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Homework 2

1a. Prove that "Implication Introduction" (the opposite of Implication Elimination) is a sound rule of inference (ROI) using a truth table. If you have a Horn clause, with 1 positive literal and n-1 negative literals, like $(\neg X \lor Z \lor \neg Y)$, you can transform it into a conjunctive rule by collecting the negative literals as positive antecedents, e.g. $X \land Y \Rightarrow Z$. It is sufficient to prove this for n-1=2 antecedents. (In fact, this is a truth-preserving operation, hence sound.)

Recall: A sound rule of inference only generates new sentences that are entailed.

X	Y	\mathbf{Z}	$\neg X$	$\neg Y$	$\neg X \vee \neg Y \vee Z$	$X{\wedge}Y$	$X \land Y \Rightarrow Z$
F	F	F	T	Τ	${f T}$	F	${f T}$
\mathbf{F}	\mathbf{F}	${\rm T}$	T	Τ	${f T}$	\mathbf{F}	${f T}$
\mathbf{F}	\mathbf{T}	F	T	\mathbf{F}	${f T}$	\mathbf{F}	${f T}$
\mathbf{F}	\mathbf{T}	\mathbf{T}	T	\mathbf{F}	${f T}$	\mathbf{F}	${f T}$
${ m T}$	F	F	F	\mathbf{T}	${f T}$	\mathbf{F}	${f T}$
${ m T}$	F	${\rm T}$	F	\mathbf{T}	${f T}$	\mathbf{F}	${f T}$
\mathbf{T}	\mathbf{T}	\mathbf{F}	F	\mathbf{F}	\mathbf{F}	${ m T}$	\mathbf{F}
${\rm T}$	\mathbf{T}	\mathbf{T}	F	\mathbf{F}	${f T}$	${ m T}$	${f T}$

Table 1: Truth Table for 1a

Since all models that satisfy the premise $\neg X \lor \neg Y \lor Z$ also satisfy the conclusion $X \land Y \Rightarrow Z$, we prove that the conclusion is entailed. Thus, implication introduction is a sound rule of inference. In fact, since the premise and the conclusion are satisfied by the exact same models, implication introduction is a truth-preserving rule of inference.

1b. Prove that $(A \land B \Rightarrow C \land D) \vdash (A \land B \Rightarrow C)$ ("conjunctive rule splitting") is a sound rule-of-inference using a truth table.

Recall: A sound rule of inference only generates new sentences that are entailed.

A	В	\mathbf{C}	D	$A \wedge B$	$C{\wedge}D$	$A{\wedge}B{\Rightarrow}C{\wedge}D$	$A{\wedge}B{\Rightarrow}C$
F	F	F	F	F	F	\mathbf{T}	\mathbf{T}
\mathbf{F}	\mathbf{F}	\mathbf{F}	${\rm T}$	F	\mathbf{F}	${f T}$	${f T}$
\mathbf{F}	F	\mathbf{T}	F	F	\mathbf{F}	${f T}$	${f T}$
\mathbf{F}	F	Τ	T	F	${ m T}$	${f T}$	${f T}$
\mathbf{F}	${\rm T}$	F	F	F	\mathbf{F}	${f T}$	${f T}$
\mathbf{F}	Τ	F	T	F	\mathbf{F}	${f T}$	${f T}$
\mathbf{F}	\mathbf{T}	\mathbf{T}	\mathbf{F}	F	\mathbf{F}	${f T}$	${f T}$
\mathbf{F}	\mathbf{T}	\mathbf{T}	${ m T}$	F	${ m T}$	${f T}$	${f T}$
\mathbf{T}	\mathbf{F}	\mathbf{F}	\mathbf{F}	F	\mathbf{F}	${f T}$	${f T}$
\mathbf{T}	\mathbf{F}	\mathbf{F}	${ m T}$	F	\mathbf{F}	${f T}$	${f T}$
\mathbf{T}	\mathbf{F}	\mathbf{T}	\mathbf{F}	F	\mathbf{F}	${f T}$	${f T}$
\mathbf{T}	\mathbf{F}	\mathbf{T}	${ m T}$	F	${ m T}$	${f T}$	${f T}$
\mathbf{T}	\mathbf{T}	\mathbf{F}	\mathbf{F}	Т	\mathbf{F}	\mathbf{F}	\mathbf{F}
T	Τ	F	T	T	\mathbf{F}	\mathbf{F}	\mathbf{F}
\mathbf{T}	\mathbf{T}	${\rm T}$	\mathbf{F}	Т	\mathbf{F}	\mathbf{F}	${ m T}$
${ m T}$	${\bf T}$	\mathbf{T}	\mathbf{T}	Т	${ m T}$	${f T}$	${f T}$

