Exploit Mitigation using Multi-Variant Execution

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Adobe Patches 10 Critical Vulnerabilities in Flash Player, Shockwave Player, and ColdFusion

Posted on April 9th, 2013 by Derek Erwin

Zero-Dav Vulnerability Bypasses Apple's Security Featur Op June 2012, 10:59 ces:

Updat Adobe Flash update closes several critical holes

28 March 2016, 8:15 am EDT By Horia Ungureanu Tech Times

Adobe Issues Emergency Updates For Zero-Day Flaw in Flash Player

Memory corruption flaw is being exploited in the wild to distribute ransomware samples like Locky and Cerber.

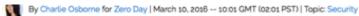
Update Flash now! Adobe releases patch, fixing critical security holes

BY GRAHAM CLULEY POSTED 22 SEP 2015 - 09:24AM

stems.

(Another) Update To Adobe Flash Addresses Latest 0-Day Vulnerability

April 15, 2011



Adobe Fixes 18 Vulnerabilities in Flash Player

By Eduard Kovacs on November 12, 2014

Critical Adobe Flash bug under active attack currently has no patch

Exploit works against the most recent version; Adobe plans update later this week.

by Dan Goodin - Jun 14, 2016 12:50pm PDT

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Adobe Releases Security Update for 19 'Critical' Vulnerabilities in Flash Player



DAVID BISSON

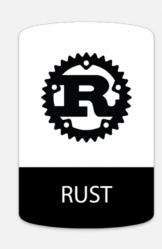
DEC 29, 2019



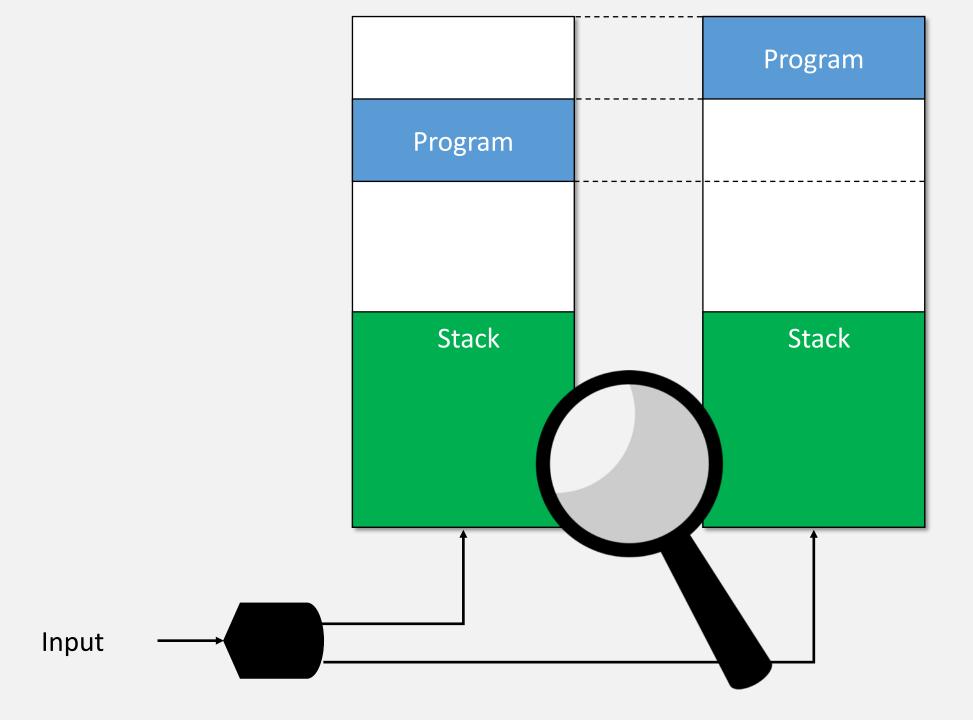
Possible Solutions

- Type-Safe Languages (e.g. Rust)
- Mitigations:
 - Integrity-Based (e.g. CFI)
 - Randomization-Based (e.g. ASLR)
- Multi-Variant Execution Environments (MVEEs)

