MPI

Spring Training 2025

Yoshihiro Izawa

2025/07/09

Contents

Contents

1. MPI Overview
1.1 What is MPI (Message Passing Interface)?
1.2 Parallel Programming Classification
1.3 MPI Features
1.4 Typical example of Usage
1.5 Comparison between implementations
1.6 Key Communication Primitives
1.7 Internal Mechanisms of MPI Communication
2. Basic Learning of MPI
2.1 MPI Simulation Website
2.2 Tutorial Programs
2.3 Minimum MPI Program
2.4 Important Terms of MPI
2.5 Point-to-Point Communication
3. Collective Communication

Contents

3.1 synchronization	31
3.2 MPI_Barrier	32
3.3 MPI_Bcast	35
3.4 MPI_Scatter	42
3.5 MPI_Gather	
3.6 MPI_Allgather	50
3.7 MPI_Reduce	
3.8 MPI_Allreduce	57
4. Application Example	60
4.1 How to Run Multi-node program on Miyabi	61
4.2 Page Rank Calculation	63
5. References	69
6. Appendix	
6.1 How to use Page Rank when requested a search query	
6.2 Page Rank Full Code	

1. MPI Overview

1.1 What is MPI (Message Passing Interface)?

- A standard API for message passing between distributed memories in parallel computing.
- MPI assumes a distributed-memory computing system
- MPI can run on shared-memory computing system
- MPI programming model (basically) uses **SPMD**(Single Program, Multiple Data).

1.2 Parallel Programming Classification

- Multi-Process: MPI(Message Passing Interface), HPF(High Performance Fortran)
- Multi-Thread: OpenMP, Pthread(POSIX Thread)

Aspect	MPI	HPF
Type	Parallel communication library	Fortran language extension
Language	C / C++ / Fortran	Fortran only
Parallelism Control	Fully manual by programmer	Mostly compiler-driven
Flexibility	Very high	Limited
Maintainability	Hard but highly tunable	Simpler but harder to tune
Learning Curve	High	Low to medium
Current Usage	Mainstream in HPC	Obsolete / deprecated

1.3 MPI Features

- Communication Model:
 - Uses message passing for communication between processes.
- Distributed Memory Support:
 - Each process has its own memory space, no shared memory.
- Multi-node Capacity:
 - Can run across multiple nodes; abstracts network communication.
- Standardized API:
 - Standardized interface in C, C++, and Fortran; highly portable.
- Multiple Implementation:
 - Available implementations include OpenMPI, MPICH, and Intel MPI, etc.
- Difficult to Debug:
 - Debugging is challenging due to concurrency and communication complexity.

1.4 Typical example of Usage

- Simulation on a supercomputer
 - ▶ Physics, Meteorology, Chemistry, etc.
- Data processing in large-scale data analysis
 - e.g., genomics, astronomy
- Machine learning training on large datasets
 - e.g., distributed deep learning

1.5 Comparison between implementations

	OpenMPI	MPICH	Intel MPI
Developer	Universities, Companies	Argonne National Laboratory	Intel Corporation
Distribution	Open source	Open source	Closed source
Optimization Target	General purpose	Lightweight, stable	Optimized for Intel architecture
Performance	Medium to high	Lightweight, stable, scalable	Best performance on Intel CPUs
Main Use	Academic clusters, general HPC	Research, education	Commercial HPC, Intel clusters

- <u>Miyabi uses OpenMPI</u> as the default MPI implementation.(Miyabiis the supercomputer system of the University of Tokyo)
- To be more specific, moior.com/mpicc.on/Miyabi is bound to nvc compiler (NVIDIA HPC SDK C compiler), which means NVIDIA HPC SDK + MPI environment is used.

1.6 Key Communication Primitives

- System function:
 - ► MPI_Init, MPI_Finalize, MPI_Comm_size, MPI_Comm_rank
- Point-to-point communication:
 - ► MPI_Send, MPI_Recv
- Collective communication:
 - ▶ MPI Bcast, MPI Reduce, MPI Alltoall
- Synchronization:
 - ► MPI_Barrier, MPI_Wait, MPI_Test
- Derived data types:
 - ► MPI_Type_create_struct, MPI_Type_vector
- Non-blocking communication:
 - ► MPI_Isend, MPI_Irecv
- Remote memory access:
 - ► MPI Put, MPI Get
- Process management:
 - ► MPI_Comm_spawn, MPI_Comm_free

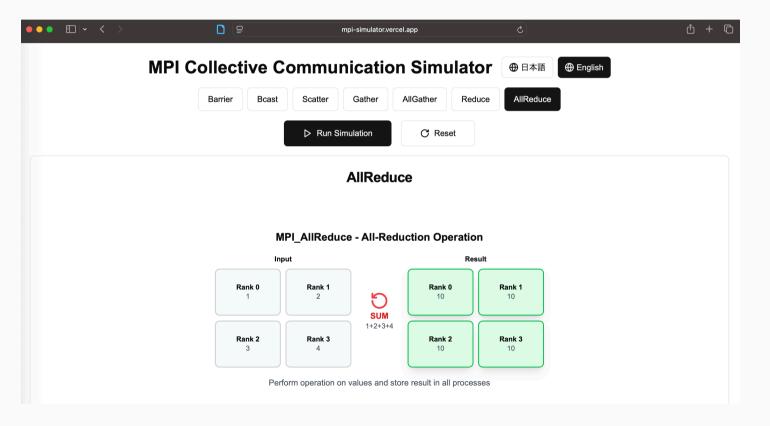
1.7 Internal Mechanisms of MPI Communication

- MPI implementations (e.g., MPICH, OpenMPI) rely on various OS system calls and low-level libraries for communication.
- Intra-node Communication (within the same node):
 - Shared memory: mmap, shm_open, memfd_create, System V shm
 - Event waiting: futex, poll, epoll, select
 - ▶ Pipes & sockets: write, read
 - UNIX domain sockets: sendmsg, recvmsg
 - Synchronization: sem_open, pthread_mutex, spinlock
- Inter-node Communication (across nodes):
 - ► TCP/IP sockets: socket, bind, listen, accept, connect, send, recv, sendto, recvfrom
 - ▶ RDMA & verbs API (e.g., InfiniBand) for zero-copy communication
- High-Performance Cluster Communication:
 - ► Libraries: libibverbs, UCX, OFI (Libfabric), XPMEM, NVLink
 - Kernel bypass with DMA (Direct Memory Access) for low-latency, high-throughput transfers

2. Basic Learning of MPI

2.1 MPI Simulation Website

- I developed a simple MPI simulation website for learning MPI.
- Please Click on: Izawa MPI Simulation Website



2.2 Tutorial Programs

• The following programs are available below.

