Gungnir's Standard Code Library

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Chapter 1 计算几何

1.1 二维

1.1.1 基础

```
typedef double DB;
   const DB eps = 1e-8;
   int sign(DB x) {
       return x < -eps ? -1 : (x > eps ? 1 : 0);
       return sign(x) > 0 ? sgrt(x) : 0;
10
11
   struct Point {
       DB x, y;
Point(): x(0), y(0) {}
12
13
       Point(DB x, DB y): x(x), y(y) {}
14
15
       Point operator+(const Point &rhs) const {
16
            return Point(x + rhs.x, y + rhs.y);
17
18
       Point operator-(const Point &rhs) const {
            return Point(x - rhs.x, y - rhs.y);
19
20
21
       Point operator*(DB k) const {
22
            return Point(x * k, y * k);
23
24
25
26
       Point operator/(DB k) const {
           assert(sign(k)):
            return Point(x / k, y / k);
27
28
       Point rotate(DB ang) const { // 逆时针旋转 ang 弧度
29
            return Point(cos(ang) *x - sin(ang) *y,
30
                    cos(ang) * y + sin(ang) * x);
31
       }
32
       Point turn90() const { // 逆时针旋转 90 度
33
            return Point(-y, x);
34
35
       Point unit() const {
36
           return *this / len();
37
38
39
   DB dot(const Point& a, const Point& b) {
       return a.x * b.x + a.y * b.y;
41
42
   DB det(const Point& a, const Point& b) {
43
       return a.x * b.y - a.y * b.x;
44
45 | #define cross(p1,p2,p3) ((p2.x-p1.x)*(p3.y-p1.y)-(p3.x-p1.x)*(p2.y-p1.y))
   #define crossOp(p1,p2,p3) sign(cross(p1,p2,p3))
   bool isLL(const Line& l1, const Line& l2, Point& p) { // 直线与直线交点 DB s1 = det(l2.b - l2.a, l1.a - l2.a), s2 = -det(l2.b - l2.a, l1.b - l2.a);
48
49
50
       if (!sign(s1 + s2)) return false;
51
52
53
       p = (l1.a * s2 + l1.b * s1) / (s1 + s2);
       return true:
54
   bool onSeg(const Line& l, const Point& p) { // 点在线段上
55
       return sign(det(p - l.a, l.b - l.a)) == 0 && sign(dot(p - l.a, p - l.b)) <= 0;
56
57
   Point projection(const Line & l, const Point& p) {
       return l.a + (l.b - l.a) * (dot(p - l.a, l.b - l.a) / (l.b - l.a).len2());
58
59
   DB disToLine(const Line& l, const Point& p) { // 点到 * 直线 * 距离
       return fabs(det(p - l.a, l.b - l.a) / (l.b - l.a).len());
61
62 }
63 | DB disToSeg(const Line& l, const Point& p) { // 点到线段距离
       return sign(dot(p - l.a, l.b - l.a)) * sign(dot(p - l.b, l.a - l.b)) == 1 ?
      \rightarrow disToLine(l, p) : std::min((p - l.a).len(), (p - l.b).len());
```

```
65 | }
    // 圆与直线交点
    bool isCL(Circle a, Line l, Point& p1, Point& p2) {
    DB x = dot(l.a - a.o, l.b - l.a),
        y = (l.b - l.a).len2(),
         d = x * x - y * ((l.a - a.o).len2() - a.r * a.r);
if (sign(d) < 0) return false;</pre>
 70
 71
 72
73
74
         Point p = l.a - ((l.b - l.a) * (x / y)), delta = (l.b - l.a) * (msqrt(d) / y);
         p1 = p + delta; p2 = p - delta;
 75 }
 76 //圆与圆的交面积
 77 DB areaCC(const Circle& c1, const Circle& c2) {
         DB d = (c1.o - c2.o).len();
if (sign(d - (c1.r + c2.r)) >= 0) return 0;
if (sign(d - std::abs(c1.r - c2.r)) <= 0) {
 80
 81
             DB r = std::min(c1.r, c2.r);
              return r * r * PI:
 82
 83
 84
         DB x = (d * d + c1.r * c1.r - c2.r * c2.r) / (2 * d),
 85
              t1 = acos(x / c1.r), t2 = acos((d - x) / c2.r);
