

Grimoire's Standard Code Library^{*}

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^{*} <https://github.com/kzoacn/Grimoire>

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Chapter 1

代数

1.1 $O(n^2 \log n)$ 求线性递推数列第 n 项

Given a_0, a_1, \dots, a_{m-1}
 $a_n = c_0 * a_{n-m} + \dots + c_{m-1} * a_0$
 a_0 is the n th element, \dots , a_{m-1} is the $n + m - 1$ th element

```
1 void linear_recurrence(long long n, int m, int a[], int c[], int p) {
2     long long v[M] = {1 % p}, u[M << 1], msk = !!n;
3     for(long long i(n); i > 1; i >= 1) {
4         msk <= 1;
5     }
6     for(long long x(0); msk; msk >= 1, x <= 1) {
7         fill_n(u, m << 1, 0);
8         int b(!!(n & msk));
9         x |= b;
10        if(x < m) {
11            u[x] = 1 % p;
12        }else {
13            for(int i(0); i < m; i++) {
14                for(int j(0), t(i + b); j < m; j++, t++) {
15                    u[t] = (u[t] + v[i] * v[j]) % p;
16                }
17            }
18            for(int i((m << 1) - 1); i >= m; i--) {
19                for(int j(0), t(i - m); j < m; j++, t++) {
20                    u[t] = (u[t] + c[j] * u[i]) % p;
21                }
22            }
23        }
24        copy(u, u + m, v);
25    }
26    //a[n] = v[0] * a[0] + v[1] * a[1] + ... + v[m - 1] * a[m - 1].
27    for(int i(m); i < 2 * m; i++) {
28        a[i] = 0;
```

```

29     for(int j(0); j < m; j++) {
30         a[i] = (a[i] + (long long)c[j] * a[i + j - m]) % p;
31     }
32 }
33 for(int j(0); j < m; j++) {
34     b[j] = 0;
35     for(int i(0); i < m; i++) {
36         b[j] = (b[j] + v[i] * a[i + j]) % p;
37     }
38 }
39 for(int j(0); j < m; j++) {
40     a[j] = b[j];
41 }
42 }

```

1.2 闪电数论变换与魔力 CRT

```

1 #define meminit(A, l, r) memset(A + (l), 0, sizeof(*A) * ((r) - (l)))
2 #define memcpy(B, A, l, r) memcpy(B, A + (l), sizeof(*A) * ((r) - (l)))
3 void DFT(int *a, int n, int f) { //f=1 逆 DFT
4     for (register int i = 0, j = 0; i < n; i++) {
5         if (i > j) std::swap(a[i], a[j]);
6         for (register int t = n >> 1; (j ^= t) < t; t >>= 1);
7     }
8     for (register int i = 2; i <= n; i <=> 1) {
9         static int exp[MAXN];
10        exp[0] = 1; exp[1] = fpm(PRT, (MOD - 1) / i, MOD);
11        if (f == 1) exp[1] = fpm(exp[1], MOD - 2, MOD);
12        for (register int k = 2; k < (i >> 1); k++) {
13            exp[k] = 1ll * exp[k - 1] * exp[1] % MOD;
14        }
15        for (register int j = 0; j < n; j += i) {
16            for (register int k = 0; k < (i >> 1); k++) {
17                register int &pA = a[j + k], &pB = a[j + k + (i >> 1)];
18                register long long B = 1ll * pB * exp[k];
19                pB = (pA - B) % MOD;
20                pA = (pA + B) % MOD;
21            }
22        }
23    }
24    if (f == 1) {
25        register int rev = fpm(n, MOD - 2, MOD);
26        for (register int i = 0; i < n; i++) {
27            a[i] = 1ll * a[i] * rev % MOD;
28            if (a[i] < 0) { a[i] += MOD; }

