

MPI

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1. MPI Overview

1.1 What is MPI (Message Passing Interface)?

- **A standard API** for message passing between distributed memories in parallel computing.
- MPI assumes a **distributed-memory computing system**
- MPI can run on **shared-memory computing system**
- MPI programming model (basically) uses **SPMD**(Single Program, Multiple Data).

1.2 Parallel Programming Classification

- **Multi-Process**: MPI(Message Passing Interface), HPF(High Performance Fortran)
- **Multi-Thread**: OpenMP, Pthread(POSIX Thread)

Aspect	MPI	HPF
Type	Parallel communication library	Fortran language extension
Language	C / C++ / Fortran	Fortran only
Parallelism Control	Fully manual by programmer	Mostly compiler-driven
Flexibility	Very high	Limited
Maintainability	Hard but highly tunable	Simpler but harder to tune
Learning Curve	High	Low to medium
Current Usage	Mainstream in HPC	Obsolete / deprecated

1.3 MPI Features

- **Communication Model:**
 - Uses message passing for communication between processes.
- **Distributed Memory Support:**
 - Each process has its own memory space, no shared memory.
- **Multi-node Capacity:**
 - Can run across multiple nodes; abstracts network communication.
- **Standardized API:**
 - Standardized interface in C, C++, and Fortran; highly portable.
- **Multiple Implementation:**
 - Available implementations include OpenMPI, MPICH, and Intel MPI, etc.
- **Difficult to Debug:**
 - Debugging is challenging due to concurrency and communication complexity.

1.4 Typical example of Usage

- **Simulation on a supercomputer**
 - Physics, Meteorology, Chemistry, etc.
- **Data processing in large-scale data analysis**
 - e.g., genomics, astronomy
- **Machine learning training on large datasets**
 - e.g., distributed deep learning

1.5 Comparison between implementations

	OpenMPI	MPICH	Intel MPI
Developer	Universities, Companies	Argonne National Laboratory	Intel Corporation
Distribution	Open source	Open source	Closed source
Optimization Target	General purpose	Lightweight, stable	Optimized for Intel architecture
Performance	Medium to high	Lightweight, stable, scalable	Best performance on Intel CPUs
Main Use	Academic clusters, general HPC	Research, education	Commercial HPC, Intel clusters

- Miyabi uses OpenMPI as the default MPI implementation. (Miyabi is the supercomputer system of the University of Tokyo)
- To be more specific, mpicc on Miyabi is bound to nvc compiler (NVIDIA HPC SDK C compiler), which means NVIDIA HPC SDK + MPI environment is used.

1.6 Key Communication Primitives

- **System function:**
 - MPI_Init, MPI_Finalize, MPI_Comm_size, MPI_Comm_rank
- **Point-to-point communication:**
 - MPI_Send, MPI_Recv
- **Collective communication:**
 - MPI_Bcast, MPI_Reduce, MPI_Alltoall
- **Synchronization:**
 - MPI_Barrier, MPI_Wait, MPI_Test
- **Derived data types:**
 - MPI_Type_create_struct, MPI_Type_vector
- **Non-blocking communication:**
 - MPI_Isend, MPI_Irecv
- **Remote memory access:**
 - MPI_Put, MPI_Get
- **Process management:**
 - MPI_Comm_spawn, MPI_Comm_free

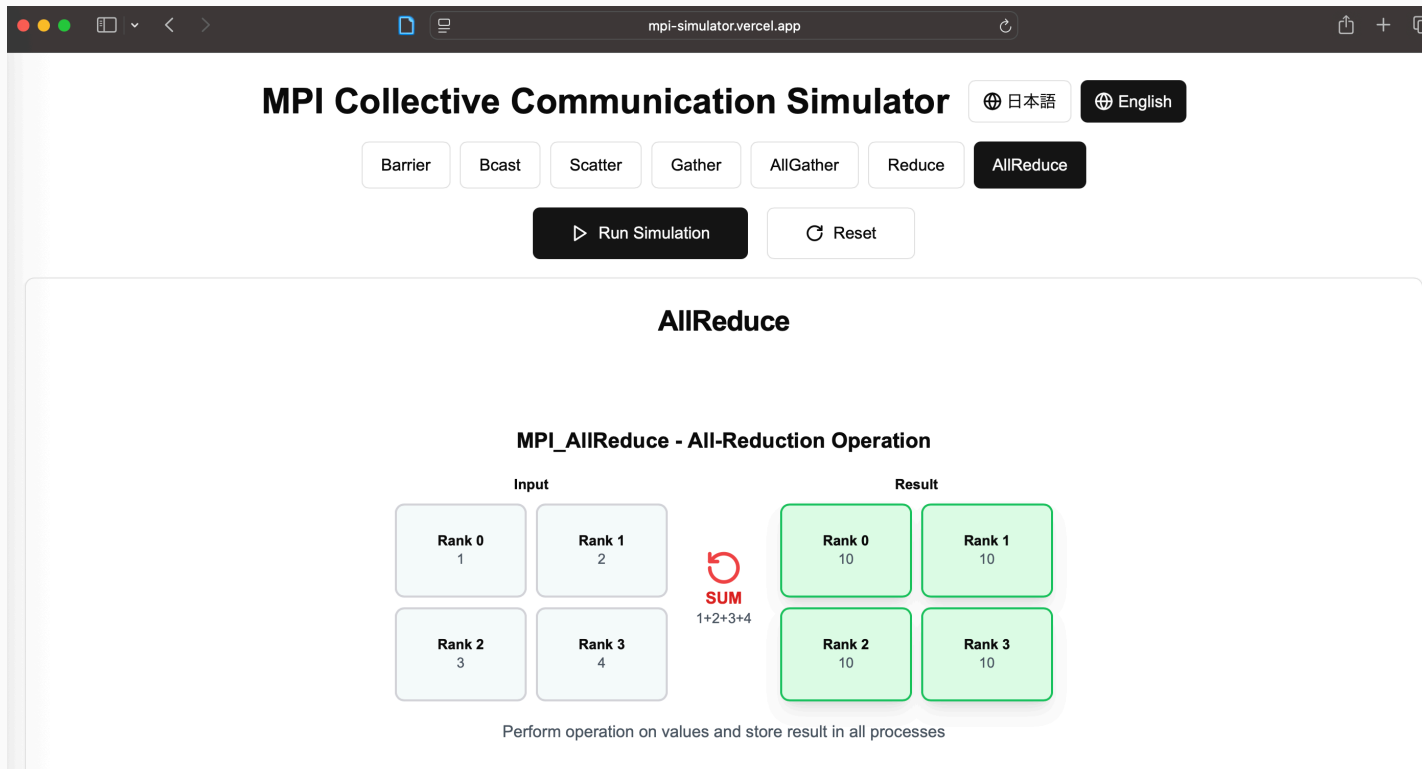
1.7 Internal Mechanisms of MPI Communication

- MPI implementations (e.g., MPICH, OpenMPI) rely on various OS system calls and low-level libraries for communication.
- **Intra-node Communication (within the same node):**
 - Shared memory: `mmap`, `shm_open`, `memfd_create`, System V `shm`
 - Event waiting: `futex`, `poll`, `epoll`, `select`
 - Pipes & sockets: `write`, `read`
 - UNIX domain sockets: `sendmsg`, `recvmsg`
 - Synchronization: `sem_open`, `pthread_mutex`, `spinlock`
- **Inter-node Communication (across nodes):**
 - TCP/IP sockets: `socket`, `bind`, `listen`, `accept`, `connect`, `send`, `recv`, `sendto`, `recvfrom`
 - RDMA & verbs API (e.g., InfiniBand) for zero-copy communication
- **High-Performance Cluster Communication:**
 - Libraries: `libibverbs`, `UCX`, `OFI` (Libfabric), `XPMMEM`, `NVLink`
 - Kernel bypass with DMA (Direct Memory Access) for low-latency, high-throughput transfers

2. Basic Learning of MPI

2.1 MPI Simulation Website

- I developed a simple MPI simulation website for learning MPI.
- Please Click on: [Izawa MPI Simulation Website](#)



2.2 Tutorial Programs

- The following programs are available below.

[MPI Tutorial GitHub Repository](#)

2.3 Minimum MPI Program

2.3.1 MPI Language Differences

	C	C++	Fortran
MPI Header	<code>#include <mpi.h></code>	<code>#include <mpi.h></code>	<code>use mpi</code> or <code>include 'mpif.h'</code>
Official MPI support	○	▲	○
Syntax intuitiveness	Explicit C syntax	Almost same as C	<code>call</code> and subroutine based
Compiler	<code>mpicc</code>	<code>mpicxx</code> or <code>mpic++</code>	<code>mpif90</code> or <code>mpifort</code>
Scientific computing	○	▲	◎

- C++
 - MPI-3.0 abolished C++ only bindings.
 - Currently, C++ also uses C interface.
- Fortran
 - Considering readability, type safety, and portability, `use mpi` is recommended.

