Geimoire'l Standard Code Library*

Shanghai Jiao Tong University

Dated: 2017 年 8 月 24 日

^{*}https://github.com/kzoacn/Grimoire

目录

1	代数	Į.
	1.1 $O(n^2 \log n)$ 求线性递推数列第 n 项	
	1.2 闪电数论变换与魔力 CRT	
	1.3 多项式求逆	
	1.4 多项式除法	
	1.5 多项式取指数取对数	
2	数论	1:
	2.1 大整数相乘取模	1
	2.2 线段下整点	1
	2.3 中国剩余定理	1
3	图论	1:
	3.1 一般图匹配	13
	3.2 一般最大权匹配	
	3.3 无向图最小割	19
4	技巧	2
	4.1 无敌的读入优化	2
	4.2 直下释放 STL 内存	

4 目录

Chapter 1

代数

1.1 $O(n^2 \log n)$ 求线性递推数列第 n 项

```
Given a_0, a_1, \cdots, a_{m-1}

a_n = c_0 * a_{n-m} + \cdots + c_{m-1} * a_0

a_0 is the nth element, \cdots, a_{m-1} is the n+m-1th element
```

```
void linear_recurrence(long long n, int m, int a[], int c[], int p) {
      long long v[M] = \{1 \% p\}, u[M << 1], msk = !!n;
2
       for(long long i(n); i > 1; i >>= 1) {
           msk <<= 1;
       for(long long x(0); msk; msk >>= 1, x <<= 1) {
6
7
           fill_n(u, m << 1, 0);
           int b(!!(n & msk));
8
           x \mid = b;
9
           if(x < m) {
               u[x] = 1 \% p;
11
12
               for(int i(0); i < m; i++) {</pre>
                    for(int j(0), t(i + b); j < m; j++, t++) {
14
                        u[t] = (u[t] + v[i] * v[j]) % p;
15
                    }
16
               }
               for(int i((m << 1) - 1); i >= m; i--) {
18
                    for(int j(0), t(i - m); j < m; j++, t++) {
19
                        u[t] = (u[t] + c[j] * u[i]) % p;
20
                    }
21
               }
22
           }
23
           copy(u, u + m, v);
24
25
       //a[n] = v[0] * a[0] + v[1] * a[1] + ... + v[m - 1] * a[m - 1].
26
      for(int i(m); i < 2 * m; i++) {</pre>
27
           a[i] = 0;
28
```

6 CHAPTER 1. 代数

```
for(int j(0); j < m; j++) {
29
                a[i] = (a[i] + (long long)c[j] * a[i + j - m]) % p;
30
31
32
       for(int j(0); j < m; j++) {
33
           b[i] = 0;
34
            for(int i(0); i < m; i++) {</pre>
35
                b[j] = (b[j] + v[i] * a[i + j]) % p;
36
            }
37
38
       for(int j(0); j < m; j++) {
39
40
           a[j] = b[j];
       }
41
42 | }
```

1.2 闪电数论变换与魔力 CRT

```
|\#define meminit(A, 1, r) memset(A + (1), 0, sizeof(*A) * ((r) - (1)))
  #define memcopy(B, A, 1, r) memcpy(B, A + (1), sizeof(*A) * ((r) - (1)))
2
  void DFT(int *a, int n, int f) { //f=1 逆 DFT
      for (register int i = 0, j = 0; i < n; i++) {
           if (i > j) std::swap(a[i], a[j]);
          for (register int t = n >> 1; (j ^= t) < t; t >>= 1);
6
7
      for (register int i = 2; i <= n; i <<= 1) {
8
           static int exp[MAXN];
9
           exp[0] = 1; exp[1] = fpm(PRT, (MOD - 1) / i, MOD);
           if (f == 1) \exp[1] = fpm(\exp[1], MOD - 2, MOD);
11
           for (register int k = 2; k < (i >> 1); k++) {
12
               \exp[k] = 111 * \exp[k - 1] * \exp[1] % MOD;
13
14
          for (register int j = 0; j < n; j += i) {
15
               for (register int k = 0; k < (i >> 1); k++) {
16
                   register int &pA = a[j + k], &pB = a[j + k + (i >> 1)];
                   register long long B = 111 * pB * exp[k];
18
                   pB = (pA - B) \% MOD;
19
                   pA = (pA + B) \% MOD;
20
               }
21
          }
22
23
      if (f == 1) {
24
          register int rev = fpm(n, MOD - 2, MOD);
25
          for (register int i = 0; i < n; i++) {
26
               a[i] = 111 * a[i] * rev % MOD;
27
               if (a[i] < 0) { a[i] += MOD; }</pre>
28
```

