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# CSE 373

Graphs 2: Dijkstra's Algorithm  
reading: Weiss 9.3

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<http://www.cs.washington.edu/373/>

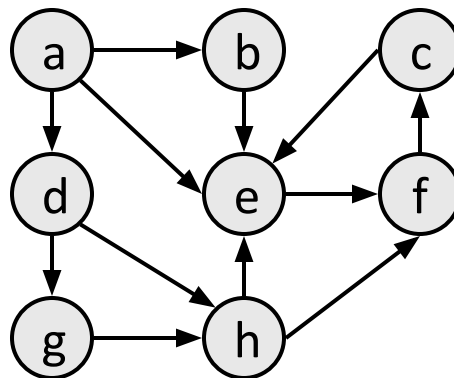
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# Recall: DFS, BFS

- **depth-first search (DFS):** Explore each possible path as far as possible before backtracking.

- Often implemented recursively.
- DFS paths from *a* to all vertices (assuming ABC edge order):

- to b: {a, b}
- to c: {a, b, e, f, c}
- to d: {a, d}
- to e: {a, b, e}
- to f: {a, b, e, f}
- to g: {a, d, g}
- to h: {a, d, g, h}



- **breadth-first search (BFS):** Take one step down all paths and then immediately backtrack.

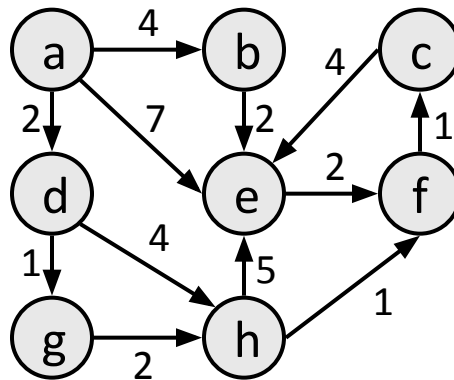
- A queue of vertices to visit.
- Always returns shortest path (one with fewest edges):

- to b: {a, b}
- to c: {a, e, f, c}
- to d: {a, d}
- to e: {a, e}
- to f: {a, e, f}
- to g: {a, d, g}
- to h: {a, d, h}

# DFS/BFS and weight

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- DFS and BFS do not consider edge weights.
- The minimum weight path is not necessarily the shortest path.
- Sometimes weight is more important:
- example: plane flight costs, network transmission (latency btwn servers)
- BFS(a,f) yields [a,e,f], but [a,d,g,h,f] has lower cost (6 vs. 9)



# Dijkstra's Algorithm

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- **Dijkstra's algorithm:** Finds the minimum-weight path between a pair of vertices in a weighted directed graph.
- Solves the "one vertex, shortest path" problem in weighted graphs.
- Made by famous computer scientist Edsger Dijkstra (look him up!)
- *basic algorithm concept:* Maintain the currently known best way to reach each vertex (cost, previous vertex), and improve it until it reaches the best solution.
- *Example:* In a graph where vertices are cities and weighted edges are roads between cities, Dijkstra's algorithm can be used to find the shortest route from one city to any other.

# Dijkstra pseudocode

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```
function dijkstra( $v_1, v_2$ ):  
  for each vertex  $v$ :                                // Initialize vertex info  
     $v$ 's cost := infinity.  
     $v$ 's previous := none.  
   $v_1$ 's cost := 0.  
   $pqueue$  := {all vertices, ordered by distance}.  
  
  while  $pqueue$  is not empty:  
     $v$  := remove vertex from  $pqueue$  with minimum cost.  
    mark  $v$  as visited.  
    for each unvisited neighbor  $n$  of  $v$ :  
       $cost$  :=  $v$ 's cost + weight of edge ( $v, n$ ).  
  
      if  $cost < n$ 's cost:  
         $n$ 's cost :=  $cost$ .  
         $n$ 's previous :=  $v$ .  
  
