

GrumpyIR Static Semantics

April 4, 2021

The GrumpyIR static semantics (i.e., the type system) is given as a five-place relation on a function typing context Δ , a variable typing context Γ , a store typing Σ , an expression e , and a type T , written $\Delta \mid \Gamma \mid \Sigma \vdash e : T$, pronounced "under Δ , Γ , and Σ , expression e has type T ". Formally, the static semantics is taken to be the smallest relation satisfying the following typing rules:

Values and names

$\frac{\text{T-N}}{\Delta \mid \Gamma \mid \Sigma \vdash n : \text{i32}}$	$\frac{\text{T-TRUE}}{\Delta \mid \Gamma \mid \Sigma \vdash \text{true} : \text{bool}}$	$\frac{\text{T-FALSE}}{\Delta \mid \Gamma \mid \Sigma \vdash \text{false} : \text{bool}}$
$\frac{\text{T-TT}}{\Delta \mid \Gamma \mid \Sigma \vdash \text{tt} : \text{unit}}$	$\frac{\text{T-LOC} \quad (l : \text{array } T) \in \Sigma}{\Delta \mid \Gamma \mid \Sigma \vdash l : \text{array } T}$	$\frac{\text{T-VAR} \quad (x : T) \in \Gamma}{\Delta \mid \Gamma \mid \Sigma \vdash x : T}$
$\frac{\text{T-FUNPTR} \quad (p : \text{fun } (T_1, T_2, \dots, T_n) T) \in \Delta}{\Delta \mid \Gamma \mid \Sigma \vdash p : \text{fun } (T_1, T_2, \dots, T_n) T}$		

Unary and binary operators

$\frac{\text{T-NEG} \quad \Delta \mid \Gamma \mid \Sigma \vdash e : \text{bool}}{\Delta \mid \Gamma \mid \Sigma \vdash (\text{neg } e) : \text{bool}}$		
$\frac{\text{T-BINOP-I32} \quad \Delta \mid \Gamma \mid \Sigma \vdash e_1 : \text{i32} \quad \Delta \mid \Gamma \mid \Sigma \vdash e_2 : \text{i32} \quad b \in \{+, *, -, /\}}{\Delta \mid \Gamma \mid \Sigma \vdash (b \ e_1 \ e_2) : \text{i32}}$		
$\frac{\text{T-BINOP-BOOL} \quad \Delta \mid \Gamma \mid \Sigma \vdash e_1 : \text{i32} \quad \Delta \mid \Gamma \mid \Sigma \vdash e_2 : \text{i32} \quad b \in \{<, ==\}}{\Delta \mid \Gamma \mid \Sigma \vdash (b \ e_1 \ e_2) : \text{bool}}$		

Let-expressions and sequencing

$$\frac{\text{T-LET} \quad \Delta \mid \Gamma \mid \Sigma \vdash e_1 : T_1 \quad \Delta \mid \Gamma, (x : T_1) \mid \Sigma \vdash e_2 : T_2}{\Delta \mid \Gamma \mid \Sigma \vdash (\text{let } x \ e_1 \ e_2) : T_2}$$

$$\frac{\text{T-SEQ} \quad \Delta \mid \Gamma \mid \Sigma \vdash e_1 : T_1 \quad \Delta \mid \Gamma \mid \Sigma \vdash e_2 : T_2}{\Delta \mid \Gamma \mid \Sigma \vdash (\text{seq } e_1 \ e_2) : T_2}$$

Arrays

$$\frac{\text{T-ALLOC} \quad \Delta \mid \Gamma \mid \Sigma \vdash e_{size} : \text{i32} \quad \Delta \mid \Gamma \mid \Sigma \vdash e_{init} : T}{\Delta \mid \Gamma \mid \Sigma \vdash (\text{alloc } e_{size} \ e_{init}) : \text{array } T}$$

$$\frac{\text{T-SET} \quad \Delta \mid \Gamma \mid \Sigma \vdash e_{arr} : \text{array } T \quad \Delta \mid \Gamma \mid \Sigma \vdash e_{ix} : \text{i32} \quad \Delta \mid \Gamma \mid \Sigma \vdash e : T}{\Delta \mid \Gamma \mid \Sigma \vdash (\text{set } e_{arr} \ e_{ix} \ e) : \text{unit}}$$

$$\frac{\text{T-GET} \quad \Delta \mid \Gamma \mid \Sigma \vdash e_{arr} : \text{array } T \quad \Delta \mid \Gamma \mid \Sigma \vdash e_{ix} : \text{i32}}{\Delta \mid \Gamma \mid \Sigma \vdash (\text{get } e_{arr} \ e_{ix}) : T}$$

Conditionals and function calls

$$\frac{\text{T-COND} \quad \Delta \mid \Gamma \mid \Sigma \vdash e_{cond} : \text{bool} \quad \Delta \mid \Gamma \mid \Sigma \vdash e_1 : T \quad \Delta \mid \Gamma \mid \Sigma \vdash e_2 : T}{\Delta \mid \Gamma \mid \Sigma \vdash (\text{cond } e_{cond} \ e_1 \ e_2) : T}$$

$$\frac{\text{T-CALL} \quad \Delta \mid \Gamma \mid \Sigma \vdash e : \text{fun } (T_1, T_2, \dots, T_n) \ T \quad \Delta \mid \Gamma \mid \Sigma \vdash e_i : T_i}{\Delta \mid \Gamma \mid \Sigma \vdash (\text{call } e \ e_1 \ e_2 \ \dots \ e_n) : T}$$