Gungnir's Standard Code Library

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Chapter 1 计算几何

1.1 二维

1.1.1 基础

```
typedef double DB;
   const DB eps = 1e-8;
    __inline int sign(DB x) {
       return x < -eps ? -1 : (x > eps ? 1 : 0);
    _inline DB msqrt(DB x) {
       return sign(x) > 0 ? sgrt(x) : 0;
11
   struct Point {
12
       DB x, y;
__inline Point(): x(0), y(0) {}
13
       <u>__inline Point(DB x, DB y): x(x), y(y) {}</u>
14
15
       __inline Point operator+(const Point &rhs) const {
16
            return Point(x + rhs.x, y + rhs.y);
17
       __inline Point operator-(const Point &rhs) const {
18
19
            return Point(x - rhs.x, y - rhs.y);
20
21
       __inline Point operator*(DB k) const {
22
23
           return Point(x * k, y * k);
24
25
26
27
       __inline Point operator/(DB k) const {
   assert(sign(k));
            return Point(x / k, y / k);
28
29
30
    _inline DB dot(const P& a, const P& b) {
31
       return a.x * b.x + a.y * b.y;
32 }
33
    _inline DB det(const P& a, const P& b) {
35
       return a.x * b.y - a.y * b.x;
```

1.1.2 凸包

```
inline void clear(std::vector<Point>& v) {
      v.clear();
      std::vector<Point>(v).swap(v);
   struct Convex {
      std::vector<Point> a, upper, lower;
       void make shell(const std::vector<Point>& p,
              std::vector<Point>& shell) { // p needs to be sorted.
10
           clear(shell); int n = p.size();
11
          12
13
14
15
              shell push back(p[i]);
16
17
18
      void make convex() {
          std::sort(a.begin(), a.end());
make_shell(a, lower);
19
20
21
22
23
          std::reverse(a.begin(), a.end());
          make_shell(a, upper);
24
          for (std::vector<Point>::iterator it = upper.begin(); it != upper.end(); it++)
25
              if (!(*it == *a.rbegin()) \&\& !(*it == *a.begin()))
26
                  a.push back(*it);
          n = a.size();
```

```
29
        void init(const std::vector<Point>& _a) {
30
            clear(a); a = _a; n = a.size();
make_convex();
31
32
33
        void read(int n) { // Won't make convex.
34
            clear(a); n = _n; a.resize(n);
            for (int i = 0; i < n; i++)
35
36
                 a[i].read():
37
38
        std::pair<DB, int> get_tangent(
39
                 const std::vecTor<Point>& convex, const Point& vec) {
            int l = 0, r = (int)convex.size() - 2;
40
41
            assert(r >= 0);
42
            for (; l + 1 < r; ) {
                 int mid = (l + r) / 2;
43
44
                 if (sign(det(convex[mid + 1] - convex[mid], vec)) > 0)
45
                     r = mid:
                 else l = mid;
46
47
            return std::max(std::make_pair(det(vec, convex[r]), r),
48
49
                     std::make pair(det(vec, convex[0]), 0));
50
51
        int binary_search(Point u, Point v, int l, int r) {
52
53
            int s1 = sign(det(v - u, a[l % n] - u));
            for (; l + 1 < r; ) {
int mid = (l + r) / 2;
54
55
                 int smid = sign(det(v - u, a[mid % n] - u));
56
                 if (smid == s1) l = mid;
57
                 else r = mid;
58
59
            return 1 % n;
60
        // 求凸包上和向量 vec 叉积最大的点,返回编号,共线的多个切点返回任意一个
61
62
        int get_tangent(Point vec) {
            std::pair<DB, int> ret = get_tangent(upper, vec);
ret.second = (ret.second + (int)lower.size() - 1) % n;
63
64
65
            ret = std::max(ret, get_tangent(lower, vec));
66
            return ret.second;
67
68
        // 求凸包和直线 u, v 的交点, 如果不相交返回 false, 如果有则是和 (i, next(i)) 的
      → 交点, 交在点上不确定返回前后两条边其中之一
       bool get_intersection(Point u, Point v, int &i0, int &i1) {
  int p0 = get_tangent(u - v), p1 = get_tangent(v - u);
  if (sign(det(v - u, a[p0] - u)) * sign(det(v - u, a[p1] - u)) <= 0) {
    if (p0 > p1) std::swap(p0, p1);
69
70
71
72
                 i0 = binary_search(u, v, p0, p1);
73
                 i1 = binary_search(u, v, p1, p0 + n);
74
75
                 return true:
76
77
            else return false;
78
        }
79 };
```

Chapter 2 数论

2.1 求逆元

```
void ex gcd(long long a, long long b, long long &x, long long &y) {
    if (b == 0) {
        x = 1;
        y = 0;
        return;
    }
    long long xx, yy;
    ex_gcd(b, a % b, xx, yy);
    y = xx - a / b * yy;
    x = yy;
}
long long inv(long long x, long long MODN) {
```

```
long long inv_x, y;
ex_gcd(x, MODN, inv_x, y);
return (inv_x % MODN + MODN) % MODN;
}
```

