

# SMART CONTRACT AUDIT REPORT

for

QuantoSwap

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# 1 Introduction

Given the opportunity to review the design document and related source code of the QuantoSwap protocol, we outline in the report our systematic approach to evaluate potential security issues in the smart contract implementation, expose possible semantic inconsistencies between smart contract code and design document, and provide additional suggestions or recommendations for improvement. Our results show that the given version of smart contracts can be further improved due to the presence of several issues related to either security or performance. This document outlines our audit results.

#### 1.1 About QuantoSwap

QuantoSwap is a DEX that is forked from the popular UniswapV2 with the further integration with MasterChef to engage protocol users. Harnessing the power of multiple liquidity pools, QuantoSwap LPs are fungible and are composable for a wide variety of utility, including on-chain collateralization, farming, as well as other yield strategies. The basic information of audited contracts is as follows:

Item	Description
Name	QuantoSwap
Website	https://quantoswap.org/
Туре	Smart Contract
Language	Solidity
Audit Method	Whitebox
Latest Audit Report	May 7, 2024

Table 1.1: Basic Information of QuantoSwap

In the following, we show the Git repositories of reviewed files and the commit hash values used in this audit.

- https://github.com/QuantoSwap/farming.git (db6d742)
- https://github.com/QuantoSwap/time-lock.git (9b1ff54)

- https://github.com/QuantoSwap/v2-periphery.git (6af45bb)
- https://github.com/QuantoSwap/v2-core.git (5db3c39)

And here are the commit IDs after all fixes for the issues found in the audit have been checked in:

- https://github.com/QuantoSwap/farming.git (7397641)
- https://github.com/QuantoSwap/time-lock.git (9b1ff54)
- https://github.com/QuantoSwap/v2-periphery.git (6ba929d)
- https://github.com/QuantoSwap/v2-core.git (5db3c39)

#### 1.2 About PeckShield

PeckShield Inc. [11] is a leading blockchain security company with the goal of elevating the security, privacy, and usability of current blockchain ecosystems by offering top-notch, industry-leading services and products (including the service of smart contract auditing). We are reachable at Telegram (https://t.me/peckshield), Twitter (http://twitter.com/peckshield), or Email (contact@peckshield.com).

High Critical High Medium

High Medium

Low

Medium Low

High Medium

Low

High Medium

Low

Likelihood

Table 1.2: Vulnerability Severity Classification

## 1.3 Methodology

To standardize the evaluation, we define the following terminology based on OWASP Risk Rating Methodology [10]:

• <u>Likelihood</u> represents how likely a particular vulnerability is to be uncovered and exploited in the wild;

- Impact measures the technical loss and business damage of a successful attack;
- Severity demonstrates the overall criticality of the risk.

Likelihood and impact are categorized into three ratings: *H*, *M* and *L*, i.e., *high*, *medium* and *low* respectively. Severity is determined by likelihood and impact, and can be accordingly classified into four categories, i.e., *Critical*, *High*, *Medium*, *Low* shown in Table 1.2.

To evaluate the risk, we go through a list of check items and each would be labeled with a severity category. For one check item, if our tool or analysis does not identify any issue, the contract is considered safe regarding the check item. For any discovered issue, we might further deploy contracts on our private testnet and run tests to confirm the findings. If necessary, we would additionally build a PoC to demonstrate the possibility of exploitation. The concrete list of check items is shown in Table 1.3.

In particular, we perform the audit according to the following procedure:

- <u>Basic Coding Bugs</u>: We first statically analyze given smart contracts with our proprietary static code analyzer for known coding bugs, and then manually verify (reject or confirm) all the issues found by our tool.
- <u>Semantic Consistency Checks</u>: We then manually check the logic of implemented smart contracts and compare with the description in the white paper.
- Advanced DeFi Scrutiny: We further review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.
- Additional Recommendations: We also provide additional suggestions regarding the coding and development of smart contracts from the perspective of proven programming practices.

To better describe each issue we identified, we categorize the findings with Common Weakness Enumeration (CWE-699) [9], which is a community-developed list of software weakness types to better delineate and organize weaknesses around concepts frequently encountered in software development. Though some categories used in CWE-699 may not be relevant in smart contracts, we use the CWE categories in Table 1.4 to classify our findings. Moreover, in case there is an issue that may affect an active protocol that has been deployed, the public version of this report may omit such issue, but will be amended with full details right after the affected protocol is upgraded with respective fixes.

