CS2323 HW3

Roll: CO22BTECH11006

Name: Om Dave

Q1 (Decimal to floating point representation)

(a) -13.25 (Single Precision)

Sign Bit: 1

Binary Conversion:

Integer (13): 1101Fraction (0.25): .01Combined: 1101.01

Normalization: 1.10101 * 2^3

Exponent Calculation:

• Actual Exponent: 3

• Bias (Single): 127

• Biased Exponent: 3 + 127 = 130

• Binary (8-bit): 10000010

Mantissa Calculation:

• From Normalization: 10101

• Padded (23-bit): 10101000000000000000000

Final Binary Representation:

1 10000010 10101000000000000000000

Hexadecimal Conversion:

• Final Hex: 0xC1540000

(b) 0.1 (Single Precision)

Sign Bit: 0

Binary Conversion:

• 0.0001100110011... (repeating 0011)

Normalization: 1.100110011.* 2^-4

Exponent Calculation:

- Actual Exponent: -4
- Bias (Single): 127
- Biased Exponent: -4 + 127 = 123
- Binary (8-bit): 01111011

Mantissa Calculation:

- Repeating Pattern: 100110011...
- Truncated/Rounded (23-bit): 1001100110011001101

Final Binary Representation:

0 01111011 1001100110011001101

Hexadecimal Conversion:

- Grouped: 0011 1101 1100 1100 1100 1100 1101
- Final Hex: 0x3DCCCCCD

(c) 156.75 (Double Precision)

Sign Bit: 0

Binary Conversion:

- Integer (156): 10011100
- Fraction (0.75): .11
- Combined: 10011100.11

Normalization: 1.001110011 * 2^7

Exponent Calculation:

- Actual Exponent: 7
- Bias (Double): 1023
- Biased Exponent: 7 + 1023 = 1030
- Binary (11-bit): 10000000110

Mantissa Calculation:

- From Normalization: 001110011
- Padded (52-bit): 001110011000...000

Final Binary Representation:

Hexadecimal Conversion:

- Grouped: 0100 0000 0110 0011 1001 1000 0000 .0000
- Final Hex: 0x4063980000000000

(d) -0.0078125 (Double Precision)

Sign Bit: 1

Binary Conversion:

• 0.0000001

Normalization: 1.0 * 2^-7

Exponent Calculation:

Actual Exponent: -7Bias (Double): 1023

• Biased Exponent: -7 + 1023 = 1016

• Binary (11-bit): 01111111000

Mantissa Calculation:

• From Normalization: 0

• Padded (52-bit): All zeros

Final Binary Representation:

Hexadecimal Conversion:

• Grouped: 1011 1111 1000 0000 0000 .0000

• Final Hex: 0xBF8000000000000

Q2 (Floating point to Decimal representation)

(a) 0xC1200000 (Single Precision)

Hex to Binary:

Deconstruct Bits (1-8-23):

• **Sign:** 1 (Negative)

• Exponent: 10000010 (Decimal: 130) -> Actual Exponent: 130 - 127 = 3

Decimal Conversion:

• Normalized Binary (1.Mantissa): 1.01

• Non-normalized Binary (Shift by Exponent 3): $1.01 * 2^3 = 1010.0$

• Binary to Decimal: 8 + 0 + 2 + 0 = 10.0

• Final Value (With Sign): -10.0

(b) 0x3F800000 (Single Precision)

Hex to Binary:

Deconstruct Bits (1-8-23):

```
• Sign: 0 (Positive)
```

• Exponent: 011111111 (Decimal: 127) -> Actual Exponent: 127 - 127 = 0

Decimal Conversion:

- Normalized Binary (1.Mantissa): 1.0
- Non-normalized Binary (Shift by Exponent 0): $1.0 \times 2^0 = 1.0$
- Binary to Decimal: 1.0
- Final Value (With Sign): 1.0
- (c) 0xBFF0000000000000 (Double Precision)

Hex to Binary:

Deconstruct Bits (1-11-52):

- **Sign:** 1 (Negative)
- Exponent: 01111111111 (Decimal: 1023) -> Actual Exponent: 1023 1023 = 0
- Mantissa: 000 . . . (All zeros)

Decimal Conversion:

- Normalized Binary (1.Mantissa): 1.0
- Non-normalized Binary (Shift by Exponent 0): $1.0 * 2^0 = 1.0$
- Binary to Decimal: 1.0
- Final Value (With Sign): -1.0
- (d) 0x4024000000000000 (Double Precision)

Hex to Binary:

Deconstruct Bits (1-11-52):

- **Sign:** 0 (Positive)
- Exponent: 10000000010 (Decimal: 1026) -> Actual Exponent: 1026 1023 = 3
- Mantissa: 010000...

