

# Grimoire's Standard Code Library<sup>\*</sup>

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<sup>\*</sup> <https://github.com/kzoacn/Grimoire>



# 目录

<b>1</b>	<b>代数</b>	<b>5</b>
1.1	$O(n^2 \log n)$ 求线性递推数列第 $n$ 项 . . . . .	5
1.2	闪电数论变换与魔力 CRT . . . . .	6
1.3	多项式求逆 . . . . .	7
1.4	多项式除法 . . . . .	8
1.5	多项式取指数取对数 . . . . .	9
<b>2</b>	<b>数论</b>	<b>11</b>
2.1	大整数相乘取模 . . . . .	11
2.2	线段下整点 . . . . .	11
2.3	中国剩余定理 . . . . .	11
<b>3</b>	<b>图论</b>	<b>13</b>
3.1	一般图匹配 . . . . .	13
3.2	无向图最小割 . . . . .	15
<b>4</b>	<b>技巧</b>	<b>17</b>
4.1	无敌的读入优化 . . . . .	17
4.2	真正释放 STL 内存 . . . . .	18



# Chapter 1

## 代数

### $O(n^2 \log n)$ 求线性递推数列第 $n$ 项

Given  $a_0, a_1, \dots, a_{m-1}$   
 $a_n = c_0 * a_{n-m} + \dots + c_{m-1} * a_0$   
 $a_0$  is the  $n$ th element,  $\dots$ ,  $a_{m-1}$  is the  $n + m - 1$ th element

```
1 void linear_recurrence(long long n, int m, int a[], int c[], int p) {
2     long long v[M] = {1 % p}, u[M << 1], msk = !!n;
3     for(long long i(n); i > 1; i >= 1) {
4         msk <= 1;
5     }
6     for(long long x(0); msk; msk >= 1, x <= 1) {
7         fill_n(u, m < 1, 0);
8         int b(!!(n & msk));
9         x |= b;
10        if(x < m) {
11            u[x] = 1 % p;
12        }else {
13            for(int i(0); i < m; i++) {
14                for(int j(0), t(i + b); j < m; j++, t++) {
15                    u[t] = (u[t] + v[i] * v[j]) % p;
16                }
17            }
18            for(int i((m < 1) - 1); i >= m; i--) {
19                for(int j(0), t(i - m); j < m; j++, t++) {
20                    u[t] = (u[t] + c[j] * u[i]) % p;
21                }
22            }
23        }
24        copy(u, u + m, v);
25    }
26    //a[n] = v[0] * a[0] + v[1] * a[1] + ... + v[m - 1] * a[m - 1].
27    for(int i(m); i < 2 * m; i++) {
28        a[i] = 0;
```

```

29     for(int j(0); j < m; j++) {
30         a[i] = (a[i] + (long long)c[j] * a[i + j - m]) % p;
31     }
32 }
33 for(int j(0); j < m; j++) {
34     b[j] = 0;
35     for(int i(0); i < m; i++) {
36         b[j] = (b[j] + v[i] * a[i + j]) % p;
37     }
38 }
39 for(int j(0); j < m; j++) {
40     a[j] = b[j];
41 }
42 }

```

## 闪电数论变换与魔力 CRT

```

1 #define meminit(A, l, r) memset(A + (l), 0, sizeof(*A) * ((r) - (l)))
2 #define memcpy(B, A, l, r) memcpy(B, A + (l), sizeof(*A) * ((r) - (l)))
3 void DFT(int *a, int n, int f) { //f=1 逆 DFT
4     for (register int i = 0, j = 0; i < n; i++) {
5         if (i > j) std::swap(a[i], a[j]);
6         for (register int t = n >> 1; (j ^= t) < t; t >>= 1);
7     }
8     for (register int i = 2; i <= n; i <=> 1) {
9         static int exp[MAXN];
10        exp[0] = 1; exp[1] = fpm(PRT, (MOD - 1) / i, MOD);
11        if (f == 1) exp[1] = fpm(exp[1], MOD - 2, MOD);
12        for (register int k = 2; k < (i >> 1); k++) {
13            exp[k] = 1ll * exp[k - 1] * exp[1] % MOD;
14        }
15        for (register int j = 0; j < n; j += i) {
16            for (register int k = 0; k < (i >> 1); k++) {
17                register int &pA = a[j + k], &pB = a[j + k + (i >> 1)];
18                register long long B = 1ll * pB * exp[k];
19                pB = (pA - B) % MOD;
20                pA = (pA + B) % MOD;
21            }
22        }
23    }
24    if (f == 1) {
25        register int rev = fpm(n, MOD - 2, MOD);
26        for (register int i = 0; i < n; i++) {
27            a[i] = 1ll * a[i] * rev % MOD;
28            if (a[i] < 0) { a[i] += MOD; }

