Gungnir's Standard Code Library

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Chapter 1 计算几何

1.1 二维

1.1.1 基础

```
typedef double DB;
   const DB eps = 1e-8;
    __inline int sign(DB x) {
       return x < -eps ? -1 : (x > eps ? 1 : 0);
    _inline DB msqrt(DB x) {
       return sign(x) > 0 ? sgrt(x) : 0;
   struct Point {
       DB x, y;
__inline Point(): x(0), y(0) {}
12
13
       <u>__inline Point(DB x, DB y): x(x), y(y) {}</u>
14
       __inline Point operator+(const Point &rhs) const {
15
16
            return Point(x + rhs.x, y + rhs.y);
17
        __inline Point operator-(const Point &rhs) const {
18
19
            return Point(x - rhs.x, y - rhs.y);
20
21
22
23
       __inline Point operator*(DB k) const {
            return Point(x * k, y * k);
24
25
26
27
28
29
30
       __inline Point operator/(DB k) const {
   assert(sign(k));
            return Point(x / k, y / k);
    _inline DB dot(const P& a, const P& b) {
31
       return a.x * b.x + a.v * b.v:
32 }
33
    _inline DB det(const P& a, const P& b) {
35
       return a.x * b.y - a.y * b.x;
```

1.1.2 凸包

```
inline void clear(std::vector<Point>& v) {
       v.clear();
       std::vector<Point>(v).swap(v);
   struct Convex {
       std::vector<Point> a, upper, lower;
       void make_shell(const std::vector<Point>& p,
               std::vector<Point>& shell) { // p needs to be sorted.
10
           clear(shell); int n = p.size();
11
           12
13
15
               shell push back(p[i]);
16
17
18
       void make convex() {
           std::sort(a.begin(), a.end());
make_shell(a, lower);
std::reverse(a.begin(), a.end());
19
20
21
22
23
           make_shell(a, upper);
24
           for (std::vector<Point>::iterator it = upper.begin(); it != upper.end(); it++)
25
               if (!(*it == *a.rbegin()) \&\& !(*it == *a.begin()))
26
                   a.push back(*it);
           n = a.size();
```

```
29
        void init(const std::vector<Point>& _a) {
30
            clear(a); a = _a; n = a.size();
make_convex();
31
32
33
        void read(int n) { // Won't make convex.
34
            clear(a); \overline{n} = \underline{n}; a.resize(n);
             for (int i = 0; i < n; i++)
35
36
                 a[i].read():
37
38
        std::pair<DB, int> get_tangent(
39
                 const std::vecTor<Point>& convex, const Point& vec) {
40
             int l = 0, r = (int)convex.size() - 2;
41
            assert(r >= 0);
42
             for (; l + 1 < r; ) {
                  int mid = (l + r) / 2;
43
                 if (sign(det(convex[mid + 1] - convex[mid], vec)) > 0)
44
45
                     r = mid:
                 else l = mid;
46
47
48
             return std::max(std::make pair(det(vec, convex[r]), r),
49
                      std::make pair(det(vec, convex[0]), 0));
50
51
        int binary_search(Point u, Point v, int l, int r) {
52
53
             int s1 = sign(det(v - u, a[l % n] - u));
            for (; l + 1 < r; ) {
int mid = (l + r) / 2;
54
55
                 int smid = sign(det(v - u, a[mid % n] - u));
                 if (smid == s1) l = mid:
57
                 else r = mid;
58
59
            return l % n:
60
        // 求凸包上和向量 vec 叉积最大的点,返回编号,共线的多个切点返回任意一个
61
62
        int get_tangent(Point vec) {
            std::pair<DB, int> ret = get_tangent(upper, vec);
ret.second = (ret.second + (int)lower.size() - 1) % n;
63
64
             ret = std::max(ret, get_tangent(lower, vec));
65
66
             return ret.second;
67
68
        // 求凸包和直线 u, v 的交点, 如果不相交返回 false, 如果有则是和 (i, next(i)) 的
      → 交点, 交在点上不确定返回前后两条边其中之一
       bool get_intersection(Point u, Point v, int &i0, int &i1) {
  int p0 = get_tangent(u - v), p1 = get_tangent(v - u);
  if (sign(det(v - u, a[p0] - u)) * sign(det(v - u, a[p1] - u)) <= 0) {
    if (p0 > p1) std::swap(p0, p1);
69
70
71
72
                 i0 = binary_search(u, v, p0, p1);
73
74
75
                 i1 = binary_search(u, v, p1, p0 + n);
                 return true:
76
77
             else return false:
78
        }
79 };
```

Chapter 2 图论

2.1 基础

```
struct Graph { // Remember to call .init()!
       int e, nxt[M], v[M], adj[N], n;
       bool base;
       __inline void init(bool _base, int _n = 0) {
5
           assert(n < N);
           n = _n; base = _base;
7
           e = \overline{0}; memset(a\overline{d}) + base, -1, sizeof(*adj) * n);
8
9
       __inline int new_node() {
           adj[n + base] = -1;
10
11
           assert(n + base + 1 < N);
12
            return n++ + base;
13
```

CHAPTER 2. 图论 3

```
__inline void ins(int u0, int v0) {    // directional
    assert(u0 < n + base && v0 < n + base);
    v[e] = v0; nxt[e] = adj[u0]; adj[u0] = e++;
    assert(e < M);
}
__inline void bi_ins(int u0, int v0) {    // bi-directional
    ins(u0, v0); ins(v0, u0);
};
</pre>
```

