Gungnir's Standard Code Library

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计算几何

1.1 二维

1.1.1 基础

```
1 typedef double DB;
2
  const DB eps = 1e-8;
3
   __inline int sign(DB x) {
4
5
       return x < -eps ? -1 : (x > eps ? 1 : 0);
  }
6
7
   __inline DB msqrt(DB x) {
8
       return sign(x) > 0 ? sqrt(x) : 0;
9 }
10
  struct Point {
11
12
      DB x, y;
       __inline Point(): x(0), y(0) {}
13
       _{inline Point(DB x, DB y): x(x), y(y) {}
14
       __inline Point operator+(const Point &rhs) const {
15
           return Point(x + rhs.x, y + rhs.y);
16
17
       __inline Point operator-(const Point &rhs) const {
18
           return Point(x - rhs.x, y - rhs.y);
19
20
      __inline Point operator*(DB k) const {
21
           return Point(x * k, y * k);
22
23
       __inline Point operator/(DB k) const {
25
           assert(sign(k));
           return Point(x / k, y / k);
26
27
  };
28
29
   __inline DB dot(const P& a, const P& b) {
30
31
       return a.x * b.x + a.y * b.y;
  }
32
33
34 __inline DB det(const P& a, const P& b) {
```

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```
35 return a.x * b.y - a.y * b.x;
36 }
```

1.1.2 凸包

```
inline void clear(std::vector<Point>& v) {
2
       v.clear();
       std::vector<Point>(v).swap(v);
3
  }
4
5
  struct Convex {
6
7
       int n;
8
       std::vector<Point> a, upper, lower;
       void make_shell(const std::vector<Point>& p,
9
               std::vector<Point>& shell) { // p needs to be sorted.
10
           clear(shell); int n = p.size();
11
           for (int i = 0, j = 0; i < n; i++, j++) {
12
               for (; j \ge 2 \&\& sign(det(shell[j-1] - shell[j-2],
13
                                p[i] - shell[j-2])) \le 0; --j) shell.pop_back();
14
15
               shell.push_back(p[i]);
           }
16
       }
17
       void make_convex() {
18
           std::sort(a.begin(), a.end());
19
           make_shell(a, lower);
20
21
           std::reverse(a.begin(), a.end());
22
           make_shell(a, upper);
           a = lower;
23
24
           for (std::vector<Point>::iterator it = upper.begin(); it != upper.end(); it++)
               if (!(*it == *a.rbegin()) && !(*it == *a.begin()))
25
                   a.push_back(*it);
26
27
           n = a.size();
       }
28
       void init(const std::vector<Point>& _a) {
29
           clear(a); a = _a; n = a.size();
30
           make_convex();
31
32
       void read(int _n) { // Won't make convex.
33
           clear(a); n = _n; a.resize(n);
34
           for (int i = 0; i < n; i++)
35
               a[i].read();
36
37
       std::pair<DB, int> get_tangent(
38
39
               const std::vector<Point>& convex, const Point& vec) {
           int l = 0, r = (int)convex.size() - 2;
40
           assert(r >= 0);
41
           for (; l + 1 < r; ) {
42
               int mid = (l + r) / 2;
43
               if (sign(det(convex[mid + 1] - convex[mid], vec)) > 0)
45
                   r = mid;
               else l = mid;
46
47
           return std::max(std::make_pair(det(vec, convex[r]), r),
48
                   std::make pair(det(vec, convex[0]), 0));
49
```

1.1. 二维