```

```

29     }
30 }
31 }
32 // 在不写高精度的情况下合并 FFT 所得结果对 MOD 取模后的答案
33 // 值得注意的是，这个东西不能最后再合并，而是应该每做一次多项式乘法就 CRT 一次
34 int CRT(int *a) {
35     static int x[3];
36     for (int i = 0; i < 3; i++) {
37         x[i] = a[i];
38         for (int j = 0; j < i; j++) {
39             int t = (x[i] - x[j] + FFT[i] -> MOD) % FFT[i] -> MOD;
40             if (t < 0) t += FFT[i] -> MOD;
41             x[i] = 1LL * t * inv[j][i] % FFT[i] -> MOD;
42         }
43     }
44     int sum = 1, ret = x[0] % MOD;
45     for (int i = 1; i < 3; i++) {
46         sum = 1LL * sum * FFT[i] -> MOD % MOD;
47         ret += 1LL * x[i] * sum % MOD;
48         if (ret >= MOD) ret -= MOD;
49     }
50     return ret;
51 }
52 for (int i = 0; i < 3; i++) // inv 数组的预处理过程，inverse(x, p) 表示求 x 在 p 下逆元
53     for (int j = 0; j < 3; j++)
54         inv[i][j] = inverse(FFT[i] -> MOD, FFT[j] -> MOD);

```

1.3 多项式求逆

Given polynomial a and n , b is the polynomial such that $a * b \equiv 1 \pmod{x^n}$

```

1 void getInv(int *a, int *b, int n) {
2     static int tmp[MAXN];
3     b[0] = fpm(a[0], MOD - 2, MOD);
4     for (int c = 2, M = 1; c < (n <= 1); c <= 1) {
5         for (; M <= 3 * (c - 1); M <= 1);
6         meminit(b, c, M);
7         meminit(tmp, c, M);
8         memcpy(tmp, a, 0, c);
9         DFT(tmp, M, 0);
10        DFT(b, M, 0);
11        for (int i = 0; i < M; i++) {
12            b[i] = 1LL * b[i] * (2LL - 1LL * tmp[i] * b[i] % MOD + MOD) % MOD;
13        }
14        DFT(b, M, 1);
15        meminit(b, c, M);

```

```

16     }
17 }

```

1.4 多项式除法

d is quotient and r is remainder

```

1 void divide(int n, int m, int *a, int *b, int *d, int *r) { // n、m 分别为多项式 A (被除数)
    ↪ 和 B (除数) 的指数 + 1
2     static int M, tA[MAXN], tB[MAXN], inv[MAXN], tD[MAXN];
3     for (; n > 0 && a[n - 1] == 0; n--);
4     for (; m > 0 && b[m - 1] == 0; m--);
5     for (int i = 0; i < n; i++) tA[i] = a[n - i - 1];
6     for (int i = 0; i < m; i++) tB[i] = b[m - i - 1];
7     for (M = 1; M <= n - m + 1; M <= 1);
8     if (m < M) meminit(tB, m, M);
9     getInv(tB, inv, M);
10    for (M = 1; M <= 2 * (n - m + 1); M <= 1);
11    meminit(inv, n - m + 1, M);
12    meminit(tA, n - m + 1, M);
13    DFT(inv, M, 0);
14    DFT(tA, M, 0);
15    for (int i = 0; i < M; i++) {
16        d[i] = 1ll * inv[i] * tA[i] % MOD;
17    }
18    DFT(d, M, 1);
19    std::reverse(d, d + n - m + 1);
20    for (M = 1; M <= n; M <= 1);
21    memcpy(tB, b, 0, m);
22    if (m < M) meminit(tB, m, M);
23    memcpy(tD, d, 0, n - m + 1);
24    meminit(tD, n - m + 1, M);
25    DFT(tD, M, 0);
26    DFT(tB, M, 0);
27    for (int i = 0; i < M; i++) {
28        r[i] = 1ll * tD[i] * tB[i] % MOD;
29    }
30    DFT(r, M, 1);
31    meminit(r, n, M);
32    for (int i = 0; i < n; i++) {
33        r[i] = (a[i] - r[i] + MOD) % MOD;
34    }
35 }

