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1. (i) and (ii) are true; (iii) and (iv) are wrong.

P.j.(i) Sn ≤t, this mean that nth arrival occurred at some time T≤t

=> N(T) = N , $N(+) \ge N(T) = N$.

=> {Sn = \ } = { N(+) = n}

 $N(t) \ni N$, this mean there are m arrivals at time t, N(t) = m. $m \ni N$.

=> Sm < t, because man

=> Sn & Sm &t.

[NHD> n] = {Snet}

Combine 0 and Q. We canget [NKt)=nj={Sn \left};

- (ii) Since (i) is true, (ii) is true by taking complement.
- (iii). When, sock(Sn+1, N(t)=n, S{N(t)=n}, But Sn <t, So (iii) is false.
- (iv) Since the not true, it when Sn <t < Sn+1, N(t) = n, So (iv) is false

2. Because there are 5 courts, the time of #1 orrival is a exponentially districted. With meath:
$$E(x) = \frac{40 \text{ min}}{5} = 8 \text{ min}$$

If there are k other pairs waiting in queue, It the new pairs can get a court at k+ 1 thmes.

3. a:
$$P_{\tau}\{N(t)=n \mid S_n=T\}=P_{\tau}\{X_{n+1}>t-T\}=1-\int_0^{t-T} \lambda e^{-\lambda X} dX=e^{-\lambda (t-T)}$$

b.
$$P_{r}[N(t)=n] = \int_{0}^{t} P_{r}[N(t)=n|S_{n}=t] \cdot f_{S_{n}}(t) dt$$

$$= \int_{0}^{t} e^{-\lambda(t-\tau)} \frac{\lambda^{n} t^{n-1} \cdot e^{-\lambda t}}{(n-1)!} dt$$

$$= \frac{\lambda^{n} \cdot e^{-\lambda t}}{(n-1)!} \cdot \int_{0}^{t} t^{n-1} dt$$

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4. Prove that Geometric distribution has the memoryless property:

Pf: The memoryless defination:

$$P_r(X > t + x) = P_r(X > t) \cdot P_r(X > x)$$

$$LHS = \sum_{i=t+x+1}^{\infty} (1-p)^{i+p}$$

$$= \sum_{i=1}^{\infty} (1-p)^{i+p} - \sum_{i=1}^{t+x} (1-p)^{i+p}$$

$$= \frac{P}{1-(1-p)} - \frac{(1-p)^{i+x} P - P}{(1-p)-1}$$

=
$$[1 + [(1-p)^{t+x} - 1]]$$

= $(1-p)^{t+x}$

RHS=
$$P_r (x>t) \cdot P_r (x>x)$$

= $(1-p)^t \cdot (1-p)^x$
= $(1-p)^{x+t}$ = LHS.

J. Because In is a Bionmint random variable,

$$P_{1}(X_{n}=i) = \binom{n}{k} P_{n}^{i} (1-P_{n})^{n-i}$$

$$= \frac{n!}{i!(n-i)!} P_{n}^{i} (1-P_{n})^{n-i}$$

$$= \frac{n \cdot (n-i)!}{i!} P_{n}^{i} (1-P_{n})^{n-i}$$

$$= \frac{n \cdot (n-i) \cdot \dots \cdot (n-i+i)}{i!} \cdot (\frac{\lambda}{n})^{i} (1-\frac{\lambda}{n})^{n-i}$$

$$= \frac{n(n-i) \cdot \dots \cdot (n-i+i)}{n^{i}} \cdot \frac{\lambda^{i}}{i!} \cdot \frac{(1-\frac{\lambda}{n})^{n}}{(1-\frac{\lambda}{n})^{i}}$$

When
$$n \to \infty$$
, $n(n-1)-(n-i+1) = \frac{n \cdot n \cdot -n}{i} = n^{i}$

$$(1-\frac{\lambda}{n})^{i} \xrightarrow{n \to \infty} 1$$

$$(1-\frac{\lambda}{n})^{n} \xrightarrow{n \to \infty} e^{\lambda}.$$

$$2HS = \frac{n^{i}}{n^{i}} \cdot \frac{\lambda^{i}}{i!} \cdot \frac{e^{\lambda}}{1} = \frac{\lambda^{i}e^{\lambda}}{i!}$$

6. ten

P.f: Let a= 181, b= 181, c= |c|= 18081 If A and B are independ, we can get:

Pr(c) = Pr(A). Pr(B)

宁= 鲁·青.

pc= a.b.

since p is prime, at least one of a and b is divisible byp because $0 \le a, b \le p$, so at least one of a and b. is equal to our p

7. Let $Y=X-\frac{\alpha}{2}$, Because $P_Y(0 \le X \le \alpha)=1$, so $0 \le X \le \alpha$

2-0-3=Y < 0-2=2 / M < 2

Since X and T only differ by a constant, Var(X)=Var(Y)

 $V_{orr}(x) = V_{orr}(x) = E(x^2) - (E[x])^2 \leq E[x^2] \leq \frac{\alpha^2}{4}$

 δ . We want to find the distribution of Mn- $\frac{1}{2}$ byn.

$$F_{X}(x) = P_{T}\{M_{n} - \frac{1}{\lambda} \log n \leq x\}$$

=
$$P_r \{ M_n \leq x + \frac{1}{\lambda} lagn \}$$
.

Because $Mn = \max\{X_1, ..., X_n\}$. If (1) is real, that means there is i.i.d. then:

We know
$$\Pr\{X_n \leq x\} = 1 - e^{-\lambda x}$$

$$PHS = \prod_{n=1}^{\infty} \left[1 - e^{-\lambda \cdot (x + \frac{1}{\lambda} \log n)}\right] = \left[1 - e^{-\lambda (x + \frac{1}{\lambda} \log n)}\right]^n$$

$$= \left[1 - \frac{e^{-\lambda x}}{n}\right]^n$$

$$= e^{-\lambda x} (n \to \infty).$$