

# 1 Probability revision

## 1.1 Common random variables

**Bernoulli distribution**  $Ber(p)$ :

$P(X = x) = p^x(1 - p)^{1-x}$ ,  $x \in \{0, 1\}$ , mean =  $p$ , variance =  $p(1 - p)$

**Binomial distribution**  $Bin(n, p)$ :

$P(X = x) = \binom{n}{x} p^x(1 - p)^{n-x}$ ,  $x \in \mathbb{Z} \cap [0, n]$ , mean =  $np$ , variance =  $np(1 - p)$

**Poisson distribution**  $Pois(\lambda)$ :

$P(X = x) = e^{-\lambda} \frac{\lambda^x}{x!}$ ,  $x \in \mathbb{Z}^{\geq 0}$ , mean =  $\lambda$ , variance =  $\lambda$

**Geometric distribution**  $Geo(p)$ :

$P(X = x) = p(1 - p)^{x-1}$ ,  $x \in \mathbb{Z}^{>0}$ , mean =  $\frac{1}{p}$ , variance =  $\frac{1-p}{p^2}$

## 1.2 Property of expectation

If  $X \geq 0$ , then  $E[X] \geq 0$

If  $X \geq 0$  and  $E[X] = 0$ , then  $P(X = 0) = 1$

If  $a$  and  $b$  are constants, then  $E[a + bX] = a + bE[X]$

If  $X$  and  $Y$  are random variables, then  $E[X + Y] = E[X] + E[Y]$

$E[X]$  is the constant that minimizes  $E[(X - c)^2]$

If  $X$  is a non-negative random variable that takes integral value, then we have  $E[X] = \sum_{i=1}^{\infty} P(X \geq i)$

## 1.3 Property of variance

$Var(X) \geq 0$

If  $Var(X) = 0$ , then  $X$  is a constant.

If  $a$  and  $b$  are constants, then  $Var(a + bX) = b^2 Var(X)$

If  $X$  and  $Y$  are random variables, then  $Var(X + Y) = Var(X) + Var(Y) + 2Cov(X, Y)$ , for independent  $X$  and  $Y$ , their covariance is 0.

$Var(X) = E[X^2] - (E[X])^2$

## 1.4 Moment generating function

It is the unique identifier of a random variable.

$M_X(t) = E(e^{tX})$

Calculate  $k$ -th moment of  $X$ :

$E[X^k] = \frac{d^k}{dt^k} M_X(t)|_{t=0}$

Linear transformation:

$M_{aX+b}(t) = e^{bt} M_X(at)$

MGF of binomial distribution:  $(pe^t + 1 - p)^n$

MGF of Poisson distribution:  $\exp[\lambda(e^t - 1)]$

MGF of geometric distribution:  $\frac{pe^t}{1 - (1 - p)e^t}$

## 1.5 Joint distribution

**Definition:**  $p_{X,Y}(x, y) = P(X = x, Y = y) = P(\{\omega : X(\omega) = x, Y(\omega) = y\})$

**Marginal distribution:**  $p_X(x) = P(X = x) = \sum_y P(X = x, Y = y) = P(\{\omega : X(\omega) = x\})$

## 1.6 Independence

Discrete random variables  $X_1, \dots, X_n$  are independent if for any  $x_1, \dots, x_n$ ,  $P(X_1 = x_1, \dots, X_n = x_n) = \prod_{i=1}^n P(X_i = x_i)$

If  $f_1, \dots, f_n$  are functions from  $\mathbf{R}$  to  $\mathbf{R}$ , then  $Y_i = f_i(X_i)$  are also random variables, and they preserve independence.

If  $X_i$ 's are independent, then  $E[\prod_{i=1}^n X_i] = \prod_{i=1}^n E[X_i]$

If  $Z = \sum_{i=1}^n X_i$ , then  $Var(Z) = \sum_{i=1}^n Var(X_i)$ , and  $M_Z(t) = \prod_{i=1}^n M_{X_i}(t)$

## 1.7 Conditional probability

**Discrete:**  $p_{X|Y}(x|y) = \frac{P(X = x, Y = y)}{P(Y = y)}$

**Continuous:**  $f_{X|Y}(x|y) = \frac{f(x, y)}{f_Y(y)}$

**Multiplication law:**  $p_{X,Y}(x, y) = p_{X|Y}(x|y)p_Y(y) = p_{Y|X}(y|x)p_X(x)$

**Law of total probability:**  $p_X(x) = \sum_y p_{X|Y}(x|y)$

**Bayes' formula:**  $p_{Y|X}(y|x) = \frac{p_{X|Y}(x|y)p_Y(y)}{\sum_y p_{X|Y}(x|y)p_Y(y)}$

**Conditional independence:**  $P(X = x, Y = y|Z = z) = P(X = x|Z = z)P(Y = y|Z = z)$

**Law of total expectation:**  $E[X] = E[E[X|Y]]$

# 2 Markov Chain preliminaries

## 2.1 Definition

A stochastic process with discrete state space  $S$  and discrete time set  $T$  satisfying the Markovian property.