MPI Tutorial GitHub Repository

2.3.1 MPI Language Differences

	С	C++	Fortran	
MPI Header	<pre>#include <mpi.h></mpi.h></pre>	<pre>#include <mpi.h></mpi.h></pre>	use mpiorinclude 'mpif.h'	
Official MPI support	0	A	0	
Syntax intuitiveness	Explicit C syntax	Almost same as C	call and subroutine based	
Compiler	mpicc	mpicxx or mpic++	mpif90 or mpifort	
Scientific computing	0	A		

- C++
 - ► MPI-3.0 abolished C++ only bindings.
 - Currently, C++ also uses C interface.
- Fortran
 - Considering readability, type safety, and portability, use mpi is recommended.

2.3.2 Hello World (C)

```
#include <mpi.h>
#include <stdio.h>
#include <stdlib.h>
int main(int argc, char *argv[])
    MPI Init(&argc, &argv);
    int num procs;
    int my rank;
    MPI Comm size(MPI COMM WORLD, &num procs);
    MPI Comm rank(MPI COMM WORLD, &my rank);
    printf("Num of Proc : %d\n", num procs);
    printf("My Rank : %d\n", my rank);
   MPI Finalize();
    return EXIT_SUCCESS;
```

```
mpicc mpi_hello.c -o mpi_hello
mpirun -np 4 ./mpi_hello
Num of Proc : 4
My Rank : 3
Num of Proc : 4
My Rank : 2
Num of Proc : 4
My Rank : 0
Num of Proc : 4
My Rank : 1
```

2.3.3 Hello World (C++)

```
#include <mpi.h>
#include <iostream>
#include <cstdlib>
int main(int argc, char *argv[])
{
    MPI Init(&argc, &argv);
    int num procs;
    int my rank;
    MPI Comm size(MPI COMM WORLD, &num procs);
    MPI Comm rank(MPI COMM WORLD, &my rank);
    std::cout << "Num of Proc : " << num procs <<</pre>
std::endl:
    std::cout << "My Rank : " << my rank << std::endl;</pre>
    MPI Finalize();
    return EXIT_SUCCESS;
```

```
mpic++ mpi_hello.cpp -o mpi_hello
mpirun -np 4 ./mpi_hello
Num of Proc : 4
My Rank : 3
Num of Proc : 4
My Rank : 1
Num of Proc : 4
My Rank : 0
Num of Proc : 4
My Rank : 0
Num of Proc : 4
My Rank : 2
```

2.3.4 Hello World (Fortran)

```
program hello mpi
 use mpi
 implicit none
 integer :: ierr, rank, size
 call MPI Init(ierr)
 call MPI Comm rank(MPI COMM WORLD, rank, ierr)
 call MPI Comm size(MPI COMM WORLD, size, ierr)
 print *, "Num of Proc:", size
 print *, "My Rank: ", rank
 call MPI Finalize(ierr)
end program hello mpi
```

```
mpif90 mpi_hello.f90 -o mpi_hello
mpirun -np 4 ./mpi_hello
Num of Proc: 4
My Rank: 2
Num of Proc: 4
My Rank: 0
Num of Proc: 4
My Rank: 3
Num of Proc: 4
My Rank: 1
```

2.4.1 Overview

- Process:
 - computing unit in parallel computing in MPI.
 - process num is determined by mpirun -np
- Group:
 - a set of processes that can communicate with each other.
- Communicator:
 - a group of processes that can communicate with each other.
- Rank:
 - unique identifier for each process in MPI.
 - ► Ranks are assigned from 0 to num_procs 1.

2.4.2 Communicator Image

- Each process belongs to some group.
- A group is associated with a **communicator**.
- Each process in a communicator has a unique rank.

```
- Example:
- Process 0, 1, 2, 3 belong to a group.
- Communicator: MPI_COMM_WORLD
- Group: [P0, P1, P2, P3]
- Rank: 0, 1, 2, 3

MPI_COMM_WORLD (communicator)

├ Group: { P0, P1, P2, P3 }

│ ├ P0 (rank 0)

│ ├ P1 (rank 1)

│ ├ P2 (rank 2)

│ └ P3 (rank 3)
```

2.4.3 Communicator and Group

- **Group** is a set(list) of processes. It is just a list, it does not have any communication capability.
- Communicator containes a group and communication capability. It is a unit of communication.
- A group can create multiple communicators.
- A communicator cannnot communicate with other communicators.

```
Group A: {a, b, c}

↓ Communicator B created from Group A

Communicator B: {a, b, c, d}

↓ Communicator C created from Group A

Communicator C: {a, b, c, e}

- ○ a, b, c ↔ d (via Communicator B)

- ○ a, b, c ↔ e (via Communicator C)

- x d ↔ e (cannot communicate)
```

2.4.4 MPI Functions for Communicator

- MPI_COMM_WORLD:
 - the default communicator that includes all processes.
 - all processes first belong to this communicator.
 - becomes the default communicator for most MPI functions.
- MPI_Comm_rank():
 - retrieves the rank of the calling process in specified communicator.
 - ▶ usually use MPI_COMM_WORLD as the communicator.
- MPI_Comm_size():
 - retrieves the number of processes in the specified communicator.
 - ▶ usually use MPI_COMM_WORLD as the communicator.
- MPI_Comm_split():
 - creates a new communicator by splitting the existing one based on a color and key.
 - in other words, create a new communicator with the same color processes.

MPI_Comm_rank and MPI_Comm_size

```
int rank;
MPI_Comm_rank(MPI_COMM_WORLD, &rank);
printf("I am process %d\n", rank);
```

MPI_Comm_size

```
int size;
MPI_Comm_size(MPI_COMM_WORLD, &size);
printf("There are %d processes\n", size);
```

MPI_Comm_split

```
// ランクの3の剰余を基に color=0, 1, 2 に分ける
int color = rank % 3;
MPI Comm new comm;
MPI Comm split(MPI COMM WORLD, color, rank, &new comm);
int new rank, new size;
MPI Comm rank(new comm, &new rank);
MPI Comm size(new comm, &new size);
printf("World Rank %d => Group %d, New Rank %d of %d\n",
       rank, color, new rank, new size);
MPI_Comm_free(&new_comm); // 新しいコミュニケータの解放
```

```
$ mpicc comm_split.c -o comm_split
$ mpirun -np 8 ./comm_split

World Rank 5 => Group 2, New Rank 1 of 2
World Rank 2 => Group 2, New Rank 0 of 2
World Rank 1 => Group 1, New Rank 0 of 3
World Rank 4 => Group 1, New Rank 1 of 3
World Rank 7 => Group 1, New Rank 2 of 3
World Rank 3 => Group 0, New Rank 1 of 3
World Rank 6 => Group 0, New Rank 2 of 3
World Rank 6 => Group 0, New Rank 2 of 3
World Rank 0 => Group 0, New Rank 0 of 3
```

