

## McLAB: A toolkit for static and dynamic compilers for MATLAB



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Intro - 1

## Overview



- Why MATLAB?
- Overview of the McLAB tools
- McVM – a Virtual Machine and Just-In-Time (JIT) compiler
- McFOR – translating MATLAB to FORTRAN95

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FORTRAN  
C/C++

C, Parallel C,  
Java, AspectJ



**MATLAB**  
PERL  
Python  
Domain-specific

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## Culture Gap

### Scientists / Engineers

- Comfortable with informal descriptions and “how to” documentation.
- Don’t really care about types and scoping mechanisms, at least when developing small prototypes.
- Appreciate libraries, convenient syntax, simple tool support, and interactive development tools.

### Programming Language / Compiler Researchers

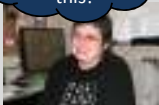
- Prefer more formal language specifications.
- Prefer well-defined types (even if dynamic) and well-defined scoping and modularization mechanisms.
- Appreciate “harder/deeper/more beautiful” programming language/compiler research problems.

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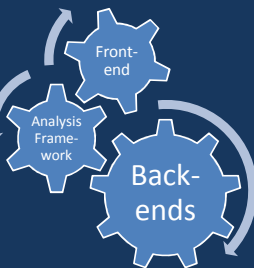
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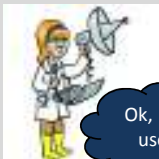
Ok, I can  
deal with  
this!



### McLAB Compiler Framework



Ok, this is  
useful!



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## Goals of the McLab Project

- Improve the understanding and documentation of the semantics of MATLAB.
- Provide front-end compiler tools suitable for MATLAB and language extensions of MATLAB.
- Provide a flow-analysis framework and a suite of analyses suitable for a wide range of compiler/soft. eng. applications.
- Provide back-ends that enable experimentation with JIT and ahead-of-time compilation.

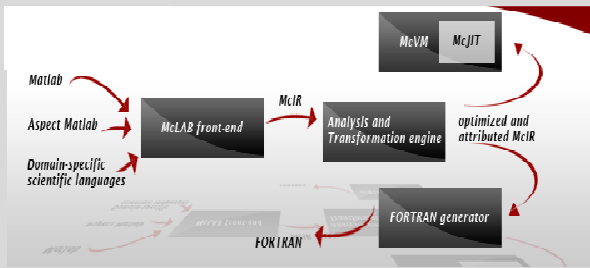
Enable PL, Compiler and SE Researchers to work on MATLAB

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## McLAB – Overall Structure

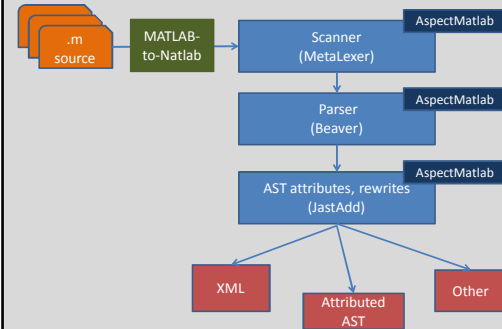


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## McLAB Extensible Front-end



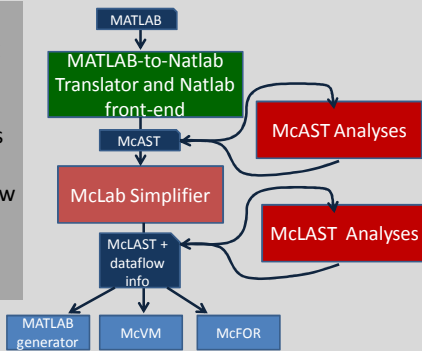
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## Analysis Engine

Analyses are written using an Analysis Framework that supports forward and backward flow analysis over McAST and McLAST.



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## How to build the back-ends?



- No official language specification.
- Closed-source implementations, including the library.

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## Back-end #1: MATLAB generator

- McLAB can be used as a source-to-source translator:
  - source-level optimizations like loop unrolling
  - source-level refactoring tool
- McLAB can generate:
  - xml files (used to communicate with McVM)
  - text .m files
- Comments are maintained to enable readable output.