2.3 Minimum MPI Program

2.3.2 Hello World (C)

```
#include <mpi.h>
#include <stdio.h>
#include <stdlib.h>

int main(int argc, char *argv[])
{
    MPI_Init(&argc, &argv);

    int num_procs;
    int my_rank;

    MPI_Comm_size(MPI_COMM_WORLD, &num_procs);
    MPI_Comm_rank(MPI_COMM_WORLD, &my_rank);

    printf("Num of Proc : %d\n", num_procs);
    printf("My Rank      : %d\n", my_rank);

    MPI_Finalize();
    return EXIT_SUCCESS;
}
```

```
mpicc mpi_hello.c -o mpi_hello
mpirun -np 4 ./mpi_hello
Num of Proc : 4
My Rank      : 3
Num of Proc : 4
My Rank      : 2
Num of Proc : 4
My Rank      : 0
Num of Proc : 4
My Rank      : 1
```


2.3 Minimum MPI Program

2.3.3 Hello World (C++)

```
#include <mpi.h>
#include <iostream>
#include <cstdlib>

int main(int argc, char *argv[])
{
    MPI_Init(&argc, &argv);

    int num_procs;
    int my_rank;

    MPI_Comm_size(MPI_COMM_WORLD, &num_procs);
    MPI_Comm_rank(MPI_COMM_WORLD, &my_rank);

    std::cout << "Num of Proc : " << num_procs <<
std::endl;
    std::cout << "My Rank      : " << my_rank << std::endl;

    MPI_Finalize();
    return EXIT_SUCCESS;
}
```

```
mpic++ mpi_hello.cpp -o mpi_hello
mpirun -np 4 ./mpi_hello
Num of Proc : 4
My Rank      : 3
Num of Proc : 4
My Rank      : 1
Num of Proc : 4
My Rank      : 0
Num of Proc : 4
My Rank      : 2
```

2.3 Minimum MPI Program

2.3.4 Hello World (Fortran)

```
program hello_mpi
  use mpi
  implicit none

  integer :: ierr, rank, size

  call MPI_Init(ierr)
  call MPI_Comm_rank(MPI_COMM_WORLD, rank, ierr)
  call MPI_Comm_size(MPI_COMM_WORLD, size, ierr)

  print *, "Num of Proc:", size
  print *, "My Rank:    ", rank

  call MPI_Finalize(ierr)
end program hello_mpi
```

```
mpif90 mpi_hello.f90 -o mpi_hello
mpirun -np 4 ./mpi_hello
Num of Proc: 4
My Rank:    2
Num of Proc: 4
My Rank:    0
Num of Proc: 4
My Rank:    3
Num of Proc: 4
My Rank:    1
```

2.4 Important Terms of MPI

2.4.1 Overview

- **Process:**
 - computing unit in parallel computing in MPI.
 - process num is determined by `mpirun -np`
- **Group:**
 - a set of processes that can communicate with each other.
- **Communicator:**
 - a group of processes that can communicate with each other.
- **Rank:**
 - unique identifier for each process in MPI.
 - Ranks are assigned from 0 to `num_procs - 1`.

2.4 Important Terms of MPI

2.4.2 Communicator Image

- Each process belongs to some **group**.
- A group is associated with a **communicator**.
- Each process in a communicator has a unique **rank**.

```
- Example:  
  - Process 0, 1, 2, 3 belong to a group.  
  - Communicator: MPI_COMM_WORLD  
  - Group: [P0, P1, P2, P3]  
  - Rank: 0, 1, 2, 3
```

```
MPI_COMM_WORLD (communicator)  
├─ Group: { P0, P1, P2, P3 }  
│   ├─ P0 (rank 0)  
│   ├─ P1 (rank 1)  
│   ├─ P2 (rank 2)  
│   └─ P3 (rank 3)
```

2.4 Important Terms of MPI

2.4.3 Communicator and Group

- **Group** is a set(list) of processes. It is just a list, it does not have any communication capability.
- **Communicator** contains a group and communication capability. It is a unit of communication.
- A group can create multiple communicators.
- A communicator cannot communicate with other communicators.

Q&A

- Q: Q: Can processes be added to an existing group?
- A: No. You cannot add processes to an existing group.

2.4 Important Terms of MPI

Attention

- A single process can belong to multiple groups and communications.
- In such cases, the process has a unique rank within each communicator.

```
MPI_COMM_WORLD
```

```
|- Group A: {0, 1, 2, 3, 4}
```

```
↓ Communicator B (created from Group A)
```

```
|- Group B {0, 1, 2, 3}
```

```
↓ Communicator C (created from Group A)
```

```
|- Group C {0, 1, 2, 4}
```

```
↓
```

- ○ 0, 1, 2 ↔ 3 (via Communicator B)
- ○ 0, 1, 2 ↔ 4 (via Communicator C)
- ○ 3 ↔ 4 (via Communicator A)
- x 3 ↔ 4 (via Communicator B and C)

- In the above example, process 0 belongs to three communicators:
 - MPI_COMM_WORLD (Group A), Communicator B (Group B), Communicator C (Group C)

2.4 Important Terms of MPI

2.4.4 MPI Functions for Communicator

- **MPI_COMM_WORLD:**
 - the default communicator that includes all processes.
 - all processes first belong to this communicator.
 - becomes the default communicator for most MPI functions.
- **MPI_Comm_rank():**
 - retrieves the rank of the calling process in specified communicator.
 - usually use MPI_COMM_WORLD as the communicator.
- **MPI_Comm_size():**
 - retrieves the number of processes in the specified communicator.
 - usually use MPI_COMM_WORLD as the communicator.
- **MPI_Comm_split():**
 - creates a new communicator by splitting the existing one based on a color and key.
 - in other words, create a new communicator with the same color processes.

2.4 Important Terms of MPI

- **MPI_Comm_rank** and **MPI_Comm_size**

```
int rank;  
MPI_Comm_rank(MPI_COMM_WORLD, &rank);  
printf("I am process %d\n", rank);
```

- **MPI_Comm_size**

```
int size;  
MPI_Comm_size(MPI_COMM_WORLD, &size);  
printf("There are %d processes\n", size);
```


2.4 Important Terms of MPI

- **MPI_Comm_split**

```
// ランクの3の剰余を基に color=0, 1, 2 に分ける
int color = rank % 3;

MPI_Comm new_comm;
MPI_Comm_split(MPI_COMM_WORLD, color, rank, &new_comm);

int new_rank, new_size;
MPI_Comm_rank(new_comm, &new_rank);
MPI_Comm_size(new_comm, &new_size);

printf("World Rank %d => Group %d, New Rank %d of %d\n",
       rank, color, new_rank, new_size);

MPI_Comm_free(&new_comm); // 新しいコミュニケータの解放
```

```
$ mpicc comm_split.c -o comm_split
$ mpirun -np 8 ./comm_split
```

```
World Rank 5 => Group 2, New Rank 1 of 2
World Rank 2 => Group 2, New Rank 0 of 2
World Rank 1 => Group 1, New Rank 0 of 3
World Rank 4 => Group 1, New Rank 1 of 3
World Rank 7 => Group 1, New Rank 2 of 3
World Rank 3 => Group 0, New Rank 1 of 3
World Rank 6 => Group 0, New Rank 2 of 3
World Rank 0 => Group 0, New Rank 0 of 3
```

2.4 Important Terms of MPI

World Rank	color(= rank % 3)	new group	new rank	new_size
0	0	{0, 3, 6}	0	3
1	1	{1, 4, 7}	0	3
2	2	{2, 5}	0	2
3	0	{0, 3, 6}	1	3
4	1	{1, 4, 7}	1	3
5	2	{2, 5}	1	2
6	0	{0, 3, 6}	2	3
7	1	{1, 4, 7}	2	3

2.5 Point-to-Point Communication

- Many MPI functions have the following signature
- **MPI_Send** sends data to a specific process.
- **MPI_Recv** receives data from a specific process.

```
MPI_Send(  
    void* data,           // data buffer address  
    int count,           // number of elements in the buffer  
    MPI_Datatype datatype, // data type of the elements  
    int destination,     // destination process rank  
    int tag,             // message tag (for filtering)  
    MPI_Comm communicator}; // communicator
```

```
MPI_Recv(  
    void* data,  
    int count,  
    MPI_Datatype datatype,  
    int source,  
    int tag,  
    MPI_Comm communicator,  
    MPI_Status* status);
```

2.5 Point-to-Point Communication

MPI Data Type	C Type
MPI_SHORT	short int
MPI_INT	int
MPI_LONG	long int
MPI_LONG_LONG	long long int
MPI_UNSIGNED_CHAR	unsigned char
MPI_UNSIGNED_SHORT	unsigned short int
MPI_UNSIGNED	unsigned int
MPI_UNSIGNED_LONG	unsigned long int
MPI_UNSIGNED_LONG_LONG	unsigned long long int
MPI_FLOAT	float
MPI_DOUBLE	double
MPI_LONG_DOUBLE	long double
MPI_BYTE	char