```


1.5 多项式取指数取对数

Given polynomial a and n , b is the polynomial such that $b \equiv e^a \pmod{x^n}$ or $b \equiv \ln a \pmod{x^n}$

```

1 void getDiff(int *a, int *b, int n) { // 多项式取微分
2     for (int i = 0; i + 1 < n; i++) {
3         b[i] = 1ll * (i + 1) * a[i + 1] % MOD;
4     }
5     b[n - 1] = 0;
6 }
7 void getInt(int *a, int *b, int n) { // 多项式取积分, 积分常数为 0
8     static int inv[MAXN];
9     inv[1] = 1;
10    for (int i = 2; i < n; i++) {
11        inv[i] = 1ll * (MOD - MOD / i) * inv[MOD % i] % MOD;
12    }
13    b[0] = 0;
14    for (int i = 1; i < n; i++) {
15        b[i] = 1ll * a[i - 1] * inv[i] % MOD;
16    }
17 }
18 void getLn(int *a, int *b, int n) {
19     static int inv[MAXN], d[MAXN];
20     int M = 1;
21     for (; M <= 2 * (n - 1); M <= 1);
22     getInv(a, inv, n);
23     getDiff(a, d, n);
24     meminit(d, n, M);
25     meminit(inv, n, M);
26     DFT(d, M, 0); DFT(inv, M, 0);
27     for (int i = 0; i < M; i++) {
28         d[i] = 1ll * d[i] * inv[i] % MOD;
29     }
30     DFT(d, M, 1);
31     getInt(d, b, n);
32 }
33 void getExp(int *a, int *b, int n) {
34     static int ln[MAXN], tmp[MAXN];
35     b[0] = 1;
36     for (int c = 2, M = 1; c < (n < 1); c <= 1) {
37         for (; M <= 2 * (c - 1); M <= 1);
38         int bound = std::min(c, n);
39         memcpy(tmp, a, 0, bound);
40         meminit(tmp, bound, M);
41         meminit(b, c, M);
42         getLn(b, ln, c);
43         meminit(ln, c, M);

```

```
44     DFT(b, M, 0);
45     DFT(tmp, M, 0);
46     DFT(ln, M, 0);
47     for (int i = 0; i < M; i++) {
48         b[i] = 111 * b[i] * (111 - ln[i] + tmp[i] + MOD) % MOD;
49     }
50     DFT(b, M, 1);
51     meminit(b, c, M);
52 }
53 }
```

Chapter 2

数论

2.1 大整数相乘取模

```
1 // x 与 y 须非负
2 long long mult(long long x, long long y, long long MODN) {
3     long long t = (x * y - (long long)((long double)x / MODN * y + 1e-3) * MODN) % MODN;
4     return t < 0 ? t + MODN : t;
5 }
```

2.2 线段下整点

solve for $\sum_{i=0}^{n-1} \lfloor \frac{a+bi}{m} \rfloor$, $n, m, a, b > 0$

```
1 LL solve(LL n, LL a, LL b, LL m){
2     if(b==0) return n*(a/m);
3     if(a>=m) return n*(a/m)+solve(n, a%m, b, m);
4     if(b>=m) return (n-1)*n/2*(b/m)+solve(n, a, b%m, m);
5     return solve((a+b*n)/m, (a+b*n)%m, m, b);
6 }
```

2.3 中国剩余定理

first is remainder, second is module

```
1 inline void fix(LL &x, LL y) {
2     x = (x % y + y) % y;
3 }
4 bool solve(int n, std::pair<LL, LL> a[],
5             std::pair<LL, LL> &ans) {
6     ans = std::make_pair(1, 1);
7     for (int i = 0; i < n; ++i) {
8         LL num, y;
9         euclid(ans.second, a[i].second, num, y);
```

```
10     LL divisor = std::__gcd(ans.second, a[i].second);
11     if ((a[i].first - ans.first) % divisor) {
12         return false;
13     }
14     num *= (a[i].first - ans.first) / divisor;
15     fix(num, a[i].second);
16     ans.first += ans.second * num;
17     ans.second *= a[i].second / divisor;
18     fix(ans.first, ans.second);
19 }
20 return true;
21 }
```

