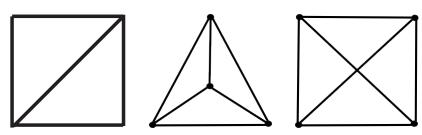
#### COMP229: Introduction to Data Science Lecture 12: Kuratowski's criterion of planarity

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## Planar graphs

**Definition 12.1**. An abstract graph that can be embedded (drawn without intersections) in the plane  $\mathbb{R}^2$  (or, equivalently  $S^2$ ) is called *planar*.



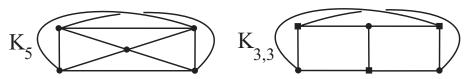
Each of the graphs in the picture has 4 vertices.

Which of them are planar by the definition above?



# Criterion for planar graphs

Claim 12.2. [Kuratowski's criterion of planarity] A graph G is planar if and only if G has no (subdivisions of) subgraphs  $K_5$  and  $K_{3,3}$ , which are called Kuratowski's *forbidden* subgraphs.



We'll prove the easier "only if" part: if a graph G contains  $K_5$  or  $K_{3,3}$ , then G isn't planar. It's enough to prove that  $K_5$  or  $K_{3,3}$  aren't planar.



# A proof by contradiction

When we need to prove a claim for all objects, e.g. any continuous map  $K_5 \to \mathbb{R}^2$  has intersections in the image, a typical way is by contradiction.

Assume that the graph  $K_5 \subset \mathbb{R}^2$  is embedded.

This embedded graph  $K_5$  has V=5 vertices and E=10 edges. Use Euler's formula to find the number F of faces (connected regions of  $\mathbb{R}^2-K_5$ ).

# Counting edges twice

Euler's formula V - E + F = 2 for the connected graph  $K_5 \subset \mathbb{R}^2$  with V = 5, E = 10 gives F = 7.

**Claim 12.3**. Each of 7 (hypothetical) faces of a graph  $K_5 \subset \mathbb{R}^2$  has at least 3 boundary edges.

*Proof.*  $K_5$  has no loops (edges connecting a vertex to itself) and no double edges (two different edges between the same pair of vertices).

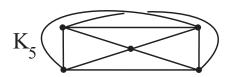
Do you already see a contradiction?

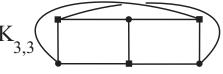


## Why is $K_5$ not planar?

If each of 7 faces of  $K_5 \subset \mathbb{R}^2$  has at least 3 boundary edges, the total number of edges is at least  $7 \times 3 = 21$ . We counted each edge twice, because every edge belongs to the boundaries of 2 faces. The double number of edges is  $2E \geq 21$ .

However,  $K_5$  has 10 edges, which contradicts  $2E \ge 21$  and proves that  $K_5$  isn't planar.

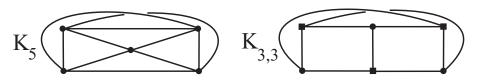




## Why is $K_{3,3}$ not planar?

Try the similar approach by contradiction.

Assume that there is an embedding  $K_{3,3} \subset \mathbb{R}^2$ .



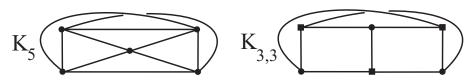
The graph has V=6 vertices, E=9 edges and F=5 faces by Euler's formula V-E+F=2.

If each of F=5 faces has at least 3 edges, the double number of edges is  $2E \ge 5 \times 3 = 15$ .



# $K_{3,3}$ is indeed not planar

 $K_{3,3}$  is bipartite, hence two types of vertices (square/round dots) alternate and any cycle has an even number (minimum 4, not 3) of edges.



The double number of edges counted over 5 faces is  $2E > 5 \times 4 = 20$ , which contradicts E = 9.

What graphs can we embed in 3 dimensions?



# 3 dimensions are enough

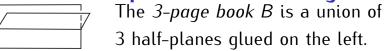
**Claim 12.4**. Any graph can be embedded in  $\mathbb{R}^3$ .

*Proof.* When we draw an intersection with one edge going over another edge, we implicitly show (a plane projection) of an embedding  $G \subset \mathbb{R}^3$ 

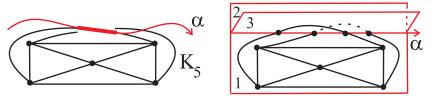
Draw any graph with only double intersections between edges. Such a drawing (in a general position) can be obtained by a slight perturbation. At every crossing, push one arc into the 3rd dimension, keep another arc in the plane.



## 3 half-planes are enough



**Claim 12.5**. Any graph G can be embedded in B.



*Proof.* Draw a curve  $\alpha$  without self-intersections that goes once through all overcrossing arcs. Deform  $\mathbb{R}^2$  to make  $\alpha$  straight. Push overcrossing arcs to the 3rd half-plane glued along  $\alpha$ .

### Your questions and the quiz

To benefit from the lecture, now you could

- ask or submit your anonymous questions to the COMP229 folder after the lecture;
- write down your summary in 2-3 phrases,
  e.g. list key concepts you have learned;
- talk to your classmates to revise the lecture.

**Question**. Are the graphs  $K_5$  and  $K_{3,3}$  not planar, because their drawings below have intersections?



# Answer to the quiz and summary

**Answer**. No, these drawings prove nothing.  $K_5$  and  $K_{3,3}$  are not planar, because they have no possible drawings without self-intersections (by Kuratowski's criterion of planarity).

- A *planar* graph: can be embedded (drawn without intersections) in the plane  $\mathbb{R}^2$ .
- A graph G is planar if and only if G contains no (subdivisions of) the graphs  $K_5$  and  $K_{3,3}$ . The "only if" part follows from Euler's formula.
- Any graph can be drawn in  $\mathbb{R}^3$ , in 3 pages.