 86
         return c1.r * c1.r * t1 + c2.r * c2.r * t2 - d * c1.r * sin(t1);
 87
    // 圆与圆交点
 88
    | DB s1 = (a.o - b.o).len();
 89
         if (sign(s1 - a.r - b.r) > 0 \mid | sign(s1 - std::abs(a.r - b.r)) < 0) return false;
         DB s2 = (a.r * a.r - b.r * b.r) / s1;
         DB aa = (s1 + s2) * 0.5, bb = (s1 - s2) * 0.5;
P o = (b.o - a.o) * (aa / (aa + bb)) + a.o;
 94
         P delta = (b.o - a.o).unit().turn90() * msqrt(a.r * a.r - aa * aa);
 95
 96
97
         p1 = o + delta, p2 = o - delta;
         return true:
 98 }
    // 求点到圆的切点,按关于点的顺时针方向返回两个点
 99
    bool tanCP(const Circle &c, const Point &p0, Point &p1, Point &p2) {
100
         double x = (p0 - c.o) \cdot len2(), d = x - c.r * c.r;
         if (d < eps) return false; // 点在圆上认为没有切点
         Point p = (p0 - c.o) * (c.r * c.r / x);
Point delta = ((p0 - c.o) * (-c.r * sqrt(d) / x)).turn90();
103
104
105
         p1 = c.o + p + delta;
         p2 = c.o + p - delta;
106
107
         return true;
108 }
    // 求圆到圆的内共切线,按关于 cl.o 的顺时针方向返回两条线
std::vector<Line> intanCC(const Circle &c1, const Circle &c2) {
111
         std::vector<Line> ret;
         Point p = (c1.0 * c2.r + c2.0 * c1.r) / (c1.r + c2.r);
112
113
         Point p1, p2, q1, q2;
114
         if (tanCP(c1, p, p1, p2) && tanCP(c2, p, q1, q2)) { // 两圆相切认为没有切线
115
              ret.push_back(Line(p1, q1));
116
              ret_push_back(Line(p2, q2));
117
118
         return ret:
119
    // 点在多边形内
120
     bool inPolygon(const Point& p, const std::vector<Point>& poly) {
         int n = polv.size():
123
         int counter = 0;
124
         for (int i = 0; i < n; ++ i) {
125
             P = poly[i], b = poly[(i + 1) % n];
126
             if (onSeg(Line(a, b), p)) return false; // 边界上不算
127
             int x = sign(det(p - a, b - a));
             int y = sign(a.y - p.y);
128
              int z = sign(b.y - p.y);
129
             if (x < 0 & x < 0 & x < 0 & x > 0) ++ counter;
if (x < 0 & x < 0 & x < 0 & x > 0) -- counter;
130
131
132
133
         return counter != 0;
134
```

```
135 // 用半平面 (q1,q2) 的逆时针方向去切凸多边形
    std::vector<Point> convexCut(const std::vector<Point>&ps, Point q1, Point q2) {
136
137
        std::vector<Point> qs; int n = ps.size();
        for (int i = 0; i < n; ++i) {
   Point p1 = ps[i], p2 = ps[(i + 1) % n];
   int d1 = crossOp(q1,q2,p1), d2 = crossOp(q1,q2,p2);
138
139
140
141
             if (d1 \ge 0) qs.push_back(p1);
142
             if (d1 * d2 < 0) qs_push_back(isSS(p1, p2, q1, q2));
143
144
        return qs;
145 }
    // 求凸包
146
147
    std::vector<Point> convexHull(std::vector<Point> ps) {
148
        int n = ps.size(); if (n <= 1) return ps;</pre>
        std::sort(ps_begin(), ps.end());
149
150
        std::vector<Point> qs;
151
        for (int i = 0; i < n; qs.push_back(ps[i ++]))</pre>
             while (qs.size() > 1 \&\& sign(det(qs[qs.size() - 2], qs.back(), ps[i])) <= 0)
152
153
154
        for (int i = n - 2, t = qs.size(); i >= 0; qs.push_back(ps[i --]))
155
             while ((int)qs.size() > t && sign(det(qs[qs.size() - 2], qs.back(), ps[i])) <=
       → 0)
156
                 qs.pop_back();
        return qs;
```