Chapter 3

图论

3.1 图论基础

```
1 struct Graph { // Remember to call .init()!
2     int e, nxt[M], v[M], adj[N], n;
3     bool base;
4     __inline void init(bool _base, int _n = 0) {
5         assert(n < N);
6         n = _n; base = _base;
7         e = 0; memset(adj + base, -1, sizeof(*adj) * n);
8     }
9     __inline int new_node() {
10         adj[n + base] = -1;
11         assert(n + base + 1 < N);
12         return n++ + base;
13     }
14     __inline void ins(int u0, int v0) { // directional
15         assert(u0 < n + base && v0 < n + base);
16         v[e] = v0; nxt[e] = adj[u0]; adj[u0] = e++;
17         assert(e < M);
18     }
19     __inline void bi_ins(int u0, int v0) { // bi-directional
20         ins(u0, v0); ins(v0, u0);
21     }
22 };
```

3.2 一般图匹配

```
1 // 0-base, match[u] is linked to u
2 vector<int> lnk[MAXN];
3 int match[MAXN], Queue[MAXN], pred[MAXN], base[MAXN], head, tail, sta, fin, nbase;
4 bool inQ[MAXN], inB[MAXN];
5 inline void push(int u) {
```

```

6   Queue[tail++] = u; inQ[u] = 1;
7   }
8   inline int pop() {
9       return Queue[head++];
10  }
11  inline int FindCA(int u, int v) {
12      static bool inP[MAXN];
13      fill(inP, inP + n, false);
14      while (1) {
15          u = base[u]; inP[u] = 1;
16          if(u == sta) break;
17          u = pred[match[u]];
18      }
19      while (1) {
20          v = base[v];
21          if (inP[v]) break;
22          v = pred[match[v]];
23      }
24      return v;
25  }
26  inline void RT(int u) {
27      int v;
28      while (base[u] != nbase) {
29          v = match[u];
30          inB[base[u]] = inB[base[v]] = 1;
31          u = pred[v];
32          if (base[u] != nbase) pred[u] = v;
33      }
34  }
35  inline void BC(int u, int v) {
36      nbase = FindCA(u, v);
37      fill(inB, inB + n, 0);
38      RT(u); RT(v);
39      if (base[u] != nbase) pred[u] = v;
40      if (base[v] != nbase) pred[v] = u;
41      for (int i = 0; i < n; ++i)
42          if (inB[base[i]]) {
43              base[i] = nbase;
44              if (!inQ[i]) push(i);
45          }
46  }
47  bool FindAP(int u) {
48      bool found = false;
49      for (int i = 0; i < n; ++i) {
50          pred[i] = -1; base[i] = i; inQ[i] = 0;
51      }

```

```

52 sta = u; fin = -1; head = tail = 0; push(sta);
53 while (head < tail) {
54     int u = pop();
55     for (int i = (int)lnk[u].size() - 1; i >= 0; --i) {
56         int v = lnk[u][i];
57         if (base[u] != base[v] && match[u] != v) {
58             if (v == sta || match[v] >= 0 && pred[match[v]] >= 0) BC(u, v);
59             else if (pred[v] == -1) {
60                 pred[v] = u;
61                 if (match[v] >= 0) push(match[v]);
62                 else {
63                     fin = v;
64                     return true;
65                 }
66             }
67         }
68     }
69 }
70 return found;
71 }
72 inline void AP() {
73     int u = fin, v, w;
74     while (u >= 0) {
75         v = pred[u]; w = match[v];
76         match[v] = u; match[u] = v;
77         u = w;
78     }
79 }
80 inline int FindMax() {
81     for (int i = 0; i < n; ++i) match[i] = -1;
82     for (int i = 0; i < n; ++i)
83         if (match[i] == -1 && FindAP(i)) AP();
84     int ans = 0;
85     for (int i = 0; i < n; ++i) {
86         ans += (match[i] != -1);
87     }
88     return ans;
89 }