Table 2: Truth Table for 1b

Since all models that satisfy the premise $A \land B \Rightarrow C \land D$ also satisfy the conclusion $A \land B \Rightarrow C$, we prove that the conclusion is entailed. Thus, conjunctive rule splitting is a sound rule-of-inference.

1c. Also prove $(A \land B \Rightarrow C \land D) \models (A \land B \Rightarrow C)$ using Natural Deduction. (hint: use 1a above)

Premise

1. $A \land B \Rightarrow C \land D$

Derivations

- 2. $\neg(A \land B) \lor C \land D$ [Implication Elimination, 1]
- 3. $(\neg A \lor \neg B) \lor C \land D$ [DeMorgan's Law, 2]
- 4. $(\neg A \lor \neg B \lor C) \land (\neg A \lor \neg B \lor D)$ [Distributivity, 3]
- 5. $(\neg A \lor \neg B \lor C)$ [And Elimination, 4]
- 6. $A \land B \Rightarrow C$ [Implication Introduction, 5]

1d. Also prove $(A \land B \Rightarrow C \land D) \models (A \land B \Rightarrow C)$ using Resolution.

Convert premise to Conjunctive-Normal Form Clauses

- 1. $A \land B \Rightarrow C \land D$
- 2. $\neg(A \land B) \lor C \land D$ [Implication Elimination, 1]
- 3. $(\neg A \lor \neg B) \lor C \land D$ [DeMorgan's Law, 2]
- 4. $(\neg A \lor \neg B \lor C) \land (\neg A \lor \neg B \lor D)$ [Distributivity, 3]
- 5. $\neg A \lor \neg B \lor C$ [And Elimination, 4] (CNF)
- 6. $\neg A \lor \neg B \lor D$ [And Elimination, 4] (CNF)

Negate The Query:

- $\neg (A \land B \Rightarrow C)$
- $\neg(\neg(A \land B) \lor C)$ [Implication Elimination]
- $\neg\neg(A \land B) \land \neg C$ [DeMorgan's Laws]
- $A \land B \land \neg C$ [Double Negation Elimination]
- A, B, $\neg C$ [And Elimination]

Apply resolution refutation by adding the negated query into the premises:

Premises (All in CNF):

- 1. $\neg A \lor \neg B \lor C$
- 2. $\neg A \lor \neg B \lor D$
- 3. A
- 4. B
- 5. ¬C

Derivations

- 6. $\neg B \lor C$ [Resolution on 1 and 3]
- 7. C [Resolution on 4 and 6]
- 8. \emptyset [Resolution on 5 and 7]

2. Sammy's Sport Shop

2a. Using these propositional symbols, write a propositional knowledge base (sammy.kb) that captures the knowledge in this domain (i.e. implications of what different observations or labels mean, as well as constraints inherent in this problem, such as that all boxes have different contents). Do it in a complete and general way, writing down all the rules and constraints, not just the ones needed to make the specific inference about the middle box. Do not include derived knowledge that depends on the particular labeling of this instance shown above; stick to what is stated in the problem description above. Your KB should be general enough to reason about any alternative scenario, not just the one given above (e.g. with different observations and labels and box contents).

