

$\textcircled{2}_{\text{a.}}$	$\begin{array}{c cc c} & L & C & R \\ \hline + & 3,1 & 2,2 & 0,0 \\ B & 2,2 & 0,0 & 3,1 \end{array}$
-------------------------------	--

For Player 2, R is Strictly dominated by L

	L	C
+	3,1	2,2
B	2,2	0,0

Now for Player 1, B is Strictly dominated by +:

	L	C
+	3,1	2,2

correct and now for Player 2, L is Strictly dominated by C, so the IDS-Strategy is (+, c)

3.

	L	C	R
+	5,7	3,2	1,6
B	2,1	3,3	4,6

No, there is no Strictly or Weakly dominant strategy for either player.

Q.

center is dominated by a mixed strategy,
but then no strategy is dominated in the reduced game

One such mixed strategy is $1/2L + 1/2R$, but there is a range of values. The best one would have slightly more L than R as the payoff is higher.

		C	R
		4, 6	0, 6
		4, 3	1, 2
C.	L	C	R
A	4, 6	0, 6	3, 5
B	4, 3	1, 2	1, 3

There is no Strictly dominated strategy. If we use IWDS, it depends on which column we eliminate first, as Both C and L weakly dominates C & R:

deleting center first results in (top, left)
deleting right first, gives (bottom, left)

The option i said is also okay, deleting the column strategies first leads to multiple outcomes

		L
		4, 6
		4, 3
A	4, 6	
B	4, 3	

It's Strictly dominated But the Payoffs from B & A are the same for Player 2 so both survive.

IWDS is also unsatisfactory because it is order-dependent and can remove Nash-equilibrium in some games.

d. Strictly dominated strategies are never best responses so they cannot be part of a Nash Equilibrium.

And if only one strategy remains after all elimination then, each player's remaining strategy must be the best response. Hence that must be a unique Nash Equilibrium.

d, e: correct

e. Find a Nash equilibrium, each player's strategy being the best response. And a best response cannot be strictly dominated as it would not be the best response then.

		L	R
		U	0, 1
		D	3, 0
U		0, 1	2, 2
D		3, 0	0, 3

(U, R) is the pure Nash equilibrium,
L is strictly dominated by R:

		R
		U
		D
U		2, 2
D		0, 3

D is strictly dominated by U. **in the reduced game**

As (U, R) is the result of IDG, it is
the unique Nash equilibrium.

		L	R
		U	0, 3
		D	3, 0
U		0, 3	2, 2
D		3, 0	0, 3

There is no pure Nash equilibrium.

Let r be the probability that Column plays R, then the payoffs for row are:

$$U_1(U) = 0 \cdot r + 0(1-r) + 2r = 2r$$

$$U_1(D) = 3(1-r) + 0r = 3 - 3r$$

If $r \leq \frac{3}{5}$, playing D is better.
Therefore, the optimal probability of row playing the move D is

$$\begin{aligned} d(r) &= 1 \text{ if } r < \frac{3}{5}, \\ d(r) &= [0, \frac{3}{5}] \text{ if } r = \frac{3}{5} \text{ and} \\ d(r) &= 0 \text{ if } r > \frac{3}{5} \end{aligned}$$

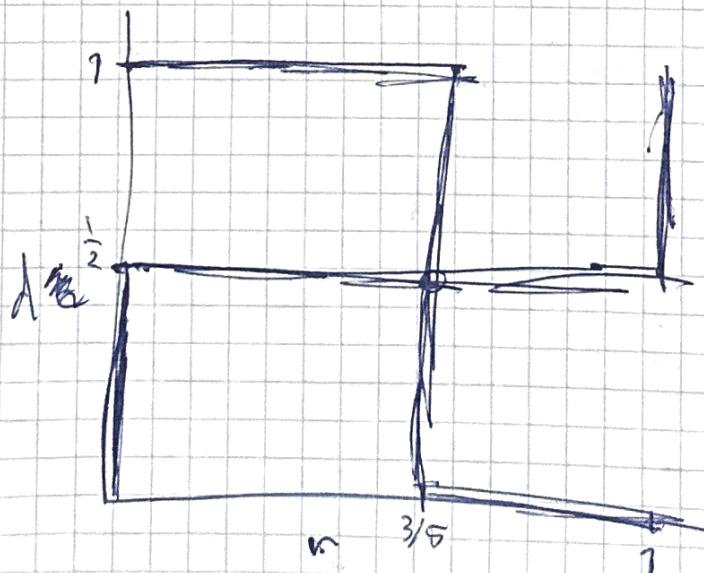
$$U_2(L) = 0d + 3(1-d) = 3 - 3d$$

~~if $d > \frac{1}{2}$~~

$$U_2(R) = 3d + 2(1-d) = 2 + d$$

$$\hat{r}(d) = 0 \text{ if } d < \frac{1}{2}, \quad \hat{r}(d) = [0, 1] \text{ if } d = \frac{1}{2}, \quad \hat{r}(d) = 1 \text{ if } d > \frac{1}{2}$$