2.5 Point-to-Point Communication

send_data()

```
int send_data[10];
for (int i = 0; i < 10; i++)
    send_data[i] = i + 1;
int data_count = 10;

printf("Rank 0: Sending data.\n");
printf("send_data: [");
for (int i = 0; i < 10; i++)
    printf(" %d", send_data[i]);
printf(" ]\n");

MPI_Send((void*)send_data, data_count, MPI_INT,
1, 0, MPI_COMM_WORLD);
```

recv_data()

```
int data[10];
int data_count = 10;
MPI_Status st;

printf("Rank 1: Receiving data.\n");
MPI_Recv((void*)data, data_count, MPI_INT, 0, 0,
MPI_COMM_WORLD, &st);
printf("recv_data: [");
for (int i = 0; i < 10; i++)
    printf(" %d", data[i]);
printf(" ]\n");
```

```
$ mpicc send_recv.c -o mpi_send_recv
$ mpirun -np 2 ./mpi_send_recv
Rank 0: Sending data.
send_data: [ 1 2 3 4 5 6 7 8 9 10 ]
Rank 1: Receiving data.
recv_data: [ 1 2 3 4 5 6 7 8 9 10 ]
```

2.5 Point-to-Point Communication

- We can combine `send_data()` and `recv_data()` into a single program.
- The program can be run with `mpirun -np 2 ./mpi_send_recv2`

```
int data_count = 10;
int numbers[data_count];
if (world_rank == 0) {
    for (int i = 0; i < 10; i++) numbers[i] = i + 1;
    printf("Send Data:");
    for (int i = 0; i < 10; i++)
        printf(" %d", numbers[i]);
    printf("\n");
    MPI_Send((void *)&numbers, data_count, MPI_INT, 1, 0, MPI_COMM_WORLD);
} else if (world_rank == 1) {
    MPI_Recv(&numbers, data_count, MPI_INT, 0, 0, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
    printf("Received Data:");
    for (int i = 0; i < 10; i++)
        printf(" %d", numbers[i]);
    printf("\n");
}
```

2.5 Point-to-Point Communication

```
$ mpicc send_recv2.c -o mpi_send_recv2  
$ mpirun -np 2 ./mpi_send_recv2
```

```
Send Data: 1 2 3 4 5 6 7 8 9 10
```

```
Received Data: 1 2 3 4 5 6 7 8 9 10
```

3. Collective Communication

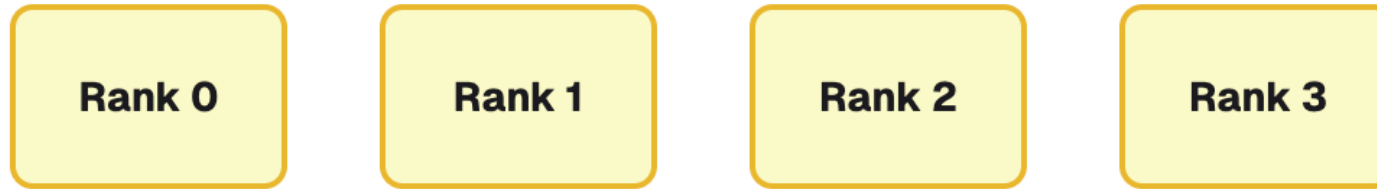
3.1 synchronization

- Collective communication is a communication method that involves all processes in a communicator.
- In collective communication, synchronization among all process is required.
- All process cannot proceed until all processes reach the same point.
- To achieve this, MPI provides several collective communication functions.

3.2 MPI_Barrier

- **MPI_Barrier** is a collective communication function that synchronizes all processes in a communicator.
- All processes must call `MPI_Barrier` to ensure that all processes reach the same point before proceeding.
- It is often used to ensure that all processes have completed their previous tasks before moving on to the next step.
- The most basic usage of `MPI_Barrier` is for precise time measurement.
- If you do not call `MPI_Barrier` in all processes, the program will block and cannot proceed.
- `MPI_Barrier(MPI_Comm communicator);`

MPI_Barrier - Synchronization



Wait until all processes reach the barrier

All processes synchronized → Proceed to next operation

3.2 MPI_Barrier

```
int rank, size;
MPI_Comm_rank(MPI_COMM_WORLD, &rank);
MPI_Comm_size(MPI_COMM_WORLD, &size);

printf("Rank %d: before barrier\n", rank);

sleep(rank);

MPI_Barrier(MPI_COMM_WORLD);
printf("Rank %d: after barrier\n", rank);
```

```
$ mpicc barrier.c -o barrier
$ mpirun -np 4 ./barrier
```

```
Rank 2: before barrier
Rank 3: before barrier
Rank 0: before barrier
Rank 1: before barrier
Rank 3: after barrier
Rank 1: after barrier
Rank 0: after barrier
Rank 2: after barrier
```

3.3 MPI_Bcast

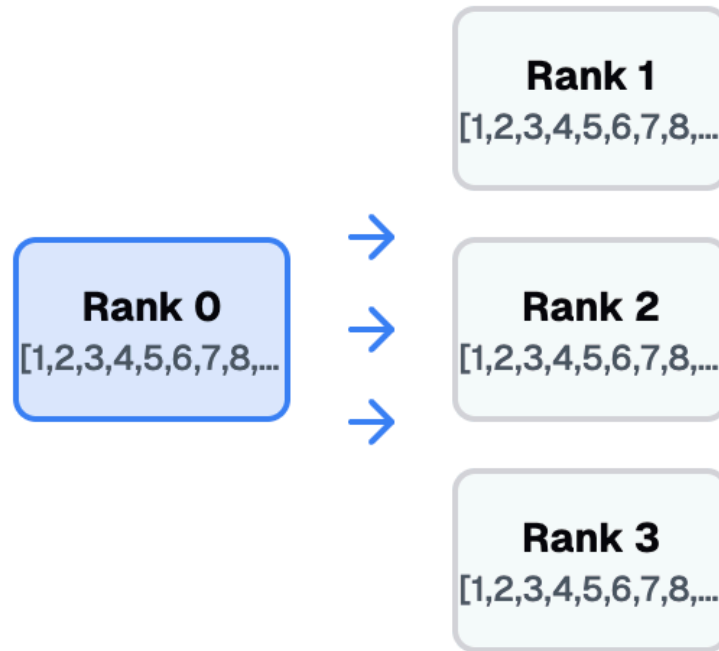
- **MPI_Bcast** is a collective communication function that broadcasts data from one process to all other processes in a communicator.
- It is used to distribute data from a root process to all other processes.
- All processes in the communicator must call `MPI_Bcast` with the same parameters.

```
MPI_Bcast(  
    void* data,           // data buffer address  
    int count,           // number of elements in the buffer  
    MPI_Datatype datatype, // data type of the elements  
    int root,            // rank of the root process  
    MPI_Comm communicator // communicator  
);
```

- Root Process: Sends the data to all other processes.
- Other Processes: Receive the data from the root process.

3.3 MPI_Bcast

MPI_Bcast - Broadcast



Send same data from Rank 0 to all processes

3.3 MPI_Bcast

```
int rank, size;
MPI_Comm_rank(MPI_COMM_WORLD, &rank);
MPI_Comm_size(MPI_COMM_WORLD, &size);

int data[10];
if (rank == 0) {
    for (int i = 0; i < 10; i++) {
        data[i] = i + 1;
    }
    printf("Rank 0: broadcasting data = [");
    for (int i = 0; i < 10; i++)
        printf(" %d", data[i]);
    printf(" ]\n");
}

MPI_Bcast(data, 10, MPI_INT, 0, MPI_COMM_WORLD);

printf("Rank %d: received data = [", rank);
for (int i = 0; i < 10; i++)
    printf(" %d", data[i]);
printf(" ]\n");
```

```
$ mpicc bcast.c -o bcast
$ mpirun -np 4 ./bcast
```

```
Rank 0: broadcasting data = [ 1 2 3 4 5 6 7 8 9
10 ]
Rank 0: received data = [ 1 2 3 4 5 6 7 8 9 10 ]
Rank 2: received data = [ 1 2 3 4 5 6 7 8 9 10 ]
Rank 3: received data = [ 1 2 3 4 5 6 7 8 9 10 ]
Rank 1: received data = [ 1 2 3 4 5 6 7 8 9 10 ]
```

3.3 MPI_Bcast

- We can implement MPI_Bcast wrapper using MPI_Send and MPI_Recv.

```
void my_bcast(void* data, int count, MPI_Datatype datatype, int root, MPI_Comm communicator) {
    int world_rank;
    MPI_Comm_rank(communicator, &world_rank);
    int world_size;
    MPI_Comm_size(communicator, &world_size);

    if (world_rank == root) {
        for (int i = 0; i < world_size; i++) {
            if (i != world_rank) {
                MPI_Send(data, count, datatype, i, 0, communicator);
            }
        }
    } else {
        MPI_Recv(data, count, datatype, root, 0, communicator, MPI_STATUS_IGNORE);
    }
}
```

- Q: Is this equivalent to MPI_Bcast?