#### 1.4 Disclaimer

Table 1.3: The Full List of Check Items

Category	Check Item
	Constructor Mismatch
	Ownership Takeover
	Redundant Fallback Function
	Overflows & Underflows
	Reentrancy
	Money-Giving Bug
	Blackhole
	Unauthorized Self-Destruct
Basic Coding Bugs	Revert DoS
Dasic Couling Dugs	Unchecked External Call
	Gasless Send
	Send Instead Of Transfer
	Costly Loop
	(Unsafe) Use Of Untrusted Libraries
	(Unsafe) Use Of Predictable Variables
	Transaction Ordering Dependence
	Deprecated Uses
Semantic Consistency Checks	Semantic Consistency Checks
	Business Logics Review
	Functionality Checks
	Authentication Management
	Access Control & Authorization
	Oracle Security
Advanced DeFi Scrutiny	Digital Asset Escrow
ravancea Ber i Geraemi,	Kill-Switch Mechanism
	Operation Trails & Event Generation
	ERC20 Idiosyncrasies Handling
	Frontend-Contract Integration
Additional Recommendations	Deployment Consistency
	Holistic Risk Management
	Avoiding Use of Variadic Byte Array
	Using Fixed Compiler Version
	Making Visibility Level Explicit
	Making Type Inference Explicit
	Adhering To Function Declaration Strictly
	Following Other Best Practices

Table 1.4: Common Weakness Enumeration (CWE) Classifications Used in This Audit

Category	Summary
Configuration	Weaknesses in this category are typically introduced during
	the configuration of the software.
Data Processing Issues	Weaknesses in this category are typically found in functional-
	ity that processes data.
Numeric Errors	Weaknesses in this category are related to improper calcula-
	tion or conversion of numbers.
Security Features	Weaknesses in this category are concerned with topics like
	authentication, access control, confidentiality, cryptography,
	and privilege management. (Software security is not security
	software.)
Time and State	Weaknesses in this category are related to the improper man-
	agement of time and state in an environment that supports
	simultaneous or near-simultaneous computation by multiple
	systems, processes, or threads.
Error Conditions,	Weaknesses in this category include weaknesses that occur if
Return Values,	a function does not generate the correct return/status code,
Status Codes	or if the application does not handle all possible return/status
	codes that could be generated by a function.
Resource Management	Weaknesses in this category are related to improper manage-
	ment of system resources.
Behavioral Issues	Weaknesses in this category are related to unexpected behav-
	iors from code that an application uses.
Business Logics	Weaknesses in this category identify some of the underlying
	problems that commonly allow attackers to manipulate the
	business logic of an application. Errors in business logic can
	be devastating to an entire application.
Initialization and Cleanup	Weaknesses in this category occur in behaviors that are used
	for initialization and breakdown.
Arguments and Parameters	Weaknesses in this category are related to improper use of
	arguments or parameters within function calls.
Expression Issues	Weaknesses in this category are related to incorrectly written
	expressions within code.
Coding Practices	Weaknesses in this category are related to coding practices
	that are deemed unsafe and increase the chances that an ex-
	ploitable vulnerability will be present in the application. They
	may not directly introduce a vulnerability, but indicate the
	product has not been carefully developed or maintained.

Note that this security audit is not designed to replace functional tests required before any software release, and does not give any warranties on finding all possible security issues of the given smart contract(s) or blockchain software, i.e., the evaluation result does not guarantee the nonexistence of any further findings of security issues. As one audit-based assessment cannot be considered comprehensive, we always recommend proceeding with several independent audits and a public bug bounty program to ensure the security of smart contract(s). Last but not least, this security audit should not be used as investment advice.



# 2 | Findings

#### 2.1 Summary

Here is a summary of our findings after analyzing the design and implementation of the QuantoSwap protocol. During the first phase of our audit, we study the smart contract source code and run our in-house static code analyzer through the codebase. The purpose here is to statically identify known coding bugs, and then manually verify (reject or confirm) issues reported by our tool. We further manually review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.

Severity	# of Findings
Critical	0
High	1
Medium	4
Low	2
Undetermined	1
Total	8

We have so far identified a list of potential issues: some of them involve subtle corner cases that might not be previously thought of, while others refer to unusual interactions among multiple contracts. For each uncovered issue, we have therefore developed test cases for reasoning, reproduction, and/or verification. After further analysis and internal discussion, we determined a few issues of varying severities need to be brought up and paid more attention to, which are categorized in the above table. More information can be found in the next subsection, and the detailed discussions of each of them are in Section 3.