  reconstruct path from  $v_2$  back to  $v_1$ , following previous pointers.
```

# Dijkstra example

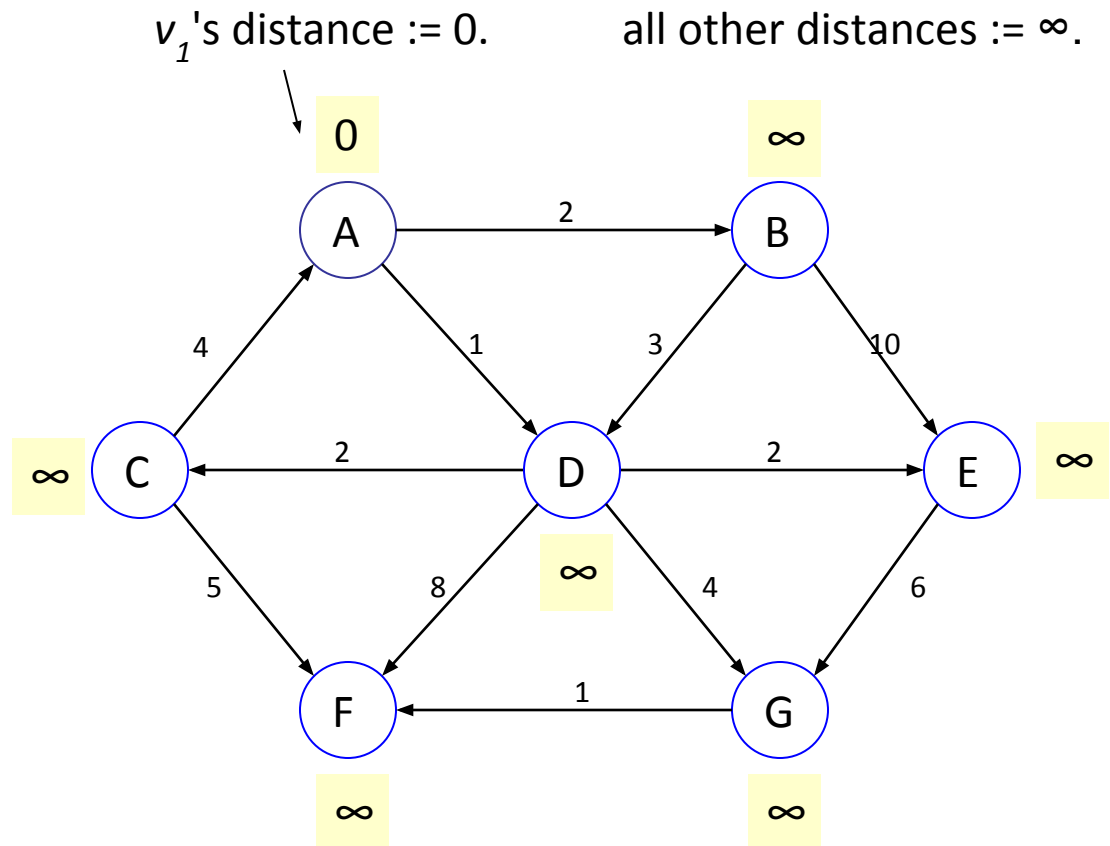
- `dijkstra(A, F);`

```
function dijkstra( $v_1, v_2$ ):
  for each vertex  $v$ : // Initialize vertex info
     $v$ 's cost := infinity.
     $v$ 's previous := none.
   $v_1$ 's cost := 0.
   $pqueue := \{\text{all vertices, by distance}\}.$ 
```

```
while  $pqueue$  is not empty:
   $v := pqueue.removeMin()$ .
  mark  $v$  as visited.
  for each unvisited neighbor  $n$  of  $v$ :
     $cost := v$ 's cost +  $edge(v, n)$ 's weight.
```

```
  if  $cost < n$ 's cost:
     $n$ 's cost :=  $cost$ .
     $n$ 's previous :=  $v$ .
```

```
reconstruct path from  $v_2$  back to  $v_1$ ,
following previous pointers.
```



$pqueue = [A:0, B:\infty, C:\infty, D:\infty, E:\infty, F:\infty, G:\infty]$

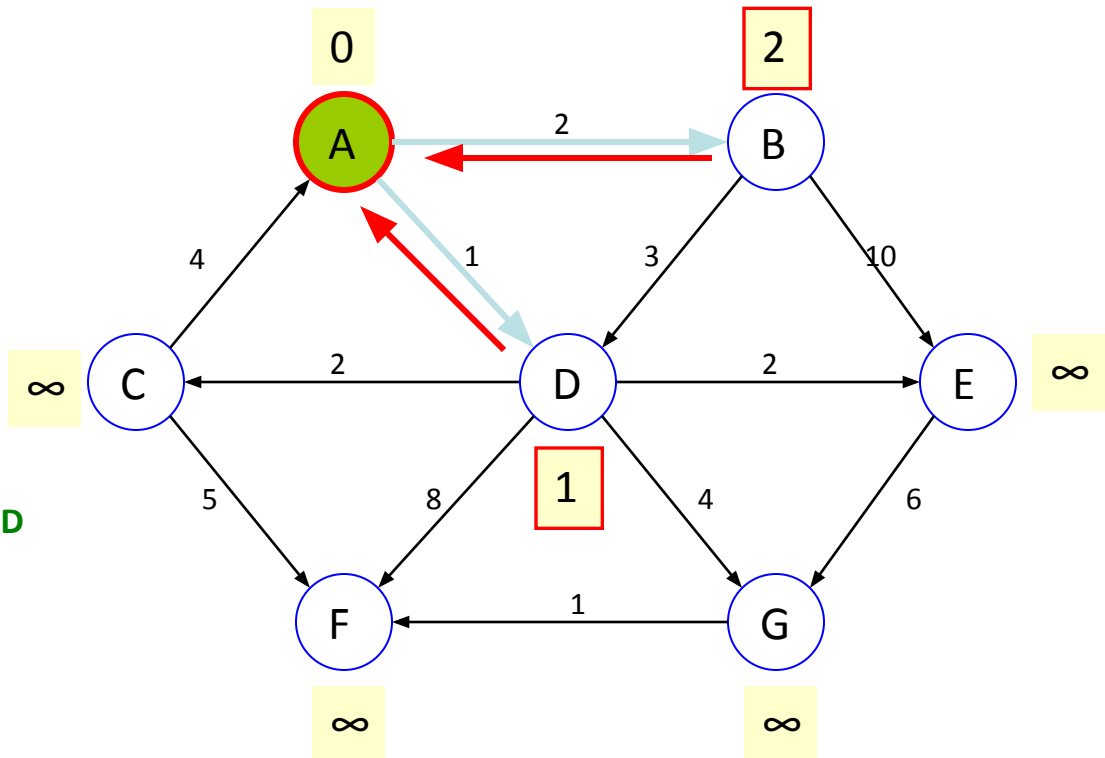
# Dijkstra example

- `dijkstra(A, F);`

