Gungnir's Standard Code Library

Shanghai Jiao Tong University

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Chapter 1 计算几何

1.1 二维

1.1.1 基础

```
typedef double DB;
   const DB eps = 1e-8;
   int sign(DB x) {
       return x < -eps ? -1 : (x > eps ? 1 : 0);
       return sign(x) > 0 ? sgrt(x) : 0;
10
11
   struct Point {
       DB x, y;
Point(): x(0), y(0) {}
12
13
       Point(DB x, DB y): x(x), y(y) {}
14
15
       Point operator+(const Point &rhs) const {
16
            return Point(x + rhs.x, y + rhs.y);
17
18
       Point operator-(const Point &rhs) const {
            return Point(x - rhs.x, y - rhs.y);
19
20
21
       Point operator*(DB k) const {
22
            return Point(x * k, y * k);
23
24
25
26
       Point operator/(DB k) const {
           assert(sign(k)):
            return Point(x / k, y / k);
27
28
       Point rotate(DB ang) const { // 逆时针旋转 ang 弧度
29
            return Point(cos(ang) *x - sin(ang) *y,
30
                    cos(ang) * y + sin(ang) * x);
31
       }
32
       Point turn90() const { // 逆时针旋转 90 度
33
            return Point(-y, x);
34
35
       Point unit() const {
36
           return *this / len();
37
38
39
   DB dot(const Point& a, const Point& b) {
       return a.x * b.x + a.y * b.y;
41
42
   DB det(const Point& a, const Point& b) {
43
       return a.x * b.y - a.y * b.x;
44
45 | #define cross(p1,p2,p3) ((p2.x-p1.x)*(p3.y-p1.y)-(p3.x-p1.x)*(p2.y-p1.y))
   #define crossOp(p1,p2,p3) sign(cross(p1,p2,p3))
   bool isLL(const Line& l1, const Line& l2, Point& p) { // 直线与直线交点 DB s1 = det(l2.b - l2.a, l1.a - l2.a), s2 = -det(l2.b - l2.a, l1.b - l2.a);
48
49
50
       if (!sign(s1 + s2)) return false;
51
52
53
       p = (l1.a * s2 + l1.b * s1) / (s1 + s2);
       return true:
54
   bool onSeg(const Line& l, const Point& p) { // 点在线段上
55
       return sign(det(p - l.a, l.b - l.a)) == 0 && sign(dot(p - l.a, p - l.b)) <= 0;
56
57
   Point projection(const Line & l, const Point& p) {
       return l.a + (l.b - l.a) * (dot(p - l.a, l.b - l.a) / (l.b - l.a).len2());
58
59
   DB disToLine(const Line& l, const Point& p) { // 点到 * 直线 * 距离
       return fabs(det(p - l.a, l.b - l.a) / (l.b - l.a).len());
61
62 }
63 | DB disToSeg(const Line& l, const Point& p) { // 点到线段距离
       return sign(dot(p - l.a, l.b - l.a)) * sign(dot(p - l.b, l.a - l.b)) == 1 ?
      \rightarrow disToLine(l, p) : std::min((p - l.a).len(), (p - l.b).len());
```

```
65 | }
    // 圆与直线交点
    bool isCL(Circle a, Line l, Point& p1, Point& p2) {
    DB x = dot(l.a - a.o, l.b - l.a),
        y = (l.b - l.a).len2(),
         d = x * x - y * ((l.a - a.o).len2() - a.r * a.r);
if (sign(d) < 0) return false;</pre>
 70
 71
 72
73
74
         Point p = l.a - ((l.b - l.a) * (x / y)), delta = (l.b - l.a) * (msqrt(d) / y);
         p1 = p + delta; p2 = p - delta;
 75 }
 76 //圆与圆的交面积
 77 DB areaCC(const Circle& c1, const Circle& c2) {
        DB d = (c1.o - c2.o).len();
if (sign(d - (c1.r + c2.r)) >= 0) return 0;
if (sign(d - std::abs(c1.r - c2.r)) <= 0) {
 80
 81
             DB r = std::min(c1.r, c2.r);
             return r * r * PI:
 82
 83
 84
         DB x = (d * d + c1.r * c1.r - c2.r * c2.r) / (2 * d),
 85
             t1 = acos(x / c1.r), t2 = acos((d - x) / c2.r);
 86
         return c1.r * c1.r * t1 + c2.r * c2.r * t2 - d * c1.r * sin(t1);
 87
    // 圆与圆交点
 88
    | DB s1 = (a.o - b.o).len();
 89
         if (sign(s1 - a.r - b.r) > 0 \mid | sign(s1 - std::abs(a.r - b.r)) < 0) return false;
         DB s2 = (a.r * a.r - b.r * b.r) / s1;
        DB aa = (s1 + s2) * 0.5, bb = (s1 - s2) * 0.5;
P o = (b.o - a.o) * (aa / (aa + bb)) + a.o;
 94
         P delta = (b.o - a.o).unit().turn90() * msqrt(a.r * a.r - aa * aa);
 95
 96
97
         p1 = o + delta, p2 = o - delta;
         return true:
 98 }
    // 求点到圆的切点,按关于点的顺时针方向返回两个点
 99
    bool tanCP(const Circle &c, const Point &p0, Point &p1, Point &p2) {
100
         double x = (p0 - c.o) \cdot len2(), d = x - c.r * c.r;
         if (d < eps) return false; // 点在圆上认为没有切点
         Point p = (p0 - c.o) * (c.r * c.r / x);
Point delta = ((p0 - c.o) * (-c.r * sqrt(d) / x)).turn90();
103
104
105
         p1 = c.o + p + delta;
         p2 = c.o + p - delta;
106
107
         return true;
108 }
    // 求圆到圆的内共切线,按关于 cl.o 的顺时针方向返回两条线
std::vector<Line> intanCC(const Circle &c1, const Circle &c2) {
111
         std::vector<Line> ret;
         Point p = (c1.0 * c2.r + c2.0 * c1.r) / (c1.r + c2.r);
112
113
         Point p1, p2, q1, q2;
114
         if (tanCP(c1, p, p1, p2) && tanCP(c2, p, q1, q2)) { // 两圆相切认为没有切线
115
             ret.push_back(Line(p1, q1));
116
             ret_push_back(Line(p2, q2));
117
118
         return ret:
119
    // 点在多边形内
120
    bool inPolygon(const Point& p, const std::vector<Point>& poly) {
         int n = polv.size():
123
         int counter = 0;
124
         for (int i = 0; i < n; ++ i) {
125
             P = poly[i], b = poly[(i + 1) % n];
126
             if (onSeg(Line(a, b), p)) return false; // 边界上不算
127
             int x = sign(det(p - a, b - a));
             int y = sign(a.y - p.y);
128
             int z = sign(b.y - p.y);
129
             if (x < 0 & & z < 0 & & y > 0) ++ counter;
if (x < 0 & & z <= 0 & & y > 0) -- counter;
130
131
132
133
         return counter != 0;
134
```

```
135 // 用半平面 (q1,q2) 的逆时针方向去切凸多边形
    std::vector<Point> convexCut(const std::vector<Point>&ps, Point q1, Point q2) {
136
137
        std::vector<Point> qs; int n = ps.size();
        for (int i = 0; i < n; ++i) {
   Point p1 = ps[i], p2 = ps[(i + 1) % n];
   int d1 = crossOp(q1,q2,p1), d2 = crossOp(q1,q2,p2);
138
139
140
141
             if (d1 \ge 0) qs.push_back(p1);
142
             if (d1 * d2 < 0) qs_push_back(isSS(p1, p2, q1, q2));
143
144
        return qs;
145 }
    // 求凸包
146
147
    std::vector<Point> convexHull(std::vector<Point> ps) {
148
        int n = ps.size(); if (n <= 1) return ps;</pre>
        std::sort(ps_begin(), ps.end());
149
150
        std::vector<Point> qs;
151
        for (int i = 0; i < n; qs.push_back(ps[i ++]))</pre>
             while (qs.size() > 1 \&\& sign(det(qs[qs.size() - 2], qs.back(), ps[i])) <= 0)
152
153
154
        for (int i = n - 2, t = qs.size(); i >= 0; qs.push_back(ps[i --]))
155
             while ((int)qs.size() > t && sign(det(qs[qs.size() - 2], qs.back(), ps[i])) <=
       → 0)
156
                 qs.pop_back();
        return qs;
```