Resolved

Resolved

Resolved

Mitigated

Coding Practices

Business Logic

**Business Logic** 

Security Features

#### 2.2 Key Findings

Overall, these smart contracts are well-designed and engineered, though the implementation can be improved by resolving the identified issues (shown in Table 2.1), including 1 high-severity vulnerabilities, 4 medium-severity vulnerability, 2 low-severity vulnerabilities, and 1 undetermined issue.

ID Title Category Severity **Status** PVE-001 Duplicate Pool Detection and Preven-Resolved Medium **Business Logic** tion in MasterChef **PVE-002** Timely massUpdatePools During Pool Resolved Low **Business Logic** Weight Changes **PVE-003** Medium Suggested Adherence of Checks-Effects-**Business Logic** Resolved Interactions (SmartChef) **PVE-004** High Sybil Attacks on QuantoSwap Token **Business Logic** Resolved Voting

Implicit Assumption Enforcement In Ad-

Inconsistent Fee Share Calculation in

Staking Incompatibility With Deflation-

QuantoSwapV2LiquidityMathLibrary

ary Tokens in SmartChef

Trust Issue Of Admin Keys

dLiquidity()

Table 2.1: Key Audit Findings

Beside the identified issues, we emphasize that for any user-facing applications and services, it is always important to develop necessary risk-control mechanisms and make contingency plans, which may need to be exercised before the mainnet deployment. The risk-control mechanisms should kick in at the very moment when the contracts are being deployed on mainnet. Please refer to Section 3 for details.

PVE-005

PVE-006

**PVE-007** 

**PVE-008** 

Low

Medium

Undetermined

Medium

# 3 Detailed Results

#### 3.1 Duplicate Pool Detection and Prevention in MasterChef

• ID: PVE-001

Severity: MediumLikelihood: Low

• Impact: Medium

• Target: MasterChef

Category: Business Logic [7]CWE subcategory: CWE-841 [4]

#### Description

The QuantoSwap protocol provides incentive mechanisms that reward the staking of supported assets with certain reward tokens. The rewards are carried out by designating a number of staking pools into which supported assets can be staked. Each pool has its allocPoint\*100%/totalAllocPoint share of scheduled rewards and the rewards for stakers are proportional to their share of LP tokens in the pool.

In current implementation, there are a number of concurrent pools that share the rewarded tokens and more can be scheduled for addition (via a proper governance procedure). To accommodate these new pools, the design has the necessary mechanism in place that allows for dynamic additions of new staking pools that can participate in being incentivized as well.

The addition of a new pool is implemented in add(), whose code logic is shown below. It turns out it did not perform necessary sanity checks in preventing a new pool but with a duplicate token from being added. Though it is a privileged interface (protected with the modifier onlyOwner), it is still desirable to enforce it at the smart contract code level, eliminating the concern of wrong pool introduction from human omissions.

```
108
             totalAllocPoint = totalAllocPoint.add(_allocPoint);
109
             poolInfo.push(
110
                 PoolInfo({
111
                      lpToken: _lpToken,
112
                      allocPoint: _allocPoint,
113
                      lastRewardBlock: lastRewardBlock,
114
                      accQNSPerShare: 0
                 })
115
116
             );
117
```

Listing 3.1: MasterChef::add()

**Recommendation** Detect whether the given pool for addition is a duplicate of an existing pool. The pool addition is only successful when there is no duplicate.

```
103
         function checkPoolDuplicate(IERC20 _lpToken) public {
104
             uint256 length = poolInfo.length;
105
             for (uint256 pid = 0; pid < length; ++pid) {</pre>
106
                 require(poolInfo[_pid].lpToken != _lpToken, "add: existing pool?");
107
             }
108
         }
109
110
         function add( uint256 _allocPoint, IERC20 _lpToken, bool _withUpdate ) public
             onlyOwner {
111
             if (_withUpdate) {
112
                 massUpdatePools();
113
114
             checkPoolDuplicate(_lpToken);
115
             uint256 lastRewardBlock = block.number > startBlock ? block.number : startBlock;
             totalAllocPoint = totalAllocPoint.add(_allocPoint);
116
117
             poolInfo.push(
118
                 PoolInfo({
119
                     lpToken: _lpToken,
120
                     allocPoint: _allocPoint,
121
                     lastRewardBlock: lastRewardBlock,
122
                     accQNSPerShare: 0
123
                 })
124
             );
125
```

Listing 3.2: Revised MasterChef::add()

We point out that if a new pool with a duplicate LP token can be added, it will likely cause a havoc in the distribution of rewards to the pools and the stakers.