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## Back-end #2: McVM with JIT compiler



- Based on Maxime Chevalier-Boisvert's M.Sc. thesis, McGill, published in CC 2010.
- Currently being extended by Nurudeen Lameed and Rahul Garg, Ph.D. candidates, McGill.

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## McVM-McJIT

- The dynamic nature of MATLAB makes it very suitable for a VM/JIT.
- MathWorks' implementation does have a JIT, although technical details are not known.
- McVM/McJIT is an open implementation aimed at supporting research into dynamic optimization techniques for MATLAB.

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## Design Choices for JITs

1. Interpreter + JIT compiler with various levels of optimizations.
  2. Fast JIT for naïve code generation + optimizing JIT with various levels of optimizations.
- McVM uses the 1<sup>st</sup> option because it simplifies adding new features, if a feature is not yet supported by the JIT it can back-up to the interpreter implementation, which is easy to provide.

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## McVM Design

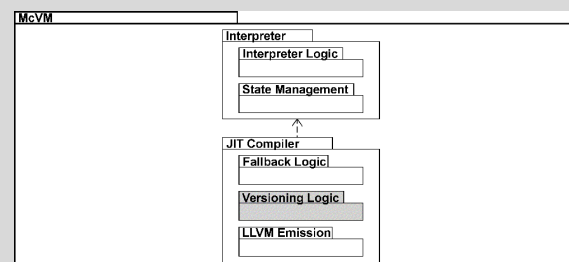
- A basic, but fast, interpreter for the MATLAB language.
- A garbage-collected JIT Compiler as an extension to the interpreter.
- Easy to add new data types and statements by modifying only the interpreter.
- Supported by the LLVM compiler framework and some numerical computing libraries.
- Written entirely in C++; interface with the McLAB front-end via a network port.

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## McVM Organization

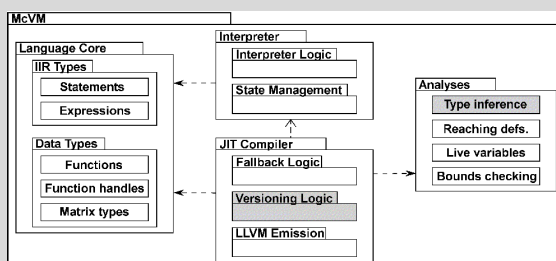


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## McVM Organization

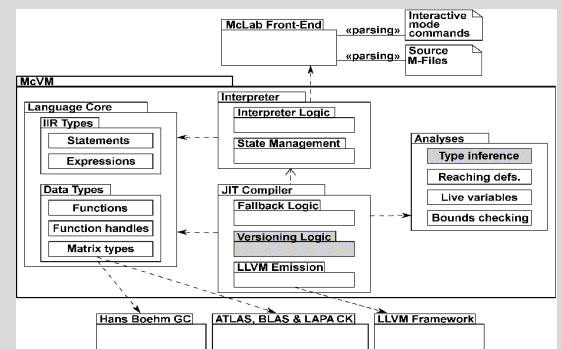


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## McVM Organization



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## MATLAB Optimization Challenges

```
float sumvals(float start, float step, float stop)
{
    float i = start;
    float s = i;

    while (i < stop)
    {
        i = i + step;
        s = s + i;
    }

    return s;
}
```

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## MATLAB Optimization Challenges

```
function s = sumvals(start, step, stop)
    i = start;
    s = i;

    while i < stop
        i = i + step;
        s = s + i;
    end

function caller()
    a = sumvals(1, 1, 10^6);
    b = sumvals([1 2], [1.5 3], [20^5 20^5]);
    c = [a b];

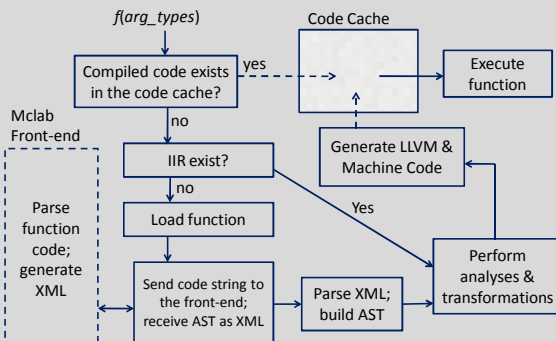
    disp(c);
end
```

