

$$a. U_1 = W_0 + kA \log(1 + e_1 + e_2) - e_1$$

$$b. \max_{e_1} U_1 = W_0 + kA \log(1 + e_1 + \bar{e}_2) - e_1$$

$$\frac{\partial U_1}{\partial \bar{e}_1} = \frac{kA}{1 + e_1 + \bar{e}_2} - 1 = 0 \quad (\text{set to 0 to optimize})$$

$$\cancel{kA} = 1 + e_1 + \bar{e}_2$$

$$e_1 = kA - 1 - \bar{e}_2$$

c. the workers are identical so $e_1 = \bar{e}_2$:

$$e_1 = kA - 1 - e_1$$

$$e_1 = \frac{kA - 1}{2}$$

Do I need to prove this? $\exists \in$ there a proof?
 (It's the only way to do it I think because if you do it the other way around and get $e_2 = kA - 1 - e_1$ it is the same equation. So 2 unknowns with 1 equation.)

d. if we merge the utility of the workers:

$$U_{\text{merge}} = 2W_0 + 2kA \log(1 + e_1 + e_2) - (e_1 + e_2)$$

max
 e_1, e_2 U_{merge}

$$\frac{\partial U_{\text{merge}}}{\partial e_1} = \frac{2kA}{1 + e_1 + e_2} - 1 = 0 \quad (\text{set to 0 to maximize})$$

$$2kA = 1 + e_1 + e_2$$

$$e_1 + e_2 = 2kA - 1$$

~~With~~ With the Nash equilibrium, $e_1 + e_2 = kA - 1$

$$U_{\text{Nash, merged}} = 2W_0 + 2kA \log(1 + kA - 1) - kA + 1$$

$$U_{\text{merged}} = 2W_0 + 2kA \log(2kA) - 2kA + 1$$

$$\begin{aligned} U_{\text{merged}} - U_{\text{Nash, merged}} &= (2kA \log(2kA) - kA) - (2kA \log(kA) - kA + 1) \\ &= 2kA \log(2) + 2kA \log(kA) - kA - 2kA \log(kA) + kA - 1 \\ &= 2kA \log(2) - kA \\ &= kA (\log(4) - 1) \\ &= kA \log(4/e) > 0 \end{aligned}$$

Because $kA > 1$ and $\log(4/e) \approx 0.386$

So $U_{\text{merged}} > U_{\text{merged, Nash}}$ so the

the workers can obtain a higher utility.

$$\textcircled{2} \quad \pi = P_F \frac{F(b)}{b} - P_F \cdot \frac{1}{4}$$

B. that mean that people will send books until
Profits are 0, so will

$$P_F \frac{F(b)}{b} - P_F \frac{1}{4} = 0$$

$$P_F \left(\frac{F(b)}{b} - \frac{1}{4} \right) = 0$$

$$P_F = 0 \quad \vee \quad \frac{F(b^*)}{b^*} = \frac{1}{4} \Rightarrow b^* = 4 F(b^*)$$

So until either the Price of Fish is zero or
until the number of Books is 4 times the
amount of Fish caught. ~~But~~ However, that also
means that eventually, everyone will make 0
Profit.

$$C \quad \max_{b^*} \pi = P_F F(b^*) - \cancel{\frac{P_F b^*}{4}} - \frac{P_F b^*}{4}$$

$$0 = \frac{\partial \pi}{\partial b^*} = P_F F'(b^*) - \frac{P_F}{4}$$

(set to 0
to maximize)

$$P_F = 0 \quad \vee \quad F'(b^*) = \frac{1}{4}$$

So it will continue sending out Books until the amount
of extra fish caught is $\frac{1}{4}$ of a ton.

d. Profits will be maximized so it is ~~the~~ Pareto-efficient.
Compared to the decentralized case, everyone will
keep sending out ~~the~~ Books until profits are 0, so
~~so~~ even if they would share the profit and go
to this Pareto-efficient case, it would be a Pareto-improvement,
everyone would be better off.

$$c.f. b^1 = 4F(b) = 4 \cdot 8 \cdot b^{\frac{1}{4}}$$

$$b^{\frac{3}{4}} = 32$$

$$b^1 \approx 101,6$$

$$F'(b^*) = \frac{1}{4}$$

$$2b^{\frac{3}{4}} = \frac{1}{4}$$

$$b^{\frac{3}{4}} = \frac{1}{8}$$

$$b^* = 16$$

F. The tax will make sure that profits are 0 when $b = 16$ so:

$$\pi = P_F \cdot \frac{8b^{\frac{1}{4}}}{b} - \frac{P_F}{4} - t = 0$$

$$P_F \cdot 8b^{-\frac{3}{4}} - \frac{P_F}{4} = t$$

$$t = P_F \cdot 8 \cdot 16^{-\frac{3}{4}} - \frac{P_F}{4}$$

$$t = \frac{3}{4} P_F$$

③ a. $m_S = 0$ and as she will maximize her utility and U_S is an increasing function,
 $D = 100$ so $U_S = 10 \cdot \sqrt{100} + 0 = 100$

B. By merging utility functions:

$$\max_D U_{\text{merge}} = 10D^{\frac{1}{2}} + m_S + 10(100-D)^{\frac{1}{2}} + m_F$$

$$\frac{\partial U_{\text{merge}}}{\partial D} = 5D^{-\frac{1}{2}} - 5(100-D)^{-\frac{1}{2}} = 0$$

$$D^{-\frac{1}{2}} = (100-D)^{-\frac{1}{2}}$$

$$D = 100 - D$$

$$D = 50$$

↖ (set to 0 to maximize)

C IF P is the amount she pays, then in order for the student to be indifferent to paying or to not paying:

$$10(100 - 50)^{\frac{1}{2}} + 100 - P = 10(100 - 100)^{\frac{1}{2}} + 100$$

$$10\sqrt{50} + 100 - P = 100$$

$$P = 10\sqrt{50} = 50\sqrt{2} \approx 70,7$$

So that is the maximum.

IF she pays 50, then:

$$U_S = 10 \cdot \sqrt{50} + 50 \approx 120,7 > 100$$

So the student would be better off by taking the money.