

HW2 (1)

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Homework 2 Resolution

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Q1. SpongeBob and Pacman (Search Formulation)

Pacman bought a car, was speeding in Pac-City, and SpongeBob wasn't able to catch him. Now Pacman has run out of gas, his car has stopped, and he is currently hiding out at an undisclosed location. In this problem, you are on SpongeBob's side, tryin' to catch Pacman!

There are still p of SpongeBob's cars in the Pac-city of dimension m by n . In this problem, **all of SpongeBob's cars can move, with two distinct integer controls: throttle and steering, but Pacman has to stay stationary**. Spongebob's cars can control both the throttle and steering for each step. Once one of SpongeBob's cars takes an action which lands it in the same grid as Pacman, Pacman will be caught and the game ends.

Throttle: $t_i \in \{1, 0, -1\}$, corresponding to {Gas, Coast, Brake}. This controls the **speed** of the car by determining its acceleration. The integer chosen here will be added to the velocity for the next state. For example, if a SpongeBob car is currently driving at 5 grid/s and chooses Gas (1) it will be traveling at 6 grid/s in the next turn.

Steering: $s_i \in \{1, 0, -1\}$, corresponding to {Turn Left, Go Straight, Turn Right}. This controls the **direction** of the car. For example, if a SpongeBob car is facing North and chooses Turn Left, it will be facing West in the next turn.

- (a) Suppose you can **only control 1 SpongeBob car**, and have absolutely no information about the remainder of $p - 1$ SpongeBob cars, or where Pacman has stopped to hide. Also, the SpongeBob cars can travel up to 6 grid/s so $0 \leq v \leq 6$ at all times.

- (i) What is the **tightest upper bound** on the size of state space, if your goal is to use search to plan a sequence of actions that guarantees Pacman is caught, no matter where Pacman is hiding, or what actions other SpongeBob cars take. Please note that your state space representation must be able to represent **all** states in the search space.

$28 \cdot (2)^{mn} \sim \text{possible speeds} \times \text{directions} \times \text{whether the node has been visited or not } (c = 0, 1)$

- (ii) What is the maximum branching factor? Your answer may contain integers, m, n .

$9 \rightarrow (s_i)_{i=1,0,-1} \times (t_i)_{i=1,0,-1} \sim 3 \cdot 3 = 9$

- (iii) Which algorithm(s) is/are guaranteed to return a path passing through all grid locations on the grid, if one exists?

☐ Depth First Tree Search ☒ Breadth First Tree Search
☒ Depth First Graph Search ☒ Breadth First Graph Search

DFS-tree can get stuck on loops

- (iv) Is Breadth First Graph Search guaranteed to return the path with the shortest number of **time steps**, if one exists?

☒ Yes ☐ No

- (b) Now let's suppose you can control **all** p SpongeBob cars at the same time (and know all their locations), but you still have no information about where Pacman stopped to hide

- (i) Now, you still want to search a sequence of actions such that the paths of p SpongeBob cars combined **pass through all $m \times n$ grid locations**. Suppose the size of the state space in part (a) was N_1 , and the size of the state space in this part is N_p . Please select the correct relationship between N_p and N_1 .

☐ $N_p = p \cdot N_1$ ☐ $N_p = p^{N_1}$ ☐ $N_p = (N_1)^p$ ☒ None of the above

Actually $N_p = (28mn)^p \cdot 2^{mn}$ as we still need only one boolean list to keep track of visited nodes

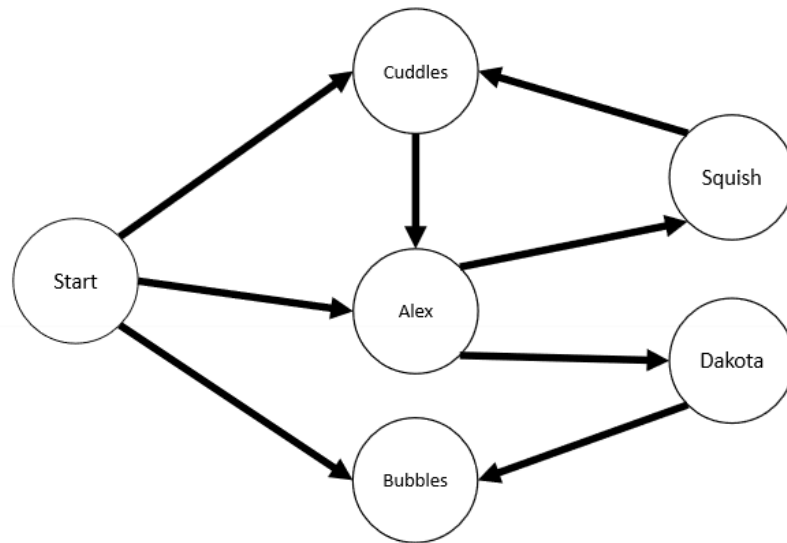
- (ii) Suppose the maximum branching factor in part (a) was b_1 , and the maximum branching factor in this part is b_p . Please select the correct relationship between b_p and b_1 .

