

CS224

Section No.: 04

Fall 2019

Lab No. 3

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1) -77.125

a)

Single precision:

$77 = 38 \cdot 2 + 1$	$0.125 \cdot 2 = 0.25$
$38 = 19 \cdot 2 + 0$	$0.25 \cdot 2 = 0.5$
$19 = 9 \cdot 2 + 1$	$0.5 \cdot 2 = 1.0$
$9 = 4 \cdot 2 + 1$	
$4 = 2 \cdot 2 + 0$	$.125 = 001$
$2 = 1 \cdot 2 + 0$	
$1 = 0 \cdot 2 + 1$	

77 = 1001101 So, 77.125 = 1001101.001 in binary form

We convert the number into scientific notation.

$$1.001101001 \times 2^6$$

Mantissa = 001101001

Sign is 1 since it is negative. Sign = 1

Exponential is 6. But this is biased.

$$127 + 6 = 133$$

In binary form 133 = 1000 0101

In single precision format the number consists of;

sign(1), exponential(8), mantissa(23)

So -77.125 is;

$$1100\ 0010\ 1001\ 1010\ 0100\ 0000\ 0000\ 0000 = (\text{C2 9A 40 00})_{16}$$

Double precision:

We get the scientific notation of the binary number using the same method as before.

$$\text{So } 77.125 = 1.001101001 \times 2^6$$

The exponent bias is 1023 in double precision conversion and the exponent is 6.

Alas, we know that 1029 = 1000 0000 101 in binary. Sign is 1 since the value is negative.

Combining these we get;

1 1000 0000 101 0011 0100 1000 0000 0.....

The hexadecimal value would be = (C0 53 48 00 00 00 00 00)₁₆

b)

Single Precision:

If the bias for single precision was 120 then only the exponential part would change so the result would be

-77.125 = 1011 1111 0001 1010 0100 0000 0000 0000 = (DF 1A 40 00)₁₆

Double Precision:

With a bias of 1020 the binary would be as follows:

1 1000 0000 010 0011 0100 1000 0000 0000 0... = (C0 23 48 00 00 00 00 00)₁₆

c) C1 A0 00 00 = 1100 0001 1010 0000 0000 0000 0000

We know that the sign consists of 1 bit and exponential part is 8 bits and the rest (23 bits) is mantissa.

So;

Sign = 1 = negative

Exponential = 1000 0011

Mantissa = 0100 0000 0000 000....0

We have to get the bias removed from the exponent part which means we must convert the binary into decimal. 1000 0011 = 131, 131 - 127 = 4.

This means that, to get the scientific notation we jumped 4 digits to the left.

1.0100000... x 2⁴ = 10100.0000...

In this case, we can clearly see that we don't have a fractional part and 10100 = 20

C1 10 00 00 = -20

2)

Okan Sen

Start of recursiveSummation

#

la \$a0, number

move \$t0, \$a0

jal recursiveSummation

result comes in \$v0

move \$a1, \$a0

move \$t0, \$v0

la \$a0, msg1

li \$v0, 4

syscall

move \$a0, \$t0

li \$v0, 1

syscall

exit:

li \$v0, 10 # system call to exit

syscall # bye bye

stop execution here by syscall

recursiveSummation:

Basically, stores all chars in stacks,

when the last string char is reached, pop all the stacks

convert each char into integers and add them in another stack

a0 = string

if (string.length() == 1) return Value

#return RS(string-1)

save in stack

addi \$sp, \$sp, -12

```
sw $ra, 0($sp)
sw $s0, 4($sp)
sw $s1, 8($sp)
```

```
move $s0, $a0      # point to string
```

```
lb $t5, 0($a0) # get byte length
beq $t5, $zero, zeroVal      # if equals one length get the value
```

```
addi $a0, $s0, 1
```

```
jal recursiveSummation
```

```
add $s1, $zero, $v0      # $s1 = SR(str - 1)
```

```
jal returnVal
```

```
add $v0, $v0, $s1
```

exitRS:

```
lw  $ra, 0($sp)      # read registers from stack
lw  $s0, 4($sp)
lw  $s1, 8($sp)
addi $sp, $sp, 12     # bring back stack pointer
jr $ra
```

returnVal:

```
lbu $t1, ($s0)
addi $t1, $t1, -48
add $v0, $v0, $t1
j exitRS
```

zeroVal:

```
j exitRS
```

.data

number: .asciiz "1204"

msg1: .asciiz "Result is: "

```
endl:                .ascii "\n"
```

```
# -----
```

```
#
```

```
# End of recursiveSummation
```

```
#
```

```
# -----
```

3)

I have included which parts should be added to test the code, without errors.