1.1.2 凸包

```
1 // 凸包中的点按逆时针方向
    struct Convex {
         std::vector<Point> a, upper, lower;
         void make_shell(const std::vector<Point>& p,
                   std::vector<Point>& shell) { // p needs to be sorted.
              clear(shell); int n = p.size();
              for (int i = 0, j = 0; i < n; i++, j++) {
    for (; j >= 2 && sign(det(shell[j-1] - shell[j-2],
        p[i] - shell[j-2])) <= 0; --j) shell.pop_back();
                   shell.push_back(p[i]);
11
12
13
14
         void make convex() {
15
              std::sort(a.begin(), a.end());
              make_shell(a, lower);
16
17
              std::reverse(a.begin(), a.end());
             make_shell(a, upper);
a = lower; a.pop_back();
a.insert(a.end(), upper.begin(), upper.end());
if ((int)a.size() >= 2) a.pop_back();
18
19
20
21
22 23
              n = a.size();
24
25
26
27
28
29
         void init(const std::vector<Point>& _a) {
              clear(a); a = _a; n = a.size();
make_convex();
         void read(int _n) { // Won't make convex.
              clear(a); n = _n; a.resize(n);
for (int i = 0; i < n; i++)
    a[i].read();</pre>
30
31
32
33
         std::pair<DB, int> get_tangent(
34
35
36
                   const std::vector<Point>& convex, const Point& vec) {
              int l = 0, r = (int)convex.size() - 2;
              assert(r >= 0);
37
              for (; l + 1 < r; ) {
38
39
40
                   int mid = (l + r) / 2;
                   if (sign(det(convex[mid + 1] - convex[mid], vec)) > 0)
                        r = mid;
41
                   else l = mid;
42
43
              return std::max(std::make_pair(det(vec, convex[r]), r),
                        std::make_pair(det(vec, convex[0]), 0));
```

```
46
         int binary_search(Point u, Point v, int l, int r) {
47
              int s1 = sign(det(v - u, a[l % n] - u));
48
              for (; l + 1 < r; ) {
49
                    int mid = (l + r) / 2:
                    int smid = sign(det(v - u, a[mid % n] - u));
50
                    if (smid == s1) l = mid;
51
52
                   else r = mid;
53
54
              return 1 % n:
55
         }
         // 求凸包上和向量 vec 叉积最大的点,返回编号,共线的多个切点返回任意一个 int get_tangent(Point vec) {
56
57
              std::pair<DB, int> ret = get_tangent(upper, vec);
ret.second = (ret.second + (int)lower.size() - 1) % n;
58
59
60
              ret = std::max(ret, get_tangent(lower, vec));
61
               return ret.second;
62
         // 求凸包和直线 u, v 的交点,如果不相交返回 false,如果有则是和 (i, next(i))的
63
       → 交点, 交在点上不确定返回前后两条边其中之一
        → 文献、文社無工不確定返目前海内赤色共平之
bool get_intersection(Point u, Point v, int &i0, int &i1) {
    int p0 = get_tangent(u - v), p1 = get_tangent(v - u);
    if (sign(det(v - u, a[p0] - u)) * sign(det(v - u, a[p1] - u)) <= 0) {
        if (p0 > p1) std::swap(p0, p1);
        if (p0 > p1) std::swap(p0, p1);
65
66
67
68
                    i0 = binary_search(u, v, p0, p1);
69
                    i1 = binary_search(u, v, p1, p0 + n);
70
                    return true:
71
72
              else return false;
73
74 | };
```