World Rank	color(= rank % 3)	new group	new rank	new_size
0	0	{0, 3, 6}	0	3
1	1	{1, 4, 7}	0	3
2	2	{2, 5}	0	2
3	0	{0, 3, 6}	1	3
4	1	{1, 4, 7}	1	3
5	2	{2, 5}	1	2
6	0	{0, 3, 6}	2	3
7	1	{1, 4, 7}	2	3

- Many MPI functions have the following signature
- MPI_Send sends data to a specific process.
- MPI_Recv receives data from a specific process.

```
MPI_Recv(
    void* data,
    int count,
    MPI_Datatype datatype,
    int source,
    int tag,
    MPI_Comm communicator,
    MPI_Status* status);
```

MPI Data Type	C Type
MPI_SHORT	short int
MPI_INT	int
MPI_LONG	long int
MPI_LONG_LONG	long long int
MPI_UNSIGNED_CHAR	unsigned char
MPI_UNSIGNED_SHORT	unsigned short int
MPI_UNSIGNED	unsigned int
MPI_UNSIGNED_LONG	unsigned long int
MPI_UNSIGNED_LONG_LONG	unsigned long long int
MPI_FLOAT	float
MPI_DOUBLE	double
MPI_LONG_DOUBLE	long double
MPI_BYTE	char

```
send_data()
```

```
int send_data[10];
for (int i = 0; i < 10; i++)
    send_data[i] = i + 1;
int data_count = 10;

printf("Rank 0: Sending data.\n");
printf("send_data: [");
for (int i = 0; i < 10; i++)
    printf(" %d", send_data[i]);
printf(" ]\n");

MPI_Send((void*)send_data, data_count, MPI_INT,
1, 0, MPI_COMM_WORLD);</pre>
```

```
recv_data()

int data[10];
int data_count = 10;
MPI_Status st;

printf("Rank 1: Receiving data.\n");
MPI_Recv((void*)data, data_count, MPI_INT, 0, 0,
MPI_COMM_WORLD, &st);
printf("recv_data: [");
for (int i = 0; i < 10; i++)
    printf(" %d", data[i]);
printf(" ]\n");</pre>
```

```
$ mpicc send_recv.c -o mpi_send_recv
$ mpirun -np 2 ./mpi_send_recv
Rank 0: Sending data.
send_data: [ 1 2 3 4 5 6 7 8 9 10 ]
Rank 1: Receiving data.
recv_data: [ 1 2 3 4 5 6 7 8 9 10 ]
```

- We can combine send_data() and recv_data() into a single program.
- The program can be run with mpirun -np 2 ./mpi_send_recv2

```
int data count = 10;
int numbers[data count];
if (world rank == 0) {
    for (int i = 0; i < 10; i++) numbers[i] = i + 1;
    printf("Send Data:");
    for (int i = 0; i < 10; i++)
        printf(" %d", numbers[i]);
    printf("\n");
    MPI Send((void *)&numbers, data count, MPI INT, 1, 0, MPI COMM WORLD);
} else if (world rank == 1) {
    MPI Recv(&numbers, data count, MPI INT, 0, 0, MPI COMM WORLD, MPI STATUS IGNORE);
    printf("Received Data:");
    for (int i = 0; i < 10; i++)
        printf(" %d", numbers[i]);
    printf("\n");
}
```

```
$ mpicc send_recv2.c -o mpi_send_recv2
$ mpirun -np 2 ./mpi_send_recv2

Send Data: 1 2 3 4 5 6 7 8 9 10
Received Data: 1 2 3 4 5 6 7 8 9 10
```

3. Collective Communication

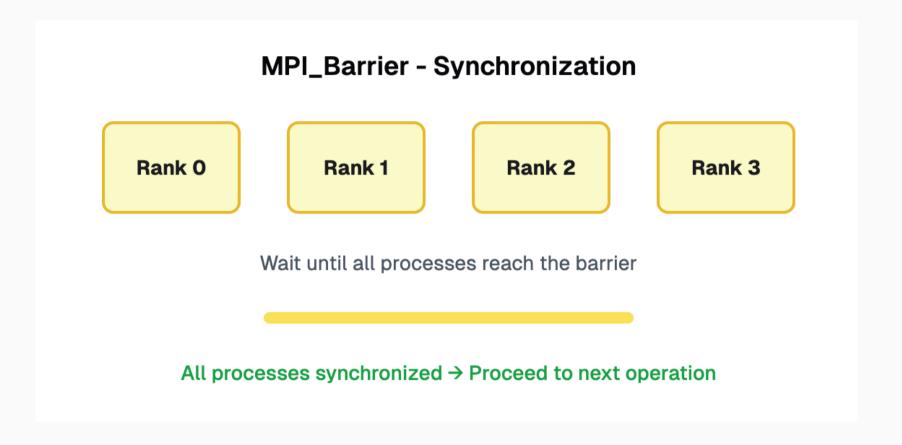
3.1 synchronization

- Collective communication is a communication method that <u>involves all processes in a communicator</u>.
- In collective communication, synchronization among all process is required.
- · All process cannot proceed until all processes reach the same point.
- To achieve this, MPI provides several collective communication functions.

3.2 MPI_Barrier

- MPI_Barrier is a collective communication function that synchronizes all processes in a communicator.
- All processes must call MPI_Barrier to ensure that all processes reach the same point before proceeding.
- It is often used to ensure that all processes have completed their previous tasks before moving on to the next step.
- The most basic usage of MPI_Barrier is for precise time measurement.
- If you do not call MPI_Barrier in all processes, the program will block and cannot proceed.
- MPI_Barrier(MPI_Comm communicator);

3.2 MPI_Barrier



3.2 MPI_Barrier

```
int rank, size;
MPI_Comm_rank(MPI_COMM_WORLD, &rank);
MPI_Comm_size(MPI_COMM_WORLD, &size);

printf("Rank %d: before barrier\n", rank);

sleep(rank);

MPI_Barrier(MPI_COMM_WORLD);
printf("Rank %d: after barrier\n", rank);
```