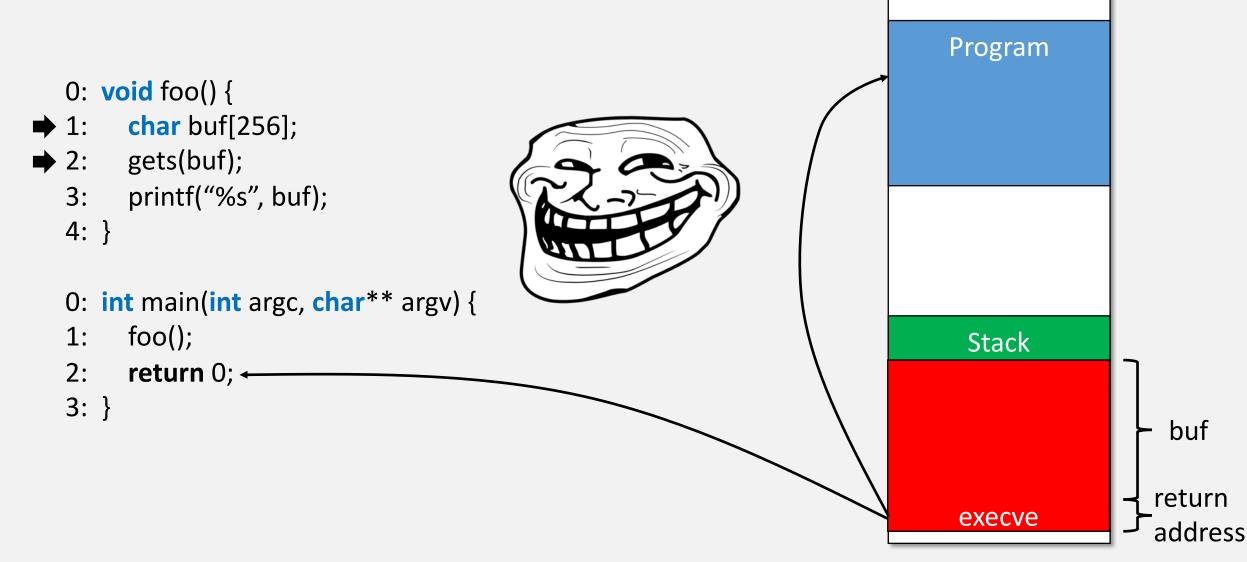


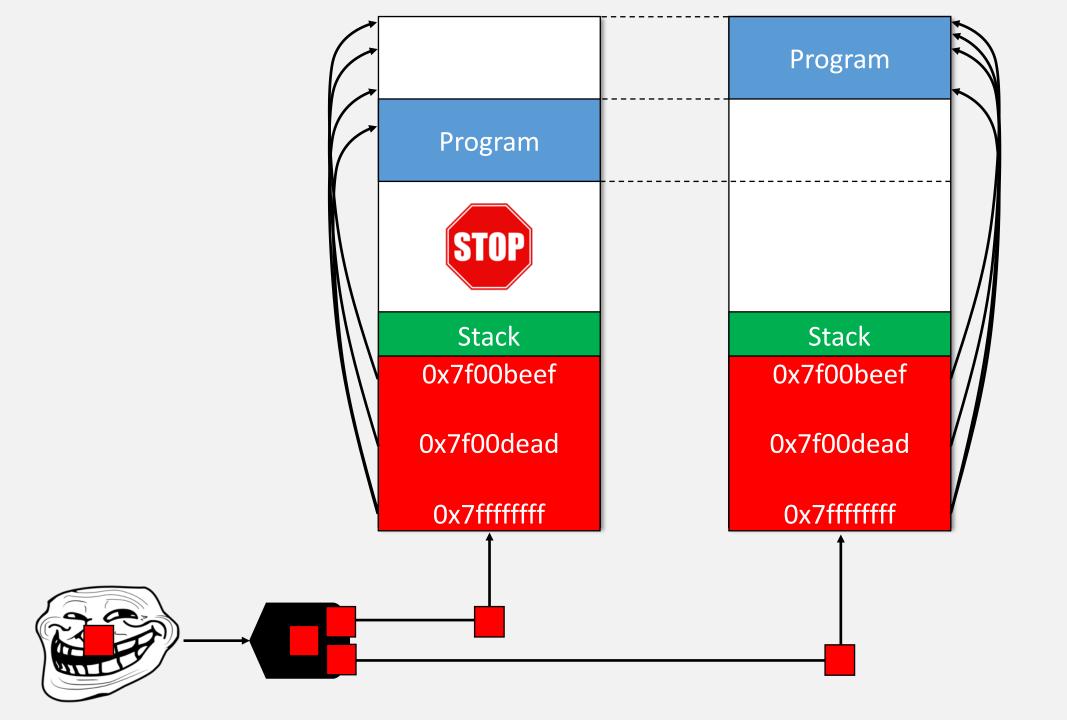




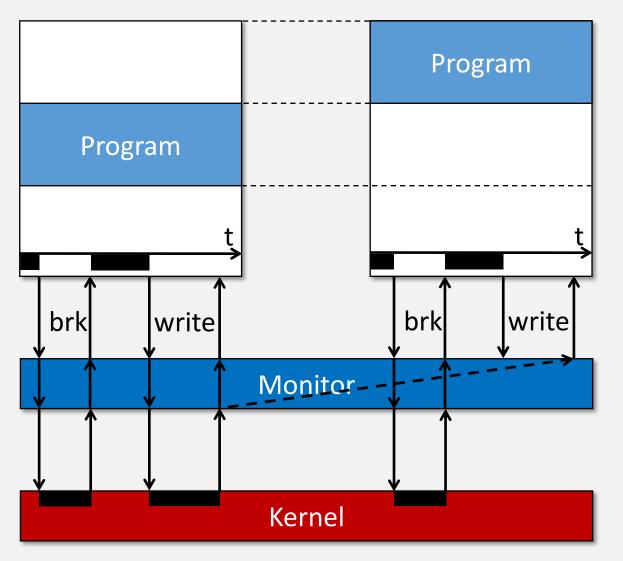


Memory Corruption Attacks





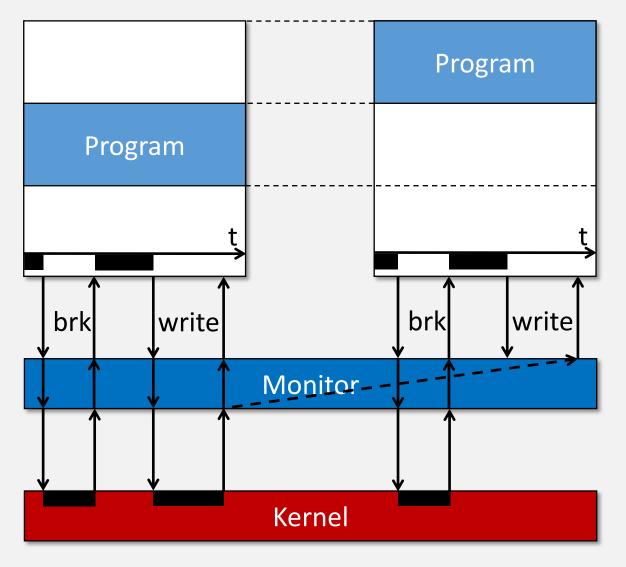
Multi-Variant Execution Environments (MVEEs)



In a nutshell:

- Run multiple program variants in parallel
- Variant system calls executed in lock-step
- Suspend them at every system call