```

3.3 一般最大权匹配

```

1 //maximum weight blossom, change g[u][v].w to INF - g[u][v].w when minimum weight blossom
  ↪ is needed
2 //type of ans is long long
3 //replace all int to long long if weight of edge is long long

```

```

4
5 struct WeightGraph {
6     static const int INF = INT_MAX;
7     static const int MAXN = 400;
8     struct edge{
9         int u, v, w;
10        edge() {}
11        edge(int u, int v, int w): u(u), v(v), w(w) {}
12    };
13    int n, n_x;
14    edge g[MAXN * 2 + 1][MAXN * 2 + 1];
15    int lab[MAXN * 2 + 1];
16    int match[MAXN * 2 + 1], slack[MAXN * 2 + 1], st[MAXN * 2 + 1], pa[MAXN * 2 + 1];
17    int flower_from[MAXN * 2 + 1][MAXN+1], S[MAXN * 2 + 1], vis[MAXN * 2 + 1];
18    vector<int> flower[MAXN * 2 + 1];
19    queue<int> q;
20    inline int e_delta(const edge &e){ // does not work inside blossoms
21        return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2;
22    }
23    inline void update_slack(int u, int x){
24        if(!slack[x] || e_delta(g[u][x]) < e_delta(g[slack[x]][x]))
25            slack[x] = u;
26    }
27    inline void set_slack(int x){
28        slack[x] = 0;
29        for(int u = 1; u <= n; ++u)
30            if(g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
31                update_slack(u, x);
32    }
33    void q_push(int x){
34        if(x <= n)q.push(x);
35        else for(size_t i = 0; i < flower[x].size(); i++)
36            q_push(flower[x][i]);
37    }
38    inline void set_st(int x, int b){
39        st[x]=b;
40        if(x > n) for(size_t i = 0; i < flower[x].size(); ++i)
41            set_st(flower[x][i], b);
42    }
43    inline int get_pr(int b, int xr){
44        int pr = find(flower[b].begin(), flower[b].end(), xr) - flower[b].begin();
45        if(pr % 2 == 1){
46            reverse(flower[b].begin() + 1, flower[b].end());
47            return (int)flower[b].size() - pr;
48        } else return pr;
49    }

```



```

50 inline void set_match(int u, int v){
51     match[u]=g[u][v].v;
52     if(u > n){
53         edge e=g[u][v];
54         int xr = flower_from[u][e.u], pr=get_pr(u, xr);
55         for(int i = 0; i < pr; ++i)
56             set_match(flower[u][i], flower[u][i ^ 1]);
57         set_match(xr, v);
58         rotate(flower[u].begin(), flower[u].begin()+pr, flower[u].end());
59     }
60 }
61 inline void augment(int u, int v){
62     for(;;){
63         int xnv=st[match[u]];
64         set_match(u, v);
65         if(!xnv)return;
66         set_match(xnv, st[pa[xnv]]);
67         u=st[pa[xnv]], v=xnv;
68     }
69 }
70 inline int get_lca(int u, int v){
71     static int t=0;
72     for(++t; u || v; swap(u, v)){
73         if(u == 0)continue;
74         if(vis[u] == t)return u;
75         vis[u] = t;
76         u = st[match[u]];
77         if(u) u = st[pa[u]];
78     }
79     return 0;
80 }
81 inline void add_blossom(int u, int lca, int v){
82     int b = n + 1;
83     while(b <= n_x && st[b]) ++b;
84     if(b > n_x) ++n_x;
85     lab[b] = 0, S[b] = 0;
86     match[b] = match[lca];
87     flower[b].clear();
88     flower[b].push_back(lca);
89     for(int x = u, y; x != lca; x = st[pa[y]]) {
90         flower[b].push_back(x),
91         flower[b].push_back(y = st[match[x]]),
92         q_push(y);
93     }
94     reverse(flower[b].begin() + 1, flower[b].end());
95     for(int x = v, y; x != lca; x = st[pa[y]]) {