1.3. 多项式求逆 7

```
}
29
      }
30
31 | }
  // 在不写高精度的情况下合并 FFT 所得结果对 MOD 取模后的答案
32
33 /// 值得注意的是,这个东西不能最后再合并,而是应该每做一次多项式乘法就 CRT 一次
34 int CRT(int *a) {
      static int x[3];
35
      for (int i = 0; i < 3; i++) {
36
          x[i] = a[i];
37
          for (int j = 0; j < i; j++) {
38
              int t = (x[i] - x[j] + FFT[i] \rightarrow MOD) \% FFT[i] \rightarrow MOD;
39
              if (t < 0) t += FFT[i] -> MOD;
40
              x[i] = 1LL * t * inv[j][i] % FFT[i] -> MOD;
41
          }
42
      }
43
      int sum = 1, ret = x[0] % MOD;
      for (int i = 1; i < 3; i ++) {
45
          sum = 1LL * sum * FFT[i - 1] \rightarrow MOD % MOD;
46
          ret += 1LL * x[i] * sum % MOD;
47
          if(ret >= MOD) ret -= MOD;
48
49
      return ret;
50
51 }
  for (int i = 0; i < 3; i++) // inv 数组的预处理过程, inverse(x, p) 表示求 x 在 p 下逆元
52
      for (int j = 0; j < 3; j++)
53
          inv[i][j] = inverse(FFT[i] -> MOD, FFT[j] -> MOD);
```

1.3 多项式求逆

Given polynomial a and n, b is the polynomial such that $a * b \equiv 1 \pmod{x^n}$

```
void getInv(int *a, int *b, int n) {
      static int tmp[MAXN];
2
      b[0] = fpm(a[0], MOD - 2, MOD);
3
      for (int c = 2, M = 1; c < (n << 1); c <<= 1) {
4
5
           for (; M \le 3 * (c - 1); M \le 1);
           meminit(b, c, M);
6
           meminit(tmp, c, M);
7
           memcopy(tmp, a, 0, c);
8
           DFT(tmp, M, 0);
9
10
           DFT(b, M, 0);
           for (int i = 0; i < M; i++) {
11
               b[i] = 111 * b[i] * (211 - 111 * tmp[i] * b[i] % MOD + MOD) % MOD;
12
13
           DFT(b, M, 1);
14
           meminit(b, c, M);
15
```

CHAPTER 1. 代数

```
16 }
17 }
```

1.4 多项式除法

d is quotient and r is remainder

```
void divide(int n, int m, int *a, int *b, int *d, int *r) { // n \in M 分别为多项式 A (被除数)
     \rightarrow 和 B (除数)的指数 + 1
      static int M, tA[MAXN], tB[MAXN], inv[MAXN], tD[MAXN];
2
      for (; n > 0 \&\& a[n - 1] == 0; n--);
3
      for (; m > 0 \&\& b[m - 1] == 0; m--);
       for (int i = 0; i < n; i++) tA[i] = a[n - i - 1];
5
       for (int i = 0; i < m; i++) tB[i] = b[m - i - 1];
6
       for (M = 1; M \le n - m + 1; M \le 1);
7
      if (m < M) meminit(tB, m, M);</pre>
8
       getInv(tB, inv, M);
9
      for (M = 1; M \le 2 * (n - m + 1); M \le 1);
10
      meminit(inv, n - m + 1, M);
11
      meminit(tA, n - m + 1, M);
12
      DFT(inv, M, 0);
13
      DFT(tA, M, 0);
14
       for (int i = 0; i < M; i++) {
15
           d[i] = 111 * inv[i] * tA[i] % MOD;
16
17
      DFT(d, M, 1);
18
19
       std::reverse(d, d + n - m + 1);
      for (M = 1; M \le n; M \le 1);
20
      memcopy(tB, b, 0, m);
21
       if (m < M) meminit(tB, m, M);</pre>
22
      memcopy(tD, d, 0, n - m + 1);
23
      meminit(tD, n - m + 1, M);
24
      DFT(tD, M, 0);
      DFT(tB, M, 0);
26
      for (int i = 0; i < M; i++) {
27
           r[i] = 111 * tD[i] * tB[i] % MOD;
28
      }
29
      DFT(r, M, 1);
30
      meminit(r, n, M);
31
       for (int i = 0; i < n; i++) {
32
           r[i] = (a[i] - r[i] + MOD) % MOD;
33
34
35 | }
```

