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1.

a)

Variables: $\{K1, K2, Q1, Q2\}$

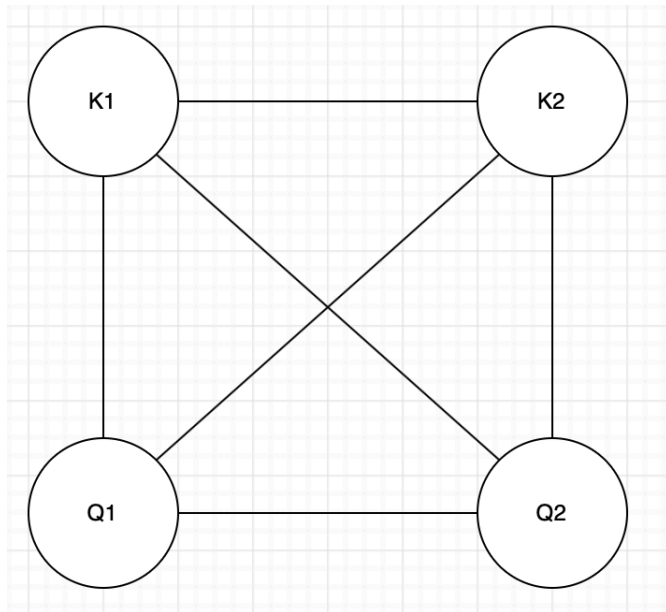
Domain (same for all variables): $\{1, 2, 3, 4\}$

Assume the horizontal axis is i and vertical axis is j .

Constraints: $Q1j = 3, Q2j = 4, K1j = 1, K2j = 2,$

$K1i \neq K2i \pm 2, K1i \neq Q1i \pm 1, K2i \neq Q2i \pm 1, K2i \neq Q1i \pm 2,$

$|Q1i - Q2i| \neq |Q1j - Q2j|, Q1i \neq Q2i, |Q_i - K_i| \neq |Q_j - K_j|, Q_i \neq K_i.$



b)

$K1 = \{1, 2, 3, 4\}$

$K2 = \{1, 2, 3, 4\}$

$Q1 = \{1, 2, 3, 4\}$

$Q2 = \{1, 2, 3, 4\}$

Queue: ~~$K1i \neq K2i \pm 2, K2i \neq K1i, K1i \neq Q1i \pm 1, Q1i \neq K1i, K2i \neq Q2i \pm 1, Q2i \neq K2i, K2i \neq Q1i \pm 2, Q1i \neq K2i, |Q1i - Q2i| \neq |Q1j - Q2j|, |Q1j - Q2j| \neq |Q1i - Q2i|, Q1i \neq Q2i, Q2i \neq Q1i, |Q_i - K_i| \neq |Q_j - K_j|, |Q_j - K_j| \neq |Q_i - K_i|, Q_i \neq K_i, K_i \neq Q_i,$~~ $Q1i \pm 1 \neq K1i,$

$$Q2i \pm 1 \neq K2i, Q1i \pm 2 \neq K2i, Q2i \neq Q1i.$$

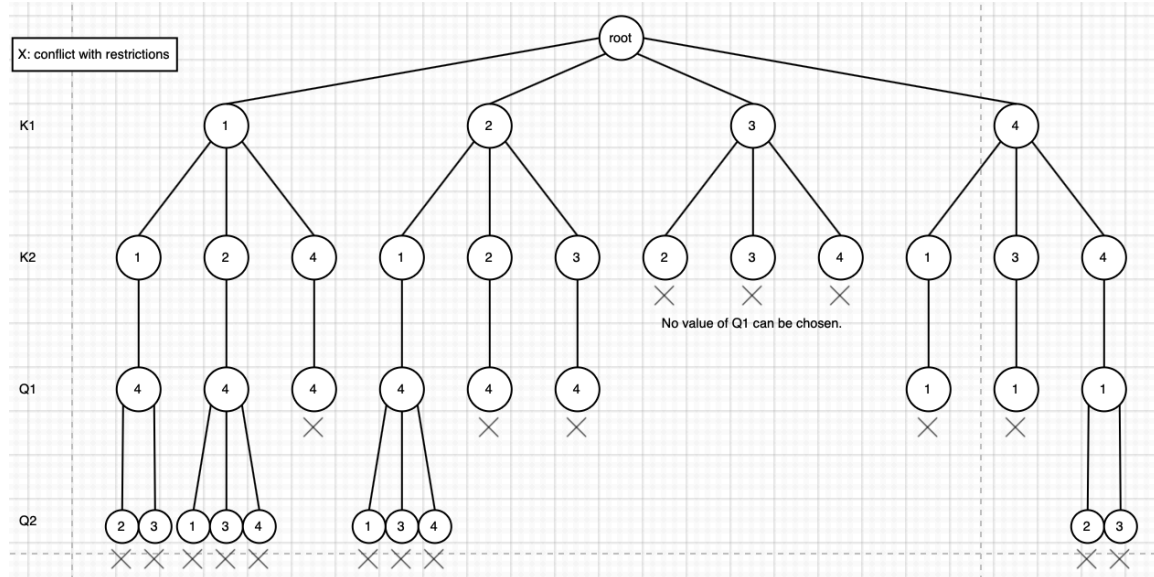
For the backtracking search, since there is no value left for each variable, a tree cannot be built using this method, as a result, no solution can be found.