1.2 三维

1.2.1 基础

```
// 三维绕轴旋转, 大拇指指向 axis 向量方向, 四指弯曲方向转 w 弧度
     Point rotate(const Point& s, const Point& axis, DB w) {
          DB x = axis.x, y = axis.y, z = axis.z;

DB s1 = x * x + y * y + z * z, ss1 = msqrt(s1),

cosw = cos(w), sinw = sin(w);
           DB a[4][4];
           memset(a, 0, sizeof a);
           a[3][3] = 1;
 9
           a[0][0] = ((y * y + z * z) * cosw + x * x) / s1;
          a[0][1] = x * y * (1 - cosw) / s1 + z * sinw / ss1;
a[0][2] = x * z * (1 - cosw) / s1 - y * sinw / ss1;
a[1][0] = x * y * (1 - cosw) / s1 - z * sinw / ss1;
10
11
12
          a[1][1] = ((x * x + z * z) * cosw + y * y) / s1;
a[1][2] = y * z * (1 - cosw) / s1 + x * sinw / ss1;
a[2][0] = x * z * (1 - cosw) / s1 + y * sinw / ss1;
13
14
15
16
           a[2][1] = y * z * (1 - cosw) / s1 - x * sinw / ss1;
           a[2][2] = ((x * x + y * y) * cos(w) + z * z) / s1;
17
          DB ans [4] = {0, 0, 0, 0}, c[4] = {s.x, s.y, s.z, 1}; for (int i = 0; i < 4; ++ i)
18
19
                for (int j = 0; j < 4; ++ j)
ans[i] += a[j][i] * c[j];
20
21
22
23
           return Point(ans[0], ans[1], ans[2]);
```

Chapter 2 数论

$O(m^2 \log n)$ 求线性递推数列第 n 项

```
Given a_0, a_1, \dots, a_{m-1}

a_n = c_0 \times a_{n-m} + \dots + c_{m-1} \times a_{n-1}

Solve for a_n = v_0 \times a_0 + v_1 \times a_1 + \dots + v_{m-1} \times a_{m-1}
```

```
void linear_recurrence(long long n, int m, int a[], int c[], int p) {
  long long v[M] = {1 % p}, u[M << 1], msk = !!n;</pre>
```

```
for(long long i(n); i > 1; i >>= 1) {
           msk <<= 1:
       for(long long x(0); msk; msk >>= 1, x <<= 1) {
            fill_n(u, m << 1, 0);
            int b(!!(n & msk));
            x \mid = b;
            if(x < m) {
                u[x] = 1 % p;
11
12
            }else {
13
                for(int i(0); i < m; i++) {
14
                    for(int j(0), t(i + b); j < m; j++, t++) {
15
                        u[t] = (u[t] + v[i] * v[j]) % p;
16
17
18
                for(int i((m << 1) - 1); i >= m; i--) {
                    for(int j(0), t(i - m); j < m; j++, t++) {
    u[t] = (u[t] + c[j] * u[i]) % p;
19
20
21
22
                }
23
24
            copy(u, u + m, v);
25
       //a[n] = v[0] * a[0] + v[1] * a[1] + ... + v[m - 1] * a[m - 1].
26
27
       for(int i(m); i < 2 * m; i++) {
28
            a[i] = 0:
29
            for(int j(0); j < m; j++) {
30
                a[i] = (a[i] + (long long)c[j] * a[i + j - m]) % p;
31
32
33
       for(int j(0); j < m; j++) {
34
35
            for(int i(0); i < m; i++) {
36
                b[j] = (b[j] + v[i] * a[i + j]) % p;
37
38
39
       for(int j(0); j < m; j++) {
           a[j] = b[j];
40
41
```

2.2 求逆元

```
void ex_gcd(long long a, long long b, long long &x, long long &y) {
    if (b == 0) {
        x = 1;
        y = 0;
        return;
    }
    long long xx, yy;
    ex_gcd(b, a % b, xx, yy);
    y = xx - a / b * yy;
    x = yy;
}
long long inv(long long x, long long MODN) {
    long long inv_x, y;
    ex_gcd(x, MODN, inv_x, y);
    return (inv_x % MODN + MODN) % MODN;
}
```

2.3 中国剩余定理

Chapter 3 代数

3.1 快速傅里叶变换

```
// n 必须是 2 的次幂
2
   void fft(Complex a[], int n, int f) {
        for (int i = 0; i < n; ++i)
  if (R[i] < i) swap(a[i], a[R[i]]);</pre>
        for (int i = 1, h = 0; i < n; i <<= 1, h++) {
             Complex wn = Complex(cos(pi / i), f * sin(pi / i));
             Complex w = Complex(1, 0):
8
             for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
9
             for (int p = i \ll 1, j = 0; j < n; j += p) {
                 for (int k = 0; k < i; ++k) {
10
11
                      Complex x = a[j + k], y = a[j + k + i] * tmp[k]; a[j + k] = x + y; a[j + k + i] = x - y;
12
13
14
            }
15
        }
16 }
```

