

Enigma

Enigma cipher machine has several rotors, and each of them has a value. Each value pair of two rotors should be coprime. To develop such a cipher, can you calculate how many ordered coprime integer pairs are there in a given range?

Naive Approach

```
for i = lo, lo + 1, ..., hi
    for j = lo, lo + 1, ..., hi
        if gcd(i, j) == 1
            count += 1
```

Analysis

```
for i = <hi - lo + 1 times>  
    for j = <hi - lo + 1 times>  
        if <log(min(i, j))-ish stuff>  
            count += 1
```

Time complexity: $O((hi - lo)^2 \log(hi))$

WHAT IF WE TRIED
MORE POWER?



Observation

$$42 = 2 * 3 * 7$$

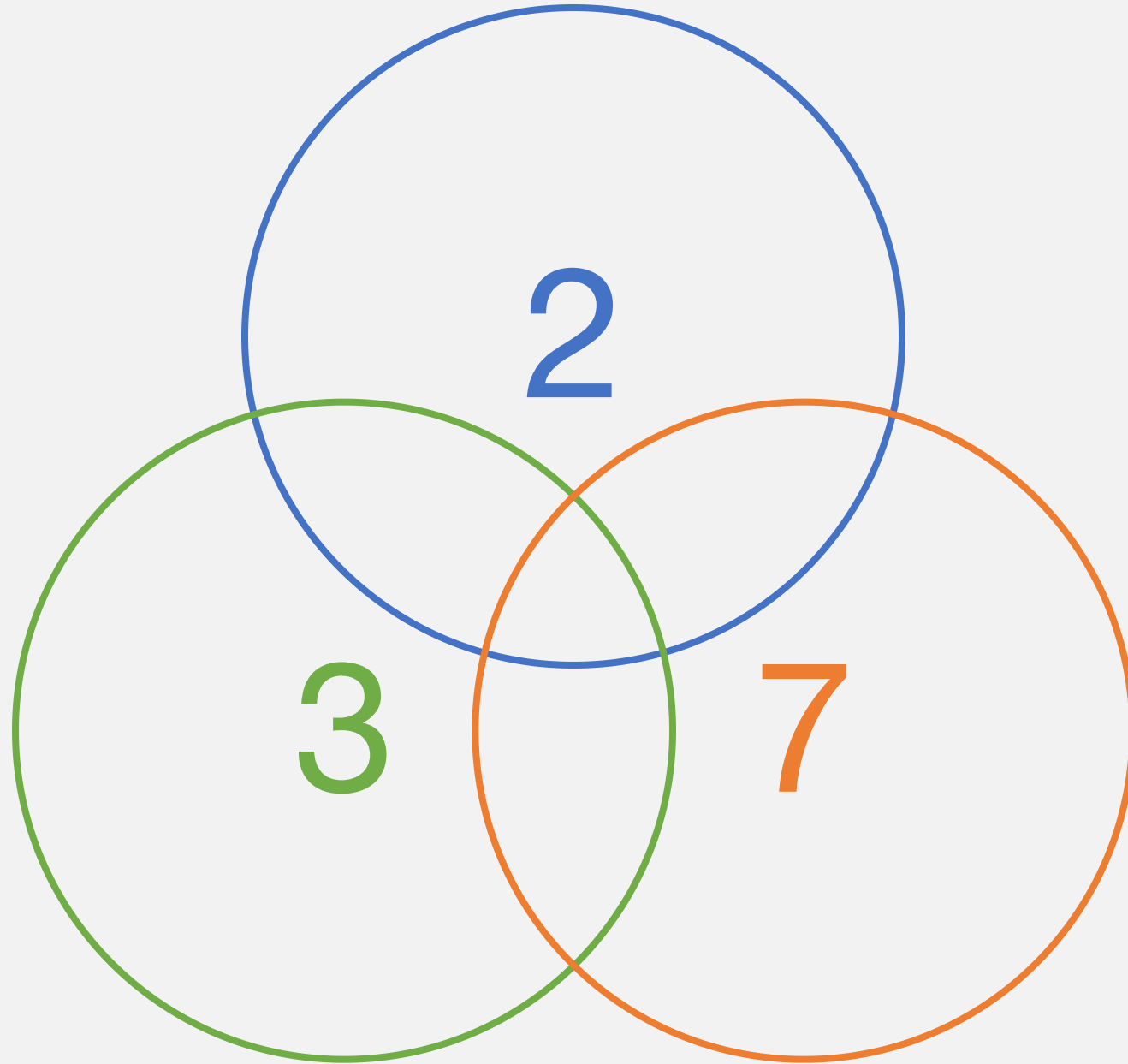
If number X and 42 are coprime...