- Compare system call numbers/arguments
- Master/Slave replication for I/O

Performance Considerations

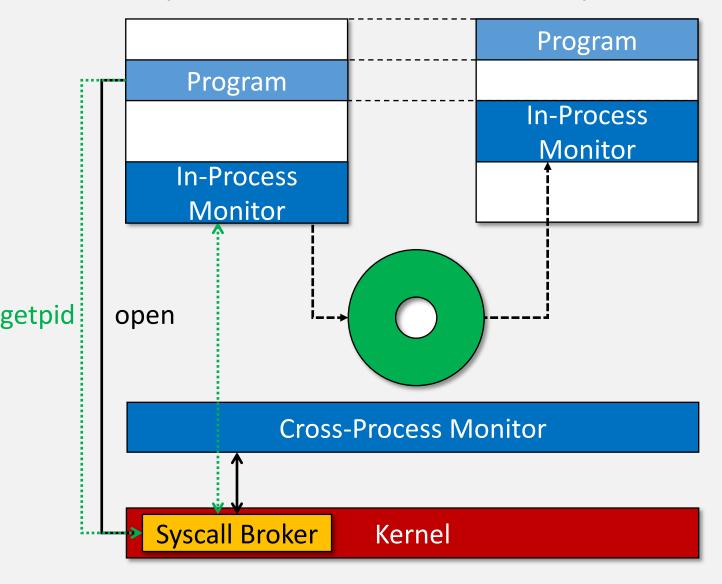


Programs can execute at **native speed** (assuming you have enough idle CPU cores and memory bandwidth)

The total system load does **not** increase by a factor of **n** (with n the number of variants)

BUT there are some problems!

