CSC 480: Artificial Intelligence

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Logistics Project

- Moodle: Artificial Intelligence (Kurfess)
 - * enrollment key: Al-MUAS-S15
 - groups being set up for project teams
 - project description
 - project overview
 - tentative schedule
 - * features, requirements, evaluation criteria





Logistics

Lab and Homework Assignments

- Lab 2 due tonight (23:59); demos during the lab time
- Lab 3: implementation of breadth-first, depth-first search

Quizzes

- * Quiz 2
 - Available Tue, April 14, all day (0:00 23:59)
 - longer availability of quizzes?





Chapter Overview Search

- Motivation
- Objectives
- Search as Problem-Solving
 - problem formulation
 - problem types
- Uninformed Search
 - breadth-first
 - depth-first
 - uniform-cost search
 - depth-limited search
 - iterative deepening
 - bi-directional search

Informed Search

- best-first search
- search with heuristics
- memory-bounded search
- iterative improvement search

Non-Traditional Search

- local search and optimization
- constraint satisfaction
- search in continuous spaces
- partially observable worlds

Important Concepts and Terms

Chapter Summary





Motivation and Objectives

Pre-Test

Motivation

Objectives

Evaluation Criteria





Motivation

- search strategies are important methods for many approaches to problem-solving
- the use of search requires an abstract formulation of the problem and the available steps to construct solutions
- search algorithms are the basis for many optimization and planning methods





Objectives

- formulate appropriate problems as search tasks
 - states, initial state, goal state, successor functions (operators), cost
- know the fundamental search strategies and algorithms
 - uninformed search
 - breadth-first, depth-first, uniform-cost, iterative deepening, bi-directional
 - informed search
 - best-first (greedy, A*), heuristics, memory-bounded, iterative improvement
- evaluate the suitability of a search strategy for a problem
 - completeness, time & space complexity, optimality





Background

What is Search in AI?
Problem-Solving and Search
Terminology





Search as Problem-Solving Strategy

- many problems can be viewed as reaching a goal state from a given starting point
 - often there is an underlying state space that defines the problem and its possible solutions in a more formal way
 - the space can be traversed by applying a successor function (operators) to proceed from one state to the next
 - if possible, information about the specific problem or the general domain is used to improve the search
 - experience from previous instances of the problem
 - strategies expressed as heuristics
 - simpler versions of the problem
 - constraints on certain aspects of the problem





Examples

getting from home to Cal Poly

- start: home on Clearview Lane
- goal: Cal Poly CSC Dept.
- operators: move one block, turn

loading a moving truck

- start: apartment full of boxes and furniture
- goal: empty apartment, all boxes and furniture in the truck
- operators: select item, carry item from apartment to truck, load item

getting settled

- start: items randomly distributed over the place
- goal: satisfactory arrangement of items
- operators: select item, move item





Problem-Solving Agents

- agents whose task it is to solve a particular problem
 - goal formulation
 - what is the goal state
 - what are important characteristics of the goal state
 - how does the agent know that it has reached the goal
 - are there several possible goal states
 - are they equal or are some more preferable
 - problem formulation
 - what are the possible states of the world relevant for solving the problem
 - what information is accessible to the agent
 - how can the agent progress from state to state





Problem Formulation

formal specification for the task of the agent

- goal specification
- states of the world
- actions of the agent

identify the type of the problem

- what knowledge does the agent have about the state of the world and the consequences of its own actions
- does the execution of the task require up-to-date information
 - sensing is necessary during the execution





Well-Defined Problems

- problems with a readily available formal specification
 - initial state
 - starting point from which the agent sets out
 - actions (operators, successor functions)
 - describe the set of possible actions
 - state space
 - set of all states reachable from the initial state by any sequence of actions
 - path
 - sequence of actions leading from one state in the state space to another
 - goal test
 - determines if a given state is the goal state





Well-Defined Problems (cont.)

- solution
 - path from the initial state to a goal state
- search cost
 - time and memory required to calculate a solution
- path cost
 - determines the expenses of the agent for executing the actions in a path
 - sum of the costs of the individual actions in a path
- total cost
 - sum of search cost and path cost
 - overall cost for finding a solution





Selecting States and Actions

- states describe distinguishable points or periods during the problem-solving process
 - dependent on the task and domain
- actions move the agent from one state to another one
 - an operator is applied to the initial state and takes the agent to the successor state
 - dependent on states, capabilities of the agent, and properties of the environment
- choice of suitable states and operators
 - can make the difference between a problem that can or cannot be solved (in principle, or in practice)





Example: Commute Home => MUAS

states

- locations:
 - obvious: buildings that contain your home, MUAS Computer Science building
 - more difficult: intermediate states
 - public transit: stops (bus, street car, U-Bahn, S-Bahn)
 - walking: ??
 - bicycling: ??
 - car: ??
- environment-centric states
 - e.g., public transit stops
- agent-centric states
 - * moving, turning, resting, ...

operators

- depend on the choice of states
- e.g. move_one_block
- abstraction is necessary to omit irrelevant details
 - valid: can be expanded into a detailed version
 - useful: easier to solve than in the detailed version





Search Example: Traveling by Car

- Route finding: Starting Point => Goal Point
 - focus on
 - communities, landmarks
 - highways, streets, intersections
 - street addresses
 - constraints
 - one-way, avoid highways/ferries/toll roads, scenic vs. direct, ...
 - traffic considerations





Search Example: Traveling by Train

- Route finding: Starting Station => Goal Station
 - focus on railway network
 - train stations
 - constraints
 - train type: ICE vs. regional
 - passenger special needs: bicycle transport
 - traffic considerations





Example Problems for Search

Toy ProblemsReal-World Problems





Example Problems

toy problems

- vacuum world
- 8-puzzle
- 8-queens
- cryptarithmetic
- vacuum agent
- missionaries and cannibals

- traveling salesperson
- VLSI layout
- robot navigation
- assembly sequencing
- Web search

real-world problems

- route finding
- touring problems







Simple Vacuum World

- states
 - two locations
 - dirty, clean
- initial state
 - any legitimate state
- successor function (operators)
 - left, right, suck
- goal test
 - all squares clean
- path cost
 - one unit per action

Properties: discrete locations, discrete dirt (binary), deterministic OCHSO

More Complex Vacuum Agent

states

- configuration of the room
 - dimensions, obstacles, dirtiness

initial state

- locations of agent, dirt
- successor function (operators)
 - move, turn, suck
- goal test
 - all squares clean
- path cost
 - one unit per action

Properties: discrete locations, discrete dirt, deterministic, $d * 2^n$ states for dirt degree d,n locations





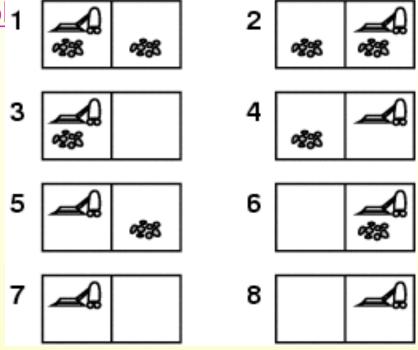
Vacuum World Example

- from the AIMA textbook slides
- conversion from PowerPoint to Keynote





Single-state, start in #5. So 1

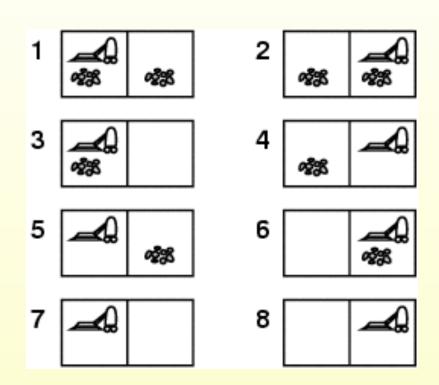






Single-state, start in #5.
Solution? [Right, Suck]

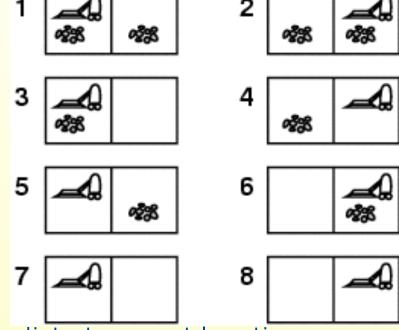
Sensorless, start in {1,2,3,4,5,6,7,8} e.g., Right goes to {2,4,6,8} Solution?







Sensorless, start in {1,2,3,4,5,6,7,8} e.g., Right goes to {2,4,6,8} Solution? [Right,Suck,Left,Suck]



Contingency

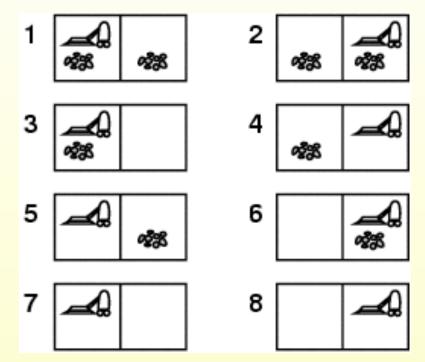
- Nondeterministic: Suck may dirty a clean carpet
- Partially observable: location, dirt at current location.
- Percept: [L, Clean], i.e., start in #5 or #7 Solution?



Sensorless, start in {1,2,3,4,5,6,7,8} e.g., Right goes to {2,4,6,8} Solution? [Right,Suck,Left,Suck]



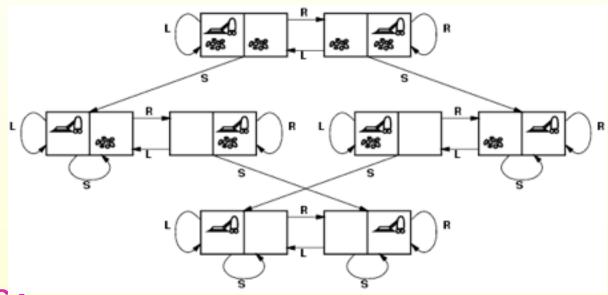
- Nondeterministic: Suck may dirty a clean carpet
- Partially observable: location, dirt at current location.
- Percept: [L, Clean], i.e., start in #5 or #7
 Solution? [Right, if dirt then Suck]



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Vacuum world state space graph



- * states:
- * actions?
- * goal test?
- * path cost?





