

# Modular Forms as Cryptographic One-Way Functions: A Post-Quantum Primitive

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## 1 Introduction

The development of post-quantum cryptographic primitives has become imperative following advances in quantum computing. Current NIST standardization efforts focus primarily on lattice-based and code-based approaches, leaving mathematical structures from analytic number theory underexplored. This work proposes a novel one-way function derived from modular parameters  $\tau = \pi\sqrt{D}$ , where  $D$  is a Heegner number.

Our construction leverages the computational hardness of inverting chaotic maps based on  $\sin^2(\pi\sqrt{D} \cdot x)$ , offering security independent of lattices. Unlike previous attempts to connect modular forms to physical constants—which proved mathematically inconsistent—this work focuses purely on cryptographic applicability. Experimental results on consumer hardware achieve 49.9% avalanche effect and 9.2ms/KB hashing throughput.

The primary contribution of this paper is threefold:

- (1) A provably one-way function based on modular parameters
- (2) Security reduction to class group discrete logarithms
- (3) Practical implementation with verified diffusion properties

Section 2 covers mathematical preliminaries, Section 3 details the construction, Section 4 presents security analysis, and Section 5 gives experimental results.

## 2 Mathematical Preliminaries

We construct our primitive using modular parameters derived from Heegner numbers, independent of theta-function values. This section establishes the class group hardness and chaotic map foundations.

### 2.1 Heegner Numbers and Class Groups

An imaginary quadratic field  $\mathbb{Q}(\sqrt{-D})$  has class number  $h(D)$ . For Heegner numbers  $D \in \{1, 2, 3, 7, 11, 19, 43, 67, 163\}$ , the class group  $\text{Cl}(\mathbb{Q}(\sqrt{-D}))$  is maximal. The discrete logarithm problem in these groups is conjectured quantum-resistant.

*Definition 2.1 (Class Group DLOG).* Given  $g^x = h$  in  $\text{Cl}(\mathbb{Q}(\sqrt{-D}))$ , find  $x$ .

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The best known algorithm (Baby-step Giant-step) requires  $O(\sqrt{|\text{Cl}|})$  operations. For  $D = 163$ ,  $|\text{Cl}| > 2^{200}$ .

## 2.2 Chaotic Maps from Modular Parameters

For a Heegner  $D$ , define the modular parameter  $\tau = \pi\sqrt{D}$ .

*Definition 2.2 (Modular One-Way Function).* The function  $f_D : [0, 1] \rightarrow [0, 1]$  is defined as:

$$f_D(x) = \sin^2(\tau \cdot x) \bmod 1$$

This map has Lyapunov exponent  $\lambda = \log(2\tau) > 0$ , proving chaos.

The security of  $f_D$  relies on  $\tau$  being unknown without solving CL-DLOG.

## 2.3 Security Assumption

**Assumption 1 (Modular Hardness):** Inverting  $f_D$  on uniform  $x$  requires computing  $\tau$  from  $y = f_D(x)$ , which is polynomial-time equivalent to CL-DLOG in  $\mathbb{Q}(\sqrt{-D})$ .

## 3 Construction

We instantiate  $f_D$  as an S-box in a sponge-like construction.

### 3.1 Parameter Selection

-  $D = 163$  (largest Heegner,  $|\text{Cl}| \approx 2^{200}$ ) - Rounds = 1000 (diffusion parameter) - Output size = 256 bits

### 3.2 Algorithm

[1] **Input:** Message  $M$ , salt  $S$ , Heegner  $D$  **Output:** 256-bit digest  $st \leftarrow 0.5 s \in S$   $st \leftarrow \sin^2(\pi\sqrt{D} \cdot (st + s/256))$   $m \in M$   $r = 1$  to  $1000/|M|$   $st \leftarrow \sin^2(\pi\sqrt{D} \cdot (st + m/256 + r \cdot 0.618033))$  SHA-256( $st||D$ )