3.3 MPI_Bcast

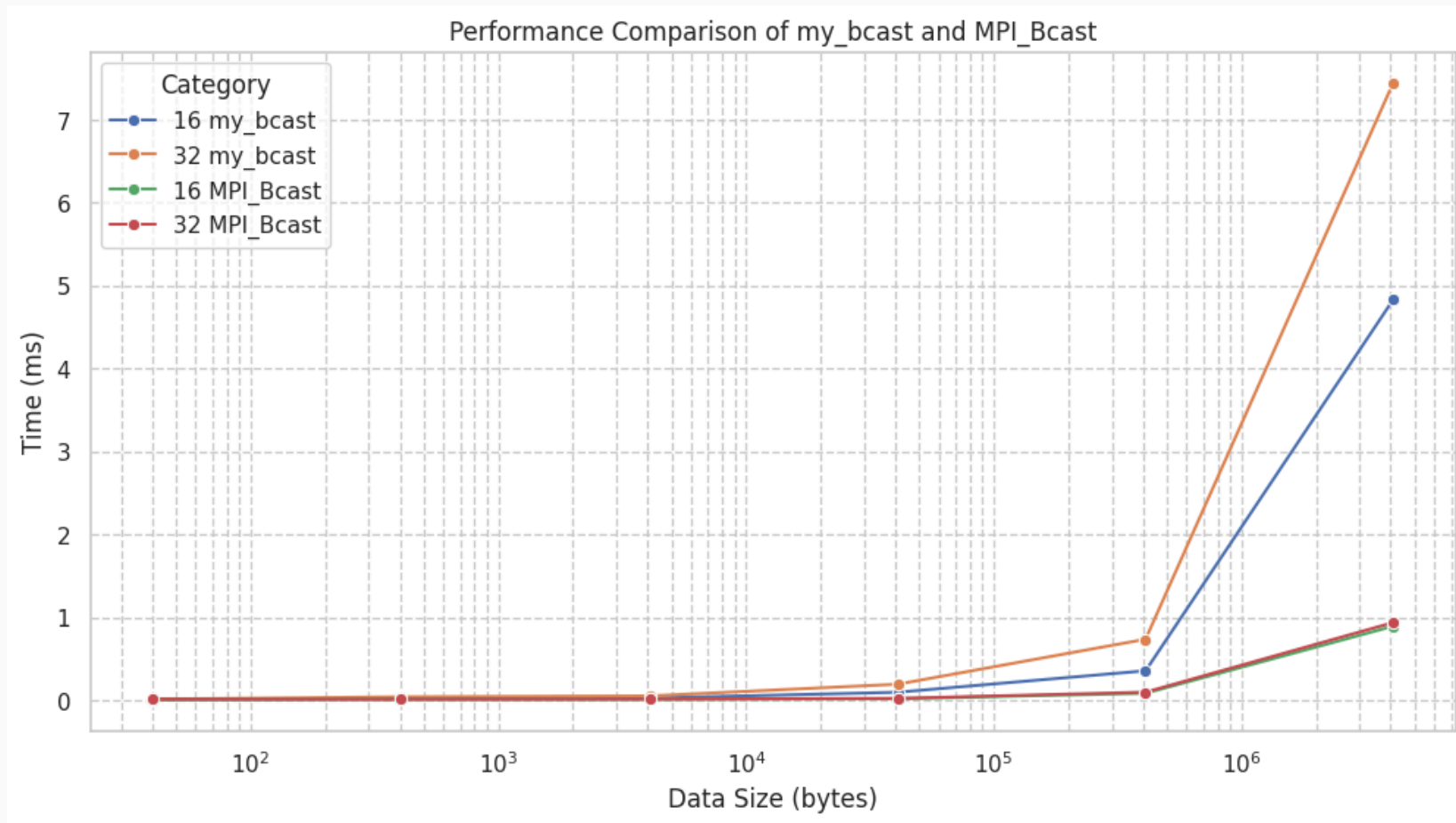
- A: **No**, it is less efficient than MPI_Bcast.
- This implementation has only one network communication link. The process with root rank sends data to all other processes one by one.
- MPI_Bcast uses Tree-based broadcast algorithm.
 1. The root process sends data to process 1.
 2. The root process sends data to process 2, process 1 sends data to process 3.
 3. The root process sends data to process 4, process 1 sends data to process 5, process 2 sends data to process 6, process 3 sends data to process 7.
- ...

3.3 MPI_Bcast

- Comparison of MPI_Bcast and my_bcast
- Average time of 10 trials

Procs	Data Size	my_bcast (ms)	MPI_Bcast (ms)
16	40	0.008	0.009
16	400	0.022	0.009
16	4k	0.026	0.010
16	40k	0.096	0.018
16	400k	0.355	0.084
16	4000k	4.832	0.893
32	40	0.012	0.012
32	400	0.041	0.011
32	4k	0.052	0.013
32	40k	0.193	0.022
32	400k	0.735	0.097
32	4000k	7.447	0.937

3.3 MPI_Bcast



3.4 MPI_Scatter

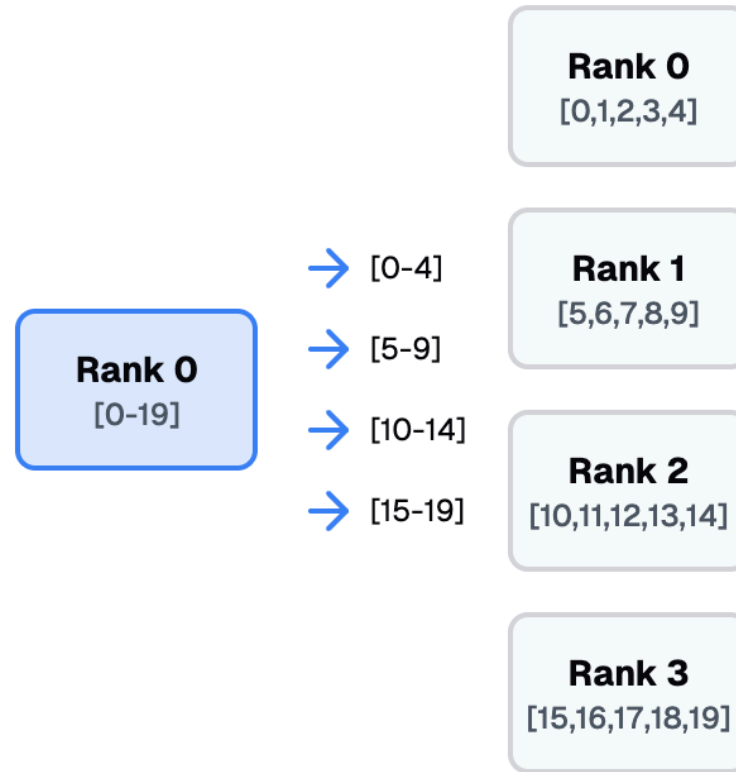
- **MPI_Scatter** is a collective communication function that distributes data from a root process to all other processes in a communicator.
- MPI_Bcast sends the same data to all processes, while MPI_Scatter sends different chunks of data to each process.

```
MPI_Scatter(  
    void* send_data,           // data buffer address to send  
    int send_count,           // number of elements to send to each process  
    MPI_Datatype send_datatype, // data type of the elements to send  
    void* recv_data,          // data buffer address to receive  
    int recv_count,           // number of elements to receive  
    MPI_Datatype recv_datatype, // data type of the elements to receive  
    int root,                 // rank of the root process  
    MPI_Comm communicator.    // communicator  
);
```

- send_count: the number of elements to send to each process.
- recv_count: the number of elements to receive from each process.

