# Guide to Convexity in Optimization

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# 1 Introduction to Convexity

Convexity is one of the most important concepts in optimization because it allows us to:

- Guarantee that local minima are global minima
- Develop efficient optimization algorithms
- Analyze problems theoretically

### 1.1 What Does "Convex" Mean?

Imagine a simple bowl shape - this is convex. A saddle shape is non-convex. The key property is that if you draw a line between any two points on a convex object, the line stays entirely within the object.

### 2 Convex Sets

#### 2.1 Basic Definition

**Definition 2.1** (Convex Set). A set  $C \subseteq \mathbb{R}^n$  is **convex** if the line segment between any two points in C lies entirely in C:

$$\forall x, y \in C, \forall \theta \in [0, 1] : \theta x + (1 - \theta)y \in C$$

### 2.2 Advanced Examples

• Linear Matrix Inequality Solution Set:

$$\left\{ x \in \mathbb{R}^k : \sum_{i=1}^k x_i A_i \preceq B \right\}$$

where  $A_i, B \in \mathbb{S}^n$ . Convexity follows because it's an affine pre-image of the PSD cone.

• Conditional Probability Distributions: The set of conditional distributions derived from convex joint distributions remains convex through linear-fractional transformations.

## 2.3 Key Theorems

**Theorem 2.2** (Separating Hyperplane). For disjoint convex sets C, D, there exists a, b such that:

$$C \subseteq \{x : a^{\top}x \le b\}, \quad D \subseteq \{x : a^{\top}x \ge b\}$$

**Theorem 2.3** (Solution Set Convexity). For convex optimization problems, the set of optimal solutions  $X_{opt}$  is convex.

*Proof.* Let  $x, y \in X_{\text{opt}}$ . For  $t \in [0, 1]$ , define z = tx + (1 - t)y. By convexity:

$$g_i(z) \le tg_i(x) + (1-t)g_i(y) \le 0$$
  
 $Az = tAx + (1-t)Ay = b$   
 $f(z) \le tf(x) + (1-t)f(y) = f^*$ 

Thus  $z \in X_{\text{opt}}$ .

### 2.4 Examples of Convex Sets

• Norm balls: All points within a certain distance from center

$$\{x: ||x|| \le r\}$$

Hyperplanes and halfspaces:

$$\{x : a^{\top}x = b\}, \quad \{x : a^{\top}x < b\}$$

• Polyhedrons: Solutions to systems of linear inequalities

$$\{x : Ax \leq b, Cx = d\}$$

• Positive semidefinite cone (for matrices):

$$\mathbb{S}^n_+ = \{X \in \mathbb{R}^{n \times n} : X = X^\top, X \succeq 0\}$$

### 2.5 Operations that Preserve Convexity

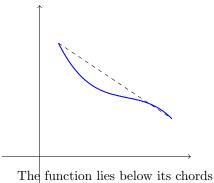
- Intersection: If  $C_1$  and  $C_2$  are convex, then  $C_1 \cap C_2$  is convex
- Affine transformations: For convex C, the set  $\{Ax + b : x \in C\}$  is convex
- Linear-fractional transformations: For convex C, the set  $\left\{\frac{Ax+b}{c^{\top}x+d}:x\in C\right\}$  is convex

#### **Convex Functions** 3

#### **Basic Definition** 3.1

**Definition 3.1** (Convex Function). A function  $f: \mathbb{R}^n \to \mathbb{R}$  is convex if its domain is convex and:

 $f(\theta x + (1 - \theta)y) \le \theta f(x) + (1 - \theta)f(y) \quad \forall x, y \in dom(f), \theta \in [0, 1]$ 



#### 3.2 **Extended Properties**

**Theorem 3.2** (Partial Minimization). If g(x,y) is convex and C is convex, then  $f(x) = \min_{y \in C} g(x, y)$  is convex.