```
50
      int binary_search(Point u, Point v, int l, int r) {
51
          int s1 = sign(det(v - u, a[l % n] - u));
52
          for (; l + 1 < r; ) {
53
              int mid = (l + r) / 2;
54
              int smid = sign(det(v - u, a[mid % n] - u));
55
              if (smid == s1) l = mid;
56
              else r = mid;
57
58
          return 1 % n;
59
60
      // 求凸包上和向量 vec 叉积最大的点,返回编号,共线的多个切点返回任意一个
61
62
      int get_tangent(Point vec) {
          std::pair<DB, int> ret = get_tangent(upper, vec);
63
          ret.second = (ret.second + (int)lower.size() - 1) % n;
64
          ret = std::max(ret, get_tangent(lower, vec));
65
          return ret.second;
66
67
68
      // 求凸包和直线 u, v 的交点,如果不相交返回 false,如果有则是和 (i, next(i)) 的交点,交在点上不确
    → 定返回前后两条边其中之一
      bool get_intersection(Point u, Point v, int &i0, int &i1) {
69
          int p0 = get_tangent(u - v), p1 = get_tangent(v - u);
70
71
          if (sign(det(v - u, a[p0] - u)) * sign(det(v - u, a[p1] - u)) <= 0) {
72
              if (p0 > p1) std::swap(p0, p1);
73
              i0 = binary_search(u, v, p0, p1);
              i1 = binary_search(u, v, p1, p0 + n);
74
              return true;
75
76
          else return false;
77
78
      }
79 };
```

8 CHAPTER 1. 计算几何

数论

2.1 求逆元

```
void ex_gcd(long long a, long long b, long long &x, long long &y) {
      if (b == 0) {
2
           x = 1;
3
           y = 0;
4
           return;
5
       long long xx, yy;
8
       ex_gcd(b, a % b, xx, yy);
      y = xx - a / b * yy;
9
10
       x = yy;
  }
11
12
  long long inv(long long x, long long MODN) {
13
       long long inv_x, y;
14
       ex_gcd(x, MODN, inv_x, y);
15
       return (inv_x % MODN + MODN) % MODN;
16
17 }
```

2.2 中国剩余定理

```
1 // 返回 (ans, M), 其中 ans 是模 M 意义下的解
  std::pair<long long, long long> CRT(const std::vector<long long>& m, const std::vector<long long, long
    → long>& a) {
      long long M = 1, ans = 0;
3
      int n = m.size();
4
      for (int i = 0; i < n; i++) M *= m[i];
5
6
      for (int i = 0; i < n; i++) {
          ans = (ans + (M / m[i]) * a[i] % M * inv(M / m[i], m[i])) % M; // 可能需要大整数相乘取模
7
8
      return std::make_pair(ans, M);
9
10 }
```

10 CHAPTER 2. 数论

图论

3.1 基础

```
struct Graph { // Remember to call .init()!
       int e, nxt[M], v[M], adj[N], n;
2
3
       bool base;
       __inline void init(bool _base, int _n = 0) {
4
           assert(n < N);</pre>
5
6
           n = _n; base = _base;
           e = 0; memset(adj + base, -1, sizeof(*adj) * n);
8
       __inline int new_node() {
9
           adj[n + base] = -1;
10
           assert(n + base + 1 < N);
11
           return n++ + base;
12
13
       __inline void ins(int u0, int v0) { // directional
14
           assert(u0 < n + base && v0 < n + base);
15
           v[e] = v0; nxt[e] = adj[u0]; adj[u0] = e++;
16
17
           assert(e < M);</pre>
18
       __inline void bi_ins(int u0, int v0) { // bi-directional
19
           ins(u0, v0); ins(v0, u0);
20
       }
21
22 };
```

3.2 KM

```
struct KM {
2
     // Truly 0(n^3)
     // 邻接矩阵,不能连的边设为 -INF, 求最小权匹配时边权取负,但不能连的还是 -INF, 使用时先对 1 -> n
3
    → 调用 hungary() ,再 get_ans() 求值
     int w[N][N];
     int lx[N], ly[N], match[N], way[N], slack[N];
5
     bool used[N];
6
     void init() {
7
         for (int i = 1; i \le n; i++) {
8
            match[i] = 0;
```

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```
lx[i] = 0;
10
11
                ly[i] = 0;
                way[i] = 0;
12
           }
13
14
       void hungary(int x) {
15
           match[0] = x;
16
           int j0 = 0;
17
           for (int j = 0; j <= n; j++) {
18
                slack[j] = INF;
19
20
                used[j] = false;
           }
21
22
           do {
23
                used[j0] = true;
24
25
                int i0 = match[j0], delta = INF, j1 = 0;
26
                for (int j = 1; j \le n; j++) {
                    if (used[j] == false) {
27
                         int cur = -w[i0][j] - lx[i0] - ly[j];
28
                         if (cur < slack[j]) {</pre>
29
                             slack[j] = cur;
30
                             way[j] = j0;
31
32
                        }
33
                         if (slack[j] < delta) {</pre>
34
                             delta = slack[j];
                             j1 = j;
35
                        }
36
                    }
37
38
                }
                for (int j = 0; j \le n; j++) {
39
                    if (used[j]) {
40
                        lx[match[j]] += delta;
41
                         ly[j] -= delta;
42
43
44
                    else slack[j] -= delta;
45
                }
                j0 = j1;
46
           } while (match[j0] != 0);
47
48
           do {
49
50
                int j1 = way[j0];
51
                match[j0] = match[j1];
52
                j0 = j1;
           } while (j0);
53
       }
54
55
       int get_ans() {
56
57
           int sum = 0;
58
           for(int i = 1; i <= n; i++) {
                if (w[match[i]][i] == -INF); // 无解
59
                if (match[i] > 0) sum += w[match[i]][i];
60
61
62
           return sum;
63
       }
```