```
function dijkstra( $v_1, v_2$ ):
  for each vertex  $v$ : // Initialize vertex info
     $v$ 's cost := infinity.
     $v$ 's previous := none.
   $v_1$ 's cost := 0.
   $pqueue := \{\text{all vertices, by distance}\}$ .

  while  $pqueue$  is not empty:
     $v := pqueue.removeMin()$ . // A
    mark  $v$  as visited.
    for each unvisited neighbor  $n$  of  $v$ : // B, D
       $cost := v$ 's cost + edge( $v, n$ )'s weight.

      if  $cost < n$ 's cost: // B's cost = 0 + 2
         $n$ 's cost :=  $cost$ . // D's cost = 0 + 1
         $n$ 's previous :=  $v$ .
```



reconstruct path from  $v_2$  back to  $v_1$ ,  
following previous pointers.

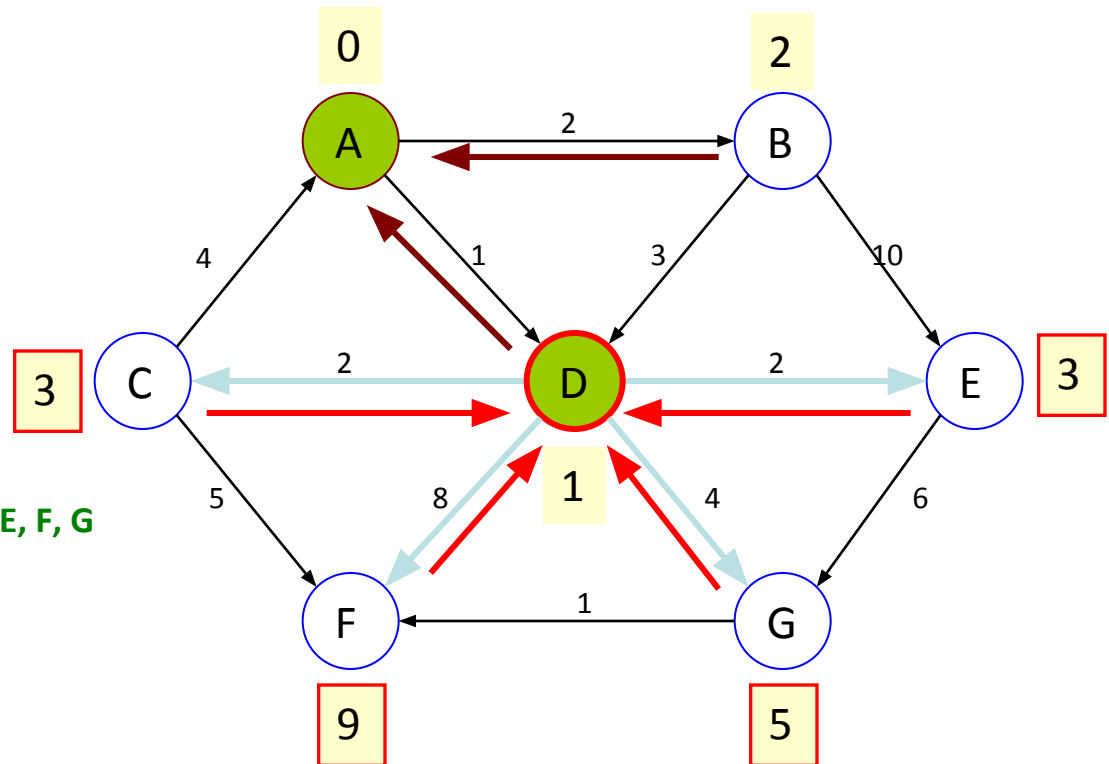
$pqueue = [D:1, B:2, C:\infty, E:\infty, F:\infty, G:\infty]$

# Dijkstra example

- dijkstra(A, F);

```
function dijkstra( $v_1, v_2$ ):
  for each vertex  $v$ : // Initialize vertex info
     $v$ 's cost := infinity.
     $v$ 's previous := none.
   $v_1$ 's cost := 0.
   $pqueue := \{\text{all vertices, by distance}\}$ .