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## McJIT: Executing a Function



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## Just-In-Time Specialization (1)

```
>> a = sumvals(1, 1, 10^6);
>> b = sumvals([1 2], [1.5 3], [20^5 20^5]);
>> a = sumvals(1, 1, 500);
>> c = [a b];
>> disp(c);
```

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## Just-In-Time Specialization (2)

```
>> a = sumvals(1, 1, 10^6);
Interpreter.runCommand("a = sumvals(1, 1, 10^6);");
Interpreter.callFunction(sumvals, [1, 1, 10^6]);
JIT.callFunction(sumvals, [1, 1, 10^6]);
```

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## Just-In-Time Specialization (3)

```
JITCompiler.callFunction(sumvals, [1, 1, 10^6]);
sumvalsJIT = JITCompiler.compileFunction(sumvals, [<scalar int>, <scalar int>, <scalar int>]);

function s = sumvals(start <scalar int>, step <scalar int>, stop <scalar int>)
    i = start;
    s = i;
    while i < stop
        i = i + step;
        s = s + i;
    end
end
```

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## Just-In-Time Specialization (4)

```
JITCompiler.callFunction(sumvals, [1, 1, 10^6]);

sumvals.JIT = JITCompiler.compileFunction(sumvals, [<scalar int>, <scalar int>, <scalar int>]);

function s <scalar int> = sumvals(start <scalar int>, step <scalar int>, stop <scalar int>)
    i <scalar int> = start;
    s <scalar int> = i;
    while i < stop
        i <scalar int> = i + step;
        s <scalar int> = s + i;
    end
end
```

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## Just-In-Time Specialization - 2<sup>nd</sup> example

```
>> a = sumvals(1, 1, 10^6);
>> b = sumvals([1 2], [1.5 3], [20^5 20^5]);
>> a = sumvals(1, 1, 500);
>> c = [a b];
>> disp(c);
```

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## JIT – second specialization (1)

```
JITCompiler.callFunction(sumvals, [[1 2], [1.5 3], [20^5 20^5]]);

sumvals.JIT2 = JITCompiler.compileFunction(sumvals, [<1x2 int>, <1x2 real>, <1x2 int>]);

function s = sumvals(start <1x2 int>, step <1x2 real>, stop <1x2 int>)
    i <1x2 int> = start;
    s <1x2 int> = i;
    while (i <1x2 int>) < (stop <1x2 int>)
        i <1x2 real> = i + step;
        s <1x2 real> = s + i;
    end
end
```

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## JIT – second specialization (2)

```
JITCompiler.callFunction(sumvals, [[1 2], [1.5 3], [20^5 20^5]]);

sumvals.JIT2 = JITCompiler.compileFunction(sumvals, [<1x2 int>, <1x2 real>, <1x2 int>]);

function s <1x2 real> = sumvals(start <1x2 int>, step <1x2 real>, stop <1x2 int>)
    i <1x2 int> = start;
    s <1x2 int> = i;
    while (i <1x2 real>) < (stop <1x2 int>)
        i <1x2 real> = i + step;
        s <1x2 real> = s + i;
    end
end
```

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## JIT – third specialization same as first

```
>> a = sumvals(1, 1, 10^6);
>> b = sumvals([1 2], [1.5 3], [20^5 20^5]);
>> a = sumvals(1, 1, 500);
>> c = [a b];
>> disp(c);
```

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## Type and Shape Inference

- In MATLAB, work with incomplete information
  - Dynamic loading: working with incomplete program
  - Dynamic typing: variables can change type
- Know argument types, what can we infer?
  - Propagate type info to deduce locals type and return type
- Forward dataflow analysis
  - Based on abstract interpretation
  - Structure-based fixed point
  - Annotates AST with type info