**r is correct,
d should be 1/4**



Looks like a swastika lol

So the mixed equilibrium is at $r = \frac{3}{5}$, $d = \frac{1}{2}$

notation: $(3/5R + 2/5L, 1/4D + 3/4D)$

c.

$$U_1(U) = 0(1-r) + 2r = 2r$$

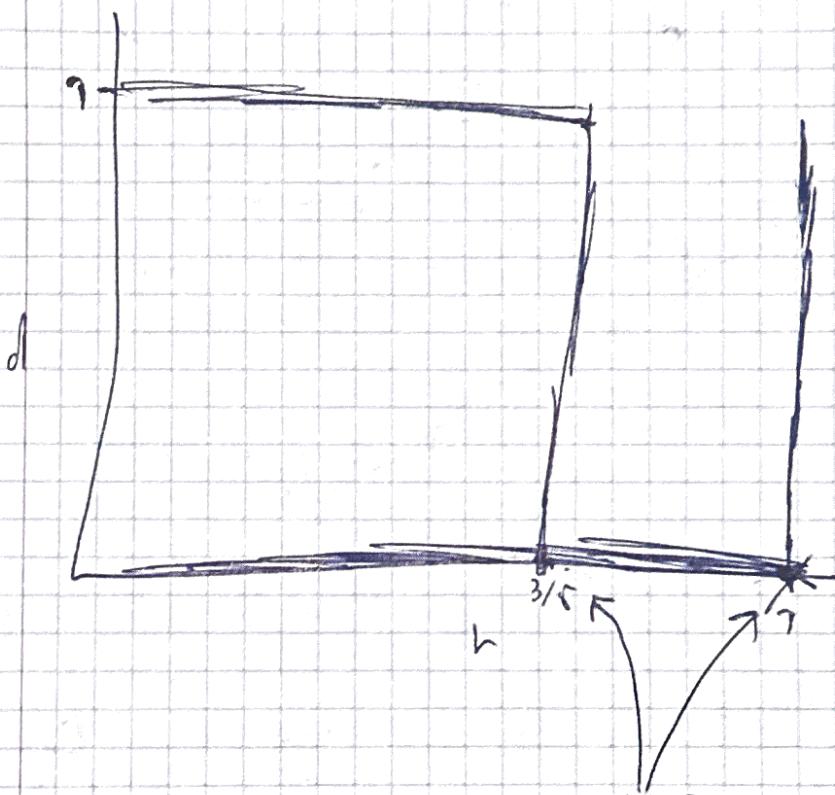
$$U_1(D) = 3(1-r) + 0r = 3 - 3r$$

So d stays the same.

$$U_2(L) = od + 2(1-d) = 2 - 2d$$

$$U_2(R) = 3d + 2(1-d) = 2 - d$$

$$\tilde{v}(d) = \text{max} \quad \text{if } d > 0, \quad \tilde{v}(d) = [0, 1] \quad \text{if } d = 0$$



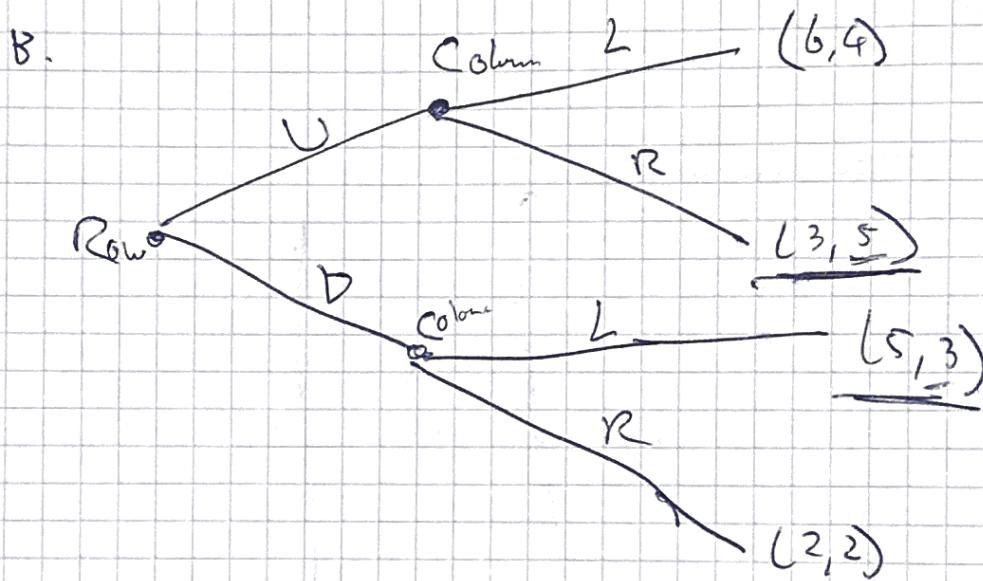
good!
Column equilibrium is when overlapping, so the Nash
Column plays R with probability between $\frac{3}{5}$ and 1

6.