Chapter 4 字符串

4.1 后缀数组

```
const int MAXN = MAXL * 2 + 1;
           int a[MAXN], x[MAXN], y[MAXN], c[MAXN], sa[MAXN], rank[MAXN], height[MAXN];
           void calc sa(int n) {
                         int m = alphabet, k = 1;
                        memset(c, 0, sizeof(*c) * (m + 1));
for (int i = 1; i <= n; ++i) c[x[i] = a[i]]++;
                         for (int i = 1; i \le m; ++i) c[i] += c[i - 1];
                         for (int i = n; i; --i) sa[c[x[i]]--] = i;
                         for (: k <= n: k <<= 1) {
10
                                       int tot = k;
                                       for (int i = n - k + 1; i \le n; ++i) y[i - n + k] = i;
11
                                       for (int i = 1; i <= n; ++i)
if (sa[i] > k) y[++tot] = sa[i] - k;
12
13
                                       memset(c, 0, sizeof(*c) * (m + 1));
14
                                     for (int i = 1; i <= n; ++i) c[x[i]]++;

for (int i = 1; i <= m; ++i) c[i] += c[i - 1];

for (int i = n; i; --i) sa[c[x[y[i]]]--] = y[i];

for (int i = 1; i <= n; ++i) y[i] = x[i];

tot = 1; x[sa[1]] = 1;
15
16
17
18
19
20
                                       for (int i = 2; i \le n; ++i) {
                                                      if (\max(sa[i], sa[i-1]) + k > n || y[sa[i]] != y[sa[i-1]] || y[sa[i] +
21
                    \rightarrow k] != y[sa[i - 1] + k]) ++tot;
                                                    x[sa[i]] = tot;
23
24
                                       if (tot == n) break; else m = tot;
25
26
27
           void calc_height(int n) {
                         for (int i = 1; i \le n; ++i) rank[sa[i]] = i;
28
                         for (int i = 1; i <= n; ++i) {
29
30
                                       height[rank[i]] = max(0, height[rank[i-1]] - 1);
31
                                       if (rank[i] == 1) continue;
32
                                       int j = sa[rank[i] - 1];
                                      while (\max(i, j) + \text{height}[\text{rank}[i]] \le n \&\& a[i + \text{height}[\text{rank}[i]]] == a[j + n \&\& a[i + \text{height}[\text{ra
33

    height[rank[i]]]) ++height[rank[i]];
34
                       }
35
```

4.2 后缀自动机

```
static const int MAXL = MAXN * 2; // MAXN is original length static const int alphabet = 26; // sometimes need changing
 3 int l, last, cnt, trans[MAXL][alphabet], par[MAXL], sum[MAXL], seq[MAXL], mxl[MAXL],
        → size[MAXL]; // mxl is maxlength, size is the size of right
    char str[MAXL]:
    inline void init() {
          l = strlen(str + 1); cnt = last = 1;
          for (int i = 0; i \le l * 2; ++i) memset(trans[i], 0, sizeof(trans[i]));
         memset(par, 0, sizeof(*par) * (l * 2 + 1));
memset(mxl, 0, sizeof(*mxl) * (l * 2 + 1));
memset(size, 0, sizeof(*size) * (l * 2 + 1));
11 }
12
    inline void extend(int pos, int c) {
         int p = last, np = last = ++cnt;
mxl[np] = mxl[p] + 1; size[np] = 1;
13
14
15
          for (; p && !trans[p][c]; p = par[p]) trans[p][c] = np;
          if (!p) par[np] = 1;
16
17
          else {
18
                int q = trans[p][c];
               if (mxl[p] + 1 == mxl[q]) par[np] = q;
19
20
                else {
21
22
                     mxl[nq] = mxl[p] + 1;
23
                     memcpy(trans[nq], trans[q], sizeof(trans[nq]));
                     par[nq] = par[q];
par[np] = par[q] = nq;
24
25
26
                     for (; trans[p][c] == q; p = par[p]) trans[p][c] = nq;
27
         }
28
29
30
    inline void buildsam() {
          for (int i = 1; i \le l; ++i) extend(i, str[i] - 'a');
31
         memset(sum, 0, sizeof(*sum) * (l * 2 + 1));
for (int i = 1; i <= cnt; ++i) sum[mxl[i]]++;
for (int i = 1; i <= l; ++i) sum[i] += sum[i - 1];
for (int i = cnt; i; --i) seq[sum[mxl[i]]--] = i;
for (int i = cnt; i; --i) size[par[seq[i]]] += size[seq[i]];</pre>
32
33
34
35
36
```