  while  $pqueue$  is not empty:
     $v := pqueue.removeMin()$ . // D
    mark  $v$  as visited.
    for each unvisited neighbor  $n$  of  $v$ : // C, E, F, G
       $cost := v$ 's cost + edge( $v, n$ )'s weight.
      // C's cost = 1 + 2
      if  $cost < n$ 's cost: // E's cost = 1 + 2
         $n$ 's cost :=  $cost$ . // F's cost = 1 + 8
         $n$ 's previous :=  $v$ . // G's cost = 1 + 4
```



reconstruct path from  $v_2$  back to  $v_1$ ,  
following previous pointers.

$pqueue = [B:2, C:3, E:3, G:5, F:9]$



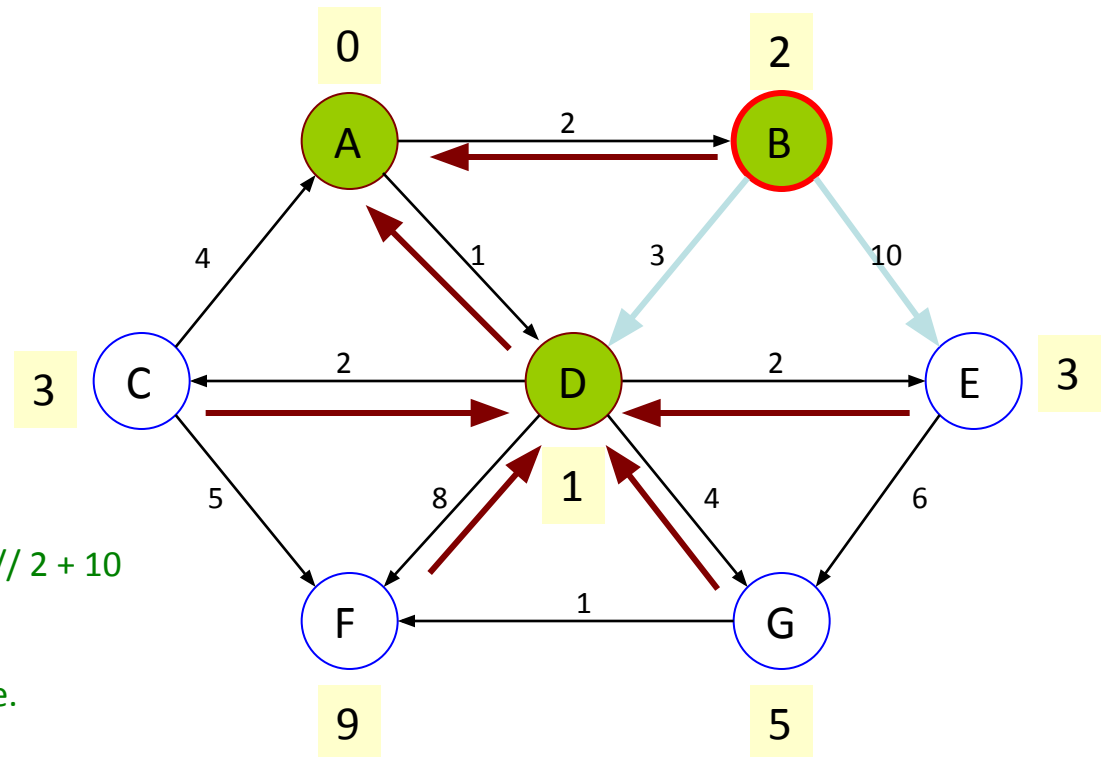
# Dijkstra example

- `dijkstra(A, F);`

```
function dijkstra( $v_1, v_2$ ):
  for each vertex  $v$ : // Initialize vertex info
     $v$ 's cost := infinity.
     $v$ 's previous := none.
   $v_1$ 's cost := 0.
   $pqueue := \{\text{all vertices, by distance}\}$ .

  while  $pqueue$  is not empty:
     $v := pqueue.removeMin()$ . // B
    mark  $v$  as visited.
    for each unvisited neighbor  $n$  of  $v$ : // E
       $cost := v$ 's cost + edge( $v, n$ )'s weight. // 2 + 10