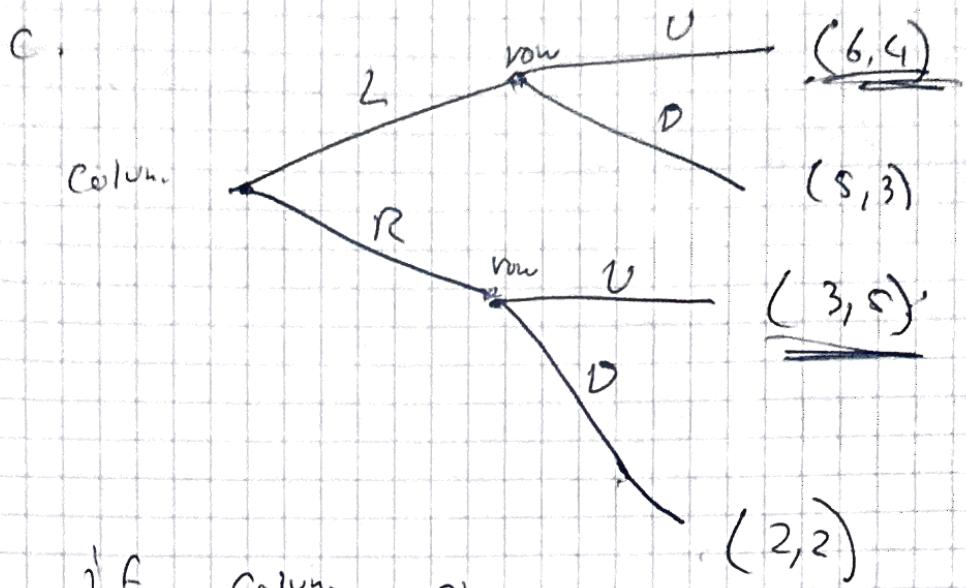
	L	R	
U	6, 4 3, 5	3, 5	2
D	S, 3 2, 2	2, 2	U 6, 9 3, 5
	D	S, 3 <u>2, 2</u>	

May be

so (U, R) as they are the best responses. say: its IDS so unique



If row plays U , column would play R ,
 if row plays D , column would play R .
 (D, L) give row a best payoff than (U, R)
 SPE: (D, RL)
 So there is the Job Game - Perfect Equilibrium.
 It is a Nash-equilibrium of the sequential game as they are the best responses,
 But not of the original simultaneous game.



If Column plays L, Row plays U,
 If Column plays R, Row also plays U.
 As ~~(UU)~~ (U, R) leads to the high
 Payoff, that is the SPE: (UU, R)
 (R, UU)

Q10

- (D, D) For both, D ~~is~~ strictly dominant.
- IF the game is played a finite "N" number of times, there is only one Nash Equilibrium and it is to ~~defect~~ always defect for both players.

In the final round, each player will defect as there is no point in cooperating, if you now step back one period this is true again as you know that both are gonna defect anyway in the next period etcetera.

c. ~~The Payoff from always Cooperating is~~

$$S + SS + SS^2 + SS^3 \dots$$

~~the~~

the Value of Cooperating forever is

$$S + SS + SS^2 + SS^3 + SS^4 \dots$$

$$= \frac{S}{1-S}$$

if a Player deviates, they get 8 that period and afterwards 4 forever!

$$8 + 48 + 48^2 + 48^3 + 48^4 \dots$$

$$= 8 + 8 \frac{4}{1-8}$$

$$\frac{5}{1-S} = 8 + 8 \frac{4}{1-S}$$

$$5 = 8 - 8S + 4S = 8 - 4S$$

$$4S = 3$$

$$S = \frac{3}{4}$$

So if ~~the~~ $8 < \frac{3}{4}$, "triggering"
is profitable and so the Nash equilibrium

d. only if indeed $\delta > \frac{3}{4}$, then cooperation is
after the Nash equilibrium and the Subgame
where one defects is also Nash equilibrium.
But if $\delta \leq \frac{3}{4}$, the first cooperation is not
a best-response, so it is not SPE as
SPE requires a Nash equilibrium in every Subgame.

Also: Because of how it is phrased, if player 1 plays D, and player 2 plays C in a round:

The strategy says play C for player 1 in the next round but that is not nash, should be D in the next round.

Though this is not very relevant because it would never happen that player 1 plays D and player 2 plays C but it is technically a subgame

So even if delta is greater or equal to 3/4, it is not SPNE, so the answer to the question is NO

IF you want to word it to get SPNE you should say: if ANY (including yourself) player plays D, you play D in the next round

- c. - Minmax Punishment: the lowest payoff a player
i can be forced to receive by others, assuming
that i is maximizing their payoff. The
strategy is the one that yields this.
- A strategy is individually rational if for every
player, the payoff is higher than the minmax
punishment payoff.

A Payoff is a feasible pair of feasible payoffs

I don't know the last one

f. The folk theorem is that ~~when~~ cooperation
is possible with a certain strategy and patience.