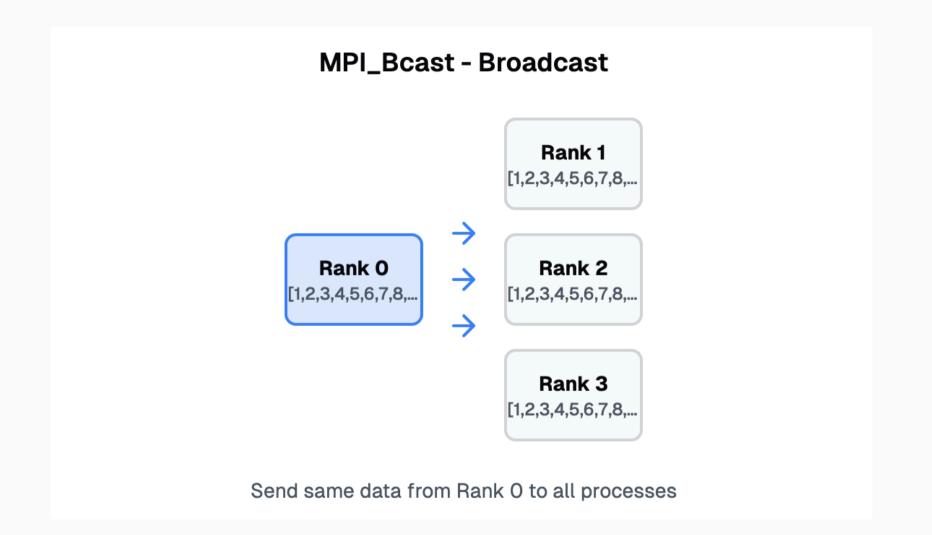
```
$ mpicc barrier.c -o barrier
$ mpirun -np 4 ./barrier

Rank 2: before barrier
Rank 3: before barrier
Rank 0: before barrier
Rank 1: before barrier
Rank 3: after barrier
Rank 0: after barrier
Rank 0: after barrier
Rank 0: after barrier
```

3.3 MPI_Bcast

- MPI_Bcast is a collective communication function that broadcasts data from one process to all other processes in a communicator.
- It is used to distribute data from a root process to all other processes.
- All processes in the communicator must call MPI_Bcast with the same parameters.

- Root Process: Sends the data to all other processes.
- Other Processes: Receive the data from the root process.



```
int rank, size;
MPI Comm rank(MPI COMM WORLD, &rank);
MPI Comm size(MPI COMM WORLD, &size);
int data[10]:
if (rank == 0) {
    for (int i = 0; i < 10; i++) {
        data[i] = i + 1;
    printf("Rank 0: broadcasting data = [");
    for (int i = 0; i < 10; i++)
        printf(" %d", data[i]);
    printf(" ]\n");
}
MPI Bcast(data, 10, MPI INT, 0, MPI COMM WORLD);
printf("Rank %d: received data = [", rank);
for (int i = 0; i < 10; i++)
    printf(" %d", data[i]);
printf(" ]\n");
```

```
$ mpicc bcast.c -o bcast
$ mpirun -np 4 ./bcast

Rank 0: broadcasting data = [ 1 2 3 4 5 6 7 8 9 10 ]
Rank 0: received data = [ 1 2 3 4 5 6 7 8 9 10 ]
Rank 2: received data = [ 1 2 3 4 5 6 7 8 9 10 ]
Rank 3: received data = [ 1 2 3 4 5 6 7 8 9 10 ]
Rank 1: received data = [ 1 2 3 4 5 6 7 8 9 10 ]
```

• We can implement MPI_Bcast wrapper using MPI_Send and MPI_Recv.

```
void my bcast(void* data, int count, MPI Datatype datatype, int root, MPI Comm communicator) {
  int world rank;
 MPI Comm rank(communicator, &world rank);
 int world size;
 MPI Comm size(communicator, &world_size);
  if (world rank == root) {
   for (int i = 0; i < world_size; i++) {</pre>
     if (i != world rank) {
       MPI Send(data, count, datatype, i, 0, communicator);
 } else {
    MPI Recv(data, count, datatype, root, 0, communicator, MPI STATUS IGNORE);
```

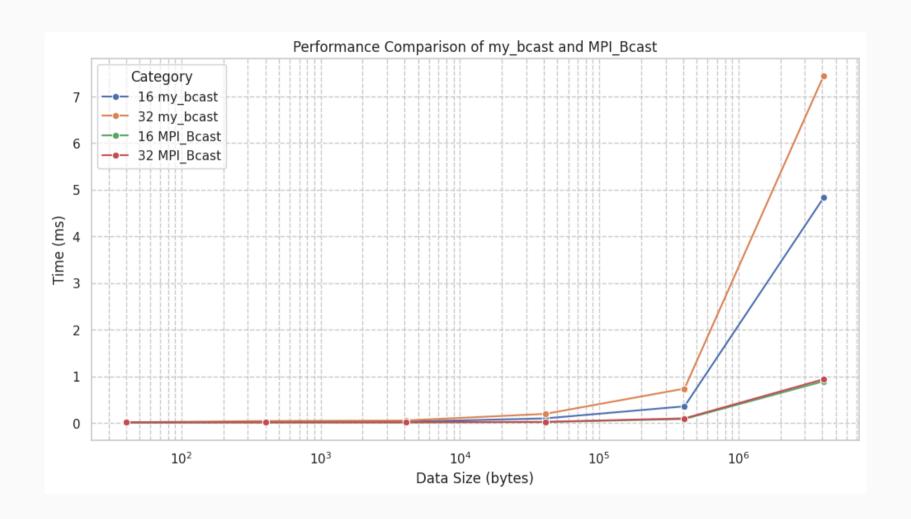
• Q: Is this equivalent to MPI_Bcast?

- A: No, it is less efficient than MPI_Bcast.
- This implementation has only one network communication link. The process with root rank sends data to all other processes one by one.
- MPI_Bcast uses Tree-based broadcast algorithm.
- 1. The root process sends data to process 1.
- 2. The root process sends data to process 2, process 1 sends data to process 3.
- 3. The root process sends data to process 4, process 1 sends data to process 5, process 2 sends data to process 6, process 3 sends data to process 7.

. . .

- Comparison of MPI_Bcast and my_bcast
- Average time of 10 trials

Procs	Data Size	my_bcast (ms)	MPI_Bcast (ms)
16	40	0.008	0.009
16	400	0.022	0.009
16	4k	0.026	0.010
16	40k	0.096	0.018
16	400k	0.355	0.084
16	4000k	4.832	0.893
32	40	0.012	0.012
32	400	0.041	0.011
32	4k	0.052	0.013
32	40k	0.193	0.022
32	400k	0.735	0.097
32	4000k	7.447	0.937



- MPI_Scatter is a collective communication function that distributes data from a root process to all other processes in a communicator.
- MPI_Bcast sends the same data to all processes, while MPI_Scatter sends different chunks of data to each process.

