

# CITS2211 Assignment 1

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## Question 1

a)  $P \vee (Q \vee \neg P)$

$P$	$Q$	$\neg P$	$Q \vee \neg P$	$P \vee (Q \vee \neg P)$
$T$	$T$	$F$	$T$	$T$
$T$	$F$	$F$	$F$	$T$
$F$	$T$	$T$	$T$	$T$
$F$	$F$	$T$	$T$	$T$

As the truth table demonstrates, the proposition is a tautology since it's always true regardless of the truth values of  $Q$  or  $P$

b)  $(P \wedge \neg P) \vee \neg Q$

$P \wedge \neg P = F$  (by contradiction)

$\neg Q$  (By absorption)

the proposition depends on the truth value of  $Q$  hence it's a contingent

c)  $Q \rightarrow (P \wedge \neg Q)$

$P$	$Q$	$\neg Q$	$P \wedge \neg Q$	$Q \rightarrow (P \wedge \neg Q)$
$T$	$T$	$F$	$F$	$F$
$T$	$F$	$T$	$T$	$T$
$F$	$T$	$F$	$F$	$F$
$F$	$F$	$T$	$F$	$T$

As the truth table demonstrates, the truth value of the proposition depends on the truth values of  $P$  and  $Q$ . Hence, the proposition is a contingent.

## Question 2

$$P \vee \neg(P \vee \neg Q) \equiv P \vee \neg Q$$

1. $P \vee \neg(P \vee \neg Q)$	Premise
2. $P \vee (\neg P \wedge Q)$	1, Demorgan's laws
3. $(P \vee \neg P) \wedge (P \vee Q)$	2, distributivity
4. $T \wedge (P \vee Q)$	3, excluded middle
5. $P \vee Q$	4, absorption

5.  $P \vee Q$   
*Q.E.D.*

### Question 3

$\forall x \exists n (x \leq n \leq x+5 \wedge (\exists a \exists b (a \neq n) \wedge (a \neq 1) \wedge (b \neq n) \wedge (b \neq 1) \wedge (a \times b = n)))$

### Question 4

let  $N(x, y)$  be "x is a neighbour of y"

a) **Anna has no neighbours**

$$\neg(\exists x.N(x, a))$$

b) **Ben has two neighbours**

$$\exists x \exists y (N(x, b) \wedge N(y, b) \wedge x \neq y \wedge \forall z (N(z, b) \rightarrow (z = x \vee z = y)))$$

c) **If somebody is a neighbour of Ben, Ben is also a neighbour of that person**

$$\forall x (N(x, b) \rightarrow N(b, x))$$

d) **Except for Anna, everyone is the neighbour of someone**

$$\forall x (x \neq a \rightarrow \exists y (N(x, y)))$$

### Question 5

a)  $\frac{P}{P}$

$P \equiv P$  by rule of identity  
hence, the inference rule is sound

b)  $\frac{P}{P \leftrightarrow Q}$

if Q is false, the axiom is true but the conclusion is false  
Hence, the inference rule is unsound

c)  $\frac{P \leftrightarrow Q}{P}$

if P and Q are false, the axiom is true but the conclusion is false  
Hence, the inference rule is unsound

d)  $\frac{P \quad Q}{P \vee Q}$

From the premise P we can infer  $P \vee Q$  by disjunction introduction rule  
Hence, the inference rule is sound

$$\text{e) } \frac{P \rightarrow Q \quad \neg\neg P}{Q}$$

1. $\neg\neg P$	Premise
2. $P \rightarrow Q$	Premise
3. $P$	1, double negation
4. $Q$	3, 2, Demorgan's laws

Hence the rule is sound

## Question 6

$$\forall x.(\neg Q(x) \vee P(x)) \vee \exists x(Q(x) \vee (P(x) \wedge R(x))) \rightarrow \exists x.R(x)$$

1. $\forall x(\neg Q(x) \wedge P(x))$	premise
2. $\exists x(Q(x) \vee (P(x) \wedge R(x)))$	premise
3. $Q(a) \vee (P(a) \wedge R(a))$	2, Exist elimination
4. $\neg Q(a) \wedge P(a)$	1, Forall elimination
5. $\neg Q(a)$	4, conjunction elimination
6. $\neg Q(a) \rightarrow (P(a) \wedge R(a))$	3, implication law
7. $P(a) \wedge R(a)$	5, 6, Modus ponens
8. $R(a)$	7, conjunction elimination
9. $\exists x.R(x)$	8, Exists introduction

## Question 7

Assuming the difference between a squidgy and a non-squidgy number multiplied by 2 produces a squidgy number

we have a non-squidgy number  $k$  and a squidgy number  $l$ ,  
 $l$  can be represented as  $\frac{p}{q}$  where  $p$  and  $q$  are integers

$$\left(\frac{p}{q} - k\right) * 2 = \frac{p'}{q'}$$

where  $p'$  and  $q'$  are integers (by the assumption that the result is a squidgy number)

$$k = \frac{2pq' - p'q}{2q'q}$$

By arithmetic

we know that the product and the difference of integers is an integer.

$2pq' - p'q$  and  $2q'q$  are integers

Hence,  $k$  can be represented in the form of  $\frac{a}{b}$  where  $a$  and  $b$  are integers,  
 which contradicts with our premise that  $k$  is a non-squidgy number.

Our assumption that the result is a squidgy number must be false.

Hence, the result of multiplying the difference between a squidgy and a non-squidgy number by 2 is a non-squidgy number

Q.E.D.

## Question 8

this is a bidirection proof so firstly, we prove that if  $x$  is odd then  $5x-1$  is even,  
 and then secondly, prove that if  $5x-1$  is even then  $x$  is odd.

**first direction: if  $x$  is odd then  $5x-1$  is even**

let  $x$  be an arbitrary odd integer then  $x$  is of the form  $2k-1$  where  $k$  is an integer  
 $x = 2k-1$

$5x = 10k - 5$  (multiplying both sides by 5)

$5x - 1 = 10k - 6$  (subtracting 1 from both sides)

$= 2(5k - 3)$  (factoring out 2)

$= 2n$  (where  $n = 5k-3$ )

which is an even number.

therefore, for any odd integer  $x$ ,  $5x - 1$  is an even integer

**second direction: if  $5x-1$  is an even integer then  $x$  is an odd integer**

The contrapositive of the statement is -

if  $x$  is not an odd integer then  $5x-1$  is not an even integer

an integer not being odd implies that the integer is even

and similarly an integer not being even implies that the integer is odd

Hence, we need to prove that if  $x$  is an even integer then  $5x-1$  is an odd integer

let  $x$  be an arbitrary even integer then  $x$  is of the form  $2n$  where  $n$  is an integer

$x = 2n$

$5x = 10n$  (multiplying both sides by 5)

$$\begin{aligned}
 5x - 1 &= 10n - 1 \text{ (subtracting 1 from both sides)} \\
 &= 2(5n) - 1 \text{ (factoring out 2)} \\
 &= 2k - 1 \text{ (where k is } 5n\text{)}
 \end{aligned}$$

Hence,  $5x - 1$  is of the form  $2k - 1$  which is an odd number

therefore, for any even integer  $x$ ,  $5x - 1$  is an odd integer **Conclusion** therefore, we have proved that if  $x$  is odd, then  $5x-1$  is even

and if  $5x-1$  is even then  $x$  is odd.

Hence, it follows that  $x$  is odd if and only if  $5x-1$  is even

## Question 9

Skill issues

## Question 10