4.3 回文自动机

```
int nT, nStr, last, c[MAXT][26], fail[MAXT], r[MAXN], l[MAXN], s[MAXN]; int allocate(int len) {
       l[nT] = len;
       r[nT] = 0;
       fail[nT] = 0;
       memset(c[nT], 0, sizeof(c[nT]));
       return nT++;
8
   void init() {
   nT = nStr = 0;
10
       int newE = allocate(0):
11
       int new0 = allocate(-1):
13
       last = newE;
       fail[newE] = new0;
14
       fail new01 = newE:
15
16
       s[0] = -1:
17 }
18
   void add(int x) {
19
       s[++nStr] = x:
20
       int now = last;
21
       while (s[nStr - l[now] - 1] != s[nStr]) now = fail[now];
22
23
       if (!c[now][x]) {
            int newnode = allocate(l[now] + 2), &newfail = fail[newnode];
24
            newfail = fail[now];
25
            while (s[nStr - l[newfail] - 1] != s[nStr]) newfail = fail[newfail];
26
            newfail = c[newfail][x];
27
            c[now][x] = newnode;
28
       last = c[now][x];
```

```
30  | r[last]++;

31  | }

void count() {

    for (int i = nT - 1; i >= 0; i--) {

        r[fail[i]] += r[i];

    35  | }
```

Chapter 5 图论

5.1 基础

```
struct Graph { // Remember to call .init()!
        int e, nxt[M], v[M], adj[N], n;
        bool base;
        __inline void init(bool _base, int _n = 0) {
            assert(n < N);
            n = _n; base = _base;
7
            e = \overline{0}; memset(a\overline{d}j + base, -1, sizeof(*adj) * n);
8
9
        __inline int new_node() {
            adj[n + base] = -1;
10
11
            assert(n + base + 1 < N);
12
            return n++ + base;
13
       inline void ins(int u0, int v0) { // directional assert(u0 < n + base && v0 < n + base);
14
15
16
            v[e] = v0; nxt[e] = adj[u0]; adj[u0] = e++;
17
            assert(e < M);
18
        __inline void bi_ins(int u0, int v0) { // bi-directional
19
20
            ins(u0, v0); ins(v0, u0);
21
22 };
```

5.2 KM

```
struct KM {
       // Truly 0(n^3)
       // 邻接矩阵,不能连的边设为 -INF, 求最小权匹配时边权取负, 但不能连的还是 -INF,
      → 使用时先对 1 -> n 调用 hungary() , 再 get_ans() 求值
       int w[N][N];
       int lx[N], ly[N], match[N], way[N], slack[N];
5
       bool used[N];
       void init() {
8
            for (int i = 1; i <= n; i++) {
9
                match[i] = 0;
                lx[i] = 0;
ly[i] = 0;
10
11
                way[i] = 0;
12
13
14
15
       void hungary(int x) {
   match[0] = x;
16
17
            int j0 = 0;
            for (int j = 0; j <= n; j++) {
    slack[j] = INF;</pre>
18
19
20
                used[j] = false;
21
22
            }
23
            do {
24
                 used[j0] = true;
25
                 int i0 = match[j0], delta = INF, j1 = 0;
26
                for (int j = 1; j <= n; j++) {
   if (used[j] == false) {</pre>
27
28
                          int cur = -w[i0][j] - lx[i0] - ly[j];
29
                          if (cur < slack[i]) {</pre>
30
                              slack[i] = cur;
31
                              way[j] = j0;
```

CHAPTER 5. 图论 6

```
32
33
                        if (slack[j] < delta) {</pre>
34
35
                            delta = slack[j];
                            j1 = j;
36
37
38
               39
40
41
42
                        ly[j] -= delta;
43
44
                   else slack[j] -= delta;
45
46
               j0 = j1;
47
48
           } while (match[j0] != 0);
49
50
               int j1 = way[j0];
51
               match[j0] = match[j1];
52
                j0 = j1;
53
           } while (j0);
54
55
56
57
       int get_ans() {
           int sum = 0;
58
           for(int i = 1; i <= n; i++) {
59
                if (w[match[i]][i] == -INF); // 无解
60
                if (match[i] > 0) sum += w[match[i]][i];
61
62
           return sum;
63
   } km;
```

