

## LISTE DER NOCH ZU ERLEDIGENDEN PUNKTE

maybe improve abstract . . . . .	xi
translate abstract . . . . .	xi
Do I need a list of tables? . . . . .	xvi
write . . . . .	2
add reference to [8], mention that Weske says pretty much the same . . . . .	4
Much like the process instance lifecycle, each activity follows a cycle - introduce activity lifecycle/control flow . . . . .	4
Abbildung: activity lifecycle . . . . .	4
references to the other languages . . . . .	6
If we speak of the bpmn spec, we always mean version 2.0 . . .	6
- introduce the other essential elements, types of events . . . .	7
Abbildung: Sample BPMN choreography model, maybe the same as before, but with 2 lanes and communication in-between . . .	7
More precisely, Events are defined as follows... . . . . .	8
Abbildung: state-transition-diagram of temporal order . . . . .	8
Abbildung: pub/sub and event processing architecture . . . . .	9
'relational' db . . . . .	9
Abbildung: stream processing . . . . .	9
more on epl, what can it do? windows, joins, filters,patterns . .	10
ref . . . . .	13
I need to work with a precise lifecycle . . . . .	13
show a cep query for that scenario to make it more precise . . .	14
Consider more possible event occurrence times to prepare for the next chapter . . . . .	15
show a cep query for that scenario to make it more precise . . .	16
improve figure: include EOS notions, save space, ? . . . . .	18
add a back reference to the examples. In example XY, events can occur before... whereas in example... . . . . .	18
by event I usually mean event from a cep pub/sub event source. The requirements relate to external intermediate message events. A limitation that is made in accordance with the motivated scenario. though some of the aspects might ap- ply to other types of eventy, like the start event. . . . .	19
which functionality should be evaluated exactly?: all occurrence scenarios, but no buffer policies. The buffer will always store the last version of the event and also deliver that version. . . . .	23

Notably, there always is a certain delay when consuming an event through an intermediate catch event. The BPMN specification itself states that <i>the handling consists of waiting for the Event to occur</i> . However, that delay can theoretically be infinitely small. <– maybe mention this at the beginning of ‘problem statement’ to differentiate between normal delay and unnecessary delay? . . . . .	24
could exemplify this at ex1.2, In the previously mentioned Example 1.2 . . . . .	25
Do I want to work with the requirements or not??? . . . . .	29
where is the code? . . . . .	32
give them short names for reference? Or make one of them Requirement 5 . . . . .	33
there could be a requirement that states that there should be no additional process elements unless there is an explicit subscription time . . . . .	33
what exactly is the issue here? . . . . .	34
write connector to next chapter: it has been shown that... therefore we do ... . . . . .	34
Introduce a new concept. three pillars: bpmn-x, process engine behavior, event handling api . . . . .	35
cite braun2015behind . . . . .	35
this is the explanation for the process designer? or is there an additional event engineer? . . . . .	35
mention that others have extended the events instead, but this proposal goes another way . . . . .	36
Abbildung: A UML diagram of Message, MessageFlow, MessageEventDefinition, Catch Event, Boundary Event, Receive Task . . . . .	36
explain that the Message is a common element in BPMN models. It is the main generic type used for communication, collaboration. The proposed extension to the Message type is only to be used by Message Receive situations in Intermediate Catch Event, the Boundary Catch Event and the Receive Task. . . . .	36
ref background . . . . .	36
Find this source; explain what they do (different) . . . . .	37
By default, there is no interference between the buffers of different messages, process instances or processes. Each buffer instance will contain the latest information as if it was the only buffer in the system. Performance improvements to avoid duplicate buffer content will be managed by the system without explicit action by the user. Section ... later introduces a shared, more complex usage scenario of the event buffers. . . . .	38
ref background . . . . .	38

reorganize this chapter . . . . .	38
For each of the options: Define exactly (according to BPMN spec or standard literature), when in the flow the subscription is executed. . . . .	38
which example to reference? . . . . .	39
show what that would look like in the example. Maybe some XML? . . . . .	39
it follows the ideas presented in ... but through referencing the message element . . . . .	39
latest concept (xsd) states that it's an extension of tTask . . . . .	39
Abbildung: Example 2 using an explicit subscription task . . . . .	39
As an improvement to the options for subscription time, there could be an option "ASAP", so that the process engine is- sues the subscription automatically as soon as the required process data becomes available . . . . .	39
this could be added in the requirements and referenced . . . . .	40
only if subscriptionTime not 'on depl' . . . . .	40
reference esper docs 5.13.1. Joining SQL Query Results . . . . .	40
missingref: dependent example . . . . .	40
What does this mean for the process designer? A model that can take a state in that a subscription shall be issued, though the data is unavailable, is invalid. When will an error oc- cur? . . . . .	40
add ref to sankalitas paper; do I elaborate on the changes made in comparison to the paper? + the reasons? . . . . .	41
ref external . . . . .	41
ref example . . . . .	41
for the reader it would not be clear, what a 'buffer instance' is . . . . .	41
write text for given keypoints . . . . .	42
Talk about an alternative solution? Table of events for a certain event source. The buffer is already there for me to pick from when designing the process. . . . .	42
missingref . . . . .	45
ref . . . . .	45
Maybe: we consider these two operations as given, because com- mon event processing platforms have these in common, but there is not common event buffering concept yet. even though e.g. esper has something which goes in that direc- tion: output clauses . . . . .	46
note that splitting this into four actions is not a new concept. actually thats how event processing platforms normally operate, only without creation of the buffer and without delivering the buffered events . . . . .	46
what kind of API should this be? REST? Java? none specifically, but the reader might want clarification when reading this. . . . .	46

be clearer about notificationPath. Specify in background and reference. also about addNotificationRecipient . . . . .	46
add table with buffer policies, their possible values and the default val . . . . .	47
show the steps with sample data from one of the examples . .	47
Abbildung: maybe a uml sequence diagram. Or other diagram format? . . . . .	47
maybe provide a swagger definition for this? . . . . .	47
Generally, extending the event processing platform is advisable?	
Or put up brief pros and cons of each of the options . . .	48
Abbildung: three options to implement the event buffering and api extension . . . . .	48
the communication overhead must be considered: if many events are sent to the buffer, but they are rarely issued to a process instance, then it will make sense to place the buffer closer to the event engine w.r.t. the the involved comm overhead. Though it must be noted, that if events occur rarely, but many process instances consume them, than there will be an additional overhead to send events from the buffer to the instance. in that scenario, it might be advantageous to implement the event buffering close to the process engine . . . . .	48
write . . . . .	48
be more precise about the time the calls should be executed (if possible). "reached"? "completed"? use the right bpmn words	50
Do I want to write about extended validity checks? soundness?	
The question would be: What happens if the model erroneous w.r.t. to the bpmn extension? . . . . .	50
Time of un-subscription also must be clarified in bpmnx . . . .	50
write conclusions that repeat how each requirement has been fulfilled . . . . .	50
write from keypoints . . . . .	51
Application to examples missing . . . . .	51
write from keypoints . . . . .	51
Abbildung: UML class diagram of: BufferedLiveQueryListener extends LiveQueryListener, BufferManager (list<EventBuffer>   createBuffer, updatebuffer, deleteBuffer) implements runnable, EventBuffer (Policies   update, maintain, retrieve), Constants	53
??? Swagger definition of the implemented rest API . . . . .	55
ref/footnote camunda documentation . . . . .	56
can i reference another part of the thesis? could be mentioned in buffered event handling . . . . .	56
that means that in the implementation we need additional configuration values -> implementation chapter . . . . .	57

could mention that registerQuery can be bound either to an instance or to a process definition. this differentiation is made based on the subscriptionDefinition which is passed to the listener on creation. . . . .	59
Abbildung: UML of all involved classes . . . . .	59
Abbildung: maybe provide pseudo code? . . . . .	59
complete xml example of a BPMN model using the extensions	66



DENNIS WOLF

FLEXIBLE EVENT SUBSCRIPTION  
IN BUSINESS PROCESSES





# FLEXIBLE EVENT SUBSCRIPTION IN BUSINESS PROCESSES

DENNIS WOLF



Digital Engineering • Universität Potsdam

< Any Subtitle? >

August 2017 – version 1

Dennis Wolf: *Flexible Event Subscription*  
*in Business Processes, < Any Subtitle? >*, © August 2017

## ABSTRACT

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Business Processes have become an essential tool in organizing, documenting and executing company workflows while Event Processing can be used as a powerful tool to increase their flexibility especially in distributed scenarios. The publish-subscribe paradigm is commonly used when communicating with complex event processing platforms, nevertheless prominent process modelling notations do not specify how to handle event subscription.

At the example of BPMN 2.0, the first part of this work illustrates the need for a flexible usage of event subscription in process models and derives new requirements for process modelling notations. An assessment of the coverage of these requirements in BPMN 2.0 is presented and shortcomings are pointed out.

Based on the identified requirements, this work presents a new concept for handling event subscription in business process management solutions, predominantly built on the notion of event buffers. The concept includes an extension to the BPMN meta model, specifies the semantics and API of a new event buffering module and describes the changes necessary to the behaviour of the process engine.

For evaluation purposes, the concept has been implemented as a reusable Camunda Process Engine Plugin that interacts with the academic Complex Event Processing Platform UNICORN.

maybe improve abstract

## ZUSAMMENFASSUNG

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translate abstract

Kurze Zusammenfassung des Inhaltes in deutscher Sprache...



## CONTENTS

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1	INTRODUCTION	1
1.1	Motivation	2
1.2	Contribution	2
1.3	Structure	2
2	BACKGROUND ON EVENT-DRIVEN BUSINESS PROCESS MANAGEMENT	3
2.1	Business Process Management	3
2.1.1	Business Process Meta Model and Activity Life-cycle	4
2.1.2	Business Process Management Systems	5
2.1.3	Business Process Model and Notation	6
2.2	Complex Event Processing	8
2.3	Event-driven Business Process Management	10
3	PROBLEM STATEMENT	13
3.1	Motivating Examples	13
3.2	Event Occurrence Scenarios	17
3.3	Requirements Definition	19
4	ASSESSMENT OF CURRENT BUSINESS PROCESS MANAGEMENT SOLUTIONS	23
4.1	BPMN Models in presence of the Event Occurrence Scenarios	23
4.2	Implementation of Early Event Subscription using standard Camunda	29
4.3	Discussion	33
4.4	Requirements extension	34
5	FLEXIBLE EVENT SUBSCRIPTION IN BPMN	35
5.1	BPMN Extension	35
5.1.1	Adding basic subscription information	36
5.1.2	The time of event subscription modeled in BPMN	37
5.1.3	Using Process Variables in Event Queries	40
5.1.4	Advanced Buffer Parameters	41
5.2	Design Decisions	42
6	AUTOMATIC SUBSCRIPTION HANDLING	45
6.1	Buffered Event Processing	45
6.2	Extended Process Engine Behavior	48
7	REFERENCE IMPLEMENTATION	51
7.1	Extending the Event Processing Platform Unicorn	51
7.1.1	Event Buffering	52
7.1.2	REST API Extension	54
7.2	Event Subscription Handling in Camunda	56
7.2.1	Employing ExecutionListeners in a Process Engine Plugin	56

	7.2.2 Managing event Subscriptions at Runtime	57
8	RELATED WORK	61
9	CONCLUSIONS	63
	9.1 Discussion	63
	9.2 Future Work	63
A	APPENDIX	65
	BIBLIOGRAPHY	67

## LIST OF FIGURES

---

Figure 1	Business process management systems architecture model (see [30], p. 120)	5
Figure 2	Simple BPMN model of issuing a quote for car rental	7
Figure 3	BPMN Model of a Logistics Process using events for route optimization (Example 1.1)	14
Figure 4	Transport via English Channel that is timed to a delivery slot (Example 1.2)	15
Figure 5	Model of a retail order management process (Example 2)	16
Figure 6	Possible event occurrence times in relation to a process execution life cycle	18
Figure 7	Generic process terminating after an intermediate message catch event	24
Figure 8	Tu duo titolo debitas latente	25
Figure 9	Event Element in parallel process flow	25
Figure 10	Event Buffering through an auxiliary buffering process	27
Figure 11	Generic Example Process in Camunda for Occurrence Scenario <i>EOS</i> <sub>3</sub>	29
Figure 12	Generic Example Process in Camunda for Occurrence Scenarios <i>EOS</i> <sub>4</sub> and <i>EOS</i> <sub>5</sub>	30
Figure 13	Auxiliary Buffering Process in the Camunda Modeler	31
Figure 14	Auxiliary Event Delivery Process in Camunda Modeler	32
Figure 15	The concept for flexible event subscription involves three modules: A BPMN extension, enhanced process engine behavior and buffered event handling.	35
Figure 16	Flight Booking process using a consuming buffer	42
Figure 17	Shared consuming buffer in Complaints Handling	42
Figure 18	Architecture for flexible event subscription in Camunda and Unicorn	52

## LIST OF TABLES

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Do I need a list of tables?

## LISTINGS

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Listing 1	Sample Query in Esper EPL	10
Listing 2	Example of a JSON notification sent by UNICORN	56
Listing 3	XSD schema of the BPMN extension for flexible event subscription	65

## ACRONYMS

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API	Application Programming Interface
BPM	Business Process Management
BPMN	Business Process Model and Notation
CEP	Complex Event Processing
EdBPM	Event-driven Business Process Management
EOS	Event Occurrence Scenario
EPL	Event Processing Language
JVM	Java Virtual Machine
HTTP	Hypertext Transfer Protocol
PEP	Process Engine Plugin
REST	Representational State Transfer
SQL	Structured Query Language
UML	Unified Modeling Language
XML	Extensible Markup Language



## INTRODUCTION

---

Given the increasing competition on the global market place, companies are seeking to improve their products while reducing costs. In many areas, Business Process Management has been chosen as one of the tools to help stay competitive. Especially large enterprises, but also small to medium businesses formalize their workflows in business process models to allow archiving, documentation and automated management and execution.

With Business Process Technology in constant progression, the opportunities that the field has to offer are ever growing. Since the recent years, many efforts have been dedicated towards bringing together business processes and Complex Event Processing. By the help of Event Processing Systems, companies are trying to get a hold of the exponentially growing amounts of data that occur in today's IT environment. Incorporated in process executions, events heavily increase their capabilities and flexibility. They can be utilized for intraorganizational communication between processes or business departments, but also allow to respond to external situations within seconds or milliseconds. The Business Process Model and Notation (BPMN), an industry standard for representing business processes both visually and textually, natively supports the use of events in a plethora of ways. Events are considered a main building block of a feature-rich process modeling language and can, for instance, be used for instantiating processes, communicating between process participants or to support decisions.

An interaction with a Complex Event Processing platform generally follows the publish-subscribe paradigm: The event consumer contacts the platform and issues a subscription to a specific subset of available events. An event producer, for example a Vehicle providing its current GPS location, publishes information to the event processing platform, which is then forwarded to every consumer that had subscribed. Intermediate Events are a basic way to implement event-based communication in BPMN and facilitate the reception of external message events. Nevertheless, the BPMN specification does not specifically consider the publish-subscribe workflow and provides limited capabilities when it comes to incorporating event subscription and un-subscription operations in business process models. This work investigates the consequences of this lack of specification and provides a design and implementation to overcome the identified shortcomings.

### 1.1 MOTIVATION

why? - cant live without complex event processing in business processes, increasing demand for using events in processes - pub/sub is a fundamental part of using events in BPs - still it is not considered in bpmn - there is research on including the subscription query in the model, but the subscription time is not further defined - the bpmn specification says "..", that leaves us with a very limited listening time

- but distributed setup -> hard to control - we need subscription before occurrence time - issues will occur when mis-used, delay or blockage, significant time and financial loss -> brief, textual example
- to ensure the efficient use of events in processes, a more flexible use of subscription is necessary

- > the problem will be further illustrated in motivating examples

### 1.2 CONTRIBUTION

Working towards a more flexible use of event subscription in business processes...

- (o) reviewing the problem from different perspectives, deriving requirements (1) assessment of the capabilities of standard bpmn (2) Proposition of a BPMN extension for flexible event subscription, its advantages (3) derived requirements to process engines and CEP platforms (4) A reference implementation using Camunda and UNICORN

### 1.3 STRUCTURE

write

## BACKGROUND ON EVENT-DRIVEN BUSINESS PROCESS MANAGEMENT

---

### 2.1 BUSINESS PROCESS MANAGEMENT

With its origins dating back to the process orientation trend of the 1990s, Business Process Management (BPM) has meanwhile become a mainstream tool to support organizations. It had been noted, that company workflows can essentially be broken down into activities that are executed in a coordinated manner by one or more parties. A certain group of activities thereby form a process which is executed within an organization. More precisely, Weske [30] defines a single *business process* as follows:

**DEFINITION 1 (BUSINESS PROCESS):** A *business process* consists of a set of activities that are performed in coordination to realize a business goal. Each business process is enacted by a single organization, but it may interact with processes performed by other organizations.