3.4 MPI_Scatter

MPI_Scatter - Data Distribution



Distribute data from Rank 0 to each process

3.4 MPI_Scatter

```
#define TOTAL_DATA 20
...
MPI_Comm_rank(MPI_COMM_WORLD, &rank);
MPI_Comm_size(MPI_COMM_WORLD, &size);

int send_data[TOTAL_DATA];
int recv_count = TOTAL_DATA / size;
int recv_data[recv_count];

if (rank == 0) {
    for (int i = 0; i < TOTAL_DATA; i++)
        send_data[i] = i;
    printf("Rank 0: Scattering data...\n");
}

MPI_Scatter(send_data, recv_count, MPI_INT,
            recv_data, recv_count, MPI_INT,
            0, MPI_COMM_WORLD);

printf("Rank %d received:", rank);
for (int i = 0; i < recv_count; i++)
    printf(" %d", recv_data[i]);
printf("\n");
```

```
// send_data == 5
```

```
$ mpicc scatter.c -o scatter
```

```
$ mpirun -np 4 ./scatter
```

```
Rank 0: Scattering data...
```

```
Rank 0 received: 0 1 2 3 4
```

```
Rank 1 received: 5 6 7 8 9
```

```
Rank 2 received: 10 11 12 13 14
```

```
Rank 3 received: 14 15 16 17 18
```

if `send_data` cannot divide by `size`, the last process will receive the remaining data.

```
// send_data == recv_data == 6
```

```
Rank 0 received: 0 1 2 3 4 5
```

```
Rank 1 received: 6 7 8 9 10 11
```

```
Rank 2 received: 12 13 14 15 16 17
```

```
Rank 3 received: 18 19 20 -875497504 65535 20
```

3.4 MPI_Scatter

- What happens if `send_count` is not the same as `recv_count`?

```
int send_data[TOTAL_DATA];
int send_count = TOTAL_DATA / size;
int recv_count = TOTAL_DATA / size;
int recv_data[recv_count];

if (rank == 0) {
    for (int i = 0; i < TOTAL_DATA; i++) {
        send_data[i] = i;
    }
    printf("Rank 0: Scattering data...\n");
}

if (rank == 3) {
    recv_count = 4;
}

MPI_Scatter(send_data, send_count, MPI_INT,
            recv_data, recv_count, MPI_INT,
            0, MPI_COMM_WORLD);
```

3.4 MPI_Scatter

- A: **It will cause an error.**

```
$ mpicc scatter2.c -o scatter2
$ mpirun -np 4 ./scatter2
```

```
Rank 0: Scattering data...
```

```
Rank 0 received: 0 1 2 3 4
```

```
Rank 1 received: 5 6 7 8 9
```

```
Rank 2 received: 10 11 12 13 14
```

```
[miyabi-g3:1978858] *** An error occurred in MPI_Scatter
```

```
[miyabi-g3:1978858] *** reported by process [1190789121,3]
```

```
[miyabi-g3:1978858] *** on communicator MPI_COMM_WORLD
```

```
[miyabi-g3:1978858] *** MPI_ERR_TRUNCATE: message truncated
```

```
[miyabi-g3:1978858] *** MPI_ERRORTS ARE_FATAL (processes in this communicator will now abort,
```

```
[miyabi-g3:1978858] *** and potentially your MPI job)
```


3.5 MPI_Gather

- **MPI_Gather** is a collective communication function that collects data from all processes in a communicator and sends it to a root process.
- It is the reverse operation of `MPI_Scatter`.
- This is used in parallel sorting, parallel searching, and other parallel algorithms.

```
MPI_Gather(  
    void* send_data,           // data buffer address to send  
    int send_count,           // number of elements to send from each process  
    MPI_Datatype send_datatype, // data type of the elements to send  
    void* recv_data,          // data buffer address to receive  
    int recv_count,           // number of elements to receive from each process  
    MPI_Datatype recv_datatype, // data type of the elements to receive  
    int root,                 // rank of the root process  
    MPI_Comm communicator     // communicator  
);
```

- Except for the root process, pass `NULL` for `recv_data` is allowed.
- `recv_count` is the number of elements to receive from each process, not the total number of elements.

3.5 MPI_Gather

MPI_Gather - Data Collection



Collect data from each process to Rank 0

3.5 MPI_Gather

```
#define ITEMS_PER_PROC 2
...
int send_data[ITEMS_PER_PROC];
send_data[0] = rank * 2;
send_data[1] = rank * 2 + 1;

int recv_data[ITEMS_PER_PROC * size];

MPI_Gather(send_data, ITEMS_PER_PROC, MPI_INT,
           recv_data, ITEMS_PER_PROC, MPI_INT,
           0, MPI_COMM_WORLD);

if (rank == 0) {
    printf("Rank 0 gathered data: ");
    for (int i = 0; i < ITEMS_PER_PROC * size;
        i++)
        printf("%d ", recv_data[i]);
    printf("\n");
} else {
    printf("Rank %d sent data: %d %d\n", rank,
           send_data[0], send_data[1]);
}
```

```
$ mpicc gather.c -o gather
$ mpirun -np 4 ./gather
```

```
Rank 1 sent data: 2 3
Rank 3 sent data: 6 7
Rank 2 sent data: 4 5
Rank 0 gathered data: 0 1 2 3 4 5 6 7
```

3.6 MPI_Allgather

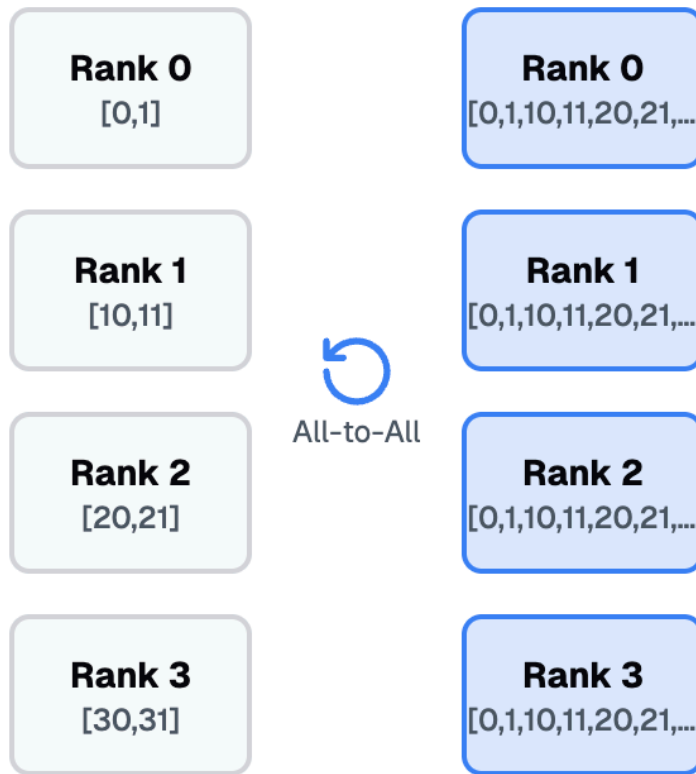
- MPI_Scatter and MPI_Gather conduct many-to-one or one-to-many communication.
- It is useful if you send data from multiple processes to multiple processes.
- **MPI_Allgather** is a collective communication function that collects data from all processes in a communicator and sends it to all other processes.
- It is like first MPI_Gather and then MPI_Bcast. Collect data by process rank order.

```
MPI_Allgather(  
    void* send_data,           // data buffer address to send  
    int send_count,           // number of elements to send from each process  
    MPI_Datatype send_datatype, // data type of the elements to send  
    void* recv_data,          // data buffer address to receive  
    int recv_count,           // number of elements to receive from each process  
    MPI_Datatype recv_datatype, // data type of the elements to receive  
    MPI_Comm communicator     // communicator  
);
```

- MPI_Allgather does not have a root process parameter.

3.6 MPI_Allgather

MPI_AllGather - All-to-All Collection



Collect and distribute data to all processes

3.6 MPI_Allgather

```
#define ITEMS_PER_PROC 2
...
int rank, size;
MPI_Comm_rank(MPI_COMM_WORLD, &rank);
MPI_Comm_size(MPI_COMM_WORLD, &size);

int send_data[ITEMS_PER_PROC];
send_data[0] = rank * 10;
send_data[1] = rank * 10 + 1;

int recv_data[ITEMS_PER_PROC * size];

MPI_Allgather(send_data, ITEMS_PER_PROC,
MPI_INT,
               recv_data, ITEMS_PER_PROC,
MPI_INT,
               MPI_COMM_WORLD);

printf("Rank %d received:", rank);
for (int i = 0; i < ITEMS_PER_PROC * size; i++)
    printf(" %d", recv_data[i]);
printf("\n");
```

```
$ mpicc allgather.c -o allgather
$ mpirun -np 4 ./allgather
```

```
Rank 2 received: 0 1 10 11 20 21 30 31
Rank 3 received: 0 1 10 11 20 21 30 31
Rank 0 received: 0 1 10 11 20 21 30 31
Rank 1 received: 0 1 10 11 20 21 30 31
```

3.7 MPI_Reduce

- `reduce` is a basic concept in functional programming. It transforms a set of numbers into a smaller set of numbers.
 - `reduce([1, 2, 3, 4, 5], sum) = 15`
 - `reduce([1, 2, 3, 4, 5], multiply) = 120`
- Collect distributed data and apply a reduction operation is a tough task. However, MPI provides a simple interface to do this.
- **MPI_Reduce** is a collective communication function that collects data from all processes in a communicator to the root process, and applies a reduction operation to the data.