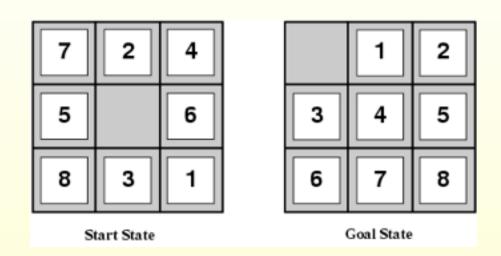
Vacuum world state space graph

stat
acti
goa
path





Example: The 8-puzzle



- states?
- actions?
- goal test?
- path cost?

• [Note: optimal solution of n-Puzzle family is NP-hard]





8-Puzzle

states

- location of tiles (including blank tile)
- properties of tiles (numbers, letters, image parts, ...)

initial state

any legitimate configuration

successor function (operators)

- move tile
- alternatively: move blank

goal test

- specific legitimate configuration of tiles
 - particular arrangement of tiles
 - may reflect ordering relations, words, image

path cost

one unit per move

Properties: abstraction leads to discrete configurations, discrete moves.

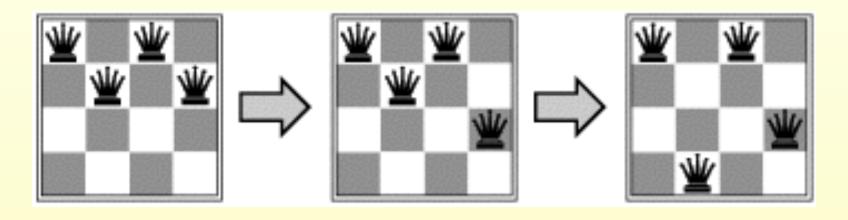
- deterministic
- 9!/2 = 181,440 reachable states





Example: n-queens

- Put n queens on a n x n board with no two queens on the same row, column, or diagonal
 - see also <u>Wikipedia Eight Queens Puzzle</u>







8-Queens Incremental Approach

- start with an empty board
- add queens one by one
 - no violation of constraints
 - different row, column, diagonal from sitting queens
- incremental formulation
 - states
 - arrangement of up to 8 queens on the board
 - initial state
 - empty board
 - successor function (operators)
 - add a queen to any square
 - goal test
 - all queens on board
 - no queen attacked
 - path cost
 - irrelevant (all solutions equally valid)
- Properties: 3*1014 possible sequences; can be reduced to 2,057



8-Queens Complete-State Approach

- start with a full board
 - all n queens placed on the board
 - conflicts are to be expected
- try to find a better configuration
 - reduced number of conflicts
- complete-state formulation
 - states
 - arrangement of 8 queens on the board
 - initial state
 - all 8 queens on board
 - successor function (operators)
 - move a queen to a different square
 - goal test
 - no queen attacked
 - path cost
 - irrelevant (all solutions equally valid)
- Properties: good strategies can reduce the number of possible sequences considerably



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8-Queens Refined

- simple solutions may lead to very high search costs
 - 64 fields, 8 queens ==> 648 possible sequences
- more refined solutions trim the search space, but may introduce other constraints
 - place queens on "unattacked" places
 - much more efficient
 - may not lead to a solution depending on the initial moves
 - move an attacked queen to another square in the same column, if possible to an "unattacked" square
 - much more efficient





Crypt-arithmetic

- states
 - puzzle with letters and digits
- initial state
 - only letters present



- successor function (operators)
 - replace all occurrences of a letter by a digit not us
- goal test
 - only digits in the puzzle
 - calculation is correct
- path cost
 - all solutions are equally valid







Missionaries and Cannibals

- states
 - number of missionaries, cannibals, and boats on the banks of a river
 - illegal states
 - missionaries are outnumbered by cannibals on either bank
- initial states
 - all missionaries, cannibals, and boats are on one bank
- successor function (operators)
 - transport a set of up to two participants to the other bank
 - {1 missionary} | { 1cannibal} | {2 missionaries} | {2 cannibals} | {1 missionary and 1 cannibal}
- goal test
 - nobody left on the initial river bank
- path cost
 - number of crossings
- also known as "goats and cabbage", "wolves and sheep", etc.





Route Finding

- states
 - locations
- initial state
 - starting point
- successor function (operators)
 - move from one location to another
- goal test
 - arrive at a certain location
- path cost
 - may be quite complex
 - money, time, travel comfort, scenery, ...





Traveling Salesperson

states

- locations / cities
- illegal states
 - each city may be visited only once
 - visited cities must be kept as state information

initial state

- starting point
- no cities visited

successor function (operators)

move from one location to another one

goal test

- all locations visited
- agent at the initial location

path cost

distance between locations





VLSI Layout

states

positions of components, wires on a chip

initial state

- incremental: no components placed
- complete-state: all components placed (e.g. randomly, manually)

successor function (operators)

- incremental: place components, route wire
- complete-state: move component, move wire

goal test

- all components placed
- components connected as specified

path cost

- may be complex
 - distance, capacity, number of connections per component





Robot Navigation

states

- locations
- position of actuators

initial state

- start position (dependent on the task)
- successor function (operators)
 - movement, actions of actuators
- goal test
 - task-dependent
- path cost
 - may be very complex
 - distance, energy consumption





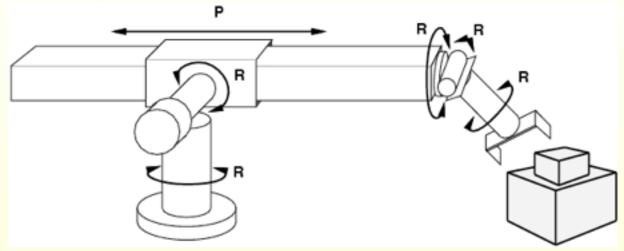
Assembly Sequencing

- states
 - location of components
- initial state
 - no components assembled
- successor function (operators)
 - place component
- goal test
 - system fully assembled
- path cost
 - number of moves





Example: robotic assembly



- states?: real-valued coordinates of robot joint angles parts of the object to be assembled
- actions?: continuous motions of robot joints
- goal test?: complete assembly
- path cost?: time to execute





Search Methods

Terminology
Search and Graphs





Searching for Solutions

traversal of the search space

- from the initial state to a goal state
- legal sequence of actions as defined by successor function (operators)

general procedure

- check for goal state
- expand the current state
 - determine the set of reachable states
 - return "failure" if the set is empty
- select one from the set of reachable states
- move to the selected state

a search tree is generated

nodes are added as more states are visited





Search Terminology

search tree

- generated as the search space is traversed
 - the search space itself is not necessarily a tree, frequently it is a graph
 - the tree specifies possible paths through the search space
- expansion of nodes
 - as states are explored, the corresponding nodes are expanded by applying the successor function
 - this generates a new set of (child) nodes
 - the fringe (frontier) is the set of nodes not yet visited
 - newly generated nodes are added to the fringe
- search strategy
 - determines the selection of the next node to be expanded
 - can be achieved by ordering the nodes in the fringe
 - e.g. queue (FIFO), stack (LIFO), "best" node w.r.t. some measure (cost)

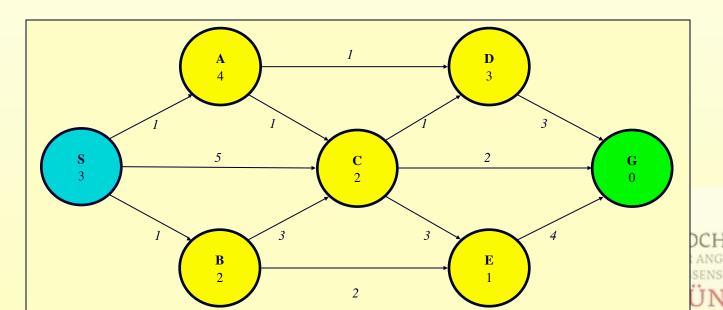




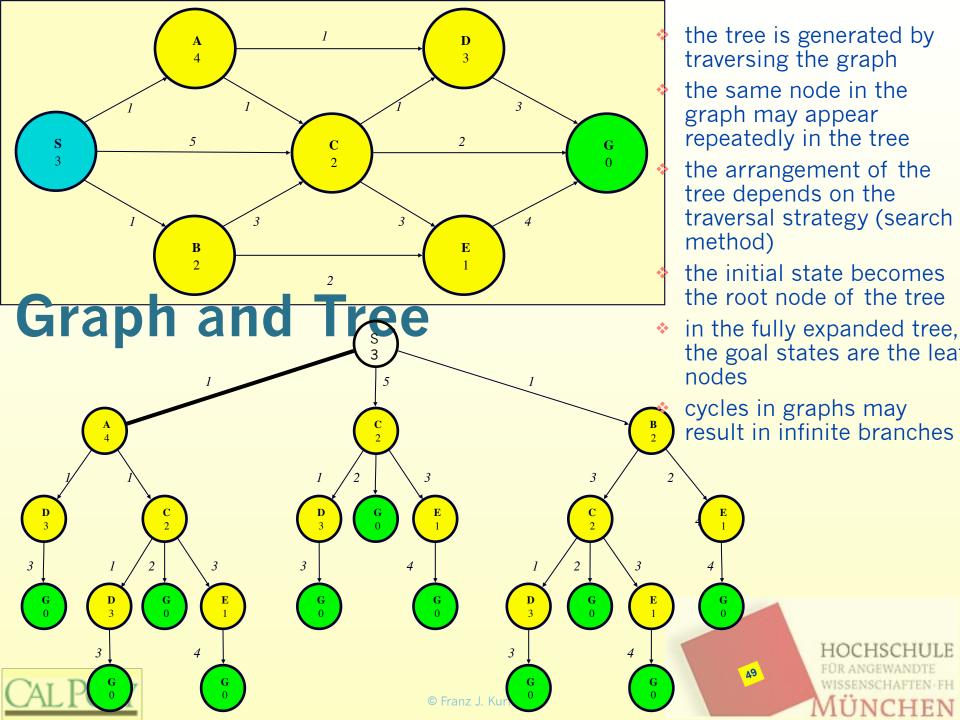
Example: Graph Search

describes the search (state) space

- each node represents one state in the search space
 - e.g. a city to be visited in a routing or touring problem
- additional information
 - names and properties for the states (e.g. S, 3)
 - links between nodes, specified by the successor function
 - properties for links (distance, cost, name, ...)