2.2 中国剩余定理

Chapter 3 图论

3.1 基础

```
struct Graph { // Remember to call .init()!
       int e, nxt[M], v[M], adj[N], n;
       bool base;
       __inline void init(bool _base, int _n = 0) {
            assert(n < N);</pre>
            n = _n; base = _base;
            e = \overline{0}; memset(a\overline{d}j + base, -1, sizeof(*adj) * n);
       __inline int new_node() {
            adj[n + base] = -1;
10
11
            assert(n + base + 1 < N);
            return n++ + base;
12
13
14
       __inline void ins(int u0, int v0) { // directional
15
            assert(u0 < n + base && v0 < n + base):
            v[e] = v0; nxt[e] = adj[u0]; adj[u0] = e++;
16
17
            assert(e < M);
18
       __inline void bi_ins(int u0, int v0) { // bi-directional
19
20
            ins(u0, v0); ins(v0, u0);
21
22 };
```

3.2 KM

```
struct KM
      // Truly 0(n^3)
      // 邻接矩阵,不能连的边设为 -INF, 求最小权匹配时边权取负, 但不能连的还是 -INF,
     → 使用时先对 1 -> n 调用 hungary() , 再 get_ans() 求值
      int w[N][N];
      int lx[N], ly[N], match[N], way[N], slack[N];
bool used[N];
      void init() {
          for (int i = 1; i <= n; i++) {
              match[i] = 0;
9
              lx[i] = 0;
10
11
              ly[i] = 0;
12
              way[i] = 0;
13
14
15
      void hungary(int x) {
16
          match[0] = x;
17
          int j0 = 0;
          for (int j = 0; j <= n; j++) {
```

```
19
                 slack[j] = INF;
20
                 used[j] = false;
21
22
            }
23
            do {
24
                 used[j0] = true;
25
26
                  int i0 = match[j0], delta = INF, j1 = 0;
                 for (int j = 1; j <= n; j++) {
   if (used[j] == false) {</pre>
27
                           int cur = -w[i0][j] - lx[i0] - ly[j];
28
29
                           if (cur < slack[j]) {</pre>
30
                               slack[i] = cur;
31
                                way[j] = j0;
32
33
                           if (slack[j] < delta) {</pre>
                               delta = slack[j];
j1 = j;
34
35
36
37
                      }
38
39
                 for (int j = 0; j <= n; j++) {
   if (used[j]) {</pre>
40
41
                           lx[match[j]] += delta;
42
                           lv[i] -= delta;
43
44
                      else slack[i] -= delta;
45
46
                  i0 = i1:
47
            } while (match[j0] != 0);
48
49
                  int j1 = way[j0];
50
51
                 match[i0] = match[i1];
52
                 j0 = j1;
53
            } while (j0);
        }
54
55
56
        int get_ans() {
57
             int sum = 0:
58
             for(int i = 1; i <= n; i++) {
59
                  if (w[match[i]][i] == -INF); // 无解
60
                  if (match[i] > 0) sum += w[match[i]][i];
61
62
             return sum;
63
64
   } km;
```

3.3 点双连通分量

bcc.forest is a set of connected tree whose vertices are chequered with cut-vertex and BCC.

```
const bool BCC_VERTEX = 0, BCC_EDGE = 1;
   struct BCC { // N = N0 + M0. Remember to call init(&raw_graph).
        Graph *g, forest; // g is raw graph ptr.
        int dfn[N], DFN, low[N];
5
        int stack[N], top;
                                  // Where edge i is expanded to in expaned graph.
6
        int expand_to[N];
        // Vertex \bar{i} expaned to i.
        int compress_to[N]; // Where vertex i is compressed to.
bool vertex_type[N], cut[N], compress_cut[N], branch[M];
10
        //std::vector<int> BCC_component[N]; // Cut vertex belongs to none.
        __inline void init(Graph *raw_graph) {
11
12
            g = raw_graph;
13
        void DFS(int u, int pe) {
   dfn[u] = low[u] = ++DFN; cut[u] = false;
14
15
16
            if (!~q->adi[u]) {
17
                 cut[u] = 1;
18
                 compress_to[u] = forest.new_node();
19
                 compress_cut[compress_to[u]] = 1;
20
```

```
21
            for (int e = g\rightarrow adj[u]; \sim e; e = g\rightarrow nxt[e]) {
22
                int v = g -> v[e];
23
                if ((e^pe) > 1 \& dfn[v] > 0 \& dfn[v] < dfn[u]) {
24
                    stack[top++] = e;
25
26
                    low[u] = std::min(low[u], dfn[v]);
27
                else if (!dfn[v]) {
28
                    stack[top++] = e; branch[e] = 1;
                    DFS(v, e);
low[u] = std::min(low[v], low[u]);
29
30
31
                    if (low[v] >= dfn[u]) {
32
                        if (!cut[u]) {
33
                             cut[u] = 1;
34
                             compress_to[u] = forest.new_node();
35
                             compress cut[compress to[u]] = 1;
36
37
                         int cc = forest.new_node();
38
                         forest.bi_ins(compress_to[u], cc);
39
                         compress_cut[cc] = 0;
                         //BCC_component[cc].clear();
40
41
                         do {
42
                             int cur_e = stack[--top];
43
                             compress_to[expand_to[cur_e]] = cc;
                             compress_to[expand_to[cur_e^1]] = cc;
44
45
                             if (branch[cur_e]) {
46
                                 int v = g->v[cur_e];
                                 if (cut[v])
47
48
                                      forest.bi_ins(cc, compress_to[v]);
49
                                 else {
50
                                     //BCC_component[cc].push_back(v);
51
                                     compress_to[v] = cc;
52
53
54
                        } while (stack[top] != e);
55
                   }
               }
56
57
58
59
       void solve() {
60
            forest.init(g->base);
61
            int n = g->n;
62
            for (int i = 0; i < g -> e; i ++) {
63
                expand_to[i] = g->new_node();
64
65
            memset(branch, 0, sizeof(*branch) * g->e);
            memset(dfn + g->base, 0, sizeof(*dfn) * n); DFN = 0;
66
67
            for (int i = 0; i < n; i++)
68
                if (!dfn[i + g->base]) {
69
                    top = 0;
70
                    DFS(i + q->base, -1);
71
72
73
74
   } bcc;
75
   bcc.init(&raw_graph);
76 | bcc.solve();
   // Do something with bcc.forest ...
```

