## 1 Parser

Once the actual code has been separated from proof code, the Parser then parses into an abstract syntax tree.

```
module Parser where
import Data.Map (Map)
import qualified Data.Map.Strict as Map
import Lexer (lexify, Token)
```

The NativeAST is, for now, a placeholder for whatever type is produced by the lexer/parser of the language being proven.

```
\begin{array}{l} \textbf{data} \  \, \textbf{NativeAST} = \textbf{NativeASTNode} \\ \\ \textbf{parseCode} \  \, :: \  \, \textbf{String} \  \, \rightarrow \  \, \textbf{AST} \\ \\ \textbf{parseCode} \  \, \textbf{code} = \textbf{transformAST} \  \, \textbf{NativeASTNode} \\ \\ \textbf{transformAST} \  \, :: \  \, \textbf{NativeAST} \  \, \rightarrow \  \, \textbf{AST} \\ \\ \textbf{transformAST} \  \, \textbf{native} = \textbf{ID} \  \, \textbf{"The code"} \end{array}
```

Once the code has been turned into a NativeAST, it is then transformed into the AST by the pluggable Transformer. Meanwhile, the proof code must also be converted into the definitions used to prove the program. These are represented by the same AST as the code, but this transformation is handled here.

```
data AST = ID String -- name

| Type AST [AST] -- ID Parameters
| TypeOf AST -- Type
| Annotation AST AST -- ID Type
| Let AST AST AST -- ID Type Body
| Function [AST] AST -- Parameters Body
| Application AST [AST] -- Function Parameters
| Exists AST AST AST -- ID Type Body
| IntroExists AST AST -- Type Value [will this need another argument?]
| ElimExists AST AST -- Exists Body [how does this work again? does it need another argument too?]
| And AST AST -- Type Type
| IntroAnd AST AST -- Left Right
```

```
ElimAndLeft AST AST -- And Body
 ElimAndRight AST AST -- And Body
 Or AST AST -- Type Type
 IntroOrLeft AST AST -- Or Value
 IntroOrRight AST AST -- Or Value
 ElimOr AST AST AST -- Or LeftBody RightBody
 Contradiction
 ElimContradiction AST AST -- Contradiction Body [does this
    have a body? contradiction usually means done]
-- value nodes
| VInteger Int -- Value
 VFloat Float -- Value [is this needed? or just define as a
    pair or in STL]
 VChar Char -- Value [is this needed? or just define as an
    int or in STL]
 VBoolean Bool -- True/False
 VCons AST AST -- Head Tail
 VEmpty -- empty list
 VSymbol String -- For
 VNull -- the empty value
VUndefined -- the non-existent value
-- induction [do these need that 4th param like last time?]
IndInteger AST AST AST -- Int BodyS BodyZ [what if it isn't
    natural, is it the same?]
-- [how to use a float? is float usage STL?]
-- [how to use a char? is char usage STL?]
| IndBoolean AST AST AST -- Bool BodyT BodyF
| IndList AST AST AST -- List BodyL BodyE [is this correct?]
-- [how to use a symbol?]
-- [how to use null?]
-- [how to use undefined?]
```

The first step in parsing the proof code is, of course, lexifying it. This step is taken on by the Lexer.

```
parseProofs :: String → AST
parseProofs proofText =
  parse $ lexify proofText

parse :: [Token] → AST
parse _ = ID "The proof"
```

Once parsing is complete the two trees are merged into one containing the actual code annotated by proof terms. This is the final tree which is returned to

the Compiler to be used by the Analyzer in assuring that the program is valid.

annotates :: AST  $\rightarrow$  AST  $\rightarrow$  AST

annotates proof code = Annotation code proof