# Gungnir'l Standard Code Library\*

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 $<sup>{\</sup>rm *https://github.com/footoredo/Gungnir}$ 

## ${\bf Contents}$

1	计算	几何																						2
		二维									 		 	 		 								. 5
		1.1.1 基础									 		 	 		 								. :
		1.1.2 凸包									 		 	 		 								. :
		1.1.3 三角刑																						
		1.1.4 半平面																						
		1.1.5 圆交面																						
	1.2	三维																						
		1.2.1 基础																						
		1.2.2 凸包	• •		٠.		٠			٠	 ٠.	٠	 	 ٠.	 ٠	 	٠	٠.	•	٠	•	 ٠	 ٠	. ;
2	数论																							
_		$O(m^2 \log n)$	求线	性迫	維	数列	第	n I	J.		 		 	 		 								
	2.2																							
	2.3	中国剩余定理	!								 		 	 		 								. (
	$^{2.4}$	素性测试									 		 	 		 								. (
	2.5	质因数分解																						
	2.6	线下整点 .									 		 	 		 						 ٠		. (
9	代数	,																						
3	3.1		抽																					,
	3.2	自适应辛普森																						
		单纯形																						
	0.0						•			•	 	·	 	 	·	 	•	•	•	•		 ·	 ·	
4	字符																							7
		后缀自动机																						
		EX 后缀自动																						
	4.4	回文自动机					٠				 	٠	 	 	 ٠	 	٠		•	٠		 ٠	 ٠	. 8
5	数据	结构																						8
	~ nH																							,
	5.1	KD-Tree									 		 	 		 								. ?
		KD-Tree Treap																						
	5.2										 		 	 		 								. 9
	5.2 5.3	Treap Link/cut Tr									 		 	 		 								. 10
	5.2 5.3 图论	Treap Link/cut Tr	ee .								 		 	 		 								. 9 . 10
	5.2 5.3 图论 6.1	Treap Link/cut Tr 基础	ee .								 		 	 		 								. 9 . 10 . 10
	5.2 5.3 图论 6.1 6.2	Treap Link/cut Treap	ee .								 		 	 	 	 							 	. 9 . 10 . 10 . 10
	5.2 5.3 <b>图论</b> 6.1 6.2 6.3	Treap Link/cut Tr 基础 KM 点双连通分量	ree .								 		 	 	 	 						 	 	. (c) . 10 . 10 . 11 . 11 . 11
	5.2 5.3 <b>图论</b> 6.1 6.2 6.3 6.4	Treap Link/cut Tr 基础  KM  点双连通分量 边双连通分量	ree .								 		 	 	 	 						 	 	. 90 . 10 . 10 . 11 . 11
	5.2 5.3 <b>图论</b> 6.1 6.2 6.3 6.4 6.5	Treap Link/cut Tr 基础 KM 点双连通分量 遗小树形图	ree .								 		 	 	 	 						 	 	. 9 . 10 . 10 . 10 . 11 . 12 . 12
	5.2 5.3 <b>图论</b> 6.1 6.2 6.3 6.4 6.5	Treap Link/cut Tr 基础 KM 点双连通分量 遗小树形图	ree .								 		 	 		 						 	 	. 9 . 10 . 10 . 11 . 11 . 12 . 15
	5.2 5.3 <b>图论</b> 6.1 6.2 6.3 6.4 6.5 6.6	Treap Link/cut Tr 基础 KM 点双连通分量 最小树形图 带花树	ree .										 									 		. 9 . 10 . 10 . 11 . 12 . 12 . 13
6	5.2 5.3 <b>图论</b> 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8	Treap Link/cut Tr 基础 KM 点双连通分量 最小树形图 带花树 Dominator 无向图最小害	ree .										 									 		. 9 . 10 . 10 . 11 . 12 . 13 . 13 . 13
6	5.2 5.3 <b>图</b> 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 其他	Treap Link/cut Tr 基础 KM 点双连通分量 最小树形图 带花树 Dominator 无向图最小害	ree .								 					 						 		. 9 . 10 . 10 . 13 . 13 . 13 . 13 . 13
6	5.2 5.3 <b>图论</b> 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 其1	Treap Link/cut Tr 基础 KM 点双连通分量 最小树形图 带花树 Dominator 无向图最小割	ree .								 													. 9 . 10 . 10 . 11 . 12 . 13 . 13 . 13 . 13
6	5.2 5.3 <b>图</b> 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 其他	Treap Link/cut Tr 基础 KM 点双连通分量 最小树形图 带花树 Dominator 无向图最小害	ree .								 													. 9 . 10 . 10 . 11 . 12 . 13 . 13 . 13 . 13
6	5.2 5.3 <b>图</b> 6.1 6.2 6.3 6.4 6.5 6.6 6.7 7.2	Treap Link/cut Tr 基础 KM 点双连通分量 最小树形图 带花树 Dominator 无向图最小害	ree .								 													. (9 10 10 10 10 10 10 10 10 10 10 10 10 10
6	5.2 5.3 <b>图论</b> 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 其1	Treap Link/cut Tr 基础 KM 点双连通分量 贵小树 Dominator 无向图最小害 Dancing Lin 蔡勒公式	tree .																					. (9 10 10 10 10 10 10 10 10 10 10 10 10 10
6	5.2 5.3 <b>图论</b> 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 <b>其他</b> 7.1 7.2 <b>技巧</b>	Treap Link/cut Tr 基础 KM 点双连通分量 最小树形图 带花树 Dominator 无向图最小害	Tree	······································			· · · · · · · · · · · · · · · · · · ·																	. (9 10 10 10 10 10 10 10 10 10 10 10 10 10
6	5.2 5.3 <b>图论</b> 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 <b>其他</b> 7.1 7.2 <b>技巧</b>	Treap Link/cut Tr 基础 KM 点双连通分量 带花树 Dominator 无向图最小害 Dancing Lin 萘勒公式	ree	・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・		· · · · · · · · · · · · · · · · · · ·																		. 9 10 10 10 10 10 10 10 10 10 10 10 10 10
6	5.2 5.3 <b>图论</b> 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 <b>其他</b> 7.1 7.2 <b>技巧</b> 8.1 8.2	Treap Link/cut Tr 基础 KM 点双连通分量 带花树 Dominator 无向图最小事 Dancing Lin 蔡勒公式	·····································	・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・		·····································																		. 9 10 10 10 10 10 10 10 10 10 10 10 10 10
6 7 8	5.2 5.3 <b>图论</b> 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 <b>其</b> 1 7.2 <b>技</b> 3 8.1 8.2 8.3 8.4	Treap Link/cut Tr 基础 KM 点边最新 通通分分量 带边通形材 Dominator 无向向 Dancing Lin 数 的的的的 转 更无无格 称数 数 例 数 数 例 数 数 例 数 数 例 数 数 例 数 数 例 数 数 例 数 数 例 数 数 例 数 数 例 数 数 例 数 数 例 数 数 例 数 数 例 数 数 例 数 数 例 数 数 例 数	·····································	・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・		·····································																		. 9 10 10 10 10 10 10 10 10 10 10 10 10 10
6 7 8	5.2 5.3 <b>图论</b> 6.1 6.2 6.3 6.4 6.5 6.6 6.7 7.2 <b>技</b> 8.1 8.2 8.3 8.4 <b>提</b> 示	Treap Link/cut Tr 基础 KM 点页连列 基面通升 Dominator 无向图最小 Dancing Lin 蔡 真无政政 转数的的的 转数 真无无 概数 数人 第 第 第 第 第 第 第 第 第 第 第 第 第 第 第 第 第 第	·····································			· · · · · · · · · · · · · · · · · · ·																		. 9 10 10 10 10 10 10 10 10 10 10 10 10 10
6 7 8	5.2 5.3 <b>图论</b> 6.1 6.2 6.3 6.4 6.5 6.6 6.7 7.2 <b>技</b> 8.1 8.2 8.3 8.4 <b>提</b> 9.1	Treap Link/cut Tr 基础 KM 连通通形 动最增越形 Dominator 无向图显示式 的的的的旋转 控数分件 的的的的旋转	·····································			·····································																		. 9 10 10 10 10 10 10 10 10 10 10 10 10 10
6 7 8	5.2 5.3 <b>图</b> 6.1 6.2 6.3 6.4 6.5 6.6 6.7 <b>技</b> 8.1 8.2 8.3 8.4 <b>提</b> 9.1 9.2	Treap Link/cut Tr 基础 KM 点边最带面通通形 Dominator 无向图最中的的的成 下面的的的成 下面的的的成 下面的的的成 转数人 有数数人 有数数人 有数数人 有数数人 有数数人 有数数人 有数数人	·····································																					
6 7 8	5.2 5.3 <b>8</b> 6.1 6.2 6.3 6.4 6.5 6.6 6.7 7.2 <b>5</b> 7 <b>5</b> 8.1 <b>5</b> 8.2 <b>5</b> 8.3 <b>6</b> 8.3 <b>7</b> 7.2 <b>7</b> 7.2 <b>7</b> 7.2 <b>9</b> 9.2 <b>9</b>	Treap Link/cut Tr 基础 KM 点边最带还通通形型 Dancing Lin 要 真无无梅 控制 cout 和 控制 cout 和 控制 cout 和 控制 cout 和 证数 数 使数 使数 使数 使数 使数 是 数 数 使数 数 数 数 数 数 数	eee																					. 9 10 10 10 11 11 11 11 11 11 11 11 11 11
6 7 8	5.2 5.3 <b>8</b> 6.1 6.2 6.3 6.4 6.5 6.6 6.7 7. 技1 8.2 8.3 8.4 <b>2</b> 9.1 9.2 9.3 9.4	Treap Link/cut Tr  基M KM 点边最带及的向 连连树树 加速连树树 加速连树树 加速连树树 加速连树树 加速连树树 加速				· · · · · · · · · · · · · · · · · · ·																		. 9 10 10 10 11 11 11 11 11 11 11 11 11 11
6 7 8	5.2 5.3 <b>1</b> 66.2 6.3 6.4 6.6 6.6 6.7 <b>1</b> 7.2 <b>1</b> 7.2 <b>1</b> 7.2 <b>1</b> 7.2 <b>1</b> 9.2 9.3 9.3 9.5	Treap Link/cut Tr 基础 KM 点边最带还通通形型 Dancing Lin 要 真无无梅 控制 cout 和 控制 cout 和 控制 cout 和 控制 cout 和 证数 数 使数 使数 使数 使数 使数 是 数 数 使数 数 数 数 数 数 数				·····································																		