- send_count: the number of elements to send to each process.
- recv_count: the number of elements to receive from each process.



```
#define TOTAL DATA 20
MPI Comm rank(MPI COMM WORLD, &rank);
MPI Comm size(MPI COMM WORLD, &size);
int send data[TOTAL DATA];
int recv count = TOTAL DATA / size;
int recv data[recv count];
if (rank == 0) {
    for (int i = 0; i < TOTAL DATA; i++)
        send data[i] = i;
    printf("Rank 0: Scattering data...\n");
}
MPI Scatter(send data, recv count, MPI INT,
            recv data, recv count, MPI INT,
            0, MPI COMM WORLD);
printf("Rank %d received:", rank);
for (int i = 0; i < recv count; i++)
    printf(" %d", recv data[i]);
printf("\n");
```

```
// send_data == 5

$ mpicc scatter.c -o scatter
$ mpirun -np 4 ./scatter

Rank 0: Scattering data...
Rank 0 received: 0 1 2 3 4
Rank 1 received: 5 6 7 8 9
Rank 2 received: 10 11 12 13 14
Rank 3 received: 14 15 16 17 18
```

if send_data cannot divide by size, the last process will receive the remaining data.

```
// send_data == recv_data == 6

Rank 0 received: 0 1 2 3 4 5
Rank 1 received: 6 7 8 9 10 11
Rank 2 received: 12 13 14 15 16 17
Rank 3 received: 18 19 20 -875497504 65535 20
```

What happens if send_count is not the same as recv_count?

```
int send_data[TOTAL_DATA];
int send count = TOTAL DATA / size;
int recv count = TOTAL DATA / size;
int recv data[recv count];
if (rank == 0) {
    for (int i = 0; i < TOTAL DATA; i++) {
        send data[i] = i;
    printf("Rank 0: Scattering data...\n");
}
if (rank == 3) {
    recv count = 4;
MPI Scatter(send data, send count, MPI INT,
            recv data, recv count, MPI INT,
            0, MPI COMM WORLD);
```

• A: It will cause an error.

```
$ mpicc scatter2.c -o scatter2
$ mpirun -np 4 ./scatter2

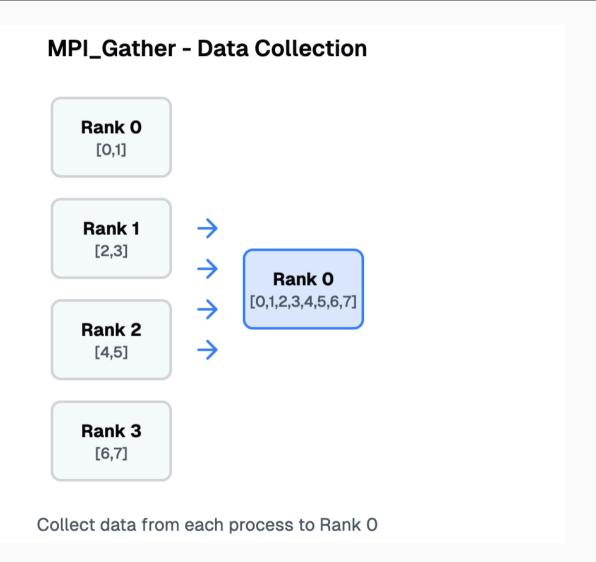
Rank 0: Scattering data...
Rank 0 received: 0 1 2 3 4
Rank 1 received: 5 6 7 8 9
Rank 2 received: 10 11 12 13 14
[miyabi-g3:1978858] *** An error occurred in MPI_Scatter
[miyabi-g3:1978858] *** reported by process [1190789121,3]
[miyabi-g3:1978858] *** on communicator MPI_COMM_WORLD
[miyabi-g3:1978858] *** MPI_ERR_TRUNCATE: message truncated
[miyabi-g3:1978858] *** MPI_ERR_TRUNCATE: message truncated
[miyabi-g3:1978858] *** mPI_ERRORS_ARE_FATAL (processes in this communicator will now abort,
[miyabi-g3:1978858] *** and potentially your MPI job)
```

3.5 MPI_Gather

- MPI_Gather is a collective communication function that collects data from all processes in a communicator and sends it to a root process.
- It is the reverse operation of MPI_Scatter.
- This is used in parallel sorting, parallel searching, and other parallel algorithms.

- Except for the root process, pass NULL for recv_data is allowed.
- recv_count is the number of elements to receive from each process, not the total number of elements.

3.5 MPI_Gather



3.5 MPI_Gather

```
#define ITEMS PER PROC 2
. . .
int send data[ITEMS PER PROC];
send data[0] = rank * 2;
send data[1] = rank * 2 + 1;
int recv data[ITEMS PER PROC * size];
MPI Gather(send data, ITEMS PER PROC, MPI INT,
            recv data, ITEMS PER PROC, MPI INT,
            0, MPI COMM WORLD);
if (rank == 0) {
    printf("Rank 0 gathered data: ");
    for (int i = 0; i < ITEMS PER PROC * size;</pre>
i++)
        printf("%d ", recv data[i]);
    printf("\n");
} else {
    printf("Rank %d sent data: %d %d\n", rank,
send data[0], send data[1]);
```

```
$ mpicc gather.c -o gather
$ mpirun -np 4 ./gather

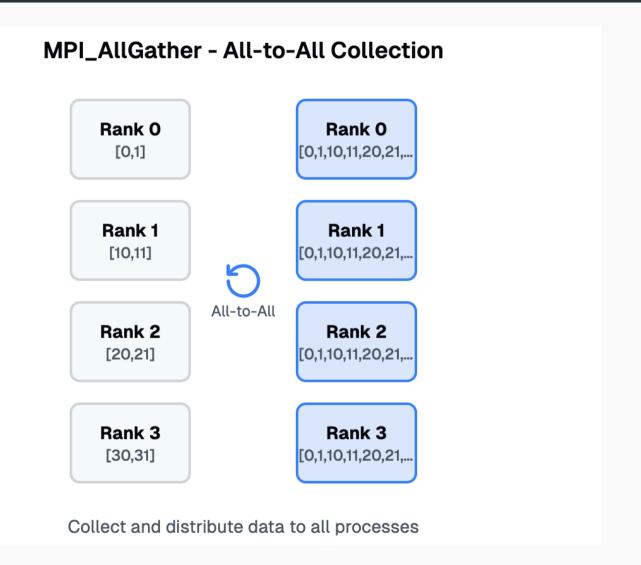
Rank 1 sent data: 2 3
Rank 3 sent data: 6 7
Rank 2 sent data: 4 5
Rank 0 gathered data: 0 1 2 3 4 5 6 7
```

3.6 MPI_Allgather

- MPI_Scatter and MPI_Gather conduct many-to-one or one-to-many communication.
- It is useful if you send data from multiple processes to multiple processes.
- MPI_Allgather is a collective communication function that collects data from all processes in a communicator and sends it to all other processes.
- It is like first MPI_Gather and then MPI_Bcast. Collect data by process rank order.

MPI_Allgather does not have a root process parameter.