## 2.2 Markovian property

Given  $X_n$ , what happened afterwards ( $t > n$ ) is independent with what happened before ( $t < n$ ). Mathematically, for any set of state  $i_0, \dots, i_{n-1}, i, j \in S$  and  $n \geq 0$ ,  $P(X_{n+1} = j|X_n = i, X_{n-1} = i_{n-1}, \dots, X_0 = i_0) = P(X_{n+1} = j|X_n = i)$

## 2.3 One-step transition probability

Let  $p_{i,j}^{n,n+1} = P(X_{n+1} = j|X_n = i)$  denote the one-step transition probability.

## 2.4 m-step transition probability

Given all one-step probabilities, we can calculate m-step probabilities.

$$p_{i,j}^{n,n+m} = p(X_{n+m} = j|X_n = i) \quad (1)$$

$$= \sum_{k \in S} p(X_{n+m} = j, X_{n+1} = k|X_n = i) \quad (2)$$

$$= \sum_{k \in S} p(X_{n+m} = j|X_{n+1} = k, X_n = i)P(X_{n+1} = k|X_n = i) \quad (3)$$

$$= \sum_{k \in S} p(X_{n+m} = j|X_{n+1} = k)p(X_{n+1} = k|X_n = i) \quad (4)$$

$$= \sum_{k \in S} p_{k,j}^{n+1,n+m} p_{i,k}^{n,n+1} \quad (5)$$

$$= \sum_{k_1, k_2, \dots, k_{m-1} \in S} [p_{i,k_1}^{n,n+1} p_{k_1,k_2}^{n+1,n+2} \dots p_{k_{m-1},j}^{n+m-1,n+m}] \quad (6)$$

Intuitively, the m-step probability is to consider the sum of probabilities over all m-length paths from state  $i$  to state  $j$  from time  $n$  to time  $n + m$ . From (3) to (4) we exploit Markovian property, from (5) to (6) is to expand iteratively according to previous steps.

## 2.5 Specification of Markov Chain

We define a Markov Chain by the index set  $T$ , the state space  $S$ , and all one-step transition probabilities  $p_{i,j}^{n,n+1}$ . The one-step transition probabilities could be denoted by a one-step transition probability matrix  $\mathbf{P}^{n,n+1} = (p_{i,j}^{n,n+1})$

## 2.6 Properties of transition probability matrix

All entries are non-negative.

Row sum is always 1.

**Chapman-Kolmogorov equations:**

$$\mathbf{p}^{n,n+m+1} = \mathbf{p}^{n,n+1} \mathbf{p}^{n+1,n+m+1} = \mathbf{p}^{n,n+m} \mathbf{p}^{n+m,n+m+1}$$

## 2.7 Stationary Markov Chain

A Markov Chain is stationary if the one-step transition probability matrix does not change when  $n$  changes. Mathematically, we have  $p_{i,j}^{n,n+1} = p_{i,j} \quad \forall n \in T$

The Chapman-Kolmogorov equation reduces to  $\mathbf{P}^{n,n+m} = \mathbf{P}^{(m)} = \mathbf{P}^m$

# 3 First step analysis

## 3.1 Initial distribution

Suppose the head of the Markov Chain  $X_0$  follows a certain distribution  $\pi_0$ , then  $X_n|X_0 \sim \pi_0$  will follow a distribution of  $\pi_0 \mathbf{P}^n$ .

## 3.2 Terminologies

**Absorbing state:** A state  $i$  which satisfies that  $p_{i,j} = 0 \quad \forall j \neq i$ .

**Stopping time:**  $T = \min\{n \geq 0 : X_n = i\}$

## 3.3 Trick: move one step forward

Consider the case of gambler's ruin with winning probability  $p$ . Intuition suggests that  $P(X_T = 0|X_1 = 2) = P(X_T = 0|X_0 = 2)$