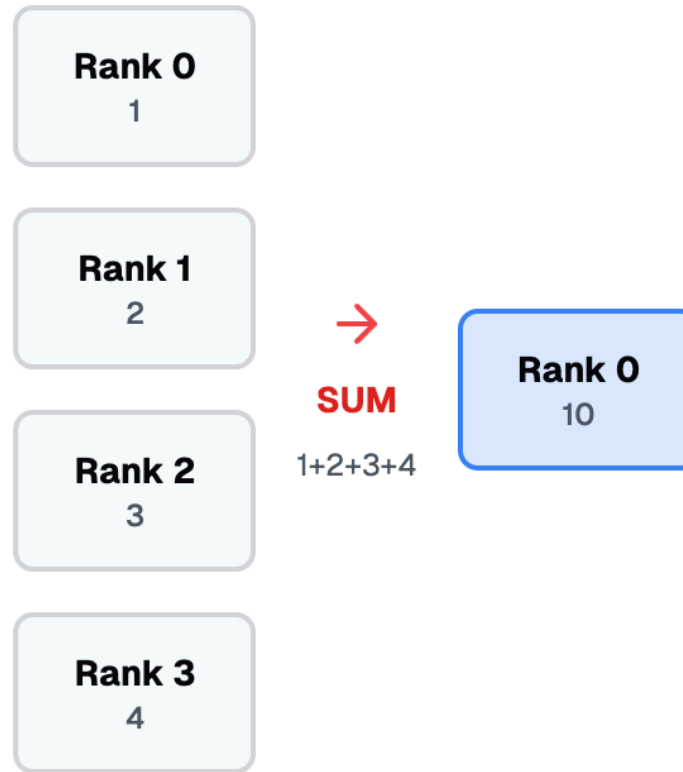
3.7 MPI_Reduce

```
MPI_Reduce(  
    void* send_data,          // data buffer address to send  
    void* recv_data,         // data buffer address to receive  
    int count,               // number of elements to send from each process  
    MPI_Datatype datatype,   // data type of the elements to send  
    MPI_Op op,               // reduction operation to apply  
    int root,                // rank of the root process  
    MPI_Comm communicator    // communicator  
);
```

- MPI Reduction Operations:
 - **MPI_MAX** - maximum value
 - **MPI_MIN** - minimum value
 - **MPI_SUM** - sum of values
 - **MPI_PROD** - product of values
 - **MPI_LAND** - logical AND of values
 - **MPI_LOR** - logical OR of values
 - **MPI_BAND** - bitwise AND of values
 - **MPI_BOR** - bitwise OR of values
 - **MPI_MAXLOC** - maximum value and its rank

3.7 MPI_Reduce

MPI_Reduce - Reduction Operation



Perform operation on values and store result in Rank 0

3.7 MPI_Reduce

```
int rank, size;
MPI_Comm_rank(MPI_COMM_WORLD, &rank);
MPI_Comm_size(MPI_COMM_WORLD, &size);

int value = rank + 1;
int sum;

MPI_Reduce(&value, &sum, 1, MPI_INT, MPI_SUM, 0,
MPI_COMM_WORLD);

double average = (double)sum / size;
if (rank == 0) {
    printf("Rank %d: sum = %d, avg = %.2f\n",
rank, sum, average);
} else {
    printf("Rank %d: sum = %d, avg = %.2f\n",
rank, sum, average);
}
```

```
$ mpicc reduce.c -o reduce
$ mpirun -np 4 ./reduce
```

```
Rank 1: sum = 4197236, avg = 1049309.00
Rank 0: sum = 10, avg = 2.50
Rank 2: sum = 4197236, avg = 1049309.00
Rank 3: sum = 4197236, avg = 1049309.00
```

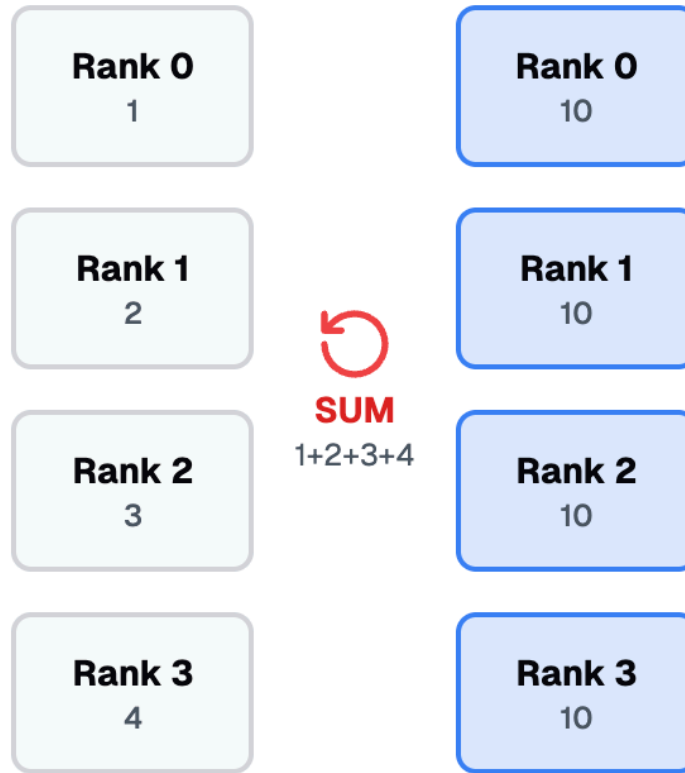
3.8 MPI_Allreduce

- **MPI_Allreduce** is a collective communication function that collects data from all processes in a communicator, applies a reduction operation to the data, and distributes the result to all processes.
- It is similar to `MPI_Reduce`, but the result is available to all processes, not just the root process.
- This is useful when all processes need to know the result of the reduction operation.

```
MPI_Allreduce(  
    void* send_data,          // data buffer address to send  
    void* recv_data,         // data buffer address to receive  
    int count,               // number of elements to send from each process  
    MPI_Datatype datatype,   // data type of the elements to send  
    MPI_Op op,               // reduction operation to apply  
    MPI_Comm communicator    // communicator  
);
```

3.8 MPI_Allreduce

MPI_AllReduce - All-Reduction Operation



Perform operation on values and store result in all processes

3.8 MPI_Allreduce

```
int rank, size;
MPI_Comm_rank(MPI_COMM_WORLD, &rank);
MPI_Comm_size(MPI_COMM_WORLD, &size);

int value = rank + 1;
int total_sum = 0;

MPI_Allreduce(&value, &total_sum, 1, MPI_INT,
MPI_SUM, MPI_COMM_WORLD);

double average = (double)total_sum / size;
printf("Rank %d: total sum = %d, average =
%.2f\n", rank, total_sum, average);
```

```
$ mpicc allreduce.c -o allreduce
$ mpirun -np 4 ./allreduce
```

```
Rank 2: total sum = 10, average = 2.50
Rank 1: total sum = 10, average = 2.50
Rank 0: total sum = 10, average = 2.50
Rank 3: total sum = 10, average = 2.50
```

4. Application Example

4.1 How to Run Multi-node program on Miyabi

- The following is an example of how to run the Page Rank calculation on a multi-node cluster.
- When you submit a job to Miyabi-G cluster, you can specify the number of node and the number of process per node using `-l select={num_nodes}:mpiprocs={num_procs_per_node}` option.
- When you run a MPI program on a multi-node cluster, you need to specify the number of processes using `mpirun -np {num_procs}` option.

Example `run.sh` script:

```
#!/bin/bash
#PBS -q debug-g
#PBS -l select=16:mpiprocs=4
#PBS -W group_list=gc64
#PBS -o latest_result.txt
#PBS -j oe
...
mpirun -np 64 ./naive/pagerank_naive
```

4.1 How to Run Multi-node program on Miyabi

- If you want to run a MPI program on another process-per-node than `mpiprocs`, you can use `--host` option or `--hostfile` option.
- `--host` option allows you to specify the hostnames and the number of processes per host.
- `--hostfile` option allows you to specify a file that contains the hostnames and the number of processes per host.