**Status** This issue has been fixed in the following commit: 7397641.

## 3.2 Timely massUpdatePools During Pool Weight Changes

• ID: PVE-002

• Severity: Low

• Likelihood: Low

• Impact: Medium

• Target: MasterChef

• Category: Business Logic [7]

• CWE subcategory: CWE-841 [4]

#### Description

As mentioned earlier, the QuantoSwap protocol provides incentive mechanisms that reward the staking of supported assets. The rewards are carried out by designating a number of staking pools into which supported assets can be staked. And staking users are rewarded in proportional to their share of LP tokens in the reward pool.

The reward pools can be dynamically added via add() and the weights of supported pools can be adjusted via set(). When analyzing the pool weight update routine set(), we notice the need of timely invoking massUpdatePools() to update the reward distribution before the new pool weight becomes effective.

Listing 3.3: MasterChef::set()

If the call to massUpdatePools() is not immediately invoked before updating the pool weights, certain situations may be crafted to create an unfair reward distribution. Moreover, a hidden pool without any weight can suddenly surface to claim unreasonable share of rewarded tokens. Fortunately, this interface is restricted to the owner (via the onlyOwner modifier), which greatly alleviates the concern.

Recommendation Timely invoke massUpdatePools() when any pool's weight has been updated. In fact, the third parameter (\_withUpdate) to the set() routine can be simply ignored or removed. Note the same issue is also applicable to other routines, including updateMultiplier() and newQNSPerBlock().

```
function set(uint256 _pid, uint256 _allocPoint, bool _withUpdate) public onlyOwner {
    massUpdatePools();
    totalAllocPoint = totalAllocPoint.sub(poolInfo[_pid].allocPoint).add(_allocPoint);
```

Listing 3.4: Revised MasterChef::set()

**Status** This issue has been fixed in the following commit: 7397641.



# 3.3 Suggested Adherence Of Checks-Effects-Interactions Pattern

• ID: PVE-003

• Severity: Low

• Likelihood: Low

Impact: Low

Target: MasterChef

• Category: Time and State [8]

• CWE subcategory: CWE-663 [3]

#### Description

A common coding best practice in Solidity is the adherence of checks-effects-interactions principle. This principle is effective in mitigating a serious attack vector known as re-entrancy. Via this particular attack vector, a malicious contract can be reentering a vulnerable contract in a nested manner. Specifically, it first calls a function in the vulnerable contract, but before the first instance of the function call is finished, second call can be arranged to re-enter the vulnerable contract by invoking functions that should only be executed once. This attack was part of several most prominent hacks in Ethereum history, including the DAO [13] exploit, and the recent Uniswap/Lendf.Me hack [12].

We notice there is an occasion where the <code>checks-effects-interactions</code> principle is violated. Using the <code>MasterChef</code> as an example, the <code>emergencyWithdraw()</code> function (see the code snippet below) is provided to externally call a token contract to transfer assets. However, the invocation of an external contract requires extra care in avoiding the above <code>re-entrancy</code>.

Apparently, the interaction with the external contract (line 261) starts before effecting the update on internal states (lines 266–267), hence violating the principle. In this particular case, if the external contract has certain hidden logic that may be capable of launching re-entrancy via the same entry function.

```
258
        function emergencyWithdraw(uint256 _pid) public {
259
             PoolInfo storage pool = poolInfo[_pid];
260
             UserInfo storage user = userInfo[_pid][msg.sender];
261
             pool.lpToken.safeTransfer(address(msg.sender), user.amount);
262
             if (_pid == 0){
263
                 depositedQNS = depositedQNS.sub(user.amount);
264
265
             emit EmergencyWithdraw(msg.sender, _pid, user.amount);
266
             user.amount = 0;
267
             user.rewardDebt = 0;
268
```

Listing 3.5: MasterChef::emergencyWithdraw()

Note that other routines share the same issue, including deposit(), withdraw(), enterStaking(), and leaveStaking().

**Recommendation** Apply necessary reentrancy prevention by utilizing the nonReentrant modifier to block possible re-entrancy.

