50006 - Compilers - (Prof Kelly) Lecture $4\,$

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Lecture Recording

Lecture recording is available here

Unbounded Register Use

We will generate code for arithmetic expressions that:

- Assumes there will always be enough registers.
- Handles the case when we run out of registers.
- Translates expressions while minimising number of registers required.

```
data Instruction
                                            = Add Reg Reg | Sub Reg Reg
    2
   3
                                                               Mul Reg Reg | Div Reg Reg --- Op r1 r2 -> r1 := r1 <Op> r2
                                                               AddImm Reg Int | SubImm Reg Int
    4
    5
                                                               \label{eq:mullmm} \text{MulImm Reg Int} \quad | \quad \text{DivImm Reg Int} \quad | \quad \text{Cop>Imm r c} \quad - \text{> r} \ := \ \text{r} \quad | \quad \text{Cop> c} \quad | 
                                                                                                                                                                               -- Load r1 n -> r1 := m
-- LoadImm r1 i -> r1 := i
                                                                                                                                                                                                                                                                                         - > r1 := mem[n]
    6
                                                                Load Reg Name
                                                               LoadImm Reg Int
    7
                                                                                                                                                                               -- Load r1 n
                                                                Store Reg Name
    8
                                                                                                                                                                                                                                                                                         -> mem[n] := r1
                                                                                                                                                                                -- Push r1
    9
                                                                 Push Reg
                                                                                                                                                                                                                                                                                           -> SP++; mem[SP] := r1
                                                                Pop Reg
                                                                                                                                                                                -- Pop r2
10
                                                                                                                                                                                                                                                                                          \rightarrow r1 := mem[SP]; SP-
                                                                                                                                                                             -- CompEq r1 r2 -> r1 := r1 - r2 = 0 ? 1 : 0

-- JTrue r1 l -> IF r1 = 1 THEN JUMP TO l
                                                                CompEq Reg Reg
11
12
                                                                JTrue Reg Label
                                                                                                                                                                                                                                                                                        \rightarrow IF r1 = 1 THEN JUMP TO 1
                                                                 JFalse Reg Label -- JFalse r1 l -> IF r1 = 0 THEN JUMP TO l
13
14
                                                                 Define Label
                                                                                                                                                                                 -- Assembler directive to set up label
```

We can take an input such as:

$$(100*3) + ((200*2) + 300) + (400 + (500*3))$$

	Instruction	Stack Slot			
	HISTIUCTION	3	2	1	0
0	PushImm 100				100
1	PushImm 3			3	100
2	Mul			3	300
3	PushImm 200			200	300
4	PushImm 2		2	200	300
5	Mul		2	400	300
6	PushImm 300		300	400	300
7	Add		300	700	300
8	Add		300	700	1000
9	PushImm 400		300	400	1000
10	PushImm 500		500	400	1000
11	PushImm 3	3	500	400	1000
12	Mul	3	1500	400	1000
13	Add	3	1500	1900	1000
14	Add	3	1500	1900	2900

We can use the placement of values in the stack (relative to the initial stack pointer) to assign registers. Assigning each slot to a register.

We want to provide the translator with the register to place the result in, it can use any higher registers.

	Instruction	Register				
	Instruction	R3	R2	R1	R0	
0	LoadImm R0 100				100	
1	LoadImm R1 3			3	100	
2	Mul R0 R1			3	300	
3	LoadImm R1 200			200	300	
4	LoadImm R2 2		2	200	300	
5	Mul R1 R2		2	400	300	
6	LoadImm R2 300		300	400	300	
7	Add R1 R2		300	700	300	
8	Add R0 R1		300	700	1000	
9	LoadImm R1 400		300	400	1000	
10	LoadImm $R2~500$		500	400	1000	
11	LoadImm R3 3	3	500	400	1000	
12	Mul R2 R3	3	1500	400	1000	
13	Add R1 R2	3	1500	1900	1000	
14	Add R0 R1	3	1500	1900	2900	

Register Improvements

We could improve our generated code if there were instructions available for in place constant application.