c)

One step forward checking:

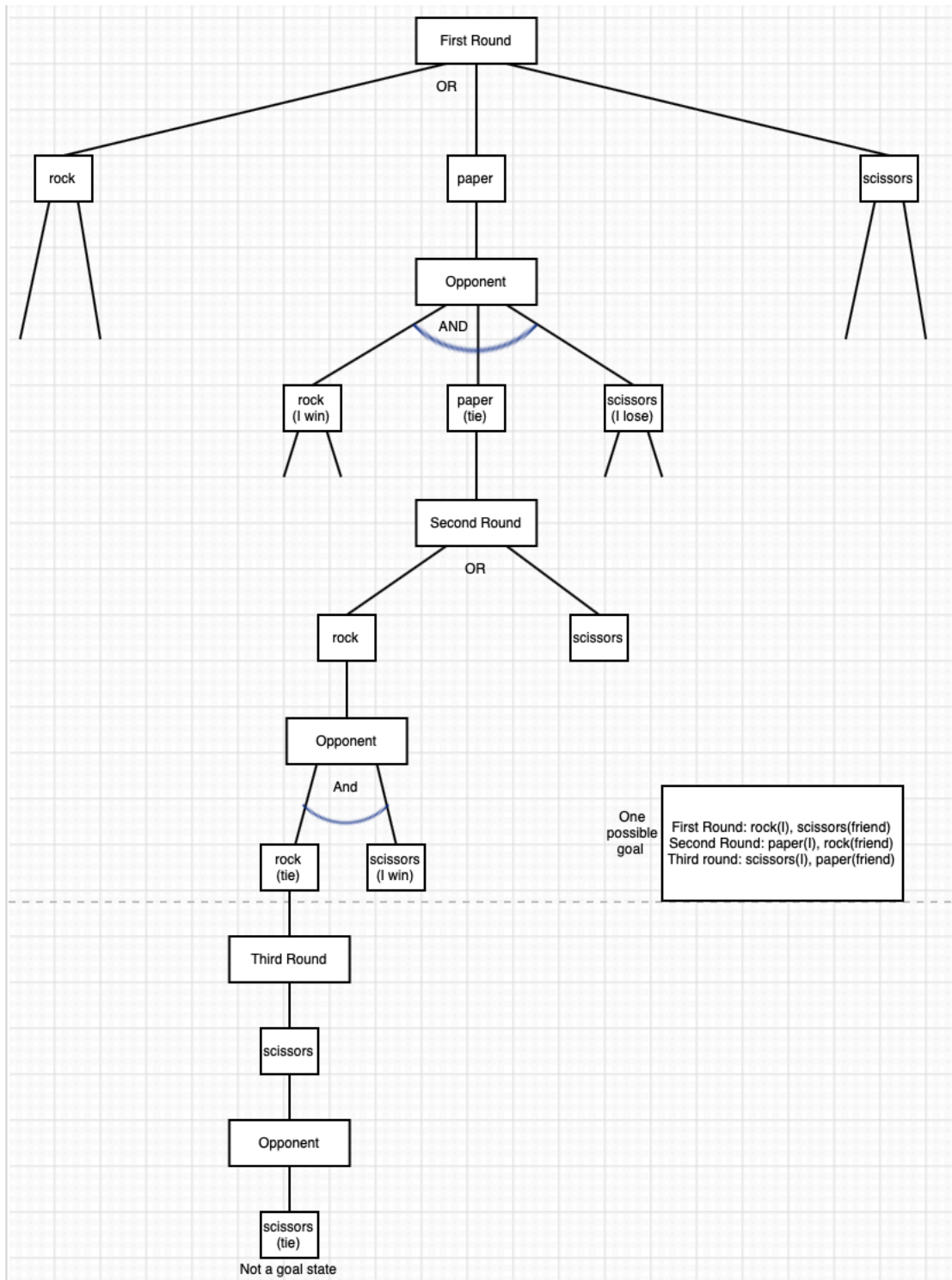
	Nothing Assigned	Assign K1 = 1	Assign K1 = 2	Assign K1 = 3	Assign K1 = 4
K1	1234	1	2	3	4
K2	1234	124	123	234	134
Q1	1234	4	4		1
Q2	1234	23	134	124	23