1.5. 多项式取指数取对数 9

1.5 多项式取指数取对数

Given polynomial a and n, b is the polynomial such that $b \equiv e^a \pmod{x^n}$ or $b \equiv \ln a \pmod{x^n}$

```
void getDiff(int *a, int *b, int n) { // 多项式取微分
1
      for (int i = 0; i + 1 < n; i++) {
2
           b[i] = 111 * (i + 1) * a[i + 1] % MOD;
3
      b[n - 1] = 0;
5
6
  |}
  void getInt(int *a, int *b, int n) { // 多项式取积分, 积分常数为 0
7
8
      static int inv[MAXN];
      inv[1] = 1;
9
      for (int i = 2; i < n; i++) {
           inv[i] = 111 * (MOD - MOD / i) * inv[MOD % i] % MOD;
11
12
      b[0] = 0;
13
      for (int i = 1; i < n; i++) {
14
           b[i] = 111 * a[i - 1] * inv[i] % MOD;
15
16
  }
17
  void getLn(int *a, int *b, int n) {
18
      static int inv[MAXN], d[MAXN];
19
      int M = 1;
20
      for (; M \le 2 * (n - 1); M \le 1);
21
22
      getInv(a, inv, n);
      getDiff(a, d, n);
23
      meminit(d, n, M);
      meminit(inv, n, M);
25
      DFT(d, M, 0); DFT(inv, M, 0);
26
      for (int i = 0; i < M; i++) {
27
           d[i] = 111 * d[i] * inv[i] % MOD;
28
29
      DFT(d, M, 1);
30
      getInt(d, b, n);
31
  }
32
33
  void getExp(int *a, int *b, int n) {
      static int ln[MAXN], tmp[MAXN];
34
      b[0] = 1;
35
      for (int c = 2, M = 1; c < (n << 1); c <<= 1) {
36
           for (; M <= 2 * (c - 1); M <<= 1);
37
           int bound = std::min(c, n);
38
           memcopy(tmp, a, 0, bound);
39
           meminit(tmp, bound, M);
40
           meminit(b, c, M);
41
           getLn(b, ln, c);
42
           meminit(ln, c, M);
43
```

10 CHAPTER 1. 代数

```
DFT(b, M, 0);
44
           DFT(tmp, M, 0);
45
           DFT(ln, M, 0);
46
           for (int i = 0; i < M; i++) {
47
               b[i] = 111 * b[i] * (111 - ln[i] + tmp[i] + MOD) % MOD;
48
           }
49
          DFT(b, M, 1);
50
           meminit(b, c, M);
51
       }
52
53 }
```

Chapter 2

数论

2.1 大整数相乘取模

```
// x 与 y 须非负
long long mult(long long x, long long y, long long MODN) {
    long long t = (x * y - (long long)((long double)x / MODN * y + 1e-3) * MODN) % MODN;
    return t < 0 ? t + MODN : t;
}
```

2.2 线段下整点

```
solve for \sum_{i=0}^{n-1} \lfloor \frac{a+bi}{m} \rfloor, n, m, a, b > 0
```

```
LL solve(LL n,LL a,LL b,LL m){
    if(b==0) return n*(a/m);
    if(a>=m) return n*(a/m)+solve(n,a%m,b,m);
    if(b>=m) return (n-1)*n/2*(b/m)+solve(n,a,b%m,m);
    return solve((a+b*n)/m,(a+b*n)%m,m,b);
}
```

2.3 中国剩余定理

first is remainder, second is module

```
inline void fix(LL &x, LL y) {
      x = (x \% y + y) \% y;
2
 }
3
 bool solve(int n, std::pair<LL, LL> a[],
                    std::pair<LL, LL> &ans) {
5
      ans = std::make_pair(1, 1);
6
7
      for (int i = 0; i < n; ++i) {
          LL num, y;
8
9
          euclid(ans.second, a[i].second, num, y);
```

