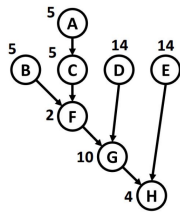


### CS3210 Tutorial 3

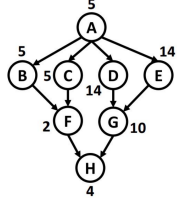
AY 25/26 Sem 2 — github/omgeta

- Q1. Instruction-Level Parallelism: pipeline, branch prediction, superscalar  
Thread-Level Parallelism: Simultaneous Multi-Threading (SMT)  
Processor-Level Parallelism: multi-core running in parallel, UPI shared memory, Omni-Path distributed memory  
IO Parallelism: Multiple DDR4 channels and PCIe lanes
- Q2. (a) SISD: single core with single instruction stream on each memory address  
(b) MIMD: parallel instruction streams on different cores can act on different data  
(c) SIMD: vector instructions act over multiple memory locations  
(d) SISD: single core with single instruction stream, pipelined instructions still act on a single memory address at a specific time step  
(e) MISD: same paper (data) operated on by different students (instructions)
- Q3. (a) False; Shared-Memory can be UMA or NUMA  
(b) False; data locality affects latency  
(c) True
- Q4. (a) False/True; binary semaphores emulate mutual exclusion but code depending on mutex ownership and semantics may be affected  
(b) False; it depends on overheads, synchronisation and workload

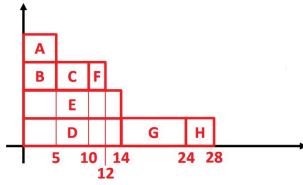
Q5. (a) Fragment 1:



Fragment 2:

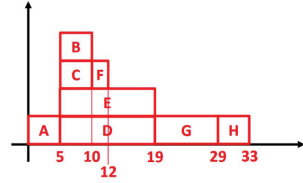


(b) Fragment 1:



Max Concurrency = 4, Average Concurrency =  $\frac{59}{28} = 2.1$

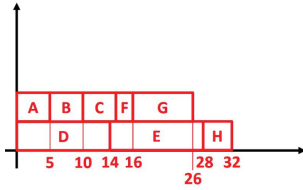
Fragment 2:



Max Concurrency = 4, Average Concurrency =  $\frac{59}{33} = 1.78$

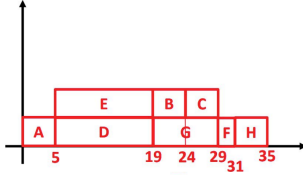
(c) Fragment 1 Speedup =  $\frac{59}{28} = 2.1$   
 Fragment 2 Speedup =  $\frac{59}{33} = 1.78$

(d) Fragment 1:



Speedup =  $\frac{59}{32} = 1.84 < 2.1$

Fragment 2:



Speedup =  $\frac{59}{35} = 1.69 < 1.78$

- Q6. (a) Data Parallelism; SIMD; each independent output element  $c_i$  is an independent dot product of the vector  $b$  with row  $a_i$
- (b) Parbegin-Parend
- Q7. (a) Task Parallelism: each stage is a task  
 Pipelining: best for cases where each stage takes a similar time
- (b) Task Parallelism: each stage is a task  
 Producer-Consumer: producer reads from socket and processes, consumer writes back to socket and there are equal number of producers as consumers
- (c) Task Parallelism: each stage is a task  
 Producer-Consumer: same as (b) but more producers than consumers
- (d) Task Parallelism: each stage is a task  
 Task Pool: pool of tasks are ready for jobs and assigned as necessary
- (e) Same; but need explicit communication
- Q8. Average CPI (Translation 1) =  $\frac{10}{5} = 2$   
 Average CPI (Translation 1) =  $\frac{9}{6} = 1.5$
- Q9. Execution Time ( $A_1$ ) =  $\frac{(5+1.2+1.3) \times 10^9}{2 \times 10^9} = 5s$   
 Execution Time ( $A_1$ ) =  $\frac{(10+1.2+1.3) \times 10^9}{2 \times 10^9} = 7.5s$

Conclusion: MIPS is not an accurate measurement of performance

- Q10. (a) Sequential Time:  $\frac{(100+100^2) \times 2}{10^9} = 20200ns$   
 Parallel Time ( $p = 10$ ):  $\frac{(100 + \frac{100^2}{10}) \times 2}{10^9} = 2200ns$  ( $9.18 \times$  speedup)  
 Parallel Time ( $p = 100$ ):  $\frac{(100 + \frac{100^2}{100}) \times 2}{10^9} = 400ns$  ( $50.6 \times$  speedup)  
 Parallel Time ( $p = \infty$ ):  $\frac{(100+0) \times 2}{10^9} = 200ns$  ( $101 \times$  speedup)
- (b) From (a),  $T_1 = 20200ns$   
 For  $p = 10$ ,  
 $\frac{(N + \frac{N_2}{10}) \times 2}{10^9} = 20200ns \implies N_{p=10} = 312 \implies S_{10}(312) = 9.669 \times$   
 For  $p = 100$ ,  
 $\frac{(N + \frac{N_2}{100}) \times 2}{10^9} = 20200ns \implies N_{p=100} = 956 \implies S_{100}(312) = 90.58 \times$