## Chapter 1 计算几何

## 1.1 二维

## 1.1.1 基础

```
int sign(DB x) {
       return (x > eps) - (x < -eps);
   DB msart(DB x) {
       return sign(x) > 0 ? sqrt(x) : 0;
   struct Point {
       DB x, y;
       Point rotate(DB ang) const { // 逆时针旋转 ang 弧度
10
           return Point(cos(ang) *x - sin(ang) *y,
11
12
                   cos(ang) * v + sin(ang) * x);
13
       Point turn90() const { // 逆时针旋转 90 度
14
15
           return Point(-y, x);
16
17
       Point unit() const {
18
           return *this / len();
19
20
21
   DB dot(const Point& a, const Point& b) {
22
       return a.x * b.x + a.y * b.y;
23
   DB det(const Point& a, const Point& b) {
25
       return a.x * b.y - a.y * b.x;
26
27 | #define cross(p1,p2,p3) ((p2.x-p1.x)*(p3.y-p1.y)-(p3.x-p1.x)*(p2.y-p1.y))
28 | #define crossOp(p1,p2,p3) sign(cross(p1,p2,p3))
   bool isLL(const Line& l1, const Line& l2, Point& p) { // 直线与直线交点
       DB s1 = det(l2.b - l2.a, l1.a - l2.a)
31
          s2 = -det(l2.b - l2.a, l1.b - l2.a);
       if (!sign(s1 + s2)) return false;
32
33
       p = (l1.a * s2 + l1.b * s1) / (s1 + s2);
34
       return true:
35 }
36 bool onSeg(const Line& l, const Point& p) { // 点在线段上
37
       return sign(det(p - l.a, l.b - l.a)) == 0 \&\& sign(dot(p - l.a, p - l.b)) <= 0;
38
39 Point projection(const Line & l, const Point& p) {
       return l.a + (l.b - l.a) * (dot(p - l.a, l.b - l.a) / (l.b - l.a).len2());
40
41
42 DB disToLine(const Line& l, const Point& p) { // 点到 * 直线 * 距离
43
       return fabs(det(p - l.a, l.b - l.a) / (l.b - l.a).len());
45 DB disToSeg(const Line& l, const Point& p) { // 点到线段距离
       return sign(dot(p - l.a, l.b - l.a)) * sign(dot(p - l.b, l.a - l.b)) == 1 ?
46
      \hookrightarrow disToLine(l, p) : std::min((p - l.a).len(), (p - l.b).len());
47 | }
48
   |// 圆与直线交点
   bool isCL(Circle a, Line l, Point& p1, Point& p2) {
   DB x = dot(l.a - a.o, l.b - l.a),
49
50
          y = (l.b - l.a).len2(),
51
52
          d = x * x - y * ((l.a - a.o).len2() - a.r * a.r);
       if (sign(d) < 0) return false;
53
       Point p = l.a - ((l.b - l.a) * (x / y)), delta = (l.b - l.a) * (msqrt(d) / y);
54
55
       p1 = p + delta; p2 = p - delta;
56
       return true;
57 }
58 / //圆与圆的交面积
59 DB areaCC(const Circle& c1, const Circle& c2) {
      DB d = (c1.o - c2.o).len();
if (sign(d - (c1.r + c2.r)) >= 0) return 0;
       if (sign(d - std::abs(c1.r - c2.r)) \le 0) {
           DB r = std::min(c1.r, c2.r);
```

```
64
            return r * r * PI:
 65
 66
        DB x = (d * d + c1.r * c1.r - c2.r * c2.r) / (2 * d),
            t1 = acos(x / c1.r), t2 = acos((d - x) / c2.r);
 67
 68
        return c1.r * c1.r * t1 + c2.r * c2.r * t2 - d * c1.r * sin(t1):
 69 }
 70 // 圆与圆交点
 71 bool isCC(Circle a, Circle b, P& p1, P& p2) {
        DB s1 = (a.o - b.o).len();
        if (sign(s1 - a.r - b.r) > 0 || sign(s1 - std::abs(a.r - b.r)) < 0) return false;
        DB s2 = (a.r * a.r - b.r * b.r) / s1;
 74
        DB aa = (s1 + s2) * 0.5, bb = (s1 - s2) * 0.5;
P o = (b.o - a.o) * (aa / (aa + bb)) + a.o;
 75
 76
        P \text{ delta} = (b.o - a.o).unit().turn90() * msgrt(a.r * a.r - aa * aa);
 77
 78
        p1 = o + delta, p2 = o - delta;
 79
        return true;
 80 }
 81 // 求点到圆的切点,按关于点的顺时针方向返回两个点
    bool tanCP(const Circle &c, const Point &p0, Point &p1, Point &p2) {
        double x = (p0 - c.o).len2(), d = x - c.r * c.r;
        if (d < eps) return false; // 点在圆上认为没有切点
        Point p = (p0 - c.o) * (c.r * c.r / x);
        Point delta = ((p0 - c.o) * (-c.r * sqrt(d) / x)).turn90();
 86
 87
        p1 = c.o + p + delta;
 88
        p2 = c.o + p - delta;
 89
        return true:
 90 }
    // 求圆到圆的外共切线,按关于 c1.o 的顺时针方向返回两条线
 91
 92
    vector<Line> extanCC(const Circle &c1, const Circle &c2) {
        vector<Line> ret;
 94
        if (sign(c1.r - c2.r) == 0) {
 95
            Point dir = c2.0 - c1.0;
 96
            dir = (dir * (c1.r / dir.len())).turn90();
 97
            ret.push_back(Line(c1.o + dir, c2.o + dir));
 98
            ret.push_back(Line(c1.o - dir, c2.o - dir));
99
        } else {
100
            Point p = (c1.0 * -c2.r + c2.0 * c1.r) / (c1.r - c2.r);
101
            Point p1, p2, q1, q2;
            if (tanCP(c1, p, p1, p2) && tanCP(c2, p, q1, q2)) {
    if (c1.r < c2.r) swap(p1, p2), swap(q1, q2);
102
103
104
                ret.push_back(Line(p1, q1));
105
                ret_push_back(Line(p2, q2));
106
107
108
        return ret;
109
110 // 求圆到圆的内共切线,按关于 c1.o 的顺时针方向返回两条线
    std::vector<Line> intanCC(const Circle &c1, const Circle &c2) {
        std::vector<Line> ret;
112
113
        Point p = (c1.0 * c2.r + c2.0 * c1.r) / (c1.r + c2.r);
114
        Point p1, p2, q1, q2;
        if (tanCP(c1, p, p1, p2) && tanCP(c2, p, q1, q2)) { // 两圆相切认为没有切线
115
116
            ret_push_back(Line(p1, q1));
117
            ret.push_back(Line(p2, q2));
118
119
        return ret:
120
    bool contain(vector<Point> polygon, Point p) { // 判断点 p 是否被多边形包含,包括落在
121
      → 边界上
        int ret = 0, n = polygon.size();
122
123
        for(int i = 0; i < n; ++ i) {
124
            Point u = polygon[i], v = polygon[(i + 1) % n];
            if (onSeg(Line(u, v), p)) return true; // Here I guess.
125
            if (sign(u,y - v,y) \le 0) swap(u, v);
126
127
            if (sign(p.y - u.y) > 0 \mid | sign(p.y - v.y) \le 0) continue;
            ret += sign(det(p, v, u)) > 0;
128
129
130
        return ret & 1;
131
132 // 用半平面 (q1,q2) 的逆时针方向去切凸多边形
```

CHAPTER 1. 计算几何 3

```
133 | std::vector<Point> convexCut(const std::vector<Point>&ps, Point q1, Point q2) {
134
         std::vector<Point> qs; int n = ps.size();
         for (int i = 0; i < n; ++i) {
   Point p1 = ps[i], p2 = ps[(i + 1) % n];
135
136
              int d1 = cross0p(q1,q2,p1), d2 = cross0p(q1,q2,p2);
137
138
              if (d1 \ge 0) qs.push_back(p1);
              if (d1 * d2 < 0) qs.push back(isSS(p1, p2, q1, q2));
139
140
         return qs;
141
142 }
143
    // 求凸包
144
    std::vector<Point> convexHull(std::vector<Point> ps) {
145
         int n = ps.size(); if (n <= 1) return ps;</pre>
         std::sort(ps.begin(), ps.end());
146
147
         std::vector<Point> as:
         for (int i = 0; i < n; qs.push_back(ps[i ++]))
    while (qs.size() > 1 && sign(det(qs[qs.size() - 2], qs.back(), ps[i])) <= 0)
148
149
150
                   qs.pop_back();
         for (int i = n - 2, t = qs.size(); i >= 0; qs.push_back(ps[i --])) while ((int)qs.size() > t && sign(det(qs[qs.size() - 2], qs.back(), ps[i])) <=
151
152
         qs.pop_back();
return qs;
```

#### 1.1.2 凸包

```
// 凸包中的点按逆时针方向
   struct Convex {
        std::vector<Point> a, upper, lower;
        void make_shell(const std::vector<Point>& p,
    std::vector<Point>& shell) { // p needs to be sorted.
             clear(shell); int n = p.size();
            shell push back(p[i]);
12
13
14
        void make_convex() {
            std::sort(a.begin(), a.end());
make_shell(a, lower);
std::reverse(a.begin(), a.end());
15
16
17
18
             make_shell(a, upper);
19
             a = Tower; a.pop_back();
            a.insert(a.end(), upper.begin(), upper.end());
if ((int)a.size() >= 2) a.pop_back();
20
21
22
23
            n = a.size():
24
25
26
        void init(const std::vector<Point>& _a) {
            clear(a); a = _a; n = a.size();
make_convex();
27
28
29
30
31
32
33
        void read(int n) { // Won't make convex.
            clear(a); n = _n; a.resize(n);
for (int i = 0; i < n; i++)
                 a[i].read();
        std::pair<DB, int> get_tangent(
34
35
36
37
                 const std::vecTor<Point>& convex, const Point& vec) {
             int l = 0, r = (int)convex.size() - 2;
             assert(r >= 0);
             for (; l + 1 < r; ) {
                 int mid = (l + r) / 2;
38
39
                 if (sign(det(convex[mid + 1] - convex[mid], vec)) > 0)
40
                 else l = mid;
41
42
             return std::max(std::make_pair(det(vec, convex[r]), r),
43
                      std::make_pair(det(vec, convex[0]), 0));
44
```

```
int binary_search(Point u, Point v, int l, int r) {
              int s1 = sign(det(v - u, a[l % n] - u));
47
48
              for (; l + 1 < r; ) {
                   int mid = (l + r) / 2;
49
50
                   int smid = sign(det(v - u, a[mid % n] - u));
                   if (smid == s1) l = mid:
51
52
                   else r = mid;
53
54
             return l % n;
55
        }
         // 求凸包上和向量 vec 叉积最大的点,返回编号,共线的多个切点返回任意一个
56
57
         int get_tangent(Point vec) {
             std::pair<DB, int> ret = get_tangent(upper, vec);
ret.second = (ret.second + (int)lower.size() - 1) % n;
58
59
60
              ret = std::max(ret, get_tangent(lower, vec));
61
              return ret.second:
62
63
         // 求凸包和直线 u, v 的交点,如果不相交返回 false,如果有则是和 (i, next(i)) 的
       → 交点, 交在点上不确定返回前后两条边其中之一
        bool get_intersection(Point u, Point v, int &i0, int &i1) {
  int p0 = get_tangent(u - v), p1 = get_tangent(v - u);
  if (sign(det(v - u, a[p0] - u)) * sign(det(v - u, a[p1] - u)) <= 0) {
    if (p0 > p1) std::swap(p0, p1);
    if (p0 > p1) std::swap(p0, p1);
65
66
67
                   i0 = binary_search(u, v, p0, p1);
i1 = binary_search(u, v, p1, p0 + n);
68
69
70
                   return true:
71
72
             else return false;
73
74 \ \ \ ;
```