3.6 MPI_Allgather



3.6 MPI_Allgather

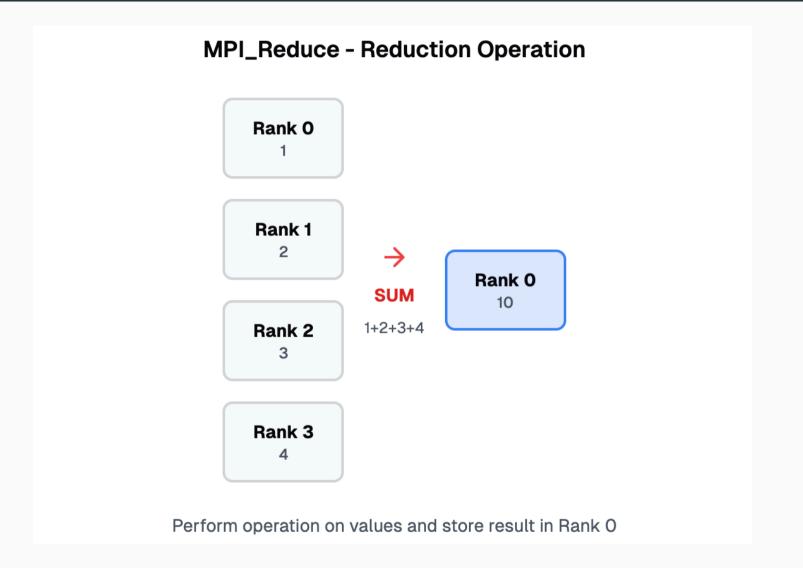
```
#define ITEMS PER PROC 2
. . .
int rank, size;
MPI Comm rank(MPI COMM WORLD, &rank);
MPI Comm size(MPI COMM WORLD, &size);
int send data[ITEMS PER PROC];
send data[0] = rank * 10;
send data[1] = rank * 10 + 1;
int recv data[ITEMS PER PROC * size];
MPI_Allgather(send data, ITEMS PER PROC,
MPI INT,
              recv data, ITEMS PER PROC,
MPI INT,
              MPI COMM WORLD);
printf("Rank %d received:", rank);
for (int i = 0; i < ITEMS PER PROC * size; i++)</pre>
    printf(" %d", recv data[i]);
printf("\n");
```

```
$ mpicc allgather.c -o allgather
$ mpirun -np 4 ./allgather

Rank 2 received: 0 1 10 11 20 21 30 31
Rank 3 received: 0 1 10 11 20 21 30 31
Rank 0 received: 0 1 10 11 20 21 30 31
Rank 1 received: 0 1 10 11 20 21 30 31
```

- reduce is a basic concept in functional programming. It transforms a set of numbers into a smaller set of numbers.
 - \rightarrow reduce([1, 2, 3, 4, 5], sum) = 15
 - ▶ reduce([1, 2, 3, 4, 5], multiply) = 120
- Collect dictributed data and apply a reduction operation is a tough task. However, MPI provides a simple interface to do this.
- MPI_Reduce is a collective communication function that collects data from all processes in a communicator to the root process, and applies a reduction operation to the data.

- MPI Reduction Operations:
 - ► MPI_MAX maximum value
 - ► MPI_MIN minimum value
 - ► MPI_SUM sum of values
 - MPI_PROD product of values
 - MPI_LAND logical AND of values
 - ► MPI_LOR logical OR of values
 - ► MPI_BAND bitwise AND of values
 - MPI_BOR bitwise OR of values
 - ► MPI_MAXLOC maximum value and its rank



```
int rank, size;
MPI Comm rank(MPI COMM WORLD, &rank);
MPI Comm size(MPI COMM WORLD, &size);
int value = rank + 1;
int sum;
MPI Reduce(&value, &sum, 1, MPI INT, MPI SUM, 0,
MPI COMM WORLD);
double average = (double)sum / size;
if (rank == 0) {
    printf("Rank %d: sum = %d, avg = %.2f\n",
rank, sum, average);
} else {
    printf("Rank %d: sum = %d, avg = %.2f\n",
rank, sum, average);
```

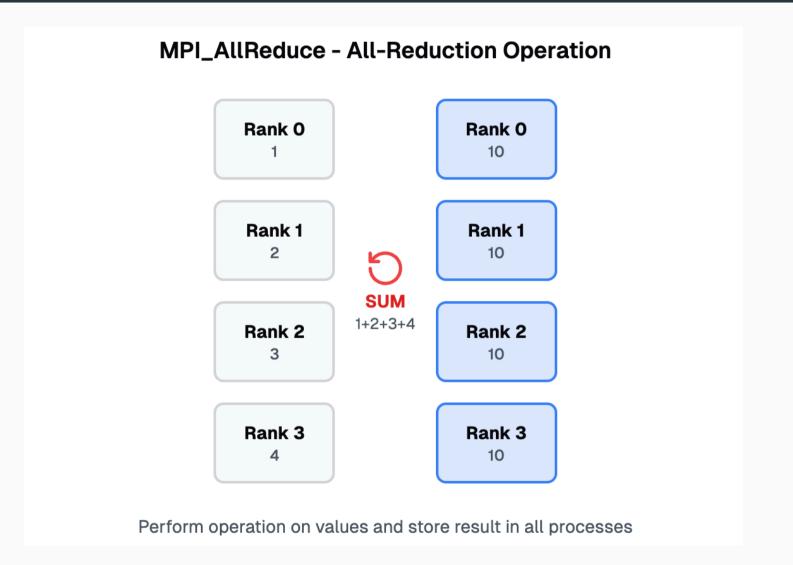
```
$ mpicc reduce.c -o reduce
$ mpirun -np 4 ./reduce

Rank 1: sum = 4197236, avg = 1049309.00
Rank 0: sum = 10, avg = 2.50
Rank 2: sum = 4197236, avg = 1049309.00
Rank 3: sum = 4197236, avg = 1049309.00
```

3.8 MPI_Allreduce

- MPI_Allreduce is a collective communication function that collects data from all processes in a communicator, applies a reduction operation to the data, and distributes the result to all processes.
- It is similar to MPI_Reduce, but the result is available to all processes, not just the root process.
- This is useful when all processes need to know the result of the reduction operation.

3.8 MPI_Allreduce



3.8 MPI_Allreduce

```
int rank, size;
MPI_Comm_rank(MPI_COMM_WORLD, &rank);
MPI_Comm_size(MPI_COMM_WORLD, &size);

int value = rank + 1;
int total_sum = 0;

MPI_Allreduce(&value, &total_sum, 1, MPI_INT,
MPI_SUM, MPI_COMM_WORLD);

double average = (double)total_sum / size;
printf("Rank %d: total sum = %d, average =
%.2f\n", rank, total_sum, average);
```

```
$ mpicc allreduce.c -o allreduce
$ mpirun -np 4 ./allreduce

Rank 2: total sum = 10, average = 2.50
Rank 1: total sum = 10, average = 2.50
Rank 0: total sum = 10, average = 2.50
Rank 3: total sum = 10, average = 2.50
```

4. Application Example

4.1 How to Run Multi-node program on Miyabi

- The following is an example of how to run the Page Rank calculation on a multi-node cluster.
- When you submit a job to Miyabi-G cluster, you can specify the number of node and the number of process per node using -l select={num_nodes}:mpiprocs={num_procs_per_node} option.
- When you run a MPI program on a multi-node cluster, you need to specify the number of processes using mpirun -np {num_procs} option.