Observation

$$42 = 2 * 3 * 7$$

If number X and 42 are **not** coprime

$$X = 2A \text{ or } X = 3B \text{ or } X = 7C$$



Inclusion-exclusion Principle

let $f(<...>)$ =

number of integers that divisible
by $<...>$ within the range $lo..hi$

$f(a \text{ or } b) =$

$f(a) + f(b) - f(a \text{ and } b)$

Inclusion-exclusion Principle

$$f(a \text{ or } b \text{ or } c) =$$

$$f(a) + f(b) + f(c)$$

$$- f(a \text{ and } b) - f(a \text{ and } c) - f(b \text{ and } c)$$

$$+ f(a \text{ and } b \text{ and } c)$$

Inclusion-exclusion Principle

$f(\text{at least one member in set } S^*) =$

sum of $f(\text{all member in } S)$

where S is an odd-size subset of S^*

- sum of $f(\text{all member in } S)$

where S is an even-size subset of S^*

How to Calculate $f(\dots)$ Components

For all member in S :

$$f(2, 3, \text{ and } 7) = f(2 * 3 * 7) = f(42)$$

$$f(\text{all member in } S) = f(\text{product of } S)$$

For a single integer:

$$f(42) = \text{floor}(hi / 42) - \text{floor}((lo - 1) / 42)$$

$$f(x) = \text{floor}(hi / x) - \text{floor}((lo - 1) / x)$$

I-e Principle Approach

```
for i = lo, lo + 1, ..., hi
    S* = factor(i)
    foreach S = subset of S*
        x = PI(S)
        if |S| % 2 == 1
            count += floor(hi / x)
                  - floor((lo - 1) / x)
        else
            count -= floor(hi / x)
                  - floor((lo - 1) / x)
```

Analysis

```
for i = <hi - lo + 1 times>  
    <sqrt(i)-ish stuff>  
    foreach S = <"count of integers  
        divisible by i" times>  
        <|S|-ish stuff>
```

Analysis

```
for i = <hi - lo + 1 times>  
    <sqrt(i)-ish stuff>  
    foreach S = <1 + 0.5 + 0.333  
        + ... = log(i)>  
        <log(log(i))-ish stuff>
```

Time complexity: $O((hi - lo)^{1.5})$

WHAT IF WE TRIED
MORE POWER?



Dynamic Programming

Factor Tree

$2 \rightarrow 1$

$3 \rightarrow 1$

$4 \rightarrow 2 \rightarrow 1$

$6 \rightarrow 2 \rightarrow 1$

$42 \rightarrow 6 \rightarrow 2 \rightarrow 1$

...

Factor Tree

```
for i = lo, lo + 1, ..., hi  
    if factor[i] is null  
        fill(i)
```

Factor Tree

```
procedure fill(i)
  for j = i, i * i, ..., hi
    for k = 1, 2, ..., hi / j
      if factor[k] < i
        put_node(i, j, k)
```

Factor Tree

```
procedure put_node(i, j, k)
    factor[k * j] = i
    remain[k * j] =
        remain[k] * factor[k]
```

Analysis

```
for i = <hi - lo + 1 times>  
    <sqrt(i) - ish stuff>  
    foreach S = <1 + 0.5 + 0.333  
        + ... = log(i)>  
        <log(log(i)) - ish stuff>
```

Time complexity:

$O(hi \log(hi) \log(\log(hi)))$

WHAT IF WE TRIED
MORE POWER?



Factor Tree

42 -> 6 -> 2 -> 1

84 -> 12 -> 4 -> 2 -> 1

126 -> 24 -> 8 -> 4 -> 2 -> 1

...

Factor Tree

42 -> 6 -> 2 -> 1

84 -> 12 -> 4 -> 2 -> 1

126 -> 24 -> 8 -> 4 -> 2 -> 1

...

Analysis

Build the tree: $O(h_i \log(\log(h_i)))$

Number of nodes: $O(h_i / \log(h_i))$

I-e Principle: $O(\log(h_i) \log(\log(h_i)))$
(per node)

Comprehensive: $O(h_i \log(\log(h_i)))$