#### 1.1.3 三角形的心

```
Point inCenter(const Point &A, const Point &B, const Point &C) { // 内心 double a = (B - C).len(), b = (C - A).len(), c = (A - B).len(), s = fabs(det(B - A, C - A)),
             r = s / p;
        return (A * a + B * b + C * c) / (a + b + c);
 6
    Point circumCenter(const Point &a, const Point &b, const Point &c) { // 外心
        Point bb = b - a, cc = c - a;
        double db = bb.len2(), dc = cc.len2(), d = 2 * det(bb, cc);
10
         return a - Point(bb.y * dc - cc.y * db, cc.x * db - bb.x * dc) / d;
11
12 Point othroCenter(const Point &a, const Point &b, const Point &c) { // 垂心 13 Point ba = b - a, ca = c - a, bc = b - c;
14
        double Y = ba.y * ca.y * bc.y,
15
                 A = ca.x * ba.y - ba.x * ca.y,
16
                 x0 = (Y + ca.x * ba.y * b.x - ba.x * ca.y * c.x) / A,
                 y0 = -ba.x * (x0 - c.x) / ba.y + ca.y;
17
18
        return Point(x0, y0);
19
```

#### 1.1.4 半平面交

```
| struct Point {
| int quad() const { return sign(y) == 1 || (sign(y) == 0 && sign(x) >= 0);}
| struct Line {
| bool include(const Point &p) const { return sign(det(b - a, p - a)) > 0; }
| Line push() const{ // 将半平面向外推 eps
| const double eps = 1e-6;
| Point delta = (b - a).turn90().norm() * eps;
| return Line(a - delta, b - delta);
| }
| bool sameDir(const Line &0, const Line &1) { return parallel(10, 11) && sign(dot(10.b - 10.a, 11.b - 11.a)) == 1; }
| bool operator < (const Point &a, const Point &b) {
```

CHAPTER 1. 计算几何 4

```
14
       if (a.quad() != b.quad()) {
15
           return a.quad() < b.quad();</pre>
16
       } else {
17
           return sign(det(a, b)) > 0;
18
19 }
20 | bool operator < (const Line &l0, const Line &l1) {
21
       if (sameDir(l0, l1)) {
22
           return l1.include(l0.a);
23
24
           return (l0.b - l0.a) < (l1.b - l1.a);
25
26 }
27 | bool check(const Line &u, const Line &v, const Line &w) { return
     vector<Point> intersection(vector<Line> &l) {
29
       sort(l.begin(), l.end());
30
       deque<Line> q;
31
       for (int i = 0; i < (int)l.size(); ++i) {
32
           if (i && sameDir(l[i], l[i - 1])) {
33
34
35
           while (q.size() > 1 \&\& !check(q[q.size() - 2], q[q.size() - 1], l[i]))
     → q.pop back();
36
           while (q.size() > 1 \& \{check(q[1], q[0], l[i])\} q.pop_front();
37
           g push back(l[i]);
38
39
       while (q.size() > 2 \&\& !check(q[q.size() - 2], q[q.size() - 1], q[0]))
       while (q.size() > 2 \& ! check(q[1], q[0], q[q.size() - 1])) q.pop_front();
40
41
       vector<Point> ret;
       for (int i = 0; i < (int)q.size(); ++i) ret.push_back(intersect(q[i], q[(i + 1) %
     \rightarrow a.size()]):
43
       return ret;
44 }
```

#### 1.1.5 圆交面积及重心

```
struct Event {
       Point p;
       double and:
       int delta;
       Event (Point p = Point(0, 0), double ang = 0, double delta = 0) : p(p), ang(ang),

    delta(delta) {}
 6 };
   bool operator < (const Event &a, const Event &b) {
       return a.ang < b.ang;
9 }
10
   void addEvent(const Circle &a, const Circle &b, vector<Event> &evt, int &cnt) {
       double d2 = (a.o - b.o).len2(),
11
               dRatio = ((a.r - b.r) * (a.r + b.r) / d2 + 1) / 3
12
               pRatio = sqrt(-(d2 - sqr(a.r - b.r)) * (d2 - sqr(a.r + b.r)) / (d2 * d2 *
13
      \rightarrow 4));
14
       Point d = b.o - a.o, p = d.rotate(PI / 2),
              q0 = a.o + d * dRatio + p * pRatio,
15
              q1 = a \cdot o + d * dRatio - p * pRatio;
16
       double ang 0 = (q0 - a.o).ang(),
17
              ang1 = (q1 - a.o).ang();
18
19
       evt.push_back(Event(q1, ang1, 1));
20
       evt.push_back(Event(q0, ang0, -1));
21
       cnt += ang1 > ang0;
22 | }
23 | bool issame(const Circle &a, const Circle &b) { return sign((a.o - b.o).len()) == 0 &&
      \rightarrow sign(a.r - b.r) == 0; }
24 | bool overlap(const Circle &a, const Circle &b) { return sign(a.r - b.r - (a.o -
      \rightarrow b.o).len()) >= 0: }
25 | bool intersect(const Circle &a, const Circle &b) { return sign((a.o - b.o).len() - a.r
      \rightarrow - b.r) < 0; }
26 Circle c[N];
27 double area[N];
                    // area[k] -> area of intersections >= k.
28 | Point centroid[N];
```

```
29 | bool keep[N];
   void add(int cnt, DB a, Point c) {
30
        area[cnt] += a;
        centroid[cnt] = centroid[cnt] + c * a;
33 | }
34 void solve(int C) {
35
        for (int i = 1; i <= C; ++ i) {
36
            area[i] = 0;
37
            centroid[i] = Point(0, 0);
38
39
        for (int i = 0; i < C; ++i) {
            int cnt = 1:
40
41
            vector<Event> evt;
            for (int j = 0; j < i; ++j) if (issame(c[i], c[j])) ++cnt; for (int j = 0; j < C; ++j) {
    if (j != i \&\& !issame(c[i], c[j]) \&\& overlap(c[j], c[i])) {
42
43
44
45
46
47
48
            for (int j = 0; j < C; ++j) {
49
                 if (j != i && !overlap(c[j], c[i]) && !overlap(c[i], c[j]) &&

    intersect(c[i], c[i])) →
50
                     addÉvent(c[i], c[j], evt, cnt);
51
52
53
            if (evt.size() == 0u) {
54
                 add(cnt, PI * c[i].r * c[i].r, c[i].o);
55
            } else {
56
                 sort(evt.begin(), evt.end());
57
                 evt.push back(evt.front());
58
                 for (int j = 0; j + 1 < (int)evt.size(); ++j) {
59
                      cnt += evt[j].delta;
60
                      add(cnt, det(evt[j].p, evt[j + 1].p) / 2, (evt[j].p + evt[j + 1].p) /
      → 3);
61
                      double ang = evt[j + 1].ang - evt[j].ang;
62
                      if (ang < 0) {
63
                          ang += PI * 2;
64
65
                      if (sign(ang) == 0) continue;
                      add(cnt, ang * c[i].r * c[i].r / 2, c[i].o +
66
                          Point(\sin(\arg 1) - \sin(\arg 0), -\cos(\arg 1) + \cos(\arg 0)) * (2 / (3 *
67
      \rightarrow ang) * c[i].r)):
                      add(cnt, -sin(ang) * c[i].r * c[i].r / 2, (c[i].o + evt[j].p + evt[j + c[i].n / 2])
      \hookrightarrow 1].p) / 3);
69
                 }
70
71
72
        for (int i = 1; i <= C; ++ i)
73
            if (sign(area[i])) {
74
                 centroid[i] = centroid[i] / area[i];
75
76 }
```

## 1.2 三维

#### 1.2.1 基础

```
// 三维绕轴旋转,大拇指指向 axis 向量方向,四指弯曲方向转 w 弧度
   Point rotate(const Point& s, const Point& axis, DB w) {
      DB x = axis.x, y = axis.y, z = axis.z;
      DB s1 = x * x + y * y + z * z, ss1 = msqrt(s1),
         cosw = cos(w), sinw = sin(w);
      DB a[4][4];
      memset(a, 0, sizeof a);
      a[3][3] = 1;
      a[0][0] = ((y * y + z * z) * cosw + x * x) / s1;
10
      a[0][1] = x * y * (1 - cosw) / s1 + z * sinw / ss1;
11
      a[0][2] = x * z * (1 - cosw) / s1 - y * sinw / ss1;
12
      a[1][0] = x * y * (1 - cosw) / s1 - z * sinw / ss1;
      a[1][1] = ((x * x + z * z) * cosw + y * y) / s1;
```

```
14
        a[1][2] = y * z * (1 - cosw) / s1 + x * sinw / ss1;
        a[2][0] = x * z * (1 - \cos w) / s1 + y * \sin w / ss1;
15
        a[2][1] = y * z * (1 - cosw) / s1 - x * sinw / ss1;
16
        a[2][2] = ((x * x + y * y) * cos(w) + z * z) / s1;
17
        DB ans[4] = \{0, 0, 0, 0\}, c[4] = \{s.x, s.y, s.z, 1\};
18
19
        for (int i = 0; i < 4; ++ i)
        for (int j = 0; j < 4; ++ j)
ans[i] += a[j][i] * c[j];
return Point(ans[0], ans[1], ans[2]);
20
21
22 23 }
```

