

# Chapter 34: Counting Sort

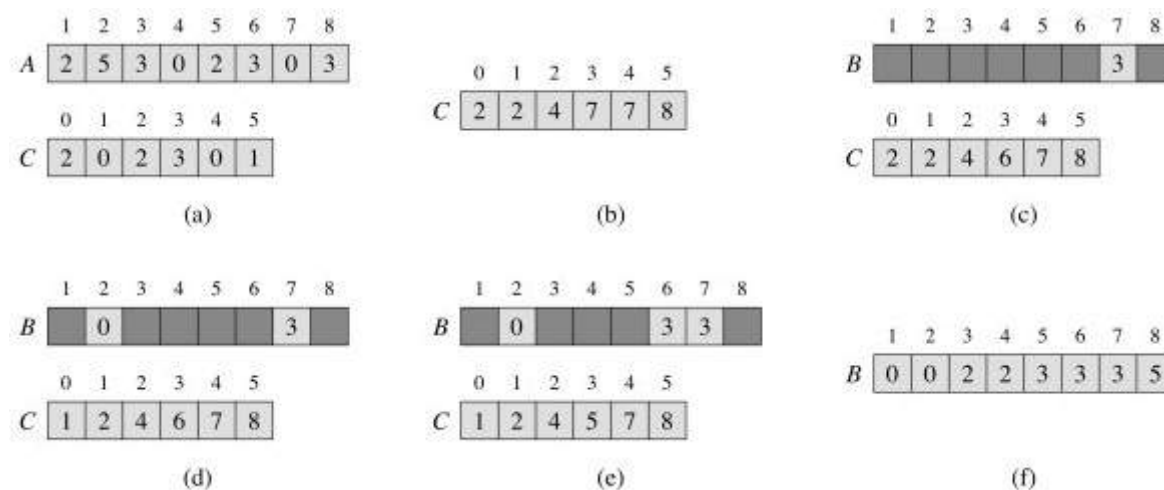
## Section 34.1: Counting Sort Basic Information

Counting sort is an integer sorting algorithm for a collection of objects that sorts according to the keys of the objects.

### Steps

1. Construct a working array  $C$  that has size equal to the range of the input array  $A$ .
2. Iterate through  $A$ , assigning  $C[x]$  based on the number of times  $x$  appeared in  $A$ .
3. Transform  $C$  into an array where  $C[x]$  refers to the number of values  $\leq x$  by iterating through the array, assigning to each  $C[x]$  the sum of its prior value and all values in  $C$  that come before it.
4. Iterate backwards through  $A$ , placing each value in to a new sorted array  $B$  at the index recorded in  $C$ . This is done for a given  $A[x]$  by assigning  $B[C[A[x]]]$  to  $A[x]$ , and decrementing  $C[A[x]]$  in case there were duplicate values in the original unsorted array.

### Example of Counting Sort



**Auxiliary Space:**  $O(n+k)$

**Time Complexity:** Worst-case:  $O(n+k)$ , Best-case:  $O(n)$ , Average-case  $O(n+k)$

## Section 34.2: Psuedocode Implementation

Constraints:

1. Input (an array to be sorted)
2. Number of element in input ( $n$ )
3. Keys in the range of  $0..k-1$  ( $k$ )
4. Count (an array of number)

Pseudocode:

```
for x in input:
    count[key(x)] += 1
total = 0
for i in range(k):
    oldCount = count[i]
    count[i] = total
    total += oldCount
```

```
for x in input:
    output[count[key(x)]] = x
    count[key(x)] += 1
return output
```