

## **Practical No. 4**

### **Title:**

Write a Program to Solve 0-1 Knapsack Problem Using Dynamic Programming

### **Aim:**

To implement a program that solves the 0-1 Knapsack problem using the dynamic programming approach to maximize total profit without exceeding the knapsack capacity.

### **Objective:**

To understand and apply the dynamic programming method to solve the 0-1 Knapsack problem by building a table that stores intermediate results to avoid redundant calculations.

### **Hardware Requirements:**

- Processor: Intel Core i3 or higher
- RAM: Minimum 2 GB
- Storage: Minimum 100 MB free space
- Input Devices: Keyboard and Mouse
- Output Device: Monitor

### **Software Requirements:**

- Operating System: Windows / Ubuntu
- Programming Language: Python 3.x (or C/C++/Java)
- IDE/Text Editor: VS Code / PyCharm / Notepad++
- Terminal or Command Prompt

## Theory:

The 0-1 Knapsack problem is a classical optimization problem where:

- There are  $n$  items, each with weight  $w[i]$  and profit  $p[i]$ .
- The knapsack has a maximum capacity  $W$ .
- The objective is to select items (either whole — 0 or 1, no fractions) to maximize profit without exceeding capacity.

Dynamic Programming solves this by constructing a 2D table where rows represent items and columns represent weight capacities, building solutions for smaller subproblems and combining them.

## Algorithm (Dynamic Programming):

1. Create a 2D array  $dp[n+1][W+1]$  where  $dp[i][j]$  represents max profit using first  $i$  items and capacity  $j$ .
2. Initialize  $dp[0][j] = 0$  for all  $j$ , and  $dp[i][0] = 0$  for all  $i$ .
3. For each item  $i$  from 1 to  $n$ :
  - For each capacity  $j$  from 1 to  $W$ :
    - If weight of  $i$ th item  $\leq j$ :  
 $dp[i][j] = \max(dp[i-1][j], \text{profit}[i-1] + dp[i-1][j - \text{weight}[i-1]])$
    - Else:  
 $dp[i][j] = dp[i-1][j]$
4. The answer is  $dp[n][W]$ .

## Conclusion:

Dynamic programming is an effective approach to solve the 0-1 Knapsack problem by solving overlapping subproblems and storing intermediate results, thus reducing time complexity from exponential to polynomial.