Backtracking search:



2.

a)



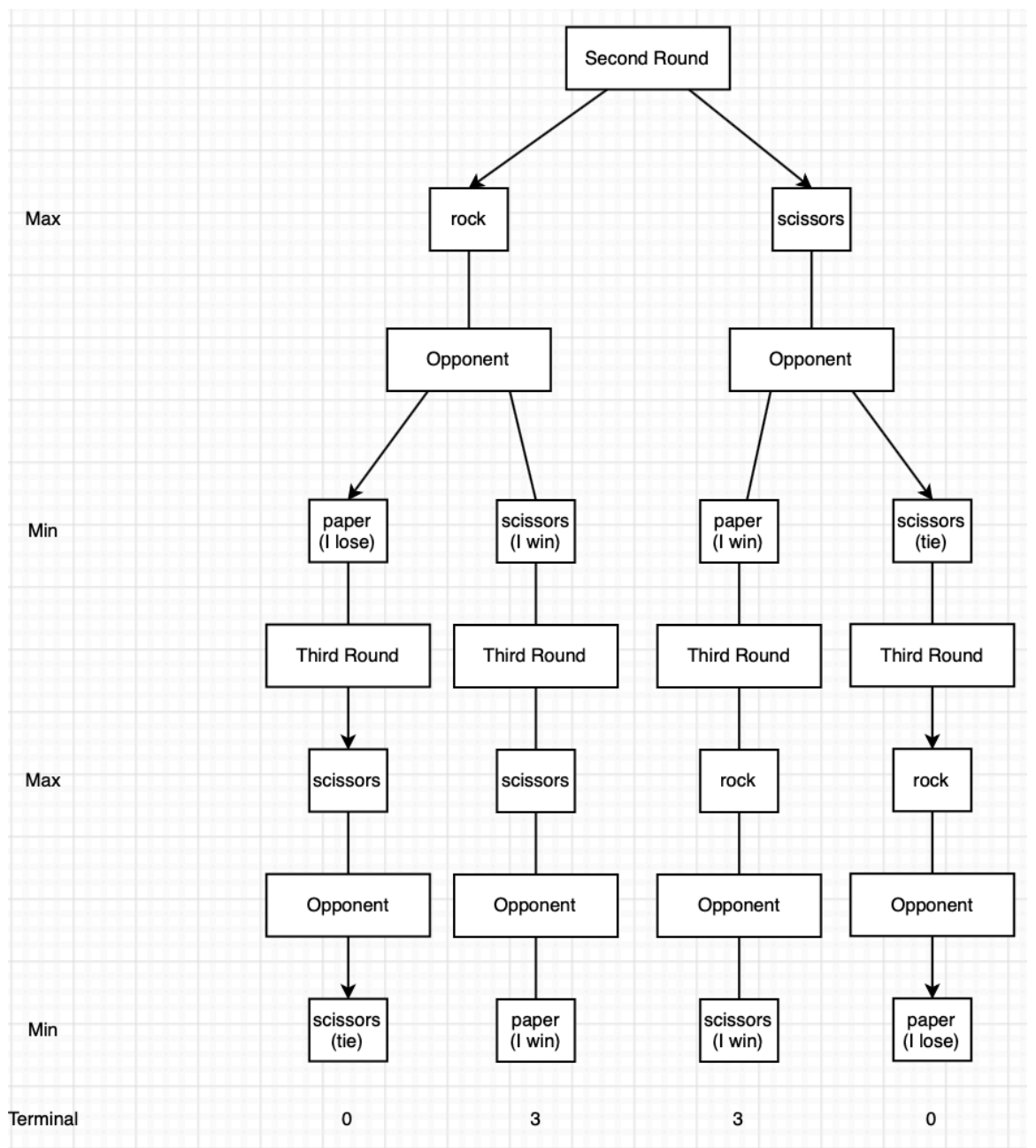
b)