12 CHAPTER 2. 数论

```
LL divisor = std::__gcd(ans.second, a[i].second);
10
          if ((a[i].first - ans.first) % divisor) {
11
               return false;
12
          }
13
          num *= (a[i].first - ans.first) / divisor;
14
          fix(num, a[i].second);
15
          ans.first += ans.second * num;
16
          ans.second *= a[i].second / divisor;
17
          fix(ans.first, ans.second);
18
19
      return true;
20
21 }
```

Chapter 3

图论

3.1 一般图匹配

```
1 // 0-base, match[u] is linked to u
vector<int> lnk[MAXN];
int match[MAXN], Queue[MAXN], pred[MAXN], base[MAXN], head, tail, sta, fin, nbase;
4 bool inQ[MAXN], inB[MAXN];
5 inline void push(int u) {
      Queue[tail++] = u; inQ[u] = 1;
  }
7
8
  inline int pop() {
      return Queue[head++];
9
10 }
  inline int FindCA(int u, int v) {
11
      static bool inP[MAXN];
12
      fill(inP, inP + n, false);
13
      while (1) {
14
           u = base[u]; inP[u] = 1;
15
           if(u == sta) break;
16
           u = pred[match[u]];
17
18
      while (1) {
19
           v = base[v];
20
           if (inP[v]) break;
21
           v = pred[match[v]];
22
23
      return v;
24
  }
25
  inline void RT(int u) {
26
      int v;
27
      while (base[u] != nbase) {
28
           v = match[u];
29
           inB[base[u]] = inB[base[v]] = 1;
30
          u = pred[v];
31
```

14 CHAPTER 3. 图论

```
if (base[u] != nbase) pred[u] = v;
32
       }
33
 }
34
  inline void BC(int u, int v) {
35
       nbase = FindCA(u, v);
36
       fill(inB, inB + n, 0);
37
      RT(u); RT(v);
38
       if (base[u] != nbase) pred[u] = v;
39
       if (base[v] != nbase) pred[v] = u;
40
       for (int i = 0; i < n; ++i)
41
           if (inB[base[i]]) {
42
43
               base[i] = nbase;
               if (!inQ[i]) push(i);
44
           }
45
  }
46
  bool FindAP(int u) {
47
       bool found = false;
48
       for (int i = 0; i < n; ++i) {
49
           pred[i] = -1; base[i] = i; inQ[i] = 0;
50
51
       sta = u; fin = -1; head = tail = 0; push(sta);
52
       while (head < tail) {</pre>
53
           int u = pop();
54
           for (int i = (int)lnk[u].size() - 1; i >= 0; --i) {
55
                int v = lnk[u][i];
56
                if (base[u] != base[v] && match[u] != v) {
57
                    if (v == sta \mid \mid match[v] >= 0 \&\& pred[match[v]] >= 0) BC(u, v);
58
                    else if (pred[v] == -1) {
59
                        pred[v] = u;
60
                        if (match[v] >= 0) push(match[v]);
61
                        else {
62
                             fin = v;
63
                             return true;
64
65
                        }
                    }
66
               }
67
           }
68
       }
69
       return found;
70
  }
71
  inline void AP() {
72
       int u = fin, v, w;
73
       while (u \ge 0) {
74
           v = pred[u]; w = match[v];
75
           match[v] = u; match[u] = v;
76
           u = w;
77
```

3.2. 一般最大权匹配 15

```
}
78
  }
79
  inline int FindMax() {
80
       for (int i = 0; i < n; ++i) match[i] = -1;
81
       for (int i = 0; i < n; ++i)
82
           if (match[i] == -1 \&\& FindAP(i)) AP();
83
       int ans = 0;
84
       for (int i = 0; i < n; ++i) {
85
           ans += (match[i] != -1);
86
87
       return ans;
88
89
  }
```