3.4 边双连通分量

```
struct BCC {
Graph *g, forest;
int dfn[N], low[N], stack[N], tot[N], belong[N], vis[N], top, dfs_clock;
// tot[] is the size of each BCC, belong[] is the BCC that each node belongs to
pair<int, int > ori[M]; // bridge in raw_graph(raw node)
bool is_bridge[M];
__inline void init(Graph *raw_graph) {
    g = raw_graph;
    memset(is_bridge, false, sizeof(*is_bridge) * g -> e);
    memset(vis + g -> base, 0, sizeof(*vis) * g -> n);
}
void tarjan(int u, int from) {
```

```
dfn[u] = low[u] = ++dfs_clock; vis[u] = 1; stack[++top] = u;
for (int p = g -> adj[u]; ~p; p = g -> nxt[p]) {
   if ((p ^ 1) == from) continue;
13
14
15
16
                   int v = g \rightarrow v[p];
                   if (vis[v]) {
17
18
                         if (vis[v] == 1) low[u] = min(low[u], dfn[v]);
19
                   } else {
20
                        tarjan(v, p);
                        low[u] = min(low[u], low[v]);
if (low[v] > dfn[u]) is_bridge[p / 2] = true;
21
22
23
24
25
              if (dfn[u] != low[u]) return;
26
              tot[forest.new node()] = 0;
27
28
                   belong[stack[top]] = forest.n;
29
                   vis[stack[top]] = 2;
30
31
                   tot[forest.n]++;
                   --top;
32
              } while (stack[top + 1] != u);
33
34
         void solve() {
35
              forest.init(g -> base);
36
37
              int n = g \rightarrow n;
for (int i = 0; i < n; ++i)
38
                   if (!vis[i + g -> base]) {
                        top = dfs_clock = 0;
39
40
                        tarjan(i + g \rightarrow base, -1);
41
42
              for (int i = 0; i < g -> e / 2; ++i)
43
                   if (is_bridge[i]) {
44
                         int e = forest.e;
45
                         forest.bi_ins(belong[g \rightarrow v[i * 2]], belong[g \rightarrow v[i * 2 + 1]], g \rightarrow
       \hookrightarrow w[i * 2]);
46
                        ori[e] = make_pair(g -> v[i * 2 + 1], g -> v[i * 2]);
                        ori[e + 1] = make_pair(g \rightarrow v[i * 2], g \rightarrow v[i * 2 + 1]);
47
48
49
50
   } bcc;
```

Chapter 4 技巧

4.1 真正的释放 STL 容器内存空间

```
// vectors for example.
std::vector<int> v;
// Do something with v...
v.clear(); // Or having erased many.
std::vector<int>(v).swap(v);
```

4.2 无敌的大整数相乘取模

Time complexity O(1).

```
// 需要保证 x 和 y 非负
long long mult(long long x, long long y, long long MODN) {
long long t = (x * y - (long long)((long double)x / MODN * y + 1e-3) * MODN) %

→ MODN;
return t < 0 ? t + MODN : t;
}
```

4.3 无敌的读入优化

```
// getchar() 读入优化 << 关同步 cin << 此优化
// 用 isdigit() 会小幅变慢
namespace Reader {
    const int L = (1 << 15) + 5;
    char buffer[L], *S, *T;
```

CHAPTER 4. 技巧 5

```
__inline void get_int(int &x) {
    char ch; bool neg = 0;
    for (; get_char(ch), ch < '0' || ch > '9'; ) neg ^= ch == '-';
    x = ch - '0';
    for (; get_char(ch), ch >= '0' && ch <= '9'; )
        x = x * 10 + ch - '0';
    if (neg) x = -x;
}
</pre>
```