The term *Business process management* describes the techniques available to develop and support processes throughout their life-cycle. It is grounded in the use of explicit process representations which ultimately allow the exchange, analysis and reproduction of the workflows. This process specification is referred to as the *business process model*, composed mainly of activities and the rules that are necessary to coordinate their execution. When a process is performed according to its model, the single execution is called *process instance*. Based on a process model, the number of possible instances is theoretically unbounded.

**DEFINITION 2 (BUSINESS PROCESS MODEL):** A *business process model* consists of a set of activity models and execution constraints between them. A *business process instance* represents a concrete case in the operational business of a company, consisting of activity instances. [30, p. 7]

The lifecycle of a business process can be described as a cycle of four phases in that numerous stakeholders interact and contribute depending on their specialization. Process development starts with a *Design & Analysis* phase which yields a refined and validated business process model. In the following *Configuration* phase, it is necessary to prepare the process implementation, select the means and an environment to run the process in. The action of making the process runnable

in the execution environment is called *process deployment*. After these preparations the process can be enacted in daily business while its current state is monitored and system maintenance is performed if necessary (*Enactment* phase). A single process execution begins with the *process instantiation*, when the process has succeeded or is aborted, we say the process is *terminated*. During the enactment, system and stakeholders can start collecting performance indicators and process execution logs to allow evaluating the quality of the process specification. If that *Evaluation* step reveals deficiencies, the lifecycle starts over by re-entering the design phase. The *process un-deployment* is performed if necessary, so that no new instances of the process can be started. [30, p. 11 ff.]

add reference to [8], mention that Weske says pretty much the same

### 2.1.1 Business Process Meta Model and Activity Lifecycle

Attempting to...

**DEFINITION 3 (BUSINESS PROCESS META MODEL):** Let  $C$  be a set of control flow constructs.  $P = (N, E, \text{type})$  is a *process model* if it consists of a set  $N$  of nodes, and a set  $E$  of edges. [30], p. 91

- $N = N_A \cup N_E \cup N_G$ , where  $N_A$  is a set of activity models,  $N_E$  is a set of event models and  $N_G$  is a set of gateway models. These sets are mutually disjoint.
- $E$  is a set of directed edges between nodes, such that  $E \subseteq N \times N$ , representing control flow.
- $\text{type} : N_G \rightarrow C$  assigns to each gateway model a control flow construct.

Much like the process instance lifecycle, each activity follows a cycle - introduce activity lifecycle/control flow



While traditionally activities are executed manually by company staff following the written process specifications, computer systems are used today to drive the execution and enforcement of business processes and organizational rules. The generic software systems utilized for that purpose are introduced in the following section.

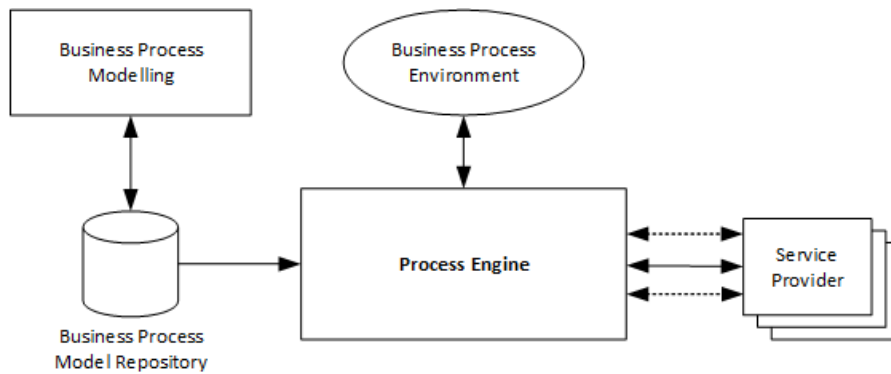


Figure 1: Business process management systems architecture model (see [30], p. 120)

### 2.1.2 Business Process Management Systems

The implementation of business processes has developed from a manual execution guided by business rules to a fully automatized execution in a specialized IT environment. One of the main reasons for BPM's growing popularity is that in today's fast-paced economy, a large part of the business activity is either supported by computers or even carried out autonomously by them. The specialized software systems that are utilized to support the enactment of business processes are referred to as *business process management systems*.

**PROCESS MANAGEMENT ARCHITECTURE** A typical IT infrastructure for driving business processes is illustrated in Figure 1. Five principal building blocks are considered which will be explained in the following.

With reference to the business process lifecycle, the visualized scenario commences with the *Business Process Modeling*. As a result of the *Design & Analysis* phase, new process models are created and refined to be stored in the *Business Process Model Repository*. The relation between the two elements includes writing new models to the repository as well as reading models for review and further modification. Given that the desired model is approved for enactment, the process gets deployed to the *Process Engine* as part of the configuration step. The process engine is the heart of the execution environment. It performs the execution of the processes from deployment until un-deployment, while the enactment and instantiation is controlled by the *Business Process Environment*. An indefinite number of *Service Providers* realize application services to support the process execution. A service provider can be a software module but also a knowledge worker performing a particular process step.

**THE CAMUNDA BUSINESS PROCESS ENGINE** A large, growing number of process engines is available on the market, including solu-

tions from IT giants like SAP, IBM and Oracle. In this work, *Camunda BPM* [12] has been chosen to illustrate implementations. As of August 2017, the software product is available in version 7.7.0 and comes in a commercial, regularly updated version and in a free, community-driven solution that is updated with every major release. Camunda is popular among the research community as the source code is openly available, the product is mature, but actively developed and offers comprehensive support for BPMN 2.0. It is designed to be extensible and easily modifiable to adapt to custom requirements. *Camunda BPM* comprises a modeling tool, the Camunda process engine core and a number of browser-based user-interfaces to control process enactment and monitor execution state. Chapter 7 will provide further details about the engine architecture and extension mechanisms.

### 2.1.3 Business Process Model and Notation

Given the general semantics of business processes, a specific modeling notation has to be selected to express an informal process description in a formal, interchangeable way. Different languages and notations have become available over the years, each serving different specializations. Kossak et al. [19] organize some of the more popular languages as follows: A subset of them are focused on the control flow of business processes, for instance Business Process Model and Notation (BPMN) [26], Yet Another Workflow Language and Petri Nets; some focus on object-orientation, like the Unified Modeling Language (UML) activity diagrams and use case diagrams; some are data-flow oriented, e. g. the Structured Analysis and Design Technique.

#### references to the other languages

Among these, the BPMN has developed into a widely-adopted industry standard, also becoming ISO-standard in 2013 [15]. The standard is developed by the Object Management Group [25] and now available in version 2.0 (January 2011) after being first released in January 2008.

#### If we speak of the bpmn spec, we always mean version 2.0

BPMN can be understood as an extension to the abstract business process meta model (Section 2.1.1) adding a comprehensive catalog of visual representations and semantic constructs on top of a meta model. Furthermore, one of the most important features of its latest version is the a standardized interchange format provided through an XML specification, as [30] points out. As emphasized by Muehlen, Recker, and Indulska [24], the increased expressiveness of modern languages like BPMN comes at the cost of an increased complexity. An aspect that, apparently, did not stop it from gaining popularity.

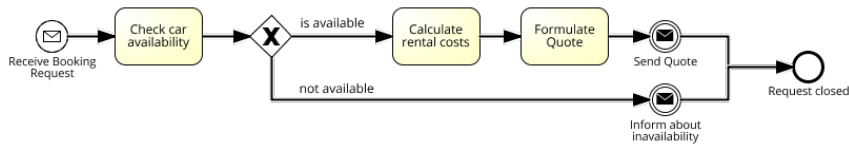



Figure 2: Simple BPMN model of issuing a quote for car rental

**ESSENTIAL ELEMENTS OF A BPMN MODEL** Following the abstract business process meta model, the core elements in any BPMN model are flow elements (nodes) and connecting objects (edges). Flow elements can be either *Events*, *Activities* or *Gateways*, each of them coming in different variations. This section will introduce a subset of the elements available through the [BPMN](#) specification to build the foundation to comprehend the thoughts presented in this work.

Figure 2 shows how a booking request might be handled in a car rental business. Circular elements represent events, diamond-shaped elements are gateways. Activities are visualized by rectangles with rounded corners. The given process gets instantiated whenever a booking request request is received from a customer, shown as a Message Start event. As a first step, the employee assigned to handle the request must check if the desired car is available. To that follows an exclusive OR-Gateway, distinguishing the further process flow depending on the availability of the car. If the car is available, the quote must be created in two sequential activities to be then sent out by the *Send Message Event*. If the car is not available, the customer is informed about the closing of his request. In either case, the process ends with an *End Event* after the customer was informed about the result of his request.

- introduce the other essential elements, types of events

**PROCESS CHOREOGRAPHIES** - talk about how several parties can interact => choreography - Message - choreographies



Fehlende  
Abbildung

Sample BPMN choreography model,  
maybe the same as before, but with 2  
lanes and communication in-between

## 2.2 COMPLEX EVENT PROCESSING

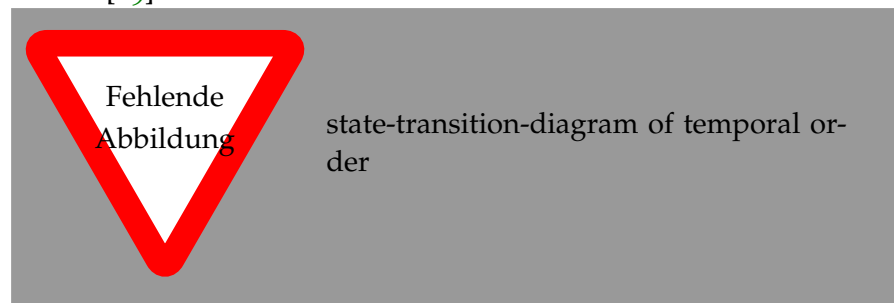
The IT world is facing an exponential increase in the amount of produced data. A significant part of this data are pieces of information about real-life occurrences, such as a current sensor value, an interaction on a website or the location of a vehicle on the road. We call this kind of strongly time-related information an *Event* and the according computer science field Complex Event Processing (CEP) [10]. - events are of a certain type

More precisely, Events are defined as follows...

from <http://www.softnet.tuc.gr/minos/Papers/rtstreams15.pdf> - Event tuples are defined by [11] as tuples  $e = \langle s, t \rangle$ , where  $e$  represents the event of interest,  $s$  refers to a list of attributes and  $t$  is a list of timestamps, the first being the start of the event and the last the end of it

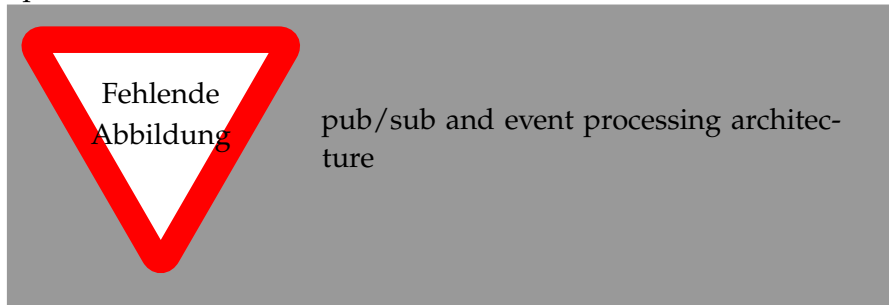
Four major components take part in an event processing network: (a) An *event producer* provides information to the system, (b) an *event agent* processes the occurring information, so that it can be delivered to the *event consumer* (c). The components are linked through *event channels* (d). Typically, a so-called *CEP Engine* (also: CEP platform) is at the heart of the system, taking the role of an event agent. Modern CEP platforms are trimmed to maximum efficiency, being able to process hundreds of thousands of events every minute. Their main purpose is to accept incoming events from event producers, filter and match them according to selection criteria and, finally, derive a new event occurrence to be sent to the registered event consumers.

**THE PUBLISH/SUBSCRIBE PRINCIPLE** In event-based architectures, communication takes place according to the *Publish/Subscribe Principle*. The concept essentially demands that an event processing middleware publishes events to processes only after they have issued a subscription for these events. Consequently, there is a strict temporal order between the actions subscribe, consume and un-subscribe. The consumption and un-subscription can only happen after the subscription. Once an un-subscription has been issued, no consumption can follow. [29]





One of the main advantages of this principle is, that the involved parties are *referentially decoupled*. They do not need to explicitly refer to each other, an aspect that is also acknowledged in [10]. Their decoupled nature facilitates the management and development of event processing networks. Event producers and consumers might change frequently. Whenever a new event source is available it can be connected to the CEP platform without considering all future consumers. Consumers can subscribe and un-subscribe without influencing the operations on the consumer side.

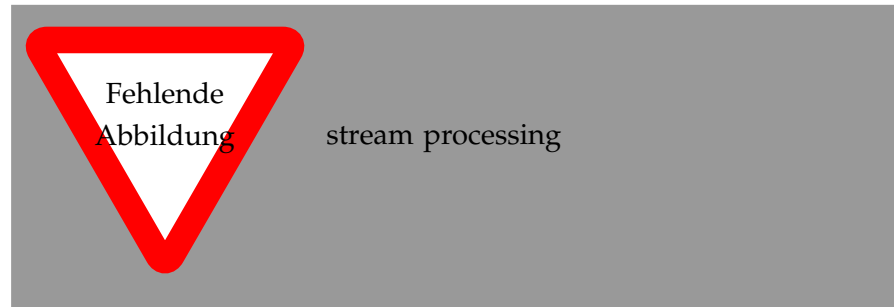


**STREAM PROCESSING** To be able to cope with potentially large amounts of data, Complex Event Processing Platforms work on the basis of *stream processing*.

'relational' db

In a traditional database, information is stored for an indefinite amount of time. When a user queries the data store, the system processes the tabular data and calculates the requested result. The advantage of this approach is that the user can access historic data at any time, as long as it is not explicitly deleted from the database. In many occasions, the amount of available data significantly surpasses the storage capacities and that concept can no longer be followed.

Stream processing addresses the mentioned challenge by largely reducing the amount of data that is persisted in the system. Instead, it is the goal to keep only those pieces of information, that are necessary to process a result for currently registered queries. Incoming data objects, or events in the case of a CEP platform, enter the event stream and are immediately evaluated against all existing query expressions. If the information is not required to process any of the queries, it is deleted instantly. As a consequence, a certain event can not be part of a query result, if that query has been registered after the occurrence of the event. In case aggregated information is demanded by the query, the stream processor will internally store aggregated information, but not keep every information that led to the aggregated value. [18]



**EVENT QUERY LANGUAGES** The subscription to an event in a CEP platform is primarily defined by an *Event Query*. Many modern event query languages are inspired by [SQL](#), but cannot be entirely compliant due to the different underlying data processing concept. When formulating event queries, it is essential to consider the stream processing principle.

For the illustration of the concepts, this thesis relies on the Esper Event Processing Language ([EPL](#)) [[14](#)], utilized in Esper-based event processing engines like the one employed in [Chapter 7](#).

more on epl, what can it do? windows, joins, filters, patterns

Listing 1: Sample Query in Esper EPL

```
SELECT time, delay, delayreason
FROM eurotunnel
WHERE delay > 30
```

### 2.3 EVENT-DRIVEN BUSINESS PROCESS MANAGEMENT

- see opheretzion p.17 and where are we now - maybe also reference baumgrass paper / GET project | Towards a Methodology for the Engineering of Event-Driven Process Applications (BG2016) - +

- there are 2 use cases for events in bp: (1) for process monitoring (=> pe is event producer) > what's it about this thesis focusses on (2) as it investigates subscription mechanisms during the execution of a process

(2) for driving processes (=> pe is event consumer) > Event-Driven Process Control (event-driven business process management: where are we now)

- interplay BPT to CEP platforms, a subscription must be issued
- no standard yet available to do subscription in bpmn - it must be assumed that the subscription is either already active or explicitly modeled in BPMN, e.g. using a service task - OR given the BPMN spec it is generally assumed that the subscription is executed as soon as an event is enabled - further analysis of this topic is provided in ...
- the time of event subscription is clear for start events
  - Correlating events to process instances

- exercise through one example



## PROBLEM STATEMENT

---

Event subscription is an essential part of event-driven architectures and must consequently be considered when using events in business processes. While the use of external events is an active topic in research, most approaches only briefly discuss the subscription mechanism and use the [BPMN](#) semantics for orientation.

ref

Though event subscription is not directly addressed in the BPMN specification, the descriptions on intermediate events state:

*"For Intermediate Events, the handling consists of waiting for the Event to occur. Waiting starts when the Intermediate Event is reached. Once the Event occurs, it is consumed."* [26], p. 440

The usual interpretation of this excerpt is that the subscription to the event takes place as soon as the event gets enabled, the un-subscription when the event terminates.

