# Functional programming

## Introduction

- High level of abstraction
- Based on lambda calculus
- Language of choice: Haskell
- Functional programming -> expressions > statements.

#### Example:

```
--Add the first ten numbers together sum[1..10]
```

• Install Hugs, Haskell interpreter (ghci is another interpreter)

## First steps

• Start hugs in terminal

#### Examples:

```
> head[1,2,3,4] --take first element
1
> tail[1,2,3,4] --remove first element
[2,3,4]
> [1,2,3,4] !! 2 --element #2
3
> take 3 [1,2,3,4] -- generalization of head
[1,2,3]
> drop 3 [1,2,3,4,5] --generalization of tail
[4,5]
> [1,2,3]++[4,5] --append
[1,2,3,4,5]
```

- First element in list has index 0
- List different to array -> indexing bad idea, not in constant time but in linear

**Function application**: function application is denoted by space. Higher priority

```
f a b + c*d --f(a,b) + c d from math
f a + b --f(a) + b from math
```

- Haskell file (script) -> .hs
- Define function in script, then open Hugs with script as argument so that functions are available. If script is changed use :reload. Also possible to load using :load script
- Infix operator: xfy --> f x y
- Naming:
  - function and parameter name must begin with lowercase
  - can use quotes (prime)
  - type has to start with uppercase
  - convention -> s at the end means list, ss list of lists
- Indentation like Python, implicit grouping
- Useful commands -> :load script, :reload, :edit script, :type expression.:?
- Comments: one line --comment, nested:

```
{-
very long
comment goes
here
-}
```

## Types and classes

- Type: name for a collection of related values. Example Bool
- Applying a function to a wrong type makes a type error
- e :: t -> e has type t
- Type inference: compiler calculates type of expression prior to execution. Haskell programs are type safe, type error never happens in run time
- :type <exp> to calculate type of expression

Table 1: Different types in Haskell

Type	Explanation
Bool	Logical value: True or False
Char	Single character, enclosed in sigle quotes: 'a'
String	String of characters, double quotes: "abc"
Int	Fixed precision integer
Integer	Arbitrary precision integer, doesn't overflow
Float	Single precision floating point number

• List: sequence of values with same type. Can be infinite. Examples:

```
[False, True, False]::[Bool] -- list of elements type Bool ['a', 'b']:: [Char] -- list of elements type char
```

• **Tuple**: sequence of values of different type. Number of elements is called *arity*. Finite number of elements because type of all have to be calculated. Example:

```
(False, 'a') :: (Bool, Char) -- length appears in type
```

• **Function**: mapping from values of a type to values of another type. Examples:

```
not :: Bool -> Bool
isDigit :: Char -> Bool
function :: t1 -> t2 -- from domain to range in general
```

- Curried function: functions that return arguments one at a time (functions can return functions): a -> (a -> a) equivalent to a -> a -> a, arrow associates to the right. Any function that returns more than one values can be curried. Useful for partially applying functions. Most functions applied in curried form, if tuples are not explicitly declared.
- Polymorphic function: functions not defined for a particular type. Example:

```
length :: [a] -> Int
```

- Price for polymorphism: type variables start with lowercase and types with uppercase
- Overloaded function: functions with same name but different types. In Haskell overloading means that there is a restriction in the type class. Example:

```
sum :: Num [a] => [a] -> Int -- only numeric values allowed
```

## **Defining functions**

#### Conditional expressions

Example:

```
abs :: Int \rightarrow Int
abs n = if n >=0 then n else -n
```

- Can be nested
- Conditional expressions must have an else branch

#### Guarded equation:

- Sequence of logical expressions
- Alternative to conditional (Haskell people prefer this)

```
abs n | n >= 0 = n -- /= such that
| otherwise = -n
```

• Can be used to make definitions involving multiple conditions

## Pattern matching

More efficient way using wildcard + lazy evaluation:

```
(&&) :: Bool -> Bool -> Bool
True && b = b -- True && something --> something
False && _ = False -- always False
```