## 1.2.2 凸包

```
inline P cross(const P& a, const P& b) {
        return P(
                  a.y * b.z - a.z * b.y
                  a.z * b.x - a.x * b.z
                  a.x * b.y - a.y * b.x
 7 }
     __inline DB mix(const P& a, const P& b, const P& c) {
        return dot(cross(a, b), c);
11
12
     _inline DB volume(const P& a, const P& b, const P& c, const P& d) {
14
        return mix(b - a, c - a, d - a);
15 }
17
    struct Face {
18
        int a, b, c;
19
         __inline Face() {}
20
        __inline Face(int _a, int _b, int _c):
21
             a(a), b(b), c(c)
22
        __inline DB area() const {
23
             return 0.5 * cross(p[b] - p[a], p[c] - p[a]).len();
24
25
        __inline P normal() const {
26
             return cross(p[b] - p[a], p[c] - p[a]).unit();
27
28
        __inline DB dis(const P& p0) const {
29
             return dot(normal(), p0 - p[a]);
30
31 };
32
33 st
   std::vector<Face> face, tmp; // Should be O(n).
34
35
    int mark[N][N], Time, n;
36
     inline void add(int v) {
37
38
        ++ Time:
39
        for (int i = 0; i < (int) face size(); ++ i) {
             int a = face[i].a, b = face[i].b, c = face[i].c;
40
             if (sign(volume(p[v], p[a], p[b], p[c])) > 0) {
   mark[a][b] = mark[b][a] = mark[a][c] =
41
42
                       mark[c][a] = mark[b][c] = mark[c][b] = Time;
43
44
45
             else {
46
                  tmp.push_back(face[i]);
47
48
49
        clear(face); face = tmp;
        for (int i = 0; i < (int)tmp.size(); ++ i) {
  int a = face[i].a, b = face[i].b, c = face[i].c;
  if (mark[a][b] == Time) face.emplace_back(v, b, a);</pre>
50
51
52
             if (mark[b][c] == Time) face.emplace_back(v, c, b); if (mark[c][a] == Time) face.emplace_back(v, a, c);
53
54
55
56
57
58
             assert(face.size() < 500u);
59 void reorder() {
```

```
for (int i = 2; i < n; ++ i) {
  P tmp = cross(p[i] - p[0], p[i] - p[1]);
  if (sign(tmp.len())) {</pre>
61
62
63
                    std::swap(p[i], p[2]);
                    for (int j = 3; j < n; ++ j)
if (sign(volume(p[0], p[1], p[2], p[j]))) {
64
66
                               std::swap(p[j], p[3]);
67
68
69
70
         }
71
72
73
    void build_convex() {
74
75
         reorder();
         clear(face);
76
         face.emplace_back(0, 1, 2);
         face emplace_back(0, 2, 1);
77
         for (int i = 3; i < n; ++ i)
79
              add(i):
80 }
```

## Chapter 2 数论

## $2.1 O(m^2 \log n)$ 求线性递推数列第 n 项

Given  $a_0, a_1, \dots, a_{m-1}$   $a_n = c_0 \times a_{n-m} + \dots + c_{m-1} \times a_{n-1}$ Solve for  $a_n = v_0 \times a_0 + v_1 \times a_1 + \cdots + v_{m-1} \times a_{m-1}$ 

```
void linear_recurrence(long long n, int m, int a[], int c[], int p) {
2 3 4
         long long v[M] = \{1 \% p\}, u[M \ll 1], msk = !!n;
         for(long long i(n); i > 1; i >>= 1) {
              msk <<= 1:
5
6
         for(long long x(0); msk; msk >>= 1, x <<= 1) {
              fill_n(u, m << 1, 0);
              int \overline{b}(!!(n \& msk));
              x \mid = b;
10
              if(x < m) {
11
                   u[x] = 1 % p;
12
              }else {
13
                   for(int i(0); i < m; i++) {
                        for(int j(0), t(i + b); j < m; j++, t++) {
    u[t] = (u[t] + v[i] * v[j]) % p;
14
15
16
17
18
                   for(int i((m << 1) - 1); i >= m; i--) {
                        for(int j(0), t(i - m); j < m; j++, t++) {
    u[t] = (u[t] + c[j] * u[i]) % p;
19
20
21
22
                  }
23
24
25
26
27
              copy(u, u + m, v);
         //a[n] = v[0] * a[0] + v[1] * a[1] + ... + v[m-1] * a[m-1].
         for(int i(m); i < 2 * m; i++) {
28
              a[i] = 0;
              for(int j(0); j < m; j++) {
    a[i] = (a[i] + (long long)c[j] * a[i + j - m]) % p;
29
30
31
32
33
34
         for(int j(0); j < m; j++) {
    b[j] = 0;
             for(int i(0); i < m; i++) {
    b[j] = (b[j] + v[i] * a[i + j]) % p;
35
36
37
38
39
        for(int j(0); j < m; j++) {
    a[j] = b[j];</pre>
40
```

```
42 }
  2.2 求逆元
   void ex_gcd(long long a, long long b, long long &x, long long &y) {
       if (b == 0) {
            x = 1;
            v = 0:
            return;
       long long xx, yy;
ex_gcd(b, a % b, xx, yy);
       y = xx - a / b * yy;
10
11 | }
12
13
   long long inv(long long x, long long MODN) {
        long long inv_x, y;
14
15
       ex_gcd(x, MODN, inv_x, y);
16
       return (inv_x % MODN + MODN) % MODN;
17 | }
```

## 2.3 中国剩余定理

## 2.4 素性测试

```
int strong_pseudo_primetest(long long n,int base) {
        long long n2=n-1, res;
        while(n2\%2==0) n2>>=1,s++;
        res=powmod(base,n2,n);
        if((res==1)||(res==n-1)) return 1;
        while(s \ge 0) {
            res=mulmod(res,res,n);
            if(res==n-1) return 1;
11
12
13
        return 0; // n is not a strong pseudo prime
14
15
   int isprime(long long n) {
       static LL testNum[]={2,3,5,7,11,13,17,19,23,29,31,37};
static LL lim[]={4,0,1373653LL,25326001LL,25000000000LL,2152302898747LL,
16
17
      → 3474749660383LL,341550071728321LL,0,0,0,0);
       if(n<2||n==3215031751LL) return 0;
18
        for(int i=0;i<12;++i){
19
            if(n<lim[i]) return 1:
21
            if(strong pseudo primetest(n,testNum[i])==0) return 0;
22
23
        return 1;
24 }
```

## 2.5 质因数分解

```
int ansn; LL ans[1000];
LL func(LL x,LL n){ return(mod_mul(x,x,n)+1)%n; }
```

```
3 | LL Pollard(LL n){
       LL i,x,y,p;
       if(Rabin_Miller(n)) return n;
       if(!(n&1)) return 2;
       for(i=1;i<20;i++){
           x=i; y=func(x,n); p=gcd(y-x,n);
9
           while(p==1) {x=func(x,n); y=func(func(y,n),n); p=gcd((y-x+n)%n,n)%n;}
10
           if(p==0||p==n) continue;
11
12
13 }
14 void factor(LL n){
15 LL x;
16
       x=Pollard(n);
17
       if(x==n){ ans[ansn++]=x; return; }
18
       factor(x), factor(n/x);
19 }
```

#### 2.6 线下整点

```
 \begin{array}{l} 1 \\ // \sum_{i=0}^{n-1} \left \lfloor \frac{a+bi}{m} \right \rfloor, \ n,m,a,b>0 \\ \text{LL solve(LL n,LL a,LL b,LL m)} \\ \text{if}(b==\emptyset) \ \text{return n*}(a/m); \\ \text{if}(a>=m) \ \text{return n*}(a/m) + \text{solve}(n,a\%m,b,m); \\ \text{if}(b>=m) \ \text{return } (n-1)*n/2*(b/m) + \text{solve}(n,a,b\%m,m); \\ \text{return solve}((a+b*n)/m,(a+b*n)\%m,m,b); \\ \end{array}
```

## Chapter 3 代数

## 3.1 快速傅里叶变换

```
// n 必须是 2 的次幂
2
   void fft(Complex a[], int n, int f) {
         for (int i = 0; i < n; ++i)
             if (R[i] < i) swap(a[i], a[R[i]]);</pre>
         for (int i = 1, h = 0; i < n; i <<= 1, h++) {
             Complex wn = Complex(cos(pi / i), f * sin(pi / i));
6
             Complex w = Complex(1, 0);
7
             for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w; for (int p = i << 1, j = 0; j < n; j += p) {
10
                  for (int k = 0; k < i; ++k) {
                       Complex x = a[j + k], y = a[j + k + i] * tmp[k];

a[j + k] = x + y; a[j + k + i] = x - y;
11
12
13
14
             }
15
        }
16 }
```

## 3.2 自适应辛普森积分

```
namespace adaptive_simpson {
        template<typename function>
        inline double area(function f, const double &left, const double &right) {
            double mid = (left + right) / 2;
            return (right - left) * (f(left) + 4 * f(mid) + f(right)) / 6;
8
        template<typename function>
        inline double simpson(function f, const double &left, const double &right, const

    double &eps, const double &area_sum) {
        double mid = (left + right) / 2;
}
10
            double area_left = area(f, left, mid);
double area_right = area(f, mid, right);
11
12
13
            double area_total = area_left + area_right;
14
            if (fabs(area total - area sum) <= 15 * eps) {
15
                 return area_total + (area_total - area_sum) / 15;
```

```
16
17
18
18
19
20
21
22
22
23
24
}
return simpson(f, left, right, eps / 2, area_left) + simpson(f, mid, right,
eps / 2, area_right);
}
template<typename function>
inline double simpson(function f, const double &left, const double &right, const

odouble &eps) {
return simpson(f, left, right, eps, area(f, left, right));
}
```

#### 3.3 单纯形

```
const double eps = 1e-8;
   // max{c * x | Ax <= b, x >= 0} 的解, 无解返回空的 vector, 否则就是解.
3 vector<double> simplex(vector<vector<double> > &A, vector<double> b, vector<double> c)
       int n = A.size(), m = A[0].size() + 1, r = n, s = m - 1;
       vector<vector<double> > D(n + 2, vector<double>(m + 1));
       vector < int > ix(n + m):
       for(int i = 0; i < n + m; i++) {
8
           ix[i] = i;
9
10
       for(int i = 0; i < n; i++) {
           for(int j = 0; j < m - 1; j++) {
    D[i][j] = -A[i][j];
11
12
13
           D[i][m-1]=1;
14
15
           D[i][m] = b[i];
            if (D[r][m] > D[i][m]) {
16
17
                r = i;
18
19
20
21
       for(int j = 0; j < m - 1; j++) {
22
           D[n][j] = c[j];
23
24
25
26
       D[n + 1][m - 1] = -1;
       for(double d; ;) {
            if (r < n) {
27
                swap(ix[s], ix[r + m]);
D[r][s] = 1. / D[r][s];
28
29
                for(int j = 0; j <= m; j++) {
30
                    if (i != s) {
31
                        D[r][j] *= -D[r][s];
32
33
34
                for(int i = 0; i \le n + 1; i++) {
35
                    if (i != r) {
36
                         for(int j = 0; j <= m; j++) {
37
                             if (j != s) {
38
                                 D[i][j] += D[r][j] * D[i][s];
39
40
41
                        D[i][s] *= D[r][s];
42
               }
43
44
45
            r = -1, s = -1;
46
            for(int j = 0; j < m; j++) {
47
                if (s < 0 | | ix[s] > ix[j]) {
48
                    if (D[n + 1][j] > eps || D[n + 1][j] > -eps && D[n][j] > eps) {
49
                        s = j;
50
51
                }
52
53
            if (s < 0) {
54
                break;
55
            for(int i = 0; i < n; i++) {
```

```
if (D[i][s] < -eps) {</pre>
                    if (r < 0) \mid (d = D[r][m] / D[r][s] - D[i][m] / D[i][s]) < -eps
58
59
                         | | d < eps & ix[r + m] > ix[i + m])  {
60
61
                         r = i:
62
                    }
63
64
65
66
           if (r < 0) {
67
                return vector<double> ();
68
69
70
       if (D[n + 1][m] < -eps) {
71
            return vector<double> ();
72
73
74
       vector<double> x(m - 1);
75
       for(int i = m; i < n + m; i++) {
76
           if (ix[i] < m - 1) {
                x[ix[i]] = D[i - m][m];
77
78
79
80
       return x;
81 }
```