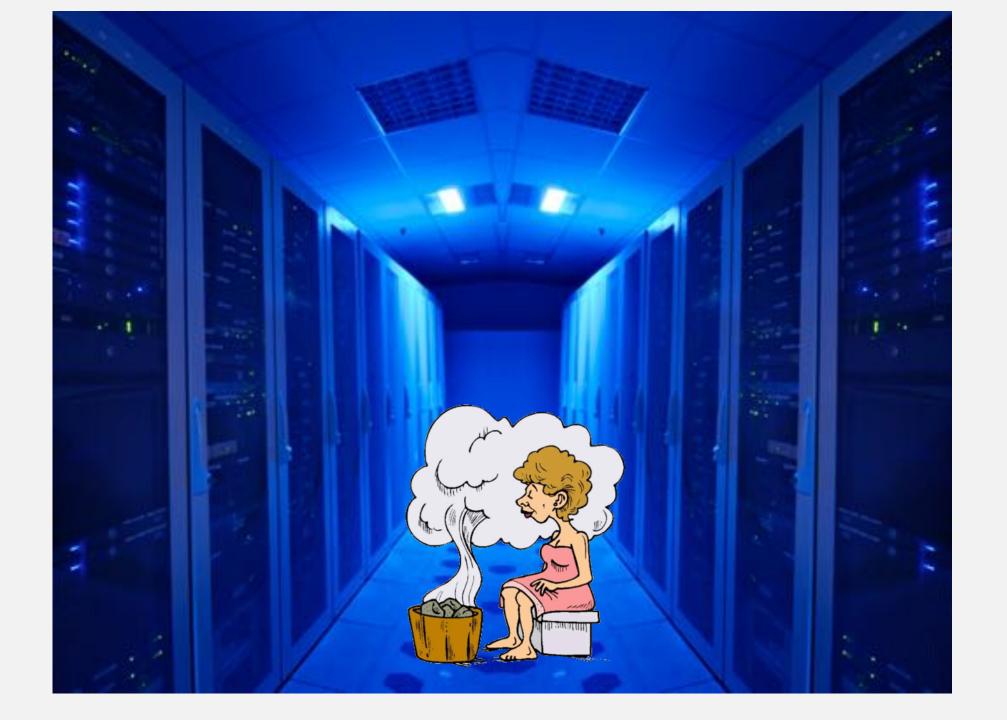
Slow System Call Interception



Split-Monitor Design:

- Handle security-sensitive system calls in Cross-Process Monitor (CP-MON)
- Handle non-sensitive system calls in In-Process Monitor (IP-MON)







Code Reuse

Program **Protects against:** return-to-libc (RILC) return-oriented programming (ROP) jump-oriented programming (JOP) Stack just-in-time code reuse (JIT-ROP) [*] counterfeit object-oriented programming (COOP) [*] 0x7f00beef 0x7f00dead 0x7ffffffff [*] Requires eXecute-only memory support

Program Stack 0x7f00beef 0x7f00dead 0x7ffffffff

Code Injection

```
\x48\x31\xd2\x48\xbb
\x2f\x2f\x62\x69\x6e
\x2f\x73\x68\x48\xc1
\xeb\x08\x53\x48\x89
\xe7\x50\x57\x48\x89
\xe6\xb0\x3b\x0f\x05
```

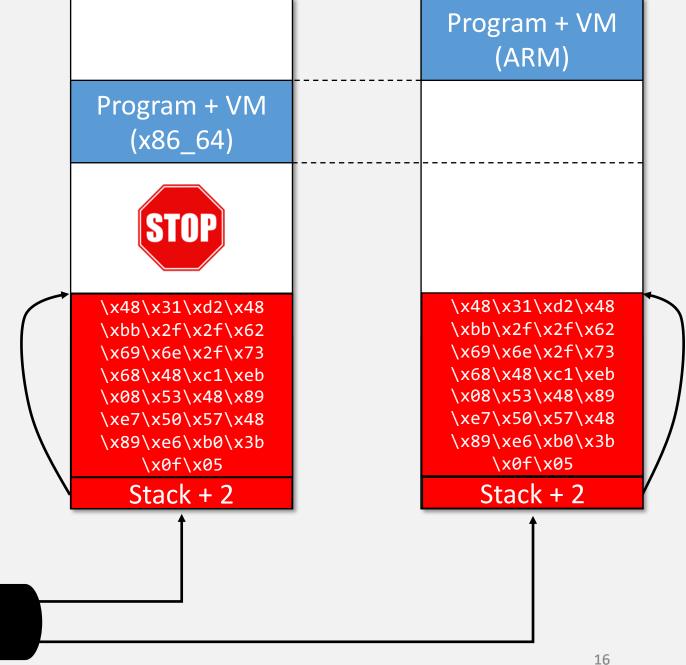


```
rdx, rdx
xor
        qword rbx, '//bin/sh'
mov
shr
        rbx, 0x8
push
       rbx
        rdi, rsp
mov
push
        rax
        rdi
push
       rsi, rsp
mov
        al, 0x3b
mov
syscall
```