- Order is important
- Patterns may not repeat variables: all the variables inside the pattern have to be different
- Lists in pattern matching: use cons definition  $(:)^1$ . Only matches not empty list. These pattern must be parethesized because function application has higher priority.

```
head :: [a] \rightarrow a
head (x : ) = x
```

<sup>&</sup>lt;sup>1</sup>Lists are constructed one element at a time from the empty list using cons operator [1,2,3] = 1:(2:(3:[]))= 1:2:3:[]

### Lambda expressions

Functions can be constructed without naming using lambda expressions:

```
\x \rightarrow x + x -- \x = \ lambda x (from lambda calculus)
```

• Useful for currying:

```
add x y = x + y
add = \x -> (\y -> x+y)
```

- For returning functions as results
- For avoiding naming functions only used once:

```
odds n = map f [0..n-1]
where
f x = x*2 +1
odds n = map (\x -> x*2 + 1)[0..n-1] --pass lambda as parameter to map
```

#### Sections

Operator written between two arguments can be used in curried way using parenthesis:

```
--Examples:
(1+) --sucessor
(1/) --reciprocate
(*2) --double
(/2) --half
```

For avoiding naming.

## List comprehensions

- Code that manipulates collections
- Favorite collection for mathematicians: sets. Problems with sets:
  - No duplication
  - Deal with equality

- Haskell has tricks to deal with sets as lists
- Set comprehensions in Math:  $\{x^2|x\in\{1,...,5\}\}$
- List comprehension in Haskell:  $[x^2 \mid x \leftarrow [1..5]]$
- $x \leftarrow [1..5]$  is called generator
- Comprehensions can have multiple generators (similar to nested loop, x is the outer loop and y is the inner):

```
> [(x,y) | x <- [1,2,3], y <- [4,5]]
> [(1,4),(1,5),(2,4),(2,5),(3,4),(3,5)]
> [(x,y) | y <- [4,5], x <- [1,2,3]]
> [(1,4),(2,4),(3,4),(1,5),(2,5),(3,5)]
```

• Generators can depend on each other, as in loops:

```
> [(x,y) | x <- [1,2,3], y <- [x..3]]
> [(1,1),(1,2),(1,3),(2,2),(2,3),(3,3)]
```

- Very concise code
- Filters (guards): [x | x <- [1..10], even x]

### The zip function

- Combines two list to a list of pairs
- Useful when programming with list comprehensions.

```
zip :: [a] -> [b]-> [(a,b)]
```

### String comprehension

- Strings = character list [Char]
- Everything that can be done in lists will work with strings

#### Recursive functions

• Tail call elimination

```
factorial 1 = 0
factorial n = n * factorial (n-1)
```

- Some functions are simpler to define using recursions
- Induction can be used to prove porperties of recursive functions
- Recursion can be used also in lists:

```
product :: [Int] -> Int
product [] = 1
product (n:ns) = n * product ns
```

Note: : appends element to list, ++ concatenates lists

- Quicksort: algorithm for sorting integers. Two rules:
  - The empty list is already sorted
  - Bolzano in the rest

```
qsort :: [Int] -> Int
qsort [] = []
qsort (x:xs)=
  qsort smaller ++ [x] ++ qsort larger
  where
    smaller = [a | a <- xs, a <= x]
    larger = [b | b <- xs, b > x]
```

## **Higher-order functions**

Higher-order functions are functions that take functions as arguments or return functions as results.

Useful for:

- Programming idioms, avoid repetition
- Domain specific languages
- Algebraic properties to reason about programs

#### Map

Applies function to every element in list

```
map :: (a -> b) -> [a] -> [b]
```

Map can be defined using list comprehension:

```
map f xs = [f x | x < - xs]
```

Or recursively (for abstraction):

```
map f [] = []
map f (x:xs) = f x : map f xs
```

### Filter

Removes elements that don't satisfy a predicate

```
filter :: (a -> Bool) -> [a] -> [a]
```

Definition using list comprehension:

```
filter p xs = [x \mid x \leftarrow xs, p x]
```

Or recursively:

#### Foldr & foldl

Homomorfism over list -> generalization of sum, product...

