P_0, \dots, P_N)$$

$$S^i p = (P_{N-(i+1)}, \dots, P_N, P_0, P_1, \dots, P_{N-i})$$

$$\text{We notice } \sum_{j=0}^N P_{ji} = \sum_{i=0}^N (S^i p)_i$$

$$= P_{N-(i+1)} + \dots + P_N + P_0 + \dots + P_{N-i}$$

$$= \sum p_i = 1 \quad \text{so } P \text{ is obviously}$$

stochastic. Now, if P is regular, then we can conclude $\pi = (\pi_0, \dots, \pi_N) \sim \text{Unif}(N)$. We don't know p , but the assumption $\alpha p \neq 0$ tells us a few things. First, there must be at least one other non-zero $p_i \neq 0$ since $\sum p_i = 1$. There are also no absorbing states, since that would require $p_j = 1$ for some j .

Let's consider the worst case $p_0, p_k \neq 0, p_i = 0 \text{ if } i \neq 0, k$

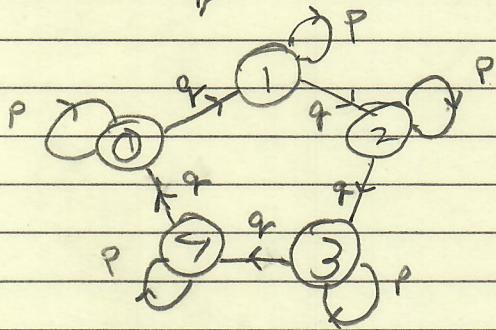
$$P = \begin{array}{c|cc|cc|cc|c} & 0 & 1 & \dots & k & k+1 & \dots & N \\ \hline p & 0 & 0 & & q & 0 & & \\ 0 & p & & & 0 & q & & \\ 0 & & \ddots & & & & & N-(k+1) \\ 0 & & & p & & & q & N-k \\ q & & \ddots & & p & & q & N-(k+1) \\ 0 & & & q & & 0 & p & N \end{array}$$

Suppose $N=4$ $P_0, P_1 \neq 0$ $P_2 = P_3 = P_4 = 0 \Rightarrow P_0 = p$ $P_1 = 1-p = q$

$$P = \begin{bmatrix} 0 & 1 & 2 & 3 & 4 \\ p & q & 0 & 0 & 0 \\ 0 & p & q & 0 & 0 \\ 0 & 0 & p & q & 0 \\ 0 & 0 & 0 & p & q \\ q & 0 & 0 & 0 & p \end{bmatrix}$$

and looking

at our transitions
we see



a similar argument
for the general case
should show all states
are communicating - with
some additional modulo k
considerations

S4.2 P

$$P(T=1) = a_1 \quad P(T=2) = a_2 \quad P(T=3) = a_3 \quad P(T=4) = a_4$$

Let X_n be age of component at n

$\{X_n\}$ is M.C. $\{X_n\}$ is a success run

M.C. with $S = \{0, 1, 2, 3, 4\}$. Here $X=0$

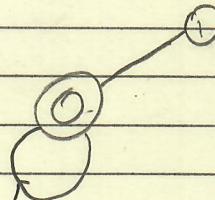
let T_k be lifetime of k^{th} component. $\sum_{k=1}^n T_k$

gives us the time the n^{th} bulb fails.

$\sum T_{k+1}$ is when the $(n+1)^{th}$ bulb is put in service.

We're going to guess $P_k = \frac{a_{k+1}}{a_{k+1} + \dots + a_N}$

$k=0, N-1$ when $k=4$ we replace component



Better! Set up as replacement policy on

$$S = \{0, 1, 2, 3, 4, 5\}$$

and we specify a new component upon reaching 5. This is consistent with since

$$P(X_{n+1}=4 | X_n=3) = ? \text{ or } .4$$

Writing

	0	1	2	3	4	5	
p_0	a_1	$1-a_1$					0
p_1		p_1	$1-p_1$				1
p_2			p_2				2
p_3			p_3				3
			p_4				4
				p_5			5
					1		

$$P_0 = \frac{a_1}{a_1 + a_2 + a_3 + a_4} = a_1$$

$$P_1 = \frac{a_2}{a_2 + a_3 + a_4} = P(X_{n+1}=2 | X_n=1)$$

? check!

$$P(T=1) = a_1$$

$$P(T \geq 1) = 1 - a_1 = \text{acc}$$

$$P(T \geq 2) = a_3 + a_4$$

4.2 86

$$P(X_{n+1} = k+1 | X_n = k) = \frac{P(X_{n+1} = k+1, X_n = k)}{P(X_n = k)}$$

$$P(X_1 = 1 | X_0 = 0) = \frac{P(X_1 = 1, X_0 = 0)}{P(X_0 = 0)}$$

$$X_0 = 0 \Rightarrow P(X_0 = 0) = 1 \quad P(X_1 = 1, X_0 = 0) = P(X_1 = 1) = a_1$$

$$P(X_2 = 2 | X_1 = 1) = P(T \geq 1) = P(T \geq 2 | T > 1)$$

$$P(X \leq 3) =$$

$$\overline{\overline{1 \ 1 \ 1}} \quad \begin{array}{lll} X \leq 1 & \varepsilon_1 \\ X \leq 2 & \varepsilon_2 \\ X \leq 3 & \varepsilon_3 \end{array} \quad \cap \varepsilon_i = \varepsilon_1$$

4
1
0

Let $P_k = \frac{a_{k+1}}{a_{k+1} \dots a_{N-1}}$ $N = 4 \text{ or } 5$
sort!

$k = 0, 1, 2,$	$(4-2)$	a_{k+1}
3	$(5-2)$	$a_{k+1} \dots a_{N-1}$

$N=5$ when

$T=4$

$\textcircled{1} - \textcircled{0} \Rightarrow$
replacement @ 5

THIS ONE!

	0	1	2	3	4
0	p_0	q_{00}	0	0	0
1	p_1	0	q_{11}	0	0
2	p_2	0	0	q_{22}	0
3	p_3	0	0	0	q_{33}
4	1	0	0	0	0

Solve this system $\pi^+ = \pi^+ \mathbf{P}$

letting $A_j = a_j + \dots a_3$

$P_0 =$

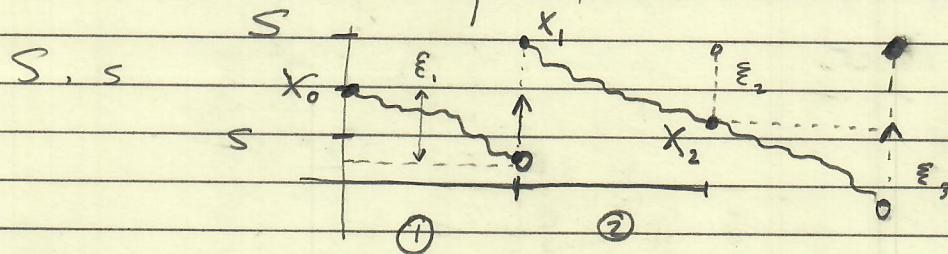
§ 9.2

$$n = 0, 1, 2, \dots$$

En demand

$$\xi_i + \xi_{n+1} (a_0, a_1, \dots)$$

$$\sum a_k = 1 \quad X_n \text{ quantity on hand}$$

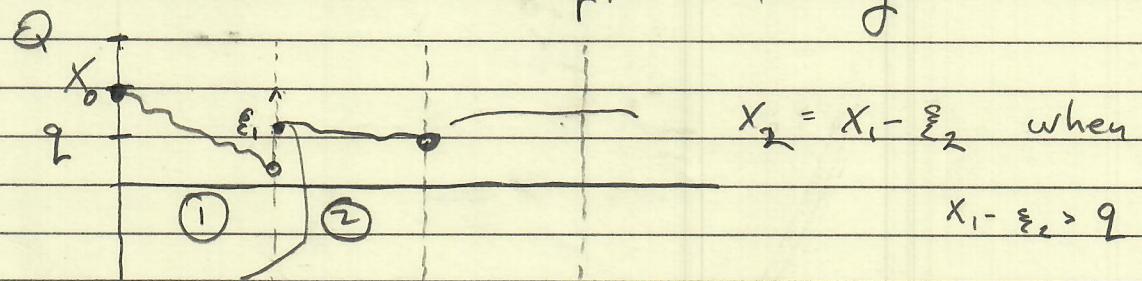


$$X_{n+1} = \begin{cases} X_n - \xi_{n+1} & S < X_n \leq S \\ S - \xi_{n+1} & X_n \leq S \end{cases} \quad \begin{matrix} \xi_i \leq S \text{ if } i \\ (\text{an assumption since we could set } X_0 = S+k) \end{matrix}$$