5.3 点双连通分量

bcc.forest is a set of connected tree whose vertices are chequered with cut-vertex and BCC.

```
const bool BCC_VERTEX = 0, BCC_EDGE = 1;
    struct BCC { // N = N0 + M0. Remember to call init(&raw_graph).
    Graph *g, forest; // g is raw graph ptr.
    int dfn[N], DFN, low[N];
    int stack[N], top;
          int expand_to[N];
                                            // Where edge i is expanded to in expaned graph.
          // Vertex \overline{i} expaned to i.
          int compress_to[N]; // Where vertex i is compressed to.
bool vertex_type[N], cut[N], compress_cut[N], branch[M];
//std::vector<int> BCC_component[N]; // Cut vertex belongs to none.
__inline void init(Graph *raw_graph) {
 8
 9
10
11
12
                g = raw_graph;
13
14
          void DFS(int u, int pe) {
    dfn[u] = low[u] = ++DFN; cut[u] = false;
15
16
                if (!~g->adj[u]) {
17
                      cut[u] = 1;
18
                      compress_to[u] = forest.new_node();
compress_cut[compress_to[u]] = 1;
19
20
21
                 for (int e = g->adj[u]; \sim e; e = g->nxt[e]) {
22
23
                       int v = g->v[e];
                       if ((e^pe) > 1 \& dfn[v] > 0 \& dfn[v] < dfn[u]) {
24
25
26
27
28
                             stack[top++] = e;
                             low[u] = std::min(low[u], dfn[v]);
                      else if (!dfn[v]) {
                             stack[top++] = e; branch[e] = 1;
                            DFS(v, e);
low[u] = std::min(low[v], low[u]);
29
30
31
                             if (low[v] >= dfn[u]) {
32
                                  if (!cut[u]) {
                                        cut[u] = 1:
```

```
34
                                compress to[u] = forest.new node():
35
                               compress_cut[compress_to[u]] = 1;
36
37
                           int cc = forest.new_node();
38
                           forest.bi ins(compress to[u], cc):
39
                           compress cut[ccl = 0:
40
                           //BCC_component[cc].clear();
41
                           do {
42
                                int cur_e = stack[--top];
                               compress_to[expand_to[cur_e]] = cc;
compress_to[expand_to[cur_e^1]] = cc;
43
44
45
                                if (branch[cur_e]) {
46
                                    int v = g->v[cur_e];
47
                                    if (cut[v])
48
                                         forest.bi ins(cc, compress to[v]);
49
                                    else {
50
                                         //BCC_component[cc].push_back(v);
51
                                         compress_to[v] = cc;
52
53
54
                           } while (stack[top] != e);
55
                     }
56
                 }
            }
57
58
59
        void solve() {
60
             forest.init(q->base);
61
             int n = g - > n;
            for (int i = 0; i < g->e; i++) {
    expand_to[i] = g->new_node();
62
63
64
65
             memset(branch, 0, sizeof(*branch) * g->e);
            memset(dfn + g->base, 0, sizeof(*dfn) * n); DFN = 0; for (int i = 0; i < n; i++)
66
67
                  if (!dfn[i + g->base]) {
68
69
                      top = 0;
70
                      DFS(i + g -> base, -1);
                 }
71
72
73
   } bcc;
75 bcc.init(&raw_graph);
76 bcc.solve():
   // Do something with bcc.forest ...
```