## Chapter 4 字符串

## 4.1 后缀数组

```
const int MAXN = MAXL * 2 + 1;
   int a[MAXN], x[MAXN], y[MAXN], c[MAXN], sa[MAXN], rank[MAXN], height[MAXN];
   void calc_sa(int n) {
        int m = alphabet, k = 1;
        memset(c, 0, sizeof(*c) * (m + 1));
        for (int i = 1; i \le n; ++i) c[x[i] = a[i]]++;
        for (int i = 1; i \le m; ++i) c[i] += c[i - 1];
8
        for (int i = n; i; --i) sa[c[x[i]]--] = i;
        for (; k <= n; k <<= 1) {
10
            int tot = k;
            for (int i = n - k + 1; i \le n; ++i) y[i - n + k] = i;
11
12
            for (int i = 1; i \le n; ++i)
13
                 if (sa[i] > k) y[++tot] = sa[i] - k;
14
            memset(c, 0, sizeof(*c) * (m + 1));
15
            for (int i = 1; i \le n; ++i) c[x[i]]++;
            for (int i = 1; i <= m; ++i) c[i] += c[i - 1];
for (int i = n; i; --i) sa[c[x[y[i]]]--] = y[i];
for (int i = 1; i <= n; ++i) y[i] = x[i];</pre>
16
17
18
            tot = 1; x[sa[1]] = 1:
19
20
            for (int i = 2; i \le n; ++i) {
                 if (max(sa[i], sa[i - 1]) + k > n || y[sa[i]] != y[sa[i - 1]] || y[sa[i] +
21
      \rightarrow k] != y[sa[i - 1] + k]) ++tot;
                 x[sa[i]] = tot;
23
24
            if (tot == n) break; else m = tot;
25
26
27
   void calc_height(int n) {
        for (int i = 1; i \le n; ++i) rank[sa[i]] = i;
28
        for (int i = 1; i <= n; ++i) {
29
30
            height[rank[i]] = max(0, height[rank[i-1]] - 1);
31
            if (rank[i] == 1) continue;
32
            int j = sa[rank[i] - 1];
33
            while (\max(i, j) + \text{height}[\text{rank}[i]] \le n \& a[i + \text{height}[\text{rank}[i]]] == a[j + n]

    height[rank[i]]]) ++height[rank[i]];
34
       }
35
```

#### 4.2 后缀自动机

```
static const int MAXL = MAXN * 2; // MAXN is original length
static const int alphabet = 26; // sometimes need changing
 3 int l, last, cnt, trans[MAXL][alphabet], par[MAXL], sum[MAXL], seq[MAXL], mxl[MAXL],

    size[MAXL]; // mxl is maxlength, size is the size of right

    char str[MAXL];
    inline void init() {
         l = strlen(str + 1); cnt = last = 1;
         for (int i = 0; i \le l * 2; ++i) memset(trans[i], 0, sizeof(trans[i]));
         memset(par, 0, sizeof(*par) * (l * 2 + 1));
        memset(mxl, 0, sizeof(*mxl) * (l * 2 + 1));
memset(size, 0, sizeof(*size) * (l * 2 + 1));
9
10
11 | }
   inline void extend(int pos, int c) {
   int p = last, np = last = ++cnt;
   mxl[np] = mxl[p] + 1; size[np] = 1;
12
13
14
         for (; p && !trans[p][c]; p = par[p]) trans[p][c] = np;
15
         if (!p) par[np] = 1;
16
17
18
              int q = trans[p][c];
19
              if (mxl[p] + 1 == mxl[q]) par[np] = q;
20
              else {
21
                    int ng = ++cnt;
22
23
                   mxl[nq] = mxl[p] + 1;
                    memcpy(trans[nq], trans[q], sizeof(trans[nq]));
24
25
26
                    par[nq] = par[q];
                    par[np] = par[q] = nq;
for (; trans[p][c] == q; p = par[p]) trans[p][c] = nq;
27
28
         }
29
30
    inline void buildsam() {
         for (int i = 1; i <= l; ++i) extend(i, str[i] - 'a');
31
         memset(sum, 0, sizeof(*sum) * (l * 2 + 1));
32
33
         for (int i = 1; i <= cnt; ++i) sum[mxl[i]]++;</pre>
         for (int i = 1; i <= l; ++i) sum[i] += sum[i - 1];
for (int i = cnt; i; --i) seq[sum[mxl[i]]--] = i;
for (int i = cnt; i; --i) size[par[seq[i]]] += size[seq[i]];
34
35
36
```

#### 4.3 EX 后缀自动机

```
inline void add node(int x, int &last) {
        int lastnode = last;
       if (c[lastnode][x]) {
            int nownode = c[lastnode][x];
            if (l[nownode] == l[lastnode] + 1) last = nownode:
            else {
                 int auxnode = ++cnt; l[auxnode] = l[lastnode] + 1;
                for (int i = 0; i < alphabet; ++i) c[auxnode][i] = c[nownode][i];</pre>
                 par[auxnode] = par[nownode]; par[nownode] = auxnode;
                 for (; lastnode && c[lastnode][x] == nownode; lastnode = par[lastnode]) {
                     c[lastnode][x] = auxnode;
12
13
                 last = auxnode;
14
15
       } else {
16
            int newnode = ++cnt; l[newnode] = l[lastnode] + 1;
            for (: lastnode && !c[lastnode][x]: lastnode = par[lastnode]) c[lastnode][x] =

→ newnode;

            if (!lastnode) par[newnode] = 1;
18
19
            else {
20
                int nownode = c[lastnode][x];
                if (l[lastnode] + 1 == l[nownode]) par[newnode] = nownode;
21
22
23
                     int auxnode = ++cnt; l[auxnode] = l[lastnode] + 1;
24
                     for (int i = 0; i < alphabet; ++i) c[auxnode][i] = c[nownode][i];</pre>
25
                     par[auxnode] = par[nownode]; par[nownode] = par[newnode] = auxnode;
for (; lastnode && c[lastnode][x] == nownode; lastnode =
      → par[lastnodel) {
```

## 4.4 回文自动机

```
int nT, nStr, last, c[MAXT][26], fail[MAXT], r[MAXN], l[MAXN], s[MAXN]; int allocate(int len) {
       l[nT] = len;
       r[nT] = 0;
5
       fail[nT] = 0;
       memset(c[nT], 0, sizeof(c[nT]));
       return nT++:
8
9
   void init() {
10
       nT = nStr = 0;
11
       int newE = allocate(0);
12
       int new0 = allocate(-1);
13
       last = newE:
14
       fail[newE] = new0;
       fail[new0] = newE;
15
16
       s[0] = -1;
17 }
18 | void add(int x) {
19
       s[++nStr] = x;
       int now = last;
20
       while (s[nStr - l[now] - 1] != s[nStr]) now = fail[now]; if (!c[now][x]) {
21
22
23
            int newnode = allocate(l[now] + 2), &newfail = fail[newnode];
24
           newfail = fail[now];
25
           while (s[nStr - l[newfail] - 1] != s[nStr]) newfail = fail[newfail]:
26
           newfail = c[newfail][x]:
27
           c[now][x] = newnode:
28
29
       last = c[now][x]:
30
       r[last]++:
31
32
   void count() {
33
       for (int i = nT - 1; i \ge 0; i--) {
           r[fail[i]] += r[i];
34
35
36
```

## Chapter 5 数据结构

#### 5.1 KD-Tree

```
long long norm(const long long &x) {
            For manhattan distance
       return std::abs(x);
           For euclid distance
       return x * x;
6
8
   struct Point {
       int x, y, id;
10
       const int& operator [] (int index) const {
11
12
           if (index == 0) {
13
               return x;
14
           } else {
15
               return y;
16
       }
17
18
19
       friend long long dist(const Point &a, const Point &b) {
20
           long long result = 0:
```

CHAPTER 5. 数据结构

```
9
```

```
for (int i = 0; i < 2; ++i) {
22
                result += norm(a[i] - b[i]);
23
24
25
26
27
28
            return result;
   } point[N];
   struct Rectangle {
29
30
        int min[2], max[2];
31
        Rectangle() {
32
33
            min[0] = min[1] = INT_MAX; // sometimes int is not enough
            \max[0] = \max[1] = INT MIN;
34
35
36
        void add(const Point &p) {
37
            for (int i = 0; i < 2; ++i) {
38
                min[i] = std::min(min[i], p[i]);
39
                max[i] = std::max(max[i], p[i]);
40
        }
41
42
43
        long long dist(const Point &p) {
44
            long long result = 0;
45
            for (int i = 0; i < 2; ++i) {
46
                // For minimum distance
47
48
                result += norm(std::min(std::max(p[i], min[i]), max[i]) - p[i]);
                // For maximum distance
49
                result += std::max(norm(max[i] - p[i]), norm(min[i] - p[i]));
50
51
            return result:
52
53
54
55
   };
   struct Node {
56
57
        Point seperator;
        Rectangle rectangle;
58
59
        int child[2];
60
        void reset(const Point &p) {
61
            seperator = p;
62
            rectangle = Rectangle();
63
            rectangle.add(p);
64
            child[0] = child[1] = 0;
   } tree[N << 1];</pre>
68
   int size, pivot;
   bool compare(const Point &a, const Point &b) {
   if (a[pivot] != b[pivot]) {
70
71
72
            return a[pivot] < b[pivot]:</pre>
73
74
75
76
        return a.id < b.id;</pre>
   // 左閉右開: build(1, n + 1)
78
   int build(int l, int r, int type = 1) {
79
        pivot = type;
80
        if (l >= r) {
81
            return 0;
82
83
84
        int mid = l + r \gg 1;
        std::nth_element(point + l, point + mid, point + r, compare);
85
86
        tree[x]. reset(point[mid]);
87
        for (int i = l; i < r; ++i) {
88
            tree[x].rectangle.add(point[i]);
89
90
        tree[x].child[0] = build(l, mid, type ^ 1);
91
92
93
        tree[x].child[1] = build(mid + 1, r, type ^ 1);
        return x;
```

```
95
    int insert(int x, const Point &p, int type = 1) {
 96
        pivot = type;
 97
        if (x == 0)
 98
            tree[++size].reset(p);
 99
             return size;
100
101
        tree[x].rectangle.add(p);
        if (compare(p, tree[x] seperator)) {
102
103
             tree[x].child[0] = insert(tree[x].child[0], p, type ^ 1);
104
105
            tree[x].child[1] = insert(tree[x].child[1], p, type ^ 1);
106
107
        return x:
108
109
110 // For minimum distance
    // For maximum: 下面递归 query 时 0, 1 换顺序;< and >;min and max
    void query(int x, const Point &p, std::pair<long long, int> &answer, int type = 1) {
        if (x == 0 | | tree[x].rectangle.dist(p) > answer.first) {
114
115
             return;
116
117
        answer = std::min(answer,
118
                  std::make_pair(dist(tree[x].seperator, p), tree[x].seperator.id));
         if (compare(p, tree[x].seperator)) {
119
            query(tree[x].child[0], p, answer, type ^ 1);
120
121
             query(tree[x].child[1], p, answer, type ^ 1);
122
        } else {
123
             query(tree[x].child[1], p, answer, type ^ 1);
124
             query(tree[x].child[0], p, answer, type ^ 1);
125
126
128
129
    std::priority_queue<std::pair<long long, int> > answer;
130
    void query(int x, const Point &p, int k, int type = 1) {
131
         if (x == 0 \mid | (int)answer.size() == k && tree[x].rectangle.dist(p) >
132
       → answer.top().first) {
133
134
135
        answer.push(std::make_pair(dist(tree[x].seperator, p), tree[x].seperator.id));
        if ((int)answer size(\overline{)} > k) {
136
137
            answer.pop();
138
        if (compare(p, tree[x].seperator)) {
139
            query(tree[x] child[0], p, k, type ^ 1);
140
141
             query(tree[x].child[1], p, k, type ^ 1);
142
        } else {
143
            query(tree[x].child[1], p, k, type ^ 1);
query(tree[x].child[0], p, k, type ^ 1);
144
145
146 }
```