Example `--host` option:

```
$ mpirun -np 4 --host node1:2,node2:2 ./your_mpi_program  
→node1 and node2 each run 2 processes.
```

Example `--hostfile` option: `hostfile.txt`:

```
node1 slots=2  
node2 slots=2/
```

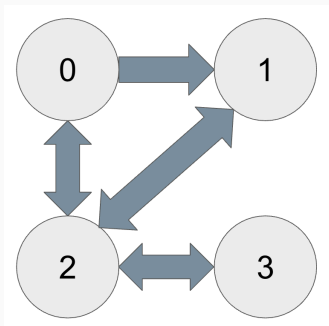
```
$ mpirun -np 4 --hostfile hostfile.txt ./your_mpi_program
```


4.2 Page Rank Calculation

- For a simple application example of MPI, I will introduce the Page Rank calculation.
- **Page Rank** is an algorithm used by Google to rank web pages in search results.
- The concept of Page Rank is that “The page is important if it is linked from many other important pages.”

$$x^{\{(k+1)\}} = dMx^{\{(k)\}} + (1 - d)v$$

- **M** is the link matrix, each element represents the link between pages,
- **x** is the Page Rank vector, each element represents the Page Rank of a page,
- **d** is the damping factor, usually set to 0.85,
- **v** is the uniform distribution vector (initial vector).



$$d = 0.85, M = \begin{pmatrix} 0 & 0 & \frac{1}{3} & 0 \\ \frac{1}{2} & 0 & \frac{1}{3} & 0 \\ \frac{1}{2} & 1 & \frac{1}{3} & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix}, v = \begin{pmatrix} \frac{1}{4} \\ \frac{1}{4} \\ \frac{1}{4} \\ \frac{1}{4} \end{pmatrix}$$

4.2 Page Rank Calculation

$$M = \begin{pmatrix} 0 & 0 & \frac{1}{3} & 0 \\ \frac{1}{2} & 0 & \frac{1}{3} & 0 \\ \frac{1}{2} & 1 & \frac{1}{3} & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix}$$

- The sum of each column in matrix M is 1 because M represents a transition probability matrix, where each column corresponds to a node and the entries indicate the probability of moving from that node to others
- The Page Rank vector is updated iteratively until convergence.

$$x^{\{1\}} = dMx^{\{0\}} + (1 - d)v$$

$$x^{\{2\}} = dMx^{\{1\}} + (1 - d)v$$

$$x^{\{3\}} = dMx^{\{2\}} + (1 - d)v$$

...

4.2 Page Rank Calculation

How to conduct Page Rank calculation in parallel?

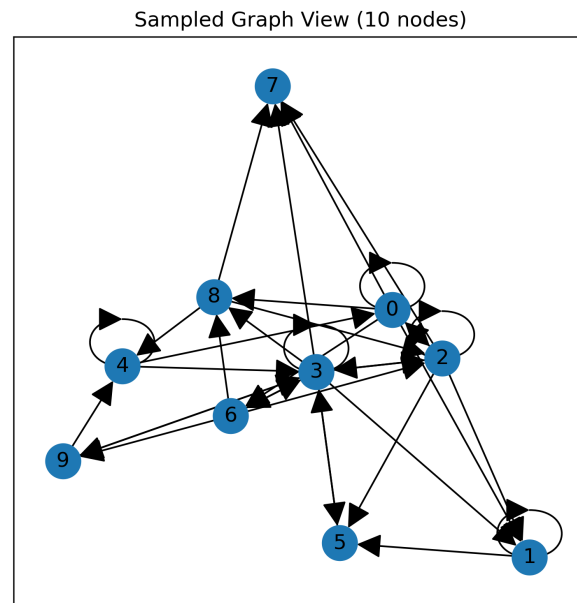
- Split the link matrix M into submatrices, each assigned to a process.
- Each process calculates the Page Rank for its submatrix.
- Use `MPI_Allreduce` to combine the results from all processes.
- Repeat until convergence.
- The following code is a simple example of Page Rank calculation using MPI.

Run Step:

1. Generate a graph data file.
2. (option) Visualize the graph data.
3. Run the Page Rank calculation.

4.2 Page Rank Calculation

```
$ cd tutorial/pagerank/  
  
// graph data generation  
$ ./create_venv.sh  
$ python3 preprocess/generate_graph.py  
$ ls data/  
  
// graph data visualization  
$ python3 preprocess/visualize.py  
  
// Page Rank calculation  
$ cd tutorial/pagerank/naive  
$ ./run.sh
```



```
$ ./run.sh
```

Final PageRank:

	0	1	2	3	4	5	6	7	8	9
PR:	0.0490	0.0817	0.0678	0.1453	0.0616	0.0748	0.0374	0.0598	0.0453	0.0384

4.2 Page Rank Calculation

- We divide the M matrix into submatrices, each assigned to a process.
- Each process only have a submatrix of M and a subvector of $x_k(*x)$, but the whole vector of $x_{k+1}(*new_x_full)$.

```
typedef struct {
    int n_nodes;
    int n_edges;
    int *out_degree;
    double **M;
} Graph;
...
Graph g;
int n = g.n_nodes;

// Data partitioning (row-wise)
int rows_per_proc = n / size;
int remainder = n % size;
int my_rows = rows_per_proc + (rank < remainder ?
1 : 0);
int my_start = rank * rows_per_proc + (rank <
remainder ? rank : remainder);
```

```
double *x = malloc(n * sizeof(double));
double *new_x_local = malloc(my_rows *
sizeof(double));

for (int i = 0; i < n; i++) x[i] = 1.0 / n;

for (int i = 0; i < size; i++) {
    recvcnts[i] = rows_per_proc + (i <
remainder ? 1 : 0);
    displs[i] = i * rows_per_proc + (i <
remainder ? i : remainder);
}

double *new_x_full = malloc(n * sizeof(double));
```

4.2 Page Rank Calculation

- `MPI_Allgatherv` is different from `MPI_Allgather` in that it allows each process to send a different number of elements.

```
for (int iter = 0; iter < MAX_ITER; iter++) {
    for (int i = 0; i < my_rows; i++) {
        int global_i = my_start + i;
        new_x_local[i] = 0.0;
        for (int j = 0; j < n; j++) {
            new_x_local[i] += g.M[global_i][j] *
x[j];
        }
        new_x_local[i] = DAMPING * new_x_local[i]
+ (1.0 - DAMPING) / n;
    }

    MPI_Allgatherv(new_x_local, my_rows,
MPI_DOUBLE,
                    new_x_full, recvcnts,
displs, MPI_DOUBLE, MPI_COMM_WORLD);
```

```
// チェック用: rank 0 が diff を計算
double diff = 0.0;
if (rank == 0) {
    for (int i = 0; i < n; i++) {
        diff += fabs(new_x_full[i] - x[i]);
    }
}

// 全rankにdiffをブロードキャスト (終了判定共有)
MPI_Bcast(&diff, 1, MPI_DOUBLE, 0,
MPI_COMM_WORLD);
if (diff < TOL) break;

// x を更新
for (int i = 0; i < n; i++)
    x[i] = new_x_full[i];
}
```

5. References

5. References

- Referenced in this tutorial:
 - [MPI「超」入門（C言語編） - 東京大学情報基盤センター](#)
 - [並列プログラミング入門](#)
 - [MPI Tutorial](#)
- For Advanced Users:
 - [行列-行列積\(2\) 非同期通信 - 東京大学情報基盤センター](#)
 - [Open MPI GitHub Repository](#)

6. Appendix

6.1 How to use Page Rank when requested a search query

- Page Rank is just a algorithm to rank web pages.
- The requested search query is just a string, so we cannot use Page Rank directly.
- One approach is:
 1. we get related web pages according to DF-IDF or BM25 algorithm
 2. then we apply Page Rank to rerank the related web pages.

6.2 Page Rank Full Code

graph_loader.h

```
#ifndef GRAPH_LOADER_H
#define GRAPH_LOADER_H

typedef struct {
    int n_nodes;
    int n_edges;
    int *out_degree;
    double **M;
} Graph;

int read_graph(const char *filename, Graph *g);
void free_graph(Graph *g);

#endif
```

graph_loader.c

```
#include "graph_loader.h"
#include <stdio.h>
#include <stdlib.h>

int read_graph(const char *filename, Graph *g) {
    FILE *fp = fopen(filename, "r");
    if (!fp) {
        perror("fopen");
        return -1;
    }

    if (fscanf(fp, "%d", &g->n_nodes) != 1) {
        fprintf(stderr, "Failed to read node count\n");
        return -1;
    }
}
```

6.2 Page Rank Full Code

```
    g->out_degree = calloc(g->n_nodes,
sizeof(int));
    g->M = malloc(g->n_nodes * sizeof(double
*));
    for (int i = 0; i < g->n_nodes; i++) {
        g->M[i] = calloc(g->n_nodes,
sizeof(double));
    }

    int from, to;
    g->n_edges = 0;
    while (fscanf(fp, "%d %d", &from, &to) == 2)
    {
        g->M[to][from] += 1.0;
        g->out_degree[from]++;
        g->n_edges++;
    }
    fclose(fp);
```

```
    // 正規化 : M[to][from] /= out_degree[from]
    for (int i = 0; i < g->n_nodes; i++) {
        for (int j = 0; j < g->n_nodes; j++) {
            if (g->out_degree[j] > 0) {
                g->M[i][j] /= g->out_degree[j];
            }
        }
    }

    return 0;
}

void free_graph(Graph *g) {
    for (int i = 0; i < g->n_nodes; i++) {
        free(g->M[i]);
    }
    free(g->M);
    free(g->out_degree);
}
```