5.4 边双连通分量

```
struct BCC {
        Graph *g, forest;
int dfn[N], low[N], stack[N], tot[N], belong[N], vis[N], top, dfs_clock;
         // tot[] is the size of each BCC, belong[] is the BCC that each node belongs to
         pair<int, int > ori[M]; // bridge in raw_graph(raw node)
         bool is_bridge[M];
         __inline void init(Graph *raw_graph) {
8
             g = raw_graph;
             memset(is_bridge, false, sizeof(*is_bridge) * g -> e);
memset(vis + g -> base, 0, sizeof(*vis) * g -> n);
10
11
12
         void tarjan(int u, int from) {
              dfn[u] = low[u] = ++dfs_clock; vis[u] = 1; stack[++top] = u;
13
             for (int p = g -> adj[u]; ^p; p = g -> nxt[p]) {
   if ((p ^ 1) == from) continue;
14
15
                  int v = g -> v[p];
if (vis[v]) {
16
17
18
                        if (vis[v] == 1) low[u] = min(low[u], dfn[v]);
19
                  } else {
                        tarjan(v, p);
low[u] = min(low[u], low[v]);
if (low[v] > dfn[u]) is_bridge[p / 2] = true;
20
21
22
23
24
```

```
if (dfn[u] != low[u]) return;
tot[forest.new_node()] = 0;
25
26
27
                         belong[stack[top]] = forest.n;
vis[stack[top]] = 2;
28
29
30
                         tot[forest.n]++;
31
                         --top;
32
                  } while (stack[top + 1] != u);
33
34
35
36
            void solve() {
                   forest.init(g -> base);
                  int n = g -> n;
for (int i = 0; i < n; ++i)
    if (!vis[i + g -> base]) {
        top = dfs_clock = 0;
        tarjan(i + g -> base, -1);
}
37
38
39
40
41
42
                  for (int i = 0; i < g -> e / 2; ++i)
    if (is_bridge[i]) {
43
                                int e = forest.e;
44
                                forest.bi_ins(belong[q \rightarrow v[i * 2]], belong[q \rightarrow v[i * 2 + 1]], q \rightarrow
         \hookrightarrow w[i * 2]);
                                ori[e] = make_pair(g -> v[i * 2 + 1], g -> v[i * 2]);
ori[e + 1] = make_pair(g -> v[i * 2], g -> v[i * 2 + 1]);
46
47
48
49
50 } bcc;
```

Chapter 6 技巧

6.1 真正的释放 STL 容器内存空间

```
template <typename T>
__inline void clear(T& container) {
    container.clear(); // 或者删除了一堆元素
    T(container).swap(container);
}
```

6.2 无敌的大整数相乘取模

Time complexity O(1).

6.3 无敌的读入优化

```
// getchar() 读入优化 << 关同步 cin << 此优化
// 用 isdigit() 会小幅变慢
// 返回 false 表示读到文件尾
namespace Reader {
```

```
const int L = (1 << 15) + 5;
char buffer[L], *S, *T;
__inline bool getchar(char &ch) {</pre>
                if (S == T) {
                     T = (S = buffer) + fread(buffer, 1, L, stdin);
                      if (S == T) {
                           ch = EOF;
11
12
                           return false;
13
14
15
16
                ch = *S++;
                return true;
17
         __inline bool getint(int &x) {
    char ch; bool neg = 0;
    for (; getchar(ch) && (ch < '0' || ch > '9'); ) neg ^= ch == '-';
18
19
20
                if (ch == EOF) return false;
21
22
23
24
25
26
                x = ch - '0';
                for (; getchar(ch), ch >= '0' && ch <= '9'; )
    x = x * 10 + ch - '0';
                if (neg) x = -x;
                return true:
27
28 }
```

6.4 控制 cout 输出实数精度

std::cout << std::fixed << std::setprecision(5);</pre>

Chapter 7 提示

7.1 线性规划转对偶

```
maximize \mathbf{c}^T \mathbf{x}
subject to \mathbf{A} \mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq 0
\Longrightarrow
minimize \mathbf{y}^T \mathbf{b}
subject to \mathbf{y}^T \mathbf{A} > \mathbf{c}^T, \mathbf{y} > 0
```

7.2 32-bit/64-bit 随机素数

32-bit	64-bit
73550053	1249292846855685773
148898719	1701750434419805569
189560747	3605499878424114901
459874703	5648316673387803781
1202316001	6125342570814357977
1431183547	6215155308775851301
1438011109	6294606778040623451
1538762023	6347330550446020547
1557944263	7429632924303725207
1981315913	8524720079480389849

7.3 NTT 素数及其原根

Prime	Primitive	root
1053818881	7	
1051721729	6	
1045430273	3	
1012924417	5	
1007681537	3	