      if  $cost < n$ 's cost: // 12 > 3; false
         $n$ 's cost :=  $cost$ . // no costs change.
         $n$ 's previous :=  $v$ .
```



reconstruct path from  $v_2$  back to  $v_1$ ,  
following previous pointers.

$pqueue = [C:3, E:3, G:5, F:9]$

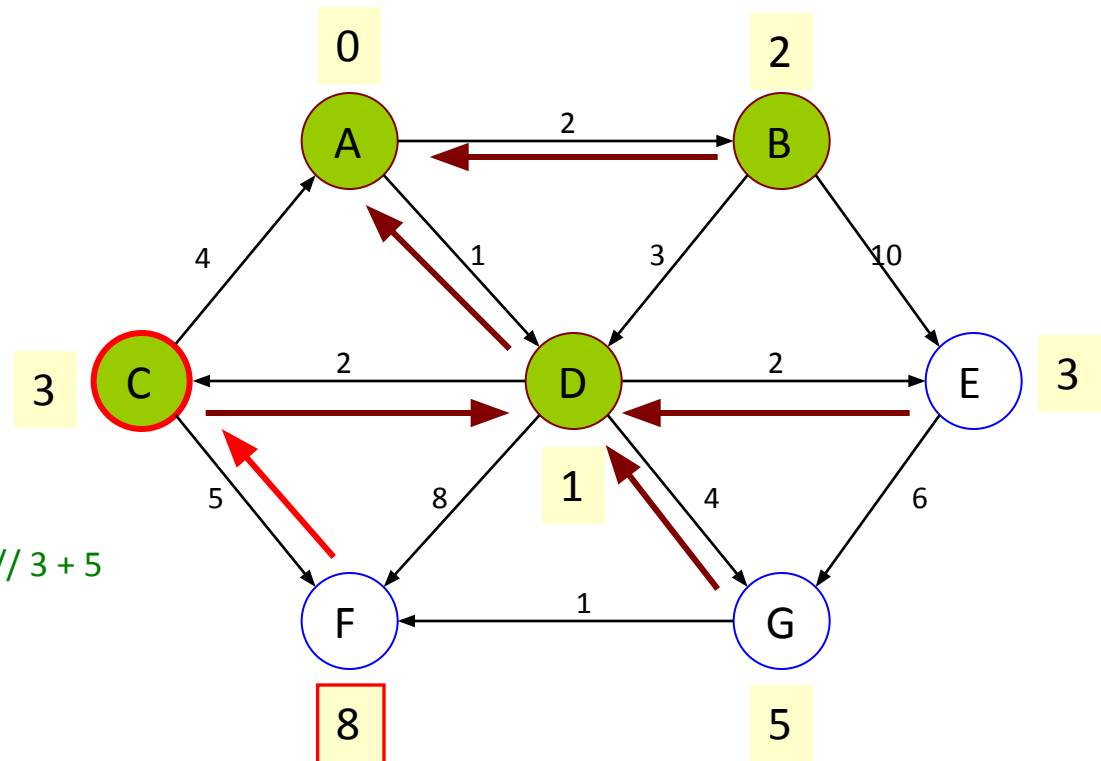
# Dijkstra example

- `dijkstra(A, F);`

```
function dijkstra( $v_1, v_2$ ):
  for each vertex  $v$ : // Initialize vertex info
     $v$ 's cost := infinity.
     $v$ 's previous := none.
   $v_1$ 's cost := 0.
   $pqueue := \{\text{all vertices, by distance}\}$ .

  while  $pqueue$  is not empty:
     $v := pqueue.removeMin()$ . // C
    mark  $v$  as visited.
    for each unvisited neighbor  $n$  of  $v$ : // F
       $cost := v$ 's cost + edge( $v, n$ )'s weight. //  $3 + 5$ 

      if  $cost < n$ 's cost: //  $8 < 9$ 
         $n$ 's cost :=  $cost$ . //  $F$ 's cost = 8
         $n$ 's previous :=  $v$ .
```



reconstruct path from  $v_2$  back to  $v_1$ ,  
following previous pointers.

$pqueue = [E:3, G:5, \mathbf{F:8}]$

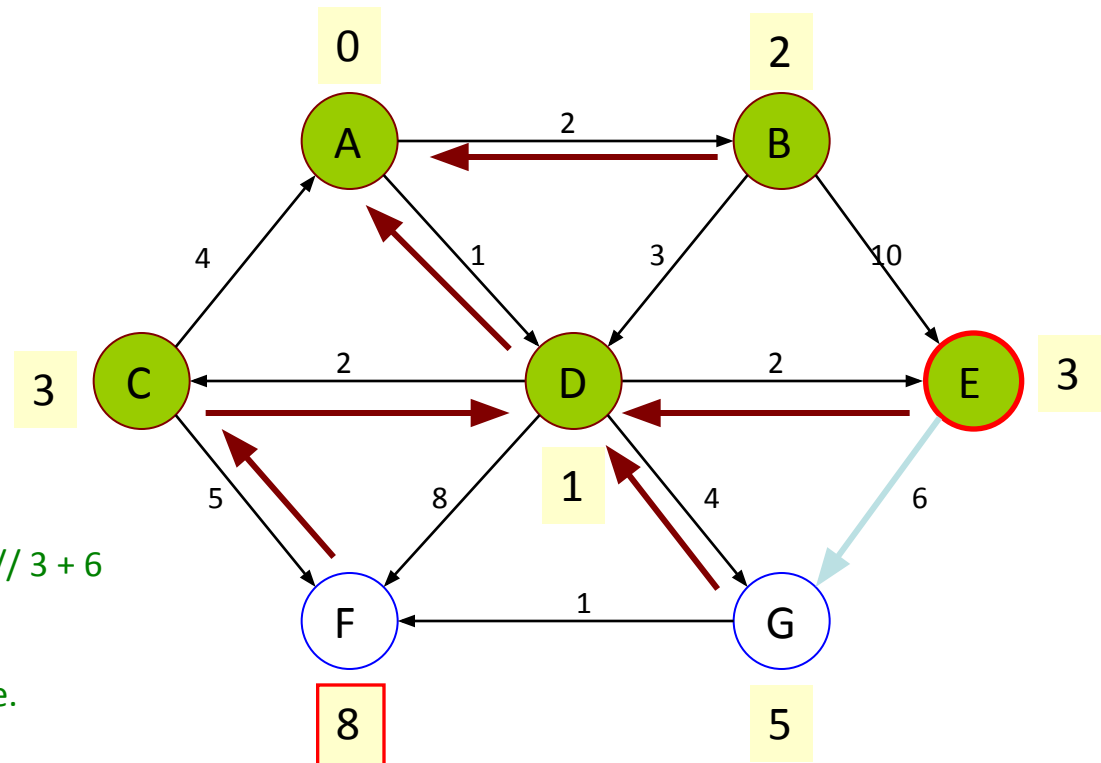
# Dijkstra example

- `dijkstra(A, F);`