I need to work with a precise lifecycle

As pointed out by Mandal et. al. [23], these subscription semantics significantly limit the flexibility of using events in business processes and can cause undesired behavior and fault. Due to the strict temporal order between event subscription, occurrence and un-subscription as required by the publish-subscribe paradigm, any events that occur before the enabling of the event element will be ignored by the process execution. That reduces the timespan for events to occur to a potentially small part of the total execution time and means that crucial events might be missed which can delay or block a process execution unnecessarily.

In the following section, the problem is further illustrated at the example of two sample scenarios. The observed situations then lead to the definition of *Event Occurrence Scenarios* and the derivation of a set of requirements that must be fulfilled by a mechanism for flexible event subscription in business processes.

### 3.1 MOTIVATING EXAMPLES

To allow a better understanding of the issue, event-driven use-cases from two different domains are presented in the following. The cases are revealed through their standard BPMN representation. It is illustrated, why the time of event subscription is of great importance which motivates to study the mechanics and implications of event subscription in business processes.

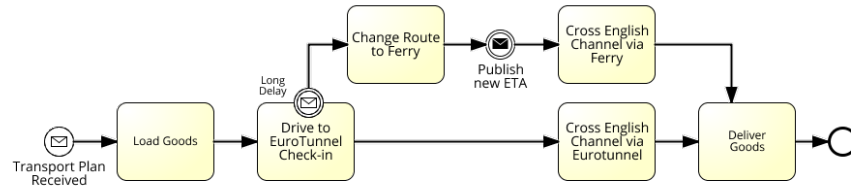


Figure 3: BPMN Model of a Logistics Process using events for route optimization (Example 1.1)

**DELAY OF A LOGISTICS PROCESS** The first example (Figure 3) is taken from the logistics domain and shows a truck transport that has to cross the English Channel. The truck driver receives the transport plan for his next tour from France to the UK. By default, the company crosses the Channel using the Eurotunnel, an underground train connection between London and Paris.

After loading the goods at the factory, the truck will head towards the check-in location of the Eurotunnel. If everything runs on schedule, the truck crosses the channel on the train and then delivers the goods in Great Britain. Alternatively, the process considers a route using the ferry from Calais (FR) to Dover (UK). The Eurotunnel administration publishes delay information approximately every 30 minutes through an RSS feed on their website. While it mostly operates on schedule, delays ranging from 15 minutes to several hours occur regularly. It can happen that new information is not published for multiple hours. Significant delay events (delay > 30 minutes) are received through a boundary catching message event attached to the activity *Drive to Eurotunnel Check-In*. The boundary event is interrupting, hence the activity is canceled if a delay occurs. The transport continues towards the ferry terminal and crosses the English Channel over sea. After crossing the channel, the goods are delivered to the recipient.

show a cep query for that scenario to make it more precise

As interpreted from the BPMN specification, the subscription to the boundary event is issued as soon as the related activity is enabled. Given that events arrive every 30 minutes, there can be a gap of up to half an hour, before the first information becomes available. In the worst cases, when data isn't published for several hours, this gap will be even bigger. Let's consider a very busy weekday; A technical fault occurred in the tunnel earlier that day and the train runs 3 hours behind schedule. The last information on the RSS feed was published at 2:35 pm. At that time the truck driver is still in the process of loading goods, finishing the activity at 2:40 pm. Following the process definition, the driver now departs towards the Eurotunnel check-in. The system publishes updated information at 3:15 pm, operations are still 2:30h behind schedule. The message gets received through the pro-

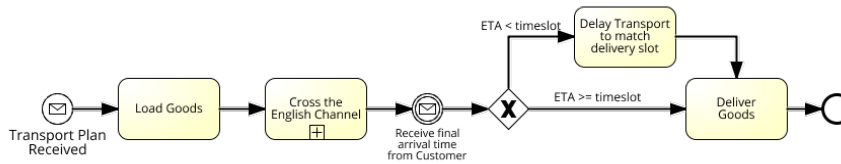


Figure 4: Transport via English Channel that is timed to a delivery slot (Example 1.2)

cess and the truck driver takes the alternative route to the ferry, but only after heading to the Eurotunnel for 35 minutes. The late change of plans causes an unnecessary delay to the shipment.

Figure 4 is an extension of the the transport process. In logistics, it is common that a delivery cannot be accepted at an arbitrary time. Instead, the receiving party assigns delivery windows to the transport company. The transport must arrive during the given time window, otherwise the delivery cannot be completed. After crossing the English Channel, the process model shows the catching of a message event containing the desired final arrival time at the factory. There is an agreement with the factory, that the delivery slots will be approved 2 hours before the expected arrival. If the current ETA of the transport is greater or equal to the arrival time, the driver will head to the drop-off point immediately. If the transport is ahead of schedule, the driver will have to delay the delivery to match the time window.

The presented process model illustrates another complexity of using events in processes. Again, the listening to the announcement of a delivery window will start when the event element is enabled, in this case after crossing the English Channel. Until an event has been received, the process will not continue. Much worse: if the receiving party sends out the arrival time information too early, i.e. while the truck is crossing the channel, the event is missed. If it is not issued again, the process cannot receive a message and will get stuck indefinitely waiting to catch the event.

Neither of the two presented catch events allow for an efficient and reliable execution of the process. They can cause unnecessary delays and even blocking of the process execution.

Consider more possible event occurrence times to prepare for the next chapter

**UP-TO-DATE SHIPPING INFORMATION FOR AN ORDER** A similar situation can be observed in the order-management process depicted in Figure 5. It describes the interaction between customer and seller in a traditional distance retail scenario: After browsing the product catalog, the customer requests a quote for the articles he or she is willing to buy. The retailer makes an offer including an approxima-

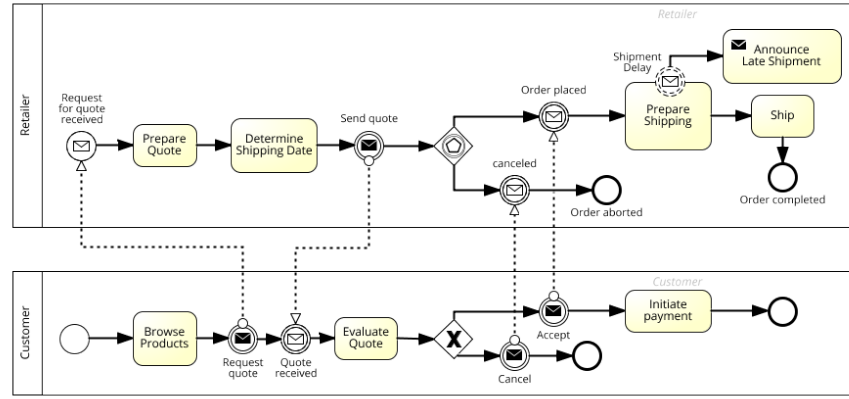


Figure 5: Model of a retail order management process (Example 2)

tion of the expected shipment date and sends it to the customer. That quote is then either accepted or not and the payment is issued if necessary. Once the retailer is informed about the placement of the order, the products are packed and shipped as soon as possible. For articles that are not currently in stock, the retailer must await the shipment from the factory. If any of the factory-shipments is delayed, the retailer cannot ship in time and will announce a delayed shipment date to the customer. This situation is modeled through a non-interrupting boundary event attached to the *Prepare Shipping* activity, which triggers the sending of the updated shipment date to the customer.

The process shows a number of similarities, but also differences in terms of event-use when compared to the logistics process in Example 1. At first we look at the three intermediate catch events, *Quote received*, *Order canceled* and *Order placed*. In each of the cases, the event to be caught is the direct response to a message that was sent right before. While the process will also enter a waiting state until the response arrives, that waiting is not to be interpreted as an unnecessary delay to the process execution. Other than in Example 1.2 (Figure 4), there is nothing useful to do before the response is received. It is furthermore worth noting that the response messages cannot be missed, because the message catch event immediately follows the message send event.

A different situation holds for the boundary event *Shipment Delay*. While the subscription to a Eurotunnel event can be issued at any time, it does not depend on any process data, the shipment delay has to be observed for each product that is part of the order. A subscription can therefore not be executed before the activity *Prepare Quote* has terminated.

show a cep query for that scenario to make it more precise

Conforming to the definition of the process, the system will listen to shipment delays once the activity *Prepare Shipping* is enabled



and therefor much later than possible. Any events that occur after the completion of *Prepare Quote* but before *Prepare Shipping*, cannot be considered in the process execution and the customer will not be informed about a possible late shipment. That will, for instance, concern events that occur while the customer makes the decision about accepting or canceling the order.

The two presented examples have illustrated the complexity of using events in business processes, especially when all possible event occurrence times are taken into consideration. Differences have been pointed out as to how exactly the event is placed in the process, if it waits for a direct response to an earlier request or if the event occurrence is unrelated to the execution of that very process. Motivated by this complexity and the possible implications, it is the goal of this work to evaluate the capabilities of BPMN and to present a concept for the flexible handling of event subscription in business processes.

### 3.2 EVENT OCCURRENCE SCENARIOS

Given the motivating examples, a generic set of *Event Occurrence Scenarios* is defined in this section. Each of the scenarios represents a real-world situation and process implementations need to be capable of handling them to avoid negative effects.

The dominant variable to consider is the event occurrence time. According to the BPMN specification, it is possible to catch an event if it occurs after the event element is enabled. As shown before, it is often impossible to control occurrence time and events do occur outside of the listening time intervals. An Event Occurrence Scenario (EOS) describes the time of event occurrence in relation to a specific step in process or engine execution. The life-cycle of a process within a process engine is explained in [Section 2.1](#) and taken as reference. It is assumed that the process and event engine are configured and running. An event is considered to always occur before or after a life-cycle step or in between two consecutive steps. Given the life cycle steps *process deployment*, *process instantiation* and *event enabling*, the following occurrence scenarios are distinguished in this work:

EOS<sub>1</sub> While the BPMN event element is enabled

EOS<sub>2</sub> The event does not occur

EOS<sub>3</sub> Between process instantiation and  
the enabling of the BPMN event

EOS<sub>4</sub> Between process deployment and  
process instantiation

EOS<sub>5</sub> Before process deployment

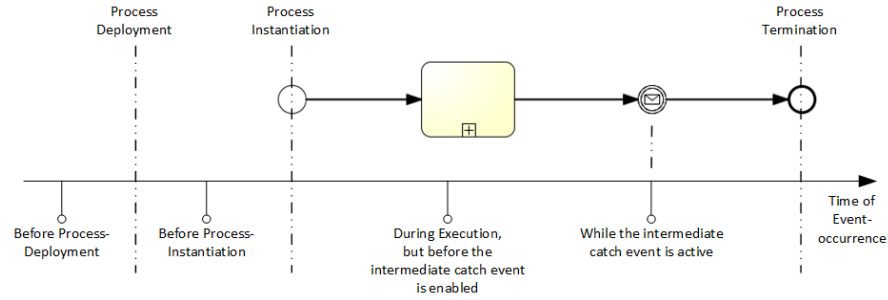


Figure 6: Possible event occurrence times in relation to a process execution life cycle

improve figure: include EOS notions, save space, ?

An overview of the EOSs and the related life-cycle steps is provided in [Figure 6](#), which uses a time-line to illustrate the possible event occurrence times. The model deliberately excludes event occurrences after termination of the event element, because that would presume that there is no more interest in the event as the execution flow continued on another branch. Otherwise, the process will remain in waiting state until the event occurs.

add a back reference to the examples. In example XY, events can occur before... whereas in example...

As introduced in [Section 2.2](#), there is a strict temporal order between event subscription, reception and un-subscription. To catch an event that occurs as described in any of the EOSs, the related subscription operation has to be executed before the start of the time interval associated with the EOS. On that basis, the latest possible event subscription time can be derived from each event occurrence scenario and is as follows: *before process deployment* (EOS<sub>5</sub>), *during process deployment* (EOS<sub>4</sub>), *at process instantiation* (EOS<sub>3</sub>) and *when the BPMN event element is enabled* (EOS<sub>1</sub>).

It is important to highlight that the subscription to an event source can depend on additional context information or process data, which can be a significant limitation to the possible subscription time. That situation is illustrated in *Example 2* ([Figure 5](#)), where the subscription to the shipment delay information cannot be executed before the *Prepare Quote* activity is completed. More generally put, a *subscription dependency* is defined as follows:

**SUBSCRIPTION DEPENDENCY** The subscription operation is *dependent*, if any of its parameters, for instance the event query, contains a variable value that has to be determined at the time of subscription. The value might reference a piece of information from the process instance context, an engine-wide piece of data or external information. In case of a dependence, the sub-

scription cannot be issued before the associated information is available.

If a subscription dependency resolves before the process execution reaches the event element, it should still be possible to flexibly chose the subscription time in the process. For that reason, in addition to the four subscription-time options presented earlier, an additional option *at any time during process execution* must be available.

### 3.3 REQUIREMENTS DEFINITION

The previous sections have exemplified how the BPMN execution semantics and notation capabilities limit users when using events in business processes. Now, these shortcomings are formalized into an initial set of requirements additional to the features offered by BPMN. They build the foundation for enhancing BPM solutions towards flexible event subscription management ([Chapter 5](#)). Before that, the formal requirements will be used to evaluate the capabilities of current Process Management Solutions ([Chapter 4](#)), which results in an additional set of requirements *R4-R6* presented in [Section 4.4](#).

by event I usually mean event from a cep pub/sub event source. The requirements relate to external intermediate message events. A limitation that is made in accordance with the motivated scenario. though some of the aspects might apply to other types of eventy, like the start event.

#### R1: FUNDAMENTAL SUBSCRIPTION INFORMATION

**R1.1 INFORMATION AVAILABILITY:** For each event that is used in a business process, all information necessary to execute an event subscription in the given execution environment must be available. That information generally includes the event query, the address of an event processing platform and auxiliary information to establish the communication, for instance authentication data.

**R1.2 VARIABLES IN EVENT QUERIES** To adapt an event subscription to execution- or time-specific information, variables can be utilized as part of the query string. At the time of event subscription, these variables must be replaced their current values before the subscription is executed.

#### R2: FLEXIBLE EVENT SUBSCRIPTION TIME

**R2.1 EXPLICITNESS:** For each event that is used in a business process, it must be possible to derive the time of event subscription from the process model. The time of subscription may either be explicitly stated or defined implicitly.

**R2.2 FLEXIBILITY:** The time of subscription can be influenced independently from the place the event element takes in the model. Timing options are made available to catch events according to any of the event occurrence scenarios EOS<sub>1</sub>, EOS<sub>3</sub>, EOS<sub>4</sub> and EOS<sub>5</sub>. Consequently, the options must include but are not limited to the subscription before process deployment, during pr. deployment, at process instantiation, subscription at an arbitrary but explicit time during process execution or when the event element is enabled. Thereby, also events that occur after the time of subscription, but before the event element is enabled shall be available to be consumed. (cf. [23])

**R2.3 AWARENESS OF SUBSCRIPTION DEPENDENCIES:** The additional flexibility provided through R2.2 is limited by the use of variables in event queries (R1.2) and the implicit dependencies on context information. The execution environment must only allow a subscription time after the resolving of all dependencies.

Dependencies can exist within a certain process instance, for example when an information from a local DataObject is referenced in a query. In more complex scenarios, the subscription might depend on data from other process instances or even other processes which must first be obtained.

### R3: EVENT BUFFERING

**R3.1 BUFFERING PRINCIPLE:** As introduced by R2.2, an indefinite time can pass between the time of subscription and the consumption of the event. While the subscription operation generally assures that the event is received, it is necessary to temporarily store it until the consuming element is reached.

This temporary storage is referred to as an *Event Buffer*. The storage must be designed to fulfill the semantics demanded by R2.2 which means that any event that occurs after the time of event subscription must be available to consume.

**R3.2 BUFFER SCOPE:** The flexible event subscription time (R2.2) allows an event subscription after process instantiation, but also before process instantiation or deployment. It can be inferred that an event buffer operates either in the scope of a process instance, a process definition, or in the scope of the complete execution environment. To implement all event subscription times stated in R2.2, the execution environment must support event buffering in each of the three scopes. The provided concept must make clear to the user, which scope is referenced at any point in time.

**R3.3 BUFFER POLICIES:** The behavior of a buffer in any occurring situation must be specific. The parameters necessary to specify

the behavior are referred to as *Buffer Policies*, following the notion chosen by Mandal et. al. [23].

- R3.3.1 LIFESPAN POLICY: Given that the time between the event subscription and its consumption may be indefinitely long, process designs can require to limit the maximum time an event is held in the buffer.
- R3.3.2 CONSUMPTION POLICY: Requirement *R3.2* implies that buffers can be shared among process instances of the same process and other processes. In this scenario, it needs to be defined if an information is consumed from the buffer upon retrieval, or if the information remains in the buffer for other participants to access.
- R3.3.3 SIZE POLICY: Provided that multiple events can occur between the time of subscription and the enabling of the event element, it must be specified for each buffer, how many events should be stored for later consumption.
- R3.3.4 RETRIEVAL ORDER POLICY: If multiple events are stored in the buffer (*R3.3.3*), it must be defined, in which order the events are provided for consumption.