IA32 with Immediate Operand

```
1 movl $3, %eax
2 imull $3, %eax
3 addl $4, %eax
```

Our Simple Assembly

	<u> </u>
	PushImm R0 3
2	PushImm R1 3
3	Mul R0 R1
4	PushImm R1 4
5	Add R0 R1

	Instruction	Register				
	Instruction	R3	R2	R1	R0	
0	LoadImm 0 3				3	
1	MulImm 0 100				300	
3	LoadImm 1 2			2	300	
4	MulImm 1 200			400	300	
5	AddImm 1 300			700	300	
6	$\mathrm{Add}\ 0\ 1$			700	1000	
7	LoadImm $1\ 3$			3	1000	
8	MulImm 1 500			1500	1000	
9	AddImm 1 400			1900	1000	
10	Add 0 1			1900	2900	

Code for Translation

```
translateOp :: Op -> (Int -> Int -> Instruction)
    translateOp Plus = Add
   translateOp Minus = Sub
3
   translateOp Times = Mul
   translateOp Divide = Div
5
   translateOpImm :: Op -> (Int -> Int -> Instruction)
8
   translateOpImm\ Plus\ =\ AddImm
    translateOpImm Minus = SubImm
10
   translateOpImm Times = MulImm
   translateOpImm Divide = DivImm
11
12
    transExp \ :: \ Exp \ -\!\!\!> \ Reg \ -\!\!\!> \ [\ Instruction\ ]
13
   transExp (Const n) r = [LoadImm r n]
14
   transExp (Ident id) r = [Load r id]
15
16
17
     - Only allow for - unary operato (e.g -3)
18
    transExp\ (Unop\ Minus\ e)\ r\ =\ transExp\ e\ r\ +\! +\ [MulImm\ r\ (-1)]
19
   transExp (Unop _ _)
20
      = error "(transExp) Only '-' unary operator supported"
21
22
     - As * and + are commutative, the order does not matter
   transExp (BinOp Times (Const n) e) r = transExp e r ++ [MulImm r n]
24
   transExp (BinOp Plus (Const n) e) r = transExp e r ++ [AddImm r n]
25
26
      Can run left hand, then do right hand with immediate operand
27
   transExp (BinOp op e (Const n)) r = transExp e r ++ [translateOpImm op r n]
     - General case for two expressions
29
   transExp (BinOp op e1 e2) r
30
     = transExp e1 r ++ transExp e2 (r+1) ++ [translateOp op r (r+1)]
```

Bounded Number of Registers

Accumulator Machine

Has a single register (Accumulator) upon which arithmetic instructions can be applied.

```
data Instruction = Add | Sub | Mul | Div
        AddImm Int | SubImm Int | MulImm Int | DivImm Int
2
3
        CompEq
                        — CompEq
                                      -> Acc :=
4
        Push
                           Push
                                      -> SP--; mem[SP] := Acc
                                     -> Acc := mem[SP]; SP++
5
        Pop
                        — Pop
        Load Name
                        — Load n
                                     \rightarrow Acc := mem[n]
7
        LoadImm Int
                       — Load i
                                      \rightarrow Acc := i
                       -- Store n -> mem[n] := Acc
8
         Store Name
                       — Jump l \rightarrow PC := l
9
        Jump Label
        JTrue Label — JTrue l -> IF Acc = 1 THEN JUMP TO l
JFalse Label — JFalse l -> IF Acc = 0 THEN JUMP TO l
10
11
        Define Label — Assembler directive to set up label
12
```

	Instruction	Acc	Stack		
	IIIsti uction		1	0	
0	LoadImm 500	500			
1	MulImm 3	1500			
3	$AddImm\ 400$	1900			
4	Push	1900		1900	
5	LoadImm 200	200		1900	
6	MulImm 2	400		1900	
7	AddImm~300	700		1900	
8	Push	700	700	1900	
9	LoadImm 100	100	700	1900	
10	MulImm 3	300	700	1900	
11	Add	1000	700	1900	
12	Add	2900	700	1900	

```
1
    transOpImm :: Op \rightarrow (Int \rightarrow Instruction)
    transOpImm Plus = AddImm
    transOpImm Minus = SubImm
    transOpImm \ Times = MulImm
    transOpImm Divide = DivImm
    transOp :: Op -> Instruction
8
    transOp Plus = Add
    transOp Minus = Sub
9
    transOp Times = Mul
    transOp Divide = Div
11
12
    transExp :: Exp -> [Instruction]
13
    transExp (Const n) = [LoadImm n]
transExp (Ident x) = [Load x]
14
15
    transExp (Unop Minus e) = transExp e ++ [MulImm (-1)]
16
17
18
     - Can only use Minus unary operator (e.g -3)
    transExp (Unop _ _)
19
      = error "(transExp) Only '-' unary operator supported"
20
21
      - If constant on the left, can use immediate operand
22
23
    transExp (BinOp op e (Const n)) = transExp e ++ [transOpImm op n]
24
25
     - With commutative operator, can switch order to use immediate operand
    transExp\ (BinOp\ Times\ (Const\ n)\ e)\ =\ transExp\ e\ +\!\!\!\!+\ [MulImm\ n]
26
    transExp \ (BinOp \ Plus \ (Const \ n) \ e) = transExp \ e \ +\!\!\!\!+ \ [AddImm \ n]
27
28
      General case for two expressions
    transExp (BinOp op e1 e2) = transExp e2 ++ Push : transExp e1 ++ [transOp op]
```