Decimal Conversion:

- Normalized Binary (1.Mantissa): 1.01
- Non-normalized Binary (Shift by Exponent 3): 1.01 * 2³ = 1010.0

- Binary to Decimal: 10.0
- Final Value (With Sign): 10.0

Q3

```
Given (hex): A = 0 \times 41480000 (single) B = 0 \times C0700000 (single)
```

Decode A (Single Precision)

- Sign Bit: 0
- Exponent: $10000010_2 = 130 \Rightarrow Actual e = 130 127 = 3$
- Mantissa (with hidden 1): 1.1001
- Value: $1.1001_2 \times 2^3 = 12.5$

Decode B (Single Precision)

- Sign Bit: 1
- Exponent: $1000000002 = 128 \Rightarrow Actual e = 1$
- Mantissa (with hidden 1): 1.111
- Value: $-(1.11112 \times 2^{1}) = -3.75$

Result (Normalized)

```
Value: 12.5 + (-3.75) = 8.75 = 1.00011_2 \times 2^3
```

Encode Result as Double Precision

- Sign Bit: 0
- Exponent Calculation: Actual e = 3, Bias (double) $1023 \Rightarrow E = 1026 \Rightarrow 10000000010$
- Mantissa (52-bit): from $1.00011 \Rightarrow$ fraction 00011 + pad to 52000110 ...
- Final Binary Representation (1–11–52): 0 10000000010 0001100...
- Final Hex: 0x402180000000000

Q4

Decode A (Double Precision)

- Hex → Binary (64): 0100 0000 0011 1001 0000 .0000
- Deconstruct (1–11–52): 0 | 10000000011 | 1001000...(zeros)...

- Sign Bit: 0
- Exponent: $10000000011_2 = 1027 \Rightarrow Actual e = 1027 1023 = 4$
- Mantissa (with hidden 1): 1.1001
- Value: $1.1001_2 \times 2^4 = 25.0$

Decode B (Double Precision)

- Hex → Binary (64): 1100 0000 0000 1000 0000 .0000
- Deconstruct (1–11–52): 1 | 10000000000 | 1000000...(zeros)...
- Sign Bit: 1
- Exponent: $1000000000002 = 1024 \Rightarrow Actual e = 1024 1023 = 1$
- Mantissa (with hidden 1): 1.1
- Value: $-(1.12 \times 2^1) = -3.0$

Multiply (Real)

• Value: 25.0 * -3 = -75.0

Normalize Product

• Value: $-75.0 = -(1001011_2) = -(1.001011_2 \times 2^6)$

Encode Result as **Single Precision**

- Sign Bit: 1
- Exponent Calculation: Actual e = 6, Biased exponent (single) ⇒ E = 127 + 6 = 133 ⇒ 10000101
- Mantissa (23-bit): from 1.001011 ⇒ fraction 001011 + pad to 23

Final Hex: 0xC2960000

Q5

Take the integer:

$$$$$
 N = $2^{24} + 1 = 16{,}777{,}217 $$$

- A 32-bit signed integer can represent all values from \$-2^{31}\$ to \$2^{31}-1\$. So \$16{,}777{,}217\$ can be represented.
- A 32-bit float (IEEE 754 single precision) has:
 - 23 fraction bits + 1 hidden bit = 24 bits of precision.
 - This means every integer up to \$2^{24}\$ can be exactly represented.
 - Once an integer requires more than 24 bits to write in binary, precision is lost.

Floating-point representation of \$2^{24}+1\$

• Binary form:

```
$ 2^{24}+1 = 1.0000\ldots000001 \times 2^{24} $$ (23 zeros then a trailing 1).
```

• Mantissa: only 23 fraction bits are stored. The trailing 1 at the 24th place does not fit.

Thus the stored value becomes:

```
$$ 1.0000\dots0000 \times 2^{24} = 2^{24} = 16{,}777{,}216 $$
```

 $16{,}777{,}217$ exists in 32-bit signed int but not in float. In float it is rounded off to $16{,}777{,}216$, losing the +1.

Q6

Let float32 numbers:

$$$$$
 a=-2^{24},\quad b=2^{24},\quad c=1. \$\$

• **Left-associate** \$(a+b)+c\$:

$$$$ ((-2^{24})+2^{24})+1 = 0+1 = 1\quad (\text{exact}). $$$$

• **Right-associate** \$a+(b+c)\$:

```
$$ b+c = 2^{24}+1 ;\xrightarrow{\text{round}}; 2^{24}\quad(\text{As seen previously}), $$ so $$ a+(b+c)=(-2^{24})+2^{24}=0. $$
```

Hence,

$$$$ (a+b)+c = 1 ; \neq (b+c) = 0. $$$$

Hence, addition in floating point numbers is not always associative.

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So \$16{,}777{,}217\$ exists in 32-bit signed int but not in float. In float it is rounded off to \$16{,}777{,}216\$, losing the \$+1\$.

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Hence,

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