In this situation, I can play **rock** or **scissors** and my friend can play **paper** or **scissors**. If I

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graph TD
    SR[Second Round] --> R[rock]
    SR --> S[scissors]
    R --> O1[Opponent]
    S --> O2[Opponent]
    O1 --> P1[paper I lose]
    O1 --> S1[scissors I win]
    O2 --> P2[paper I win]
    O2 --> S2[scissors tie]
    P1 --> TR1[Third Round]
    S1 --> TR2[Third Round]
    P2 --> TR3[Third Round]
    S2 --> TR4[Third Round]
    TR1 --> S3[scissors]
    TR2 --> S4[scissors]
    TR3 --> R2[rock]
    TR4 --> R3[rock]
    S3 --> O3[Opponent]
    S4 --> O4[Opponent]
    R2 --> O5[Opponent]
    R3 --> O6[Opponent]
    O3 --> S5[scissors tie]
    O4 --> P3[paper I win]
    O5 --> S6[scissors I win]
    O6 --> P4[paper I lose]
    S5 --- T1[0]
    P3 --- T2[3]
    S6 --- T3[3]
    P4 --- T4[0]

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The above graph is a game tree starting from second round. The connection with arrow is the direction of flow after applying minimax algorithm. This terminal result supports the idea that I cannot guarantee a win overall.

3)

a)

i)

A	B	C	$A \wedge B$	$(A \wedge B) \vee C$
T	T	T	T	T
T	T	F	T	T
T	F	T	F	T
T	F	F	F	F
F	T	T	F	T
F	T	F	F	F
F	F	T	F	T
F	F	F	F	F

5 models

ii)

A	B	$A \Rightarrow B$	$A \wedge (A \Rightarrow B)$	$A \wedge (A \Rightarrow B) \wedge \neg B$
T	T	T	T	F
T	F	F	F	F
F	T	T	F	F
F	F	T	F	F

0 models

iii)

We assume the statement is $\neg f1 \vee f2$. In order to ensure this logic to be true, $f1$ should be false and $f2$ should be true. So for $f1$, exclude all values are true, all other combinations will lead to a false. As a result, we only need to consider that exempt condition. When $f1$

is true, $B \wedge C$ is true. Thus the final result is still true for this situation. In conclusion, there are $2^5 = 32$ models.

b)

i)

A	B	C	f1: $A \Rightarrow B$	f2: $A \vee C$	f3: $f1 \wedge f2$	f4: $B \vee C$	f3 \Rightarrow f4
T	T	T	T	T	T	T	T
T	T	F	T	T	T	T	T
T	F	T	F	T	F	T	T
T	F	F	F	T	F	F	T
F	T	T	T	T	T	T	T
F	T	F	T	F	F	T	T
F	F	T	T	T	T	T	T
F	F	F	T	F	F	F	T

This is valid.

ii)

A	B	C	f1: $A \vee B$	f2: $f1 \Rightarrow C$	f3: $f2 \Rightarrow C$	f4: $B \Rightarrow C$	f5: $f3 \wedge f4$	f2 \Rightarrow f5
T	T	T	T	T	T	T	T	T
T	T	F	T	F	F	F	F	T
T	F	T	T	T	T	T	T	T
T	F	F	T	F	F	T	F	T
F	T	T	T	T	T	T	T	T
F	T	F	T	F	T	F	F	T
F	F	T	F	T	T	T	T	T
F	F	F	F	T	T	T	T	T

This is also valid.