## ASSESSMENT OF CURRENT BUSINESS PROCESS MANAGEMENT SOLUTIONS

The lack of flexibility in handling event subscription in business processes has been outlined in the previous chapters, and a set of additional requirements  $R1$  to  $R3$  for process management solutions have been presented. In this section, we take a closer look at the capabilities of current solutions with regards to the event occurrence scenarios (see [Section 3.2](#)) to get a better understanding of the issues that arise when working with event subscription in business processes.

The assessment will be carried out on the basis of the [BPMN](#) and Camunda, a state-of-the-art and widely adopted business process engine ([Section 2.1.2](#)). Both tools shall be used *as they are*, that means without utilizing their comprehensive extension mechanisms. The main goal is to identify and illustrate the shortcomings of the current process technology stack by attempting to implement the requirements identified previously. We therefore first investigate options to handle the [EOSs](#) on the BPMN model level ([Section 4.1](#)) and then proceed to implementing early event subscription and event buffering in a set of Camunda processes ([Section 4.2](#)).

Finally, the findings are discussed in [Section 4.3](#) which leads to the presentation of three additional requirements,  $R4$  to  $R6$ . The extended requirements will be referenced in [Chapter 5](#) to develop a more refined subscription handling model.

which functionality should be evaluated exactly?: all occurrence scenarios, but no buffer policies. The buffer will always store the last version of the event and also deliver that version.

### 4.1 BPMN MODELS IN PRESENCE OF THE EVENT OCCURRENCE SCENARIOS

According to the BPMN specification, the listening for an event starts only when the event element is enabled. The start of the listening is interpreted as the time of event subscription. The motivating examples illustrate that process executions can be delayed or even blocked because of a late time of event subscription ([Section 3.1](#)). The subscription-time is put in relation to the time of event occurrence by the help of the event occurrence scenarios.

In this section, I first describe for each [EOS](#) how a basic event implementation behaves in presence of the given scenario. I then evaluate if it is possible to create a BPMN model that is free from unnecessary delay in these situations. The investigations are carried out on

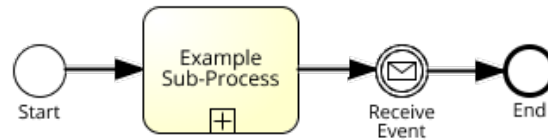


Figure 7: Generic process terminating after an intermediate message catch event

the basis of the abstract process model shown in [Figure 7](#). It uses an intermediate message catch event which follows an arbitrary process flow, represented through a collapsed sub-process. The process terminates after the event is received. It is assumed, that the event can occur independently from the preceding subprocess and therefore in accordance with any of the event occurrence scenarios. That enables us to investigate each EOS using this simple model. The event behavior is comparable with the *Eurotunnel Delay Event* considered in *Example 1.1*. Depending on the discussed EOS, the motivating examples will be used for further illustration of the situations.

*EOS<sub>1</sub>: The event occurs while the catch element is enabled*

The first scenario represents the case that is assumed by the BPMN 2.0 specification: The event occurs after the event element has been enabled and before it has terminated. It will be consumed by the catch event and the process flow can proceed without unnecessary delay. In this scenario, the use of a simple intermediate catch event is sufficient.

Notably, there always is a certain delay when consuming an event through an intermediate catch event. The BPMN specification itself states that *the handling consists of waiting for the Event to occur*. However, that delay can theoretically be infinitely small. <— maybe mention this at the beginning of ‘problem statement’ to differentiate between normal delay and unnecessary delay?

*EOS<sub>2</sub>: The event does not occur*

In certain situations an event might not occur at all. Given a basic event implementation like in [Figure 7](#), the process flow will get to a halt once it reaches the intermediate catch event and will not be able to proceed. While, depending on the process design, this might be the desired behavior, in many situations this is not acceptable. In *Example 1.2*, the process relies on the *final arrival time* information from the customer. If the event does not occur, for example due to a mistake in a factory work-flow, the process execution waits for the event indefi-



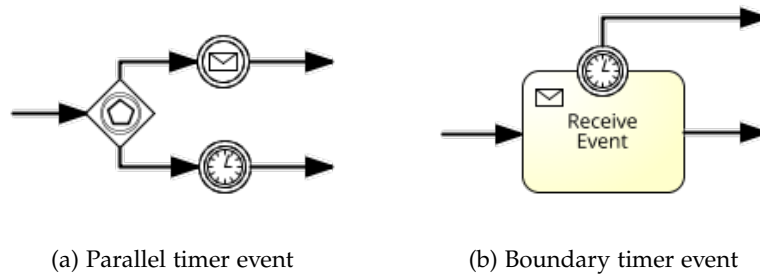


Figure 8: Receipt of an intermediate message event, either interrupted by a boundary event (a) or parallel timer event (b).

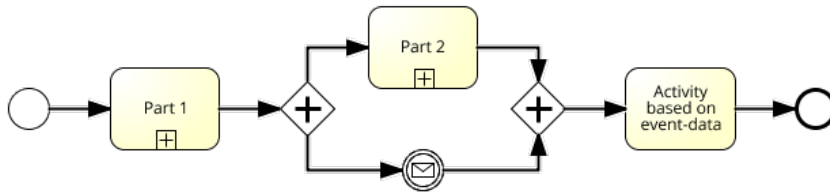


Figure 9: Event Element in parallel process flow

nately. To improve the process design, the logistics company decides that after waiting for the event for 30 minutes, an arbitrary delivery window shall be assumed. The waiting for the event shall be terminated so that the goods can be delivered immediately.

Figure 8 shows how this behavior can be implemented in two ways: (a) By the help of an event-based gateway which puts a timer event in parallel to the intermediate catch event: If the timer event occurs before the message event is consumed, the message event is terminated immediately and the process flow proceeds from the timer event. An alternative solution with equivalent semantics is depicted in (b), using a receive task and an interrupting boundary timer event. Once the timer event fires, the receive activity is canceled and the process continues along the outgoing flow of the time event. Any of the two model improvements will make sure that a process does not run into a lock if the expected event does not occur.

*EOS3: Occurrence between Process instantiation and the enabling of the BPMN event*

In case the event occurs during process execution, but before the BPMN event element is enabled and thus listening for events, the occurrence will not be available for consumption. The execution will come to a halt at the event element as if the event did not happen at all.

could exemplify this at ex1.2, In the previously mentioned *Example 1.2*

To avoid a lock in this scenario, the intermediate catch event can be placed in parallel to the rest of the process flow using a parallel gateway. This is illustrated in [Figure 9](#). The time of subscription to the event can be controlled by the position of the parallel split: To implement an event subscription right after process instantiation, the Parallel Gateway has to be the first element after the Start Event (that means *Part 1* in the illustration is empty). To implement the event subscription at a specific point during process execution, part of the process must execute before reaching the Parallel Gateway. As modeled in [Figure 9](#), the event may occur at any time during the execution of the collapsed sub-process *Part 2*.

#### *EOS<sub>4</sub> and EOS<sub>5</sub>: Before Process Instantiation*

Any Events that happen before process instantiation will not be considered in a standard intermediate catch event. That applies to both scenarios, the occurrence between deployment and instantiation (*EOS<sub>4</sub>*) and an occurrence time before the deployment of the process in the Process Engine (*EOS<sub>5</sub>*).

To create a process model that allows to catch an event before the process instance exists, three new elements are introduced in addition to the existing process referred to as *target-* or *original process*. The elements are: (1) An additional *Auxiliary Buffering Process* that can catch an incoming event independently from the target process; (2) an *Event Buffer*, a temporary data-store that keeps event data until it is ready for consumption; (3) an *Auxiliary Event Delivery Process*, that retrieves events from the buffer and makes them available to the target process. By introducing an additional process for listening to the event, the subscription time to the *Event Engine* can be before instantiation or even deployment of the original process. [Figure 10](#) shows the interaction of the original process, the two auxiliary processes and the data-store.

**THE EXECUTION FLOW** To start listening for an event, the auxiliary buffering process has to be instantiated through a message start event containing the information necessary for the event subscription. The process subscribes to the event engine through a throwing event. It then waits for an event to occur in a receive activity *Receive external event* and outputs the received event data to a *DataObject*. That object is then written to a persistent, global data-store *Event Buffer* by the service task *Write to buffer*. The design is able to handle multiple event occurrences, because, after writing to the buffer, the receiving activity is executed again. The buffering process terminates only once the *Unsubscribe* event is received. The implementation of the service task

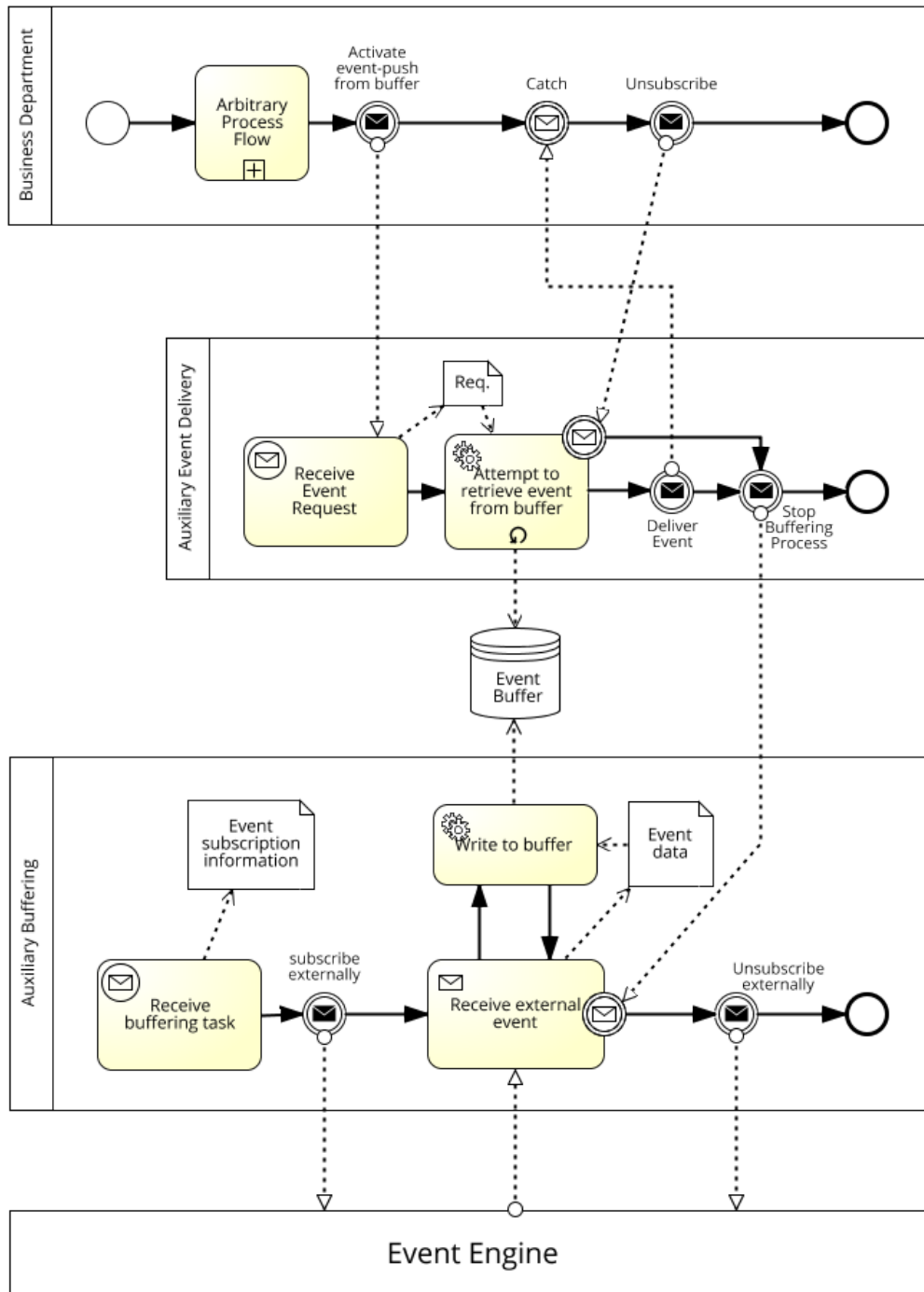


Figure 10: Event Buffering through an auxiliary buffering process

decides about the exact semantics of the buffering and hence about requirement *R3.3*. For the sake of simplicity, it is assumed that the buffer can hold a maximum of 1 element (*Size Policy*) and the service task overwrites that element whenever a new one is available, specifying the retrieval order (*R3.3.4*). An implementation of the *Lifetime Policy* is not provided. The consumption behavior (*R3.3.2*) is defined by the service task of the auxiliary delivery process: After the retrieval of an event from the buffer it remains available.

After the buffering process is running, the original process can be instantiated. In the presented model, its intermediate catch event has been explicitly split into three events to fulfill the publish/subscribe paradigm: An initial send event to request events, a catch event to receive and a final send event to signal that no events shall be received anymore. The initial send event instantiates the *Auxiliary Event Delivery Process*, which tries to read from the event buffer and delivers the event to the original process if there is one available. The central looping activity will retry reading from the buffer until data becomes available and will only be terminated once the *Unsubscribe*-event occurs.

**CONCLUSIONS** Thanks to the complex interaction of the three processes, the original process can consume the desired event, even though the event occurs before process instantiation. Consequently, *EOS4* can be handled by the model. Moreover, the *Auxiliary Buffering Process* is not bound to a specific event, it works generically with any event information that is passed to it. For that reason it is also not bound to a specific process deployment and can buffer events even before a process has been deployed, so it handles scenario *EOS5*. The buffering process can alternatively be started using an explicit message send event during process execution, similar to the *Explicit Subscription Task* introduced in [23]. By that means, *EOS1* and *EOS3* are also supported, which results in the fulfillment of requirement *R2.2*.

It remains to be highlighted that this approach will only work for *EOS4* and *EOS5* if the buffering process is instantiated before the associated event occurs. As there are no other systems in place to automatically instantiate the process, it must be assumed that the process is manually started for each event and each process model that requires an early event subscription. Furthermore, the number of processes running in the execution environment increases significantly. An instance of the buffering process is required for each event element that ought to make use of an event subscription time before deployment. For each instance of the original process and each buffered event, an instance of the auxiliary event delivery process must be started. That puts additional load on the process engine which can endanger the reliability of business-critical processes and makes the management

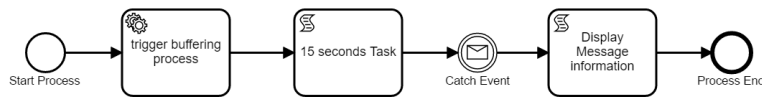


Figure 11: Generic Example Process in Camunda for Occurrence Scenario *EOS<sub>3</sub>*

of processes and their instances more complex. The downsides of the presented approaches are discussed in more detail in [Section 4.3](#).

Do I want to work with the requirements or not???

#### 4.2 IMPLEMENTATION OF EARLY EVENT SUBSCRIPTION USING STANDARD CAMUNDA

The previous chapter has shown that it is possible to create BPMN models to match each of the Event Occurrence Scenarios, though for the scenarios *EOS<sub>4</sub>* and *EOS<sub>5</sub>* the solution becomes increasingly complex. In the next step I investigate the capabilities of Camunda as an example for a state-of-the-art business process engine. Camunda shall be used without any code customization, that means as offered through the deployment package. The solution presented in [Section 4.1](#) (*EOS<sub>4</sub>* and *EOS<sub>5</sub>*) has proven capable enough to handle all event occurrence scenarios, it is therefor the goal to implement a prototypical solution in Camunda in strong accordance with the BPMN model shown in [Figure 10](#).

Two generic sample processes have been modeled for evaluate the functionality of the system. [Figure 11](#) shows a simple process with an explicit subscription activity to represent the listening to the event after process instantiation but before reaching the Catch Event (Scenario *EOS<sub>3</sub>*). It follows a sample activity that takes 15 seconds (implemented using a *Script Task*), the intermediate catch event and another script task that displays the content of the received message. The example for scenarios *EOS<sub>4</sub>* and *EOS<sub>5</sub>* ([Figure 12](#)) comprises the following elements: After the start event follows an intermediate catch event, then an activity that prints the message of the event to console and last the process end event. Both figures show screen captures from the Camunda Modeler.

**AUXILIARY BUFFERING PROCESS** The task of this process is to subscribe to a CEP Platform using a provided event query and start listening for events. Any incoming event must be stored in a data-store (*Event Buffer*). A local MySQL database has been chosen for persisting

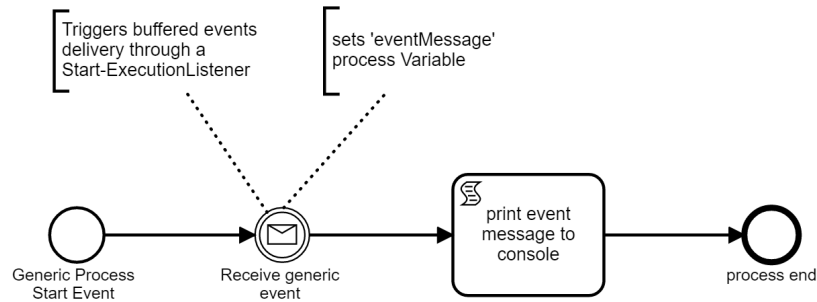


Figure 12: Generic Example Process in Camunda for Occurrence Scenarios *EOS4* and *EOS5*

the event data because it is freely available, easy to set up, offers standardized access via SQL queries and Java connectors. As complex event processing functionality itself is not required to demonstrate the use-case, the role of the *Event Engine* is taken by a basic web-service which was specifically implemented in Python. It exposes a *subscribe* method via [HTTP](#) and can be used to trigger event messages associated to registered subscriptions.