General Tree Search Algorithm

- generate the first node from the initial state of the problem
- repeat
 - return failure if there are no more nodes in the fringe
 - examine the current node; if it's a goal, return the solution
 - expand the current node, and add the new nodes to the fringe

Note: This method is called "General-Search" in earlier AIMA editions





General Tree Search Algorithm

```
function TREE-SEARCH(problem, fringe) returns solution
  fringe := INSERT(MAKE-NODE(INITIAL-STATE[problem]),
  fringe)
   loop do
       if EMPTY?(fringe) then return failure
       node := REMOVE-FIRST(fringe)
       if GOAL-TEST[problem] applied to STATE[node]
  succeeds
          then return SOLUTION(node)
       fringe := INSERT-ALL(EXPAND(node, problem),
  fringe)
```





Evaluation Criteria

completeness

if there is a solution, will it be found

optimality

the best solution will be found

time complexity

- time it takes to find the solution
- does not include the time to perform actions

space complexity

- memory required for the search
- main factors for complexity considerations:
 - branching factor b
 - depth d of the shallowest goal node
 - maximum path length m





Search Cost and Path Cost

- the search cost indicates how expensive it is to generate a solution
 - time complexity (e.g. number of nodes generated) is usually the main factor
 - sometimes space complexity (memory usage) is considered as well
- path cost indicates how expensive it is to execute the solution found in the search
 - distinct from the search cost, but often related
- total cost is the sum of search and path costs





Search Strategies

Uninformed
Informed
Local
Others





Selection of a Search Strategy

- selection of an appropriate search strategy for a given problem
 - uninformed search (blind search)
 - number of steps, path cost unknown
 - agent knows it is at a goal only after it reaches a goal
 - goals are not "visible" from a distance
 - informed search (heuristic search)
 - agent has background information about the problem
 - map, costs of actions
 - hints about the location of the goal
 - evaluating hints can be costly as well





Search Strategies

Uninformed Search

- breadth-first
 - uniform-cost search
- depth-first
 - depth-limited search
 - iterative deepening
- bi-directional search

Informed Search

- best-first search
- search with heuristics
- memory-bounded search
- iterative improvement search

Local Search and Optimization

- hill-climbing
- simulated annealing
- local beam search
- genetic algorithms
- constraint satisfaction
- Search in Continuous Spaces
- Non-deterministic Actions

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- Partial Observations
- Online Search





Breadth-First

plain breadth-first uniform cost





Breadth-First

- all the nodes reachable from the current node are explored first
 - achieved by the TREE-SEARCH method by appending newly generated nodes at the end of the search queue

function BREADTH-FIRST-SEARCH(problem) returns solution

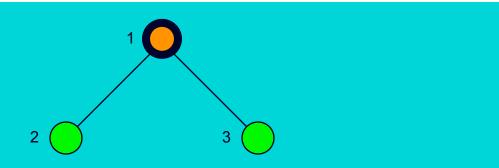
return TREE-SEARCH(problem, FIFO-QUEUE())

Time Complexity	b ^{d+1}
Space Complexity	b ^{d+1}
Completeness	yes (for finite b)
Optimality	yes (for non-negative path costs)