2.2 KM

```
1 struct KM
        // Trulv 0(n^3)
        // 邻接矩阵,不能连的边设为 -INF, 求最小权匹配时边权取负, 但不能连的还是 -INF,
      → 使用时先对 1 -> n 调用 hungary() , 再 get ans() 求值
        int w[N][N];
        int lx[N], ly[N], match[N], way[N], slack[N];
        bool used[N];
        void initialization() {
8
             for(int i = 1; i <= n; i++) {
                 match[i] = 0;
                 lx[i] = 0;
10
                 lv[i] = 0;
11
12
                 way[i] = 0;
13
14
15
16
        void hungary(int x) { // for i(1 -> n) : hungary(i);
            match[0] = x;
17
             int j0 = 0;
18
            for(int j = 0; j <= n; j++){
    slack[j] = INF;</pre>
19
20
                 used[j] = false;
21
22
23
            do {
                 used[j0] = true;
int i0 = match[j0], delta = INF, j1;
24
25
26
                 for(int j = 1; j <= n; j++) {
   if(used[j] == false) {</pre>
27
                          int cur = -w[i0][j] - lx[i0] - ly[j];
if(cur < slack[j]) {
28
29
30
                               slack[j] = cur;
31
                               way[j] = j0;
32
33
                          if(slack[j] < delta) {</pre>
34
                               delta = slack[j];
35
                               j1 = j;
36
37
38
39
                 for(int j = 0; j <= n; j++) {
   if(used[j]) {</pre>
40
                          lx[match[j]] += delta;
41
42
                          ly[j] -= delta;
43
44
                     else slack[j] -= delta;
45
46
                 i0 = i1;
47
48
            } while (match[i0] != 0);
49
50
                 int j1 = way[j0];
                 match[j0] = match[j1];
51
52
53
                 j0 = j1;
            } while(j0);
54
55
56
57
        int get_ans() { // maximum ans
            int sum = 0;
             for(int i = 1; i<= n; i++)
```

```
59 | if(match[i] > 0) sum += -w[match[i]][i]; return sum; 62 | };
```

2.3 点双连通分量

dcc.forest is a set of connected tree whose vertices are chequered with cut-vertex and DCC.

```
const bool DCC_VERTEX = 0, DCC_EDGE = 1;
   struct DCC { \overline{//} N = N0 + M0. Remember to call init(&raw graph).
        Graph *g, forest; // g is raw graph ptr.
int dfn[N], DFN, low[N];
5
        int stack[N], top;
        int expand_to[N];
                                    // Where edge i is expanded to in expaned graph.
6
        // Vertex \bar{i} expaned to i.
        int compress_to[N]; // Where vertex i is compressed to.
bool vertex_type[N], cut[N], compress_cut[N], branch[M];
//std::vector<int> DCC_component[N]; // Cut vertex belongs to none.
8
9
10
        __inline void init(Graph *raw_graph) {
11
12
             g = raw_graph;
13
14
        void DFS(int u, int pe) {
15
             dfn[u] = low[u] = ++DFN; cut[u] = false;
16
             if (!~g->adj[u]) {
17
                  cut[u] = 1:
                  compress_to[u] = forest.new_node();
18
                  compress_cut[compress_to[u]] = 1;
19
20
21
             for (int e = g\rightarrow adj[u]; \sim e; e = g\rightarrow nxt[e]) {
22
                  int v = g->v[e];
if ((e ^ pe) > 1 && dfn[v] > 0 && dfn[v] < dfn[u]) {
23
24
                      stack[top++] = e;
25
                       low[u] = std::min(low[u], dfn[v]);
26
27
                  else if (!dfn[v]) {
28
                      stack[top++] = e; branch[e] = 1;
                      DFS(v, e);
low[u] = std::min(low[v], low[u]);
if (low[v] >= dfn[u]) {
29
30
31
32
                           if (!cut[u]) {
33
                                cut[u] = 1;
                                compress to[u] = forest.new node();
34
35
                                compress cut[compress to[u]] = 1;
36
37
                            int cc = forest.new_node();
                           forest.bi_ins(compress_to[u], cc);
38
39
                           compress_cut[cc] = 0;
40
                           //DCC_component[cc].clear();
41
                           do {
42
                                int cur_e = stack[--top];
                                compress_to[expand_to[cur_e]] = cc;
43
44
                                compress_to[expand_to[cur_e^1]] = cc;
45
                                if (branch[cur_e]) {
46
                                     int v = g-v[cur_e];
                                     if (cut[v])
47
48
                                          forest.bi_ins(cc, compress_to[v]);
49
                                     else {
50
                                          //DCC_component[cc].push_back(v);
51
                                          compress to [v] = cc;
52
53
54
                           } while (stack[top] != e);
55
56
                 }
             }
57
58
59
        void solve() {
60
             forest.init(g->base);
61
             int n = g -> n;
```

Chapter 3 技巧

3.1 释放 STL 容器内存空间

```
// vectors for example.
std::vector<int> v;
// Do something with v...
v.clear(); // Or having erased many.
std::vector<int>(v).swap(v);
```

3.2 大整数相乘取模

Time complexity O(1).

```
1 // 需要保证 x 和 y 非负
2 long long mult(long long x, long long y, long long MODN) {
3 long long t = (x * y - (long long)((long double)x / MODN * y + 1e-3) * MODN) %

→ MODN;
return t < 0 ? t + MODN : t;
}
```