### Part A

## -- Parsing and initialising the structure of the network --

The graph.txt is parsed to create a ring-structure for the nodes. The nodes list is used to store this. Initially, only the neighbour ids of a node are stored. Only after successful parsing of the graph.txt file, the respective "myNeighbours" list of each node is populated with help from the stored neighbour ids. Besides, it's made sure that the ring structure is bi-directional. Election.txt is parsed, the electionRounds HashMap (<electionRound, <nodes>>) is populated.

## -- Starting elections and the procedure at each round--

The first round is initiated and incremented until it is an election round. At this point, the relevant nodes announce the start of the election. *Each node has an outgoing message queue* to ensure only one message is sent from one node to another in each round, but also the remaining messages are not lost and are eventually delivered in the subsequent rounds. *addMessages()* is called after each election round is announced, this method checks for any node with outgoing messages and add them to the network's *msgToDeliver* structure. Then the *deliverMessages()* method uses *msgToDeliver* to deliver the messages to each node's "right" neighbour (i.e. messages are delivered only clockwise). It ensures only one message is sent per round per neighbour. *Elections conclude once the electionRounds and msgToDeliver are empty*.

# -- Robert Chang Algorithm and logging to the log file--

This logic is implemented in the Node class in the <u>receiveMessage()</u> method.

### Part B

The fail.txt is parsed in a manner similar to election.txt and the <u>failureRounds</u> data structure is populated. For this part, additional logic is added to the <u>NetSimulator()</u> method. The simulation now only ever finishes after there's no more election rounds, no more messages to deliver and no more failure rounds. Whenever the current round happens to be a <u>failure round</u>, all the neighbours of the failed node are informed about the node failure and the node is removed from its neighbours' <u>myNeighbours</u> list. Then the <u>isNetworkConnected()</u> method checks if there will still be a path between the left and right neighbours of the failed node. If yes, the node is removed from the ring and the left and right nodes simply become each others' neighbours, the reason why it's implemented this way is to avoid unnecessary network traffic and computational cost by moving the left and right neighbouring nodes around in the ring to reconstruct the network. If the network is not connected, the program is terminated right away.

All the other logic remains the same.