b branching factor
d depth of the tree

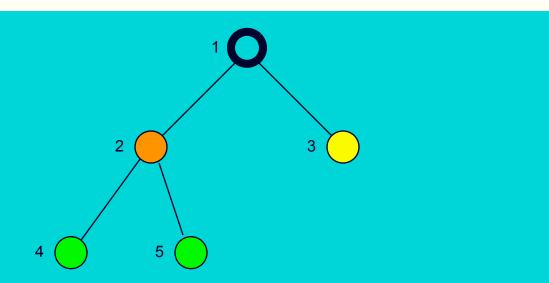






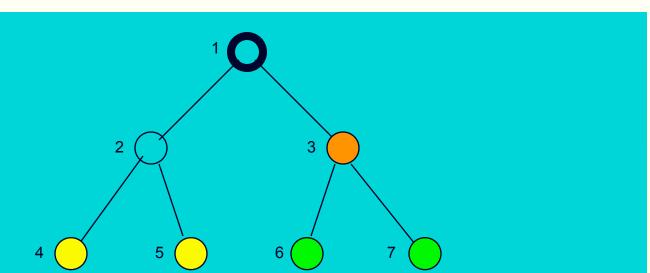




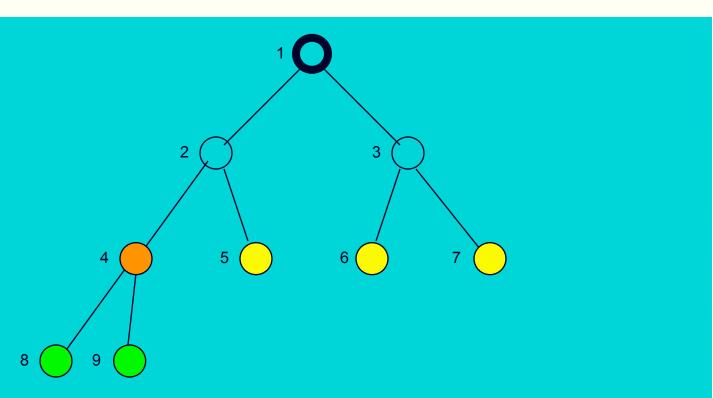






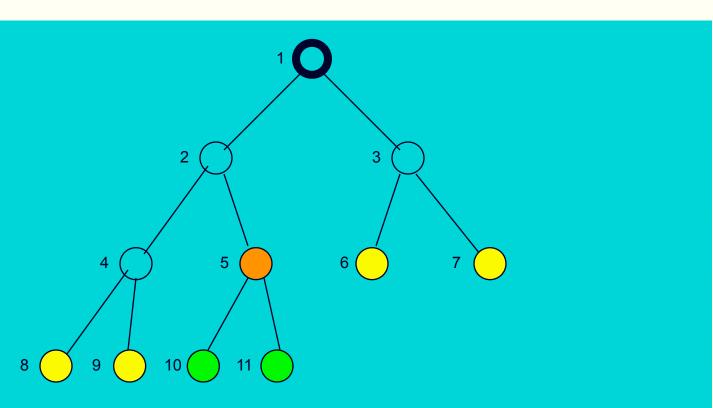






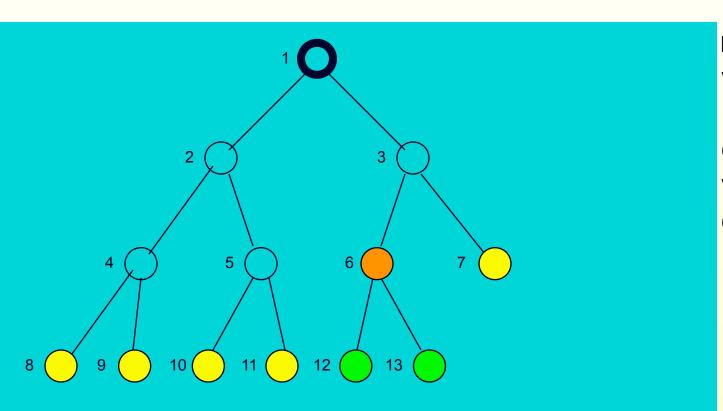






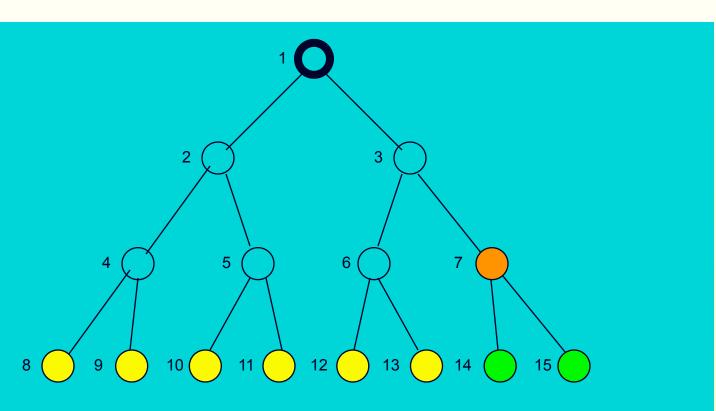






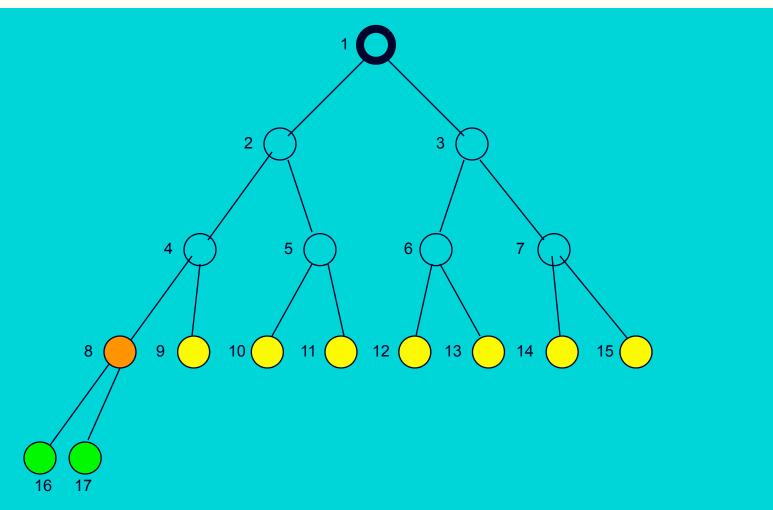






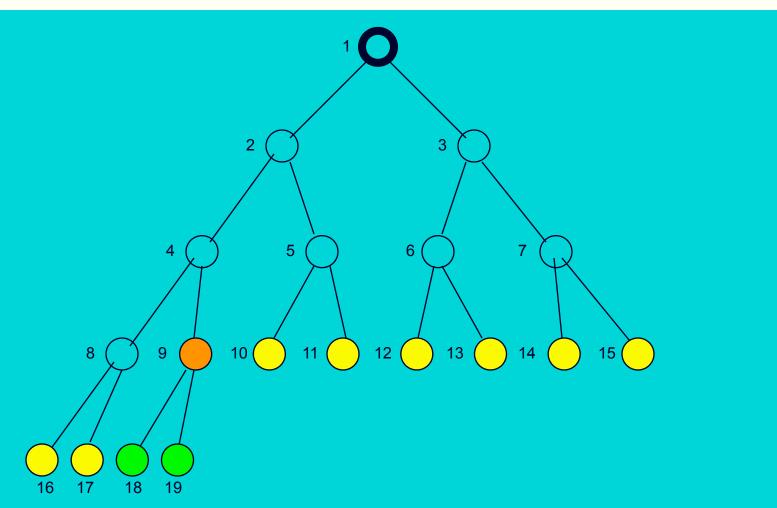




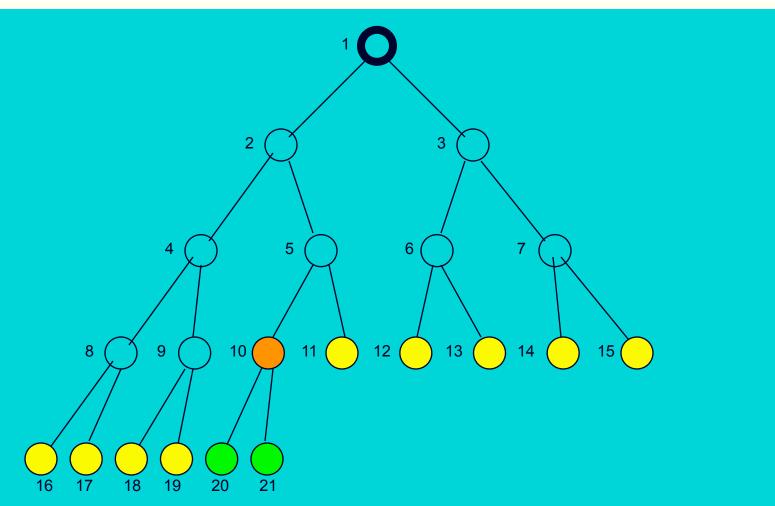




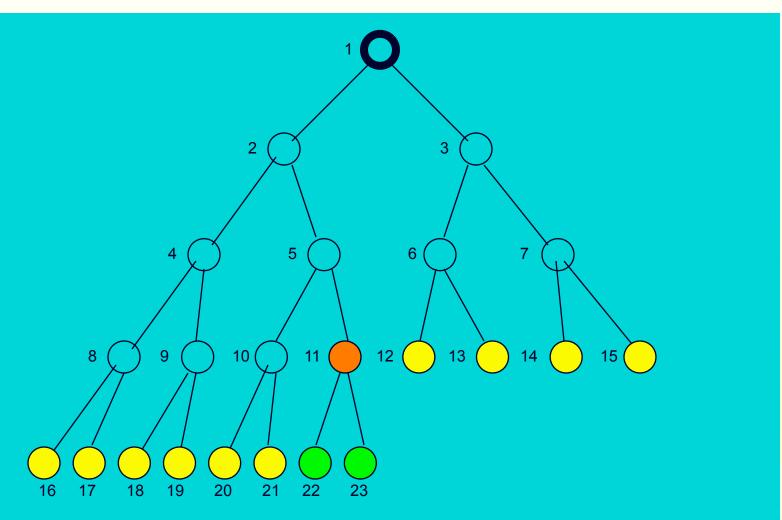




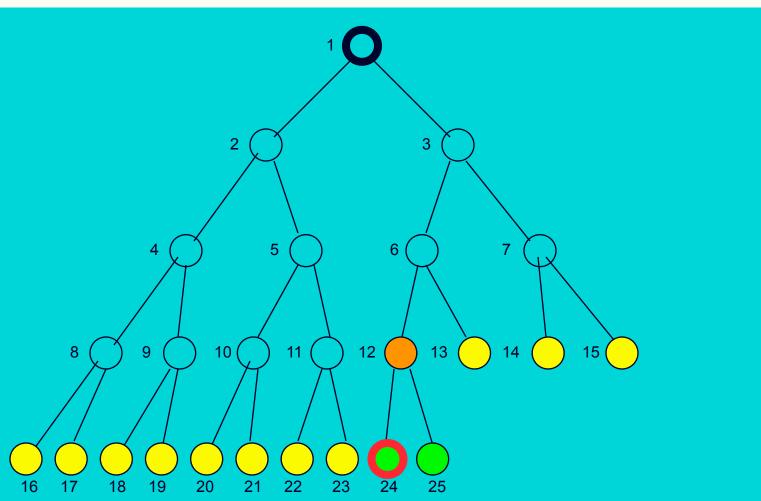








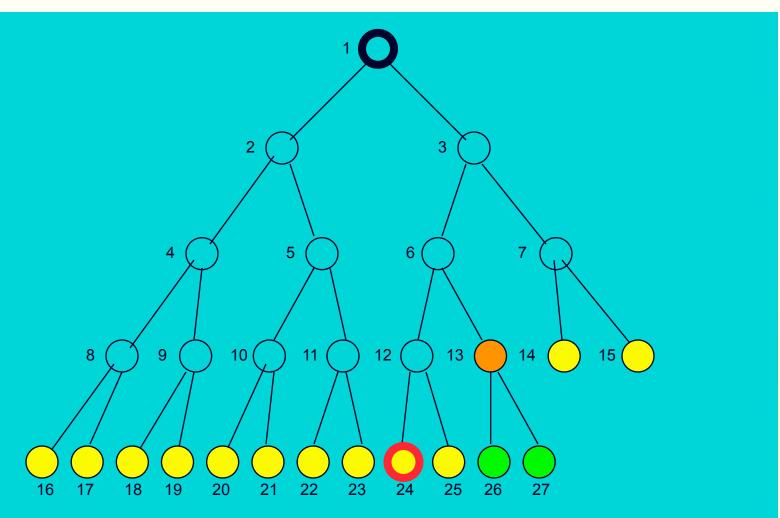




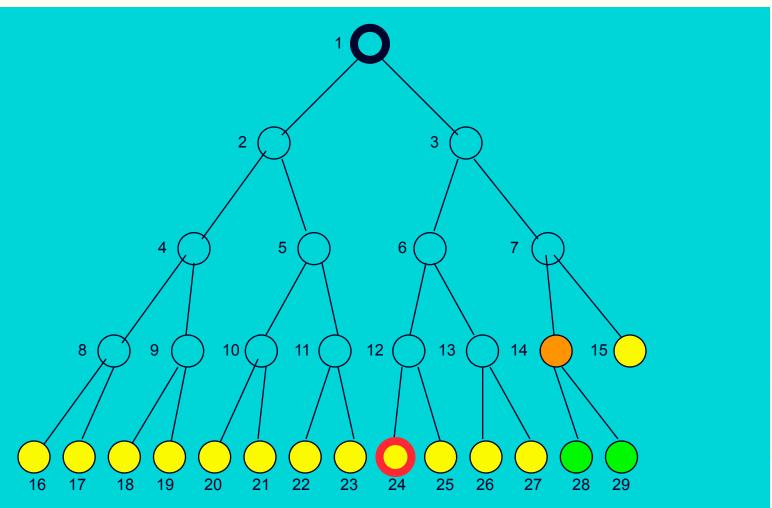
Initial
Visited
Fringe
Current
Visible
Goal

Note:
The goal node
is "visible"
here, but we
can not
perform the
goal test yet.

