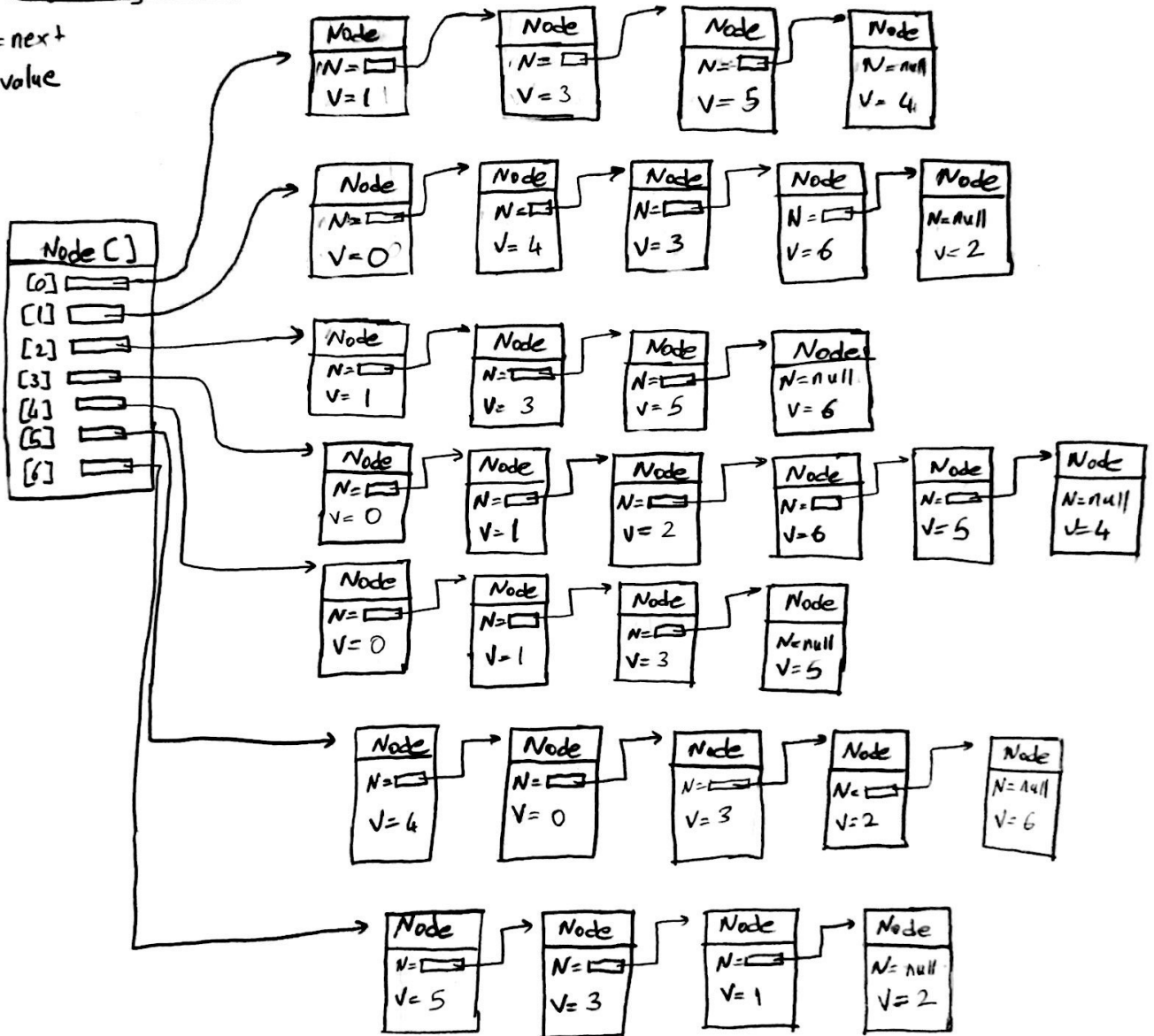


### Adjacency List

N = next  
V = value



## Adjacency Matrix

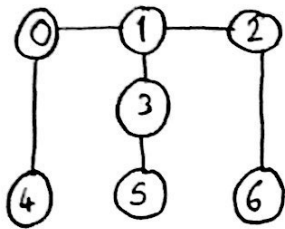
	[0]	[1]	[2]	[3]	[4]	[5]	[6]
[0]		1.0		1.0	1.0	1.0	
[1]	1.0		1.0	1.0	1.0		1.0
[2]		1.0		1.0		1.0	1.0
[3]	1.0	1.0	1.0		1.0	1.0	1.0
[4]	1.0	1.0		1.0		1.0	
[5]	1.0		1.0	1.0	1.0		1.0
[6]		1.0	1.0	1.0		1.0	

$$|V| = 7 \quad |E| = 14 \quad \text{Density} = \frac{2|E|}{|V|(|V|-1)} = \frac{2 \cdot 14}{7 \cdot 6} = \underline{\underline{0.66}}$$

This is dense graph.