☐ $b_p = p \cdot b_1$ ☐ $b_p = p^{b_1}$ ☒ $b_p = (b_1)^p$ ☐ None of the above

We have "combination" of choices for each p -th Spongebob car

Q2. Search: Snail search for love

Scorpborg the snail is looking for a mate. It can visit different potential mates based on a trail of ooze to nearby snails, and then test them for chemistry, as represented in the below graph, where each node represents a snail. In all cases, nodes with equal priority should be visited in alphabetical order.



(a) Simple search

In this part, assume that the only match for Scorpborg is Squish (i.e. Squish is the goal state). Which of the following are true **when searching the above graph**?

- (i) BFS Tree Search expands more nodes than DFS Tree Search ☒ True ☐ False
- (ii) DFS Tree Search finds a path to the goal for this graph ☒ True ☐ False **No loops**
- (iii) DFS Graph Search finds the shortest path to the goal for this graph ☒ True ☐ False
- (iv) If we remove the connection from Cuddles → Alex, can DFS Graph Search find a path to the goal for the altered graph? ☒ Yes ☐ No **Still reachable**

Visited nodes by strategy

BFS

- 1 Start
- 2 Alex
- 3 Bubbles
- 4 Cuddles
- 5 Dakota
- 6 Squish

DFS

- 1 Start
- 2 Alex
- 3 Dakota
- 4 Bubbles
- 5 ² Squish

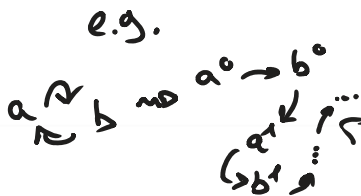
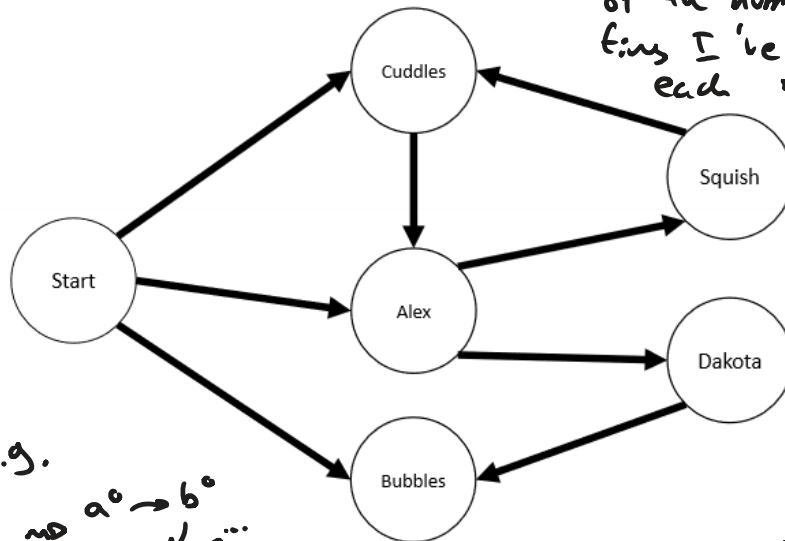
(b) **Third Time's A Charm**

Now we assume that Scorpblorg's mate preferences have changed. The new criteria she is looking for in a mate is that she has **visited the mate twice before** (i.e. when she visits any state for the third time, she has found a path to the goal).

(i) What should the most simple yet sufficient new state space representation include?

- ☒ The current location of Scorpblorg → Needed in previous code too
- ☐ The total number of edges travelled so far
- ☐ An array of booleans indicating whether each snail has been visited so far
- ☒ An array of numbers indicating how many times each snail has been visited so far
- ☐ The number of distinct snails visited so far

↓
To keep track of the number of times I've visited each mate



No loops since that the visited node "changes" each time I visit it (counter)

(The graph is copied for your convenience)

- (ii) DFS Tree Search finds a path to the goal for this graph ☒ True ☐ False
- (iii) BFS Graph Search finds a path to the goal for this graph ☒ True ☐ False Completeness
- (iv) If we remove the connection from Cuddles → Alex, can DFS Graph Search find a path to the goal for the altered graph? ☐ Yes ☒ No No loops to let a mate be visited ≥ 3 times

We continue as in part (b) where the goal is still to find a mate who is visited for the third time.

(c) **Costs for visiting snails**

Assume we are using Uniform cost search and we can now add costs to the actions in the graph.

- (i) Can one assign (non-negative) costs to the actions in the graph such that the goal state returned by UCS (Tree-search) changes? ☒ Yes ☐ No
- (ii) Can one assign (potentially negative) costs to the actions in the graph such that UCS (Tree-search) will never find a goal state? ☐ Yes ☒ No

either Alex or Cuddles the appropriate by setting weights I can find either Alex or Cuddles

3

Either way nodes will be revisited, determining a solution