Limited Register Set

One solution is to combine the register and accumulator strategies

```
1 data Instruction
2 = Add Reg Reg | Sub Reg Reg
3 | Mul Reg Reg | Div Reg Reg -- Op r1 r2 -> r1 := r1 <Op> r2
4 | AddImm Reg Int | SubImm Reg Int
5 | MulImm Reg Int | DivImm Reg Int -- <Op>Imm r c -> r := r <Op> c
```

```
AddStack Reg | SubStack Reg
6
7
             MulStack Reg | DivStack Reg
                                                     \langle Op \rangle Imm \ r \ c \ -> \ r := \ r \ \langle Op \rangle \ mem[SP]; \ SP
                                                       \rightarrow r1 := mem[n]
                                   -- Load r1 n
            Load Reg Name
8
                                   -- LoadImm r1 i -> r1 := i
9
            LoadImm Reg Int
                                                      -> mem[n] := r1
                                   -- Load r1 n
10
             Store Reg Name
            Push Reg
                                   -- Push r1
                                                        \rightarrow SP++; mem[SP] := r1
11
12
            Pop Reg
                                   -- Pop r2
                                                        \rightarrow r1 := mem[SP]; SP—
            CompEq Reg Reg — CompEq r1 r2 \rightarrow r1 := r1 \rightarrow r2 = 0 ? 1 : ( JTrue Reg Label — JTrue r1 l \rightarrow IF r1 = 1 THEN JUMP TO l
13
14
             JFalse Reg Label — JFalse r1 l \rightarrow IF r1 = 0 THEN JUMP TO l
15
                                  -- Assembler directive to set up label
16
             Define Label
```

- When free register sremain, use the register machine strategy.
- When the limit is reached (one register left to use as accumulator), switch to accumulator strategy.

This results in most expressions using full benefit of registers, while very large expressions can still be correctly executed.

```
translateOp :: Op -> (Int -> Int -> Instruction)
    translateOp Plus = Add
    translateOp Minus = Sub
    translateOp Times = Mul
 5
    translateOp Divide = Div
 6
    translateOpImm :: Op -> (Int -> Int -> Instruction)
8
    translateOpImm Plus = AddImm
9
    translateOpImm Minus = SubImm
10
    translateOpImm Times = MulImm
    translateOpImm Divide = DivImm
11
12
    translateOpStack :: Op -> (Int -> Instruction)
13
    translateOpStack Plus = AddStack
14
    translateOpStack Minus = SubStack
15
    translateOpStack Times = MulStack
16
17
    translateOpStack Divide = DivStack
18
    maxReg :: Int
19
20
    maxReg = 10
21
22
    transExp :: Exp -> Reg -> [Instruction]
23
24
     - No need to check maxreg as we only use one register (a reg or the last one)
25
    transExp (Const n) r = [LoadImm r n]
    transExp (Ident id) r = [Load r id]
    transExp \ (Unop \ Minus \ e) \ r \ = \ transExp \ e \ r \ +\!\!\!\! + \ [MulImm \ r \ (-1)]
27
28
    transExp (Unop _ _)
     = error "(transExp) Only '-' unary operator supported"
29
30
     - As * and + are commutative, the order does not matter, only use one register so no
31
       → need to check maxReg
32
    transExp (BinOp Times (Const n) e) r = transExp e r ++ [MulImm r n]
33
    transExp (BinOp Plus (Const n) e) r = transExp e r ++ [AddImm r n]
34
35
     - Can run left hand, then do right hand with immediate operand
36
    transExp (BinOp op e (Const n)) r = transExp e r ++ [translateOpImm op r n]
37
     - General case for two expressions, we need to take into account the registers used.
    transExp (BinOp op e1 e2) r
```

```
| r == maxReg = transExp e2 r ++ Push r : transExp e1 r ++ [translateOpStack op r] | otherwise = transExp e1 r ++ transExp e2 (r+1) ++ [translateOp op r (r+1)]
```