#### 5.2 Treap

```
struct Node{
         int mn, key, size, tag;
         bool rev;
        Node* ch[2];
        Node(int mn, int key, int size): mn(mn), key(key), size(size), rev(0), tag(0){}
         void downtag():
         Node* update(){
             mn = min(ch[0] \rightarrow mn, min(key, ch[1] \rightarrow mn));

size = ch[0] \rightarrow size + 1 + ch[1] \rightarrow size;
9
10
             return this:
11
12
13 typedef pair<Node*, Node*> Pair;
14 Node *null, *root;
15 | void Node::downtag(){
```

```
16
        if(rev){
17
             for(int i = 0; i < 2; i++)
18
                  if(ch[i] != null){
                      ch[i] -> rev ^= 1;
19
20
                      swap(ch[i] \rightarrow ch[0], ch[i] \rightarrow ch[1]);
21
22
23
             rev = 0:
24
        if(tag){
25
             for(int i = 0; i < 2; i++)
26
                 if(ch[i] != null){
                      ch[i] -> key += tag;
27
28
                      ch[i] -> mn += tag;
29
                      ch[i] \rightarrow tag += tag;
30
31
32
             tag = 0;
33
34
   int r(){
35
        static int s = 3023192386;
36
        return (s += (s << 3) + 1) & (\sim0u >> 1);
37
38
   bool random(int x, int y){
39
        return r() % (x + y) < x;
40
41
   Node* merge(Node *p, Node *q){
42
        if(p == null) return q;
43
        if(q == null) return p;
44
        p -> downtag();
45
        q -> downtag();
46
        if(random(p -> size, q -> size)){
             p \rightarrow ch[1] = merge(p \rightarrow ch[1], q);
47
48
             return p -> update();
49
50
             q \rightarrow ch[0] = merge(p, q \rightarrow ch[0]);
51
             return q -> update();
52
53
54 | Pair split(Node *x, int n){
55
        if(x == null) return make_pair(null, null);
56
        x -> downtag();
        if(n <= x -> ch[0] -> size){
    Pair ret = split(x -> ch[0], n);
57
58
             x \rightarrow ch[0] = ret.second;
59
             return make_pair(ret.first, x -> update());
60
61
62
        Pair ret = split(x \rightarrow ch[1], n - x \rightarrow ch[0] \rightarrow size - 1);
63
        x \rightarrow ch[1] = ret.first;
64
        return make_pair(x -> update(), ret.second);
65
66
   pair<Node*, Pair> get_segment(int l, int r){
67
        Pair ret = split(root, l - 1);
68
        return make_pair(ret.first, split(ret.second, r - l + 1));
69
70
        null = new Node(INF, INF, 0);
null -> ch[0] = null -> ch[1] = null;
71
72
73
        root = null:
74 | }
```

## 5.3 Link/cut Tree

```
inline void reverse(int x) {
    tr[x].rev ^= 1; swap(tr[x].c[0], tr[x].c[1]);
}
inline void rotate(int x, int k) {
    int y = tr[x].fa, z = tr[y].fa;
    tr[x].fa = z; tr[z].c[tr[z].c[1] == y] = x;
    tr[tr[x].c[k ^ 1]].fa = y; tr[y].c[k] = tr[x].c[k ^ 1];
    tr[x].c[k ^ 1] = y; tr[y].fa = x;
```

```
10 | }
11
12
   inline void splay(int x, int w) {
13
        int z = x; pushdown(x);
        while (tr[x] fa != w) {
14
15
            int y = tr[x].fa; z = tr[y].fa;
16
            if (z == w)
                pushdown(z = y); pushdown(x);
rotate(x, tr[y].c[1] == x);
update(y); update(x);
17
18
19
20
            } else {
                pushdown(z); pushdown(y); pushdown(x);
21
22
                 int t1 = tr[y] c[1] == x, t2 = tr[z] c[1] == y;
23
                if (t1 == t2) rotate(y, t2), rotate(x, t1);
24
                else rotate(x, t1), rotate(x, t2);
25
                update(z); update(x);
26
27
28
       update(x);
29
        if (x != z) par[x] = par[z], par[z] = 0;
30 }
31
32
   inline void access(int x) {
33
        for (int y = 0; x; y = x, x = par[x]) {
            splay(x, 0);
            if (tr[x].c[1]) par[tr[x].c[1]] = x, tr[tr[x].c[1]].fa = 0;
            tr[x].c[1] = y; par[y] = 0; tr[y].fa = x; update(x);
37
        }
38
39
40
   inline void makeroot(int x) {
41
42
        access(x); splay(x, 0); reverse(x);
43
   inline void link(int x, int y) {
45
        makeroot(x); par[x] = y;
46
47
48
   inline void cut(int x, int y) {
       access(x); splay(y, 0);
if (par[y] != x) swap(x, y), access(x), splay(y, 0);
49
51
        par[y] = 0;
52
53
54
   inline void split(int x, int y) { // x will be the root of the tree
55
        makeroot(y); access(x); splay(x, 0);
```

## Chapter 6 图论

## 6.1 基础

```
struct Graph { // Remember to call .init()!
       int e, nxt[M], v[M], adj[N], n;
       bool base;
       __inline void init(bool _base, int _n = 0) {
           assert(n < N);</pre>
           n = n; base = _base;
e = 0; memset(adj + base, -1, sizeof(*adj) * n);
       __inline int new_node() {
10
           adj[n + base] = -1;
11
           assert(n + base + \dot{1} < N);
12
13
           return n++ + base;
14
       __inline void ins(int u0, int v0) { // directional
15
           assert(u0 < n + base && v0 < n + base);
16
           v[e] = v0; nxt[e] = adj[u0]; adj[u0] = e++;
17
           assert(e < M);
18
19
       __inline void bi_ins(int u0, int v0) { // bi-directional
```

CHAPTER 6. 图论 11

```
20 | ins(u0, v0); ins(v0, u0); 21 | };
```

#### 6.2 KM

```
struct KM {
       // Trulv 0(n^3)
        // 邻接矩阵,不能连的边设为 -INF, 求最小权匹配时边权取负, 但不能连的还是 -INF,
      → 使用时先对 1 -> n 调用 hungary() , 再 get_ans() 求值
       int w[N][N];
       int lx[N], ly[N], match[N], way[N], slack[N];
bool used[N];
       void init() {
            for (int i = 1; i <= n; i++) {
                match[i] = 0;
10
                 lx[i] = 0;
11
                 lv[i] = 0;
12
                wav[i] = 0;
13
14
15
        void hungary(int x) {
            match[0] = x;
16
17
            int j0 = 0;
18
            for (int j = 0; j <= n; j++) {
19
                 slack[j] = INF;
20
                used[j] = false;
21
22
23
            do {
24
                 used[i0] = true:
25
                 int i0 = match[j0], delta = INF, j1 = 0;
                for (int j = 1; j <= n; j++) {
   if (used[j] == false) {</pre>
26
27
                         int cur = -w[i0][j] - lx[i0] - ly[j];
if (cur < slack[j]) {
28
29
30
                              slack[j] = cur;
31
                              way[j] = j0;
32
33
                         if (slack[j] < delta) {</pre>
34
                              delta = slack[j];
35
                              j1 = j;
36
37
                     }
38
                for (int j = 0; j <= n; j++) {
   if (used[j]) {</pre>
39
40
                         lx[match[j]] += delta;
41
42
                         lv[i] -= delta;
43
44
                     else slack[j] -= delta;
45
46
                 i0 = j1;
47
48
            } while (match[j0] != 0);
49
                 int j1 = way[j0];
50
51
                 match[j0] = match[j1];
52
                 i0 = i1;
53
            } while (j0);
54
55
56
57
       int get_ans() {
            int sum = 0;
58
            for(int i = 1; i <= n; i++) {
59
                 if (w[match[i]][i] == -INF); // 无解
60
                 if (match[i] > 0) sum += w[match[i]][i];
61
62
            return sum;
       }
63
```

64 | } km;