1.1.2 凸包

```
1 // 凸包中的点按逆时针方向
    struct Convex {
         std::vector<Point> a, upper, lower;
         void make_shell(const std::vector<Point>& p,
                   std::vector<Point>& shell) { // p needs to be sorted.
              clear(shell); int n = p.size();
              for (int i = 0, j = 0; i < n; i++, j++) {
    for (; j >= 2 && sign(det(shell[j-1] - shell[j-2],
        p[i] - shell[j-2])) <= 0; --j) shell.pop_back();
                   shell.push_back(p[i]);
11
12
13
14
         void make convex() {
15
              std::sort(a.begin(), a.end());
              make_shell(a, lower);
16
17
              std::reverse(a.begin(), a.end());
             make_shell(a, upper);
a = lower; a.pop_back();
a.insert(a.end(), upper.begin(), upper.end());
if ((int)a.size() >= 2) a.pop_back();
18
19
20
21
22 23
              n = a.size();
24
25
26
27
28
29
         void init(const std::vector<Point>& _a) {
              clear(a); a = _a; n = a.size();
make_convex();
         void read(int _n) { // Won't make convex.
              clear(a); n = _n; a.resize(n);
for (int i = 0; i < n; i++)
    a[i].read();</pre>
30
31
32
33
         std::pair<DB, int> get_tangent(
34
35
36
                   const std::vector<Point>& convex, const Point& vec) {
              int l = 0, r = (int)convex.size() - 2;
              assert(r >= 0);
37
              for (; l + 1 < r; ) {
38
39
40
                   int mid = (l + r) / 2;
                   if (sign(det(convex[mid + 1] - convex[mid], vec)) > 0)
                        r = mid;
41
                   else l = mid;
42
43
              return std::max(std::make_pair(det(vec, convex[r]), r),
                        std::make_pair(det(vec, convex[0]), 0));
```

```
46
         int binary_search(Point u, Point v, int l, int r) {
47
              int s1 = sign(det(v - u, a[l % n] - u));
48
              for (; l + 1 < r; ) {
49
                    int mid = (l + r) / 2:
                    int smid = sign(det(v - u, a[mid % n] - u));
50
                    if (smid == s1) l = mid;
51
52
                   else r = mid;
53
54
              return 1 % n:
55
         }
         // 求凸包上和向量 vec 叉积最大的点,返回编号,共线的多个切点返回任意一个 int get_tangent(Point vec) {
56
57
              std::pair<DB, int> ret = get_tangent(upper, vec);
ret.second = (ret.second + (int)lower.size() - 1) % n;
58
59
60
              ret = std::max(ret, get_tangent(lower, vec));
61
               return ret.second;
62
         // 求凸包和直线 u, v 的交点,如果不相交返回 false,如果有则是和 (i, next(i))的
63
       → 交点, 交在点上不确定返回前后两条边其中之一
        → 文献、文社無工不確定返目前海内赤色共平之
bool get_intersection(Point u, Point v, int &i0, int &i1) {
    int p0 = get_tangent(u - v), p1 = get_tangent(v - u);
    if (sign(det(v - u, a[p0] - u)) * sign(det(v - u, a[p1] - u)) <= 0) {
        if (p0 > p1) std::swap(p0, p1);
        if (p0 > p1) std::swap(p0, p1);
65
66
67
68
                    i0 = binary_search(u, v, p0, p1);
69
                    i1 = binary_search(u, v, p1, p0 + n);
70
                    return true:
71
72
              else return false;
73
74 | };
```