$$S = \{S, S-1, \dots, 0, -1, \dots, -S+1, \dots\}$$

S is determined by S and a_k 's

Let's draw the q, Q policy



$$X_1 = X_0 - \xi_1 + Q \quad \text{when } X_0 - \xi_1 \leq q$$

$$X_{n+1} = \begin{cases} X_n - \xi_{n+1} + Q & X_n - \xi_{n+1} \leq q \\ X_n - \xi_{n+1} & X_n - \xi_{n+1} > q \end{cases}$$

$$j = \begin{cases} i - \xi_{n+1} + q & i - \xi_{n+1} \leq q \\ i - \xi_n & i - \xi_n > q \end{cases}$$

$$\xi_n = \begin{cases} i - (j + q) & i - \xi_{n+1} \leq q \\ i - j & i - \xi_n > q \end{cases}$$

$$j-i = \xi_n + Q$$

$$\xi_n = i-j-Q$$

§ 4.2 P 1

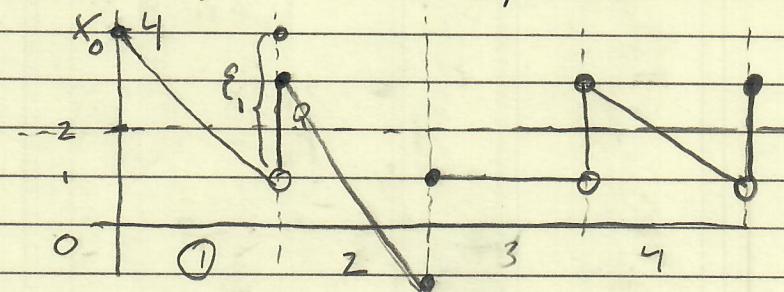
p. 4

a) Suppose $X_0 = 4$ $\xi_1 = 3$ $\xi_2 = 4$ $\xi_3 = 0$ $\xi_4 = 2$

Find X_1, X_2, X_3, X_4 for the q, Q process

with $q = Q = 2$

Let's draw!



$$X_1 = 3 \quad X_2 = X_1 - \xi_1 + Q = 1 \quad X_3 = X_2 - \xi_2 + Q = 3 \quad X_4 = X_3 - \xi_3 + Q = 3$$

Now let ξ_i i.i.d. $P(\xi = k) = (a_1, a_2, a_3, a_4)$

$$P_{ij} = P(X_{n+1} = j | X_n = i) = \begin{cases} P(\xi_{n+1} = i-j) & i > q \\ P(\xi_{n+1} = i - (j+Q)) & i \leq q \end{cases}$$

What about assumption That demand that is not filled is lost (i.e. no backordering?) does this mean $X_n \geq 0$? If that's true we have 3 cases

to worry about $X_n - \xi_{n+1} \geq q$, $X_n - \xi_{n+1} \leq q$

$X_n - \xi_{n+1} + Q < 0$ the last one \Rightarrow

$$X_n - \xi_{n+1} + Q \quad X_n - \xi_{n+1} \leq q,$$

$$X_{n+1} = \begin{cases} Q & \text{when } X_n - \xi_{n+1} < 0 \\ X_n - \xi_{n+1} & X_n - \xi_{n+1} \geq q \end{cases}$$

§ 4.2 p 2

Let $(x, y) = (\# \text{ computers operating},$ $\begin{cases} 1 & \text{one day of labor} \\ 0 & \text{on broken one} \\ 0 & \text{otherwise} \end{cases}$)

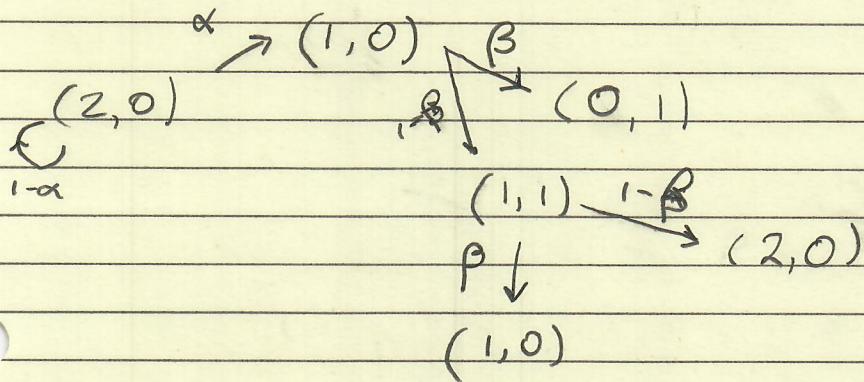
Then $S = \{(2, 0), (1, 0), (1, 1), (0, 1)\}$