Label Implications	20. $\neg C1W \land \neg C1B \Leftrightarrow C1Y$	Contains Implications
1. $L1W \Rightarrow \neg C1W$	21. $\neg C1Y \land \neg C1W \Leftrightarrow C1B$	40. C1W⇔ ¬C2W∧¬C3W
2. $L1Y \Rightarrow \neg C1Y$	22. $\neg C2Y \land \neg C2B \Leftrightarrow C2W$	41. C1Y⇔ ¬C2Y∧¬C3Y
3. $L1B \Rightarrow \neg C1B$	23. $\neg C2W \land \neg C2B \Leftrightarrow C2Y$	42. C1B⇔ ¬C2B∧¬C3B
4. $L2W \Rightarrow \neg C2W$	24. $\neg C2Y \land \neg C2W \Leftrightarrow C2B$	43. C2W⇔ ¬C1W∧¬C3W
5. $L2Y \Rightarrow \neg C2Y$	25. ¬C3Y∧¬C3B⇔C3W	
6. L2B⇒ ¬C2B	26. ¬C3W∧¬C3B⇔C3Y	44. $C2Y \Leftrightarrow \neg C1Y \land \neg C3Y$
7. L3W⇒ ¬C3W	27. ¬C3Y∧¬C3W⇔C3B	45. $C2B \Leftrightarrow \neg C1B \land \neg C3B$
8. L3Y⇒ ¬C3Y	Observation Implications	46. $C3W \Leftrightarrow \neg C2W \land \neg C1W$
9. L3B⇒ ¬C3B	28. O1W⇒C1W∨C1B	47. $C3Y \Leftrightarrow \neg C2Y \land \neg C1Y$
Color Implications	29. O1Y⇒C1Y∨C1B	48. C3B $\Leftrightarrow \neg$ C2B $\land \neg$ C1B
10. L1W $\Leftrightarrow \neg$ L2W $\land \neg$ L3W	30. $O2W \Rightarrow C2W \lor C2B$	49. C1W⇒O1W
11. $L1Y \Leftrightarrow \neg L2Y \land \neg L3Y$	31. O2Y⇒C2Y∨C2B	50. C1Y⇒O1Y
12. $L1B \Leftrightarrow \neg L2B \land \neg L3B$	32. O3W⇒C3W∨C3B	51. C1B⇒O1W∨O1Y
13. L2W $\Leftrightarrow \neg L1W \land \neg L3W$	33. O3Y⇒C3Y∨C3B	52. C2W⇒O2W
14. $L2Y \Leftrightarrow \neg L1Y \land \neg L3Y$	34. $O1W \Rightarrow \neg C1Y$	
15. L2B⇔ ¬L1B∧¬L3B	35. $O1Y \Rightarrow \neg C1W$	53. C2Y⇒O2Y
16. $L3W \Leftrightarrow \neg L2W \land \neg L1W$	36. $O2W \Rightarrow \neg C2Y$	54. C2B⇒O2W∨O2Y
17. $L3Y \Leftrightarrow \neg L2Y \land \neg L1Y$	37. $O2Y \Rightarrow \neg C2W$	55. C3W⇒O3W
18. L3B⇔ ¬L2B∧¬L1B	38. O3W⇒ ¬C3Y	56. C3Y⇒O3Y
19. ¬C1Y∧¬C1B⇔C1W	39. O3Y⇒ ¬C3W	57. C3B⇒O3W∨O3Y

2b. Prove that box 2 must contain white balls (C2W) using Natural Deduction.

Initial Facts

- 58. L1W
- 59. L2Y
- 60. L3B
- 61. O1Y
- 62. O2W
- 63. O3Y

Derivations

- 64. ¬C3B [Modus Ponens with rules 60 and 9]
- 65. C3Y∨C3B [Modus Ponens with rules 33 and 63]
- 66. C3Y [Resolution with rules 64 and 65]
- 67. ¬C2Y∧¬C1Y [Modus Ponens with rules 66 and 47]
- 68. ¬C2Y[And Elimination on rule 67]
- 69. ¬C1Y[And Elimination on rule 67]
- 70. ¬C1W [Modus Ponens with rules 1 and 58]
- 71. \neg C1Y $\land \neg$ C1W [And introduction with rules 69 and 70]
- 72. C1B [Modus Ponens with rules 71 21]
- 73. $\neg C2B \land \neg C3B$ [Modus Ponens with rules 72 and 42]
- 74. ¬C2B [And Elimination on rule 73]
- 75. $\neg C2Y \land \neg C2B$ [And Introduction on rules 68 and 74]
- 76. C2W [Modus Ponens on rules 75 and 22]

2c. Convert your KB to CNF

Label Implications	23. C2WVC2YVC2B	Contains Implications
1. \neg L1W $\lor \neg$ C1W	24. C3W\C3Y\C3B	46. $C1W \lor C2W \lor C3W$
2. $\neg L1Y \lor \neg C1Y$	25. $\neg C1W \lor \neg C1Y$	47. C1Y\C2Y\C3Y
3. ¬L1B∨¬C1B	26. ¬C1W∨¬C1B	48. C1B∨C2B∨C3B
4. $\neg L2W \lor \neg C2W$	27. ¬C1Y∨¬C1B	49. ¬C1W∨¬C2W
5. $\neg L2Y \lor \neg C2Y$	28. $\neg C2W \lor \neg C2Y$	50. ¬C1W∨¬C3W
6. ¬L2B∨¬C2B	29. ¬C2W∨¬C2B	
7. ¬L3W∨¬C3W	$30. \neg C2Y \lor \neg C2B$	51. ¬C1Y∨¬C2Y
8. ¬L3Y∨¬C3Y	31. ¬C3W∨¬C3Y	52. ¬C1Y∨¬C3Y
9. ¬L3B∨¬C3B	32. ¬C3W∨¬C3B	53. ¬C1B∨¬C2B
Color Implications	33. ¬C3Y∨¬C3B	54. ¬C1B∨¬C3B
10. $L1W \lor L2W \lor L3W$	Observation Implications	55. $\neg C2W \lor \neg C3W$
11. $L1Y \lor L2Y \lor L3Y$	34. ¬O1W∨C1W∨C1B	56. ¬C2Y∨¬C3Y
12. $L1B\lor L2B\lor L3B$	35. ¬O1Y∨C1Y∨C1B	57. ¬C2B∨¬C3B
13. ¬L1W∨¬L2W	36. $\neg O2W \lor C2W \lor C2B$	58. ¬C1W∨O1W
14. ¬L1W∨¬L3W	37. ¬O2Y∨C2Y∨C2B	59. ¬C1Y∨O1Y
15. ¬L1Y∨¬L2Y	38. ¬O3W∨C3W∨C3B	60. ¬C1B∨O1Y∨O1W
16. ¬L1Y∨¬L3Y	39. ¬O3Y∨C3Y∨C3B	
17. ¬L1B∨¬L2B	40. ¬O1W∨¬C1Y	61. ¬C2W∨O2W
18. ¬L1B∨¬L3B	41. ¬O1Y∨¬C1W	62. $\neg C2Y \lor O2Y$
19. ¬L2W∨¬L3W	42. ¬O2W∨¬C2Y	63. $\neg C2B \lor O2Y \lor O2W$
20. ¬L2Y∨¬L3Y	43. ¬O2Y∨¬C2W	64. ¬C3W∨O3W
21. ¬L2B∨¬L3B	44. ¬O3W∨¬C3Y	65. ¬C3Y∨O3Y
22. C1W\C1Y\C1B	45. ¬O3Y∨¬C3W	66. ¬C3B∨O3Y∨O3W

2d. Prove C2W using Resolution

Premises

- 67. L1W
- 68. L2Y
- 69. L3B
- 70. O1Y
- 71. O2W
- 72. O3Y

Derivations

- 73. ¬C2W [Negated query]
- 74. C2YVC2B [Resolution between 73 and 23]
- 75. \neg L2Y \lor C2B [Resolution with 74 and 5]
- 76. C2B [Resolution with 75 and 68]
- 77. ¬C1B [Resolution with 76 and 53]
- 78. C1WVC1Y [Resolution with 77 and 22]
- 79. \neg L1W \lor C1Y [Resolution with 78 and 1]
- 80. C1Y [Resolution with 79 and 67]
- 81. \neg C3Y [Resolution with 80 and 52]
- 82. ¬O3Y∨C3B [Resolution with 81 and 39]
- 83. C3B [Resolution with 82 and 72]
- 84. ¬L3B [Resolution with 83 and 9]
- 85. \emptyset [Resolution with 84 and 69]

Since we were able to derive the empty clause \emptyset from the premise $\neg C2W$, we have demonstrated that the clause $\neg C2W$ causes a contradiction in Sammy.kb, which consequently proves C2W.

3. Do Forward Chaining for the CanGetToWork KB below. You don't need to follow the formal FC algorithm (with agenda/queue and counts array). Just indicate which rules are triggered (in any order), and keep going until all consequences are generated. Show the final list of all **Inferred** propositions at the end. Is CanGetToWork among them?

```
KB = \{ a. CanBikeToWork \rightarrow CanGetToWork \}
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- b. $CanDriveToWork \rightarrow CanGetToWork$
- c. $CanWalkToWork \rightarrow CanGetToWork$
- d. HaveBike \land WorkCloseToHome \land Sunny \rightarrow CanBikeToWork
- e. HaveMountainBike \rightarrow HaveBike
- f. HaveTenSpeed \rightarrow HaveBike
- g. OwnCar \rightarrow CanDriveToWork
- h. OwnCar \rightarrow MustGetAnnualInspection
- i. OwnCar \rightarrow MustHaveValidLicense
- j. Can
RentCar \rightarrow CanDriveToWork
- k. HaveMoney \wedge CarRentalOpen \rightarrow CanRentCar
- l. Hertz Open \rightarrow Car
Rental Open
- m. AvisOpen \rightarrow CarRentalOpen
- n. EnterpriseOpen \rightarrow CarRentalOpen
- o. CarRentalOpen \rightarrow IsNotAHoliday
- p. HaveMoney \wedge TaxiAvailable \rightarrow CanDriveToWork
- q. Sunny \land WorkCloseToHome \rightarrow CanWalkToWork
- r. Have Umbrella \land Work Close To
Home \rightarrow CanWalk To
Work
- s. Sunny \rightarrow StreetsDry }

Facts:{Rainy, HaveMoutainBike, EnjoyPlayingSoccer, WorkForUniversity, WorkCloseToHome, Have-Money, HertzClosed, AvisOpen, McDonaldsOpen}

Note: For the purpose of this problem, I will discard all of the literals in **Facts** that do not appear in the knowledge base. I was given permission to do this in Campuswire Post #55.

- 1. Initialize Queue and Inferred data structures.
 - Queue = {HaveMountainBike, WorkCloseToHome, HaveMoney, AvisOpen}
 - Inferred $= \{\}$
- 2. Pop HaveMountainBike. HaveMountainBike triggers **rule e**. Add HaveMountainBike to **Inferred** and HaveBike to the **Queue**.
 - Queue = {WorkCloseToHome, HaveMoney, AvisOpen, HaveBike}
 - **Inferred** = {HaveMountainBike}
- 3. Pop WorkCloseToHome. WorkCloseToHome and our current inference list does not trigger any other rules.
 - Queue = {HaveMoney, AvisOpen, HaveBike}
 - **Inferred** = {HaveMountainBike, WorkCloseToHome}

- 4. Pop HaveMoney. HaveMoney and our current inference list does not trigger any other rules.
 - Queue = {AvisOpen, HaveBike}
 - **Inferred** = {HaveMountainBike, WorkCloseToHome, HaveMoney}
- 5. Pop AvisOpen. AvisOpen and HaveMoney trigger **rule k**. Add AvisOpen to **Inferred** and CanRentCar to the **Queue**.
 - Queue = {HaveBike, CanRentCar}
 - Inferred = {HaveMountainBike, WorkCloseToHome, HaveMoney, AvisOpen}
- 6. Pop HaveBike. HaveBike and our current inference list does not trigger any other rules.
 - Queue = {CanRentCar}
 - Inferred = {HaveMountainBike, WorkCloseToHome, HaveMoney, HaveBike}
- 7. Pop CanRentCar. CanRentCar triggers **rule j**. Add CanRentCar to **Inferred** and CanDrive-ToWork to the **Queue**.
 - Queue = {CanDriveToWork}
 - **Inferred** = {HaveMountainBike, WorkCloseToHome, HaveMoney, HaveBike, CanRent-Car}
- 8. Pop CanDriveToWork. CanDriveToWork triggers **rule b**. Add CanDriveToWork to **Inferred** and CanGetToWork to the **Queue**.
 - Queue = {CanGetToWork}