### • Composition Rules:

- Convex + non-decreasing  $\circ$  convex = convex
- Convex + non-increasing  $\circ$  concave = convex
- Affine composition preserves convexity: f(Ax + b)

#### 3.3 Important Properties

**Theorem 3.3** (Epigraph Characterization). f is convex  $\iff$  its epigraph  $epi(f) = \{(x, t) : f(x) \le t\}$  is convex.

**Theorem 3.4** (First-Order Characterization). For differentiable f:

$$f \ convex \iff f(y) \ge f(x) + \nabla f(x)^{\top} (y - x) \quad \forall x, y$$

Theorem 3.5 (Second-Order Characterization). For twice differentiable f:

$$f \ convex \iff \nabla^2 f(x) \succeq 0 \quad \forall x$$

## 3.4 Important Function Classes

• Indicator Function:

$$I_C(x) = \begin{cases} 0 & x \in C \\ \infty & \text{otherwise} \end{cases}$$

Convex when C is convex

• Support Function:

$$I_C^*(x) = \sup_{y \in C} x^\top y$$

Always convex (even for non-convex C)

### 3.5 Examples of Convex Functions

• Affine functions:  $f(x) = a^{T}x + b$ 

• Quadratic functions:  $f(x) = \frac{1}{2}x^{T}Qx + b^{T}x + c$  (when  $Q \succeq 0$ )

• Norms:  $||x||_p$  for  $p \ge 1$ 

• Exponential:  $e^{ax}$ 

• Log-sum-exp:  $\log(\sum e^{x_i})$ 

## 3.6 Strong Convexity

**Definition 3.6** (Strongly Convex Function). A function f is  $\mu$ -strongly convex if:

$$f(y) \ge f(x) + \nabla f(x)^{\top} (y - x) + \frac{\mu}{2} ||y - x||^2$$

This means the function grows at least as fast as a quadratic function.

**Theorem 3.7** (Hessian Characterization). For twice differentiable f:

$$f \text{ $\mu$-strongly convex} \iff \nabla^2 f(x) \succeq \mu I \quad \forall x$$

### 3.7 Algorithmic Implications

• Enables linear convergence rates

• Provides error bounds:

$$\frac{1}{2\mu}\|\nabla f(x)\|^2 \geq f(x) - f^*$$

• Guarantees unique global minimum

• Leads to faster convergence in optimization algorithms

• For twice differentiable f, equivalent to  $\nabla^2 f(x) \succeq \mu I$ 

# 4 Convex Optimization Problems

### 4.1 Standard Form

A convex optimization problem has the form:

$$\min_{x} \quad f(x)$$
 s.t.  $g_{i}(x) \leq 0, \quad i = 1, ..., m$  
$$a_{i}^{\top} x = b_{j}, \quad j = 1, ..., p$$

where:

- f and  $g_i$  are convex functions
- Equality constraints are affine

#### 4.2 Problem Transformations

- Eliminating Equality Constraints: Express  $x = My + x_0$  where M spans nullspace of A
- Introducing Slack Variables: Convert  $g_i(x) \le 0$  to  $g_i(x) + s_i = 0$  with  $s_i \ge 0$
- **Geometric Programming**: Non-convex posynomials become convex after log-transform:

$$\min_{y} \log \sum e^{a_k^{\top} y + b_k} \quad \text{s.t.} \quad \log \sum e^{c_i^{\top} y + d_i} \le 0$$

## 4.3 Key Properties

**Theorem 4.1** (Global Optimality). For convex problems, any local minimum is global.

**Theorem 4.2** (Solution Uniqueness). If f is strictly convex, the solution is unique.

**Theorem 4.3** (Optimality condition). : For differentiable f,  $x^*$  is optimal iff

$$\nabla f(x^*)^{\top} (y - x^*) \ge 0 \quad \forall y \ feasible$$

#### 4.4 Solution Characteristics

- Feasibility: A point x is feasible if  $x \in \bigcap \text{dom}(g_i)$ ,  $g_i(x) \leq 0$ , and Ax = b
- Active Constraints: Inequality  $g_i$  is active at x if  $g_i(x) = 0$
- $\epsilon$ -Suboptimality : x is  $\epsilon$ -suboptimal if  $f(x) \leq f^* + \epsilon$

