

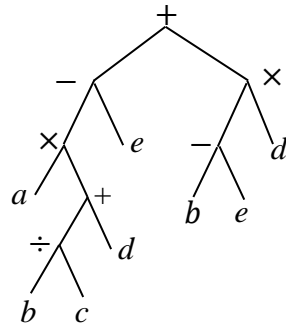
BMMS2633 ADVANCED DISCRETE MATHEMATICS

Academic year 2023/24

Session 202401

Q1. a) (i) Fully parenthesised form:
 $((a \times ((b \div c) + d)) - e) + ((b - e) \times d)$

(ii) Binary tree

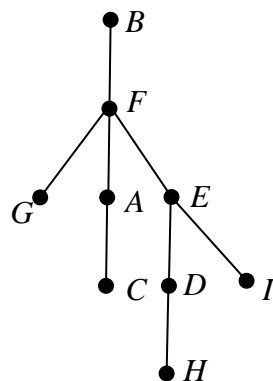


(iii) Preorder search : + - x a + ÷ b c d e x - b e d

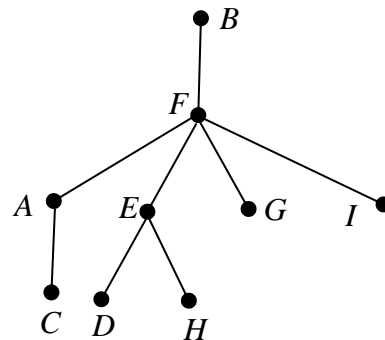
b) (i)

Vertex	List of adjacencies
<i>A</i>	<i>C F</i>
<i>B</i>	<i>F</i>
<i>C</i>	<i>A</i>
<i>D</i>	<i>E H</i>
<i>E</i>	<i>D F H I</i>
<i>F</i>	<i>A B E G I</i>
<i>G</i>	<i>F</i>
<i>H</i>	<i>D E</i>
<i>I</i>	<i>E F</i>

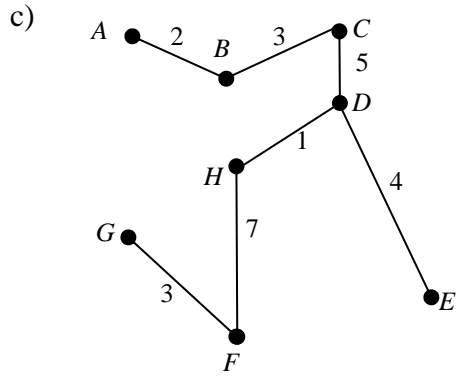
(ii) Depth-first search



Breadth-first search



BMMS2633 ADVANCED DISCRETE MATHEMATICS

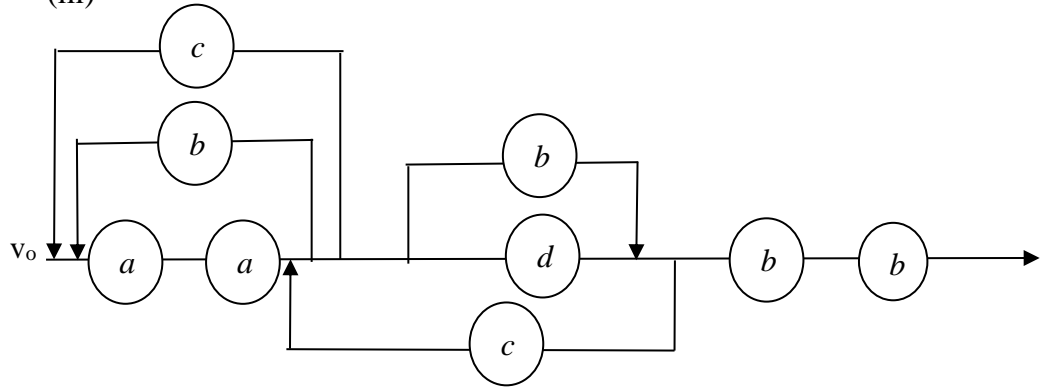


Minimum total weight = 25

- Q2. a) (i) $\langle v_0 \rangle ::= aa\langle v_1 \rangle$
 $\langle v_1 \rangle ::= b\langle v_0 \rangle \mid c\langle v_0 \rangle \mid b\langle v_2 \rangle \mid d\langle v_2 \rangle$
 $\langle v_2 \rangle ::= c\langle v_1 \rangle \mid bb$

(ii) *aacaabcb* is not a syntactically correct sentence.

(iii)



(iv) Regular expression = $a^2((b \vee c)a^2)^*((b \vee d)c((b \vee c)a^2)^*)^*(b \vee d)b^2$

- (b) (i) State transition table of f_{00101}

	f_{00101}
S_a	S_g
S_b	S_b
S_c	S_b
S_d	S_d
S_e	S_b
S_f	S_b
S_g	S_b

BMMS2633 ADVANCED DISCRETE MATHEMATICS

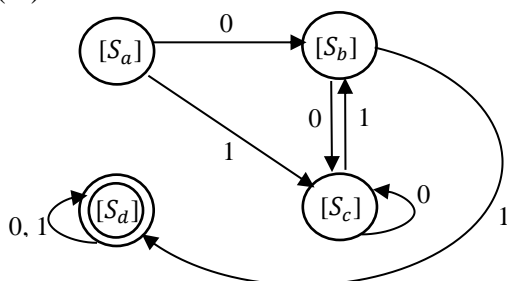
(ii)

	S_a	S_c	S_e	S_f	S_b	S_g	S_d
0	S_g	S_c	S_f	S_f	S_c	S_e	S_d
1	S_e	S_b	S_g	S_b	S_d	S_d	S_d

(iii)

	0	1
$[S_a]$	$[S_b]$	$[S_c]$
$[S_b]$	$[S_c]$	$[S_d]$
$[S_c]$	$[S_c]$	$[S_b]$
$[S_d]$	$[S_d]$	$[S_d]$

(iv)



(v)

$$f_{1001}(S_a) = S_b$$

$$f_{1001}([S_a]) = [S_b]$$

Input string 1001 is not accepted by M and M/R .

- Q3. a) (i) The description of $*$ on \mathbb{Z} is a valid definition of a binary operation.
- (ii) Description of $*$ on \mathbb{Z} is commutative.
- (iii) Description of $*$ on \mathbb{Z} is associative.
- (iv) Since $e = 0 \in \mathbb{Z}$, description of $*$ on \mathbb{Z} has an identity.
- b) (i) Subset A is not a subgroup of G .
- (ii) Subset Y is not a subgroup of G .

BMMS2633 ADVANCED DISCRETE MATHEMATICS

Q4 a) (i)

w	e(w)
00	0000
01	0110
10	1011
11	1101

(ii) The minimum distance of this (2, 4) encoding function $e_H = 2$

(iii)

\oplus	0000	0110	1011	1101
0001	0001	0111	1010	1100
0010	0010	0100	1001	1111
1000	1000	1110	0011	0101

(iv) (1) $d(0011) = 10$ (2) $d(1110) = 01$

b) (i)

Letter, (x_i)	Y	O	D	N	B	E
Number of occurrences	160	480	180	360	280	540
Probability, $P(x_i)$	0.08	0.24	0.09	0.18	0.14	0.27

(ii)

Letter, (x_i)	Y	O	D	N	B	E
Codeword, C_i	0001	10	0000	11	001	01

(iii) Average code length, $L(C) = 2.48$ bitsEntropy, $H(x) = 2.4507$

The efficiency of this code is 98.82%.