```

```

29     }
30 }
31 }
32 // 在不写高精度的情况下合并 FFT 所得结果对 MOD 取模后的答案
33 // 值得注意的是，这个东西不能最后再合并，而是应该每做一次多项式乘法就 CRT 一次
34 int CRT(int *a) {
35     static int x[3];
36     for (int i = 0; i < 3; i++) {
37         x[i] = a[i];
38         for (int j = 0; j < i; j++) {
39             int t = (x[i] - x[j] + FFT[i] -> MOD) % FFT[i] -> MOD;
40             if (t < 0) t += FFT[i] -> MOD;
41             x[i] = 1LL * t * inv[j][i] % FFT[i] -> MOD;
42         }
43     }
44     int sum = 1, ret = x[0] % MOD;
45     for (int i = 1; i < 3; i++) {
46         sum = 1LL * sum * FFT[i] -> MOD % MOD;
47         ret += 1LL * x[i] * sum % MOD;
48         if (ret >= MOD) ret -= MOD;
49     }
50     return ret;
51 }
52 for (int i = 0; i < 3; i++) // inv 数组的预处理过程，inverse(x, p) 表示求 x 在 p 下逆元
53     for (int j = 0; j < 3; j++)
54         inv[i][j] = inverse(FFT[i] -> MOD, FFT[j] -> MOD);

```

## 多项式求逆

Given polynomial  $a$  and  $n$ ,  $b$  is the polynomial such that  $a * b \equiv 1 \pmod{x^n}$

```

1 void getInv(int *a, int *b, int n) {
2     static int tmp[MAXN];
3     b[0] = fpm(a[0], MOD - 2, MOD);
4     for (int c = 2, M = 1; c < (n << 1); c <= 1) {
5         for (; M <= 3 * (c - 1); M <= 1);
6         meminit(b, c, M);
7         meminit(tmp, c, M);
8         memcpy(tmp, a, 0, c);
9         DFT(tmp, M, 0);
10        DFT(b, M, 0);
11        for (int i = 0; i < M; i++) {
12            b[i] = 1LL * b[i] * (2LL - 1LL * tmp[i] * b[i] % MOD + MOD) % MOD;
13        }
14        DFT(b, M, 1);
15        meminit(b, c, M);

```

```

16     }
17 }

```

## 多项式除法

d is quotient and r is remainder

```

1 void divide(int n, int m, int *a, int *b, int *d, int *r) { // n、m 分别为多项式 A (被除数)
   ↪ 和 B (除数) 的指数 + 1
2     static int M, tA[MAXN], tB[MAXN], inv[MAXN], tD[MAXN];
3     for (; n > 0 && a[n - 1] == 0; n--);
4     for (; m > 0 && b[m - 1] == 0; m--);
5     for (int i = 0; i < n; i++) tA[i] = a[n - i - 1];
6     for (int i = 0; i < m; i++) tB[i] = b[m - i - 1];
7     for (M = 1; M <= n - m + 1; M <= 1);
8     if (m < M) meminit(tB, m, M);
9     getInv(tB, inv, M);
10    for (M = 1; M <= 2 * (n - m + 1); M <= 1);
11    meminit(inv, n - m + 1, M);
12    meminit(tA, n - m + 1, M);
13    DFT(inv, M, 0);
14    DFT(tA, M, 0);
15    for (int i = 0; i < M; i++) {
16        d[i] = 1ll * inv[i] * tA[i] % MOD;
17    }
18    DFT(d, M, 1);
19    std::reverse(d, d + n - m + 1);
20    for (M = 1; M <= n; M <= 1);
21    memcpy(tB, b, 0, m);
22    if (m < M) meminit(tB, m, M);
23    memcpy(tD, d, 0, n - m + 1);
24    meminit(tD, n - m + 1, M);
25    DFT(tD, M, 0);
26    DFT(tB, M, 0);
27    for (int i = 0; i < M; i++) {
28        r[i] = 1ll * tD[i] * tB[i] % MOD;
29    }
30    DFT(r, M, 1);
31    meminit(r, n, M);
32    for (int i = 0; i < n; i++) {
33        r[i] = (a[i] - r[i] + MOD) % MOD;
34    }
35 }