```
function dijkstra( $v_1, v_2$ ):
  for each vertex  $v$ : // Initialize vertex info
     $v$ 's cost := infinity.
     $v$ 's previous := none.
   $v_1$ 's cost := 0.
   $pqueue := \{\text{all vertices, by distance}\}$ .

  while  $pqueue$  is not empty:
     $v := pqueue.removeMin()$ . // E
    mark  $v$  as visited.
    for each unvisited neighbor  $n$  of  $v$ : // G
       $cost := v$ 's cost + edge( $v, n$ )'s weight. // 3 + 6

      if  $cost < n$ 's cost: // 9 > 5; false
         $n$ 's cost :=  $cost$ . // no costs change.
         $n$ 's previous :=  $v$ .
```



reconstruct path from  $v_2$  back to  $v_1$ ,  
following previous pointers.

$pqueue = [G:5, F:8]$

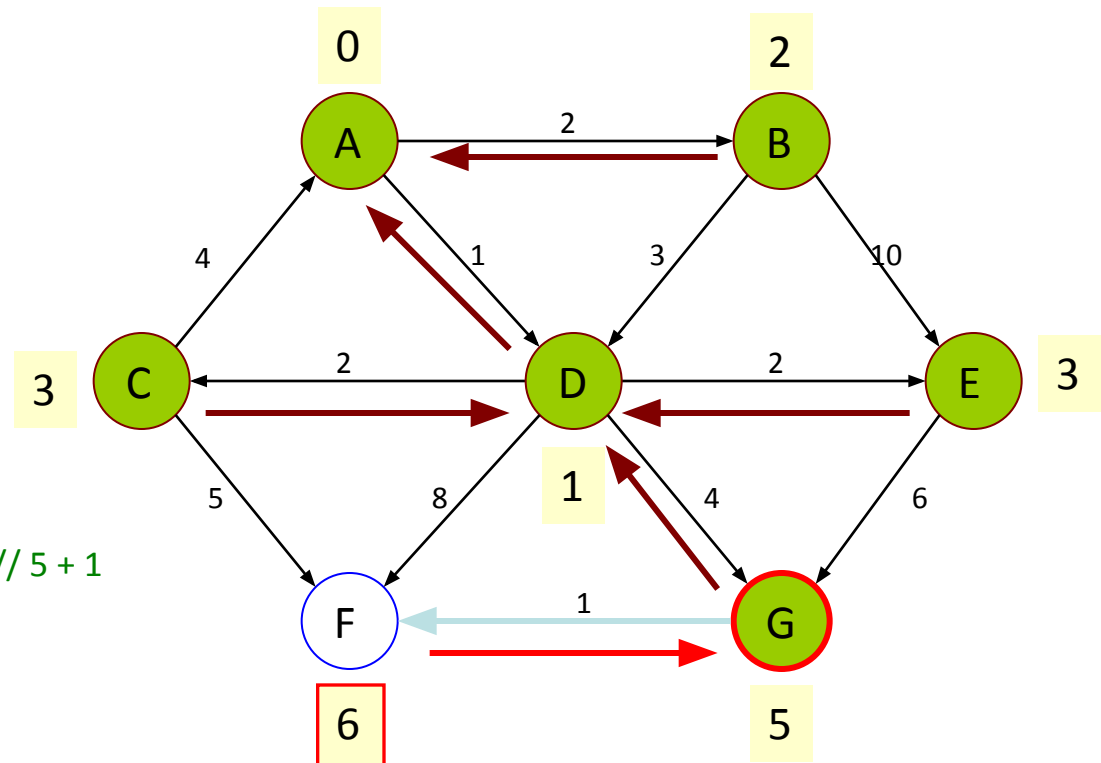
# Dijkstra example

- `dijkstra(A, F);`

