

50003 - Models of Computation - Lecture 7

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Lecture Recording

Lecture recording is available here

Note for reader

We will reference to state by set $State \triangleq (Var \rightarrow \mathbb{N})$.

Lemmas

Lemma

A small proven proposition that can be used in a proof. Used to make the proof smaller.

Also know as an "auxiliary theorem" or "helper theorem".

Corollary

A theorem connected by a short proof to another existing theorem.

If B is can be easily deduced from A (or is evident in A's proof) then B is a corollary of A.

Lemmas

1. $\forall r \in \mathbb{N}. \forall E_1, E'_1, E_2 \in SimpleExp. [E_1 \rightarrow^r E'_1 \Rightarrow (E_1 + E_2) \rightarrow^r (E'_1 + E_2)]$
2. $\forall r, n \in \mathbb{N}. \forall E_2, E'_2 \in SimpleExp. [E_2 \rightarrow^r E'_2 \Rightarrow (n + E_2) \rightarrow^r (n + E'_2)]$

Corollaries

1. $\forall n_1 \in \mathbb{N}. \forall E_1, E_2 \in SimpleExp. [E_1 \rightarrow^* n_1 \Rightarrow (E_1 + E_2) \rightarrow^* (n_1 + E_2)]$
2. $\forall n_1, n_2 \in \mathbb{N}. \forall E_2 \in SimpleExp. [E_2 \rightarrow^* n_2 \Rightarrow (n_1 + E_2) \rightarrow^* (n_1 + n_2)]$
3. $\forall n, n_1, n_2, \in \mathbb{N}. \forall E_1, E_2 \in SimpleExp. [E_1 \rightarrow^* n_1 \wedge E_2 \rightarrow^* n_2 \wedge n = n_1 + n_2 \Rightarrow (E_1 + E_2) \rightarrow^* n]$

Connecting \Downarrow and \rightarrow^* for SimpleExp

$$\forall E \in SimpleExp, n \in \mathbb{N}. [E \Downarrow n \Leftrightarrow E \rightarrow^* n]$$

We prove each direction of implication separately. Fir we prove by induction over E using the property P :

$$P(E) =_{def} \forall n \in \mathbb{N}. [E \Downarrow n \Rightarrow E \rightarrow^* n]$$

Base Case

Take arbitrary $n \in \mathbb{N}$ to show $P(n) - n \Downarrow n \Rightarrow n \rightarrow^* n$.

- (1) $n \rightarrow^0 n$ (n is in the normal form)
- (2) $n \Downarrow n$ (n is in the normal form)
- (3) $n \Downarrow n \wedge n \rightarrow^* n$ (By 1 & 2)
- (4) $n \Downarrow n \Rightarrow n \rightarrow^* n$ (By 3)

Inductive Step

Take some arbitrary E, E_1, E_2 such that $E = E_1 + E_2$.

Inductive Hypothesis

$$\forall n_1 \in \mathbb{N}. [E_1 \Downarrow n_1 \Rightarrow E_1 \rightarrow^* n_1]$$

$$\forall n_2 \in \mathbb{N}. [E_2 \Downarrow n_2 \Rightarrow E_2 \rightarrow^* n_2]$$

To show $P(E)$: $\forall n \in \mathbb{N}. [(E_1 + E_2) \Downarrow n \Rightarrow (E_1 + E_2) \rightarrow^* n]$.

- (1) Assume $(E_1 + E_2) \Downarrow n$
- (2) $\exists n_1, n_2 \in \mathbb{N}. [E_1 \Downarrow n_1 \wedge E_2 \Downarrow n_2]$ (By 1 & definition of B-ADD)
- (3) $E_1 \rightarrow^* n_1$ (By 2 & IH)
- (4) $E_2 \rightarrow^* n_2$ (By 2 & IH)
- (5) Chose some $n \in \mathbb{N}$ such that $n = n_1 + n_2$
- (6) $(E_1 + E_2) \rightarrow^* n$ (By 3,4,5 Corollary 3)
- (7) $E \rightarrow^* n$ (By 6, definition of E)

Hence assuming $E \Downarrow n$ implies $E \rightarrow^* n$, so $P(E)$.

