1 Parser

Once the actual code has been separated from proof code, the Parser then parses into an abstract syntax tree.

```
module Parser where
import Data.Map (Map)
import qualified Data.Map.Strict as Map
import qualified Lexer
```

The NativeAST is, for now, a placeholder for whatever type is produced by the lexer/parser of the language being proven.

```
\begin{array}{l} \textbf{data} \  \, \textbf{NativeAST} = \textbf{NativeASTNode} \\ \\ \textbf{parseCode} \  \, :: \  \, \textbf{String} \  \, \rightarrow \  \, \textbf{AST} \\ \\ \textbf{parseCode} \  \, \textbf{code} = \textbf{transformAST} \  \, \textbf{NativeASTNode} \\ \\ \textbf{transformAST} \  \, :: \  \, \textbf{NativeAST} \  \, \rightarrow \  \, \textbf{AST} \\ \\ \textbf{transformAST} \  \, \textbf{native} = \textbf{ID} \  \, \textbf{"The code"} \  \, \textbf{VEmpty} \end{array}
```

Once the code has been turned into a NativeAST, it is then transformed into the AST by the pluggable Transformer. Meanwhile, the proof code must also be converted into the definitions used to prove the program. These are represented by the same AST as the code, but this transformation is handled here.

```
ElimAndLeft AST AST -- And Body
 ElimAndRight AST AST -- And Body
 Or AST AST -- Type Type
 IntroOrLeft AST AST -- Or Value
 IntroOrRight AST AST -- Or Value
 ElimOr AST AST AST -- Or LeftBody RightBody
 Contradiction
 ElimContradiction AST AST -- Contradiction Body [does this
    have a body? contradiction usually means done]
-- value nodes
| VInteger Int -- Value
 VFloat Float -- Value [is this needed? or just define as a
    pair or in STL]
 VChar Char -- Value [is this needed? or just define as an
    int or in STL]
 VBoolean Bool -- True/False
 VCons AST AST -- Head Tail
 VEmpty -- empty list
 VSymbol String -- For
 VNull -- the empty value
| VUndefined -- the non-existent value
-- induction [do these need that 4th param like last time?]
IndInteger AST AST AST -- Int BodyS BodyZ [what if it isn't
    natural, is it the same?]
-- [how to use a float? is float usage STL?]
-- [how to use a char? is char usage STL?]
| IndBoolean AST AST AST -- Bool BodyT BodyF
| IndList AST AST AST -- List BodyL BodyE [is this correct?]
-- [how to use a symbol?]
-- [how to use null?]
-- [how to use undefined?]
```

The first step in parsing the proof code is, of course, lexifying it. This step is taken on by the Lexer. After that, we move on to parsing, using the LR(1) algorithm. I think. That's what I'm going for anyway.

```
parseProofs :: String → AST
parseProofs proofText = parse $ Lexer.lexify proofText

parse :: [Lexer.Token] → AST
parse tokens = fst $ stateStart $ map Left tokens

type StateFn = [Either Lexer.Token AST] → (AST, [Either Lexer.Token AST])

stateStart :: StateFn
```

```
stateStart (Left Lexer.BOF : rest) = stateDefns rest
stateDefns :: StateFn
stateDefns (Left Lexer.Type : rest) = stateTypeDef rest
{\tt stateDefns} (Left Lexer.Let : rest) = {\tt stateLet} rest
stateTypeDef :: StateFn
stateTypeDef rest =
 let (name, rest) = stateID rest in
    let (args, rest) = stateType rest in
      let (body, rest) = stateDefns rest in
        (Let name args body, rest)
stateLet :: StateFn
stateLet rest =
 let (name, rest) = stateID rest in
   let (typ, rest) = stateStart rest in
      let (body, rest) = stateDefns rest in
        (Let name typ body, rest)
stateID :: StateFn
stateID (Left (Lexer.ID name) : rest) =
 let (args, rest) = stateTypeArgs rest in
    (ID name args, rest)
stateTypeArgs :: StateFn
stateTypeArgs (Left Lexer.LBrack : rest) = stateTypeArgsList rest
stateTypeArgs rest = (VEmpty, rest)
stateTypeArgsList :: StateFn
\verb|stateTypeArgsList (Left Lexer.RBrack : rest) = (VEmpty, rest)|\\
stateTypeArgsList rest =
 let (ann, rest) = stateAnnotation rest in
    let (end, rest) = stateTypeArgsList rest in
      (ArgumentList ann end, rest)
stateAnnotation :: StateFn
stateAnnotation rest =
 let (name, rest) = stateID rest in
    let (typ, rest) = stateType rest in
      (Annotation name typ, rest)
-- TODO: this one very complex
stateType :: StateFn
stateType _ = (ID "The type!" VEmpty, [])
```

Once parsing is complete the two trees are merged into one containing the actual code annotated by proof terms. This is the final tree which is returned to

the Compiler to be used by the Analyzer in assuring that the program is valid.

annotates :: AST \rightarrow AST \rightarrow AST

annotates proof code = Annotation code proof