```

```

196         flower[b].push_back(x),
197         flower[b].push_back(y = st[match[x]]),
198         q_push(y);
199     }
200     set_st(b, b);
201     for(int x = 1; x <= n_x; ++x) g[b][x].w = g[x][b].w = 0;
202     for(int x = 1; x <= n; ++x) flower_from[b][x] = 0;
203     for(size_t i = 0; i < flower[b].size(); ++i){
204         int xs = flower[b][i];
205         for(int x = 1; x <= n_x; ++x)
206             if(g[b][x].w == 0 || e_delta(g[xs][x]) < e_delta(g[b][x]))
207                 g[b][x] = g[xs][x], g[x][b] = g[x][xs];
208         for(int x = 1; x <= n; ++x)
209             if(flower_from[xs][x]) flower_from[b][x] = xs;
210     }
211     set_slack(b);
212 }
213 inline void expand_blossom(int b){ // S[b] == 1
214     for(size_t i = 0; i < flower[b].size(); ++i)
215         set_st(flower[b][i], flower[b][i]);
216     int xr = flower_from[b][g[b][pa[b]].u], pr = get_pr(b, xr);
217     for(int i = 0; i < pr; i += 2){
218         int xs = flower[b][i], xns = flower[b][i + 1];
219         pa[xs] = g[xns][xs].u;
220         S[xs] = 1, S[xns] = 0;
221         slack[xs] = 0, set_slack(xns);
222         q_push(xns);
223     }
224     S[xr] = 1, pa[xr] = pa[b];
225     for(size_t i = pr + 1; i < flower[b].size(); ++i){
226         int xs = flower[b][i];
227         S[xs] = -1, set_slack(xs);
228     }
229     st[b] = 0;
230 }
231 inline bool on_found_edge(const edge &e){
232     int u = st[e.u], v = st[e.v];
233     if(S[v] == -1){
234         pa[v] = e.u, S[v] = 1;
235         int nu = st[match[v]];
236         slack[v] = slack[nu] = 0;
237         S[nu] = 0, q_push(nu);
238     }else if(S[v] == 0){
239         int lca = get_lca(u, v);
240         if(!lca) return augment(u, v), augment(v, u), true;
241         else add_blossom(u, lca, v);

```

```

142     }
143     return false;
144 }
145 inline bool matching(){
146     memset(S + 1, -1, sizeof(int) * n_x);
147     memset(slack + 1, 0, sizeof(int) * n_x);
148     q = queue<int>();
149     for(int x = 1; x <= n_x; ++x)
150         if(st[x] == x && !match[x]) pa[x]=0, S[x]=0, q_push(x);
151     if(q.empty())return false;
152     for(;;){
153         while(q.size()){
154             int u = q.front();q.pop();
155             if(S[st[u]] == 1)continue;
156             for(int v = 1; v <= n; ++v)
157                 if(g[u][v].w > 0 && st[u] != st[v]){
158                     if(e_delta(g[u][v]) == 0){
159                         if(on_found_edge(g[u][v]))return true;
160                     }else update_slack(u, st[v]);
161                 }
162         }
163         int d = INF;
164         for(int b = n + 1; b <= n_x; ++b)
165             if(st[b] == b && S[b] == 1)d = min(d, lab[b]/2);
166         for(int x = 1; x <= n_x; ++x)
167             if(st[x] == x && slack[x]){
168                 if(S[x] == -1)d = min(d, e_delta(g[slack[x]][x]));
169                 else if(S[x] == 0)d = min(d, e_delta(g[slack[x]][x])/2);
170             }
171         for(int u = 1; u <= n; ++u){
172             if(S[st[u]] == 0){
173                 if(lab[u] <= d)return 0;
174                 lab[u] -= d;
175             }else if(S[st[u]] == 1)lab[u] += d;
176         }
177         for(int b = n+1; b <= n_x; ++b)
178             if(st[b] == b){
179                 if(S[st[b]] == 0) lab[b] += d * 2;
180                 else if(S[st[b]] == 1) lab[b] -= d * 2;
181             }
182         q=queue<int>();
183         for(int x = 1; x <= n_x; ++x)
184             if(st[x] == x && slack[x] && st[slack[x]] != x && e_delta(g[slack[x]][x]) ==
185                ↪ 0)
186                 if(on_found_edge(g[slack[x]][x]))return true;
187         for(int b = n + 1; b <= n_x; ++b)