Next we work the other way, to show:

$$\forall E \in \text{SimpleExp}. \forall n \in \mathbb{N}. [E \rightarrow^* n \Rightarrow E \Downarrow n]$$

- (1) Take arbitrary $E \in \text{SimpleExp}$ such that $E \rightarrow^* n$ (Initial setup)
- (2) Take some $m \in \mathbb{N}$ such that $E \Downarrow m$ (By totality of \Downarrow)
- (3) $n = m$ (By 1,2 & uniqueness of result for \rightarrow)
- (4) $E \Downarrow n$ (By 3)

It is also possible to prove this without using normalisation and determinacy, by induction on E .

Multi-Step Reductions

Lemma:

$$\forall r \in \mathbb{N}. \forall E_1, E'_1, E_2. [E_1 \rightarrow^r E'_1 \Rightarrow (E_1 + E_2) \rightarrow^r (E'_1 + E_2)]$$

To prove $\forall r \in \mathbb{N}. [P(r)]$ by induction on r :

Base Case

- Base case is $r = 0$.
- Prove that $P(0)$ holds.

Inductive Step

- Inductive Case is $r = k + 1$ for arbitrary $k \in \mathbb{N}$.
- Inductive hypothesis is $P(k)$.
- Prove $P(k + 1)$ using inductive hypothesis.

Proof of the Lemma

By induction on r : **Base Case:** Take some arbitrary $E_1, E'_1, E_2 \in \text{SimpleExp}$ such that $E_1 \rightarrow^0 E'_1$.

- (1) $E_1 = E'_1$ (By definition of \rightarrow^0)
- (2) $(E_1 + E_2) = (E'_1 + E_2)$ (By 1)
- (3) $(E_1 + E_2) \rightarrow^0 (E'_1 + E_2)$ (By definition of \rightarrow^0)

Inductive Step: Take arbitrary $k \in \mathbb{N}$ such that $P(k)$

- (1) Take arbitrary E_1, E'_1, E_2 such that $E_1 \rightarrow E'_1$ (Initial setup)
- (2) Take arbitrary E''_1 such that $E''_1 \rightarrow E'_1$
- (3) $(E_1 + E_2) \rightarrow^k (E''_1 + E_2)$ (By 2 & IH)
- (4) $(E''_1 + E_2) \rightarrow (E'_1 + E_2)$ (By 2 & rule S-LEFT)
- (5) $(E_1 + E_2) \rightarrow^{k+1} (E'_1 + E_2)$ (3,4, definition of \rightarrow^{k+1})

Determinacy of \rightarrow for Exp

We extend simple expressions configurations of the form $\langle E, s \rangle$.

$$E \in \text{Exp} ::= n|x|E + E| \dots$$

Determinacy:

$$\forall E, E_1, E_2 \in \text{Exp}. \forall s, s_1, s_2 \in \text{State}. [\langle E, s \rangle \rightarrow \langle E_1, s_1 \rangle \wedge \langle E, s \rangle \rightarrow \langle E_2, s_2 \rangle \Rightarrow \langle E_1, s_1 \rangle = \langle E_2, s_2 \rangle]$$

We prove this using property P :

$$P(E, s) \triangleq \forall E_1, E_2 \in \text{Exp}. \forall s_1, s_2 \in \text{State}. [\langle E, s \rangle \rightarrow \langle E_1, s_1 \rangle \wedge \langle E, s \rangle \rightarrow \langle E_2, s_2 \rangle \Rightarrow \langle E_1, s_1 \rangle = \langle E_2, s_2 \rangle]$$

Base Case: $E = x$

Take arbitrary $n \in \mathbb{N}$ and $s \in \text{State}$ to show $P(n, s)$

- (1) take $E_1 \in \text{Exp}, s_1 \in \text{State}$ such that $\langle n, s \rangle \rightarrow \langle E_1, s_1 \rangle$ (Initial setup)
- (2) take $E_2 \in \text{Exp}, s_2 \in \text{State}$ such that $\langle n, s \rangle \rightarrow \langle E_2, s_2 \rangle$ (Initial setup)
- (3) $n = E_1 \wedge s = s_1$ (By 1 & inversion on definition of E.NUM)
- (4) $n = E_2 \wedge s = s_2$ (By 2 & inversion on definition of E.NUM)
- (5) $E_1 = E_2 \wedge s_1 = s_2$ (By 3 & 4)
- (6) $\langle E_1, s_1 \rangle = \langle E_2, s_2 \rangle$ (By 5 & definition of configurations)