3.2 一般最大权匹配

```
//maximum weight blossom, change g[u][v].w to INF - g[u][v].w when minimum weight blossom is
     \hookrightarrow needed
  //type of ans is long long
  //replace all int to long long if weight of edge is long long
  struct WeightGraph {
       static const int INF = INT_MAX;
6
      static const int MAXN = 400;
7
8
      struct edge{
           int u,v,w;
9
           edge(){}
           edge(int u,int v,int w):u(u),v(v),w(w){}
11
      };
12
       int n,n_x;
13
       edge g[MAXN*2+1][MAXN*2+1];
       int lab[MAXN*2+1];
15
      int match[MAXN*2+1],slack[MAXN*2+1],st[MAXN*2+1],pa[MAXN*2+1];
16
       int flower_from[MAXN*2+1][MAXN+1],S[MAXN*2+1],vis[MAXN*2+1];
17
      vector<int> flower[MAXN*2+1];
18
      queue<int> q;
19
      inline int e_delta(const edge &e){ // does not work inside blossoms
20
           return lab[e.u]+lab[e.v]-g[e.u][e.v].w*2;
21
22
       inline void update_slack(int u,int x){
23
           if(!slack[x]||e_delta(g[u][x])<e_delta(g[slack[x]][x]))slack[x]=u;</pre>
24
25
      inline void set_slack(int x){
26
           slack[x]=0;
27
           for(int u=1;u<=n;++u)</pre>
28
               if(g[u][x].w>0\&\&st[u]!=x\&\&S[st[u]]==0)update_slack(u,x);
29
```

16 CHAPTER 3. 图论

```
}
30
       void q_push(int x){
31
           if(x \le n)q.push(x);
32
           else for(size_t i=0;i<flower[x].size();i++)q_push(flower[x][i]);</pre>
33
34
       inline void set_st(int x,int b){
35
           st[x]=b;
36
           if(x>n)for(size_t i=0;i<flower[x].size();++i)</pre>
37
                    set_st(flower[x][i],b);
38
39
       inline int get_pr(int b,int xr){
40
           int pr=find(flower[b].begin(),flower[b].end(),xr)-flower[b].begin();
41
           if(pr%2==1){
42
               reverse(flower[b].begin()+1,flower[b].end());
43
                return (int)flower[b].size()-pr;
44
           }else return pr;
45
46
       inline void set_match(int u,int v){
47
           match[u]=g[u][v].v;
48
           if(u>n){
49
                edge e=g[u][v];
50
                int xr=flower_from[u][e.u],pr=get_pr(u,xr);
51
                for(int i=0;i<pr;++i)set_match(flower[u][i],flower[u][i^1]);</pre>
52
                set_match(xr,v);
53
                rotate(flower[u].begin(),flower[u].begin()+pr,flower[u].end());
54
           }
55
       }
56
       inline void augment(int u,int v){
57
           for(;;){
58
               int xnv=st[match[u]];
59
                set_match(u,v);
60
                if(!xnv)return;
61
                set_match(xnv,st[pa[xnv]]);
62
                u=st[pa[xnv]],v=xnv;
63
           }
64
       }
65
       inline int get_lca(int u,int v){
66
           static int t=0;
67
           for(++t;u||v;swap(u,v)){
68
                if (u==0) continue;
69
                if(vis[u]==t)return u;
70
               vis[u]=t;
71
               u=st[match[u]];
72
                if(u)u=st[pa[u]];
73
74
           return 0;
75
```