execve("/bin/sh", ["/bin/sh"], NULL)

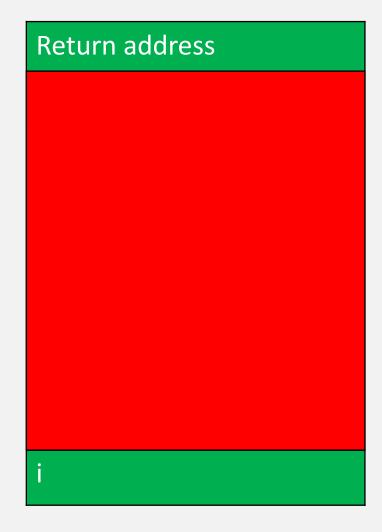
Code Injection Program Program $x48\x31\xd2\x48$ $x48\x31\xd2\x48$ $\xbb\x2f\x2f\x62$ $\xbb\x2f\x2f\x62$ $x69\x6e\x2f\x73$ $x69\x6e\x2f\x73$ $x68\x48\xc1\xeb$ $x68\x48\xc1\xeb$ \x08\x53\x48\x89 \x08\x53\x48\x89 \xe7\x50\x57\x48 $xe7\x50\x57\x48$ \x89\xe6\xb0\x3b $x89\xe6\xb0\x3b$ \x0f\x05 \x0f\x05 Stack + 2 Stack + 2 15

Code Injection



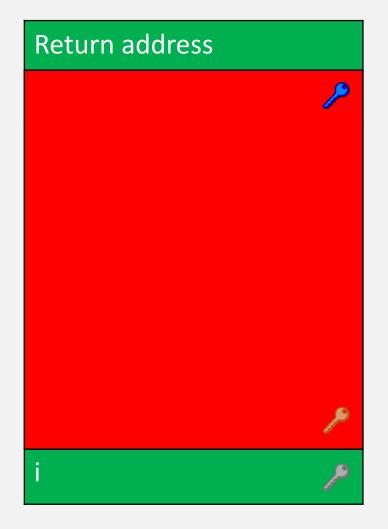
Non-Control Data Attacks

```
1: void ProcessConnection(connection* c) {
2: cred_t user;
3: char message[1024];
4: int i = 0;
5:
6: auth_user(&user, c);
7: while (!end_of_message(c))
    message[i++] = get_next_char(c);
9:
10: seteuid(user.user_id);
11: ExecuteRequest(message);
12:}
```



Non-Control Data Attacks

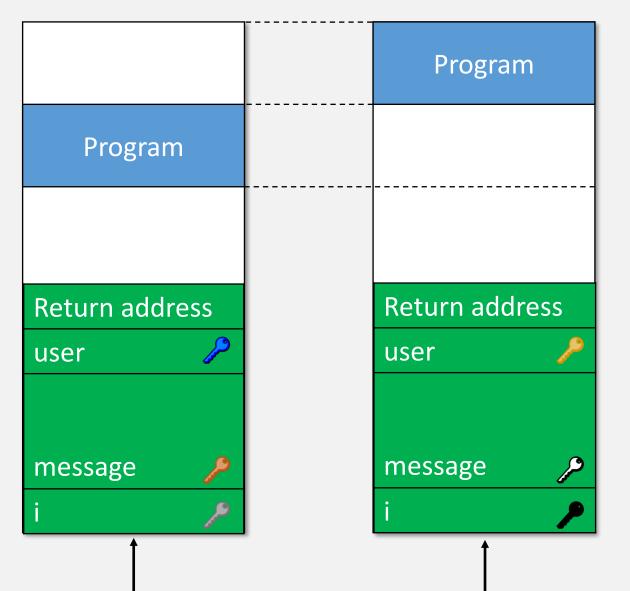
```
1: void ProcessConnection(connection* c) {
4: vorepetdes sconnection (connection* c) {
3: chen message[1024];
4: ithdr-neestyge (9024);
\frac{54}{4}: int i = 0;
6: auth_user(&user, c);
7: While (Leangle of emessage (c)) {
8: Whatetaring topported safety (j. P);
    message[tmp]=genchust(getanext_char(c), /);
10: i = encrypt(tmp + 1, P);
10: seteuid(user.user_id);
14: ExecuteRequest(message);
13: seteuid(decrypt(user.user_id, ));
14: ExecuteRequest(message);
15:}
```