$= \{\text{both operational, one is working + second needs 2 days work}$

$\text{one is working + second needs one day work}$

a)

$\text{one needs one day work + other needs 2 days}$

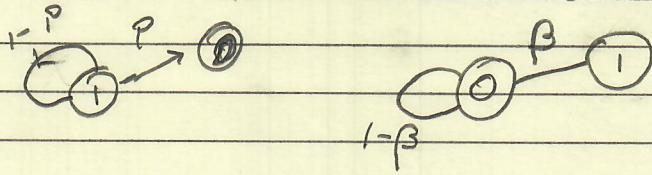


$$P = \begin{bmatrix} (2,0) & (1,0) & (1,1) & (0,1) \\ \begin{matrix} 1-\alpha & \alpha & 0 & 0 \\ 0 & 0 & 1-\beta & \beta \\ (1-\beta) & \beta & 0 & 0 \\ 0 & 1 & 0 & 0 \end{matrix} \end{bmatrix} \begin{matrix} (2,0) \\ (1,0) \\ (1,1) \\ (0,1) \end{matrix}$$

§ 4.2 p 6

Let

$$X_n = \begin{cases} 1 & \text{computer operational} \\ 0 & \text{otherwise} \end{cases}$$



We recognize $N \sim \text{Geom}(\beta)$ which we know to be memoryless - $P(N=1) = 1 - P(N \neq 1) = \beta$

i.e. success on first trial. If $0 \rightarrow 0$ the process starts over.

We can construct \mathbf{P} from our transitions now

$$\mathbf{P} = \begin{bmatrix} 0 & 1 \\ 1-\beta & \beta \\ \beta & 1-p \end{bmatrix}$$

The long run proportion of time the computer is operational is given by π , which can be obtained by solving $\pi^+ = \pi^+ \mathbf{P}$

$$\pi_0(1-\beta) + \pi_1 p = \pi_0 \Rightarrow \pi_0 = \frac{p}{\beta} \pi_1$$

$$\pi_0 \beta + \pi_1 (1-p) = \pi_1$$

$$\pi_0 + \pi_1 = 1 \Rightarrow \pi_1 = \frac{1-p}{1+\beta} \Rightarrow \pi_0 = \frac{\beta}{1+\beta} = \frac{1}{1+\frac{p}{\beta}}$$

This makes qualitative sense; β small

$\Rightarrow N$ big $\Rightarrow \pi_1$ smaller as expected.

§ 4.3 § 3

a) By inspection

$$① \leftrightarrow ②$$

$$① \leftrightarrow ③$$

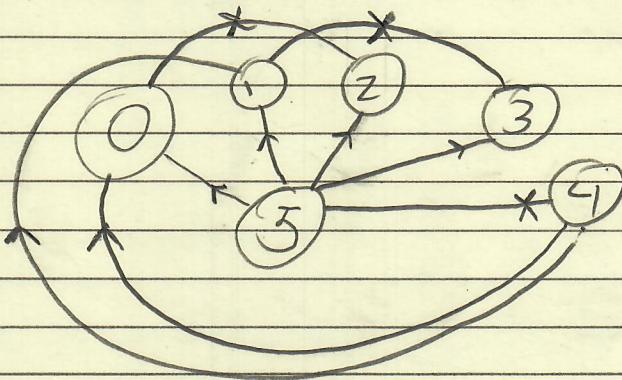
$$⑤ \leftrightarrow ④$$

$$0 \rightarrow 1 \quad 0 \rightarrow 3 \quad 0 \rightarrow 4 \quad 0 \rightarrow 5$$

$$1 \rightarrow \{0, 2, 4, 5\} \quad 2 \rightarrow \{1, 3, 4, 5\} \quad 3 \rightarrow \{0, 2, 4, 5\}$$

$$4 \rightarrow \{0, 1, 5\}$$

	0	1	2	3	4	5
0	$\frac{1}{3}$	0	$\frac{2}{3}$	0	0	0
1	0	$\frac{1}{4}$	0	$\frac{3}{4}$	0	0
2	$\frac{2}{3}$	0	$\frac{1}{3}$	0	0	0
3	0	$\frac{1}{5}$	0	$\frac{4}{5}$	0	0
4	$\frac{1}{4}$	$\frac{1}{4}$	0	0	$\frac{1}{4}$	$\frac{1}{4}$
5	$\frac{1}{6}$	$\frac{1}{6}$	$\frac{1}{6}$	$\frac{1}{6}$	$\frac{1}{6}$	$\frac{1}{6}$



ignoring P_{ii} transitions

$$C_1 = \{0, 2\} \quad C_2 = \{1, 3\} \quad C_3 = \{4, 5\}$$

Are the communicating classes.