Define the process  $Y_n = X_{n+1}$ , hence  $P(X_T = 0|X_1 = 2) = P(Y_{T-1} = 0|Y_0 = 2)$ , then we have  $Y_{T-1} = X_T = 0 \implies T_Y = T - 1$ . Since  $\{Y_n\}$  and  $\{X_n\}$  have the same probabilistic structure, we have  $P(X_T = 0|X_1 = 2) = P(Y_{T-1} = 0|Y_0 = 2) = P(Y_{T_Y} = 0|Y_0 = 2) = P(X_T = 0|X_0 = 2)$ , which completes our proof. By induction, the identity could be generalised from 1 step to any finite  $k$  steps. Similar process could be applied to derive the relation  $E[T|X_0 = i] + 1 = E[T|X_1 = i]$

The implication of this technique is that now we can express our quantity of interest as a linear combination of other related quantities, differentiated by the initial condition. We then will obtain a system of linear equations by moving one step forward, and solve the quantity.

## 3.4 General analysis procedure

1. Define  $|S|$  terms  $a_i = E[\sum_{n=0}^T g(X_n)|X_0 = i]$  for each  $i \in S$
2. Apply the law of total expectation to  $a_i$ , conditional on  $X_1$ , we have  $a_i = \sum_{k \in S} (g(i) + E[\sum_{n=1}^T g(X_n)|X_0 = i, X_1 = k])P(X_1 = k|X_0 = i)$
3.  $P(X_1 = k|X_0 = i) = p_{i,k}$  in  $\mathbf{P}$ .
4. Consider the process  $\{Y_n\}$ . It is stochastically equivalent with  $\{X_n\}$ . Further by Markovian property, we have  $E[\sum_{n=1}^T g(X_n)|X_0 = i, X_1 = k] = E[\sum_{n=1}^T g(X_n)|X_1 = k] = E[\sum_{n=0}^{T_Y} g(Y_n)|Y_0 = k] = E[\sum_{n=0}^T g(X_n)|X_0 = k] = a_k$
5. Combine the two parts, we have  $a_i = \sum_{k \in S} (g(i) + a_k)p_{i,k} = g(i) + \sum_{k \in S} a_k p_{i,k}$ .
6. Solve the system and get the result.

## 3.5 Delve into the gambler's ruin

**General case:** The gambler will stop with  $N$  dollars in hand, or he is broke.

For each game, the winning probability is  $p$ . Denote  $q = 1 - p$ .

The gambler starts with  $0 < k < N$  in hand.

**Chance of getting broke:**  $u_k = P(X_T = 0|X_0 = k)$

If game is fair, then  $p/q = 1$ , solve and we obtain  $u_k = 1 - \frac{k}{N}$ , which agrees with our intuition that when  $k$  increases, the probability of getting broke is smaller. When  $N$  increases, the probability of getting broke converges to 1 even if the game is fair. It implies that suppose the gambler plays without stopping ( $N = +\infty$ ), then he will get broke.

If game is not fair, solve and we obtain  $u_k = 1 - \frac{1-(q/p)^k}{1-(q/p)^N}$

If  $q > p$ , then it is approximately  $1 - \frac{1}{(q/p)^{N-k}}$ , again, probability converges to 1 if  $N \rightarrow +\infty$ .

If  $q < p$ , then it is approximately  $(q/p)^k - (q/p)^N$ , when  $N \rightarrow +\infty$ , probability converges to  $(q/p)^k$ , which means we have a chance of not getting broke.

In both cases, probability reduces with  $k$  increasing. It agrees with our intuition that with more one has at the beginning, the smaller probability that one will get broke.

In both cases, probability reduces with  $p$  increasing. It agrees with our intuition that with larger winning probability, the less likely one will get broke.

**Quitting time:**  $v_k = E[T|X_0 = k]$

If game is fair, solve and we obtain  $v_k = k(N - k)$ , although we know that  $P(stop) \geq P(broke) = 1$ , the expectation of time to stop still goes to infinity.

If game is not fair, solve and we obtain  $v_k = \frac{1}{p-q}[\frac{N(1-(q/p)^k)}{1-(q/p)^N} - k]$ .

If  $p > q$ ,  $v_k \rightarrow \frac{N-k}{p-q} \rightarrow +\infty$  when  $N \rightarrow +\infty$ . It agrees with our intuition that with higher winning probability, it takes more time to stop the game. With  $k$  approaching  $N$ ,  $v_k$  decreases. It agrees with our intuition that with higher  $k$ , we are more likely to reach  $N$  early and quit the game.

If  $p < q$ ,  $v_k \rightarrow \frac{k}{q-p}$  when  $N \rightarrow +\infty$ , which implies that when the winning probability is less than  $\frac{1}{2}$ , the game is expected to stop in finite time. When  $k$  increases,  $v_k$  increases. It agrees with our intuition that with higher  $k$ , one will sustain longer before getting broke.

