## 1)

## **Load Balancing**

Problem Input: m identical machines, n jobs with ith job having processing time ti

Goal: Schedule jobs to computers such that

- ☐ Jobs run contiguously on a machine
- ☐ A machine processes only one job a time
- ☐ Makespan or maximum load on any machine is minimized

Definition Let A(i) be the set of jobs assigned to machine i.

The load on i is Ti = Summation of time of every job executed on i.

The makespan of A is T = maxi Ti

Load Balancing is NP-Complete.

## 1. Greedy Algorithm

- ☐ Consider the jobs in some fixed order
- ☐ Assign job j to the machine with lowest load so far

Consider 6 jobs whose processing times is given as follows

Jobs 1 2 3 4 5 6

ti 2 3 4 6 2 2

2 
$$job 2(3) - - - job 5(2) - - : 5$$

The loads are: T1 = 8, T2 = 5, and T3 = 6.

So makespan of schedule is 8

Is the greedy algorithm optimal?

Modified Greedy Sort the jobs in descending order of processing time, and process jobs using greedy algorithm

Consider 6 jobs whose processing times is given as follows

Jobs 1 2 3 4 5 6

ti 2 3 4 6 2 2

Jobs	4	3	2	1	5	6
time	6	4	3	2	2	2

Machine 1 J4 - - - - : 6 Machine 2 J3 - - - J2 - : 6 Machine 3 J2 - - - J1 - - J3 - : 7

Makespan: 7

Machine 1 J1-----J4 : 8 Machine 2 J2-----J5 : 10

1 2 3 4 5 6 7 7 7 6 6 5 5 5 Machine 1 J1----J5----J7 :17

Machine 2 J2----J6 :12 Machine 3 J3----J4 :12

Is the modified greedy algorithm always optimal?

Greedy algorithm:

Jobs	1	2	3	4	5	6	7	8	9
Time	5	8	5	7	7	1	5	8	1

Makespan: 20

Modified greedy algorithm:

Jobs	2	8	4	5	1	3	7	6	9
Time	8	8	7	7	5	5	5	1	1

Machine 1 J2 - - - - J1 - - - J7 - - - : 18 Machine 2 J8 - - - - J3 - - - J6 - : 14 Machine 3 J4 - - - - J5 - - - - J9 - : 15 Makespan: 18

Machine 1	5	5	7		: 17
Machine 2	8	7	1	1	: 17
Machine 3	8	5			: 13

Greedy algorithm:

Jobs Time	1 6	2 9	3 6	4 9	5 9	6 2	7 6	8	9
6 9 6	9 2 9	9 6 2		24 17 17					
9	9	9	9	6	6	6	2	2	
9 9 9	9 6 6	2 6 2	20 21 17						
9	9	9	9	6	6	6	2	2	
9 6 9	9 6 9	2 6 2	20 18 20						

Makespan: 20