

# EE5903 RTS

## Chapter 6

### Real-Time Process Synchronization

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# Hardware Solutions

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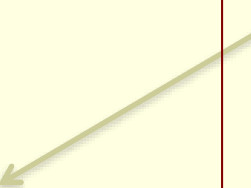
- Many systems provide hardware support for implementing the critical section code.
- Almost all solutions listed are based on idea of **locking** – That is, protecting critical regions via **locks**
- **Uniprocessors** – Disabling interrupts is one way – this means preemption is completely avoided;
- **Multiprocessors** – Not a scalable solution for multiprocessor systems;

# Hardware solutions

## General working logic:

```
do {  
    acquire lock;  
    enter CS;  
    release lock;  
    remainder section;  
} while(True);
```

Only one process will  
acquire the  
lock and others busy  
wait on that lock



Modern architectures provide atomic hardware instructions – non-interruptible;

- Test&Set()
- Compare&Swap() / Swap()

# Hardware Solution – Test&Set()

- **T&S()** is executed as an **atomic** operation; This does 2 things –
  - (i) Test the current value of the lock
  - (ii) Set the lock to **true** after acquiring

Both the above two steps are executed as though it is a **single instruction** – hence the name “atomic”;

- We say that a process **has acquired** the lock when it receives a **return value** as **FALSE** after T&S() execution;
- If the **return value** is TRUE then this means some other process has acquired the lock (and still has it);

# Hardware Solution – T&S() (cont'd)

Definition:

```
boolean test_and_set (boolean *target)
{
    boolean rv = *target;
    *target = true;
    return rv;
}
```

**Q 6.6:** Any problem with this Solution?



Process i

```
while (true) {
    while (test_and_set(&lock)) /* Entry section */
        ; // wait until lock becomes equal false
    // here lock=true

    /* critical section */
    lock = false; /* Release lock */
    /* remainder section */
}
```

# T&S() - Assembly codes

```
EnS: // Entry Section
      TandS Reg,Lock
      CMP Reg, #0
      JNE EnS
      RET
```

```
ExS: // Exit section
      MOV Lock, #0
      RET
```

```
main:
      Call EnS
      Execute CS
      Call ExS
```

# Hardware Solution – Swap()

```
void Swap(boolean *a, boolean *b)
{
    boolean temp = *a;
    *a = *b;
    *b = temp;
    return;
}
```

Q: How to make use of this swap() for controlling lock acquisition?

To be executed atomically!

## Implementation of Entry and Exit sections:

- We have a **shared boolean variable** – lock (all processes use this)
- Each process uses a **local boolean variable** - key

# Hardware Solution – Swap() (Cont'd...)

do{

key = True;

while(key == True)

(\*) Swap(&lock, &key);

Critical Section;

lock = False; /\* releasing the lock \*/

Remainder section;

} while(True);

Entry Section

Shared variable lock  
is initialized to FALSE  
at t=0

**Note:**  
lock – SHARED var  
key – LOCAL var

**Note:** After swapping, in the next iteration while() becomes False and hence it enters CS.



# Hardware Solutions - observations

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## ■ What do we observe in the two solutions?

- Mutual Exclusion property is satisfied
- Bounded Waiting time – Not satisfied;

Thus, a process **k** may be waiting but process **q** may be repeatedly executing CS indefinitely;

This can happen when its CPU time quantum does not expire even when it is in its remainder section and hence it can enter CS again!

*How do we achieve this bounded waiting time property in the above two solutions?*

## Hardware Solution – Bounded waiting time solution in T&S()

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**Q 6.7:** How do you bring in the **bounded waiting time feature** in the T&S() procedure? Show your implementation clearly.

Discuss clearly all possible cases on how a process acquires a lock, enters the CS, and releases it. Identify and convince yourself on the fairness of your implementation.

# Hardware Solution – Bounded waiting time solution in T&S()

do {

```
waiting[i] = True;
key = True; // local lock var
while(waiting[i] && key)
    key = T&S(&lock);
waiting[i] = False;
```

Critical Section;

```
j = (i+1)%n; // who is the next to enter CS?
while( ((j != i)&& !waiting[j]) // Find the next process waiting
        j=(j+1)%n
if(j == i) lock = False; // no one is waiting
else waiting[j] = False;
```

}while(True);

- First refer to the original T&S() soln

Entry Section

- waiting[] is boolean shared array;
- key is a local var to store the current lock status;

Exit Section

# Hardware Solution – Bounded waiting time solution in T&S()

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## Key observations:

- Initially lock = False;
- Until a currently running process sees that there is no one waiting, the lock is “locked” (i.e., held at the state “True” forever;
- Thus, if some process  $j$  is waiting (= True) then that processes' waiting status is changed to “False”, allowing it to enter CS section;
- If no one waits, the lock is released;
- More than one process waiting? – Choose as per a strategy that is in place! (priority driven, sequence #, largest load first, etc)
- A process once executed the CS is forced to leave thus enabling other waiting processes to enter CS, guaranteeing a bounded waiting time for others!

## Hardware Solution – Bounded waiting time solution in Swap()

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**Q 6.8:** How do you bring in the **bounded waiting time feature** in the **Swap()** procedure? Show your implementation clearly.

Discuss clearly all possible cases on how a process acquires a lock, enters the CS, and releases it. Identify and convince yourself on the fairness of your implementation.

# Mutex Locks

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- Previous solutions are complicated and generally inaccessible to application programmers; OS designers build software tools to solve critical section problem
- Simplest is **mutex** lock - Protect a critical section by:
  - First: **acquire()** a lock;
  - Then: **release()** the lock;
  - Boolean variable indicating if lock is available or not;

# Mutex Locks

## ■ **acquire()**

```
{  
    while(!available)  
        ; /* busy wait */  
    available = false;  
}
```

## **release()**

```
{  
    available = true;  
}
```

```
while(true) {  
    acquire()  
    critical section  
    release()  
    remainder section  
}
```

*Remarks:* Calls to **acquire()** and **release()** must be atomic; Usually implemented via hardware atomic instructions;

But this solution requires **busy waiting**;  
This lock therefore referred to as a **spinlock**;

## Avoiding busy waiting – *How?*

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Is there a solution which will not allow a process to *busy wait* if it is not in CS?

We can force a process to “sleep” so long as it cannot enter CS. If another process needs CS then it is “awaken” and it will be allowed to enter the CS.

Such a solution is achieved by the use of “*Semaphores*”