Example run.sh script:

```
#!/bin/bash
#PBS -q debug-g
#PBS -l select=16:mpiprocs=4
#PBS -W group_list=gc64
#PBS -o latest_result.txt
#PBS -j oe
...
mpirun -np 64 ./naive/pagerank_naive
```

4.1 How to Run Multi-node program on Miyabi

- If you want to run a MPI program on another process-per-node than mpiprocs, you can use --host option or --hostfile option.
- --host option allows you to specify the hostnames and the number of processes per host.
- --hostfile option allows you to specify a file that contains the hostnames and the number of processes per host.

Example -- host option:

```
$ mpirun -np 4 --host node1:2,node2:2 ./your_mpi_program
→node1 and node2 each run 2 processes.
```

Example -- hostfile option: hostfile.txt:

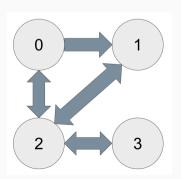
```
node1 slots=2
node2 slots=2/

$ mpirun -np 4 --hostfile hostfile.txt ./your_mpi_program
```

- For a simple application example of MPI, I will introduce the Page Rank calculation.
- Page Rank is an algorithm <u>used by Google to rank web pages</u> in search results.
- The concept of Page Rank is that "The page is important if it is linked from many other important pages."

$$x^{\{(k+1)\}} = dMx^{\{(k)\}} + (1-d)v$$

- M is the link matrix, each element represents the link between pages,
- x is the Page Rank vector, each element represents the Page Rank of a page,
- d is the damping factor, usually set to 0.85,
- v is the uniform distribution vector (initial vector).



$$d = 0.85, M = \begin{pmatrix} 0 & 0 & \frac{1}{3} & 0 \\ \frac{1}{2} & 0 & \frac{1}{3} & 0 \\ \frac{1}{2} & 1 & \frac{1}{3} & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix}, v = \begin{pmatrix} \frac{1}{4} \\ \frac{1}{4} \\ \frac{1}{4} \\ \frac{1}{4} \\ \frac{1}{4} \end{pmatrix}$$

$$M = \begin{pmatrix} 0 & 0 & \frac{1}{3} & 0 \\ \frac{1}{2} & 0 & \frac{1}{3} & 0 \\ \frac{1}{2} & 1 & \frac{1}{3} & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix}$$

- The sum of each column in matrix M is 1 because M represents a transition probability matrix, where each column corresponds to a node and the entries indicate the probability of moving from that node to others
- The Page Rank vector is updated iteratively until convergence.

$$\begin{split} x^{\{(1)\}} &= dM x^{\{(0)\}} + (1-d)v \\ x^{\{(2)\}} &= dM x^{\{(1)\}} + (1-d)v \\ x^{\{(3)\}} &= dM x^{\{(2)\}} + (1-d)v \end{split}$$

• • •

How to cunduct Page Rank calculation in parallel?

- Split the link matrix M into submatrices, each assigned to a process.
- Each process calculates the Page Rank for its submatrix.
- Use MPI_Allreduce to combine the results from all processes.
- Repeat until convergence.
- The following code is a simple example of Page Rank calculation using MPI.

Run Step:

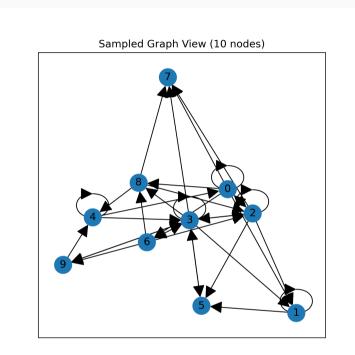
- 1. Generate a graph data file.
- 2. (option) Visualize the graph data.
- 3. Run the Page Rank calculation.

```
$ cd tutorial/pagerank/

// graph data generation
$ ./create_venv.sh
$ python3 preprocess/generate_graph.py
$ ls data/

// graph data visualization
$ python3 preprocess/visualize.py

// Page Rank calculation
$ cd tutorial/pagerank/naive
$ ./run.sh
```



```
$ ./run.sh

Final PageRank:
     0    1    2    3    4    5    6    7    8    9
PR: 0.0490 0.0817 0.0678 0.1453 0.0616 0.0748 0.0374 0.0598 0.0453 0.0384
```

- We divide the M matrix into submatrices, each assigned to a process.
- Each process only have a submatrix of M and a subvector of $x_k(*x)$, but the whole vector of $x_{k+1}(*new_x_full)$.

```
typedef struct {
    int n nodes;
    int n edges;
    int *out degree;
    double **M;
} Graph;
Graph g;
int n = g.n nodes;
// Data partitioning (row-wise)
int rows per proc = n / size;
int remainder = n % size;
int my rows = rows per proc + (rank < remainder ?</pre>
1:0):
int my start = rank * rows per proc + (rank <</pre>
remainder ? rank : remainder):
```

```
double *x = malloc(n * sizeof(double));
double *new x local = malloc(my rows *
sizeof(double));
for (int i = 0; i < n; i++) x[i] = 1.0 / n;
for (int i = 0; i < size; i++) {
    recvcounts[i] = rows per proc + (i <</pre>
remainder ? 1 : 0);
    displs[i] = i * rows per proc + (i <</pre>
remainder ? i : remainder);
double *new x full = malloc(n * sizeof(double));
```

• MPI_Allgatherv is different from MPI_Allgather in that it allows each process to send a different number of elements.

```
for (int iter = 0; iter < MAX ITER; iter++) {</pre>
    for (int i = 0; i < my rows; i++) {
        int global i = my start + i;
        new x local[i] = 0.0;
        for (int j = 0; j < n; j++) {
            new x local[i] += q.M[qlobal i][i] *
x[j];
        new x local[i] = DAMPING * new x local[i]
+ (1.0 - DAMPING) / n;
    MPI_Allgatherv(new_x_local, my rows,
MPI DOUBLE,
                    new x full, recvcounts,
displs, MPI DOUBLE, MPI COMM WORLD);
```

```
// チェック用: rank 0 が diff を計算
   double diff = 0.0:
   if (rank == 0) {
       for (int i = 0; i < n; i++) {
           diff += fabs(new x full[i] - x[i]);
   // 全rankにdiffをブロードキャスト(終了判定共有)
   MPI Bcast(&diff, 1, MPI DOUBLE, 0,
MPI COMM WORLD);
   if (diff < TOL) break;</pre>
   // x を更新
   for (int i = 0; i < n; i++)
       x[i] = new x full[i];
```