3.2. 一般最大权匹配 17

```
76
       inline void add_blossom(int u,int lca,int v){
77
            int b=n+1;
78
            while (b \le n_x \& st[b]) + +b;
79
            if(b>n_x)++n_x;
80
            lab[b]=0,S[b]=0;
81
            match[b]=match[lca];
82
            flower[b].clear();
83
            flower[b].push_back(lca);
84
            for(int x=u,y;x!=lca;x=st[pa[y]])
85
                flower[b].push_back(x),flower[b].push_back(y=st[match[x]]),q_push(y);
86
87
            reverse(flower[b].begin()+1,flower[b].end());
            for(int x=v,y;x!=lca;x=st[pa[y]])
88
                flower[b].push_back(x),flower[b].push_back(y=st[match[x]]),q_push(y);
89
            set_st(b,b);
90
            for(int x=1;x\leq n_x;++x)g[b][x].w=g[x][b].w=0;
91
            for(int x=1;x<=n;++x)flower_from[b][x]=0;</pre>
92
            for(size_t i=0;i<flower[b].size();++i){</pre>
93
                int xs=flower[b][i];
94
                for(int x=1;x\leq n_x;++x)
95
                     if(g[b][x].w==0||e_delta(g[xs][x]) < e_delta(g[b][x]))
96
                         g[b][x]=g[xs][x],g[x][b]=g[x][xs];
97
                for(int x=1;x\leq n;++x)
98
                     if(flower_from[xs][x])flower_from[b][x]=xs;
            }
100
            set_slack(b);
101
       }
       inline void expand_blossom(int b){ // S[b] == 1
103
            for(size_t i=0;i<flower[b].size();++i)</pre>
104
                set_st(flower[b][i],flower[b][i]);
105
            int xr=flower_from[b][g[b][pa[b]].u],pr=get_pr(b,xr);
106
            for(int i=0;i<pr;i+=2){</pre>
                int xs=flower[b][i],xns=flower[b][i+1];
108
                pa[xs]=g[xns][xs].u;
109
                S[xs]=1,S[xns]=0;
                slack[xs]=0,set_slack(xns);
111
                q_push(xns);
112
            }
113
            S[xr]=1,pa[xr]=pa[b];
114
            for(size_t i=pr+1;i<flower[b].size();++i){</pre>
                int xs=flower[b][i];
                S[xs]=-1, set_slack(xs);
117
            }
118
            st[b]=0;
119
       }
120
       inline bool on_found_edge(const edge &e){
121
```

18 CHAPTER 3. 图论

```
int u=st[e.u],v=st[e.v];
122
            if(S[v]==-1){
123
                 pa[v]=e.u,S[v]=1;
124
                 int nu=st[match[v]];
125
                 slack[v]=slack[nu]=0;
126
                 S[nu]=0,q_push(nu);
127
            else if(S[v]==0){
128
                 int lca=get_lca(u,v);
129
                 if(!lca)return augment(u,v),augment(v,u),true;
130
                 else add_blossom(u,lca,v);
131
            }
132
133
            return false;
       }
134
        inline bool matching(){
135
            memset(S+1,-1,sizeof(int)*n_x);
136
            memset(slack+1,0,sizeof(int)*n_x);
137
            q=queue<int>();
138
            for(int x=1;x<=n_x;++x)</pre>
139
                 if (st[x]==x\&\&!match[x])pa[x]=0,S[x]=0,q_push(x);
140
            if(q.empty())return false;
141
            for(;;){
142
                while(q.size()){
143
                     int u=q.front();q.pop();
144
                     if(S[st[u]]==1)continue;
                     for(int v=1; v<=n;++v)</pre>
146
                          if(g[u][v].w>0&&st[u]!=st[v]){
                              if(e_delta(g[u][v])==0){
148
                                   if(on_found_edge(g[u][v]))return true;
                              }else update_slack(u,st[v]);
150
                          }
151
                 }
152
                 int d=INF;
                 for(int b=n+1;b<=n_x;++b)</pre>
154
                     if(st[b]==b\&\&S[b]==1)d=min(d,lab[b]/2);
                 for(int x=1;x\leq n_x;++x)
156
                     if(st[x]==x&&slack[x]){
                          if(S[x]==-1)d=min(d,e_delta(g[slack[x]][x]));
158
                          else if(S[x]==0)d=min(d,e_delta(g[slack[x]][x])/2);
159
                     }
160
                 for(int u=1;u<=n;++u){</pre>
                     if(S[st[u]]==0){
162
                          if(lab[u]<=d)return 0;</pre>
163
                          lab[u]-=d;
164
                     }else if(S[st[u]]==1)lab[u]+=d;
165
                 }
166
                 for(int b=n+1;b<=n_x;++b)</pre>
167
```

