

III) (I won't do a Partial Because it'll need to
do a lot of assumptions)

$P \ Q \ R$	$(P \leftrightarrow (Q \rightarrow R)) \leftrightarrow ((R \rightarrow Q) \leftrightarrow R)$							
T T T		T	T				T	T
T T F		F	F				T	F
T F T		F	F				F	F
T F F		T	T				F	T
F T T		F	F				F	F
F F T		T	F				F	T
F F F		T	F				T	T
P F F		F	T				T	F

→ tautology

IV

$P \ Q$	$\neg(P \rightarrow Q) \leftrightarrow (P \wedge \neg Q)$								
1.1 + F	<u>t₁</u>	<u>t₅</u>	<u>F₃</u>	<u>F₄</u>	+ T	+ t ₆	<u>t₂</u>	<u>t₇</u>	<u>F₈</u>
1.2									
2.1	<u>t₇</u>	<u>t₅</u>	<u>F₃</u>	<u>F₄</u>	F	+ t ₆	<u>F₂</u>	<u>F₇</u>	?
2.2	<u>F₁</u>	?	<u>t₃</u>	<u>F₂</u>	F	+ t ₄	<u>t₂</u>	<u>t₅</u>	<u>F₆</u>

→ tautology

2.2

III) $| \phi \vee \psi |_A = F$ iff $|\phi|_A = F$ and $|\psi|_A = F$

IV) $|\phi \rightarrow \psi |_A = F$ iff $|\phi|_A = T$ and $|\psi|_A = F$.

V) $|\phi \leftrightarrow \psi |_A = F$ iff $|\phi|_A \neq |\psi|_A$

2.8

According to definition 2.9 $I \models \phi$ iff there is no L_1 -structure in which all sentences ~~other than~~ true in I are true and ϕ is false.

According to definition 2.11 a set of sentences is consistent iff there is an L_1 -structure under which all sentences in the structure are true.

It's so, if a set is inconsistent

there is no L_1 -structure under which all sentences ^{in that set} are true.

If the set containing all sentences $\vdash I$ and $\neg \phi$ is inconsistent then there is no counterexample as in definition 2.10. therefore, the argument must be valid.