1.2 三维

1.2.1 基础

```
// 三维绕轴旋转, 大拇指指向 axis 向量方向, 四指弯曲方向转 w 弧度
     Point rotate(const Point& s, const Point& axis, DB w) {
          DB x = axis.x, y = axis.y, z = axis.z;

DB s1 = x * x + y * y + z * z, ss1 = msqrt(s1),

cosw = cos(w), sinw = sin(w);
           DB a[4][4];
           memset(a, 0, sizeof a);
           a[3][3] = 1;
 9
           a[0][0] = ((y * y + z * z) * cosw + x * x) / s1;
          a[0][1] = x * y * (1 - cosw) / s1 + z * sinw / ss1;
a[0][2] = x * z * (1 - cosw) / s1 - y * sinw / ss1;
a[1][0] = x * y * (1 - cosw) / s1 - z * sinw / ss1;
10
11
12
          a[1][1] = ((x * x + z * z) * cosw + y * y) / s1;
a[1][2] = y * z * (1 - cosw) / s1 + x * sinw / ss1;
a[2][0] = x * z * (1 - cosw) / s1 + y * sinw / ss1;
13
14
15
16
           a[2][1] = y * z * (1 - cosw) / s1 - x * sinw / ss1;
           a[2][2] = ((x * x + y * y) * cos(w) + z * z) / s1;
17
          DB ans [4] = {0, 0, 0, 0}, c[4] = {s.x, s.y, s.z, 1}; for (int i = 0; i < 4; ++ i)
18
19
                for (int j = 0; j < 4; ++ j)
ans[i] += a[j][i] * c[j];
20
21
22
23
           return Point(ans[0], ans[1], ans[2]);
```