## 4 Classification of states

### 4.1 Accessibility

**Definition:** State  $j$  is accessible from state  $i$  if  $\mathbf{P}_{i,j}^m > 0$  for some  $m > 0$ . Note that state  $i$  is accessible from itself since we define  $\mathbf{P}^0 = \mathbf{I}$ . If state  $i$  is also accessible from state  $j$ , then they communicate with each other, and we denote by  $i \leftrightarrow j$ . Communication is an equivalence relation, which means that it exhibits reflexivity, symmetry and transitivity.

### 4.2 Communication class

Each communication class is an equivalence class defined by the communication relation. It contains all states that are communicating with each other. Suppose we have two classes  $\mathcal{C}_1, \mathcal{C}_2$ , and  $i \in \mathcal{C}_1 \rightarrow j \in \mathcal{C}_2$ , then any state in  $\mathcal{C}_2$  is accessible from any state in  $\mathcal{C}_1$ . Note that converse cannot hold in this case, otherwise they will become one communication class.

### 4.3 Reducible chain

**Definition:** An MC is irreducible if all states communicate with each other, otherwise reducible.

### 4.4 Return probability

Probability that starting from state  $i$ , and revisit state  $i$  at the  $n$ -th step, which is  $P_{i,i}^n$ . If state  $i$  is transient, then  $P_{i,i}^n \rightarrow 0$  when  $n \rightarrow +\infty$ , which means that in long term we will never revisit state  $i$ .

### 4.5 First return probability

The probability that starting from state  $i$ , the first revisit to state  $i$  occurs at the  $n$ -th step. Mathematically, we denote by  $f_{ii}^n = P(X_1 \neq i, X_2 \neq i, \dots, X_{n-1} \neq i, X_n = i|X_0 = i)$ . We define  $f_{ii}^0 = 0$ . Since first revisit at different steps are disjoint events, and if a state is transient, we may not return to it in long run, we naturally have  $P(\text{first revisit at step } n < +\infty) = \sum_{n=1}^{\infty} f_{ii}^n \leq 1$ , and  $f_{ii}^n \leq P_{ii}^n \ \forall n \in \mathbf{Z}^{>0}$ .

Condition on the first revisit, we have  $P_{ii}^n = \sum_{k=0}^n f_{ii}^k P_{ii}^{n-k}$ .

### 4.6 Recurrent and transient states

Consider  $f_{ii} = \sum_{n=0}^{\infty} f_{ii}^n = \lim_{N \rightarrow \infty} \sum_{n=0}^N f_{ii}^n$ . It denotes the probability of revisiting  $i$  in the future.

**Definition:** A state  $i$  is recurrent if  $f_{ii} = 1$ , and transient if  $f_{ii} < 1$ .

If the flow will revisit state  $i$  for sure, then state  $i$  is recurrent. Otherwise, the flow will finally leave state  $i$  and never come back.

Let  $N_i$  denote the number of times that the flow revisits state  $i$ .  $N_i = \sum_{n=1}^{\infty} I(X_n = i)$ , then recurrence means  $E[N_i|X_0 = i] = \infty$ , and transient means  $E[N_i|X_0 = i] < \infty$ .

**Theorem:** For a state  $i$ , consider the expected number of visits to  $i$ ,  $E[N_i|X_0 = i]$ , then if  $f_{ii} < 1$ (i.e. $i$  is transient), then  $E[N_i|X_0 = i] = \frac{f_{ii}}{1-f_{ii}}$ . If  $f_{ii} = 1$ (i.e. $i$  is recurrent), then  $E[N_i|X_0 = i] = \infty$ .

Note that  $E[N_i|X_0 = i] = E[\sum_{n=1}^{\infty} I(X_n = i|X_0 = i)] = \sum_{n=1}^{\infty} E[I(X_n = i|X_0 = i)] = \sum_{n=1}^{\infty} P(X_n = i|X_0 = i) = \sum_{n=1}^{\infty} P_{ii}^n$ .

**Theorem:** State  $i$  is recurrent iff  $\sum_{n=1}^{\infty} P_{ii}^n = \infty$  For recurrent states, it means a significant probability to return. For transient states, the series will converge to 0, which means that the probability of coming back vanishes in the long run.

### 4.7 Summary of recurrence

Recurrent  $\iff f_{ii} = 1 \iff \sum_{n=1}^{\infty} P_{ii}^n = \infty \iff E[N_i|X_0 = i] = \infty$

Transient  $\iff f_{ii} < 1 \iff \sum_{n=1}^{\infty} P_{ii}^n < \infty \iff E[N_i|X_0 = i] < \infty$

### 4.8 States of the same class

**Theorem:** States in the same communication class are either all transient or all recurrent. We could interpret transient/recurrent status as a status for the whole class. An MC with finite states must have at least one recurrent class.