```
function dijkstra( $v_1, v_2$ ):
  for each vertex  $v$ : // Initialize vertex info
     $v$ 's cost := infinity.
     $v$ 's previous := none.
   $v_1$ 's cost := 0.
   $pqueue := \{\text{all vertices, by distance}\}$ .

  while  $pqueue$  is not empty:
     $v := pqueue.removeMin()$ . // G
    mark  $v$  as visited.
    for each unvisited neighbor  $n$  of  $v$ : // F
       $cost := v$ 's cost + edge( $v, n$ )'s weight. // 5 + 1

      if  $cost < n$ 's cost: // 6 < 8
         $n$ 's cost :=  $cost$ . // F's cost = 6.
         $n$ 's previous :=  $v$ .
```



`pqueue = [F:6]`

reconstruct path from  $v_2$  back to  $v_1$ ,  
following previous pointers.

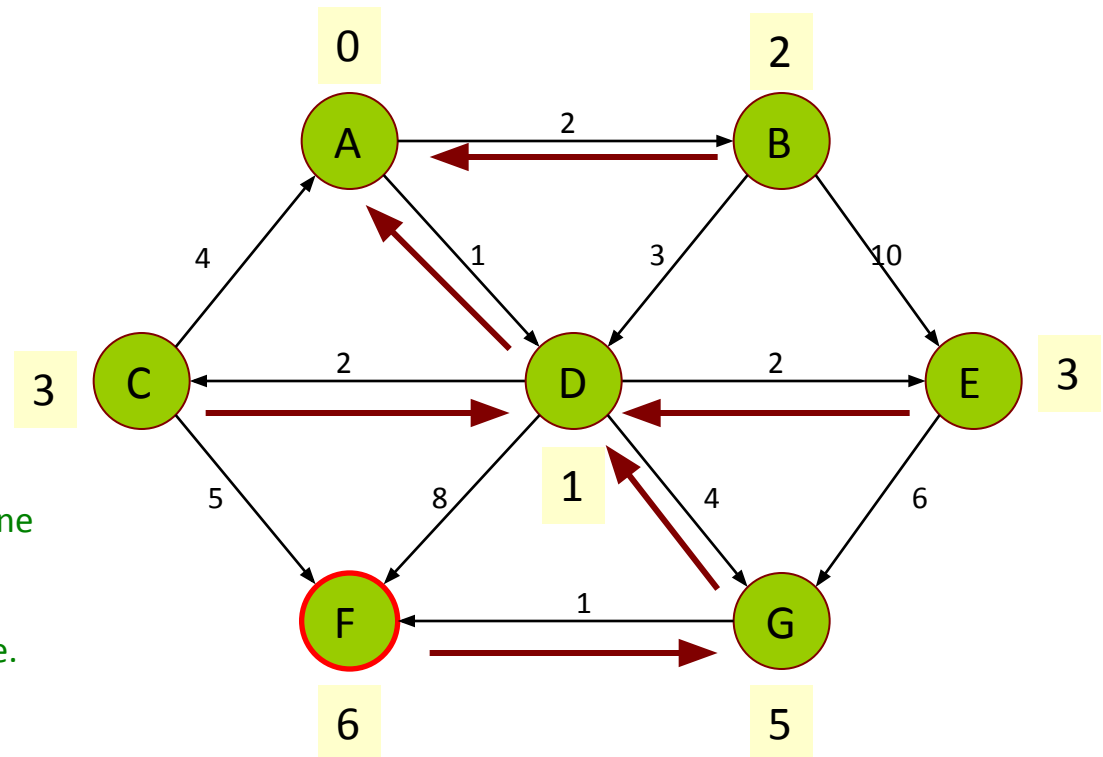
# Dijkstra example

- `dijkstra(A, F);`

```
function dijkstra( $v_1, v_2$ ):
  for each vertex  $v$ : // Initialize vertex info
     $v$ 's cost := infinity.
     $v$ 's previous := none.
   $v_1$ 's cost := 0.
   $pqueue := \{\text{all vertices, by distance}\}$ .

  while  $pqueue$  is not empty:
     $v := pqueue.removeMin()$ . // F
    mark  $v$  as visited.
    for each unvisited neighbor  $n$  of  $v$ : // none
       $cost := v$ 's cost + edge( $v, n$ )'s weight.