#### 6.3 点双连通分量

bcc.forest is a set of connected tree whose vertices are chequered with cut-vertex and BCC.

```
const bool BCC_VERTEX = 0, BCC_EDGE = 1;
   struct BCC {    // N = N0 + M0. Remember to call init(&raw_graph).
    Graph *g, forest; // g is raw graph ptr.
    int dfn[N], DFN, low[N];
        int stack[N], top;
        int expand to[N];
                                     // Where edge i is expanded to in expaned graph.
        // Vertex \bar{i} expaned to i.
        int compress_to[N]; // Where vertex i is compressed to.
bool vertex_type[N], cut[N], compress_cut[N], branch[M];
//std::vector<int> BCC_component[N]; // Cut vertex belongs to none.
__inline void init(Graph *raw_graph) {
8
9
10
11
12
             g = raw_graph;
13
14
15
        void DFS(int u, int pe) {
   dfn[u] = low[u] = ++DFN; cut[u] = false;
16
             if (!~g->adj[u]) {
17
                  cut[u] = 1;
18
                  compress_to[u] = forest.new_node();
19
                  compress_cut[compress_to[u]] = 1;
20
21
             for (int e = g\rightarrow adj[u]; \sim e; e = g\rightarrow nxt[e]) {
                   int v = g->v[e];
22
23
                  if ((e^pe) > 1 && dfn[v] > 0 && dfn[v] < dfn[u]) {
24
                       stack[top++] = e;
25
                       low[u] = std::min(low[u], dfn[v]);
26
27
                  else if (!dfn[v]) {
28
                       stack[top++] = e; branch[e] = 1;
                       DFS(v, e);
29
30
                       low[u] = std::min(low[v], low[u]);
                       if (low[v] >= dfn[u]) {
31
32
                            if (!cut[u]) {
33
                                 cut[u] = 1;
                                 compress_to[u] = forest_new_node();
34
35
                                 compress_cut[compress_to[u]] = 1;
36
37
                            int cc = forest.new_node();
38
                            forest.bi ins(compress to[u], cc);
39
                            compress cut[cc] = 0;
40
                            //BCC_component[cc].clear();
41
                            do {
42
                                 int cur_e = stack[--top];
43
                                 compress_to[expand_to[cur_e]] = cc;
44
                                 compress_to[expand_to[cur_e^1]] = cc;
45
                                 if (branch[cur e]) {
46
                                      int v = g - v[cur_e];
                                      if (cut[v])
47
48
                                           forest.bi_ins(cc, compress_to[v]);
49
                                      else {
50
                                           //BCC_component[cc].push_back(v);
51
                                           compress to[v] = cc:
52
53
54
                            } while (stack[top] != e);
55
56
                  }
57
             }
58
59
        void solve() {
             forest.init(q->base):
60
61
             int n = g -> n;
             for (int i = 0; i < g->e; i++)
62
63
                  expand_to[i] = g->new_node();
64
65
             memset(branch, 0, sizeof(*branch) * g->e);
```

CHAPTER 6. 图论 12

```
66
            memset(dfn + g->base, 0, sizeof(*dfn) * n); DFN = 0;
67
            for (int i = 0; i < n; i++)
68
                if (!dfn[i + g->base]) {
69
70
                    top = 0;
                    DFS(i + g -> base, -1);
71
72
73
74
   } bcc;
75
   bcc.init(&raw_graph);
76 | bcc.solve();
  // Do something with bcc.forest ...
```

## 6.4 边双连通分量

```
struct BCC {
         Graph *q, forest;
         int dfn[N], low[N], stack[N], tot[N], belong[N], vis[N], top, dfs\_clock; // tot[] is the size of each BCC, belong[] is the BCC that each node belongs to
         pair<int, int > ori[M]; // bridge in raw_graph(raw node)
         bool is_bridge[M];
         __inline void init(Graph *raw_graph) {
              g = raw_graph;
              memset(is_bridge, false, sizeof(*is_bridge) * g -> e);
10
              memset(vis + g \rightarrow base, 0, sizeof(*vis) * g \rightarrow n);
11
12
         void tarjan(int u, int from)
13
              dfn[u] = low[u] = ++dfs\_clock; vis[u] = 1; stack[++top] = u;
14
              for (int p = g -> adj[u]; ~p; p = g -> nxt[p]) {
   if ((p ^ 1) == from) continue;
15
                   int v = g \rightarrow v[p];
16
17
                   if (vis[v]) {
18
                        if (vis[v] == 1) low[u] = min(low[u], dfn[v]);
19
                   } else {
20
                        tarjan(v, p);
                        low[u] = min(low[u], low[v]);
if (low[v] > dfn[u]) is_bridge[p / 2] = true;
21
22
23
24
25
              if (dfn[u] != low[u]) return;
26
27
28
              tot[forest.new node()] = 0:
                   belong[stack[top]] = forest.n;
vis[stack[top]] = 2;
29
30
31
                   tot[forest.n]++;
                   --top;
              } while (stack[top + 1] != u);
32
33
34
35
         void solve() {
              forest.init(g -> base);
              int n = g \rightarrow n;
36
37
              for (int i = 0; i < n; ++i)
38
                   if (!vis[i + g -> base]) {
39
                        top = dfs_clock = 0;
40
                        tarjan(i + g \rightarrow base, -1);
41
42
              for (int i = 0; i < g -> e / 2; ++i)
    if (is_bridge[i]) {
43
44
                        int e = forest.e;
45
                        forest.bi_ins(belong[g \rightarrow v[i * 2]], belong[g \rightarrow v[i * 2 + 1]], g \rightarrow
       \hookrightarrow w[i * 2]);
46
                        ori[e] = make_pair(g -> v[i * 2 + 1], g -> v[i * 2]);
                        ori[e + 1] = make_pair(g \rightarrow v[i * 2], g \rightarrow v[i * 2 + 1]);
47
48
49
   } bcc;
```

## 6.5 最小树形图

```
const int MAXN,INF;// INF >= sum( W_ij )
int from[MAXN + 10] [MAXN * 2 + 10],n,m,edge[MAXN + 10] [MAXN * 2 + 10];
```

```
3 | int sel[MAXN * 2 + 10], fa[MAXN * 2 + 10], vis[MAXN * 2 + 10];
   int getfa(int x){if(x == fa[x]) return x; return fa[x] = getfa(fa[x]);}
   void liuzhu(){ // 1-base: root is 1, answer = (sel[i], i) for i in [2..n]
        for(int i = 2; i \le n; ++i){
            sel[i] = 1: fa[i] = i:
            for(int j = 1; j <= n; ++j) if(fa[j] != i)
9
                 if(from[j][i] = i, edge[sel[i]][i] > edge[j][i]) sel[i] = j;
10
11
12
        int limit = n;
13
        while(1){
            int prelimit = limit; memset(vis, 0, sizeof(vis)); vis[1] = 1;
for(int i = 2; i <= prelimit; ++i) if(fa[i] == i && !vis[i]){</pre>
14
15
                 int j = i; while(!vis[j]) vis[j] = i, j = getfa(sel[j]);
if(j == 1 || vis[j] != i) continue; vector<int> C; int k = j;
16
17
                 do C.push_back(k), k = getfa(sel[k]); while(k != j);
18
19
20
                 for(int i = 1; i <= n; ++i){
21
                      edge[i][limit] = INF, from[i][limit] = limit;
22
23
                 fa[limit] = vis[limit] = limit;
24
                 for(int i = 0; i < int(C.size()); ++i){
25
                      int x = C[i], fa[x] = limit;
26
                      for(int j = 1; j <= n; ++j)
27
                           if(edge[j][x] != INF && edge[j][limit] > edge[j][x] -
      \hookrightarrow edge[sel[x]][x]){
28
29
                               edge[j][limit] = edge[j][x] - edge[sel[x]][x];
                               from[j][limit] = x;
30
31
32
                 for(int j=1;j<=n;++j) if(getfa(j)==limit) edge[j][limit] = INF;</pre>
33
                 sel[limit] = 1;
34
                 for(int j = 1; j <= n; ++j)
                      if(edge[sel[limit]][limit] > edge[j][limit]) sel[limit] = j;
35
36
37
38
            if(prelimit == limit) break;
39
        for(int i = limit; i > 1; --i) sel[from[sel[i]][i]] = sel[i];
40
```

#### 6.6 带花树

```
vector<int> link[maxn];
   int n,match[maxn],Queue[maxn],head,tail;
   int pred[maxn], base[maxn], start, finish, newbase;
   bool InQueue[maxn], InBlossom[maxn];
   void push(int u){ Queue[tail++]=u;InQueue[u]=true; }
   int pop(){ return Queue[head++]; }
   int FindCommonAncestor(int u,int v){
       bool InPath[maxn];
       for(int i=0;i<n;i++) InPath[i]=0;</pre>
       while(true){ u=base[u];InPath[u]=true;if(u==start) break;u=pred[match[u]]; }
10
       while(true){ v=base[v];if(InPath[v]) break;v=pred[match[v]]; }
11
12
       return v;
13
14
   void ResetTrace(int u){
15
16
       while(base[u]!=newbase){
17
           v=match[u];
18
           InBlossom[base[u]]=InBlossom[base[v]]=true;
19
           u=pred[v];
20
           if(base[u]!=newbase) pred[u]=v;
21
22
23
   void BlossomContract(int u,int v){
24
25
       newbase=FindCommonAncestor(u,v);
       for (int i=0;i<n;i++)</pre>
26
       InBlossom[i]=0;
27
       ResetTrace(u); ResetTrace(v)
28
       if(base[u]!=newbase) pred[u]=v;
```

CHAPTER 6. 图论 13

```
29
         if(base[v]!=newbase) pred[v]=u;
30
         for(int i=0;i<n;++i</pre>
31
         if(InBlossom[base[i]]){
32
              base[i]=newbase;
33
              if(!InOueue[i]) push(i):
34
35
36
    bool FindAugmentingPath(int u){
         bool found=false;
37
38
        for(int i=0;i<n;++i) pred[i]=-1,base[i]=i;
for (int i=0;i<n;i++) InQueue[i]=0;</pre>
39
40
         start=u;finish=-1; head=tail=0; push(start);
41
         while(head<tail){</pre>
             int u=pop();
for(int i=link[u].size()-1;i>=0;i--){
    int v=link[u][i];
    int v=link[u][i];
42
43
44
                   if(base[u]!=base[v]&&match[u]!=v)
   if(v==start||(match[v]>=0&&pred[match[v]]>=0))
        BlossomContract(u,v);
45
46
47
                        else if(pred[v]==-1){
48
                             pred[v]=u:
49
50
                              if(match[v]>=0) push(match[v]);
51
                              else{ finish=v; return true; }
52
53
54
55
56
         return found;
57
58
    void AugmentPath(){
         int u=finish,v,w;
59
         while(u>=0){ v=pred[u];w=match[v];match[v]=u;match[u]=v;u=w; }
60
61
    void FindMaxMatching(){
         for(int i=0:i<n:++i) match[i]=-1:
62
         for(int i=0;i<n;++i) if(match[i]==-1) if(FindAugmentingPath(i)) AugmentPath();</pre>
63
```