**Status** This issue has been fixed in the following commit: 7397641.



## 3.4 Sybil Attacks on QuantoSwap Token Voting

• ID: PVE-004

• Severity: High

• Likelihood: Medium

• Impact: High

• Target: QuantoSwapToken

• Category: Business Logic [7]

• CWE subcategory: CWE-841 [4]

#### Description

The QuantoSwapToken tokens can be used for governance in allowing for users to cast and record the votes. Moreover, the QuantoSwapToken contract allows for dynamic delegation of a voter to another, though the delegation is not transitive. When a submitted proposal is being tallied, the number of votes are counted via getPriorVotes().

Our analysis shows that the current governance functionality is vulnerable to a new type of so-called Sybil attacks. For elaboration, let's assume at the very beginning there is a malicious actor named Malice, who owns 100 QNS tokens. Malice has an accomplice named Trudy who currently has 0 balance of QNS. This Sybil attack can be launched as follows:

```
188
        function _delegate(address delegator, address delegatee)
189
             internal
190
        {
191
             address currentDelegate = _delegates[delegator];
192
             uint256 delegatorBalance = balanceOf(delegator); // balance of underlying QNSs (
                not scaled);
193
             _delegates[delegator] = delegatee;
194
195
             emit DelegateChanged(delegator, currentDelegate, delegatee);
196
197
             _moveDelegates(currentDelegate, delegatee, delegatorBalance);
198
        }
199
        function _moveDelegates(address srcRep, address dstRep, uint256 amount) internal {
200
201
            if (srcRep != dstRep && amount > 0) {
202
                 if (srcRep != address(0)) {
203
                     // decrease old representative
204
                     uint32 srcRepNum = numCheckpoints[srcRep];
205
                     uint256 srcRepOld = srcRepNum > 0 ? checkpoints[srcRep][srcRepNum - 1].
                         votes : 0:
206
                     uint256 srcRepNew = srcRepOld.sub(amount);
207
                     _writeCheckpoint(srcRep, srcRepNum, srcRepOld, srcRepNew);
208
                 }
209
210
                 if (dstRep != address(0)) {
211
                     // increase new representative
212
                     uint32 dstRepNum = numCheckpoints[dstRep];
```

Listing 3.6: QuantoSwapToken::\_delegate()

- 1. Malice initially delegates the voting to Trudy. Right after the initial delegation, Trudy can have 100 votes if he chooses to cast the vote.
- 2. Malice transfers the full 100 balance to  $M_1$  who also delegates the voting to Trudy. Right after this delegation, Trudy can have 200 votes if he chooses to cast the vote. The reason is that the QuantoSwapToken contract's transfer() does NOT \_moveDelegates() together. In other words, even now Malice has 0 balance, the initial delegation (of Malice) to Trudy will not be affected, therefore Trudy still retains the voting power of 100 QNS. When  $M_1$  delegates to Trudy, since  $M_1$  now has 100 QNS, Trudy will get additional 100 votes, totaling 200 votes.
- 3. We can repeat by transferring  $M_i$ 's  $100~{\rm QNS}$  balance to  $M_{i+1}$  who also delegates the votes to  ${\tt Trudy}$ . Every iteration will essentially add  $100~{\rm voting}$  power to  ${\tt Trudy}$ . In other words, we can effectively amplify the voting powers of  ${\tt Trudy}$  arbitrarily with new accounts created and iterated!

**Recommendation** To mitigate, it is necessary to accompany every single transfer() and transferFrom() with the \_moveDelegates() so that the voting power of the sender's delegate will be moved to the destination's delegate. By doing so, we can effectively mitigate the above Sybil attacks.

**Status** This issue has been fixed in the following commit: 7397641.

## 3.5 Implicit Assumption Enforcement In AddLiquidity()

• ID: PVE-005

• Severity: Low

• Likelihood: Low

• Impact: Low

• Target: QuantoSwapV2Router02

• Category: Coding Practices [6]

• CWE subcategory: CWE-628 [2]