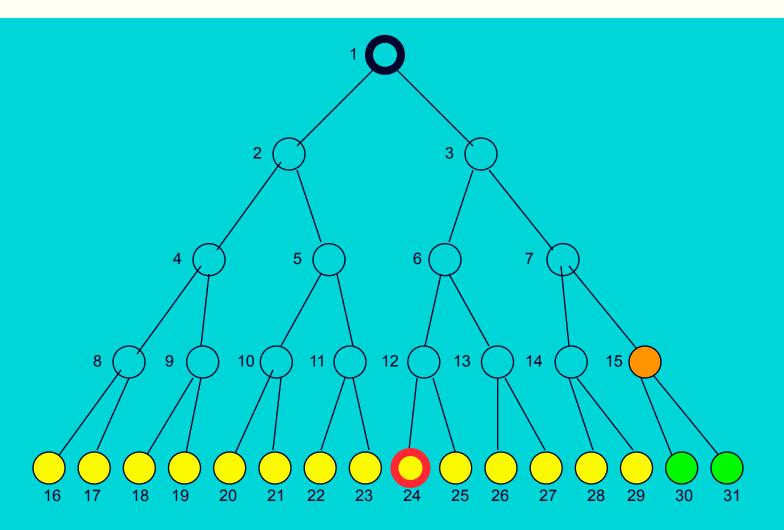






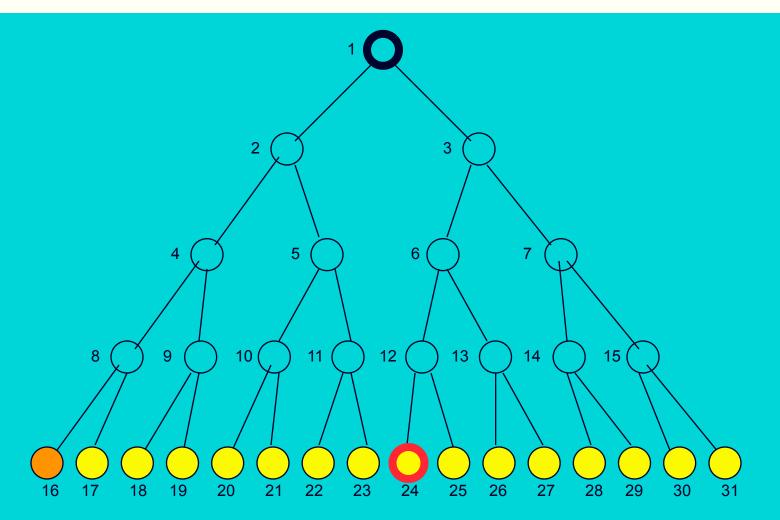






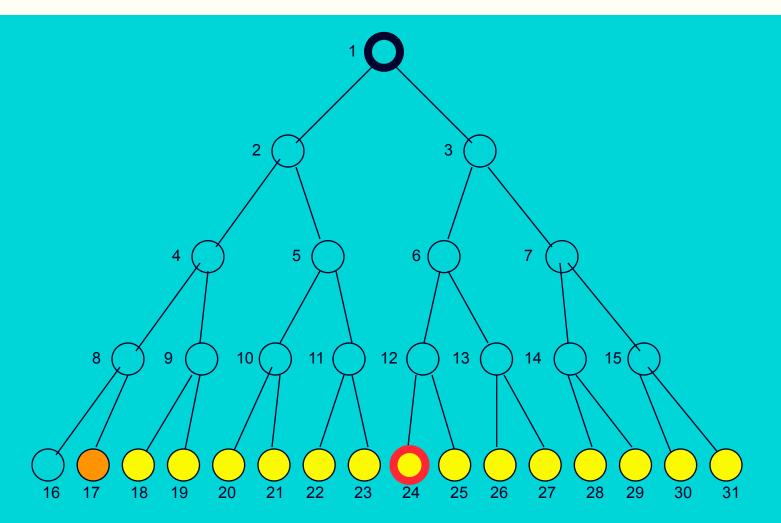
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Initial
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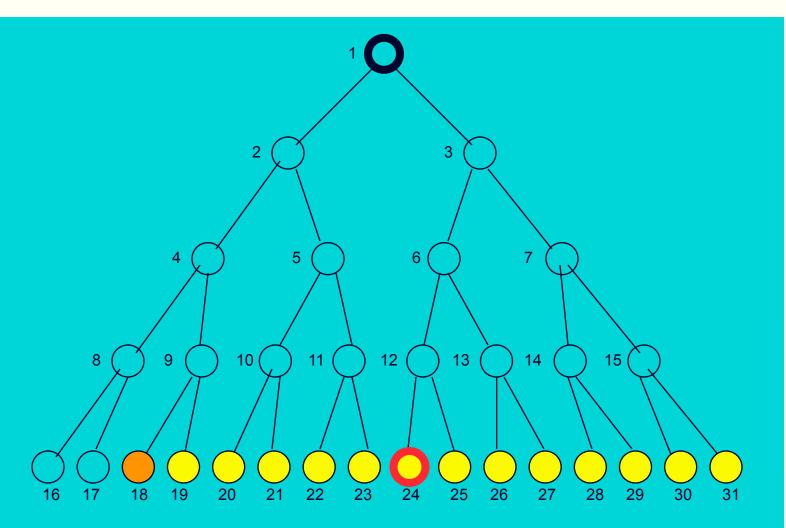




Initial
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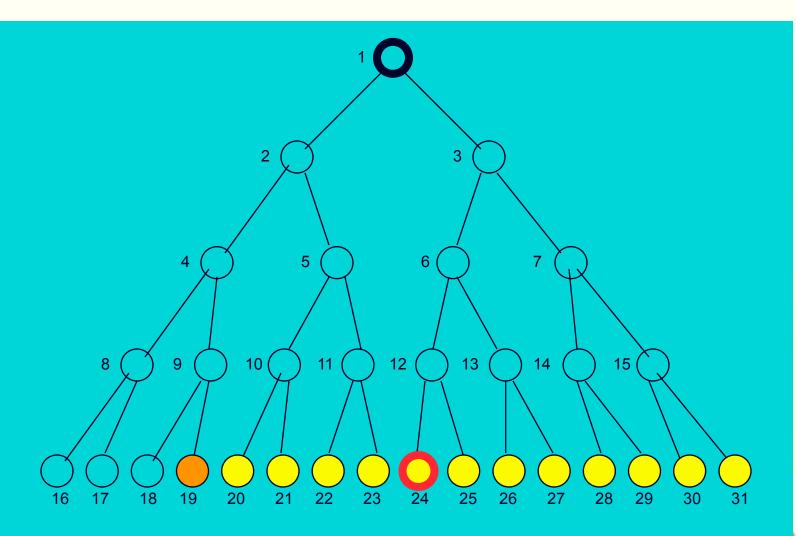


Fringe: [18,19,20,21,22,23,24,25,26,27,28,29,30,319]



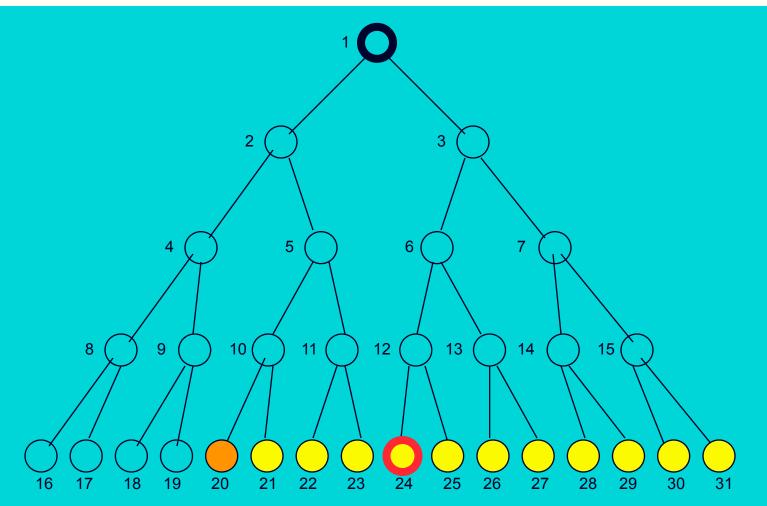
Initial
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Fringe
Current
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Goal





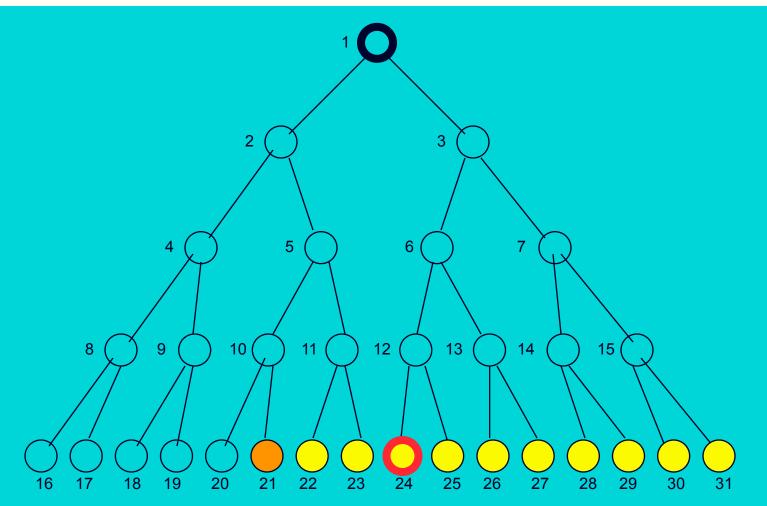
Initial Visited Fringe Current Visible Goal





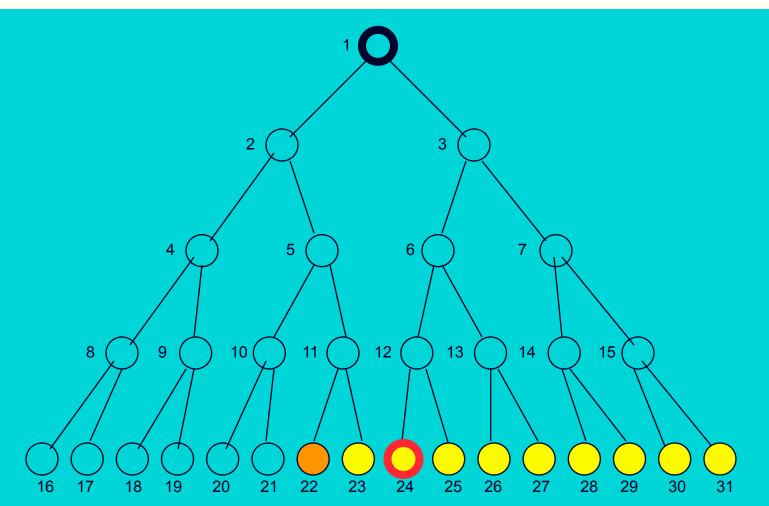
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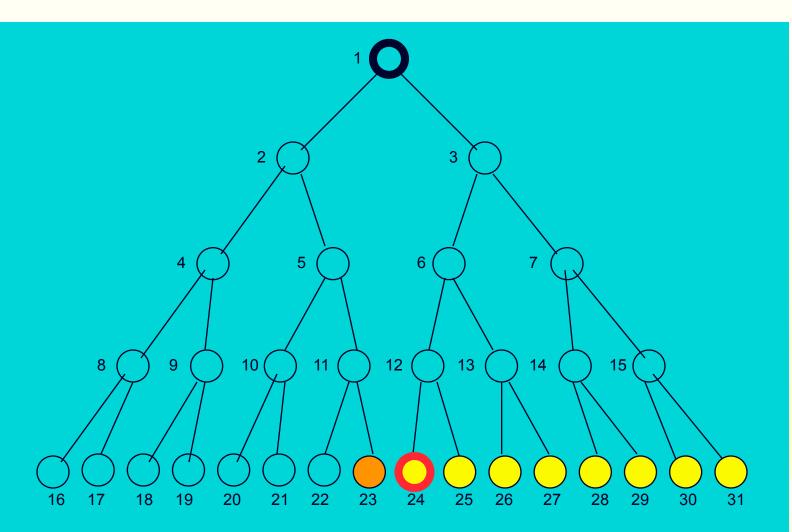
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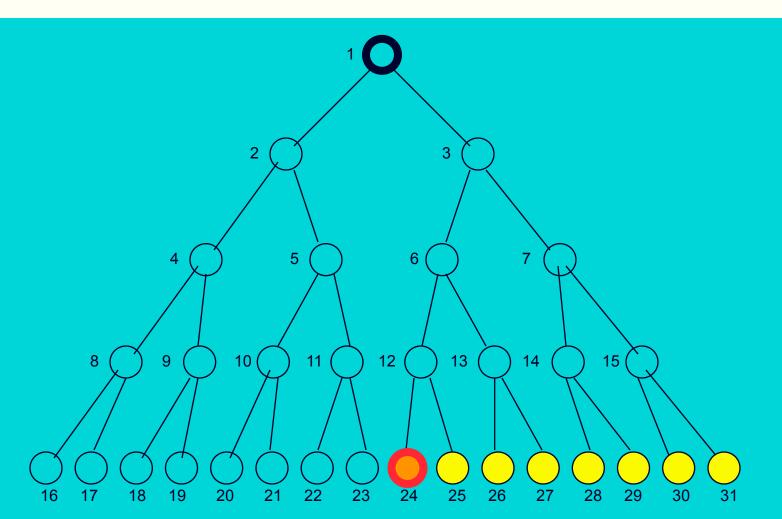
Initial
Visited
Fringe
Current
Visible
Goal





Initial
Visited
Fringe
Current
Visible
Goal





Initial
Visited
Fringe
Current
Visible
Goal

Note: The goal test is positive for this node, and a solution is found in 24 steps.