      if  $cost < n$ 's cost: // no costs change.
         $n$ 's cost :=  $cost$ .
         $n$ 's previous :=  $v$ .
```



`pqueue = []`

reconstruct path from  $v_2$  back to  $v_1$ ,  
following previous pointers.

# Dijkstra example

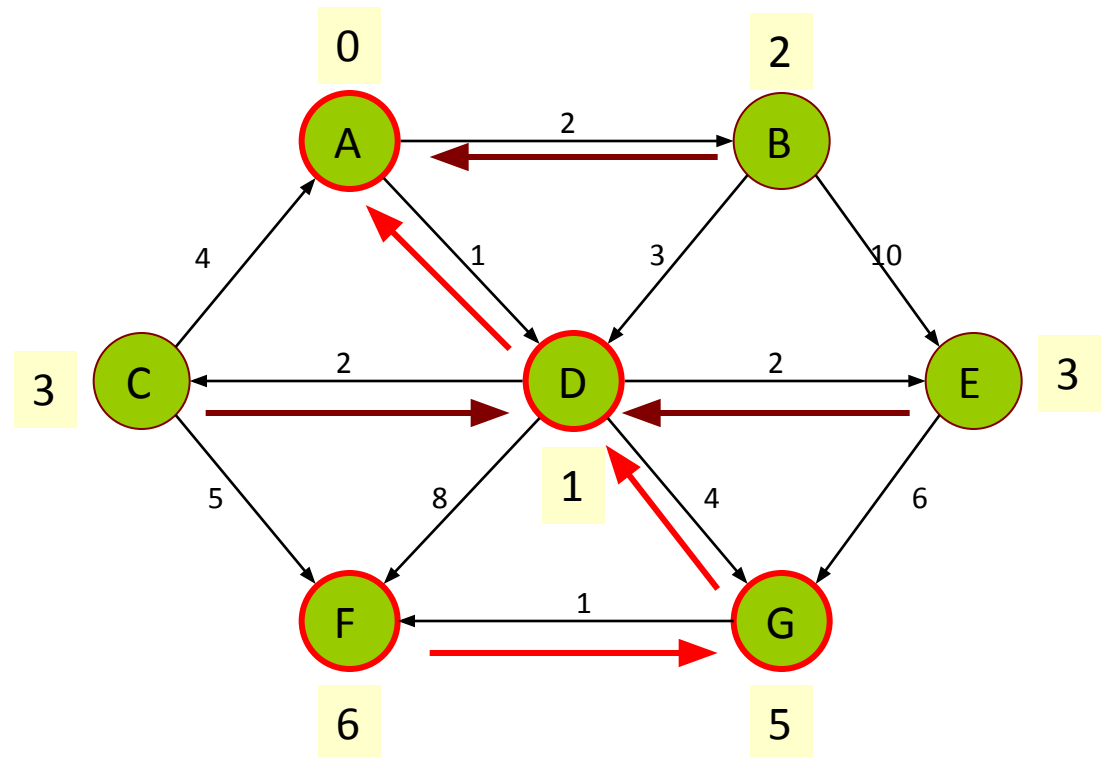
- dijkstra(A, F);

```
function dijkstra( $v_1, v_2$ ):  
  for each vertex  $v$ : // Initialize vertex info  
     $v$ 's cost := infinity.  
     $v$ 's previous := none.  
   $v_1$ 's cost := 0.  
   $pqueue := \{\text{all vertices, by distance}\}.$ 
```

```
while  $pqueue$  is not empty:  
   $v := pqueue.removeMin().$   
  mark  $v$  as visited.  
  for each unvisited neighbor  $n$  of  $v$ :  
     $cost := v$ 's cost +  $edge(v, n)$ 's weight.
```

```
  if  $cost < n$ 's cost:  
     $n$ 's cost :=  $cost.$   
     $n$ 's previous :=  $v.$ 
```

```
reconstruct path from  $v_2$  back to  $v_1$ , // path = [A, D, G, F]  
following previous pointers.
```



# Algorithm properties

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- Dijkstra's algorithm is a *greedy algorithm*:
  - Make choices that currently seem the best.
  - Locally optimal does not always mean globally optimal.
- It is correct because it maintains the following two properties:
  - 1) for every marked vertex, the current recorded cost is the lowest cost to that vertex from the source vertex.
  - 2) for every unmarked vertex  $v$ , its recorded distance is shortest path distance to  $v$  from source vertex, considering only currently known vertices and  $v$ .

# Dijkstra's runtime

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- For sparse graphs, (i.e. graphs with much less than  $|V|^2$  edges) Dijkstra's is implemented most efficiently with a priority queue.
  - initialization:  $O(|V|)$
  - while loop:  $O(|V|)$  times
    - remove min-cost vertex from  $pq$ :  $O(\log |V|)$
    - potentially perform  $|E|$  updates on cost/previous
    - update costs in  $pq$ :  $O(\log |V|)$
  - reconstruct path:  $O(|E|)$
- Total runtime:  $O(|V| \log |V| + |E| \log |V|)$ 
  - =  **$O(|E| \log |V|)$** , because  $|V| = O(|E|)$  if graph is connected
  - if a list is used instead of a  $pq$ :  $O(|V|^2 + |E|) = O(|V|^2)$



# Dijkstra exercise

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- Use Dijkstra's algorithm to determine the lowest cost path from vertex A to all of the other vertices in the graph.
  - Keep track of previous vertices so that you can reconstruct the path.