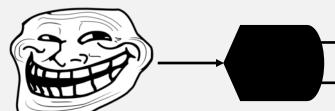


Non-Control Data Attacks

BONUS: Information Leakage Protection









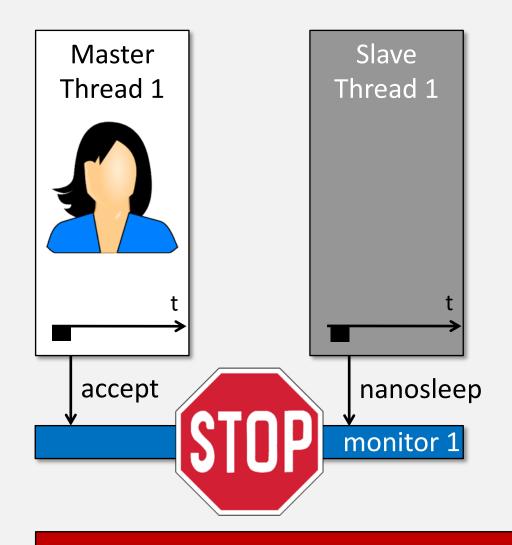
Multithreading

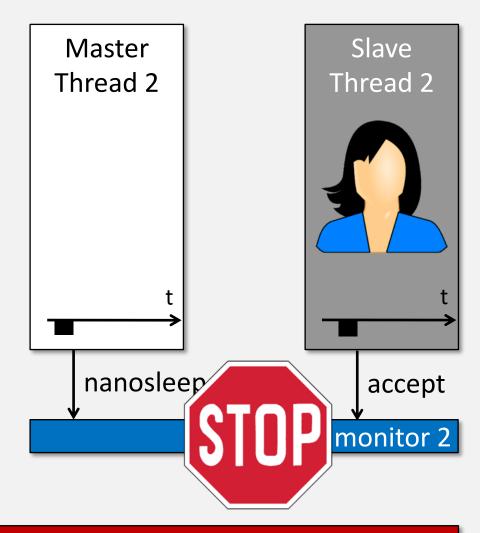
Master Thread 1 Master Thread 2 Slave Thread 1 Slave Thread 2