#### Description

In the QuantoSwapV2Router02 contract, the addLiquidity() routine (see the code snippet below) is provided to add amountADesired amount of tokenA and amountBDesired amount of tokenB into the pool as liquidity via the QuantoSwapV2Router02::addLiquidity() routine. To elaborate, we show below the related code snippet.

```
32
        function _addLiquidity(
33
            address tokenA,
34
            address tokenB,
35
            uint amountADesired,
36
            uint amountBDesired,
37
            uint amountAMin,
38
            uint amountBMin
39
        ) internal virtual returns (uint amountA, uint amountB) {
40
            // create the pair if it doesn't exist yet
41
            if (IQuantoSwapV2Factory(factory).getPair(tokenA, tokenB) == address(0)) {
42
                IQuantoSwapV2Factory(factory).createPair(tokenA, tokenB);
43
            }
44
            (uint reserveA, uint reserveB) = QuantoSwapV2Library.getReserves(factory, tokenA
                , tokenB);
45
            if (reserveA == 0 && reserveB == 0) {
46
                (amountA, amountB) = (amountADesired, amountBDesired);
47
            } else {
48
                uint amountBOptimal = QuantoSwapV2Library.quote(amountADesired, reserveA,
                    reserveB);
                if (amountBOptimal <= amountBDesired) {</pre>
49
50
                    require(amountBOptimal >= amountBMin, 'QuantoSwapV2Router:
                        INSUFFICIENT_B_AMOUNT');
51
                    (amountA, amountB) = (amountADesired, amountBOptimal);
52
53
                    uint amountAOptimal = QuantoSwapV2Library.quote(amountBDesired, reserveB
                        , reserveA);
54
                    assert(amountAOptimal <= amountADesired);</pre>
55
                    require(amountAOptimal >= amountAMin, 'QuantoSwapV2Router:
                        INSUFFICIENT_A_AMOUNT');
56
                    (amountA, amountB) = (amountAOptimal, amountBDesired);
57
                }
58
            }
59
       function addLiquidity(
```

```
address tokenA,
62
            address tokenB.
63
            uint amountADesired,
64
            uint amountBDesired,
65
            uint amountAMin,
            uint amountBMin,
66
67
           address to,
68
           uint deadline
69
       ) external virtual override ensure(deadline) returns (uint amountA, uint amountB,
           uint liquidity) {
70
            (amountA, amountB) = _addLiquidity(tokenA, tokenB, amountADesired,
                amountBDesired, amountAMin, amountBMin);
71
            address pair = QuantoSwapV2Library.pairFor(factory, tokenA, tokenB);
72
            TransferHelper.safeTransferFrom(tokenA, msg.sender, pair, amountA);
73
            TransferHelper.safeTransferFrom(tokenB, msg.sender, pair, amountB);
74
            liquidity = IQuantoSwapV2Pair(pair).mint(to);
75
```

Listing 3.7: QuantoSwapV2Router02::addLiquidity()

It comes to our attention that the QuantoSwapV2RouterO2 has implicit assumptions on the \_addLiquidity () routine. The above routine takes two amounts: amountXDesired and amountXMin. The first amount amountXDesired determines the desired amount for adding liquidity to the pool and the second amount amountXMin determines the minimum amount of used assets. There are two implicit conditions, i.e., amountADesired >= amountAMin and amountBDesired >= amountBMin. However, if these two conditions are not met, current logic will not trigger reverts because the code above performs asymmetric checks for these amounts. Hence, without stating these assumptions, slippage control for some trades on SquadSwap V2 Router may not be checked and may not be taken into account at all in certain scenarios.

Recommendation Make the requirement of amountADesired >= amountAMin and amountBDesired >= amountBMin explicitly in the addLiquidity() function.

**Status** This issue has been fixed in the following commit: 6ba929d.

## 3.6 Inconsistent Fee Share Calculation in QuantoSwapV2LiquidityMathLibrary

• ID: PVE-006

• Severity: Medium

• Likelihood: High

• Impact: Medium

• Target: QuantoSwapV2LiquidityMathLibrary

• Category: Business Logic [7]

• CWE subcategory: CWE-841 [4]

#### Description

In this section, we examine a specific QuantoSwapV2LiquidityMathLibrary library that is designed to provide a number of convenience functions, e.g. computing their exact value in terms of the underlying tokens. Our analysis of this library exposes a specific function computeLiquidityValue() for improvement.

To elaborate, we show below this computeLiquidityValue() routine. This routine implements a rather straightforward logic in computing the liquidity value given all six parameters of the pair, i.e., reservesA, reservesB, totalSupply, liquidityAmount, feeOn, and kLast. Notice that this routine uses 1/3 of collected swap fee for protocol fee while default 1/2 of collected swap fee, if turned on, will be collected for protocol fee.

```
75
        function computeLiquidityValue(
76
            uint256 reservesA ,
77
            uint256 reservesB,
78
            uint256 totalSupply,
79
            uint256 liquidityAmount,
80
            bool feeOn,
81
            uint kLast
        ) internal pure returns (uint256 tokenAAmount, uint256 tokenBAmount) {
82
83
            if (feeOn && kLast > 0) {
84
                uint rootK = Babylonian.sqrt(reservesA.mul(reservesB));
85
                uint rootKLast = Babylonian.sqrt(kLast);
86
                if (rootK > rootKLast) {
87
                    uint numerator1 = totalSupply;
88
                    uint numerator2 = rootK.sub(rootKLast);
89
                    uint denominator = rootK.mul(2).add(rootKLast);
90
                    uint feeLiquidity = FullMath.mulDiv(numerator1, numerator2, denominator)
91
                    totalSupply = totalSupply.add(feeLiquidity);
                }
92
93
            }
94
            return (reservesA.mul(liquidityAmount) / totalSupply, reservesB.mul(
                liquidityAmount) / totalSupply);
95
```

