

# Lists

Principles of Functional Programming

#### Lists

The list is a fundamental data structure in functional programming.

A list having  $x_1, ..., x_n$  as elements is written List $(x_1, ..., x_n)$ 

# Example

```
val fruit = List("apples", "oranges", "pears")
val nums = List(1, 2, 3, 4)
val diag3 = List(List(1, 0, 0), List(0, 1, 0), List(0, 0, 1))
val empty = List()
```

There are two important differences between lists and arrays.

- Lists are immutable the elements of a list cannot be changed.
- Lists are recursive, while arrays are flat.

# Lists

```
val fruit = List("apples", "oranges", "pears")
val diag3 = List(List(1, 0, 0), List(0, 1, 0), List(0, 0, 1))
```

(Drawing on Blackboard)

# The List Type

Like arrays, lists are homogeneous: the elements of a list must all have the same type.

The type of a list with elements of type T is written scala.List[T] or shorter just List[T]

#### Example

```
val fruit: List[String] = List("apples", "oranges", "pears")
val nums : List[Int] = List(1, 2, 3, 4)
val diag3: List[List[Int]] = List(List(1, 0, 0), List(0, 1, 0), List(0, 0, 1))
val empty: List[Nothing] = List()
```

# Constructors of Lists

#### All lists are constructed from:

- the empty list Nil, and
- the construction operation :: (pronounced cons):
  x :: xs gives a new list with the first element x, followed by the elements of xs.

### For example:

```
fruit = "apples" :: ("oranges" :: ("pears" :: Nil))
nums = 1 :: (2 :: (3 :: (4 :: Nil)))
empty = Nil
```

# Right Associativity

Convention: Operators ending in ":" associate to the right.

```
A :: B :: C is interpreted as A :: (B :: C).
```

We can thus omit the parentheses in the definition above.

# Example

```
val nums = 1 :: 2 :: 3 :: 4 :: Nil
```

Operators ending in ":" are also different in the they are seen as method calls of the *right-hand* operand.

So the expression above is equivalent to

```
Nil.::(4).::(3).::(2).::(1)
```

# Operations on Lists

All operations on lists can be expressed in terms of the following three operations:

```
head the first element of the list
tail the list composed of all the elements except the first.
isEmpty 'true' if the list is empty, 'false' otherwise.
```

These operations are defined as methods of objects of type list. For example:

```
fruit.head == "apples"
fruit.tail.head == "oranges"
diag3.head == List(1, 0, 0)
empty.head == throw NoSuchElementException("head of empty list")
```

# List Patterns

It is also possible to decompose lists with pattern matching.

```
Nil The Nil constant
p :: ps A pattern that matches a list with a head matching p and a tail matching ps.
List(p1, ..., pn) same as p1 :: ... :: pn :: Nil
```

#### Example

Consider the pattern x :: y :: List(xs, ys) :: zs.

What is the condition that describes most accurately the length L of the lists it matches?

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# Sorting Lists

Suppose we want to sort a list of numbers in ascending order:

- One way to sort the list List(7, 3, 9, 2) is to sort the tail List(3, 9, 2) to obtain List(2, 3, 9).
- ► The next step is to insert the head 7 in the right place to obtain the result List(2, 3, 7, 9).

This idea describes *Insertion Sort*:

```
def isort(xs: List[Int]): List[Int] = xs match
  case List() => List()
  case y :: ys => insert(y, isort(ys))
```

Complete the definition insertion sort by filling in the ???s in the definition below:

```
def insert(x: Int, xs: List[Int]): List[Int] = xs match
  case List() => ???
  case y :: ys => ???
```

What is the worst-case complexity of insertion sort relative to the length of the input list N?

0 the sort takes constant time
0 proportional to N
0 proportional to N log(N)
0 proportional to N \* N

Complete the definition insertion sort by filling in the ???s in the definition below:

```
def insert(x: Int, xs: List[Int]): List[Int] = xs match
  case List() => List(x)
  case y :: ys =>
    if x < y then x :: xs else y :: insert(x, ys)</pre>
```

What is the worst-case complexity of insertion sort relative to the length of the input list N?

```
0      the sort takes constant time
X      proportional to N
0      proportional to N * log(N)
0      proportional to N * N
```