```

```

187         if(st[b] == b && S[b] == 1 && lab[b] == 0) expand_blossom(b);
188     }
189     return false;
190 }
191 inline pair<long long, int> solve(){
192     memset(match + 1, 0, sizeof(int) * n);
193     n_x = n;
194     int n_matches = 0;
195     long long tot_weight = 0;
196     for(int u = 0; u <= n; ++u) st[u] = u, flower[u].clear();
197     int w_max = 0;
198     for(int u = 1; u <= n; ++u)
199         for(int v = 1; v <= n; ++v){
200             flower_from[u][v] = (u == v ? u : 0);
201             w_max = max(w_max, g[u][v].w);
202         }
203     for(int u = 1; u <= n; ++u) lab[u] = w_max;
204     while(matching()) ++n_matches;
205     for(int u = 1; u <= n; ++u)
206         if(match[u] && match[u] < u)
207             tot_weight += g[u][match[u]].w;
208     return make_pair(tot_weight, n_matches);
209 }
210 inline void init(){
211     for(int u = 1; u <= n; ++u)
212         for(int v = 1; v <= n; ++v)
213             g[u][v] = edge(u, v, 0);
214 }
215 };

```

3.4 无向图最小割

```

1  /*
2   * Stoer Wagner 全局最小割  $O(V^3)$ 
3   * lbase, 点数 n, 邻接矩阵 edge[MAXN][MAXN]
4   * 返回值为全局最小割
5   */
6
7  int StoerWagner() {
8      static int v[MAXN], wage[MAXN];
9      static bool vis[MAXN];
10
11      for (int i = 1; i <= n; ++i) v[i] = i;
12
13      int res = INF;

```

```

14
15     for (int nn = n; nn > 1; --nn) {
16         memset(vis, 0, sizeof(bool) * (nn + 1));
17         memset(wage, 0, sizeof(int) * (nn + 1));
18
19         int pre, last = 1; // vis[1] = 1;
20
21         for (int i = 1; i < nn; ++i) {
22             pre = last; last = 0;
23             for (int j = 2; j <= nn; ++j) if (!vis[j]) {
24                 wage[j] += edge[v[pre]][v[j]];
25                 if (!last || wage[j] > wage[last]) last = j;
26             }
27             vis[last] = 1;
28         }
29
30         res = std::min(res, wage[last]);
31
32         for (int i = 1; i <= nn; ++i) {
33             edge[v[i]][v[pre]] += edge[v[last]][v[i]];
34             edge[v[pre]][v[i]] += edge[v[last]][v[i]];
35         }
36         v[last] = v[nn];
37     }
38     return res;
39 }

```

3.5 必经点 dominator tree

```

1 //solve(s, n, raw_g): s is the root and base accords to base of raw_g
2 //idom[x] will be x if x does not have a dominator, and will be -1 if x is not reachable from
   ↪ s.
3
4 struct dominator_tree {
5     int base, dfn[N], sdom[N], idom[N], id[N], f[N], fa[N], smin[N], stamp;
6     Graph *g;
7     void predfs(int u) {
8         id[dfn[u] = stamp++] = u;
9         for (int i = g->adj[u]; ~i; i = g->nxt[i]) {
10             int v = g->v[i];
11             if (dfn[v] < 0) {
12                 f[v] = u;
13                 predfs(v);
14             }
15         }
16     }
17 }

