50001 - Algorithm Analysis and Design - Lecture $10\,$

Oliver Killane

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Lecture Recording
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Lecture recording is available here

List lookup

As you ca see !! costs O(n) as it may traverse the entire list.

If we want this access to be faster, we can use trees:

```
data Tree a = Tip | Leaf a | Node Int (Tree a) (Tree a)
3
    node :: Tree a \rightarrow Tree a \rightarrow Tree a
4
    node l r = Node (size l + size r) l r
6
    instance List Tree where
7
      toList :: Tree a -> [a]
      toList Tip = []
8
      toList (Leaf n) = [n]
9
10
      toList (Fork n l r) = toList l ++ toList r
11
12
       - Invariant: size Node n a b = n = size a + size b
      length :: Tree a -> Int
length Tip = 0
13
14
      length (Leaf _) = 1;
15
      length (Node n _{-} _{-}) = n;
16
17
      (!!) :: Tree a -> Int -> a
18
      (Leaf x) !! 0 = x
19
      (Node n l r) !! k
20
         k < m = 1!!k
21
          otherwise = r!! (k-m)
22
23
        where m = length l
24
25
         case for Tip !! n or Leaf !! >0
      _ !! _ = error "(!!): Cannot get list index"
```

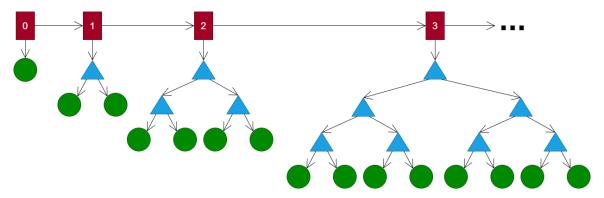
This costs $O(\log n)$ as each recursive call acts on half of the remaining list.

However we have difficulty with insertion:

```
Insert Quickly n: t = node(Leafn)t Effectively becomes a linked list, \log n search time ruined.
Insert Slowly Rebalance tree (e.g AVL tree) Complex and no longer O(1) insert.
```

Random Access Lists

A list containing elements that are either nothing, or a perfect tree with size the same as 2^{index} .



The empty tree can be represented by a Tip value (from the notes), or using type Maybe(Tree) (from the lecture) where $Tree\ a = Leaf\ x \mid Node\ n\ l\ r$.

When we add to a tree, we add to the first element of the **RAList**, if the invariant is breached (no longer perfect tree of size 2^0) it can be combined with the next list over (if empty, place, else combine and repeat).

This way while the worst case insert is O(n), our amortized complexity is O(1) much as with increment.

```
data Tree a = Tip | Leaf a | Node Int (Tree a) (Tree a)
 3
     type RAList a = [Tree a]
 4
     instance List RAList where
        toList :: RAList a -> [a]
toList (RaList ls) = concatMap toList ls
 6
 7
 8
        \begin{array}{lll} fromList & :: & [\,a\,] \ -\!\!\!\!> & RAList \ a \\ fromList & = & foldr \ (:) & empty \end{array}
 9
10
11
        empty :: RAList a
12
13
        empty = []
14
        (:) :: a \rightarrow RAList a \rightarrow RAList a
15
        n : [] = [Leaf n]
16
        n : (RAList ls) = RAList (insertTree (Leaf n) ls)
17
18
          where
19
             insertTree :: Tree a -> [Tree a] -> [Tree a]
             insertTree t ([])
20
             \begin{array}{lll} insertTree & t & ([]) & = [t] \\ insertTree & t & (Tip:ls) & = t:ls \end{array}
21
             insertTree t (t ': ts) = Tip: insertTree (node t t') ts
22
23
24
        length :: RAList a -> Int
25
        length (RAList ls) = foldr ((+) . length) 0 ls
26
27
        (!!) :: RAList a -> Int -> a
        (RAList []) !! _ = error "(!!): empty list" (RAList (x:xs)) !! n
28
29
            isEmpty x = (RAList ts) !! k
30
31
             n < m
                         = x !! n
32
             otherwise = (RAList xs) !! (n-m)
          where m = length x
33
```