5. References

5. References

- Referenced in this tutorial:
 - ▶ MPI「超」入門(C言語編)- 東京大学情報基盤センター
 - ▶ 並列プログラミング入門
 - MPI Tutorial
- For Advanced Users:
- ・ 行列-行列積(2) 非同期通信 東京大学情報基盤センター
- Open MPI GitHub Repository

6. Appendix

6.1 How to use Page Rank when requested a search query

- Page Rank is just a algorithm to rank web pages.
- The requested search query is just a string, so we cannot use Page Rank directly.
- One approach is:
 - 1. we get related web pages according to DF-IDF or BM25 algorithm
 - 2. then we apply Page Rank to rerank the related web pages.

```
graph_loader.h
```

```
#ifndef GRAPH LOADER H
#define GRAPH LOADER H
typedef struct {
    int n nodes;
    int n edges;
    int *out degree;
    double **M;
} Graph;
int read graph(const char *filename, Graph *g);
void free graph(Graph *q);
#endif
```

graph_loader.c

```
#include "graph loader.h"
#include <stdio.h>
#include <stdlib.h>
int read graph(const char *filename, Graph *g) {
    FILE *fp = fopen(filename, "r");
    if (!fp) {
        perror("fopen");
        return -1;
    if (fscanf(fp, "%d", &g->n nodes) != 1) {
        fprintf(stderr, "Failed to read node
count\n");
        return -1;
```

```
g->out degree = calloc(g->n nodes,
sizeof(int));
    g->M = malloc(g->n nodes * sizeof(double
*));
    for (int i = 0; i < q -> n \text{ nodes}; i++) {
        g->M[i] = calloc(g->n nodes,
sizeof(double));
    int from, to;
    q - n edges = 0;
    while (fscanf(fp, "%d %d", &from, &to) == 2)
{
        g \rightarrow M[to][from] += 1.0;
        g->out degree[from]++;
        g->n edges++;
    fclose(fp);
```

```
// 正規化:M[to][from] /= out degree[from]
    for (int i = 0; i < g > n \text{ nodes}; i++) {
        for (int j = 0; j < q > n \text{ nodes}; j++) {
             if (g->out degree[j] > 0) {
                 g->M[i][j] /= g->out degree[j];
    return 0:
void free graph(Graph *g) {
    for (int i = 0; i < q -> n \text{ nodes}; i++) {
        free(g->M[i]);
    free(q->M);
    free(g->out degree);
```

main.c

```
#include <stdio.h>
#include <stdlib.h>
#include <math.h>
#include <mpi.h>
#include "graph loader.h"
#define DAMPING 0.85
#define TOL 1e-6
#define MAX ITER 100
int main(int argc, char **argv) {
    MPI Init(&argc, &argv);
    int rank, size;
    MPI Comm rank(MPI COMM WORLD, &rank);
    MPI Comm size(MPI COMM WORLD, &size);
```

```
if (argc < 2) {
    if (rank == 0) {
        fprintf(stderr, "Usage: %s <n node>\n",
argv[0]);
    MPI Abort(MPI COMM WORLD, 1);
int n node = atoi(argv[1]);
char graph path[256], pagerank path[256];
snprintf(graph path, sizeof(graph path), "data/
%d/graph.txt", n node);
snprintf(pagerank path, sizeof(pagerank path),
"data/%d/pagerank.txt", n node);
Graph g;
if (read graph(graph path, &g) != 0) {
    MPI Abort(MPI COMM WORLD, 1);
```

```
int n = g.n nodes;
// データ分割(行分割)
int rows per proc = n / size;
int remainder = n % size;
int my rows = rows per proc + (rank <</pre>
remainder ? 1 : 0);
int my start = rank * rows per proc + (rank <</pre>
remainder ? rank : remainder):
// PageRankベクトルxと各プロセスが担当する新しいベクトル
new x local(len(my rows)次元)を確保
double *x = malloc(n * sizeof(double));
double *new x local = malloc(my rows *
sizeof(double));
// 全プロセスで保持するの新しいPageRankベクトル
new x full(n次元)を確保
double *new x full = malloc(n * sizeof(double));
int *recvcounts = malloc(size * sizeof(int));
int *displs = malloc(size * sizeof(int));
for (int i = 0; i < n; i++) x[i] = 1.0 / n;
```

```
for (int i = 0; i < size; i++) {
    recvcounts[i] = rows per proc + (i <</pre>
remainder ? 1 : 0);
    displs[i] = i * rows per proc + (i <</pre>
remainder ? i : remainder);
if (rank == 0) {
    printf("Starting PageRank computation with
%d processes...\n", size);
printf("[Rank %d] Local rows: %d, Start index:
%d\n", rank, my rows, my start);
double t start = MPI Wtime();
for (int iter = 0; iter < MAX ITER; iter++) {</pre>
    for (int i = 0; i < my rows; i++) {
        int global i = my start + i;
        new x local[i] = 0.0;
        for (int j = 0; j < n; j++) {
            new x local[i] += g.M[global i][j] *
x[j];
        new x local[i] = DAMPING *
new x local[i] + (1.0 - DAMPING) / n;
```

```
MPI Allgatherv(new x local, my rows,
MPI DOUBLE,
                   new x full, recvcounts,
displs, MPI DOUBLE, MPI COMM WORLD);
    // チェック用: rank 0 が diff を計算
    double diff = 0.0:
    if (rank == 0) {
       for (int i = 0; i < n; i++) {
            diff += fabs(new x full[i] - x[i]);
       printf("[Iter %d] diff = %f\n", iter +
1, diff);
    // 全rankにdiffをブロードキャスト(終了判定共有)
    MPI Bcast(&diff, 1, MPI DOUBLE, 0,
MPI COMM WORLD);
    if (diff < TOL) break;</pre>
   // x を更新
    for (int i = 0; i < n; i++) x[i] =
new x full[i];
```

```
double t end = MPI Wtime();
if (rank == 0) {
    printf("Main Loop Time: %.6f seconds\n",
t end - t start);
free(new x full);
free(recvcounts);
free(displs);
// 結果表示
if (rank == 0) {
    printf("\n\nFinal PageRank:\n");
    for (int i = 0; i < n; i++) {
        printf("Node %6d: %f\n", i, x[i]);
    printf("\n");
```

```
// 結果をファイルに書き出す
if (rank == 0) {
    FILE *fp = fopen(pagerank path, "w");
    if (fp == NULL) {
       perror("Failed to open result file");
       MPI Abort(MPI COMM WORLD, 1);
    fprintf(fp, "%d\n", n);
   for (int i = 0; i < n; i++) {
       fprintf(fp, "%d %f\n", i, x[i]);
    fclose(fp);
    printf("PageRank results written to %s\n", pagerank_path);
free(x);
free(new x local);
free graph(&g);
MPI Finalize();
return 0;
```