Add to data part

```
msg1171: .ascii "7 - delete the nodes between a given value x \n"
```

```
msg7.1:   .ascii "Enter the value: "
```

```
msg7.2:   .ascii "Number of deleted nodes is: "
```

```
endl:     .ascii "\n"
```

Change if wanted(just copy paste on top of it, 7 is now 8):

```
msg118: .ascii "8 - exit this program \n"
```

COPY PASTE STARTING FROM T7

```
T7:      bne $s1,7, T8 # if s1 = 7, do these things. Else go to T8
```

```
la $a0,msg7.1 # put msg address into a0
```

```
li $v0,4      # system call to print
```

```
syscall
```

```
li $v0,5 # system call to read
```

```
syscall      # in the integer
```

```
move $a1, $v0      # put x value into a1 before the call
```

```
move $a0, $s0      # put pointer to linked list in a0 before the call
```

```
li $v0, 0
```

```
jal deleteAfter_x
```

```
move $t0, $v0
```

```
la $a0, msg7.2
```

```
li $v0, 4
```

```
syscall
```

```
move $a0, $t0      # then put its value in a0 to print it out
```

```
li $v0, 1
```

```
syscall
```

```
la $a0, endl
```

```
li $v0, 4
```

```
syscall
```

```
la $a0, endl
```

```
li $v0, 4
```

```
syscall
```

```
T8:   bne $s1,8, T8no    # if s1 = 8, do these things. Else go to T8no
```

```
la $a0,msg127      # put msg127 address into a0
```

```
li $v0,4# system call to print
```

```
syscall      # out the thank you string
```

```
li $v0,10
```

```
# the exit syscall is 10
```

```
syscall      # goodbye...
```

```
T8no:
```

```
la $a0,msg128      # put msg128 address into a0
```

```
li $v0,4# system call to print
```

```
syscall           # out the msg128 string
```

```
j EnterChoice # go to the place to enter the choice
```

COPY PASTE INTO METHODS PART

```
deleteAfter_x:
```

```
li $v0, 0  
move $t0, $a0  
lw $t1, 4($t0)
```

```
DAnext:      beq $t1, $a1, DАfirst1      # if the input value is seen get to the  
first label
```

```
lw $t0, 0($t0) # goes to next node  
beq $t0, $zero, DАdone      # we might reach the end of the list
```

```
lw $t1, 4($t0) # loads the value in t1
```

```
j DАnext
```

```
DАfirst1:
```

```
lw $t2, 0($t0) # get next in t2  
beq $t2, $zero, DАdone      # we might reach the end of the list  
lw $t1, 4($t2) # get current value
```

```
DАfirst2:
```

```
beq $t1, $a1, DАokay # if we see the second occurrence of the input  
value, we can proceed with the process
```

```
lw $t2, 0($t2)      # get to next
```


beq \$t2, \$zero, DAdone # there might not be a second occurrence
and if we reach the end of the list finish method

lw \$t1, 4(\$t2) # get current value
j DAfirst2

DAokay:

lw \$t2, 0(\$t0) # get next in t2
lw \$t1, 4(\$t2) # get next value

beq \$t1, \$a1, DALdone # if we still don't see the second occurrence
of X, we will continue deleting nodes until we see it

lw \$t3, 0(\$t2) # next's next is t3
sw \$t3, 0(\$t0) # set current's next t3

li \$t2, 0
addi \$v0, \$v0, 1

j DAokay

DALdone:

move \$t0, \$t2 # t0 now points to next
lw \$t0, 0(\$t0) # and now t0 points to next's next for another run in case
we have more occurrences of X
beq \$t0, \$zero, DAdone # we might reach the end of the list
lw \$t1, 4(\$t0)

j DAnext

DAdone:

li \$t0, 0
li \$t1, 0
li \$t2, 0
li \$t3, 0

jr \$ra

I was not able to get the deleted node back to the heap. In other words, I can only make the nodes to point to necessary nodes after deletion but I can never delete a node. They actually stay in the heap in random memories, still having unnecessary data. In the long run, the program would have moderate memory leaks. The problem here is; there is no proper way to delete the node from the heap. Otherwise, the program works perfectly as intended. Deleting all the nodes between given X value, and with multiple occurrences; such as 2, 7, 4, 8, 4, 9, 4, 10, 15, 4, 7 will return when X is 4; 2, 7, 4, 4, 9, 4, 4, 7.