# Vidyavardhini's College of Engineering & Technology Department of Artificial Intelligence and Data Science

#### **EXPERIMENT 03**

Aim: Implementation of Diffie Hellman Key exchange algorithm Theory:

### • Diffie-Hellman key exchange algorithm:

The Diffie-Hellman key exchange algorithm is a method for two parties to establish a shared secret over an insecure channel, which can then be used for secure communication. It allows two parties, typically named Alice and Bob, to jointly compute a shared secret without ever exchanging it directly. The key idea behind Diffie-Hellman is the use of modular exponentiation in a finite field.

Here's a brief description of how the algorithm works:

- 1. Both Alice and Bob agree on a large prime number 'p' and a primitive root 'g' modulo 'p'. These parameters 'p' and 'g' are public and can be shared over the insecure channel.
- 2. Alice chooses a private key 'a', computes 'A =  $g^a \mod p$ ', and sends 'A' to Bob.3. Bob chooses a private key 'b', computes 'B =  $g^b \mod p$ ', and sends 'B' to Alice.

5. The shared secret obtained by both Alice and Bob can be used as a symmetric key for secure

- 4. Now, both Alice and Bob can compute the shared secret: Alice computes `shared\_secret = B^a mod p`.
  - Bob computes `shared secret = A^b mod p`.

Since `B^a mod  $p = (g^b)^a \mod p = g^(ba) \mod p$ ` and `A^b mod  $p = (g^a)^b \mod p = g^(ab) \mod p$ `, both parties end up with the same shared secret `g^(ab) mod p`.

communication, allowing them to encrypt and decrypt messages exchanged between them. Diffie-Hellman key exchange provides a way for two parties to establish a shared secret even if an eavesdropper intercepts all communications between them. This is because the shared secret is never exchanged directly and relies on the computational difficulty of solving the discrete logarithm problem in a finite field.

#### Code:

```
def prime checker(p):
if p < 1:
    return -1 elif p >
1:
       if p == 2:
return 1
             for i in
range(2, p):
                    if p
% i == 0:
return -1
              return 1
def primitive_check(g, p, L):
for i in range(1, p):
    L.append(pow(g, i) % p)
for i in range(1, p):
```

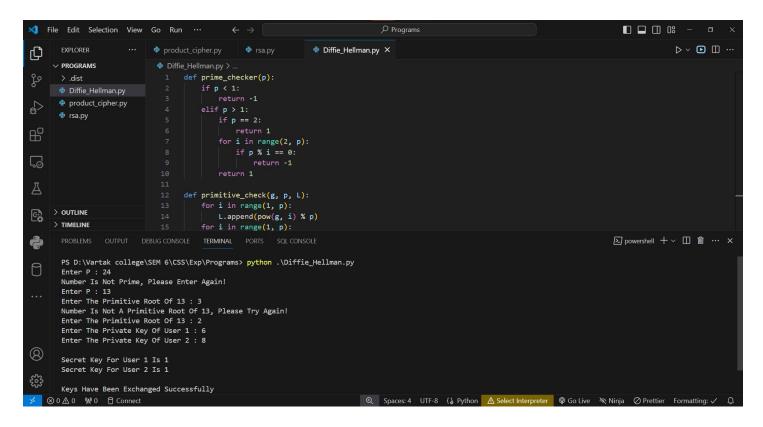


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```
L.count(i) > 1:
                                                                 L.clear()
                 return -1 return 1
                 I = [] while
                 True:
                      P = int(input("Enter P:")) if prime checker(P) == -
                                  print("Number Is Not Prime, Please Enter
                 Again!")
                                               continue
                                                                          break
                 while True:
                      G = int(input(f"Enter The Primitive Root Of {P}: ")) if
                                                                                                 print(f"Number Is Not A Primitive
                 primitive check(G, P, I) == -1:
                 Root Of {P}, Please Try Again!")
                                                                                                     continue break
                 # Private Keys x1, x2 = int(input("Enter The Private Key Of User 1:")), int(input("Enter The Private Key
                 Of User 2:
                 ")) while
                 True:
                      if x1 >= P \text{ or } x2 >= P:
                           print(f"Private Key Of Both The Users Should Be Less Than {P}!")
                 continue break y1, y2 = pow(G, x1) \% P, pow(G, x2) \% P k1, k2 = pow(y2, y2) \% P k1
                 x1) % P, pow(y1, x2) % P print(f"\nSecret Key For User 1 Is {k1}\nSecret
                 Key For User 2 Is {k2}\n")
                 if k1 == k2: print("Keys Have Been Exchanged
                 Successfully") else:
                      print("Keys Have Not Been Exchanged Successfully")
Output:
```



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## **Conclusion:**

The Diffie-Hellman key exchange algorithm was implemented, allowing two parties to establish a shared secret over an insecure channel. The algorithm relies on the computational difficulty of solving the discrete logarithm problem in a finite field, ensuring security even if an eavesdropper intercepts all communications. By agreeing on a large prime number p and a primitive root g modulo p, both parties compute their public keys and derive a shared secret using modular exponentiation. The shared secret can then be used as a symmetric key for secure communication.