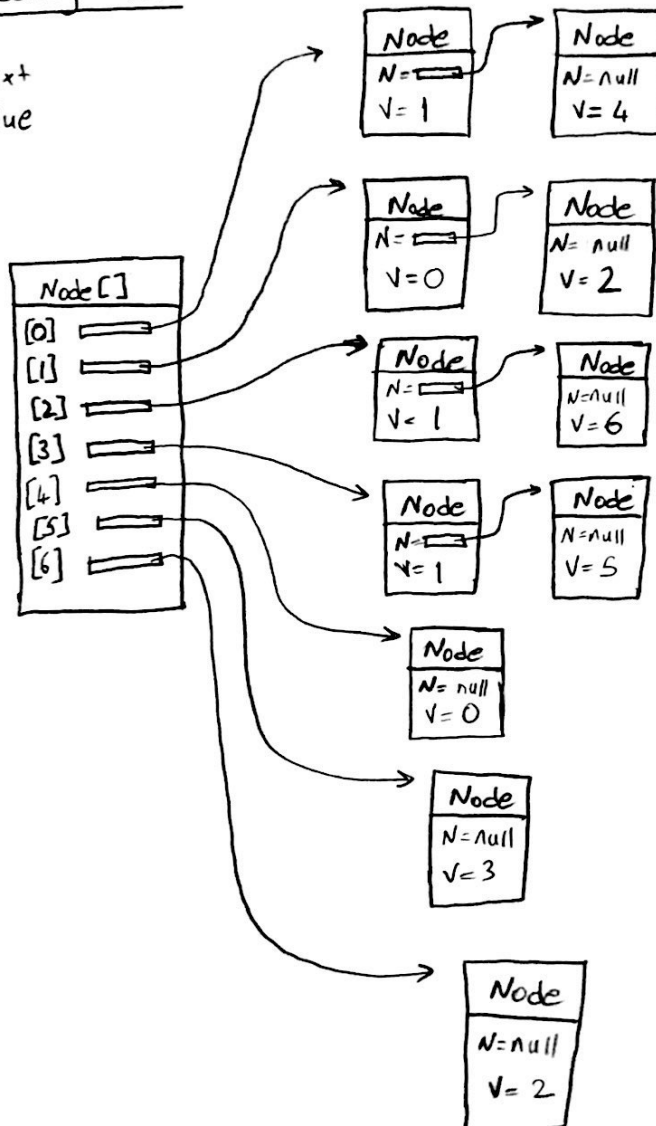
1. for each vertex  $u$  in some subset of the vertices
2. for each vertex  $v$  in some subset of the vertices
3. if  $(u,v)$  is an edge
4. Do something with edge  $(u,v)$

Step 3 for adjacency matrix rep.  $O(1)$  and overall algorithm is  $O(V^2)$   
for adjacency list rep.  $O(|E|)$  and the overall algorithm is  $O(V|E|)$ .  
So adjacency matrix is better for this graph.



### Adjacency List

N = next  
V = value



## Adjacency Matrix

	[0]	[1]	[2]	[3]	[4]	[5]	[6]
[0]		1.0			1.0		
[1]	1.0		1.0	1.0			
[2]		1.0					1.0
[3]		1.0				1.0	
[4]	1.0						
[5]				1.0			
[6]			1.0				

$$|V| = 7 \quad |E| = 6 \quad \text{Density} = \frac{2 \cdot |E|}{|V|(|V|-1)} = \frac{2 \cdot 6}{7 \cdot 6} = \underline{\underline{0.28}}$$