3.3. 点双连通分量 13

64 } km;

3.3 点双连通分量

bcc.forest is a set of connected tree whose vertices are chequered with cut-vertex and BCC.

```
const bool BCC_VERTEX = 0, BCC_EDGE = 1;
  struct BCC { // N = N0 + M0. Remember to call init(&raw_graph).
2
       Graph *g, forest; // g is raw graph ptr.
3
       int dfn[N], DFN, low[N];
4
5
       int stack[N], top;
6
       int expand_to[N];
                               // Where edge i is expanded to in expaned graph.
       // Vertex i expaned to i.
       int compress_to[N]; // Where vertex i is compressed to.
8
       bool vertex_type[N], cut[N], compress_cut[N], branch[M];
9
       //std::vector<int> BCC_component[N]; // Cut vertex belongs to none.
10
       __inline void init(Graph *raw_graph) {
11
           g = raw_graph;
12
13
       }
       void DFS(int u, int pe) {
14
           dfn[u] = low[u] = ++DFN; cut[u] = false;
15
           if (!\sim g->adj[u]) {
16
17
               cut[u] = 1;
               compress_to[u] = forest.new_node();
18
19
               compress_cut[compress_to[u]] = 1;
20
           for (int e = g->adj[u]; \sim e; e = g->nxt[e]) {
21
22
               int v = g -> v[e];
               if ((e^p) > 1 \& dfn[v] > 0 \& dfn[v] < dfn[u]) {
23
                   stack[top++] = e;
24
25
                    low[u] = std::min(low[u], dfn[v]);
               }
26
               else if (!dfn[v]) {
                    stack[top++] = e; branch[e] = 1;
28
                   DFS(v, e);
29
30
                    low[u] = std::min(low[v], low[u]);
31
                    if (low[v] >= dfn[u]) {
                        if (!cut[u]) {
32
                            cut[u] = 1;
33
                            compress_to[u] = forest.new_node();
34
                            compress_cut[compress_to[u]] = 1;
35
                        }
36
37
                        int cc = forest.new_node();
38
                        forest.bi_ins(compress_to[u], cc);
                        compress_cut[cc] = 0;
39
                        //BCC_component[cc].clear();
40
41
                        do {
                            int cur_e = stack[--top];
42
43
                            compress_to[expand_to[cur_e]] = cc;
                            compress_to[expand_to[cur_e^1]] = cc;
44
                            if (branch[cur_e]) {
45
                                int v = g->v[cur_e];
46
                                if (cut[v])
47
```

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```
forest.bi_ins(cc, compress_to[v]);
48
                                 else {
49
                                      //BCC_component[cc].push_back(v);
50
                                      compress to[v] = cc;
51
52
53
                        } while (stack[top] != e);
54
                    }
55
               }
56
           }
57
       }
58
59
       void solve() {
           forest.init(g->base);
60
           int n = g->n;
61
           for (int i = 0; i < g -> e; i ++) {
62
                expand_to[i] = g->new_node();
63
64
           memset(branch, 0, sizeof(*branch) * g->e);
65
           memset(dfn + g->base, 0, sizeof(*dfn) * n); DFN = 0;
66
           for (int i = 0; i < n; i++)
67
                if (!dfn[i + g->base]) {
68
                    top = 0;
69
                    DFS(i + g -> base, -1);
70
71
                }
72
73
  } bcc;
74
75
  bcc.init(&raw_graph);
76 bcc.solve();
  // Do something with bcc.forest ...
```

3.4 边双连通分量