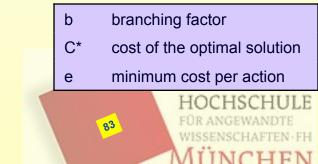
Uniform-Cost -First

- the nodes with the lowest cost are explored first
 - similar to BREADTH-FIRST, but with an evaluation of the cost for each reachable node
 - g(n) = path cost(n) = sum of individual edge costs to reach the current node

function UNIFORM-COST-SEARCH(problem) returns solution

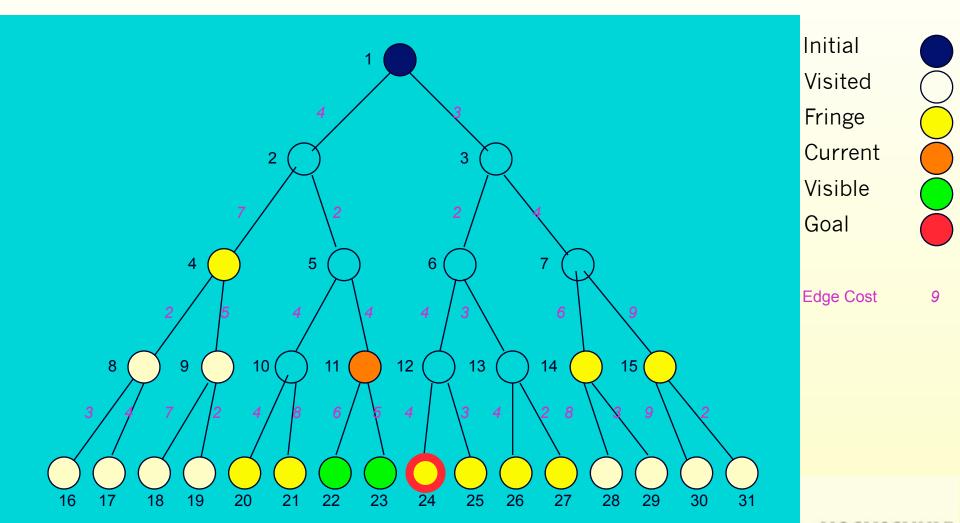
return TREE-SEARCH(problem, COST-FN, FIFO-QUEUE())

Time Complexity	b ^{C*/e}
Space Complexity	b ^{C*/e}
Completeness	yes (finite b, step costs >= e)
Optimality	yes





Uniform-Cost Snapshot



HOCHSCHULE FOR ANGEWANDTE WISSENSCHAFTEN-FH

Fringe: [27(10), 4(11), 25(12), 26(12), 14(13), 24(13), 20(14), 15(16), 21(18)] + [22(16), 23(15)]

Uniform Cost Fringe Trace

- 1. [1(0)]
- 2. [3(3), 2(4)]
- 3. **[2(4), 6(5), 7(7)]**
- **4**. **[6(5)**, **5(6)**, **7(7)**, **4(11)**]
- **5.** [5(6), 7(7), 13(8), 12(9), 4(11)]
- 6. [7(7), 13(8), 12(9), 10(10), 11(10), 4(11)]
- **7.** [13(8), 12(9), 10(10), 11(10), 4(11), 14(13), 15(16)]
- 8. [12(9), 10(10), 11(10), 27(10), 4(11), 26(12), 14(13), 15(16)]
- 9. [10(10), 11(10), 27(10), 4(11), 26(12), 25(12), 14(13), 24(13), 15(16)]
- 10. [11(10), 27(10), 4(11), 25(12), 26(12), 14(13), 24(13), 20(14), 15(16), 21(18)]
- 11. [27(10), 4(11), 25(12), 26(12), 14(13), 24(13), 20(14), 23(15), 15(16), 22(16), 21(18)]
- 12. [4(11), 25(12), 26(12), 14(13), 24(13), 20(14), 23(15), 15(16), 23(16), 21(18)]
- 13. [25(12), 26(12), 14(13), 24(13),8(13), 20(14), 23(15), 15(16), 23(16), 9(16), 21(18)]
- 14. [26(12), 14(13), 24(13),8(13), 20(14), 23(15), 15(16), 23(16), 9(16), 21(18)]
- 15. [14(13), 24(13),8(13), 20(14), 23(15), 15(16), 23(16), 9(16), 21(18)]
- 16. [24(13),8(13), 20(14), 23(15), 15(16), 23(16), 9(16), 29(16),21(18), 28(21)]

Goal reached!

Notation: [Bold+Yellow: Current Node; White: Old Fringe Node; Green+Italics: New Fringe Node].

Assumption: New nodes with the same cost as existing nodes are added after the existing node.





Breadth-First vs. Uniform-Cost

- breadth-first always expands the shallowest node
 - only optimal if all step costs are equal
- uniform-cost considers the overall path cost
 - optimal for any (reasonable) cost function
 - non-zero, positive
 - gets bogged down in trees with many fruitless, short branches
 - low path cost, but no goal node
- both are complete for non-extreme problems
 - finite number of branches
 - strictly positive cost function





Review: Breadth-First vs. Uniform-Cost

basic idea:

- breadth-first
- uniform-cost

properties

- completeness
- optimality
- time complexity
- space complexity





Depth-First

plain depth-first limited depth iterative deepening





Depth-First

- continues exploring newly generated nodes
 - achieved by the TREE-SEARCH method by appending newly generated nodes at the beginning of the search queue
 - utilizes a Last-In, First-Out (LIFO) queue, or stack

function DEPTH-FIRST-SEARCH(problem) returns solution

return TREE-SEARCH(problem, LIFO-QUEUE())

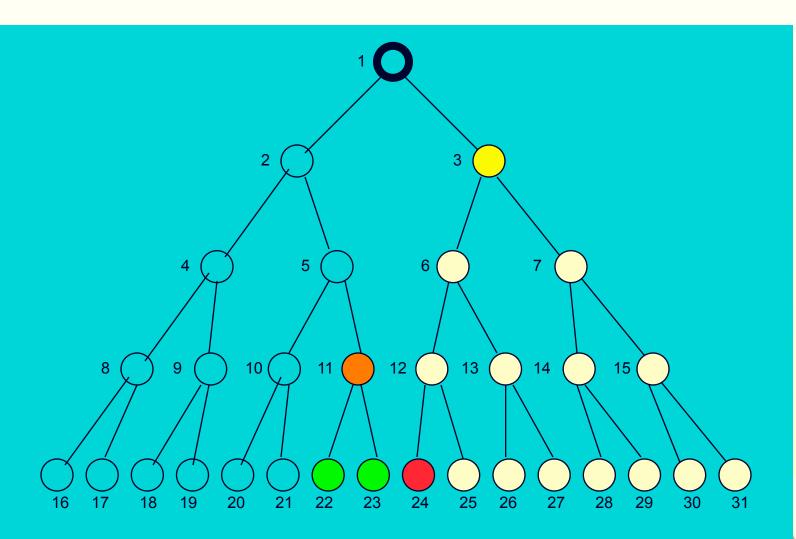
Time Complexity	b ^m
Space Complexity	b*m
Completeness	no (for infinite branch length)
Optimality	no

b branching factor
m maximum path length





Depth-First Snapshot



Initial
Visited
Fringe
Current
Visible
Goal



Depth-First vs. Breadth-First

- depth-first goes off into one branch until it reaches a leaf node
 - not good if the goal is on another branch
 - neither complete nor optimal
 - uses much less space than breadth-first
 - much fewer visited nodes to keep track of
 - smaller fringe
- breadth-first is more careful by checking all alternatives
 - complete and optimal
 - under most circumstances
 - very memory-intensive





Backtracking Search

- variation of depth-first search
 - only one successor node is generated at a time
 - even better space complexity: O(m) instead of O(b*m)
 - even more memory space can be saved by incrementally modifying the current state, instead of creating a new one
 - only possible if the modifications can be undone
 - * this is referred to as backtracking
 - frequently used in planning, theorem proving





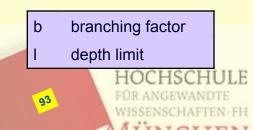
Depth-Limited Search

- similar to depth-first, but with a limit
 - overcomes problems with infinite paths
 - sometimes a depth limit can be inferred or estimated from the problem description
 - in other cases, a good depth limit is only known when the problem is solved
 - based on the TREE-SEARCH method
 - must keep track of the depth

function DEPTH-LIMITED-SEARCH(problem, depth-limit) returns solution

return TREE-SEARCH(problem, depth-limit, LIFO-QUEUE())

Time Complexity	p _l
Space Complexity	b*I
Completeness	no (goal beyond I, or infinite branch length)
Optimality	no



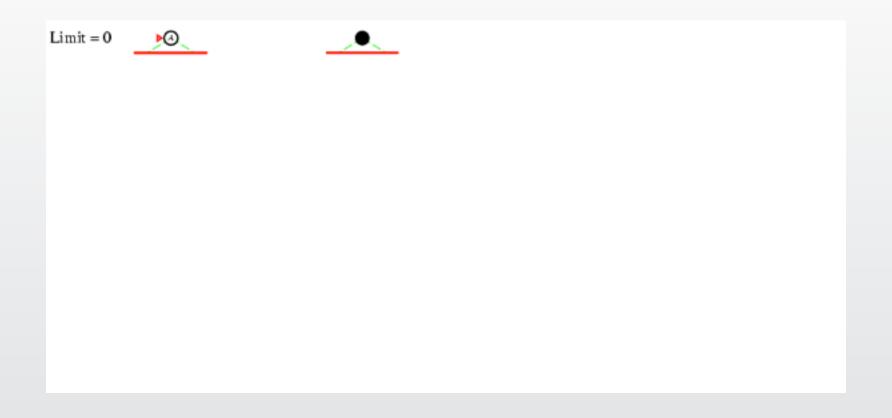


Iterative Deepening

- applies LIMITED-DEPTH with increasing depth limits
 - combines advantages of BREADTH-FIRST and DEPTH-FIRST methods
 - many states are expanded multiple times
 - doesn't really matter because the number of those nodes is small
 - in practice, one of the best uninformed search methods
 - for large search spaces, unknown depth