This is sparse matrix

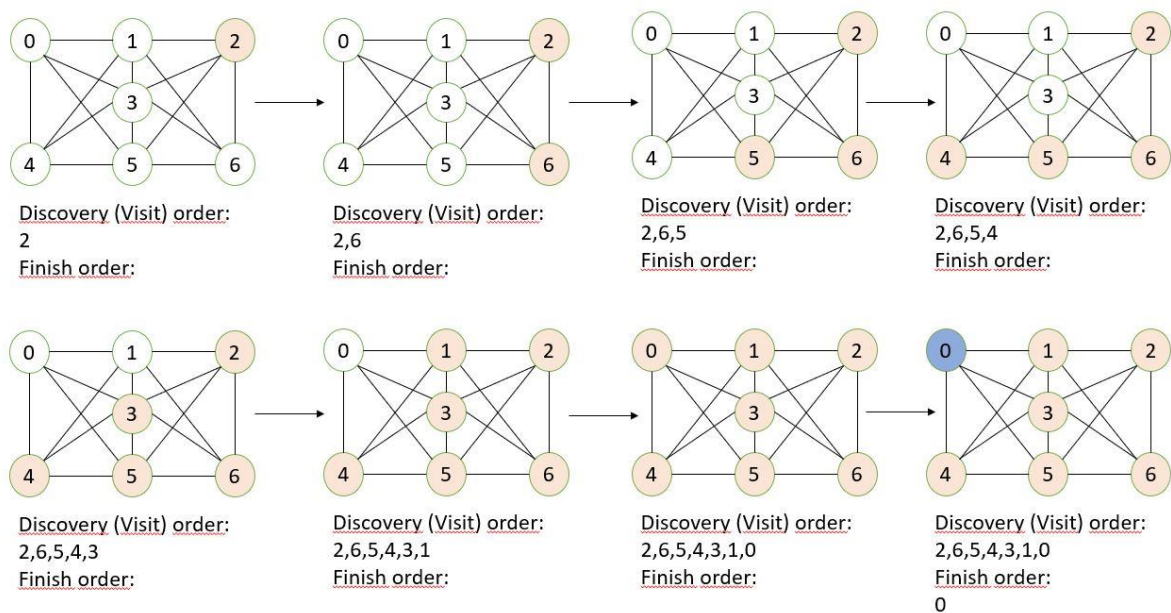
For time efficiency if the graph is sparse, the adjacency list representation is better.

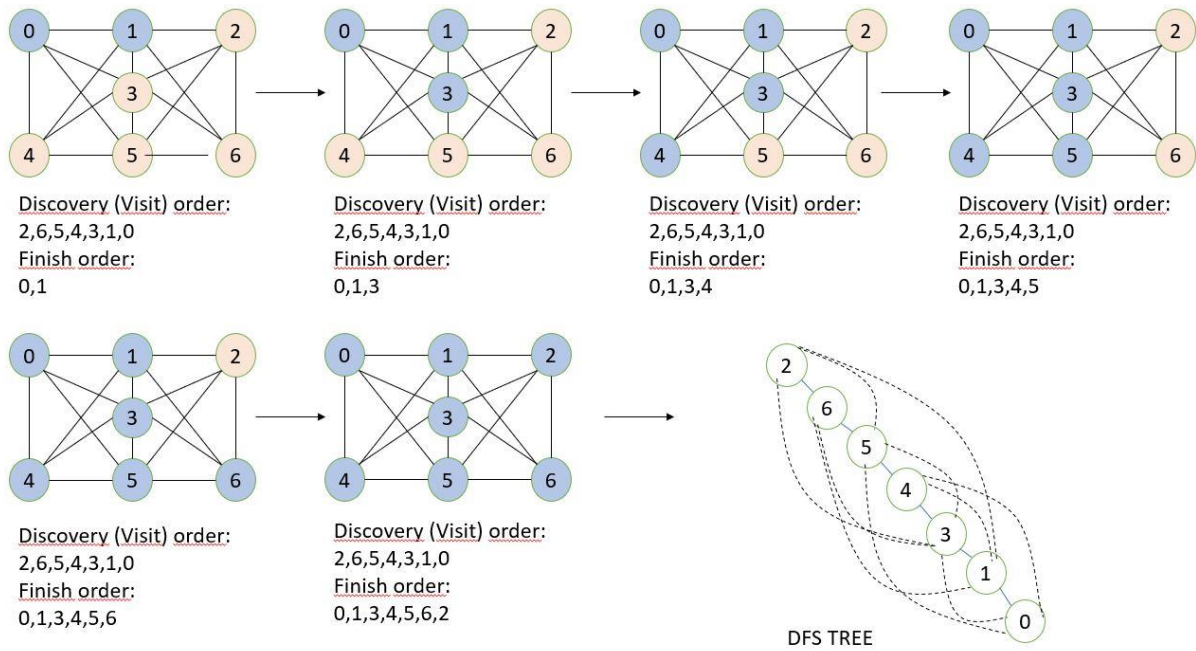
# DEPTH – FIRST SEARCH

## Algorithm for Depth-First Search

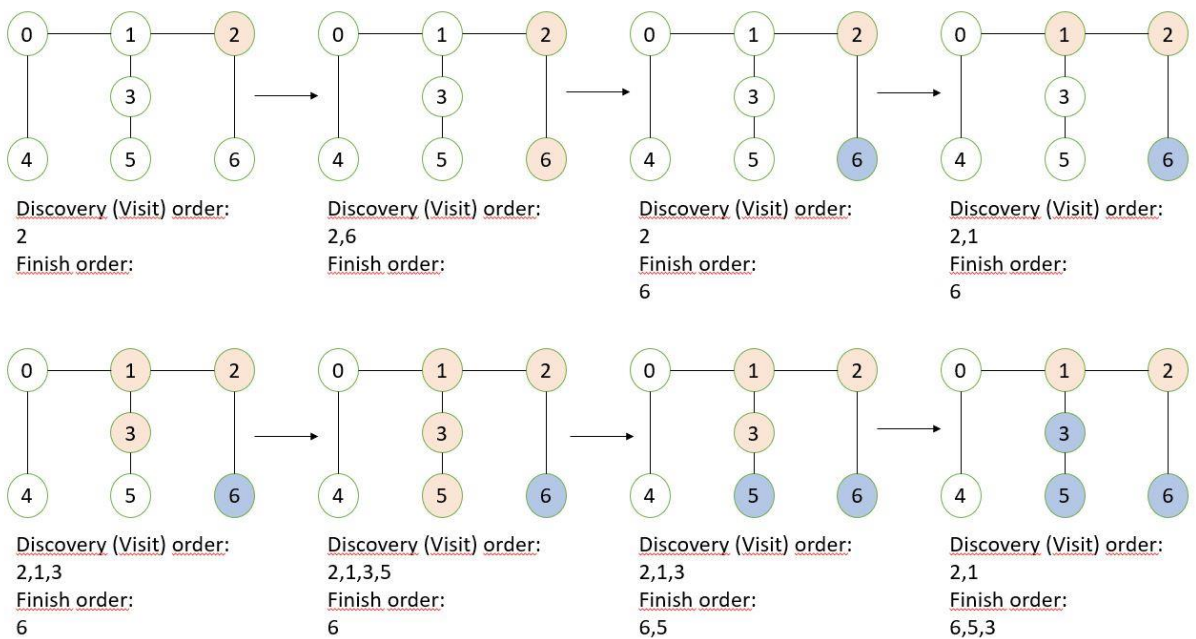
1. Mark the current vertex,  $u$ , visited (color it light orange), and enter it in the discovery order list.
2. for each vertex,  $v$ , adjacent to the current vertex,  $u$
3.     if  $v$  has not been visited
4.         Set parent of  $v$  to  $u$ .
5.         Recursively apply this algorithm starting at  $v$ .
6. Mark  $u$  finished (color it light blue ) and enter  $u$  into the finish order list

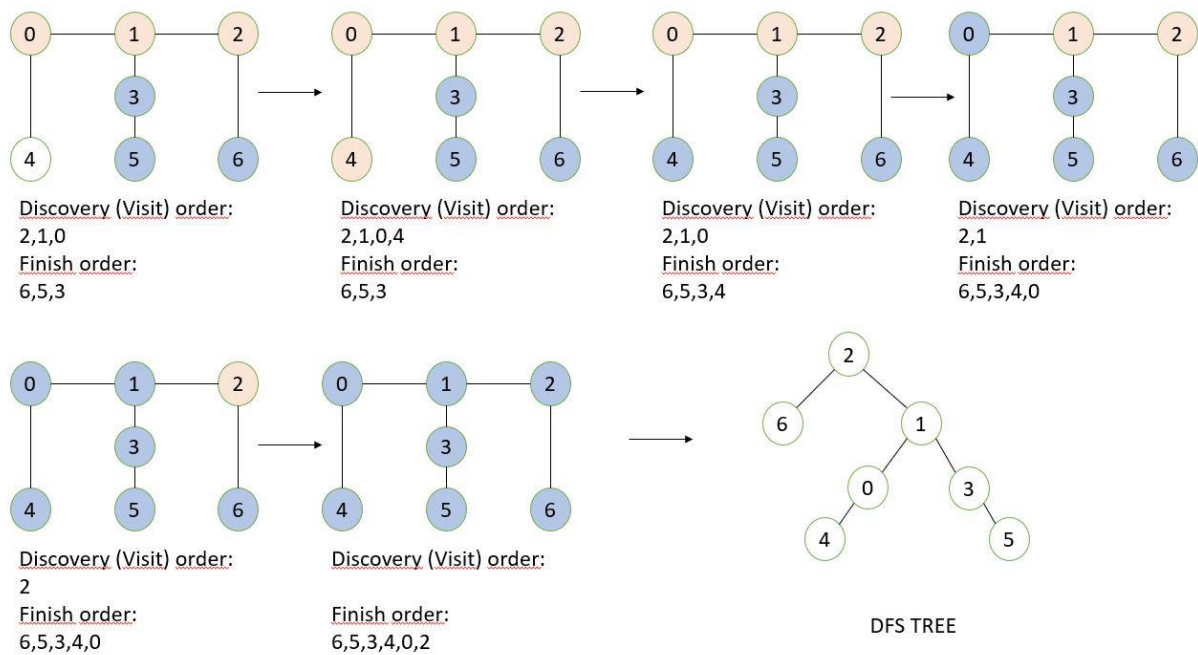
### GRAPH A)





## GRAPH B)



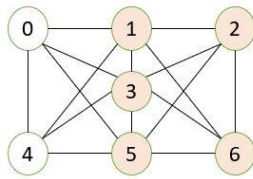


## BREADTH – FIRST SEARCH

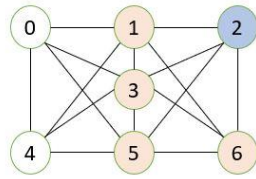
### Algorithm for Breadth-First Search

1. Take an arbitrary start vertex, mark it identified (color it light gray), and place it in a queue.
2. while the queue is not empty
  3. Take a vertex,  $u$ , out of the queue and visit  $u$ .
  4. for all vertices,  $v$ , adjacent to this vertex,  $u$ 
    5. if  $v$  has not been identified or visited
      6. Mark it identified (color it light gray).
      7. Insert vertex  $v$  into the queue.
  8. We are now finished visiting  $u$  (color it dark gray).

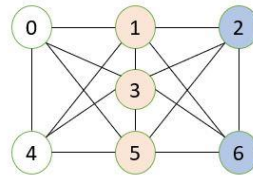
## GRAPH A)



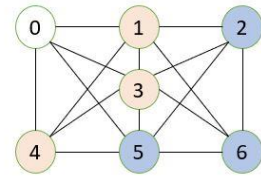
Queue:  
6,5,3,1  
Visit sequence:  
2



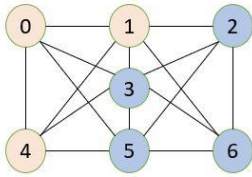
Queue:  
6,5,3,1  
Visit sequence:  
2



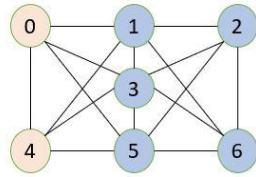
Queue:  
5,3,1  
Visit sequence:  
2,6



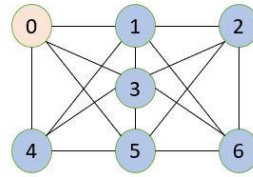
Queue:  
3,1,4  
Visit sequence:  
2,6,5



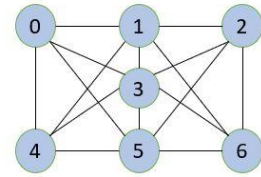
Queue:  
1,4,0  
Visit sequence:  
2,6,5,3



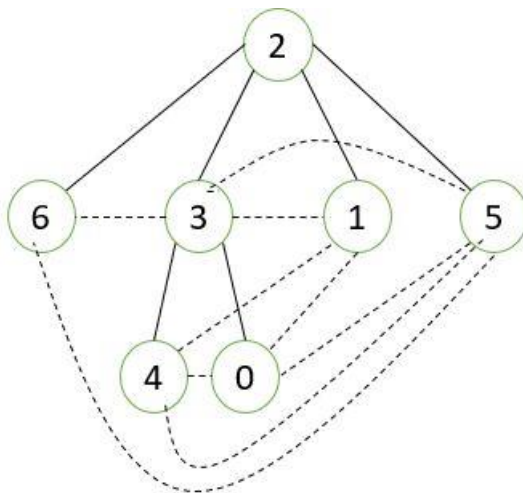
Queue:  
4,0  
Visit sequence:  
2,6,3,1



Queue:  
0  
Visit sequence:  
2,6,3,1,4



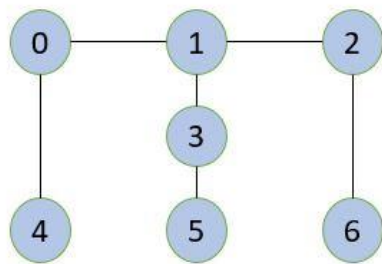
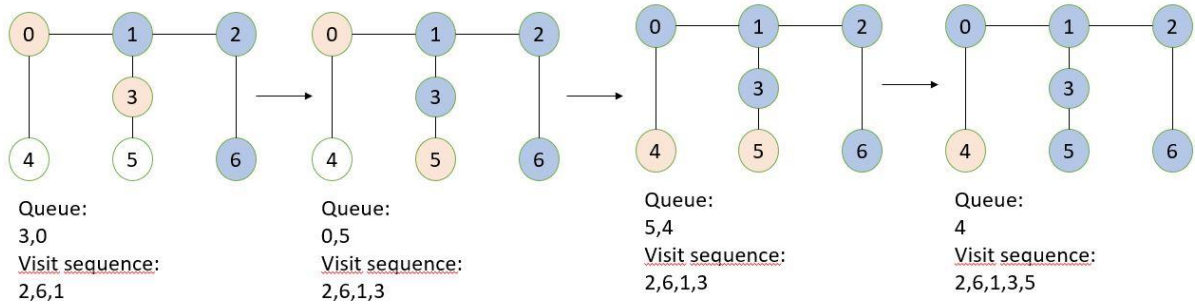
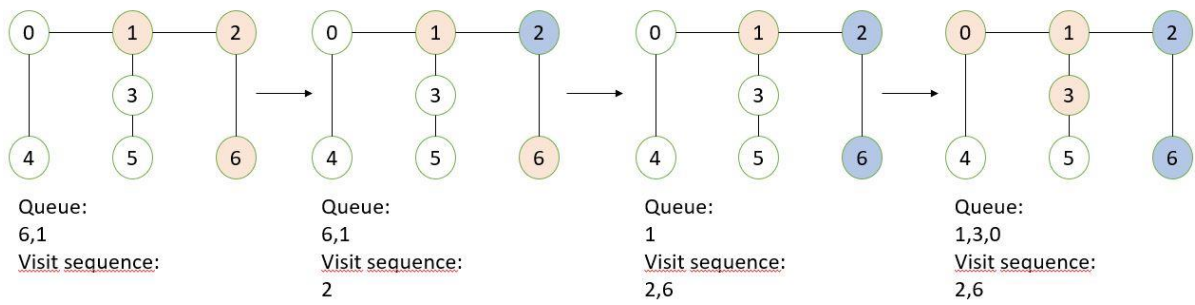
Queue:  
null  
Visit sequence:  
2,6,3,1,4,0



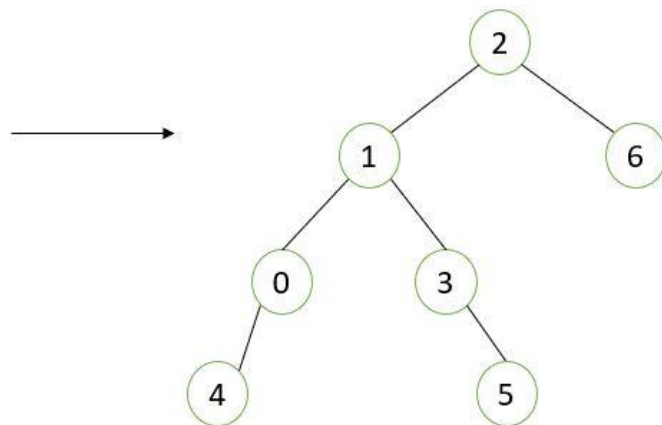
BFS TREE



## GRAPH B)



Queue:  
null  
Visit sequence:  
2,6,1,3,5,4



BFS TREE