monitor 1 monitor 2

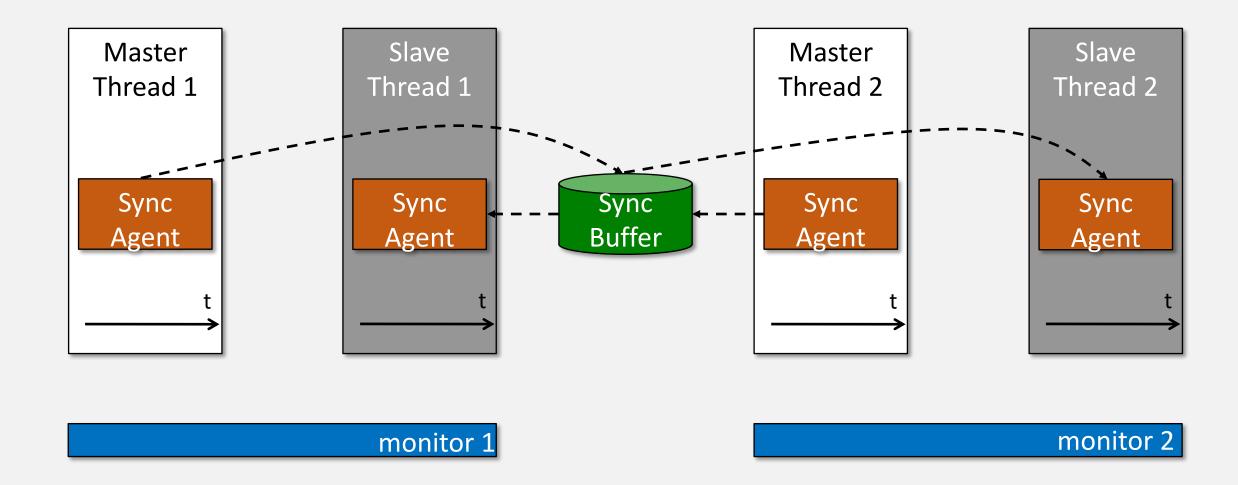
kernel

Multithreading

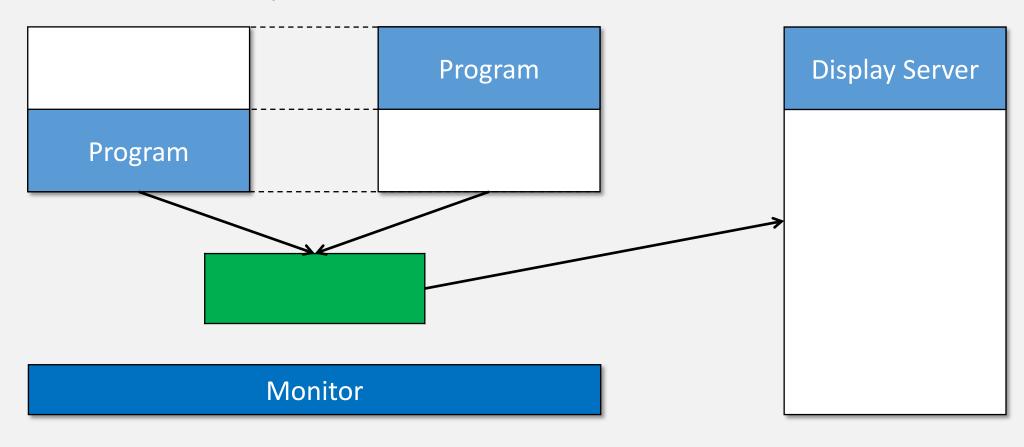




Multithreading



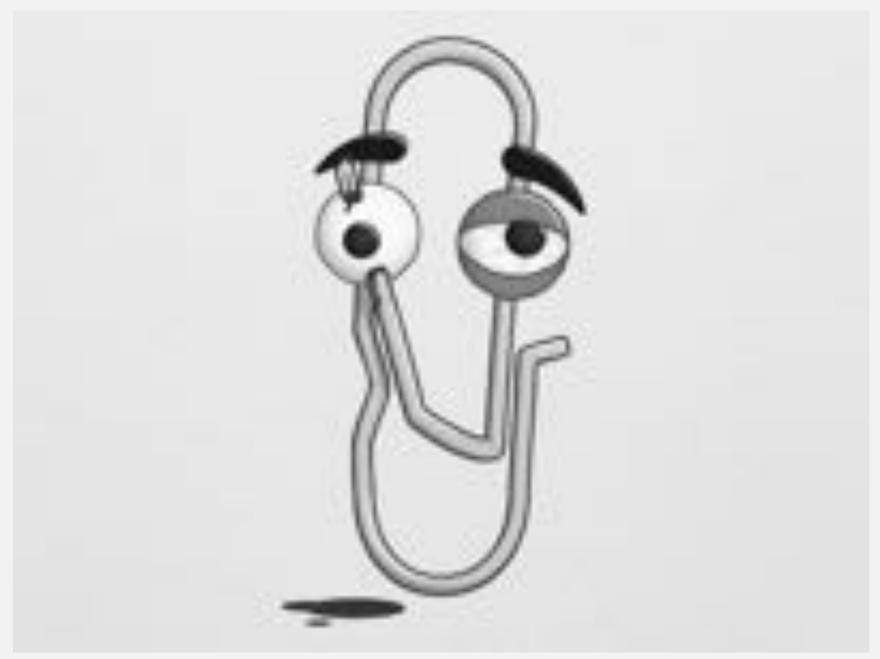
Shared Memory



Kernel

Other problems

- Data races
- vdso/vsyscall pages
- RDTSC/RDTSCP
- Address Dependence



https://github.com/stijn-volckaert/ReMon