

P1:

$$A: 3^{1500} \bmod 11$$

$$1500(10) = 10111011100(2)$$

$$= 1 \times 2^1 + 1 \times 2^3 + 1 \times 2^4 + 1 \times 2^6 + 1 \times 2^7 + 1 \times 2^8 + 1 \times 2^{10}$$

$$= 4 + 8 + 16 + 64 + 128 + 256 + 1024$$

So,

$$3^{1500} = 3^4 \times 3^8 \times 3^{16} \times 3^{64} \times 3^{128} \times 3^{256} \times 3^{512}$$

$$3^1 \equiv 3 \bmod 11$$

$$3^2 \equiv 9 \bmod 11$$

$$3^4 \equiv 9^2 \equiv 4 \bmod 11$$

$$3^8 \equiv 4^2 \equiv 5 \bmod 11$$

$$3^{16} \equiv 5^2 \equiv 3 \bmod 11$$

$$3^{32} \equiv 3^2 \equiv 9 \bmod 11$$

$$3^{64} \equiv 9^2 \equiv 4 \bmod 11$$

...

$$3^{1024} \equiv 9^2 \equiv 4 \bmod 11$$

So,

$$3^{1500} \equiv 4 \times 5 \times 3 \times 4 \times 5 \times 3 \times 4 \bmod 11$$

$$\equiv 14400 \bmod 11$$

$$\equiv 1$$

$$B: 5^{4858} \bmod 10$$

Likewise, we use the same idea as the first one.

$$\begin{aligned} 4358_{10} &= 1000100000110_2 \\ &= 1 \times 2^1 + 1 \times 2^2 + 1 \times 2^8 + 1 \times 2^{12} \\ &= 2 + 4 + 256 + 4096 \end{aligned}$$

$$\text{So, } 5^{4358} = 5^2 \times 5^4 \times 5^{256} \times 5^{4096}$$

$$5^1 \equiv 5 \pmod{10}$$

$$5^2 \equiv 5 \pmod{10}$$

$$5^4 \equiv 5 \pmod{10}$$

$$5^8 \equiv 5 \pmod{10}$$

...

$$5^{256} \equiv 5 \pmod{10}$$

$$5^{4096} \equiv 5 \pmod{10}$$

$$\text{So, } 4358 = 4096 + 256 + 4 + 2$$

$$\begin{aligned} 5^{4358} \pmod{10} &\equiv 5 \times 5 \times 5 \times 5 \pmod{10} \\ &\equiv 5 \end{aligned}$$

$$C: 6^{22345} \pmod{7}$$

$$22345_{10} = 101011101001001_2$$

$$= 2 + 8 + 64 + 256 + 512 + 1024 + 4096 + 16384$$

So,

$$6^{22345} \pmod{7} = 6^{2+8+64+256+512+1024+4096+16384} \pmod{7}$$

$$6^{22545} = 6^4 \cdot 6^8 \cdot 6^{16} \cdot 6^{32} \cdot 6^{64} \cdot 6^{128} \cdot 6^{256} \cdot 6^{512} \cdot 6^{1024} \cdot 6^{2048} \cdot 6^{4096}$$

$$6^1 \equiv 6 \pmod{7}$$

$$6^2 \equiv 1 \pmod{7}$$

$$6^4 \equiv 1 \pmod{7}$$

$$6^8 \equiv 1 \pmod{7}$$

$$6^{16} \equiv 1 \pmod{7}$$

$$6^{32} \equiv 1 \pmod{7}$$

$$\dots$$

$$6^{4096} \equiv 1 \pmod{7}$$

$$\text{So, } 6^{22545} \equiv (1 \times 1 \times 1 \dots \times 6) \pmod{7}$$

$$\equiv 6$$

P2:

$$\begin{aligned} \text{a. } \text{GCD}(648, 124) &= \text{GCD}(124, 28) \\ &= \text{GCD}(28, 12) \\ &= \text{GCD}(12, 4) \\ &= \text{GCD}(4, 0) \\ &= 4 \end{aligned}$$

$$\begin{aligned} \text{b. } \text{GCD}(123456789, 123456788) \\ &= \text{GCD}(123456788, 1) \\ &= \text{GCD}(1, 0) \end{aligned}$$

$$\begin{aligned}
 &= 1 \\
 C &= \text{GCD}(2^{300} \cdot 3^{200}, 2^{200}) \\
 &= \text{GCD}(2^{100} \cdot 2^{200} \cdot 3^{200}, 2^{200}) \\
 &= \text{GCD}(2^{200}, 0) \\
 &= 2^{200}
 \end{aligned}$$

P3:

Let's say: we have modulus P and the base g , Alice and Bob want to send their own secret number a, b to each other, they will perform a function like

$$A = g^a \text{ mod } P,$$

P is a prime number, at least 2048 bits long,

g is a small number,

Let's take an example to clarify that.

Bob and Alice has secret number $a=3, b=6$,

$$P=17, g=4.$$

then Alice:

$$A = g^a \text{ mod } P$$

$$A = 64 \text{ mod } 17$$

$$A = 13.$$

same perform for Bob:

$$B = 4^6 \text{ mod } 17$$

$$B = 4296 \bmod 17$$

$$B = 16$$

Alice then send her result A to Bob, while Bob sends his result B to Alice. Alice then calculates the shared secret s using the number she received from Bob and her number a using the formula:

$$s = B^a \bmod p$$

$$s = 16^3 \bmod 17$$

$$s = 16$$

then do same job for Bob:

$$s = A^b \bmod p$$

$$s = 13^6 \bmod p$$

$$s = 4826809 \bmod 17$$

$$s = 16$$

Another example:

We can define a decode function like $\text{encode} = s \cdot 3$

if Bob wants to send $b = 9$.

the encode-number

$$= s \cdot 3 = 16 \cdot 3 = 48;$$

Bob will send 48 to Alice;

then Alice will decode it

using same encode function.

then, s is a shared key that is only known for Bob and Alice, and they can setup for symmetric encryption, allowing them to safely send info between themselves in a way that only they can access it.

P4:

```
public static int criticalEvents(int[] arr, int t){  
    int cnt = 0;  
    for(int i = 0; i < arr.length; i++){  
        for(int j = 0; j < arr.length; j++){  
            if(arr[i] > arr[j] * t){  
                cnt++;  
            }  
        }  
    }  
    return cnt;  
}
```