```



## 多项式取指数取对数

Given polynomial  $a$  and  $n$ ,  $b$  is the polynomial such that  $b \equiv e^a \pmod{x^n}$  or  $b \equiv \ln a \pmod{x^n}$

```

1 void getDiff(int *a, int *b, int n) { // 多项式取微分
2     for (int i = 0; i + 1 < n; i++) {
3         b[i] = 1ll * (i + 1) * a[i + 1] % MOD;
4     }
5     b[n - 1] = 0;
6 }
7 void getInt(int *a, int *b, int n) { // 多项式取积分, 积分常数为 0
8     static int inv[MAXN];
9     inv[1] = 1;
10    for (int i = 2; i < n; i++) {
11        inv[i] = 1ll * (MOD - MOD / i) * inv[MOD % i] % MOD;
12    }
13    b[0] = 0;
14    for (int i = 1; i < n; i++) {
15        b[i] = 1ll * a[i - 1] * inv[i] % MOD;
16    }
17 }
18 void getLn(int *a, int *b, int n) {
19     static int inv[MAXN], d[MAXN];
20     int M = 1;
21     for (; M <= 2 * (n - 1); M <= 1);
22     getInv(a, inv, n);
23     getDiff(a, d, n);
24     meminit(d, n, M);
25     meminit(inv, n, M);
26     DFT(d, M, 0); DFT(inv, M, 0);
27     for (int i = 0; i < M; i++) {
28         d[i] = 1ll * d[i] * inv[i] % MOD;
29     }
30     DFT(d, M, 1);
31     getInt(d, b, n);
32 }
33 void getExp(int *a, int *b, int n) {
34     static int ln[MAXN], tmp[MAXN];
35     b[0] = 1;
36     for (int c = 2, M = 1; c < (n <= 1); c <= 1) {
37         for (; M <= 2 * (c - 1); M <= 1);
38         int bound = std::min(c, n);
39         memcpy(tmp, a, 0, bound);
40         meminit(tmp, bound, M);
41         meminit(b, c, M);
42         getLn(b, ln, c);
43         meminit(ln, c, M);

```

```
44     DFT(b, M, 0);
45     DFT(tmp, M, 0);
46     DFT(ln, M, 0);
47     for (int i = 0; i < M; i++) {
48         b[i] = 111 * b[i] * (111 - ln[i] + tmp[i] + MOD) % MOD;
49     }
50     DFT(b, M, 1);
51     meminit(b, c, M);
52 }
53 }
```

## Chapter 2

# 数论

### 大整数相乘取模

```
1 // x 与 y 须非负
2 long long mult(long long x, long long y, long long MODN) {
3     long long t = (x * y - (long long)((long double)x / MODN * y + 1e-3) * MODN) % MODN;
4     return t < 0 ? t + MODN : t;
5 }
```

### 线段下整点

solve for  $\sum_{i=0}^{n-1} \lfloor \frac{a+bi}{m} \rfloor$ ,  $n, m, a, b > 0$

```
1 LL solve(LL n, LL a, LL b, LL m){
2     if(b==0) return n*(a/m);
3     if(a>=m) return n*(a/m)+solve(n, a%m, b, m);
4     if(b>=m) return (n-1)*n/2*(b/m)+solve(n, a, b%m, m);
5     return solve((a+b*n)/m, (a+b*n)%m, m, b);
6 }
```

### 中国剩余定理

first is remainder, second is module

```
1 inline void fix(LL &x, LL y) {
2     x = (x % y + y) % y;
3 }
4 bool solve(int n, std::pair<LL, LL> a[],
5             std::pair<LL, LL> &ans) {
6     ans = std::make_pair(1, 1);
7     for (int i = 0; i < n; ++i) {
8         LL num, y;
9         euclid(ans.second, a[i].second, num, y);
```

```
10     LL divisor = std::__gcd(ans.second, a[i].second);
11     if ((a[i].first - ans.first) % divisor) {
12         return false;
13     }
14     num *= (a[i].first - ans.first) / divisor;
15     fix(num, a[i].second);
16     ans.first += ans.second * num;
17     ans.second *= a[i].second / divisor;
18     fix(ans.first, ans.second);
19 }
20 return true;
21 }
```