Figure 13 shows the final Buffering Process modeled in the Camunda Process Modeler. The process can be instantiated by issuing a *Buffering Task* message. This message must contain three data fields: *processDefinitionId*, to know which process definition the buffered messages belong to; *messageName*, the name of the message event within the process; *query*, the event query in the Esper Query Language. Camunda will make the message data automatically available in the process instance as process variables, so they can be used during the execution of the Buffering Process. After instantiation, the process reaches the activity *Subscribe to Event Source*, a *Java Service Task* that executes a HTTP call to the event engine. That call registers the event query in the platform, providing the process instance identifier and the message name. Based on this information, Camunda can correlate received events back to a specific process instance and message element. After the subscription, the process reaches the receiving activity *Wait for unsubscribe event* that will terminate the process as soon as the *Unsubscribe* event has been received. As long as this activity is active, events can be received through the attached non-interrupting boundary event. Incoming events have a field *eventBody*, which contains the event information and becomes available through a process variable with the same name. The boundary event triggers the service task *Write eventBody to datastore*, which takes the data from the process variable and writes it to the MySQL Database (*Global Event Buffer*).

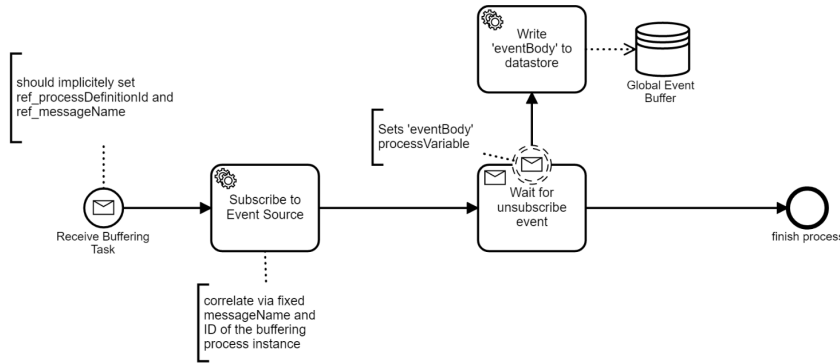


Figure 13: Auxiliary Buffering Process in the Camunda Modeler

Notably, the Camunda process model appears less complex than the one presented in [Section 4.1](#) while still implementing the same functionality. That is possible because in Camunda, the information received in intermediate message events is written to process variables for further use, a feature that the BPMN specification only considers for the receive task.

**AUXILIARY EVENT DELIVERY PROCESS** The delivery process (see [Figure 14](#)) reads the latest data from the buffer and sends it to the process instance. It can be started with a message that contains the *processInstanceId* and the *processDefinitionId* of the requesting process and the *messageName* of the message event that is requested from the buffer. A timer event *Delay Timer* has been inserted to make sure that the requesting process is already listening for an event, when the delivery process sends the message. The outgoing execution flow is treated asynchronously. It follows the service task *Retrieve event from buffer*, which executes Java code to read from the MySQL Database *Event Buffer* and store the event information in a process variable named *eventMessage*. The content of that process variable is sent to the original process in the send event, afterwards the execution is finished.

**INTERACTION OF THE PROCESSES** To initiate the subscription at the event source, the Auxiliary Event Buffering process has to be started. For scenario *EOS3*, this happens through an extra activity (*Trigger Buffering Process*) during process execution, so that events after process instantiation are received by the Buffering Process. In scenarios *EOS4* and *EOS5*, the subscription and thus the instantiation of the Buffering Process must happen before the instantiation of the Target Process. As there is no such mechanism in the standard Camunda Process Engine, the Buffering Process must be started manually, providing the *processDefinitionId*, the *messageName* and the *eventQuery*.



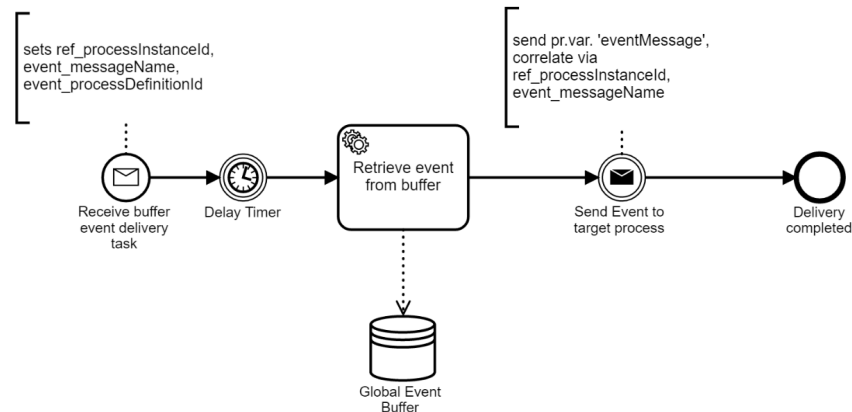


Figure 14: Auxiliary Event Delivery Process in Camunda Modeler

Now that the *Buffering Process* is running, any events matching the query will be stored to the buffer. When the Target Process reaches the Catch Event, a request for buffered events is sent as a message to trigger the *Auxiliary Event Delivery Process*. This message is sent using a short piece of Java code that gets executed when the Catch Event is reached. The code is invoked by a Start ExecutionListener attached to the Catch Event. ExecutionListeners are offered by Camunda to execute own Java programs before or after relevant events during process execution, like the execution of an element in the process. While the Original Process will now start listening for the desired events, the Event Delivery Process will send the buffered events as messages to the Original Process.

If no events have been received yet, all the involved processes remain active: the Buffering Process will keep listening for an external event. The Delivery Process will send an event to the Original Process as soon as there is one in the buffer. The Catch Event in the Original Process will keep listening for an Event. The necessary communication for the termination of the processes and the unsubscribe-operation have not been implemented in this prototypical implementation, as they would not add particular value to the evaluation.

**CONCLUSION** The described implementation of the BPMN buffering concept presented in [Section 4.1](#) serves as a brief evaluation of the model. As illustrated by the help of the two examples, the resulting set of process applications allows to issue an event subscription flexibly according to the event occurrence scenarios. While a thorough analysis was not the target of this section, it still provides an introduction to the capabilities of the Camunda process engine, which will take an important role again in [7](#).

where is the code?



Note that this is an investigative implementation that matches exactly the given use-case and is not meant to be used in production. It is neither flexible nor robust enough for that purpose, but suits very well in understanding the capabilities and the shortcomings of BPMN and Camunda when it comes to handling the Event Occurrence Scenarios.

#### 4.3 DISCUSSION

The goal of this chapter was to get a better understanding of the capabilities of the tools when it comes to covering all event occurrence scenarios. Even though it has proven possible to implement a flexible event subscription time using standard BPMN 2.0 and Camunda, the success comes at a cost. The downsides of the presented approach are presented in the following.

It was necessary to create two generic auxiliary processes for event buffering, to connect to a MySQL data-store and use ExecutionListeners to execute custom Java code in Camunda to cover all scenarios, *O1* to *O5*.

give them short names for reference? Or make one of them Requirement 5

##### MISSING AUTOMATIC SUBSCRIPTION HANDLING

In the presented process models, separate process elements had to be added to handle event subscription and initiate event delivery. That conflicts with requirement *R2*, which states that the subscription and un-subscription must be automatically handled by the process engine. For the scenarios *O4* and *O5* the Buffering Process has to be triggered manually, because it must be executed before the target process is running. Camunda does not handle external event subscription itself, especially not before the process is running.

##### ADDITIONAL MODEL COMPLEXITY

As additional process elements have to be added to handle event subscription and delivery, the models become more complicated and are less concentrated on the business case.

there could be a requirement that states that there should be no additional process elements unless there is an explicit subscription time

##### BUFFERING IS AN IT TASK

The auxiliary processes are not business tasks and are thus not suited to be modeled in BPMN. Desired functionality can be put into Camunda BPMN models thanks to its flexibility to use Java code in

Service Tasks or Event Listeners, but naturally the full functionality of the Event Buffer cannot be expressed using BPMN.

what exactly is the issue here?

#### ADDED LOAD ON THE PROCESS ENGINE

Because of the aux processes, two additional processes have to be deployed in the process engine and are potentially running in parallel to any given process instance. For each Event Element used in a process the engine has to run an instance of the Buffering Process and, eventually, an instance of the Buffer Delivery Process. That puts additional load on the process engine, which might prevent business critical processes from executing delay-free.

Even when the number of deployed and running auxiliary processes can be reduced through further optimizations there remains an event-management overhead as every event has to be handled twice: once when it is stored in the buffer and once when it's delivered to the target process.

#### HIDDEN PERFORMANCE LIMITATIONS OF THE PROCESS ENGINE

Given the large amount and high frequency in that events can occur in reality, optimal performance is required for an event-buffering module. Running essential parts of the buffering within the process engine might pose performance limitations that cannot be influenced without tempering with the process engine code.

write connector to next chapter: it has been shown that... therefor we do ...

### 4.4 REQUIREMENTS EXTENSION

**R2 AUTOMATIC SUBSCRIPTION HANDLING:** The subscription to event sources is handled implicitly by the process execution environment based on the modeled subscription information as required by *R1* and *R2*. Similarly, the removal of a subscription is performed as soon as a subscription becomes unnecessary.

+ model complexity + no influence on performance

Introduce a new concept. three pillars: bpmn-x, process engine behavior, event handling api

### 5.1 BPMN EXTENSION

cite braun2015behind

this is the explanation for the process designer? or is there an additional event engineer?

Given the additional requirements and the shortcomings identified in the previous sections, the following two chapters present an extension to the BPMN event handling model. At first, an extension to the Business Process Model and Notation (BPMN) is described, which aims at providing the Process Designer with more flexible Event Handling capabilities according to Requirement *R1*. Afterwards, [Chapter 6](#) clarifies the changes necessary to the event handling platform and the process engine to cover Requirements *R2* and *R3*. While the presented concepts are kept as general as possible, they are grounded in an analysis of the Esper-based CEP Platform Unicorn and the open-source process engine Camunda.

The extensibility mechanism allows to extend standard BPMN elements with additional attributes, while maintaining a valid BPMN core. Extensions are specified through an external definitions file and can be included into a BPMN process model by reference. To allow

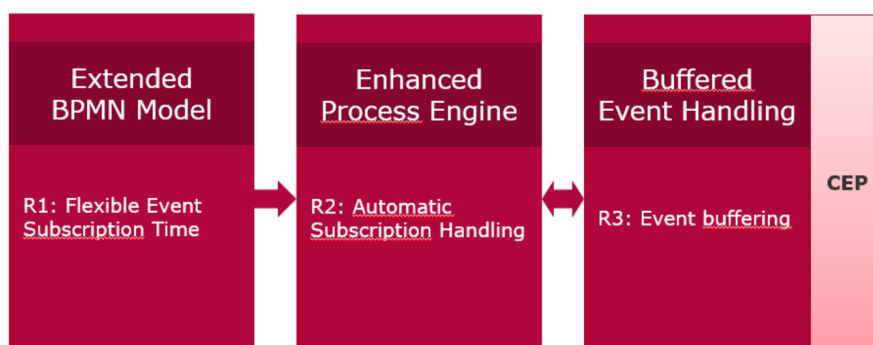
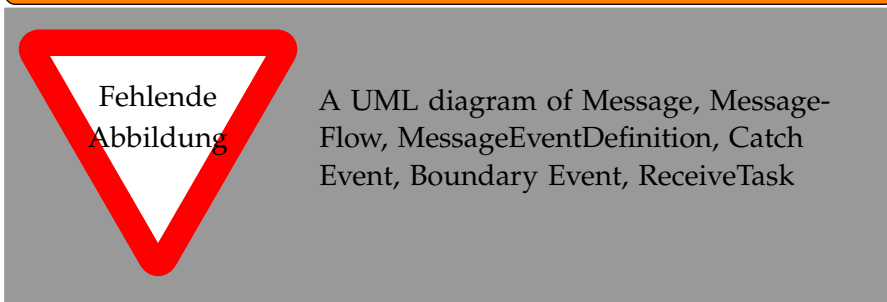


Figure 15: The concept for flexible event subscription involves three modules: A BPMN extension, enhanced process engine behavior and buffered event handling.

the flexible use of event subscription in BPMN models, a number of additional attributes must be added to BPMN process models.

The presented extension is designed for use with intermediate catch events, boundary events and receive tasks. Start events and End events are not considered because flexible subscription time is not relevant in these cases. Each of the three elements references a BPMN *Message*, the common denominator for communication within and across business processes. Semantically, the *Message* type is most suited to be extended with subscription information.

mention that others have extended the events instead, but this proposal goes another way



explain that the Message is a common element in BPMN models. It is the main generic type used for communication, collaboration. The proposed extension to the Message type is only to be used by Message Receive situations in Intermediate Catch Event, the Boundary Catch Event and the Receive Task.

By specification, *tMessage* comprises an attribute *name*, the name of the message, and *itemRef*, the reference to a BPMN *ItemDefinition*. Additionally, it inherits all attributes from the BPMN *RootElement*. In the following, the required additional attributes will be explained one after the other. A complete list is available in [Table 1](#). The goal is to retain a stand-alone model that contains all information necessary to execute the subscription to the event source.

#### 5.1.1 Adding basic subscription information

For a basic event subscription, an event query, the platform address and optionally authorization information of the CEP Platform is required.

ref background

It is assumed that only one CEP platform is in use, whose access information is configured centrally for the current process execution environment. Consequently, there is no need to specify these two parameters on message level. The event query instead needs to be specified for every message and is added to the model as an extension at-

ATTRIBUTE NAME	VALUE OPTIONS ( <u>DEFAULT</u> )	OPTIONAL
eventQuery	any string	n
subscriptionTime	process-deployment, process-instantiation, <u>event-reached</u>	y
bufferPolicies	Complex Type (see below)	y
<b>BUFFERPOLICIES</b>		
LifetimePolicy	string in ISO time-span format OR ' <u>infinite</u> '	y
ConsumptionPolicy	<u>Reuse</u> , Bounded-Reuse(n), Consume	y
SizePolicy	int (< 1 for infinite), <u>1</u>	y
OrderPolicy	<u>FIFO</u> , LIFO	y

Table 1: Available attributes in the BPMN extension for flexible event subscription

tribute *eventQuery* of type *String*, which should contain the full query as interpretable by the CEP platform.

A similar approach has been taken by X and Y, who aim at enriching BPMN models with subscription information without considering the time of subscription specifically.

Find this source; explain what they do (different)

Given this first fundamental part of the BPMN extension, it is possible to execute the subscription, but the time of subscription cannot be influenced.

### 5.1.2 The time of event subscription modeled in BPMN

This section specifically addresses the requirement *R1.1*, aiming to provide a flexible event subscription time to be selected for each BPMN message when designing an event-driven process. Two different tools are to be offered to support all subscription times demanded by *R1.1*: Firstly, the subscription can happen in the background. Alternatively, the subscription can be modeled explicitly as a flow-element in the process. It is the task of the process designer to elaborate the correct time of subscription necessary for her use case.

The subscription will be executed automatically by the system based on the information given in the BPMN model. Further information on the exact execution flow is provided in chapter XY.

**INTRODUCING EVENT BUFFERS** Any event message that occurs before reaching the event element but after the time of subscription will be kept in a buffer by the system. In its simplest version, the buffer is of length 1, that means it stores exactly one message received from a CEP platform. It always stores the latest message. When a newer message arrives, the old one is replaced in the buffer. [Section 5.1.4](#) introduces a set of advanced buffer policies to adapt this behavior further.

By default, there is no interference between the buffers of different messages, process instances or processes. Each buffer instance will contain the latest information as if it was the only buffer in the system. Performance improvements to avoid duplicate buffer content will be managed by the system without explicit action by the user. Section ... later introduces a shared, more complex usage scenario of the event buffers.