Listing 3.8: QuantoSwapV2LiquidityMathLibrary::computeLiquidityValue()

**Recommendation** Revise the above computeLiquidityValue() routine to be consistent in collecting the percentage of swap fee for protocol fee.

**Status** This issue has been fixed in the following commit: 6ba929d.

## 3.7 Staking Incompatibility With Deflationary Tokens in SmartChef

• ID: PVE-007

• Severity: Undetermined

• Likelihood: N/A

• Impact: N/A

• Target: SmartChef, MasterChef

• Category: Business Logic [7]

• CWE subcategory: CWE-841 [4]

#### Description

In the QuantoSwap protocol, the SmartChef contract is designed to take users' assets and deliver rewards depending on their share. In particular, one interface, i.e., deposit(), accepts asset transfer-in and records the depositor's balance. Another interface, i.e., withdraw(), allows the user to withdraw the asset with necessary bookkeeping under the hood. For the above two operations, i.e., deposit() and withdraw(), the contract using the safeTransfer()/safeTransferFrom() routines to transfer assets into or out of its pool. This routine works as expected with standard ERC20 tokens: namely the pool's internal asset balances are always consistent with actual token balances maintained in individual ERC20 token contract.

```
710
         function deposit(uint256 _amount) public {
711
             PoolInfo storage pool = poolInfo[0];
712
             UserInfo storage user = userInfo[msg.sender];
714
             require(user.amount.add(_amount) <= limitAmount, 'Limit amount');</pre>
716
             updatePool(0);
717
             if (user.amount > 0) {
718
                 uint256 pending = user.amount.mul(pool.accQNSPerShare).div(PRECISION_FACTOR)
                     .sub(user.rewardDebt);
719
                 if(pending > 0) {
720
                     rewardToken.safeTransfer(address(msg.sender), pending);
721
                 }
722
             }
723
             if(_amount > 0) {
724
                 pool.lpToken.safeTransferFrom(address(msg.sender), address(this), _amount);
725
                 user.amount = user.amount.add(_amount);
726
             }
727
             user.rewardDebt = user.amount.mul(pool.accQNSPerShare).div(PRECISION_FACTOR);
```

```
730 emit Deposit(msg.sender, _amount);
731 }
```

Listing 3.9: SmartChef::deposit())

However, there exist other ERC20 tokens that may make certain customization to their ERC20 contracts. One type of these tokens is deflationary tokens that charge certain fee for every transfer. As a result, this may not meet the assumption behind asset-transferring routines. In other words, the above operations, such as <code>deposit()</code> and <code>withdraw()</code>, may introduce unexpected balance inconsistencies when comparing internal asset records with external ERC20 token contracts. Apparently, these balance inconsistencies are damaging to accurate and precise portfolio management of the pool and affects protocol-wide operation and maintenance.

Specially, if we take a look at the updatePool() routine. This routine calculates pool.accQNSPerShare via dividing the reward by lpSupply, where the lpSupply is derived from pool.lpToken.balanceOf(address(this)) (line 692). Because the balance inconsistencies of the pool, the lpSupply could be 1 Wei and thus may yield a huge pool.accQNSPerShare as the final result, which dramatically inflates the pool's reward.

```
687
         function updatePool(uint256 _pid) public {
688
             PoolInfo storage pool = poolInfo[_pid];
689
             if (block.number <= pool.lastRewardBlock) {</pre>
690
                 return:
691
             }
692
             uint256 lpSupply = pool.lpToken.balanceOf(address(this));
693
             if (lpSupply == 0) {
694
                 pool.lastRewardBlock = block.number;
695
696
             }
697
             uint256 multiplier = getMultiplier(pool.lastRewardBlock, block.number);
698
             uint256 QNSReward = multiplier.mul(rewardPerBlock).mul(pool.allocPoint).div(
                 totalAllocPoint);
699
             pool.accQNSPerShare = pool.accQNSPerShare.add(QNSReward.mul(PRECISION_FACTOR).
                 div(lpSupply));
700
             pool.lastRewardBlock = block.number;
701
```

