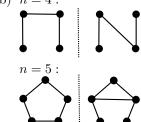
CS1231S Tutorial 11

AY 24/25 Sem 1 — github/omgeta

Q1. (a)



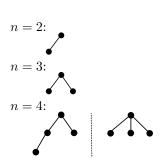
(b) n = 4:



For $n = 3, 6, K_n$ has odd edges and cannot be divided into two halves

Q2. $4 \times 3 = 12$

Q3. (a) n = 1:



(b) n = 1 has 1, n = 2 has 1, n = 3 has $\frac{3!}{2} = 3, n = 4$ has $\frac{4!}{2} + 4 = 12 + 4 = 16$

- Q4. (a) 1. Suppose G = (V, E) is a connected, simple, undirected graph
 - 2. There is spanning tree T = (V, E'), where $E' \subseteq E$ (Proposition 10.7.1)
 - 3. Since T is a tree, |E'| = |V| 1
 - 4. Therefore, $|E| \ge |V| 1$

(b) No; Counterexample:



- Q5. (a) 1. Suppose G = (V, E) is an acyclic, simple, undirected graph
 - 2. Take all complete subgraphs of G, $H_1(V_1, E_1), \cdots, H_n(V_n, E_n)$ which also form trees
 - 3. Then, $|E_1| = |V_1 1|, \dots, |E_n| = |V_n| 1$

(Theorem 10.5.2)

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- 4. $|E| = |E_1| + \dots + |E_n| = |V| n$
- 5. Therefore, $|E| \leq |V| 1$

(b) No; Counterexample:



- Q6. 1. Prove G is tree \rightarrow there is exactly one path between every pair of vertices:
 - 1.1. Suppose G = (V, E) is a tree
 - 1.2. G is connected and acyclic

(Definition of tree)

1.3. Any two vertices have a path between them

(Definition of connected)

- 1.4. Suppose there are vertices with two or more paths connecting them:
 - 1.4.1. Then, G is cyclic

(Lemma 10.5.5)

- 1.4.2. This contradicts 1.2
- 1.5. Hence, the supposition is false, and there is exactly one path between every pair of
- 2. Prove there is exactly one path between every pair of vertices $\to G$ is a tree:
 - 2.1. Suppose G = (V, E) is a graph with exactly one path between every pair of vertices
 - 2.2. G is connected

(Definition of connected)

- 2.3. Suppose G is cyclic:
 - 2.3.1. There is a cycle C in G

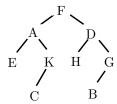
(Definition of cyclic)

- 2.3.2. Any two vertices in C have two paths connecting them
- 2.3.3. This contradicts 2.1
- 2.4. Hence, the supposition is false, and G is acyclic
- 2.5. G is a tree

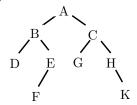
(Definition of tree)

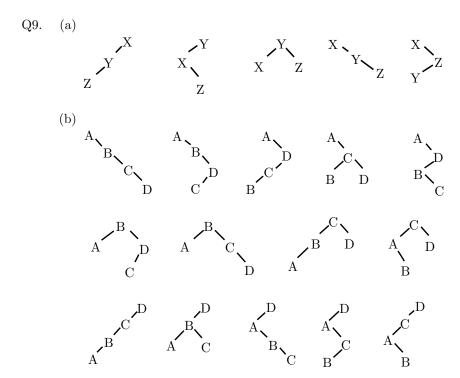
- 3. \therefore G is a tree \leftrightarrow there is exactly one path between every pair of vectors
- Q7. 1. Suppose G = (V, E) is a graph where each complete subgraph is a group of stones
 - 2. Initially, we have K_n which has $\frac{n(n-1)}{2}$ edges
 - 3. With each splitting into complete subgraphs of k_1, k_2 vertices, we remove $k_1 \times k_2$ edges

 - 4. Finally, we stop when there are exactly n subgraphs, each with 1 vertex with no edges 5. At this point, we will have removed all $\frac{n(n-1)}{2}$ edges, which is also the maximum earnt
- (a) Post-order: E C K A H B G D F Q8.



(b) Pre-order: A B D E F C G H K





Q10.