#### 6.7 Dominator Tree

```
vector<int> prec[N], succ[N];
 vector<int> ord;
 3 int stamp, vis[N];
   int num[N);
   int fa[N];
   void dfs(int u) {
       vis[u] = stamp;
num[u] = ord.size();
        ord.push_back(u);
10
        for (int^i = 0; i < (int)succ[u].size(); ++i) {
11
            int v = succ[u][i];
12
            if (vis[v] != stamp) {
13
                 fa[v] = u;
14
                dfs(v);
15
16
17
18
   int fs[N], mins[N], dom[N], sem[N];
int find(int u) {
19
20
        if (u != fs[u])
21
            int v = fs[u];
22
            fs[u] = find(fs[u]);
23
            if (\min[v] != -1 \&\& num[sem[mins[v]]] < num[sem[mins[u]]]) {
24
                mins[u] = mins[v];
25
26
27
        return fs[u];
28
29 void merge(int u, int v) { fs[u] = v; }
30 vector<int> buf[N];
31 int buf2[N];
32 void mark(int source) {
```

```
ord.clear():
34
       ++stamp;
35
       dfs(source):
36
        for (int i = 0: i < (int) \text{ ord.size}(): ++i) {
37
            int u = ord[i];
38
            fs[u] = u, mins[u] = -1, buf2[u] = -1;
39
40
        for (int i = (int) ord.size() - 1; i > 0; --i) {
            int u = ord[i], p = fa[u];
41
            sem[u] = p;
42
43
            for (int j = 0; j < (int)prec[u].size(); ++j) {</pre>
                 int v = prec[u][j];
44
45
                 if (use[v] != stamp) continue;
46
                if (num[v] > num[u]) {
                     find(v); v = sem[mins[v]];
47
48
49
                if (num[v] < num[sem[u]]) {</pre>
50
                     sem[u] = v;
51
52
53
            buf[sem[u]].push_back(u);
54
            mins[u] = u;
55
            merge(u, p);
            while (buf[p] size()) {
56
57
58
                int v = buf[p].back();
buf[p].pop_back();
                find(v);
59
                    (sem[v] == sem[mins[v]]) {
60
                     dom[v] = sem[v];
61
62
                } else {
63
                     buf2[v] = mins[v];
64
65
            }
66
67
       dom[ord[0]] = ord[0];
       for (int i = 0; i < (int)ord.size(); ++i) {</pre>
68
            int u = ord[i];
69
            if (~buf2[u]) {
70
                dom[u] = dom[buf2[u]];
71
72
73
       }
74
```

## 无向图最小割

```
int cost[maxn] [maxn], seq[maxn], len[maxn], n, m, pop, ans;
    bool used[maxn];
 3
    void Init(){
         int i, j,a,b,c;
         for(i=0;i<n;i++) for(j=0;j<n;j++) cost[i][j]=0;
        for(i=0;i<m;i++){
             scanf("%d %d %d",&a,&b,&c); cost[a][b]+=c; cost[b][a]+=c;
 9
        pop=n; for(i=0;i<n;i++) seq[i]=i;
10 }
11
12
   void Work(){
        ans=inf; int i,j,k,l,mm,sum,pk;
while(pop > 1){
13
14
             for(i=1;i<pop;i++) used[seq[i]]=0; used[seq[0]]=1;</pre>
15
             for(i=1;i<pop;i++) len[seq[i]]=cost[seq[0]][seq[i]];
pk=0; mm=-inf; k=-1;</pre>
16
17
              for(i=1;i<pop;i++) if(len[seq[i]] > mm){ mm=len[seq[i]]; k=i; }
18
             for(i=1;i<pop;i++){
    used[seq[l=k]]=1;</pre>
19
20
                  if(i==pop-2) pk=k;
                  if(i==pop-1) break;
21
22
                  mm=-inf;
                  for(j=1;j<pop;j++) if(!used[seq[j]])
    if((len[seq[j]]+=cost[seq[l]][seq[j]]) > mm)
23
24
25
                            mm=len[seq[j]], k=j;
```

## Chapter 7 其他

## 7.1 Dancing Links

```
1 | struct Node {
        Node *1, *r, *u, *d, *col;
         int size, line_no;
        Node() {
              size = 0; line_no = -1;
              l = r = \dot{u} = d = col = \dot{N}ULL:
   } *root;
   void cover(Node *c) {
10
11
        c -> l -> r = c -> r; c -> r -> l = c -> l;
         for (Node *u = c->d: u != c: u = u->d)
12
13
             for (Node *v = u->r; v != u; v = v->r) {
14
                  v\rightarrow d\rightarrow u = v\rightarrow u;
15
                  v->u->d = v->d;
16
                  -- v->col->size;
17
18
19
    void uncover(Node *c) {
21
         for (Node *u = c->u; u != c; u = u->u) {
22
23
24
25
26
27
28
              for (Node *v = u -> 1; v != u; v = v -> 1) {
                  ++ v->col->size;
                  v->u->d = v:
                  v\rightarrow d\rightarrow u = v;
        c -> l -> r = c; c -> r -> l = c;
29
30
31
   std::vector<int> answer:
   bool search(int k) {
33
         if (root->r == root) return true;
34
35
        Node *r = NULL;
        for (Node *u = root->r; u != root; u = u->r)
36
             if (r == NULL || u->size < r->size)
37
38
        if (r == NULL || r->size == 0) return false;
39
        else {
40
              cover(r);
41
             bool succ = false;
             for (Node *u = r->d; u != r && !succ; u = u->d) {
   answer.push_back(u->line_no);
42
43
44
                  for (Node *\overline{v} = u \rightarrow r; v != u; v = v \rightarrow r) // Cover row
45
                       cover(v->col);
                  succ |= search(k + 1);
46
                  for (Node *v = u -> 1; v != u; v = v -> 1)
47
48
                       uncover(v->col);
49
                  if (!succ) answer pop back();
50
51
52
             uncover(r);
             return succ;
53
54
55
56 bool entry[CR][CC];
57 | Node *who[CR][CC];
58 int cr, cc;
```

```
void construct() {
 60
         root = new Node();
 61
         Node *last = root;
 62
 63
         for (int i = 0; i < cc; ++ i) {
             Node *u = new Node();
 64
 65
              last->r = u; u->l = last;
 66
             Node *v = u; u \rightarrow line_no = i;
 67
              last = u;
             for (int j = 0; j < cr; ++ j)
    if (entry[j][i]) {</pre>
 68
 69
 70
                       ++ u->size;
 71
                      Node *cur = new Node();
 72
73
                       who[j][i] = cur;
                       cur->line_no = j;
 74
                       cur->col = u:
 75
76
                       cur->u = v; v->d = cur;
                       v = cur:
 77
 78
             v->d = u; u->u = v;
 79
 80
         last->r = root; root->l = last;
 81
         for (int j = 0; j < cr; ++ j) {
             Node *last = NULL;
 82
 83
             for (int i = cc - 1; i >= 0; -- i)
   if (entry[j][i]) {
 84
 85
                       last = who[j][i];
                       break:
 87
 88
              for (int i = 0; i < cc; ++ i)
                  if (entry[j][i]) {
 89
                       last->r = who[j][i];
who[j][i]->l = last;
 90
 91
 92
                       last = who[j][i];
 93
 94
         }
 95
 96
 97
     void destruct() {
         for (Node *u = root->r; u != root; ) {
              for (Node *v = u->d; v != u; ) {
100
                  Node *nxt = v->d:
101
                  delete(v);
102
                  v = nxt;
103
104
             Node *nxt = u->r;
105
             delete(u); u = nxt;
106
107
         delete root;
108
```

## 7.2 蔡勒公式

```
int zeller(int y, int m, int d) {
    if (m<=2) y--, m+=12; int c=y/100; y%=100;
    int w=((c>>2)-(c<<1)+y+(y>>2)+(13*(m+1)/5)+d-1)%7;
    if (w<0) w+=7; return(w);
}</pre>
```

## Chapter 8 技巧

## 8.1 真正的释放 STL 容器内存空间

```
template <typename T>
__inline void clear(T& container) {
    container.clear(); // 或者删除了一堆元素
    T(container).swap(container);
}
```

## 8.2 无敌的大整数相乘取模

Time complexity O(1).

```
1 // 需要保证 x 和 y 非负 long long mult(long long x, long long y, long long MODN) { long long t = (x * y - (long long)((long double)x / MODN * y + 1e-3) * MODN) % → MODN; return t < 0 ? t + MODN : t; }
```

## 8.3 无敌的读入优化

```
1 // getchar() 读入优化 << 关同步 cin << 此优化
  2 // 用 isdigit() 会小幅变慢
3 // 返回 false 表示读到文件尾
4 namespace Reader {
          const int L = (1 << 15) + 5;
          char buffer[L], *S, *T;
          __inline bool getchar(char &ch) {
               if (S == T) {
                    T = (S = buffer) + fread(buffer, 1, L, stdin);
                    if (S == T) {
                         ch = EOF;
11
12
                         return false:
13
14
               ch = *S++;
15
16
               return true;
17
         __inline bool getint(int &x) {
    char ch; bool neg = 0;
    for (; getchar(ch) && (ch < '0' || ch > '9'); ) neg ^= ch == '-';
18
19
20
21
22
23
24
25
26
27
28
               if (ch == EOF) return false;
               x = ch - '0';
               for (; getchar(ch), ch >= '0' && ch <= '9'; )
    x = x * 10 + ch - '0';
               if (neg) x = -x;
               return true:
```

#### 8.4 梅森旋转算法

High quality pseudorandom number generator, twice as efficient as rand() with -02. C++11 required.

```
#include <random>
int main() {
    std::mt19937 g(seed); // std::mt19937_64
    std::cout << g() << std::endl;
}</pre>
```

## Chapter 9 提示

## 9.1 控制 cout 输出实数精度

```
std::cout << std::fixed << std::setprecision(5);</pre>
```

#### 9.2 vimrc

## 9.3 让 make 支持 c ++ 11

In .bashrc or whatever:

export CXXFLAGS='-std=c++11 -Wall'

## 9.4 tuple 相关

```
mytuple = std::make_tuple (10, 2.6, 'a');
std::tie (myint, std::ignore, mychar) = mytuple;
std::get<I>(mytuple) = 20;
std::cout << std::get<I>(mytuple) << std::endl;
// get the Ith(const) element</pre>
```

## 9.5 线性规划转对偶

```
\begin{array}{l} \text{maximize } \mathbf{c}^T \mathbf{x} \\ \text{subject to } \mathbf{A} \mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq 0 \end{array} \Longleftrightarrow \begin{array}{l} \text{minimize } \mathbf{y}^T \mathbf{b} \\ \text{subject to } \mathbf{y}^T \mathbf{A} \geq \mathbf{c}^T, \mathbf{y} \geq 0 \end{array}
```

## 9.6 32-bit/64-bit 随机素数

32-bit	64-bit
73550053	1249292846855685773
148898719	1701750434419805569
189560747	3605499878424114901
459874703	5648316673387803781
1202316001	6125342570814357977
1431183547	6215155308775851301
1438011109	6294606778040623451
1538762023	6347330550446020547
1557944263	7429632924303725207
1981315913	8524720079480389849
+ .	V 4 - 10

## 9.7 NTT 素数及其原根

Prime	Primitive root
1053818881	7
1051721729	6
1045430273	3
1012924417	5
1007681537	3