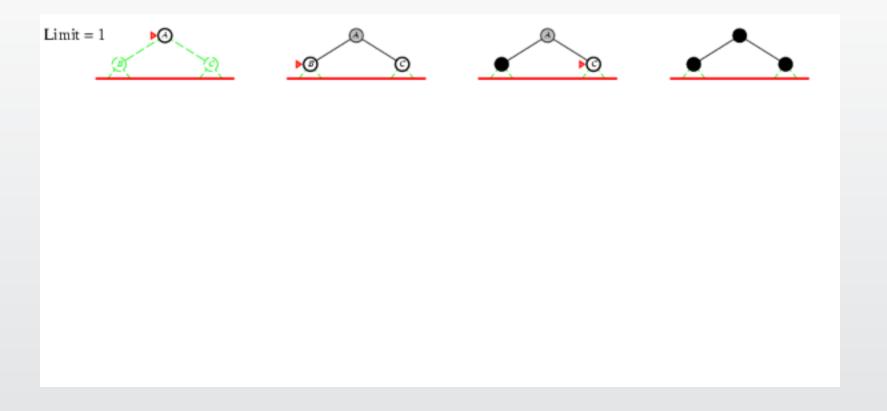
	Time Complexity	b ^d
	Space Complexity	b*d
	Completeness	yes (finite b)
7	Optimality	yes (all step costs identical)

		_
b	branching factor	
d	tree depth	CHULI
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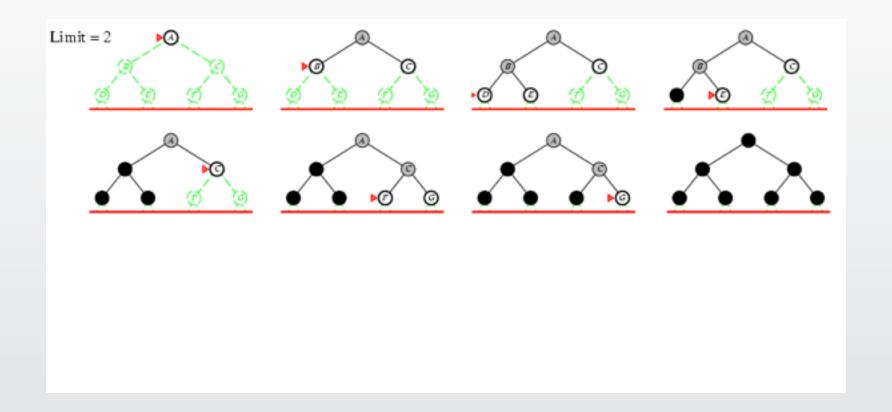






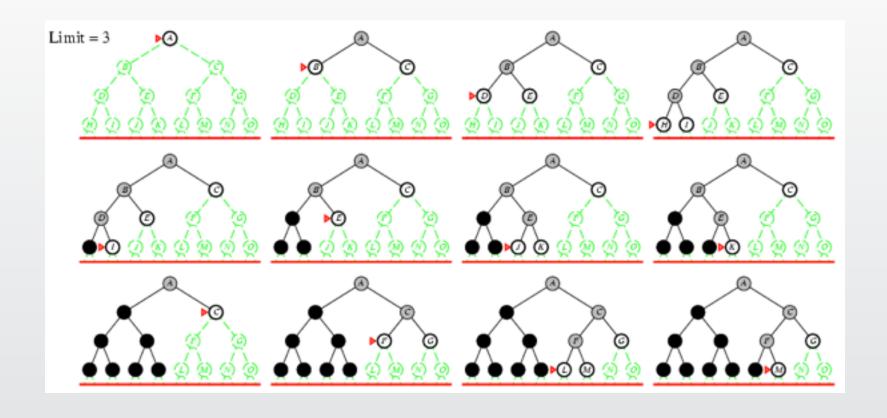
















Number of nodes generated in a depth-limited search to depth d with branching factor b:

$$N_{DIS} = b^0 + b^1 + b^2 + ... + b^{d-2} + b^{d-1} + b^d$$

Number of nodes generated in an iterative deepening search to depth d with branching factor b:

$$N_{IDS} = (d+1)b^{0} + db^{1} + (d-1)b^{2} + ... + 3b^{d-2} + 2b^{d-1} + 1b^{d}$$

• For b = 10, d = 5,

$$\bullet$$
 N_{DLS} = 1 + 10 + 100 + 1,000 + 10,000 + 100,000 = 111,111

$$N_{IDS} = 6 + 50 + 400 + 3,000 + 20,000 + 100,000 = 123,456$$

Overhead = (123,456 - 111,111)/111,111 = 11%

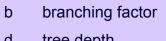




Bi-directional Search

- search simultaneously from two directions
 - forward from the initial and backward from the goal state
- may lead to substantial savings if it is applicable
- has severe limitations
 - predecessors must be generated, which is not always possible
 - search must be coordinated between the two searches
 - one search must keep all nodes in memory

	Time Complexity	b ^{d/2}
	Space Complexity	b ^{d/2}
	Completeness	yes (b finite, breadth-first for both directions)
1	Optimality	yes (all step costs identical, breadth-first for both directions)







Improving Search Methods

assumption for improvements

- remember information about the search so far
 - all nodes visited so far
 - path to the current node

make algorithms more efficient

- avoiding repeated states
- utilizing memory efficiently

use additional knowledge about the problem

- properties ("shape") of the search space
 - more interesting areas are investigated first
- pruning of irrelevant areas
 - areas that are guaranteed not to contain a solution can be discarded





Avoiding Repeated States

- in many approaches, states may be expanded multiple times
 - e.g. iterative deepening
 - problems with reversible actions
- eliminating repeated states may yield an exponential reduction in search cost
 - e.g. some n-queens strategies
 - place queen in the left-most non-threatening column
 - rectangular grid
 - ❖ 4^d leaves, but only 2^{d2} distinct states





Informed Search

Best-first

Greedy best-first

A*

A* modifications





Informed Search

- relies on additional knowledge about the problem or domain
 - frequently expressed through heuristics ("rules of thumb")
- used to distinguish more promising paths towards a goal
 - may be mislead, depending on the quality of the heuristic
- in general, performs much better than uninformed search
 - but frequently still exponential in time and space for realistic problems





Best-First Search

- relies on an evaluation function that gives an indication of how useful it would be to expand a node
 - family of search methods with various evaluation functions
 - usually gives an estimate of the distance to the goal
 - often referred to as heuristics in this context
- the node with the lowest value is expanded first
 - the name is a little misleading: the node with the lowest value for the evaluation function is not necessarily one that is on an optimal path to a goal





Greedy Best-First Search

- minimizes the estimated cost to a goal
 - expand the node that seems to be closest to a goal
 - utilizes a heuristic function as evaluation function
 - f(n) = h(n) = estimated cost from the current node to a goal
 - heuristic functions are problem-specific
 - often straight-line distance for route-finding and similar problems
 - often better than depth-first, although worst-time complexities are equal or worse (space)

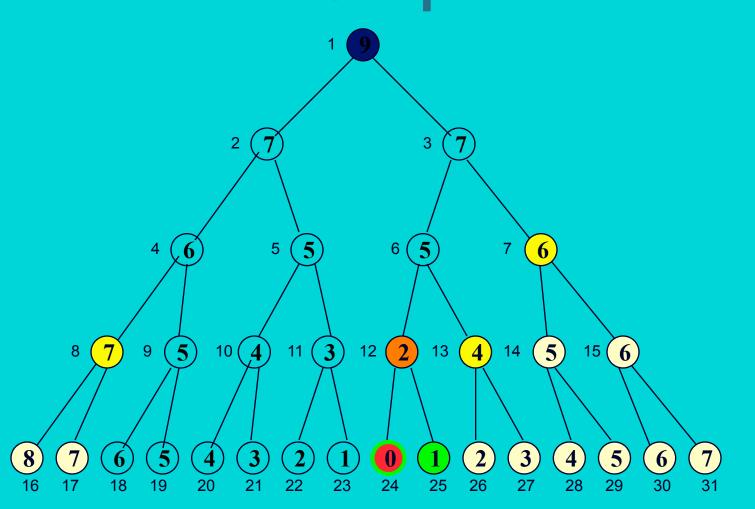
function GREEDY-SEARCH(problem) returns solution
 return BEST-FIRST-SEARCH(problem, h)

Completeness	Time Complexity	Space Complexity	Optimality
no	b ^m	b ^m	no





Greedy Best-First Search
Snapshot



Initial
Visited
Fringe
Current
Visible
Goal

Heuristics



A* Search

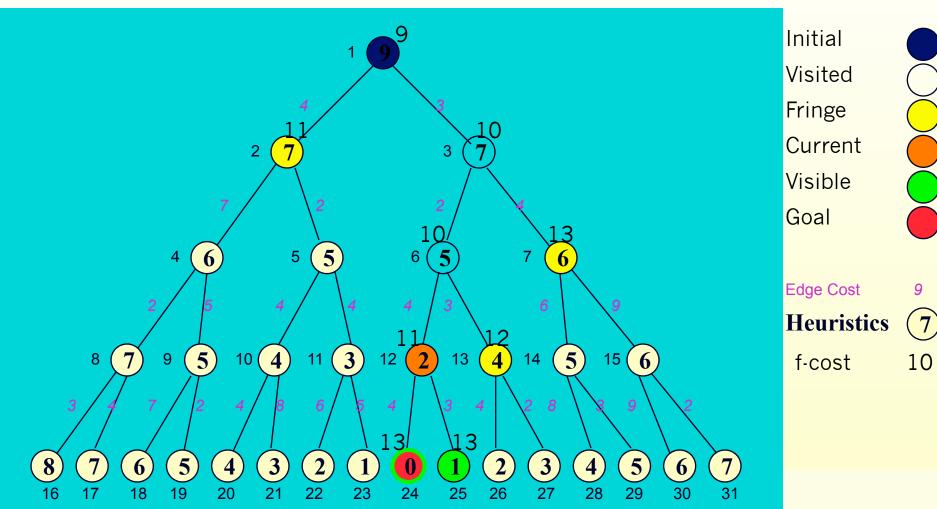
- combines greedy and uniform-cost search to find the (estimated) cheapest path through the current node
 - f(n) = g(n) + h(n)= path cost + estimated cost to the goal
 - heuristics must be admissible
 - never overestimate the cost to reach the goal
 - very good search method, but with complexity problems

```
function A*-SEARCH(problem) returns solution
return BEST-FIRST-SEARCH(problem, g+h)
```