Note $C_3 \rightarrow C_1$ and $C_3 \rightarrow C_2$ by $4 \rightarrow 0, 4 \rightarrow 1$ and $C_3 \rightarrow C_1, C_3 \rightarrow C_2$ by $5 \rightarrow 0, 5 \rightarrow 1$. All of the states are recurrent except for ④ and ⑤.

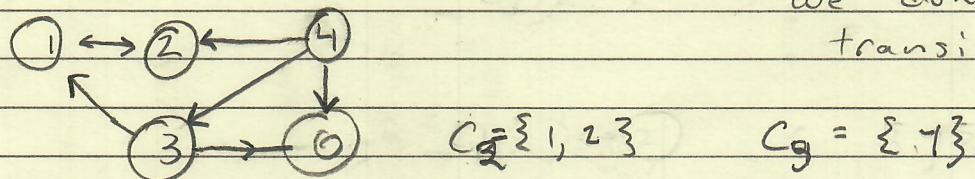
§ 4.3 E3

b)	0	1	2	3	4	5	
$P =$	1	0	0	0	0	0	0
	0	$\frac{3}{4}$	$\frac{1}{4}$	0	0	0	1
	0	$\frac{1}{8}$	$\frac{7}{8}$	0	0	0	2
	$\frac{1}{4}$	$\frac{1}{4}$	0	$\frac{1}{8}$	$\frac{3}{8}$	0	3
	$\frac{1}{3}$	0	$\frac{1}{6}$	$\frac{1}{4}$	$\frac{1}{4}$	0	4
	0	0	0	0	0	1	5

By inspection (our other choice is to calculate $P^{(n)}$)

$C_0 = \{0\}$ $C_1 = \{5\}$ are absorbing - hence not transient.

we don't draw P_{ii} transitions



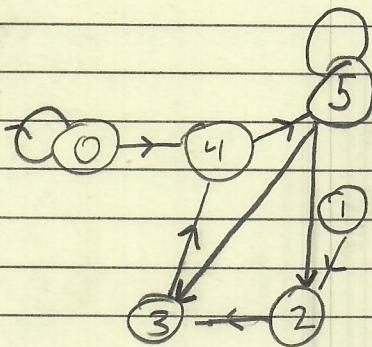
$\textcircled{1} \rightarrow \{2, 3, 4, 0\}$

$C_0 = \{0\}$ $C_1 = \{5\}$ $C_2 = \{1, 2\}$ $C_3 = \{3, 4\}$ are the communicating classes

$\textcircled{1}, \textcircled{3}$ are transient $\textcircled{0} \textcircled{2} \textcircled{4} \textcircled{5}$ are recurrent

§ 9.3 Ex

	0	1	2	3	4	5	
$P =$	$\frac{1}{2} 0$	0 0	0 0	$\frac{1}{2} 0$	0 0	0 0	0
	0 0	1 0	0 0	0 0	0 0	0 0	1
	0 0	0 1	0 0	0 0	0 0	0 0	2
	0 0	0 0	0 1	0 0	0 0	0 0	3
	0 0	0 0	0 0	0 1	0 0	0 0	4
	0 0	$\frac{1}{3}$ $\frac{1}{3}$	$\frac{1}{3}$ $\frac{1}{3}$	0 0	$\frac{1}{3}$ $\frac{1}{3}$	$\frac{1}{3}$ $\frac{1}{3}$	5



$$C_1 = \{0\} \quad C_2 = \{1\}$$

$$C_3 = \{2, 3, 4, 5\}$$

We only have to compute $d(0)$, $d(1)$ and $d(i)$ for one $i \in \{2, 3, 4, 5\}$

Since $P_{11} > 0$ $P_{55} > 0$ we have $d(0) = 1$ and $d(5) = 1$

Since periodicity is constant within a communicating class - $d(2) = d(3) = d(4) = d(5) = 1$

Finally for $d(1)$ notice the second column

of $P^{(n)}$ is always 0 - i.e. $P_{11}^{(n)} = 0 \forall n \geq 1$

$$\therefore d(1) = 0$$