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## Flow Analysis Summary

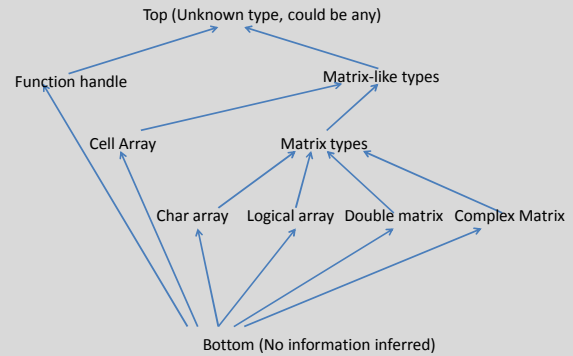
- Start from known argument types
- Propagate type information forward
- Use a transfer function for each expression
  - Transfer functions provided for all primitive operators
  - Library functions provide their own transfer functions
  - Function calls resolved, recursively inferred
- Assignment statements can change var. types
- Merge operator
  - Union + filtering

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## Lattice of McVM types



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Backends-32

## Type Abstraction Properties

- Collection of simple abstractions
  - Specific features computed in parallel
- Represent variable types with 8-tuples:

```
<overallType, is2D, isScalar, isInteger,
  sizeKnown, size, handle, cellTypes>
```

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## Type Abstraction Properties

```
a = [1 2];
```

```
type(a) =
  <overallType = double, is2D = T, isScalar = F,
  isInteger = T, sizeKnown = T, size = (1,2),
  handle = null, cellTypes = {}>
```

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## Variable Types: Type Sets

```
if (c1)
  a = [1 2];
elseif (c2)
  a = 1.5;
else
  a = 'a';
end
```

```
type(a) = {
  <overallType = double, is2D = T, isScalar = F,
  isInteger = T, sizeKnown = T, size = (1,2),
  handle = null, cellTypes = {}>,
  <overallType = double, is2D = T, isScalar = T,
  isInteger = F, sizeKnown = T, size = (1,1),
  handle = null, cellTypes = {}>,
  <overallType = char, is2D = T, isScalar = T,
  isInteger = T, sizeKnown = T, size = (1,1),
  handle = null, cellTypes = {}>
}
```

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## Type Set Filtering

```
type(a) = {
  <overallType = double, is2D = T, isScalar = F,
  isInteger = T, sizeKnown = T, size = (1,2),
  handle = null, cellTypes = {}>,
  <overallType = double, is2D = T, isScalar = T,
  isInteger = F, sizeKnown = T, size = (1,1),
  handle = null, cellTypes = {}>
}
```

Becomes:

```
type(a) = {
  <overallType = double, is2D = T, isScalar = F,
  isInteger = F, sizeKnown = F, size = (),
  handle = null, cellTypes = {}>
}
```

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## Transfer Functions

```
a = 5 * 10e11;
b = 1.0e12 * [0.8533 1.7067];
c = [a b];
```

```
type(a) = {
  <overallType = double, is2D = T, isScalar = T, isInteger = T,
  sizeKnown = T, size = (1,1), handle = null, cellTypes = {}>
}

type(b) = {
  <overallType = double, is2D = T, isScalar = F, isInteger = F,
  sizeKnown = T, size = (1,2), handle = null, cellTypes = {}>
}

type([a b]) = {
  <overallType = double, is2D = T, isScalar = F, isInteger = F,
  sizeKnown = T, size = (1,3), handle = null, cellTypes = {}>
}
```

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## Experimental Results

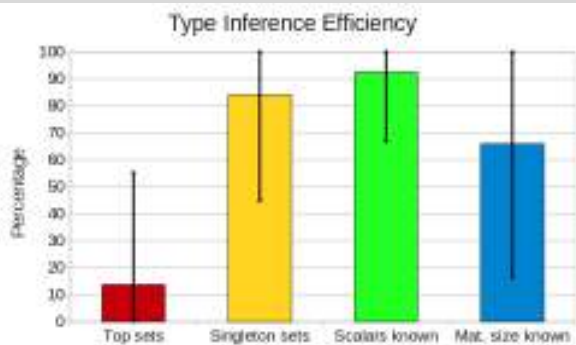
- 20 benchmark programs
  - FALCON, OTTER, etc. Some made by McLAB
- Measured
  - Dynamic availability of type info.
  - Number of versions compiled
  - Compilation time
    - 0.55s per benchmark, on average

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## Results of Type Analysis

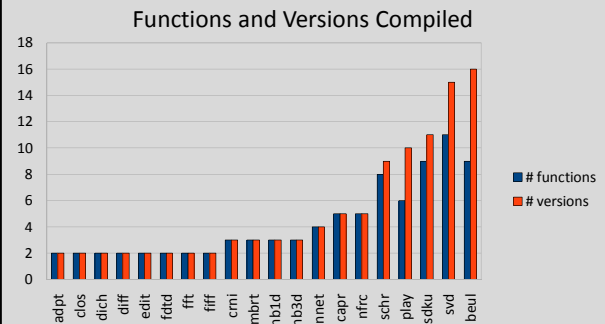


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## How many versions...