### 4.5 Important Examples

• Linear Programs (LP):

$$\min_{x} c^{\top} x$$
 s.t.  $Ax \le b$ 

• Quadratic Programs (QP):

$$\min_{x} \frac{1}{2} x^{\top} Q x + c^{\top} x \quad \text{s.t.} \quad A x \le b$$

• Semidefinite Programs (SDP):

$$\min_{X} \langle C, X \rangle$$
 s.t.  $\mathcal{A}(X) = b, X \succeq 0$ 

# 5 Applications

### 5.1 Lasso Regression

$$\min_{\beta} \|y - X\beta\|_2^2 + \lambda \|\beta\|_1$$

- $\bullet$  Convex despite non-differentiable  $\ell_1\text{-norm}$
- Promotes sparse solutions

### 5.2 Robust Lasso Variant

With Huber loss for outlier resistance:

$$\min_{\beta} \sum_{i=1}^{n} \rho(y_i - x_i^{\top} \beta) + \lambda \|\beta\|_1$$

where  $\rho(z) = \begin{cases} \frac{1}{2}z^2 & |z| \leq \delta \\ \delta |z| - \frac{1}{2}\delta^2 & \text{otherwise} \end{cases}$ . Remains convex as Huber loss is convex.

### 5.3 Support Vector Machines

$$\min_{w,b} \frac{1}{2} \|w\|_2^2 + C \sum_{i=1}^n \max(0, 1 - y_i(w^\top x_i + b))$$

- Convex objective with piecewise linear constraints
- Can be rewritten as a quadratic program

### 5.4 SVM Duality

The dual SVM formulation reveals convex structure:

$$\max_{\alpha} \sum_{i=1}^{n} \alpha_i - \frac{1}{2} \sum_{i,j} \alpha_i \alpha_j y_i y_j x_i^{\top} x_j \quad \text{s.t.} \quad 0 \le \alpha_i \le C$$

Demonstrates convex quadratic programming structure.

# 6 Why Convexity Matters

- Reliable optimization: No need to worry about getting stuck in local minima
- Efficient algorithms: Specialized methods (gradient descent, interior point) work well
- Theoretical guarantees: Can prove convergence rates and complexity bounds
- Recognizable structure: Many problems can be reformulated as convex programs

# 7 Algorithmic Implications

### 7.1 Why Convexity Helps Optimization

- No Spurious Local Minima: Gradient descent won't get stuck in poor solutions
- Predictable Convergence:
  - $-O(1/\epsilon)$  iterations for convex + smooth
  - Linear convergence  $(O(\log 1/\epsilon))$  for strongly convex
- Simple Stopping Criteria:  $\|\nabla f(x)\| < \epsilon$  guarantees near-optimality

### 7.2 Hierarchy of Convex Programs

Linear Programentic Programment Programment (CP)

Each class extends the previous one while maintaining convexity properties.

# 8 Recognizing Convexity

To verify if a problem is convex:

- 1. Check if the objective is convex (use definitions or derivative tests)
- 2. Verify all inequality constraints are convex
- 3. Ensure equality constraints are affine
- 4. Confirm the domain is convex
- **Significant Counterexample**: Geometric programs appear non-convex but become convex under logarithmic transformation

### 9 Common Mistakes

- Assuming all quadratics are convex (must check  $Q \succeq 0$ )
- Forgetting to verify convexity of constraints
- Treating equality constraints as inequalities
- Overlooking domain issues (e.g., log(x) requires x ¿ 0)

### 10 Conclusion

Convexity provides a powerful framework for optimization problems:

- Clear geometric interpretation
- Strong theoretical guarantees
- Wide range of applications
- Systematic approach to problem formulation