Listing 3.10: SmartChef::updatePool()

One mitigation is to measure the asset change right before and after the asset-transferring routines. In other words, instead of bluntly assuming the amount parameter in safeTransfer() or safeTransferFrom() will always result in full transfer, we need to ensure the increased or decreased amount in the pool before and after the safeTransfer() or safeTransferFrom() is expected and aligned well with our operation. Though these additional checks cost additional gas usage, we consider they are necessary to deal with deflationary tokens or other customized ones if their support is deemed

necessary.

Another mitigation is to regulate the set of ERC20 tokens that are permitted into QuantoSwap for support. However, certain existing stable coins may exhibit control switches that can be dynamically exercised to convert into deflationary.

**Recommendation** Check the balance before and after the safeTransfer() or safeTransferFrom() call to ensure the book-keeping amount is accurate. An alternative solution is using non-deflationary tokens as collateral but some tokens (e.g., USDT) allow the admin to have the deflationary-like features kicked in later, which should be verified carefully.

**Status** This issue has been resolved as the team confirms no plan to support deflationary tokens.

## 3.8 Trust Issue of Admin Keys

• ID: PVE-008

• Severity: Medium

• Likelihood: Medium

• Impact: Medium

• Target: Multiple Contracts

• Category: Security Features [5]

• CWE subcategory: CWE-287 [1]

#### Description

In the QuantoSwap protocol, there is a privileged owner account that plays a critical role in governing and regulating the system-wide operations (e.g., configure various parameters, set the migrator, manage reward pools, and execute privileged operations). It also has the privilege to control or govern the flow of assets managed by this protocol. Our analysis shows that the privileged account needs to be scrutinized. In the following, we examine the privileged account and the related privileged accesses in current contracts.

```
139
         function updateMultiplier(uint256 multiplierNumber) public onlyOwner {
140
             BONUS_MULTIPLIER = multiplierNumber;
141
        }
142
         function newDevAddress(address _devaddr) public onlyOwner {
143
144
             devaddr = _devaddr;
145
        7
146
147
         function newQNSPerBlock(uint256 newAmount) public onlyOwner {
148
             require(newAmount > 1, 'Bad per block');
149
             QNSPerBlock = newAmount;
150
         }
151
```

```
152
         function set( uint256 _pid, uint256 _allocPoint, bool _withUpdate) public onlyOwner
153
             if (_withUpdate) {
154
                 massUpdatePools();
155
156
             totalAllocPoint = totalAllocPoint.sub(poolInfo[_pid].allocPoint).add(_allocPoint
157
             poolInfo[_pid].allocPoint = _allocPoint;
        }
158
159
160
         // Set the migrator contract. Can only be called by the owner.
161
         function setMigrator(IMigratorChef _migrator) public onlyOwner {
162
             migrator = _migrator;
163
```

Listing 3.11: Example Privileged Functions in MasterChef

Note that if the privileged owner account is a plain EOA account, this may be worrisome and pose counter-party risk to the exchange users. A multi-sig account could greatly alleviate this concern, though it is still far from perfect. Specifically, a better approach is to eliminate the administration key concern by transferring the role to a community-governed DAO. In the meantime, a timelock-based mechanism can also be considered as mitigation.

**Recommendation** Promptly transfer the privileged account to the intended DAO-like governance contract. All changed to privileged operations may need to be mediated with necessary timelocks. Eventually, activate the normal on-chain community-based governance life-cycle and ensure the intended trustless nature and high-quality distributed governance.

**Status** This issue has been mitigated as the team makes use of a multisig to act as the privileged owner.

# 4 Conclusion

In this audit, we have analyzed the design and implementation of the QuantoSwap protocol, which is a DEX that is forked from the popular UniswapV2 with the further integration with MasterChef to engage protocol users. Harnessing the power of multiple liquidity pools, QuantoSwap LPs are fungible and are composable for a wide variety of utility, including on-chain collateralization, farming, as well as other yield strategies. The current code base is well structured and neatly organized. Those identified issues are promptly confirmed and addressed.

Meanwhile, we need to emphasize that Solidity-based smart contracts as a whole are still in an early, but exciting stage of development. To improve this report, we greatly appreciate any constructive feedbacks or suggestions, on our methodology, audit findings, or potential gaps in scope/coverage.

# References

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