```
1
   struct BCC {
2
       Graph *g, forest;
       int dfn[N], low[N], stack[N], tot[N], belong[N], vis[N], top, dfs_clock;
3
       // tot[] is the size of each BCC, belong[] is the BCC that each node belongs to
4
       pair<int, int > ori[M]; // bridge in raw_graph(raw node)
5
       bool is_bridge[M];
6
       __inline void init(Graph *raw_graph) {
7
8
           g = raw_graph;
9
           memset(is_bridge, false, sizeof(*is_bridge) * g -> e);
10
           memset(vis + g \rightarrow base, 0, sizeof(*vis) * g \rightarrow n);
11
       void tarjan(int u, int from) {
12
           dfn[u] = low[u] = ++dfs\_clock; vis[u] = 1; stack[++top] = u;
13
            for (int p = g \rightarrow adj[u]; \sim p; p = g \rightarrow nxt[p]) {
14
                if ((p ^ 1) == from) continue;
15
16
                int v = g \rightarrow v[p];
                if (vis[v]) {
17
                     if (vis[v] == 1) low[u] = min(low[u], dfn[v]);
18
                } else {
19
                    tarjan(v, p);
20
```

3.4. 边双连通分量 15

```
low[u] = min(low[u], low[v]);
21
                     if (low[v] > dfn[u]) is_bridge[p / 2] = true;
22
                }
23
           }
24
           if (dfn[u] != low[u]) return;
25
26
           tot[forest.new_node()] = 0;
           do {
27
                belong[stack[top]] = forest.n;
28
                vis[stack[top]] = 2;
29
                tot[forest.n]++;
30
31
                --top;
           } while (stack[top + 1] != u);
32
       }
33
       void solve() {
34
           forest.init(g -> base);
35
            int n = g \rightarrow n;
36
            for (int i = 0; i < n; ++i)
37
                if (!vis[i + g -> base]) {
38
                    top = dfs_clock = 0;
39
                    tarjan(i + g \rightarrow base, -1);
40
                }
41
            for (int i = 0; i < g -> e / 2; ++i)
42
                if (is_bridge[i]) {
43
44
                    int e = forest.e;
                    forest.bi_ins(belong[g \rightarrow v[i * 2]], belong[g \rightarrow v[i * 2 + 1]], g \rightarrow w[i * 2]);
45
                    ori[e] = make_pair(g -> v[i * 2 + 1], g -> v[i * 2]);
46
                    ori[e + 1] = make_pair(g -> v[i * 2], g -> v[i * 2 + 1]);
47
48
49
50
  } bcc;
```

16 CHAPTER 3. 图论

技巧

4.1 真正的释放 STL 容器内存空间

```
// vectors for example.
std::vector<int> v;
// Do something with v...
v.clear(); // Or having erased many.
std::vector<int>(v).swap(v);
```

4.2 无敌的大整数相乘取模

Time complexity O(1).

```
// 需要保证 x 和 y 非负
long long mult(long long x, long long y, long long MODN) {
    long long t = (x * y - (long long)((long double)x / MODN * y + 1e-3) * MODN) % MODN;
    return t < 0 ? t + MODN : t;
}
```

4.3 无敌的读入优化

```
1 // getchar() 读入优化 << 关同步 cin << 此优化
  |// 用 isdigit() 会小幅变慢
  namespace Reader {
3
      const int L = (1 << 15) + 5;
      char buffer[L], *S, *T;
      __inline void get_char(char &ch) {
6
          if (S == T) {
7
              T = (S = buffer) + fread(buffer, 1, L, stdin);
8
9
              if (S == T) {
                  ch = EOF;
10
11
                  return ;
12
13
          ch = *S++;
14
15
```

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```
__inline void get_int(int &x) {
    char ch; bool neg = 0;
    for (; get_char(ch), ch < '0' || ch > '9'; ) neg ^= ch == '-';
    x = ch - '0';
    x = ch - '0';
    for (; get_char(ch), ch >= '0' && ch <= '9'; )
        x = x * 10 + ch - '0';
    if (neg) x = -x;
}</pre>
```