 - Inferred = {HaveMountainBike, WorkCloseToHome, HaveMoney, HaveBike, CanRentCar, CanDriveToWork}
- 9. Pop CanGetToWork. CanGetToWork is our query, so we are done!

4. Do **Backward Chaining** for the CanGetToWork KB. In this case, you should follow the BC algorithm closely (the pseudocode for the propositional version of Back-chaining is given in the lecture slides).

Important: when you pop a subgoal (proposition) from the goal stack, you should systematically go through all rules that can be used to prove it IN THE ORDER THEY APPEAR IN THE KB. In some cases, this will lead to back-tracking, which you should show.

Also, the sequence of results depends on order in which antecedents are pushed onto the stack. If you have a rule like $A \land B \Rightarrow C$, and you pop C off the stack, push the antecedents in reverse order, so B goes in first, then A; in the next iteration, A would be the next subgoal popped off the stack.

Note: In order for the tracing to terminate with the conclusion, we must push the following known facts into the KB:

- {t. HaveMountainBike, u. WorkCloseToHome, v. HaveMoney, w. AvisOpen}
 - 1. Initialize the goal stack and push the query CanGetToWork.
 - Goal Stack = {CanGetToWork}
 - 2. Pop CanGetToWork. We branch by pushing the antecedents in rule a (Could have used a, b, or c)
 - Goal Stack = {CanBikeToWork}
 - 3. Pop CanBikeToWork. We push the antecedents in rule d
 - Goal Stack = {HaveBike, WorkCloseToHome, Sunny}
 - 4. Pop HaveBike. Push the antecedents in rule e (could have used e or f)
 - Goal Stack = {WorkCloseToHome, Sunny, HaveMountainBike}
 - 5. Pop WorkCloseToHome. It is a known fact, so we don't push any antecedents.
 - Goal Stack = {Sunny, HaveMountainBike}
 - 6. Pop Sunny. Sunny not provable, so we backtrack our choice in step 4 and push the antecedents of rule f instead.
 - Goal Stack = {WorkCloseToHome, Sunny, HaveTenSpeed}
 - 7. Pop WorkCloseToHome. It is a known fact (rule u), so we don't push any antecedents.
 - Goal Stack = {Sunny, HaveTenSpeed}
 - 8. Pop Sunny. Sunny is not provable and since we have no other antecedents to push for HaveBike, we backtrack to **Step 2** and pick a different rule for the antecedents of CanGetToWork. We select rule b and push CanDriveToWork to the Goal Stack.
 - Goal Stack = {CanDriveToWork}

- 9. Pop CanDriveToWork. We branch by pushing the antecedents in rule g (could have used g, j, p, or q)
 - Goal Stack = {OwnCar}
- 10. Pop OwnCar. OwnCar is not provable, so we backtrack to **Step 9** and push the antecedents in rule j instead.
 - Goal Stack = {CanRentCar}
- 11. Pop CanRentCar. Push the antecedents in rule k.
 - Goal Stack = {HaveMoney, CarRentalOpen}
- 12. Pop HaveMoney. HaveMoney is a known fact (rule v), so we continue.
 - Goal Stack = {CarRentalOpen}
- 13. Pop CarRentalOpen. We branch by pushing the antecedents in rule l (could have used l, m, or n)
 - Goal Stack = {HertzOpen}
- 14. Pop HertzOpen. HertzOpen is not provable, so we backtrack to **Step 13** and push the antecedents in rule m instead.
 - Goal Stack = {AvisOpen}
- 15. Pop AvisOpen. AvisOpen is a known fact, so we continue.
 - Goal Stack $= \{\}$
- 16. Since the Goal Stack is empty, we have demonstrated that we "CanGetToWork" with our knowledge base and facts.