Base Case: $E = x$

Take arbitrary $x \in Var$ and $s \in State$ to show $P(n, s)$

- | | | |
|-----|---|---|
| (1) | take $E_1 \in \mathbb{N}$, $s_1 \in State$ such that $\langle x, s \rangle \rightarrow \langle E_1, s_1 \rangle$ | (Initial setup) |
| (2) | take $E_2 \in \mathbb{N}$, $s_2 \in State$ such that $\langle x, s \rangle \rightarrow \langle E_2, s_2 \rangle$ | (Initial setup) |
| (3) | $E_1 = s(x) \wedge s_1 = s$ | (By 1 & inversion on definition of E.VAR) |
| (3) | $E_2 = s(x) \wedge s_2 = s$ | (By 2 & inversion on definition of E.VAR) |
| (5) | $E_1 = E_2 \wedge s_1 = s_2$ | (By 3 & 4) |
| (6) | $\langle E_1, s_1 \rangle = \langle E_2, s_2 \rangle$ | (By 5 & definition of configurations) |

... Inductive Step ...

Syntax of Commands

$$C \in Com ::= x := E | \text{if } B \text{ then } C \text{ else } C | C; C | \text{skip} | \text{while } B \text{ do } C$$

- **Determinacy**

$$\forall C, C_1, C_2 \in Com. \forall s, s_1, s_2 \in State. [\langle C, s \rangle \rightarrow_c \langle C_1, s_1 \rangle \wedge \langle C, s \rangle \rightarrow_c \langle C_2, s_2 \rangle \Rightarrow \langle C_1, s_1 \rangle = \langle C_2, s_2 \rangle]$$

- **Confluence**

$$\forall C, C_1, C_2 \in Com. \forall s, s_1, s_2 \in State. [\langle C, s \rangle \rightarrow_c^* \langle C_1, s_1 \rangle \wedge \langle C, s \rangle \rightarrow_c^* \langle C_2, s_2 \rangle \Rightarrow \exists C' \in Com. \exists s' \in State. [\langle C_1, s_1 \rangle \rightarrow C', s'] \wedge \langle C_2, s_2 \rangle \rightarrow C', s']]$$

- **Unique Answer**

$$\forall C \in Com. s_1 s_2 \in State. [\langle C, s \rangle \rightarrow_c^* \langle \text{skip}, s_1 \rangle \wedge \langle C, s \rangle \rightarrow_c^* \langle \text{skip}, s_2 \rangle \Rightarrow s_1 = s_2]$$

- **No Normalisation**

There exist derivations of infinite length for while.

Connecting \Downarrow and \rightarrow^* for While

1. $\forall E, n \in Exp. \forall s, s' \in State. [\langle E, s \rangle \Downarrow_e \langle n, s' \rangle \Leftrightarrow \langle E, s \rangle \rightarrow_e^* \langle n, s' \rangle]$
2. $\forall B, b \in Bool. \forall s, s' \in State. [\langle B, s \rangle \Downarrow_b \langle b, s' \rangle \Leftrightarrow \langle B, s \rangle \rightarrow_b^* \langle b, s' \rangle]$
3. $\forall C \in Com. \forall s, s' \in State. [\langle C, s \rangle \Downarrow_c \langle s' \rangle \Leftrightarrow \langle C, s \rangle \rightarrow_c^* \langle \text{skip}, s' \rangle]$

For *Exp* and *Bool* we have proofs by induction on the structure of expressions/booleans.

For \Downarrow_c it is more complex as the $\Downarrow_c \Leftarrow \rightarrow_c^*$ cannot be proven using totality. Instead **complete/strong induction** on length of \rightarrow_c^* is used.