```

```

16 }
17 int getfa(int u) {
18     if (fa[u] == u) return u;
19     int ret = getfa(fa[u]);
20     if (dfn[sdom[smin[fa[u]]]] < dfn[sdom[smin[u]]])
21         smin[u] = smin[fa[u]];
22     return fa[u] = ret;
23 }
24 void solve (int s, int n, Graph *raw_graph) {
25     g = raw_graph;
26     base = g -> base;
27     memset(dfn + base, -1, sizeof(*dfn) * n);
28     memset(idom + base, -1, sizeof(*idom) * n);
29     static Graph pred, tmp;
30     pred.init(base, n);
31     for (int i = 0; i < n; ++i) {
32         for (int p = g -> adj[i + base]; ~p; p = g -> nxt[p])
33             pred.ins(g -> v[p], i + base);
34     }
35     stamp = 0; tmp.init(base, n); predfs(s);
36     for (int i = 0; i < stamp; ++i) {
37         fa[id[i]] = smin[id[i]] = id[i];
38     }
39     for (int o = stamp - 1; o >= 0; --o) {
40         int x = id[o];
41         if (o) {
42             sdom[x] = f[x];
43             for (int i = pred.adj[x]; ~i; i = pred.nxt[i]) {
44                 int p = pred.v[i];
45                 if (dfn[p] < 0) continue;
46                 if (dfn[p] > dfn[x]) {
47                     getfa(p);
48                     p = sdom[smin[p]];
49                 }
50                 if (dfn[sdom[x]] > dfn[p]) sdom[x] = p;
51             }
52             tmp.ins(sdom[x], x);
53         }
54         while (~tmp.adj[x]) {
55             int y = tmp.v[tmp.adj[x]];
56             tmp.adj[x] = tmp.nxt[tmp.adj[x]];
57             getfa(y);
58             if (x != sdom[smin[y]]) idom[y] = smin[y];
59             else idom[y] = x;
60         }
61         for (int i = g -> adj[x]; ~i; i = g -> nxt[i])

```

```
62         if (f[g -> v[i]] == x) fa[g -> v[i]] = x;
63     }
64     idom[s] = s;
65     for (int i = 1; i < stamp; ++i) {
66         int x = id[i];
67         if (idom[x] != sdom[x]) idom[x] = idom[idom[x]];
68     }
69 }
70 };
```


Chapter 4

技巧

4.1 无敌的读入优化

```
1 // getchar() 读入优化 << 关同步 cin << 此优化
2 // 用 isdigit() 会小幅变慢
3 // 返回 false 表示读到文件尾
4 namespace Reader {
5     const int L = (1 << 15) + 5;
6     char buffer[L], *S, *T;
7     __inline bool getchar(char &ch) {
8         if (S == T) {
9             T = (S = buffer) + fread(buffer, 1, L, stdin);
10            if (S == T) {
11                ch = EOF;
12                return false;
13            }
14        }
15        ch = *S++;
16        return true;
17    }
18    __inline bool getint(int &x) {
19        char ch; bool neg = 0;
20        for (; getchar(ch) && (ch < '0' || ch > '9'); ) neg ^= ch == '-';
21        if (ch == EOF) return false;
22        x = ch - '0';
23        for (; getchar(ch), ch >= '0' && ch <= '9'; )
24            x = x * 10 + ch - '0';
25        if (neg) x = -x;
26        return true;
27    }
28 }
```

4.2 真正释放 STL 内存

```
1 template <typename T>
2 __inline void clear(T& container) {
3     container.clear(); // 或者删除了一堆元素
4     T(container).swap(container);
5 }
```