**INTERPLAY OF EVENT QUERIES AND BUFFERS** Modern event query languages are feature-rich and offer a large set of expressions to filter events from incoming streams.

ref background

The introduced basic event buffer can be used in connection with any desired event query and will store the latest output of that query. These two features together suffice to implement even more complex use-cases: Query windows of length  $n$  can be used to keep multiple events in the buffer, filter expressions allow to keep a subset of all events based on their attribute values, multiple streams can be joined together. As soon as the process flow reaches an event element, the latest CEP message is retrieved from the buffer. It is not consumed, that means a second event element that references the same BPMN Message will reuse the information from the buffer. If no information is available in the buffer, the flow element will remain in the waiting state until a message is received. Then, the process flow proceeds as usual.

reorganize this chapter

#### SUBSCRIPTION TIME AS PART OF THE BPMN MESSAGE ELEMENT

To provide the Process Designer with a simple but powerful tool to influence the time of event subscription, a field *subscriptionTime* is added to the BPMN message element. The field can take one of the following three values: *Process Deployment*, *Process Instantiation*, *Event reached*. The last option is the default option, coming closest to the BPMN specification. Note that a *subscriptionTime* set to *Event reached* will remain without effect if an explicit subscription task for the same event was executed before the event is reached.

For each of the options: Define exactly (according to BPMN spec or standard literature), when in the flow the subscription is executed.

In motivating Example Ex, it is necessary to issue the subscription as early as possible, to make sure that data is available and the process execution is not delayed.

which example to reference?

Using the BPMN extension, the use case can be implemented by defining the event query and the subscription time in the BPMN model.

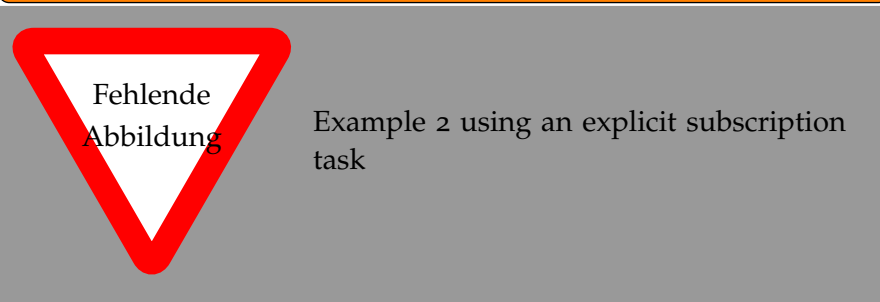
show what that would look like in the example. Maybe some XML?

#### THE EXPLICIT SUBSCRIPTION TASK

As an alternative to specifying the subscription time using the extension field *subscriptionTime* of *tMessage*, an extension to the BPMN ServiceTask is proposed.

it follows the ideas presented in ... but through referencing the message element

latest concept (xsd) states that it's an extension of tTask



The extended task is used to execute the subscription explicitly as part of the process flow. A field *messageId* is added to the service task to establish a reference between the activity and the message definition. As introduced in section [Section 5.1.1](#), the extended BPMN Message definition contains the information necessary to issue the subscription to an event source. Once the Explicit Subscription Task is activated, the subscription for the referenced message is issued.

Modeling the event subscription in an explicit task can be necessary when the subscription depends on the result of another activity. In that case, the subscription cannot be issued on process instantiation, because the necessary information is not yet available. Instead, an early subscription can be implemented using the extended service task. Apart from this particular use case, the explicit subscription task enables the Process Designer to place the subscription flexibly in the process flow and give her full control over the time of subscription.

As an improvement to the options for subscription time, there could be an option "ASAP", so that the process engine issues the subscription automatically as soon as the required process data becomes available

If both tools, the extension field *subscriptionTime* and the explicit subscription task, are used for a single BPMN message, the earlier subscription of the two will be executed, the second subscription will have no effect. That means for example if the *subscriptionTime* is set to *event reached* and an explicit subscription task is inserted before the event element, then the subscription will be executed at the time the explicit subscription task is active. If *subscriptionTime* is set to *Process Deployment*, then the subscription will happen at that time and the explicit subscription task will remain without effect. In case neither of the two is used, the system falls back to the BPMN default and executes the subscription when the event element is reached.

### 5.1.3 Using Process Variables in Event Queries

this could be added in the requirements and referenced

only if *subscriptionTime* not 'on depl'

As shown in example Z, it can be the case that the values of process variables shall be dynamically used in an event query. Therefore, the name of the process variable should be part of the event query. At the time of subscription, the mentioned variable is dynamically replaced by its current value. The exact notation for including process variables in event queries can vary depending on the applied query language as it may not interfere with any existing notation schemes. For the use with the Esper query language, the following is suggested: The exact name of the variable has to be surrounded by curly brackets and preceded by a # character: *#{VARIABLENAME}*. This notation is inspired by the usage of substitution parameters in SQL queries that are embedded in Esper. They take the form *#{expression}*.

reference esper docs 5.13.1. Joining SQL Query Results

In the example, the process uses the latest GPS position for a certain truck. The truck is identified by its unique ID which is part of the query: *SELECT lat, lng from GPSUPDATE where truckid = #{truckid}* .

missingref: dependent example

The use of dynamic process variable values introduces an additional complexity: Depending on the time of event subscription, the value of the process variable might not yet be available.

What does this mean for the process designer? A model that can take a state in that a subscription shall be issued, though the data is unavailable, is invalid. When will an error occur?



#### 5.1.4 Advanced Buffer Parameters

Inspired by Mandal, Weidlich, and Weske [23], a number of advanced buffer attributes are available through an extension attribute *buffer-Policies*.

add ref to sankalitas paper; do I elaborate on the changes made in comparison to the paper? + the reasons?

##### LIFE-TIME OF BUFFERED EVENTS

The *LifetimePolicy* allows to specify after which timespan elements in the buffer should be deleted. Timespans shall be defined using ISO timespan format.

ref external

The default value is *infinite*. Example A has been implemented by setting the *subscriptionTime* to *Process Deployment*, which means that there can be an infinite time difference between the action of subscribing to the event source and the reaching of the event element in one of the instances.

ref example

In case events are not published in a longer time, for example due to technical fault at the event producer, the buffer will contain older events that might not be relevant anymore. Using the *LifetimePolicy*, the process designer can express, that events should be deleted from the buffer after a certain period of time and thus avoid outdated information. The buffer is maintained automatically by the system. That of course comes at the price that the process has to remain in waiting state until a new event message arrives.

##### CONSUMPTION BEHAVIOR

for the reader it would not be clear, what a 'buffer instance' is

So far, the event buffers can be used isolated from each other. There is no interference between buffer instances and events are not removed from the buffer after retrieval. While for most use-cases this behavior is sufficient, more detailed control over the buffer can be desirable when a given message shall be used multiple times. Not always is it wanted, that events remain in the buffer after retrieval. An additional parameter *ConsumptionPolicy* is introduced which can take the values *Consume*, *Reuse*(default) and *Bounded Reuse*(*n*). While *Reuse* denotes the behavior that is already known, *Bounded Reuse*(*n*) will allow an element to be retrieved exactly *n* times. *n* has to be replaced by an integer value greater 0. The option *Consume* will remove an element from the buffer immediately after it has been retrieved for the first time, it is therefor equivalent to *Bounded Reuse*(1).

- given the option to consume from the buffer, it will now make a difference if the same buffer is accessed multiple times. - there are

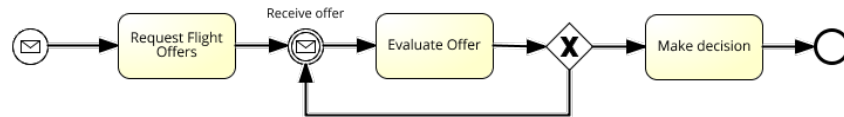


Figure 16: Flight Booking process using a consuming buffer

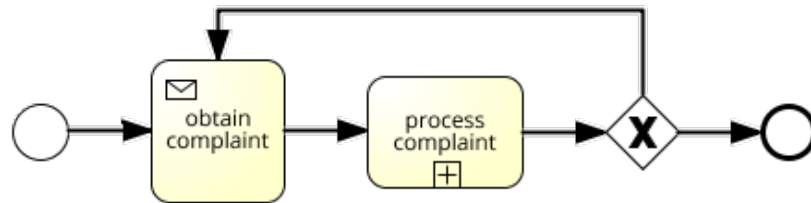


Figure 17: Shared consuming buffer in Complaints Handling

two scenarios to access the same buffer: (1) multiple times in the same instance, (2) multiple times because of parallel instances, (3) multiple times because of a shared buffer across processes - before proceeding, we need to be clear about the buffer scope: it depends on the time of subscription. (1) after instantiation: buffer only instance-wide; (2) on pr depl: buffer reused across all instances; (3) system start: buffer reused across all processes

#### BUFFER SIZE AND ORDER POLICY

- two additional buffer policies: *SizePolicy*, *OrderPolicy*, default values

write text for given keypoints

- what if messages are defined in different ways? => I want to reuse the buffer between processes, but the message definition (especially cep query) are not the same. how will the system behave?

## 5.2 DESIGN DECISIONS

The target functionality of the BPMN extension was clearly defined by the identified requirements. To implement that functionality, there were a number of options to consider and design decisions to make. This chapter provides background information on the decisions that have influenced the presented concept for flexible event subscription.

Talk about an alternative solution? Table of events for a certain event source. The buffer is already there for me to pick from when designing the process.

THE TIME OF SUBSCRIPTION IS A QUESTION OF PROCESS DESIGN  
 -> process designer - this is why the bpmn extension is presented first  
 - hide complexity from the designer

THE ACTUAL BUFFER IS MOSTLY HIDDEN FROM THE USER - Buffers are implicitly defined through the BPMN model - keep the look and feel of the message catch event - minimal changes to existing models, backwards compatibility

AVOID ADDITIONAL USER INTERFACES - the user will only use the bpmn model - we don't want any other element because of complexity - the process should be self-contained, contain all necessary information for subscription and buffering. -> single point of contact

BUFFERS ARE CLOSELY LINKED TO PROCESS MODELS Messages are only buffered as soon as they are explicitly required by a model - we don't just buffer n messages because we might need them in the future. That would be a fuzzy, incalculable performance overhead. - instead we keep as little as possible in the buffer

THE BPMN EXTENSION IS BASED ON THE MESSAGE ELEMENT - if I want to talk about related work that goes another way - why are the policies a parameter of the message and not the catch event element?

REUSE EXISTING TECHNOLOGY -> we assume that a cep is present and that basic features of event queries can be used - if not present, then only very basic buffer functionality is available - but we don't want to start designing another event processing layer with duplicated functionality



## AUTOMATIC SUBSCRIPTION HANDLING

---

After defining the notation options and functionality provided to the user of flexible event subscription, this chapter describes the changes necessary to the software infrastructure that is used for event-driven business process management. The concept requires that all subscription and event handling is executed by the system itself, without further interaction by the user. All necessary information for that purpose is provided by the BPMN model.

As described in REF, an event-driven process management setup primarily consists of ...

missingref

Changes are necessary to both, the Event Processing Module and the Process Engine. This chapter attempts to keep the change descriptions general so they can be applied to any common process engine and event processing platform. The first section describes the necessary extension to the event processing to support early subscription and event buffering. The following [Section 6.2](#) specifies the changes necessary to the behavior of the process engine as the connecting element between the BPMN model and the event processing platform.

### 6.1 BUFFERED EVENT PROCESSING

When reduced to the basics, a standard event processing platform works as follows: The user subscribes to events providing an event query and a notification-path. The platform responds with a unique identifier for that subscription. Whenever an event occurs that matches the provided query, the platform issues a notification to the notification-path. Subscriptions can be deleted through their unique identifier. These two operations, *subscribe* and *unsubscribe*, make the fundamental API of a CEP platform.

ref

EVALUATION OF COMMON CEP PLATFORMS - looking at the new bpmn extension reveals that a different behavior is required: notifications need to be kept in a buffer until they are requested by an entity - three event platforms were studied to check if the required functionality can be implemented natively: wso2, esper and ... - wso2: ? - esper: - ? : ? - a window will achieve something similar, but not exactly the same - moreover, it is desired that the same full-featured event queries as before can be used, no restrictions. - as functional-

ity wasn't available in any of the three, it was decided to specify an extended api.

Maybe: we consider these two operations as given, because common event processing platforms have these in common, but there is not common event buffering concept yet. even though e.g. es-per has something which goes in that direction: output clauses

**AN API FOR BUFFERED EVENT PROCESSING** The novel BPMN extension for flexible event subscription allows to issue a subscription for buffering well before the events ought to be delivered via the notification-path. The introduction of an event buffer as a separate entity between the varying list of notification-recipients and the event query makes an extension of the API necessary.

note that splitting this into four actions is not a new concept. actually thats how event processing platforms normally operate, only without creation of the buffer and without delivering the buffered events

what kind of API should this be? REST? Java? none specifically, but the reader might want clarification when reading this.

Firstly, the *subscribe* operation has to be divided into two steps:

A. *registerQuery(queryString[, bufferPolicies]): queryId*

The call registers an event query in the CEP platform and instantiates a buffer. Matching events will be held in the buffer according to the specified policies. It returns a unique identifier to that new query registration and hence for the connected buffer. That identifier must be used to modify the query later.

*bufferPolicies* is an optional parameter which is provided as an object with four possible fields: *LifetimePolicy*, *ConsumptionPolicy*, *SizePolicy*, *OrderPolicy*. Refer to [Section 5.1.4](#) for a detailed specification of the semantics of the parameters. If *bufferPolicies* is not or only partly specified, the system should fall back to the default values.

B. *subscribe(queryId, notificationPath): subscriptionId*

Initiates the delivery of notifications for a given *queryId* to a notification recipient. The recipient is specified through the *notificationPath*, the full address of the entity that is supposed to receive the message. Notifications are delivered asynchronously as soon as they are available. If the buffer is not empty, a message will be sent right after the *requestEvents* call. A similar operation, *addNotificationRecipient*, is available in existing CEP platforms. The difference is in the delivery of the first buffered message: *requestEvents* sends out the message from the buffer, *addNotificationRecipient* will send out notifications only for future query output.

be clearer about notificationPath. Specify in background and reference. also about addNotificationRecipient

add table with buffer policies, their possible values and the default val

A similar situation holds for the un-subscribe operation: Traditionally, a subscription is canceled through a unique identifier that is obtained as a result of the subscribe operation. After cancellation, no more notifications are delivered, the query is removed from the system. Given flexible event subscription, this operation must be split into two parts as well:

C. *unsubscribe(subscriptionId)*

Removes a notification-recipient for a given query-id. Note that the buffer and query instance remain intact, so that other recipients can still subscribe.

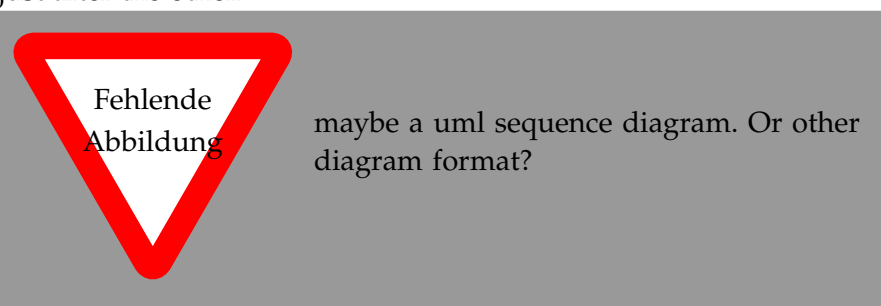
D. *removeQuery(queryId)*

Completely deletes the query and its buffer, so that no notifications are sent out any longer.

All four methods must be available to execute a subscription on process deployment. The query must be registered using *registerQuery* before the process instance, i. e. notification-path, is available. For each process instance, a subscription can be issued individually using *subscribe* and thereafter, the notification-recipient can be removed with *unsubscribe*. The query and its buffer will remain active even after any single instance has terminated. When the process gets un-deployed, the query can be deleted using *removeQuery*. [Section 6.2](#) describes the steps in detail.

show the steps with sample data from one of the examples

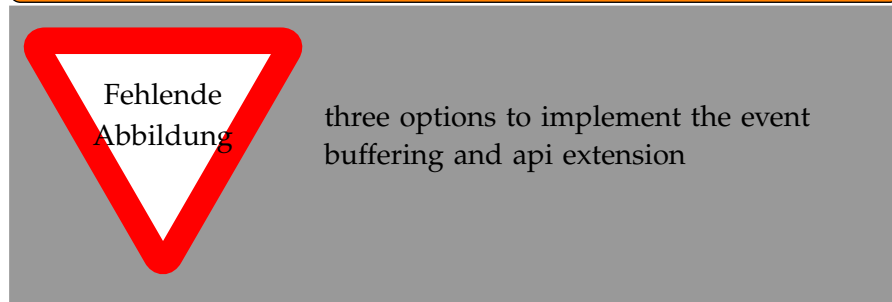
- to achieve the default behavior, each two steps have to be executed just after the other.



maybe provide a swagger definition for this?

ARCHITECTURAL OPTIONS TO IMPLEMENT THE EXTENDED API  
That extension can either be implemented by adopting the CEP platform itself, by implementing a separate middleware between process engine and CEP platform, or by implementing a buffering module as part of the process engine. Which of the three options suits best has to be evaluated for the given use-case and existing infrastructure. In some cases it might not be possible to adapt the code of process engine or event processing platform, which leaves a separate middleware as the only choice.