- r: from the right
- l: from the left

```
f [] = v

f (x:xs) = x (+operator) f xs
```

Examples:

```
sum = foldr (+) 0
product = foldr (*) 1
and = folder (&&) True
```

Replaces the empty list by v and cons by f

Useful for:

- For defining recursive functions
- Properties of functions defined usig foldr can be proved using algebracic properties
- Program optimization

### Other library functions

Composition: combines two functions into one

```
(.) :: (b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow (a \rightarrow c)

f . g = \x \rightarrow f(g \x) --first apply g, apply f to result

Example:

odd :: Int -> Bool

odd = not . even
```

Use it sparingly because it's difficult to read

all: decides if every element in a list satisfies condition any: decides if every any in a list satisfies condition takeWhile: takes elements while condition is true dropWhile: drop elements while condition is true

## Functional parsers and monads

Parser: program hat analyses piece of test and determines its syntantic structure

```
type Parser = String -> Tree

--If there is unused output
type Parser = String -> (Tree, String)

--If there are more than one option of parsing
type Parser = String -> [(Tree, String)]

--Parsers can produce any type
type Parser = String -> [(a, String)]
```

### Basic parsers

Examples

```
--Parsers single character

item :: Parser Char
item = \inp -> case inp of

[] -> []

(x:xs) -> [(x,xs)]
```

```
--case for pattern matching in the body of definition
--Always succeeds
return :: a -> Parser a
return v = \langle inp -\rangle [(v, inp)]
--Always fails
failure :: Parser a
failure = \inp -> []
The function parse applies parser to input:
```

```
parse :: Parser a -> String -> [(a,String)]
parse p inp = p inp
```

## Sequencing

A sequence of parsers can be combined using do:

```
p :: Parser (Char, Char)
p = do x \leftarrow item
        item
        y \leftarrow item
        return (x,y)
```

- Parser can be combined using +++(else), if the first one fails, apply the second and so on
- Layout rule!
- If one parser fails, all fail

## Derived primitives

Using the three basic parsers we can define parser that returns single characters that return characters that satisfy a given predicate:

```
sat :: (Char -> Bool) -> Parser Char
sat p = do x < - item
           if p x then return x else failure
```

Using sat and different predicates we can define parsers for digits, lower-case letters...

## Interactive programs

Interaction with keyboard and screen

**Problem**: Haskell programs have no side effects. Same arguments give same results. Readline, for example, does not give the same result all the time!

**Solution**: new type IO a, actions that have side effects. Function that return void: IO () (empty tuple)

#### Basic actions

Actions in the standard library. Examples:

```
getChar :: IO Char --reads character from standard input
putChar :: IO Char --writes character to standard output
```

### Sequencing

Works like parsing

### Derived primitives

We can read a string from standard input basing on getChar, for instance.

## Declaring types and classes

- We can define new types in Haskell
- Things in common with objects in Object Oriented Programming

#### Type declarations

Using type synonym between types can be defined. Example:

```
type String = [Char]
```

- Type definition can be used for better readability
- Type declarations can be nested, but they can't be recursive<sup>2</sup>.

<sup>&</sup>lt;sup>2</sup>Recursive declaration is only valid in nominal types (see Data declaration)

#### Data declarations

Using data completely new types can be defined giving their values. Example:

```
data Bool = False | True --Boolean types can be True or False
```

- These can be declared recursively, because are data and types at the same time.
- The new values of the type are called **constructors** in this case.
- Types and constructors names must begin with capital letter.
- The same constructor name cannot be used in more than one type.
- Values of new types can be used as built-in types
- The construtors can also have arguments
- Data declarations can be parameterised

#### Recursive types

New types can be declared in terms of themselves, if declared using data:

```
data Nat = Zero | Succ Nat --Creates an infinite sequence of values
```

#### Class and instance declaration

Classes are declared using class. Example:

For a type **a** to be an instance of a given class it has to comply – with the conditions in the class definition. For instancing:

• Only types declared using data can become classes.

- Classes can also be extended to make new classes (inheritance) using =>
- For making a type instance of a built-in type use deriving. Example: