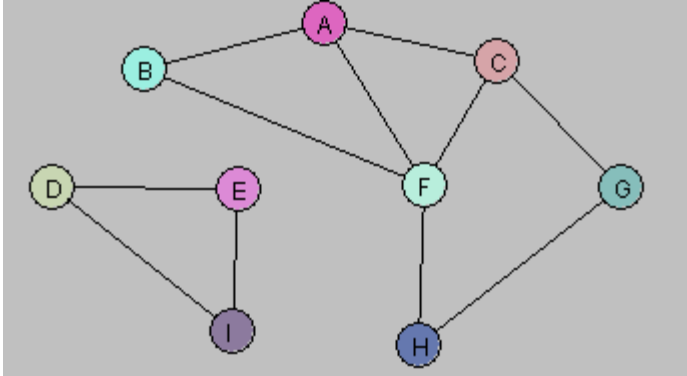
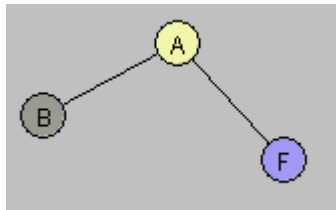


Lab 8

1. *Induced Graphs*. Answer questions about the graph $G = (V, E)$ displayed below.



- A. Let $U = \{A, B\}$. Draw $G[U]$.
- B. Let $W = \{A, C, G, F\}$. Draw $G[W]$.
- C. Let $Y = \{A, B, D, E\}$. Draw $G[Y]$.
- D. Consider the following subgraph H of G :



- Is there a subset X of the vertex set V so that $H = G[X]$? Explain.
- E. Find a way to partition the vertex set V into two subsets V_1, V_2 so that each of the induced graphs $G[V_1]$ and $G[V_2]$ is connected and $G = G[V_1] \cup G[V_2]$.

2. *Graph Implementation.* Use the BFS class to solve the following problems. Implement by implementing the unimplemented methods in the Graph class.

- ◆ Given two vertices, is there a path that joins them?
- ◆ Is the graph connected? If not, how many connected components does it have?
- ◆ Does the graph contain a cycle?

Hint: For the third problem, you may use the following Fact (which we will prove in tomorrow's class):

Fact: Suppose G is a graph and T is a spanning tree/forest for G . Then G has a cycle if and only if G has more edges than T .

3. *Graph Exercises.*

A. Suppose $G = (V, E)$ is a connected simple graph. Suppose V_1, V_2, \dots, V_k are disjoint subsets of V and that $V_1 \cup V_2 \cup \dots \cup V_k = V$. Show that there is an edge (x, y) in E such that for some $i \neq j$, x is in V_i and y is in V_j .

B. In class it was shown that a graph $G = (V, E)$ is connected whenever the following is true,

$$\epsilon > \binom{\nu - 1}{2} \quad (*)$$

where ν is the number of vertices and ϵ is the number of edges. Is the following true or false?

Every connected graph satisfies the inequality ().*

Prove your answer.

C. Suppose G is a graph with two vertices. What is the minimum number of edges it must have in order to be a connected graph? Suppose instead G has three vertices; what is the minimum number of edges it must have in order to be connected? Fill in the blank with a reasonable conjecture:

If G has n vertices, G must have at least _____ edges in order to be connected.

4. *IsPrime Problem, Revisited.* The goal of this exercise is to devise a feasible algorithm that decides whether an input integer is prime. The key fact that you will make use of is the following:

Fact: There is a function f , which runs in $O(\log n)$ (that is, $O(\text{length}(n))$), such that for any odd positive integer n and any a chosen randomly in $[1, n - 1]$, if $f(a, n) = 1$, then n is composite, but if $f(a, n) = 0$, n is “probably” prime, but is in fact composite with probability $< 1/2$.

A first try at such an algorithm would be:

Algorithm **FirstTry**:

Input: A positive integer n

Output: TRUE if n is prime, FALSE if n is composite

```
if  $n \% 2 = 0$  return FALSE
 $a \leftarrow$  random number in  $[1, n-1]$ 
if  $f(a, n) = 1$ 
    return FALSE
return TRUE
```

Notice that **FirstTry** runs in $O(\log n)$. It also produces a correct result more than half the time.

What could be done to improve the degree of correctness of **FirstTry** but still preserve a reasonably good running time? Explain.