## 5 Long run performance

### 5.1 Periodicity

**Definition:** For a state  $i$ , consider  $\{n \geq 1 : P_{ii}^n > 0\}$ , we define the period of state  $i$ ,  $d(i)$ , as the greatest common divisor of the set. If the set is empty, we define  $d(i) = 0$ . If  $d(i) = 1$ , we say that state  $i$  is aperiodic.

**Theorem:** If  $i \leftrightarrow j$ , then  $d(i) = d(j)$ .

**Theorem:**  $\exists N$  such that  $P_{ii}^{Nd(i)} > 0$ , and  $\forall n \geq N$ ,  $P_{ii}^{nd(i)} > 0$ .

**Theorem:** If  $\exists m > 0$  such that  $P_{ji}^m > 0$ , then for sufficiently large  $n$ , we have

$$P_{ji}^{m+nd(i)} > 0.$$

### 5.2 Regular Markov Chain

**Definition:** An MC is regular if  $\exists k > 0$ , such that all elements of  $\mathbf{P}^k$  are strictly positive.

It means that the flow can achieve any state from any state at step  $k$ .

**Theorem:** If a Markov Chain is irreducible, aperiodic and with finite states, then it is regular.

### 5.3 Main theorem

**Theorem:** Suppose  $\mathbf{P}$  is a regular transition probability matrix with states  $S = \{1, 2, \dots, N\}$ , then

- $\lim_{n \rightarrow \infty} p_{ij}^n$  exists.
- The limit does not depend on  $i$ . Hence, we can denote it by  $\pi_j = \lim_{n \rightarrow \infty} p_{ij}^n$ .
- $\sum_{k=1}^N \pi_k = 1$ . We call it as the limiting distribution.
- The limits  $\pi = (\pi_1, \pi_2, \dots, \pi_N)$  are the solutions of the system  $\pi_j = \sum_{k=1}^N \pi_k P_{kj}$ ,  $j = 1, 2, 3, \dots, N$ ,  $\sum_{k=1}^N \pi_k = 1$ . In matrix form, it is to solve  $\pi P = \pi$ ,  $\sum_{k=1}^N \pi_k = 1$
- The limiting distribution  $\pi$  is unique.

We should interpret  $\pi_j$  as the marginal probability that the flow is in state  $j$  for a long run. In other words, suppose we observe the flow at some time, then the probability that it stays in state  $j$  is  $\pi_j$ . That probability is not related to the original state and the time point. It gives the limit of  $\mathbf{P}^n$ , in long run, each row of  $\mathbf{P}^n$  is the same, which is  $\pi$ .  $\pi$  can be seen as the long run proportion of time in every state.

### 5.4 Stationary distribution

When the MC is not regular, it is possible that  $|S| = \infty$  and  $\pi = 0$ , but it is also possible that we can still find a non-trivial  $\pi$ . It means, if the initial states have a distribution  $\pi$ , then after any steps, the chain also has a distribution  $\pi$  on the states, which is called the stationary distribution.

**Definition:** Consider a Markov Chain with state space  $S = \{1, 2, \dots\}$  and the transition probability matrix  $P$ . A distribution  $(p_1, p_2, \dots)$  on  $S$  is called a stationary distribution, if it satisfies that if  $P(X_n = i) = p_i, i = 1, 2, \dots$ , then  $P(X_{n+1} = i) = p_i, i = 1, 2, \dots$ .

It means that if the initial states have a distribution  $\pi$ , then after any steps, the chain still has a distribution  $\pi$  on the states. Note that a irregular MC may have more than one stationary distributions. To solve for possible  $\pi$ , note that here we also have to solve the system  $\pi P = \pi$

## 6 Helpers

**Binomial expansion:**  $(a + b)^n = \sum_{i=0}^n \binom{n}{i} a^{n-i} b^i$

**Maclaurin expansion:**  $f(x) = \sum_{i=0}^{\infty} \frac{x^i}{i!} f^{(i)}(0)$

**Arithmetic series:**  $\sum_{i=1}^n a_i = \frac{n(a_1 + a_n)}{2}$

**Geometric series:**  $\sum_{i=1}^n a_i = \frac{a(1 - r^n)}{1 - r} (r \neq 1) \ \sum_{i=1}^{\infty} a_i = \frac{a}{a - r} (|r| < 1)$