Generally, extending the event processing platform is advisable?  
Or put up brief pros and cons of each of the options



A reference implementation for an extended complex event processing platform is presented in [Section 7.1](#) at the example of the Esper-based CEP platform *Unicorn*. It also explains, why extending the event platform was the preferred choice in the given scenario.

the communication overhead must be considered: if many events are sent to the buffer, but they are rarely issued to a process instance, then it will make sense to place the buffer closer to the event engine w.r.t. the the involved comm overhead. Though it must be noted, that if events occur rarely, but many process instances consume them, than there will be an additional overhead to send events from the buffer to the instance. in that scenario, it might be advantageous to implement the event buffering close to the process engine

PERFORMANCE CONSIDERATIONS - performance improvements through shared windows - due to its performance optimizations, an extension of the cep itself would make perfect sense - given the extended api, it is now possible to implement early event subscription from the process engine

write

## 6.2 EXTENDED PROCESS ENGINE BEHAVIOR

It is the task of the Business Process Engine to interpret and execute process models and connect to an event processing platform in



event-driven setups. From the three relevant parts, two have already been defined, the BPMN extension and the buffered event processing module. Out of the box, a process engine like Camunda will ignore any proprietary BPMN extensions and the subscription to an event source must be especially implemented. An example for such an implementation is provided in [Section 4.2](#). One goal of this work is to automatize the handling of event subscriptions solely based on the information available through the extended BPMN model. Additional process elements should not be required. This section will clarify, which operations need to be executed by the process engine to enable the automatic subscription handling. [Section 7.2](#) demonstrates the implementation of automatic subscription handling at the example of Camunda.

#### PARSING ADDITIONAL INFORMATION FROM THE BPMN MODEL

It is required that the process engine is able to read the additional information from the BPMN extension (see [Section 5.1](#)) so that it is available during process deployment and execution. This affects the BPMN message element, which can contain the additional attributes *eventQuery*, *subscriptionTime* and *bufferPolicies*. Secondly, the *Explicit Subscription Task* has to be processed. It contains a reference to a Message entity within the same model element. The process engine might have to be adopted to read all relevant data from the extended model.

**MANAGING SUBSCRIPTION AND UN-SUBSCRIPTION** As defined in the BPMN extension for flexible event subscription, the action of subscribing to an event source can happen at different times during process deployment and execution. The options and the implicit timing of subscription and un-subscription are specified in [Section 5.1.2](#). The process engine must communicate with the process engine using the four calls *registerQuery*, *requestEvents*, *unsubscribe* and *deleteQuery*, that were presented in the previous chapter. For each possible subscription time, the following briefly enumerates which operations must be executed when.

**IN EVERY CASE:** The return-value of *registerQuery*, a unique identifier of that query, must be stored for the related BPMN Message. The id is later necessary to execute the other three API methods. When an event element is reached, a call to *requestEvents* must be issued. When the execution of that event element is finished, call *unsubscribe*.

**SUBSCR. ON PROCESS DEPLOYMENT:** When a process gets deployed, the process engine must check if subscription information is in the model. For every Message element that is set as *subscribe on pr. deployment*, a call to *registerQuery* must be issued as part of

the deployment process. A call to *deleteQuery* is executed when the process gets un-deployed for the same messages.

**SUBSCR. ON PROCESS INSTANTIATION:** When a process gets instantiated, *registerQuery* must be executed for each *Message* that is set to *subscribe on pr. instantiation*. *deleteQuery* can be called when the last reachable event element for a *Message* has finished executing or no connected event can be reached anymore. The deletion happens at the latest when the process instance terminates.

**SUBSCR. THROUGH EXPLICIT SUBSCRIPTION TASK:** If the control flow reaches a subscription task, the process engine executes *registerQuery* for the referenced *Message*. The execution of *deleteQuery* follows the same rules as in the preceding case.

**SUBSCR. WHEN THE EVENT ELEMENT IS REACHED:** Once the event element is reached, *registerQuery* must be executed for any *Message* that is not covered by one of the prior cases. *deleteQuery* must be called when the event element is finished.

be more precise about the time the calls should be executed (if possible). "reached"? "completed"? use the right bpmn words

Do I want to write about extended validity checks? soundness?  
The question would be: What happens if the model erroneous w.r.t. to the bpmn extension?

- handling subscription dependencies: - when are process variables replaced by their actual values - The use of dynamic process variable values introduces an additional complexity: Depending on the time of event subscription, the value of the process variable might not yet be available. - reference BPMN data elements: process INSTANCE variable -> the variable value might only be available during instance execution -> can we find an exact definition of this in the spec? - see BPMN2 spec pp.211+ : Process and Activity can have DataInput and DataOutput. DataInput can have an 'optional' attribute - during execution the variable data might or might not be available. Related Work: Francesca? -> too complex, we need a simplification for this. - what happens if the data is not available?

Time of un-subscription also must be clarified in bpmnx

write conclusions that repeat how each requirement has been fulfilled

## REFERENCE IMPLEMENTATION

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[Chapter 5](#) has presented a concept that involves the model layer, the process engine and the event engine. It enables flexible event subscription to overcome the issues revealed in ... By explicitly choosing the time of event subscription and bridging the gap between event reception and consumption by the process instance it is possible to ...

write from keypoints

the model extension is formalized in [Section 5.1](#), hence BPMN models extended by all information necessary for flexible event subscription can be created.

to evaluate our results and provide a reference implementation, we enhance the business process engine Camunda and the event processing platform UNICORN following the procedures described in [Section 6.2](#) and [Section 6.1](#) respectively. Camunda is extended by providing a Process Engine Plugin, Unicorn by adapting the source code. The resulting illustration of an event-driven architecture is evaluated by implementing the examples shown in [Section 3.1](#).

Application to examples missing

### 7.1 EXTENDING THE EVENT PROCESSING PLATFORM UNICORN

UNICORN is an event engine developed for academic purposes at the Business Process Technology chair of Hasso-Plattner-Institute, Potsdam. [13] It focuses on event processing for BPM, ... Based on Esper, that means, what is esper what unicorn adds to esper: graphical and rest user-interface to ..., event persistence store, bp execution and correlation, recently also an event generator/replayer ==> it translated user interactions into the according Esper method calls

- currently unicorn uses version 5.3.0 of esper, therefor all future considerations are made based on the documentation of that version

write from keypoints

It has been decided to extend the event engine in this scenario, for several reasons. Firstly, the performance and operation of the process engine shall not be jeopardized. Event Processing features potentially require a lot of performance to handle a large number of requests in a short amount of time. By implementing the event Buffering module loosely coupled to the Process Engine, we ensure that its performance is not influenced. Implementing a separate middleware was avoided in the course of this work, to keep the architecture easy and concise

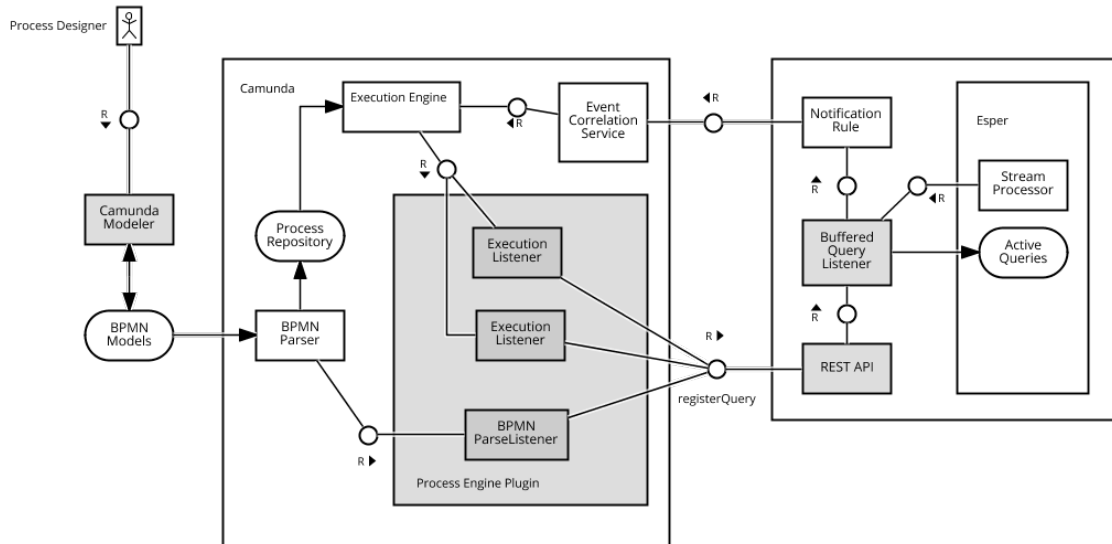


Figure 18: Architecture for flexible event subscription in Camunda and Unicorn

to describe. Unicorn is an academic prototype in constant development and therefore predestined to be directly extended for this kind of use-case. Moreover, the chosen architecture promises performance advantages thanks to running the event buffering in the same [JVM](#) as Esper.

#### 7.1.1 Event Buffering

To allow the delayed delivery of events, buffering functionality is added to Unicorn. Whereas, originally, events that match a certain query in esper are directly forwarded to the notification-recipients, they will instead be held in a buffer until requested by recipient. If a recipient is already subscribed, than the behavior will be equivalent to the original scenario, because the event will be delivered instantly.

**ENGINE-SPECIFIC IMPLEMENTATION OPTIONS** When investigating the options for implementing event buffers in UNICORN, solutions that are specific to the event processing platform were considered. Notably, the concept of a *Window* in event processing languages is very much comparable with a buffer as described in this work. Windows hold a variable number of events in memory according to specified window properties. That way, windows can act like size- or time-oriented buffers and allow access to older events whenever a new event occurs. *Named Windows* in Esper allow to create global data windows that can be modified and read from multiple statements.

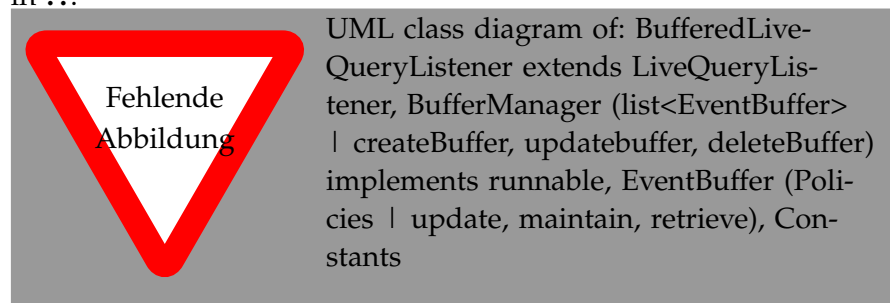
Similar functionality is provided by *Tables*, which follow a relational approach by primary key.<sup>1</sup>

If a query window is connected with the Esper-specific *Output Clause*, it is possible to delay query output with timing constraints or trigger it on the change of a global variable.<sup>2</sup> When putting together these features, an event buffer could be implemented as follows: The call *registerQuery* instantiates the window. If the *Consume-policy* is used, we use a named window that we share among queries and delete from after receiving the event. At *subscribe*, a notification-recipient is added to a query and output is triggered by causing a variable change in Esper. *Unsubscribe* and *removeQuery* remove a notification-recipient and de-register the query respectively.

A major drawback of this approach is, that the event queries have to be adopted to use the specific features which makes additional knowledge necessary when designing the processes. Alternatively, the original queries submitted by the process engine can be transformed before registering them in Unicorn, for example by encapsulating them in a sub-query. That way they make use of the mentioned features, but certain expressions will not be possible anymore due to limitations of the Esper EPL with regards to sub-queries.

Apart from this approach, it is worth noting that Unicorn offers a persistent storage of historic events, which can theoretically be used to implement the buffering functionality. However it would be necessary to abuse the relational approach to implement a persistent stream buffer on top of the SQL database, which is not desirable.

**THE GENERIC BUFFERING SOLUTION** The investigation of engine-specific solutions did not bring up an optimal solution to implement an event buffer with existing mechanisms. Furthermore, this reference implementation should not be entirely limited on one specific event engine. For that reason a new generic buffering module has been implemented within Unicorn. A UML class diagram is provided in ??.



<sup>1</sup> see Chapter 6. EPL Reference: Named Windows And Tables,  
<http://www.espertech.com/esper/release-5.3.0/esper-reference/html/nwtable.html>

<sup>2</sup> see Chapter 5. EPL Reference: Clauses,  
[http://www.espertech.com/esper/release-5.3.0/esper-reference/html/epl\\_clauses.html](http://www.espertech.com/esper/release-5.3.0/esper-reference/html/epl_clauses.html)

The module is built around the *EventBuffer*-class. Objects of the class are managed by the *BufferManager* and represent a single buffer entity, holding events of one query. Esper provides the query output as an object of type *EventBean*, which is stored in a list in *EventBuffer*. The buffer behaves according to the value of its policies, which influence the length of the list, the order in that items are retrieved from the buffer and the consumption behavior. The lifetime-policy, which requires that events are deleted from the buffer after a certain time, is ensured by a maintenance-thread that runs from the *BufferManager* class and iterates over all *EventBuffer* objects in a specified time interval. The default value is 5 seconds.

Unicorn uses the *LiveQueryListener* to react on new event-occurrences. It implements the Esper *UpdateListener*-interface and is registered in Esper as callback-function for a new query. Whenever an event matches that query, the object of *LiveQueryListener* is notified. Originally, the *QueryListener* notifies all notification-recipients that are known for that query and then drops the event. For the event buffering, a new class *BufferedLiveQueryListener* is created which extends the behavior of the standard query listener. Instead of only notifying all recipients, it also updates the associated buffer instance with the latest query output. The provided implementation serves for demonstration and reference purposes. It is not optimized for performance, for example in the case of overlapping data between buffers. Considering an event subscription issued on process instantiation, multiple *EventBuffer* instances will essentially store the same data, if multiple instances of the process are started at the same time or with minimal difference.

The presented buffering module supports the temporary storage of query results as demanded by *??*. After query creation, matching events can be stored in a buffer to be retrieved when necessary. All buffer policies described in [Section 5.1.4](#) have been implemented. The following section explains how the [REST API](#) of Unicorn is extended to make use of the buffering module.

### 7.1.2 REST API Extension

Unicorn offers a webservice that allows users to interact with the event processing engine via the Hypertext Transfer Protocol ([HTTP](#)). The restful API comprises the most basic functionality, query registration, query deletion and obtaining query strings by the subscription identifier. An interaction generally works as follows: The user executes a POST to *<platform>/EventQuery/REST*, providing an event query and information about the desired notification-recipient (*notification path*). Unicorn registers the query in Esper and returns a unique identifier to the user. Whenever an event matches the query, the platform sends a notification to the specified recipient. a DELETE call to

`<platform>/EventQuery/REST/{eventQueryUuid}` triggers the removal of the query.

In accordance with the API requirements introduced in [Section 6.1](#), additional functionality has been added to the Unicorn Webservice. The methods were added under a new path `/BufferedEventQuery` to make sure that the existing features remain. They are as follows:

**REGISTER QUERY** POST to `/BufferedEventQuery`, returns `queryId`

Payload: JSON (eventQuery[, bufferPolicies]) with  
bufferPolicies:(lifetime, consumption, size, order)

Description: The provided eventQuery is registered in Esper using a BufferedLiveQueryListener. The BufferManager is used to instantiate a new EventBuffer object. The payload JSON object bufferPolicies is optional, and will be passed to the EventBuffer if available. Otherwise, the system falls back to the default values ([Table 1](#)). A unique identifier of that query and associated buffer is returned.

**SUBSCRIBE** POST to `/BufferedEventQuery/{queryId}`

returns `subscriptionId`

Payload: JSON (notificationPath) with  
notificationPath:(notificationAddress, processInstanceId, message-Name)

Description: An new subscription is added to the selected query, so that a notification is issued based on the current buffer content and whenever an event matches the query. The notification-path is specified including the id of the target process instance and the message name. This enables Camunda to automatically correlate the issued notification to the right process execution and message.

**UNSUBSCRIBE** DELETE to `/BufferedEventQuery/{queryId}/{subscriptionId}`

Description: Remove the specified subscription from the list of subscriptions of the selected query.

**REMOVE QUERY** DELETE to `/BufferedEventQuery/{queryId}`

Description: Remove the query and the associated buffer from the system.

### ??? Swagger definition of the implemented rest API

Using instances of *NotificationRule*, Unicorn sends notifications to the recipients (see [Figure 18](#)). In our scenario, the messages must be sent to Camunda, more specifically to a specific process instance within the engine. The *Event Correlation Service* within Camunda is responsible for relating incoming message events to process instances. One way to enforce the correlation is by inserting the process instance



identifier and the message name inside the message. For that reason, message name and process instance id must be provided when adding a subscription in the event engine. Unicorn has been adopted to send notifications to Camunda including the required correlation information. A sample notification for a eurotunnel delay event is shown in [Listing 2](#).

[ref/footnote camunda documentation](#)

Listing 2: Example of a JSON notification sent by UNICORN

```
{
  "messageName": "eurotunnelDelay",
  "processInstanceId": "274a876f-aed7-4a1a-916b-e85a0c2416f7",
  "processVariables": {
    "eurotunnelDelay": {"value": "60", "type": "Integer"}
  }
}
```

can i reference another part of the thesis? could be mentioned in buffered event handling

After implementing the necessary extensions to the event engine, it is the task of the process engine to connect the extended process model to the buffered event handling API. The adaptations to the process engine are presented in the following section.

## 7.2 EVENT SUBSCRIPTION HANDLING IN CAMUNDA

Through the BPMN extension for flexible event subscription, the information necessary to register event queries in a CEP platform can be made available in process models. [Section 6.2](#) outlines how the business process engine must be adapted to execute the operations for subscription handling. Camunda is an open-source business process engine with support for the latest version of the BPMN and used to exemplify the modification process. Further information about Camunda is provided in [Section 2.1.2](#).

[Figure 18](#) depicts a simplified architecture of Camunda, highlighting modifications in gray. Our workflow for flexible event subscription mainly involves the core components execution engine, model repository, correlation service and the Camunda modeler.

### 7.2.1 Employing *ExecutionListeners* in a Process Engine Plugin

Being a community-driven open-source project, the Camunda project offers numerous options for customizing and extending its behavior. A core concept to trigger the execution of custom code during a process execution are *ExecutionListeners*. They can be incorporated in BPMN models using the properties window of the Camunda Mod-



eler as demonstrated in [Chapter 4](#). In that case they have a textual representation in the model file, a specification can be found in the product documentation <sup>3</sup>.

However, one of the main goals of this thesis is to provide an alternative solution to explicitly modeling the subscription, unless an explicit subscription task shall be used). Thus the additional necessity attach execution listeners to BPMN elements shall be avoided. Instead, subscriptions shall be managed automatically by the process engine solely based on the subscription definition provided in the extended message element. To achieve this kind of behavior, Camunda offers the concept of a Process Engine Plugin (PEP) to intercept significant engine operations and introduce custom code. By this means, execution listeners can be added programatically. The plugin is a separate software module that implements the interface *ProcessEnginePlugin* <sup>4</sup>. It is activated by adding a *plugin* entry in the process engine configuration.

[23] and [27] have chosen to directly adapt the source code of camunda. More precisely, they propose to modify the *Behavior* class of the Camunda core to execute additional code when a BPMN element starts executing. In this work, we implement a Process Engine Plugin as it allows a clearer, more understandable approach to adopting the execution behavior. That holds especially when it is only necessary to execute additional operations and not modify or delete existing code. Moreover, the PEP facilitates the re-usability across environments and different versions of Camunda.

from [Section 5.1.1](#)

that means that in the implementation we need additional configuration values -> implementation chapter

### 7.2.2 Managing event Subscriptions at Runtime

The Process Engine Plugin enables us to execute custom JAVA code at predefined points during engine execution. An entry point to the modification of the execution behavior is provided through the implementation of the *ProcessEnginePlugin* interface. It allows to intercept *preInit*, *postInit* and *postProcessEngineBuild*, that means at three different points during the engine bootstrapping. We chose to provide a custom implementation for the *preInit* method, which has an object of *ProcessEngineConfigurationImpl* as parameter. With the engine con-

<sup>3</sup> Camunda BPMN Extension Elements, <https://docs.camunda.org/manual/7.7/reference/bpmn20/custom-extensions/extension-elements/#executionlistener>

<sup>4</sup> Process Engine Plugins, <https://docs.camunda.org/manual/7.7/user-guide/process-engine/process-engine-plugins/>

figuration available, a plethora of custom handlers, validators and listeners can be registered<sup>5</sup>.

We implement a *BPMNParseListener* that is to be executed after a BPMN element finished parsing (*customPostBPMNParseListeners*). The *BPMNParseListener* interface allows us to react to the parsing of single elements based on their type by making a separate method for every BPMN element available. A helpful example for using that listener is also provided on GitHub<sup>6</sup>. The BPMN extension affects the intermediate and boundary message catch event, the receive task, the explicit subscription task and can also be used to specify subscription information for a message start event. It was attempted to utilize the separate methods of every element, but it is not possible to access the subscription information in the associated message elements from that point. Apart from the provided method arguments, the engine's *RepositoryService* can normally be used to view model information, but only after the deployment was finished. To overcome this limitation, we implement the method *parseRootElement*, to which a representation of the root element *bpmn:definitions* is passed, including all message and extension elements. Before the content of this method is described, two more aspects need to be introduced, the *SubscriptionEngine* and the *ExecutionListeners*.

**SUBSCRIPTIONENGINE** A new class *SubscriptionEngine* has been introduced, that encapsulates most of the functionality needed to communicate with the API of the event engine. It offers method representations of all four API-calls, *registerQuery*, *subscribe*, *unsubscribe* and *removeQuery*, translating them into the according HTTP calls which are then executed. The second, more important functionality is the *SubscriptionRepository*, a list of all available query and subscription identifiers and a mapping to the related process definitions or instances. As described in [Section 7.1.2](#), active queries and subscriptions can only be deleted using their unique identifiers. When issuing a subscription or registering a query, these identifiers are stored in the repository until the removal call is executed.

**EXECUTIONLISTENERS** The concept involves four execution listeners, one for each API-call, which are attached to activities through the BPMN parse listener. Execution listeners can be triggered by the camunda-internal end- or start-event of an activity and can therefore be used to implement the event subscription handling. Their imple-

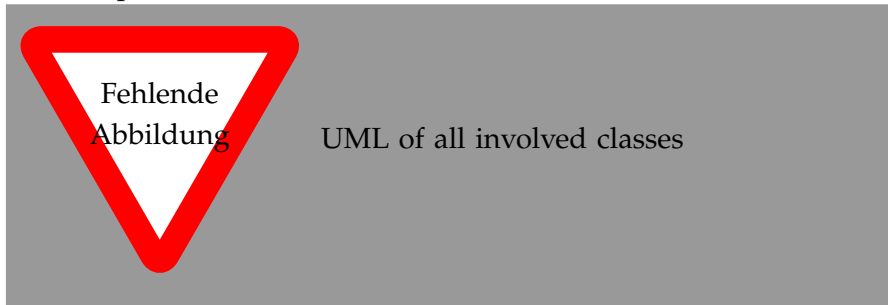
<sup>5</sup> Some arbitrary examples are pre-/post-deployers, event-handlers, post-variable-serializers, command-interceptors. Full list available from: <https://docs.camunda.org/javadoc/camunda-bpm-platform/7.7/org/camunda/bpm/engine/impl/cfg/ProcessEngineConfigurationImpl.html>

<sup>6</sup> GitHub: [camunda-bpm-examples/process-engine-plugin/bpmn-parse-listener, https://github.com/camunda/camunda-bpm-examples/tree/master/process-engine-plugin/bpmn-parse-listener](https://github.com/camunda/camunda-bpm-examples/tree/master/process-engine-plugin/bpmn-parse-listener)

mentation is concise: each of them is a separate class implementing the *ExecutionListener* interface. They only have one method *notify*, which is called by the engine when the end- or start-event fires. Within the method, each listener class issues an API-call using the functionality exposed by *SubscriptionManager*.

could mention that *registerQuery* can be bound either to an instance or to a process definition. this differentiation is made based on the *subscriptionDefinition* which is passed to the listener on creation.

All implemented classes are visualized in ... .

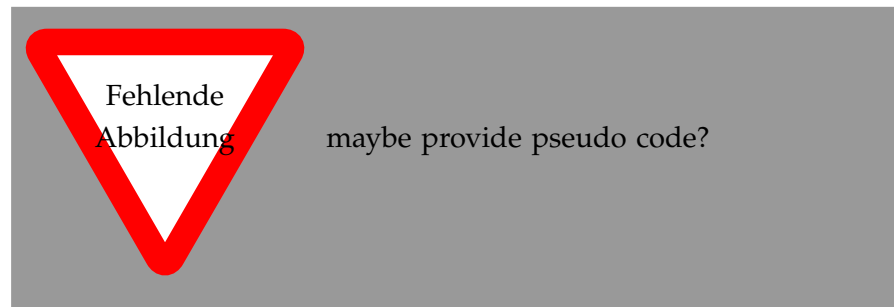


**IMPLEMENTATION OF PARSEROOTELEMENT** After introducing the *SubscriptionEngine* and the *ExecutionListeners*, this paragraph describes how the listeners are attached to the activities in the course of the execution of *parseRootElement*. By the help of the model representation passed to the method as an argument of type *Element*, all relevant xml elements can be extracted based on the element names. This results in a list of all intermediate and boundary message catch events and receive-tasks and the messages that they reference. The subscription definition is available as a child element of each message. For each element in the list, the execution now proceeds as outlined in [Section 6.2](#). Depending on the specified subscription time, *ExecutionListeners* are added at the beginning or end of activity executions or process executions. If the specified subscription time is *Process Deployment*, the provided query is registered immediately using the *SubscriptionManager*.

This can be exemplified by an intermediate message catch event with subscription time *on process instantiation*:

- An instance of *RegisterQueryListener* is added to the start event of the process. The *subscriptiondefinition* is provided to that listener.
- The *SubscribeListener* is attached to the start event of the activity representation of the intermediate message event.
- To the same activity, objects of *UnsubscribeListener* and *RemoveQueryListener* are registered to be triggered by the end event.

The complete workflow is shown as pseudo code in ??.



It is worth noting that the available code-base can easily be used to implement the event subscription for message start events, so that processes can be automatically instantiated by events delivered from the event engine. Considering that process definitions in Camunda are rarely removed completely, the removal of queries on process undeployment was not implemented. The lack of an obvious intercept point to inject code at that time makes an implementation cumbersome.

Altogether, the enhanced CEP platform and business process engine enable the automatic handling of information provided through the BPMN extension for flexible event subscription. The event engine exposes functionality for buffered event handling which is accessed through execution listeners during the process execution in Camunda. The Camunda extension can be flexibly re-used across environments thanks to its implementation in the form of a Process Engine Plugin. Given these extended features, process designers can conveniently incorporate external message events in their BPMN models.

## RELATED WORK

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The field of Event-driven Business Process Management (**EdBPM**) has developed from connecting **CEP** and **BPM** in an attempt to increase the quality and performance of business processes. **EdBPM** must be viewed from different perspectives. In event-driven business process monitoring and business activity monitoring, the process engine is solely an event producer that publishes information to an event engine. A considerable amount of research was undertaken to make use of CEP to detect process violations and execution performance. ?? BAM, process mining... [16] In [20], the authors did a comprehensive survey of research efforts in the area event-driven business process management and conclude that... > they point out that most research is in this area

...

This work has investigated how a flexible event subscription handling can be incorporated in BPMN models. It is therefor contributing towards a facilitation

roughly: event-driven business process management (anything that correlates bpm and cep)

Research in the area of event-driven business process management has been undertaken for several years now.

[27] makes use of events to control process executions and considers event subscription as an essential step in the process. However, the subscription time cannot be chosen flexibly.

[23] has been the starting point for this thesis, not only acknowledging the discrepancy between the event occurrence time in the real world and the event subscription time as interpreted from the BPMN standard, but also incorporating event subscription in the BPMN model. Inspired by their work, this thesis revisits the topic from scratch however:

[6]: investigations on the connection between events and process models, leading to [3]: extension of bpmn for receiving events, process variables in the query

similar to chapter 4: [2]: includes the streams and stream processing right into the process model same: [4]

(1)

> business activity monitoring

(2) - edpc: baumgraß, re-eval decisions - eda [7]: Process Instantiation, talks a lot about subscription, but no flexible subscr time, nobuffer, only for instantiation

[1]: Integrating Complex Events for Collaborating and Dynamically Changing Business Processes; talk about un/subscription, extend wsbpel with subscription, reporting, patterns, expression language, ..; scope, time of subscription and event buffering not discussed also: [17]

[21]: acknowledge the lack of usability in CEP and address it by applying bpmn as graphical support for the definition of cep patterns [11]: generate queries from new graphical notation, anbieterunabhängige Modellierung von EdBPM

- use case implementations [9]: event-driven manufacturing process;

- correlating events to processes

- maybe: persisting events in cep platforms | or delayed delivery of events > [28]: Event data warehousing for complex event processing and [5]:Event-Driven services: Integrating production, logistics and transportation also [22]: historic data in pub/sub > but if i want to have that historic information, than I would have to issue a subscription some time before. I integrate everything in the process model

- event buffering for the application in a BPM environment may not be confused with internal buffering techniques used within event engines to perform load shedding, even out short term load peaks.

## CONCLUSIONS

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By implementing the flexible event subscription concept in Unicorn and Camunda it could be demonstrated that state-of-the-art BPM and CEP technology is flexible enough to handle

### 9.1 DISCUSSION

- discussion: > it remains the problem, that epl knowledge is not available in process design, a problem that has been addressed in ... > while the event buffering can now be controlled from the process model level, events must still be acquired by the event engine before (set up event types and make sure that the events are pushed to the engine) > the concept tries to bridge the gap between event persistence/historic events and real-time event processing by arguing that event occurrences are kept for a limited time and can therefor still be treated as events. A contra argumentation might be that an event becomes a simple piece of information as soon as it is not instantaneously consumed. hence it should not be modeled as an event but simply as a data object, using a data store > many design decisions have been begründet with an improved usability though no empirical basis was available. Reasoning was only based on literature and chats with a small group of fellow researchers

- it can be argued that at process design time you dont care about the subscription time itself, but about the maximum age you can accept for your events. this would connect the lifetimepolicy and the subscription time into one value. but make the processing more complex and more fuzzy

### 9.2 FUTURE WORK

- modeling subscription dependencies as data objects and automatically evaluating the earliest possible time of subscription - (Andreas:) using the bpmn extension with multiple event engines - this work attempts to provide a standard for handling event subscription in bpm architectures. given the necessity and the repeated attempts to address to topic, it would be a reasonable next step to discuss available solutions and start providing a foundation that is accepted and reused across the bpm community. No matter if based on the solutions of this work or not. > it must be evaluated if that's the way event buffering shall be used. > prove the value of flexible event subscription in an industry case study - future work: event data from other

events could be used as historic data to allow the access even before process deployment - time of subscription could also be defined as "at termination of activity xy"



## APPENDIX

Listing 3: XSD schema of the BPMN extension for flexible event subscription

```

<xsd:schema xmlns:flexsub="http://www.some.url/" xmlns:xsd="http://
  www.w3.org/2001/XMLSchema" targetNamespace="http://www.some.url
  /">
  <xsd:element name="explicitSubscriptionTask" type="
    tExplicitSubscriptionTask" />
  <xsd:complexType name="tExplicitSubscriptionTask">
    <xsd:complexContent>
      <xsd:extension base="tTask">
        <xsd:attribute name="messageRef" type="xsd:QName" use="required"/>
      </xsd:extension>
    </xsd:complexContent>
  </xsd:complexType>

  <xsd:element name="subscriptionDefinition">
    <xsd:complexType>
      <xsd:sequence>
        <xsd:element type="xsd:string" name="eventQuery" minOccurs="1"
          maxOccurs="1" />
        <xsd:element type="tSubscriptionTime" name="subscriptionTime"
          minOccurs="0" maxOccurs="1" default="Element Execution"/>
        <xsd:element name="bufferPolicies" minOccurs="0" maxOccurs="1">
          <xsd:complexType>
            <xsd:sequence>
              <xsd:element type="xsd:string" name="lifetimePolicy" minOccurs="0"
                maxOccurs="1" default="infinite"/>
              <xsd:element type="xsd:string" name="consumptionPolicy" minOccurs="0"
                maxOccurs="1" default="Reuse"/>
              <xsd:element type="xsd:integer" name="sizePolicy" minOccurs="0"
                maxOccurs="1" default="-1"/>
              <xsd:element type="tOrderPolicy" name="orderPolicy" minOccurs="0"
                maxOccurs="1" default="FIFO"/>
            </xsd:sequence>
          </xsd:complexType>
        </xsd:element>
      </xsd:sequence>
    </xsd:complexType>
  </xsd:element>

  <xsd:simpleType name="tSubscriptionTime">
    <xsd:restriction base="xsd:string">
      <xsd:enumeration value="Process Deployment"/>
      <xsd:enumeration value="Process Instantiation"/>
      <xsd:enumeration value="Manual"/>
      <xsd:enumeration value="Element Execution"/>
    </xsd:restriction>
  </xsd:simpleType>

  <xsd:simpleType name="tOrderPolicy">
    <xsd:restriction base="xsd:string">
      <xsd:enumeration value="LIFO"/>
      <xsd:enumeration value="FIFO"/>
    </xsd:restriction>
  </xsd:simpleType>
</xsd:schema>

```

complete xml example of a BPMN model using the extensions

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## DECLARATION

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Put your declaration here.

*Potsdam, August 2017*

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Dennis Wolf





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