3.3. 无向图最小割 19

```
if(st[b]==b){
168
                           if(S[st[b]]==0)lab[b]+=d*2;
169
                           else if(S[st[b]]==1)lab[b]-=d*2;
170
                      }
171
                 q=queue<int>();
172
                 for(int x=1;x\leq x;++x)
173
                      if(st[x]==x\&\&slack[x]\&\&st[slack[x]]!=x\&\&e_delta(g[slack[x]][x])==0)
174
                           if(on_found_edge(g[slack[x]][x]))return true;
175
                 for(int b=n+1;b<=n_x;++b)</pre>
176
                      if(st[b]==b\&\&S[b]==1\&\&lab[b]==0)expand_blossom(b);
177
            }
178
            return false;
179
        }
180
        inline pair<long long,int> solve(){
181
            memset(match+1,0,sizeof(int)*n);
182
            n_x=n;
183
            int n_matches=0;
184
             long long tot_weight=0;
185
            for(int u=0;u<=n;++u)st[u]=u,flower[u].clear();</pre>
186
             int w_max=0;
187
            for(int u=1;u<=n;++u)</pre>
188
                 for(int v=1; v<=n; ++v){</pre>
189
                      flower_from[u][v]=(u==v?u:0);
190
                      w_{max}=max(w_{max},g[u][v].w);
191
192
            for(int u=1;u<=n;++u)lab[u]=w_max;</pre>
193
             while(matching())++n_matches;
194
             for(int u=1;u<=n;++u)</pre>
195
                 if (match[u] &&match[u] <u)</pre>
196
                      tot_weight+=g[u][match[u]].w;
197
            return make_pair(tot_weight,n_matches);
198
        }
        inline void init(){
200
            for(int u=1;u<=n;++u)</pre>
                 for(int v=1; v<=n; ++v)</pre>
202
                      g[u][v]=edge(u,v,0);
203
        }
204
   };
205
```

3.3 无向图最小割

```
/*

* Stoer Wagner 全局最小割 O(V ^ 3)

* 1base, 点数 n, 邻接矩阵 edge[MAXN][MAXN]

* 返回值为全局最小割
```

20 CHAPTER 3. 图论

```
5 */
6
  int StoerWagner() {
7
8
       static int v[MAXN], wage[MAXN];
       static bool vis[MAXN];
9
10
       for (int i = 1; i <= n; ++i) v[i] = i;</pre>
11
12
      int res = INF;
13
14
       for (int nn = n; nn > 1; --nn) {
15
           memset(vis, 0, sizeof(bool) * (nn + 1));
16
           memset(wage, 0, sizeof(int) * (nn + 1));
17
18
           int pre, last = 1; // vis[1] = 1;
19
20
           for (int i = 1; i < nn; ++i) {
21
               pre = last; last = 0;
22
               for (int j = 2; j <= nn; ++j) if (!vis[j]) {</pre>
23
                    wage[j] += edge[v[pre]][v[j]];
24
                    if (!last || wage[j] > wage[last]) last = j;
25
26
               vis[last] = 1;
27
           }
28
29
           res = std::min(res, wage[last]);
30
31
           for (int i = 1; i <= nn; ++i) {
32
               edge[v[i]][v[pre]] += edge[v[last]][v[i]];
33
               edge[v[pre]][v[i]] += edge[v[last]][v[i]];
34
35
           v[last] = v[nn];
36
37
38
       return res;
39 }
```

Chapter 4

技巧

4.1 无敌的读入优化

```
1 // getchar() 读入优化 << 关同步 cin << 此优化
  |// 用 isdigit() 会小幅变慢
3 // 返回 false 表示读到文件尾
  namespace Reader {
      const int L = (1 << 15) + 5;
5
      char buffer[L], *S, *T;
6
      __inline bool getchar(char &ch) {
7
          if (S == T) {
8
              T = (S = buffer) + fread(buffer, 1, L, stdin);
9
              if (S == T) {
10
                   ch = EOF;
11
                   return false;
12
              }
13
          }
          ch = *S++;
15
          return true;
16
17
      __inline bool getint(int &x) {
18
          char ch; bool neg = 0;
19
          for (; getchar(ch) && (ch < '0' || ch > '9'); ) neg ^= ch == '-';
20
          if (ch == EOF) return false;
21
          x = ch - '0';
22
          for (; getchar(ch), ch >= '0' && ch <= '9'; )
23
              x = x * 10 + ch - '0';
24
          if (neg) x = -x;
25
          return true;
26
      }
27
28 }
```

22 CHAPTER 4. 技巧

4.2 真正释放 STL 内存

```
template <typename T>
__inline void clear(T& container) {
    container.clear(); // 或者删除了一堆元素
    T(container).swap(container);
}
```