6.2 Page Rank Full Code

main.c

```
#include <stdio.h>
#include <stdlib.h>
#include <math.h>
#include <mpi.h>
#include "graph_loader.h"

#define DAMPING 0.85
#define TOL 1e-6
#define MAX_ITER 100

int main(int argc, char **argv) {
    MPI_Init(&argc, &argv);

    int rank, size;
    MPI_Comm_rank(MPI_COMM_WORLD, &rank);
    MPI_Comm_size(MPI_COMM_WORLD, &size);
```

```
    if (argc < 2) {
        if (rank == 0) {
            fprintf(stderr, "Usage: %s <n_node>\n",
                argv[0]);
        }
        MPI_Abort(MPI_COMM_WORLD, 1);
    }
    int n_node = atoi(argv[1]);
    char graph_path[256], pagerank_path[256];
    snprintf(graph_path, sizeof(graph_path), "data/
%d/graph.txt", n_node);
    snprintf(pagerank_path, sizeof(pagerank_path),
        "data/%d/pagerank.txt", n_node);

    Graph g;

    if (read_graph(graph_path, &g) != 0) {
        MPI_Abort(MPI_COMM_WORLD, 1);
    }
```

6.2 Page Rank Full Code

```
int n = g.n_nodes;

// データ分割 (行分割)
int rows_per_proc = n / size;
int remainder = n % size;
int my_rows = rows_per_proc + (rank <
remainder ? 1 : 0);
int my_start = rank * rows_per_proc + (rank <
remainder ? rank : remainder);

// PageRankベクトルxと各プロセスが担当する新しいベクトル
new_x_local(len(my_rows)次元)を確保
double *x = malloc(n * sizeof(double));
double *new_x_local = malloc(my_rows *
sizeof(double));

// 全プロセスで保持するの新しいPageRankベクトル
new_x_full(n次元)を確保
double *new_x_full = malloc(n * sizeof(double));
int *recvcnts = malloc(size * sizeof(int));
int *displs = malloc(size * sizeof(int));

for (int i = 0; i < n; i++) x[i] = 1.0 / n;
```

```
for (int i = 0; i < size; i++) {
    recvcnts[i] = rows_per_proc + (i <
remainder ? 1 : 0);
    displs[i] = i * rows_per_proc + (i <
remainder ? i : remainder);
}
if (rank == 0) {
    printf("Starting PageRank computation with
%d processes...\n", size);
}
printf("[Rank %d] Local rows: %d, Start index:
%d\n", rank, my_rows, my_start);
double t_start = MPI_Wtime();

for (int iter = 0; iter < MAX_ITER; iter++) {
    for (int i = 0; i < my_rows; i++) {
        int global_i = my_start + i;
        new_x_local[i] = 0.0;
        for (int j = 0; j < n; j++) {
            new_x_local[i] += g.M[global_i][j] *
x[j];
        }
        new_x_local[i] = DAMPING *
new_x_local[i] + (1.0 - DAMPING) / n;
    }
}
```

6.2 Page Rank Full Code

```
MPI_Allgatherv(new_x_local, my_rows,
MPI_DOUBLE,
                new_x_full, recvcunts,
displs, MPI_DOUBLE, MPI_COMM_WORLD);
// チェック用: rank 0 が diff を計算
double diff = 0.0;
if (rank == 0) {
    for (int i = 0; i < n; i++) {
        diff += fabs(new_x_full[i] - x[i]);
    }
    printf("[Iter %d] diff = %f\n", iter +
1, diff);
}

// 全rankにdiffをブロードキャスト (終了判定共有)
MPI_Bcast(&diff, 1, MPI_DOUBLE, 0,
MPI_COMM_WORLD);
if (diff < TOL) break;

// x を更新
for (int i = 0; i < n; i++) x[i] =
new_x_full[i];
}
```

```
double t_end = MPI_Wtime();
if (rank == 0) {
    printf("Main Loop Time: %.6f seconds\n",
t_end - t_start);
}

free(new_x_full);
free(recvcunts);
free(displs);

// 結果表示
if (rank == 0) {
    printf("\n\nFinal PageRank:\n");
    for (int i = 0; i < n; i++) {
        printf("Node %6d: %f\n", i, x[i]);
    }
    printf("\n");
}
```

6.2 Page Rank Full Code

```
// 結果をファイルに書き出す
if (rank == 0) {
    FILE *fp = fopen(pagerank_path, "w");
    if (fp == NULL) {
        perror("Failed to open result file");
        MPI_Abort(MPI_COMM_WORLD, 1);
    }
    fprintf(fp, "%d\n", n);
    for (int i = 0; i < n; i++) {
        fprintf(fp, "%d %f\n", i, x[i]);
    }
    fclose(fp);
    printf("PageRank results written to %s\n", pagerank_path);
}

free(x);
free(new_x_local);
free_graph(&g);

MPI_Finalize();
return 0;
}
```


6.3 Example code of Communication and Group

6.3 Example code of Communication and Group

- Example code of p.21 slide

```
MPI_Group world_group;
MPI_Comm_group(MPI_COMM_WORLD, &world_group);
// === Group A: {0,1,2,3} ===
int ranks_a[] = {0, 1, 2, 3};
MPI_Group group_a;
MPI_Group_incl(world_group, 4, ranks_a,
&group_a);
MPI_Comm comm_a;
MPI_Comm_create_group(MPI_COMM_WORLD, group_a, 0,
&comm_a);
// === Group B: {0,1,2,4} ===
int ranks_b[] = {0, 1, 2, 4};
MPI_Group group_b;
MPI_Group_incl(world_group, 4, ranks_b,
&group_b);
MPI_Comm comm_b;
MPI_Comm_create_group(MPI_COMM_WORLD, group_b, 1,
&comm_b);
```

```
// Rank 0 in each comm sends a message to others
if (comm_a != MPI_COMM_NULL) {
    int comm_a_rank;
    MPI_Comm_rank(comm_a, &comm_a_rank);
    if (comm_a_rank == 0) {
        int msg = 100;
        for (int i = 1; i < 4; i++)
            MPI_Send(&msg, 1, MPI_INT, i, 0,
comm_a);
    } else {
        int recv;
        MPI_Recv(&recv, 1, MPI_INT, 0, 0, comm_a,
MPI_STATUS_IGNORE);
        printf("comm_a: World Rank %d received %d
from Rank 0\n", world_rank, recv);
    }
}
```

6.3 Example code of Communication and Group

```
if (comm_b != MPI_COMM_NULL) {
    int comm_b_rank;
    MPI_Comm_rank(comm_b, &comm_b_rank);
    if (comm_b_rank == 0) {
        int msg = 200;
        for (int i = 1; i < 4; i++)
            MPI_Send(&msg, 1, MPI_INT, i, 0,
comm_b);
    } else {
        int rcv;
        MPI_Recv(&rcv, 1, MPI_INT, 0, 0, comm_b,
MPI_STATUS_IGNORE);
        printf("comm_b: World Rank %d received %d
from Rank 0\n", world_rank, rcv);
    }
}
```

```
// Rank 3 sends a message to Rank 4 via (failed)
if (world_rank == 3) {
    int msg = 999;
    int rc = MPI_Send(&msg, 1, MPI_INT, 4, 99,
MPI_COMM_WORLD);
    printf("Rank 3 tried to send to Rank 4 via
MPI_COMM_WORLD (rc=%d)\n", rc);
}

if (world_rank == 4) {
    int rcv, flag;
    MPI_Status status;
    MPI_Iprobe(3, 99, MPI_COMM_WORLD, &flag,
&status);
    if (flag) {
        MPI_Recv(&rcv, 1, MPI_INT, 3, 99,
MPI_COMM_WORLD, MPI_STATUS_IGNORE);
        printf("Rank 4 received unexpected
message from Rank 3: %d\n", rcv);
    } else {
        printf("Rank 4: No message from Rank 3
(as expected)\n");
    }
}
```

6.3 Example code of Communication and Group

```
// Cleanup
if (comm_a != MPI_COMM_NULL)
MPI_Comm_free(&comm_a);
if (comm_b != MPI_COMM_NULL)
MPI_Comm_free(&comm_b);
MPI_Group_free(&group_a);
MPI_Group_free(&group_b);
MPI_Group_free(&world_group);
```

```
$ mpicc comm_group.c -o comm_group
$ mpirun -np 5 ./comm_group
```

```
comm_a: World Rank 1 received 100 from Rank 0
comm_b: World Rank 1 received 200 from Rank 0
comm_a: World Rank 3 received 100 from Rank 0
Rank 3 tried to send to Rank 4 via MPI_COMM_WORLD
(rc=0)
comm_a: World Rank 2 received 100 from Rank 0
comm_b: World Rank 2 received 200 from Rank 0
comm_b: World Rank 4 received 200 from Rank 0
Rank 4: No message from Rank 3 (as expected)
```