## CS2040S Tutorial 2

AY 24/25 Sem 2—github/omgeta

Q1. (a.) 
$$T(n) = O(n)$$

(b.) 
$$T(n) = T(\frac{n}{2}) + O(n) = O(n)$$

(c.) 
$$T(n) = O(n^2)$$

(d.) 
$$T(n) = 2T(\frac{n}{2}) + n = O(n\log n)$$

(e.) 
$$T(n) = O(\phi^n)$$

(f.) 
$$T(n) = \log(1) + \log(2) + \dots + \log(n) < \log(n) + \log(n) + \dots + \log(n) = O(n \log n)$$

(g.) 
$$T(n) = O(n^2)$$

```
Q2. (a.) T(n) = T(n-1) + O(n) \le c1 + c2 + cn = \frac{cn(n+1)}{2} = O(n^2)
        private static void sort(int[] A, int n) {
             if (n > 0) {
                  sort(A, n-1);
                  int x = A[n];
                  int j = n-1;
                  while (j \ge 0 \&\& A[j] > x)  {
                       A[j+1] = A[j];
                       j--;
                  A[j+1] = x;
             }
        }
```

- 1. First, use SelectionSort to sort by b.
  - 2. Second, use MergeSort to sort by a.
  - 3. Since MergeSort is stable, it preserves the order of the pairs being sorted by b in 1.

```
(c.) T(n) = O(n \log n), S(n) = O(n)
   private static void sort(int[] a) {
     if (a.length <= 1) return;</pre>
     int[] tmp = new int[a.length];
     int n = 1;
     while (n < a.length) {</pre>
          for (int left = 0; left < a.length; left += 2 * n) {
              int mid = Math.min(left + n, a.length);
              int right = Math.min(left + 2 * n, a.length);
              merge(a, tmp, left, mid, right);
          for (int i = 0; i < a.length; i++) a[i] = tmp[i];</pre>
          n *= 2;
     }
   }
   private static void merge(int[] src, int[] dst, int left, int
      mid, int right) {
     int i = left, j = mid, k = left;
     while (i < mid && j < right)</pre>
          if (src[i] <= src[j]) dst[k++] = src[i++];</pre>
          else dst[k++] = src[j++];
     while (i < mid) dst[k++] = src[i++];
     while (j < right) dst[k++] = src[j++];
   }
```

- Q3. (a.) Use a pointer to keep track of where the head/tail of the stack/queue is
  - (b.) Keep a pointer of both the head and tail of the deque.
  - (c.) We would need to throw an error or fail silently if attempts are made to remove from an empty ADT, or add to a full ADT
  - (d.) Solution:

```
private static boolean check(String s) {
   Stack stack = new Stack();
   for (int i = 0; i < s.length(); i++) {
      char c = s.charAt(i);
      if (c == '(') stack.push(c);
      else if (c == ')')
        if (stack.isEmpty() || stack.pop() != '(') return false;
   }
   return stack.isEmpty();
}</pre>
```

(e.) Solution:

```
private static boolean check(String s) {
  Stack stack = new Stack();
  for (int i = 0; i < s.length(); i++) {</pre>
    char c = s.charAt(i);
    if (c == '(' || c == '{' || c == '[') stack.push(c);
    else if (c == ')' || c == '}' || c == ']')
      if (stack.isEmpty() || stack.pop() == getOpening(c))
   return false;
 }
  return stack.isEmpty();
private static char getOpening(char c) {
  if (c == ')') return '(';
  else if (c == '}') return '{';
 else if (c == ']') return '[';
  else return null;
}
```

- Q4. Solution: Traverse the elements, adding them to the stack. If the current element ≤ previous element, remove them from the stack. If the stack is empty, then the current visibility is infinite, else print the number of elements in the stack.
- Q5. Perform an iterative MergeSort in the same way as before but using the two queues.