Completeness	Time Complexity	Space Complexity	Optimality
yes	þ ^d	þ ^d	yes
1 1 0 1 1	.1 (.1 1 .1	1 1 61 1 1	1 1 1 1



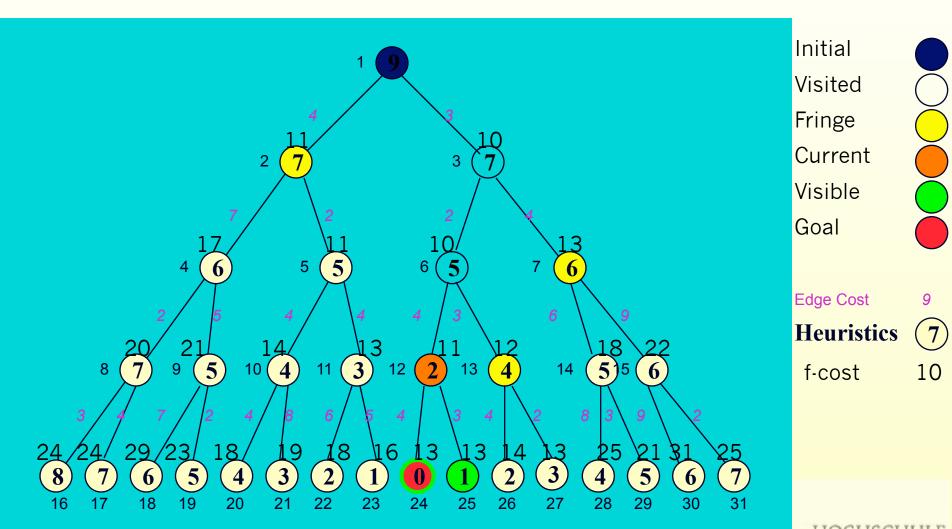
A* Snapshot



HOCHSCHULE
FÜR ANGEWANDTE
WISSENSCHAFTEN-FH

Fringe: [2(4+7), 13(3+2+3+4), 7(3+4+6)] + [24(3+2+4+4+0), 25(3+2+4+3+1)] 109

A* Snapshot with all f-Costs







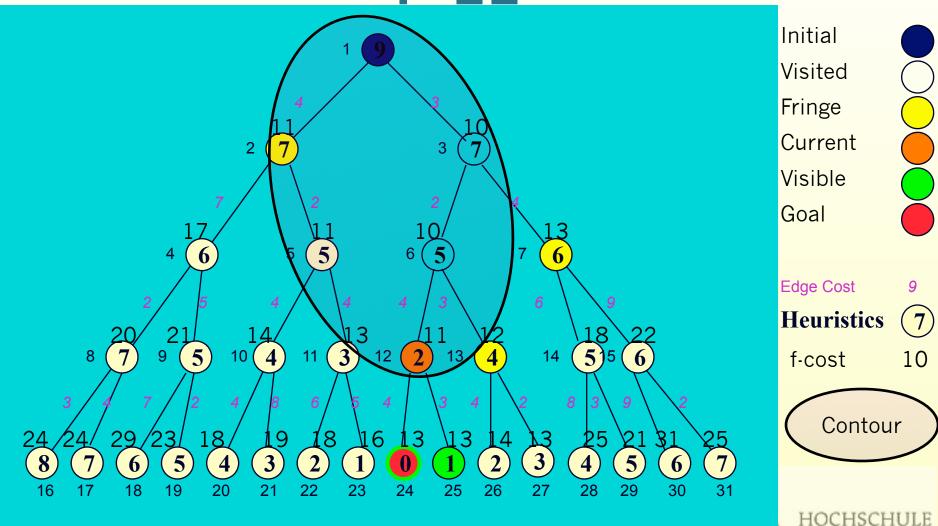
A* Properties

- the value of f never decreases along any path starting from the initial node
 - also known as monotonicity of the function
 - almost all admissible heuristics show monotonicity
 - those that don't can be modified through minor changes
- this property can be used to draw contours
 - regions where the f-cost is below a certain threshold
 - with uniform cost search (h = 0), the contours are circular
 - the better the heuristics h, the narrower the contour around the optimal path



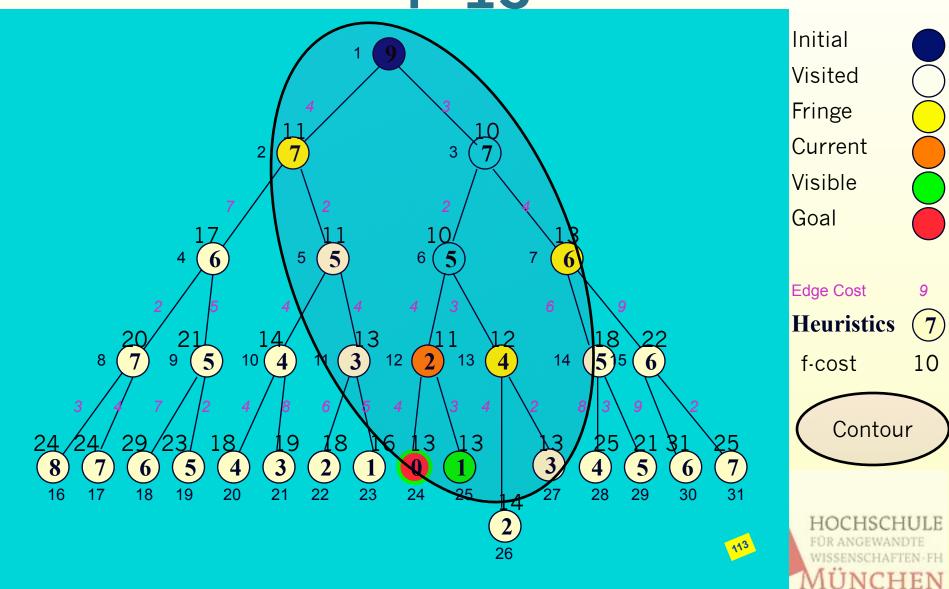


A* Snapshot with Contour f=11





A* Snapshot with Contour f=13



Optimality of A*

- A* will find the optimal solution
 - the first solution found is the optimal one
- A* is optimally efficient
 - no other algorithm is guaranteed to expand fewer nodes than A*
- A* is not always "the best" algorithm
 - optimality refers to the expansion of nodes
 - other criteria might be more relevant
 - it generates and keeps all nodes in memory
 - improved in variations of A*





Complexity of A*

- the number of nodes within the goal contour search space is still exponential
 - with respect to the length of the solution
 - better than other algorithms, but still problematic
- frequently, space complexity is more severe than time complexity
 - A* keeps all generated nodes in memory





Memory-Bounded Search

- search algorithms that try to conserve memory
- most are modifications of A*
 - iterative deepening A* (IDA*)
 - simplified memory-bounded A* (SMA*)





Iterative Deepening A* (IDA*)

- explores paths within a given contour (f-cost limit) in a depthfirst manner
 - this saves memory space because depth-first keeps only the current path in memory
 - but it results in repeated computation of earlier contours since it doesn't remember its history
 - was the "best" search algorithm for many practical problems for some time
 - does have problems with difficult domains
 - contours differ only slightly between states
 - algorithm frequently switches back and forth
 - similar to disk thrashing in (old) operating systems





Recursive Best-First Search

- similar to best-first search, but with lower space requirements
 - O(bd) instead of O(bm)
- it keeps track of the best alternative to the current path
 - best f-value of the paths explored so far from predecessors of the current node
 - if it needs to re-explore parts of the search space, it knows the best candidate path
 - still may lead to multiple re-explorations





Simplified Memory-Bounded A* (SMA*)

- uses all available memory for the search
 - drops nodes from the queue when it runs out of space
 - those with the highest f-costs
 - avoids re-computation of already explored area
 - keeps information about the best path of a "forgotten" subtree in its ancestor
 - complete if there is enough memory for the shortest solution path
 - often better than A* and IDA*
 - but some problems are still too tough
 - trade-off between time and space requirements





Heuristics for Searching

- for many tasks, a good heuristic is the key to finding a solution
 - prune the search space
 - move towards the goal
- relaxed problems
 - fewer restrictions on the successor function (operators)
 - its exact solution may be a good heuristic for the original problem





8-Puzzle Heuristics

level of difficulty

- around 20 steps for a typical solution
- branching factor is about 3
 - exhaustive search would be 320 = 3.5 * 109
- 9!/2 = 181,440 different reachable states
 - distinct arrangements of 9 squares

candidates for heuristic functions

- number of tiles in the wrong position
- sum of distances of the tiles from their goal position
 - city block or Manhattan distance

generation of heuristics

possible from formal specifications

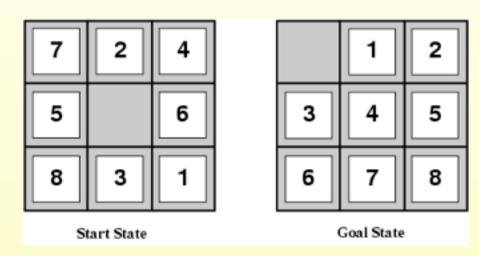




Admissible heuristics

E.g., for the 8-puzzle:

- h1(n) = number of misplaced tiles
- h2(n) = total Manhattan distance
- (i.e., no. of squares from desired location of each tile)



$$+ h2(S) = ?3+1+2+2+3+3+2 = 18$$





Important Concepts and Terms

- agent
- A* search
- best-first search
- bi-directional search
- breadth-first search
- depth-first search
- depth-limited search
- completeness
- constraint satisfaction
- depth-limited search
- genetic algorithm
- general search algorithm
- goal
- goal test function
- greedy best-first search
- heuristics
- initial state

- iterative deepening search
- iterative improvement
- local search
- memory-bounded search
- operator
- optimality
- path
- path cost function
- problem
- recursive best-first search
- search
- space complexity
- * state
- state space
- time complexity
- uniform-cost search





Chapter Summary

- tasks can often be formulated as search problems
 - initial state, successor function (operators), goal test, path cost
- various search methods systematically comb the search space
 - uninformed search
 - breadth-first, depth-first, and variations
 - informed search
 - best-first, A*, iterative improvement
- the choice of good heuristics can improve the search dramatically
 - task-dependent