## Chapter 3

# 图论

### 一般图匹配

```
1 // 0-base, match[u] is linked to u
2 vector<int> lnk[MAXN];
3 int match[MAXN], Queue[MAXN], pred[MAXN], base[MAXN], head, tail, sta, fin, nbase;
4 bool inQ[MAXN], inB[MAXN];
5 inline void push(int u) {
6     Queue[tail++] = u; inQ[u] = 1;
7 }
8 inline int pop() {
9     return Queue[head++];
10 }
11 inline int FindCA(int u, int v) {
12     static bool inP[MAXN];
13     fill(inP, inP + n, false);
14     while (1) {
15         u = base[u]; inP[u] = 1;
16         if(u == sta) break;
17         u = pred[match[u]];
18     }
19     while (1) {
20         v = base[v];
21         if (inP[v]) break;
22         v = pred[match[v]];
23     }
24     return v;
25 }
26 inline void RT(int u) {
27     int v;
28     while (base[u] != nbase) {
29         v = match[u];
30         inB[base[u]] = inB[base[v]] = 1;
31         u = pred[v];
```

```

32     if (base[u] != nbase) pred[u] = v;
33 }
34 }
35 inline void BC(int u, int v) {
36     nbase = FindCA(u, v);
37     fill(inB, inB + n, 0);
38     RT(u); RT(v);
39     if (base[u] != nbase) pred[u] = v;
40     if (base[v] != nbase) pred[v] = u;
41     for (int i = 0; i < n; ++i)
42         if (inB[base[i]]) {
43             base[i] = nbase;
44             if (!inQ[i]) push(i);
45         }
46 }
47 bool FindAP(int u) {
48     bool found = false;
49     for (int i = 0; i < n; ++i) {
50         pred[i] = -1; base[i] = i; inQ[i] = 0;
51     }
52     sta = u; fin = -1; head = tail = 0; push(sta);
53     while (head < tail) {
54         int u = pop();
55         for (int i = (int)lnk[u].size() - 1; i >= 0; --i) {
56             int v = lnk[u][i];
57             if (base[u] != base[v] && match[u] != v) {
58                 if (v == sta || match[v] >= 0 && pred[match[v]] >= 0) BC(u, v);
59                 else if (pred[v] == -1) {
60                     pred[v] = u;
61                     if (match[v] >= 0) push(match[v]);
62                     else {
63                         fin = v;
64                         return true;
65                     }
66                 }
67             }
68         }
69     }
70     return found;
71 }
72 inline void AP() {
73     int u = fin, v, w;
74     while (u >= 0) {
75         v = pred[u]; w = match[v];
76         match[v] = u; match[u] = v;
77         u = w;

```

```

78     }
79 }
80 inline int FindMax() {
81     for (int i = 0; i < n; ++i) match[i] = -1;
82     for (int i = 0; i < n; ++i)
83         if (match[i] == -1 && FindAP(i)) AP();
84     int ans = 0;
85     for (int i = 0; i < n; ++i) {
86         ans += (match[i] != -1);
87     }
88     return ans;
89 }

```

## 无向图最小割

```

1  /*
2   * Stoer Wagner 全局最小割  $O(V^3)$ 
3   * lbase, 点数 n, 邻接矩阵 edge[MAXN][MAXN]
4   * 返回值为全局最小割
5   */
6
7  int StoerWagner() {
8      static int v[MAXN], wage[MAXN];
9      static bool vis[MAXN];
10
11     for (int i = 1; i <= n; ++i) v[i] = i;
12
13     int res = INF;
14
15     for (int nn = n; nn > 1; --nn) {
16         memset(vis, 0, sizeof(bool) * (nn + 1));
17         memset(wage, 0, sizeof(int) * (nn + 1));
18
19         int pre, last = 1; // vis[1] = 1;
20
21         for (int i = 1; i < nn; ++i) {
22             pre = last; last = 0;
23             for (int j = 2; j <= nn; ++j) if (!vis[j]) {
24                 wage[j] += edge[v[pre]][v[j]];
25                 if (!last || wage[j] > wage[last]) last = j;
26             }
27             vis[last] = 1;
28         }
29
30         res = std::min(res, wage[last]);

```

```
31
32     for (int i = 1; i <= nn; ++i) {
33         edge[v[i]][v[pre]] += edge[v[last]][v[i]];
34         edge[v[pre]][v[i]] += edge[v[last]][v[i]];
35     }
36     v[last] = v[nn];
37 }
38 return res;
39 }
```



## Chapter 4

# 技巧

### 无敌的读入优化

```
1 // getchar() 读入优化 << 关同步 cin << 此优化
2 // 用 isdigit() 会小幅变慢
3 // 返回 false 表示读到文件尾
4 namespace Reader {
5     const int L = (1 << 15) + 5;
6     char buffer[L], *S, *T;
7     __inline bool getchar(char &ch) {
8         if (S == T) {
9             T = (S = buffer) + fread(buffer, 1, L, stdin);
10            if (S == T) {
11                ch = EOF;
12                return false;
13            }
14        }
15        ch = *S++;
16        return true;
17    }
18    __inline bool getint(int &x) {
19        char ch; bool neg = 0;
20        for (; getchar(ch) && (ch < '0' || ch > '9'); ) neg ^= ch == '-';
21        if (ch == EOF) return false;
22        x = ch - '0';
23        for (; getchar(ch), ch >= '0' && ch <= '9'; )
24            x = x * 10 + ch - '0';
25        if (neg) x = -x;
26        return true;
27    }
28 }
```

## 真正释放 STL 内存

```
1 template <typename T>
2 __inline void clear(T& container) {
3     container.clear(); // 或者删除了一堆元素
4     T(container).swap(container);
5 }
```