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## Performance Results

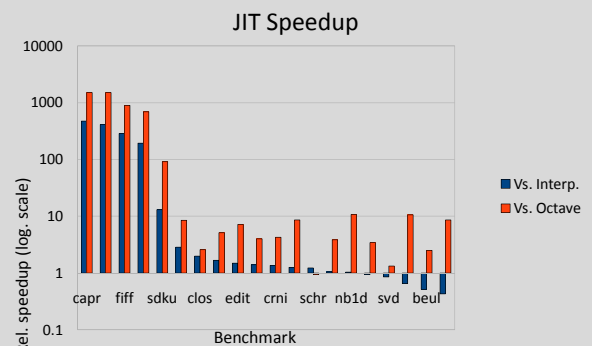
- Experimental setup
  - Core 2 Quad Q6600, 4GB RAM
  - Ubuntu 9.10, kernel 2.6.31, 32-bit
  - All timings averaged over 10 runs
- Comparing
  - McVM interpreter, McVM JIT w/ spec.
  - MATLAB R2009a
  - GNU Octave 3.0.5
  - McFor / GNU Fortran 4.4.1

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## Results (Speed-up vs Interpreters)



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## Value Range Propagation

### MATLAB

```
n = floor(11.5); // n=[n,n]
for i = 1:n      // i=[1,n]
    x = 1+2*i;   // x=[3,1+2*n]

A(x) = i;       // Size(A)=
end             1+2*n
n = fix(n/2);   // n = [n,n]

A(n+1) = n;
```

### FORTTRAN

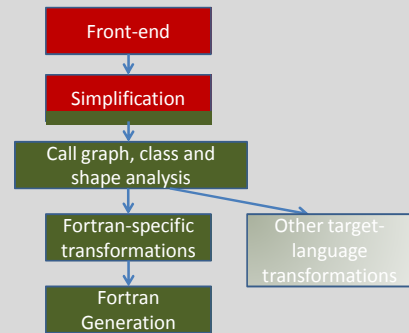
```
n = floor(10.5);
DO i = 1,n
    x = (1+(2*i));
    IF((.NOT.ALLOCATED(A))) THEN
        ALLOCATE(A((1+(2*n))));
    END IF
    A(x) = i;
END DO
n = fix(n/2);
ARRAYBOUNDSCHECKING(A,[n+1]);
A(n+1) = n;
```

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## Basic structure of 2<sup>nd</sup> generation McFOR



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## Related Work

- Procedure cloning: Cooper et al. (1992)
- MATLAB type inference: Joisha & Banerjee (2001)
  - Suggested for error detection
- MATLAB Partial Evaluator: Elphick et al. (2003)
  - Source-to-source transformation
- MaJIC: JIT compilation and offline code cache (2002)
  - Speculative compilation MATLAB to C/Fortran
- Psyco: Python VM with specialization by need (2004)
- TraceMonkey: JIT optimization of code traces (2009)

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## Ongoing Work

- McVM:
  - profile-guided optimization and re-optimization with on-stack replacement
  - target GPU/multi-core
- McFOR:
  - "decompile" to more programmer-friendly FORTRAN95
  - refactoring toolkit to help restructure "dynamic" features to "static" features

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## Conclusions

- McLAB is a toolkit to enable PL, Compiler and SE research for MATLAB
- front-end for language extensions
- analysis framework
- three back-ends including McVM and McFOR

<http://www.sable.mcgill.ca/mclab>

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Intermediate Representation

## EXTRA SLIDES

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## IIR: A Simple MATLAB Program

## McVM Project Class Hierarchy (C++ Classes)

## Supported Types