Chapter 2 数论

$O(m^2 \log n)$ 求线性递推数列第 n 项

```
Given a_0, a_1, \dots, a_{m-1}

a_n = c_0 \times a_{n-m} + \dots + c_{m-1} \times a_{n-1}

Solve for a_n = v_0 \times a_0 + v_1 \times a_1 + \dots + v_{m-1} \times a_{m-1}
```

```
void linear_recurrence(long long n, int m, int a[], int c[], int p) {
  long long v[M] = {1 % p}, u[M << 1], msk = !!n;</pre>
```

```
for(long long i(n); i > 1; i >>= 1) {
           msk <<= 1:
       for(long long x(0); msk; msk >>= 1, x <<= 1) {
            fill_n(u, m << 1, 0);
            int b(!!(n & msk));
            x \mid = b;
            if(x < m) {
                u[x] = 1 % p;
11
12
            }else {
13
                for(int i(0); i < m; i++) {
14
                    for(int j(0), t(i + b); j < m; j++, t++) {
15
                        u[t] = (u[t] + v[i] * v[j]) % p;
16
17
18
                for(int i((m << 1) - 1); i >= m; i--) {
                    for(int j(0), t(i - m); j < m; j++, t++) {
    u[t] = (u[t] + c[j] * u[i]) % p;
19
20
21
22
                }
23
24
            copy(u, u + m, v);
25
       //a[n] = v[0] * a[0] + v[1] * a[1] + ... + v[m - 1] * a[m - 1].
26
27
       for(int i(m); i < 2 * m; i++) {
28
            a[i] = 0:
29
            for(int j(0); j < m; j++) {
30
                a[i] = (a[i] + (long long)c[j] * a[i + j - m]) % p;
31
32
33
       for(int j(0); j < m; j++) {
34
35
            for(int i(0); i < m; i++) {
36
                b[j] = (b[j] + v[i] * a[i + j]) % p;
37
38
39
       for(int j(0); j < m; j++) {
           a[j] = b[j];
40
41
```

2.2 求逆元

```
void ex_gcd(long long a, long long b, long long &x, long long &y) {
    if (b == 0) {
        x = 1;
        y = 0;
        return;
    }
    long long xx, yy;
    ex_gcd(b, a % b, xx, yy);
    y = xx - a / b * yy;
    x = yy;
}
long long inv(long long x, long long MODN) {
    long long inv_x, y;
    ex_gcd(x, MODN, inv_x, y);
    return (inv_x % MODN + MODN) % MODN;
}
```

2.3 中国剩余定理

Chapter 3 代数

3.1 快速傅里叶变换

```
// n 必须是 2 的次幂
2
   void fft(Complex a[], int n, int f) {
        for (int i = 0; i < n; ++i)
  if (R[i] < i) swap(a[i], a[R[i]]);</pre>
        for (int i = 1, h = 0; i < n; i <<= 1, h++) {
             Complex wn = Complex(cos(pi / i), f * sin(pi / i));
             Complex w = Complex(1, 0):
8
             for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
9
             for (int p = i \ll 1, j = 0; j < n; j += p) {
                 for (int k = 0; k < i; ++k) {
10
11
                      Complex x = a[j + k], y = a[j + k + i] * tmp[k]; a[j + k] = x + y; a[j + k + i] = x - y;
12
13
14
            }
15
        }
16 }
```

Chapter 4 字符串

4.1 后缀数组

```
const int MAXN = MAXL * 2 + 1;
           int a[MAXN], x[MAXN], y[MAXN], c[MAXN], sa[MAXN], rank[MAXN], height[MAXN];
           void calc sa(int n) {
                         int m = alphabet, k = 1;
                        memset(c, 0, sizeof(*c) * (m + 1));
for (int i = 1; i <= n; ++i) c[x[i] = a[i]]++;
                         for (int i = 1; i \le m; ++i) c[i] += c[i - 1];
                         for (int i = n; i; --i) sa[c[x[i]]--] = i;
                         for (: k <= n: k <<= 1) {
10
                                       int tot = k;
                                       for (int i = n - k + 1; i \le n; ++i) y[i - n + k] = i;
11
                                       for (int i = 1; i <= n; ++i)
if (sa[i] > k) y[++tot] = sa[i] - k;
12
13
                                       memset(c, 0, sizeof(*c) * (m + 1));
14
                                     for (int i = 1; i <= n; ++i) c[x[i]]++;

for (int i = 1; i <= m; ++i) c[i] += c[i - 1];

for (int i = n; i; --i) sa[c[x[y[i]]]--] = y[i];

for (int i = 1; i <= n; ++i) y[i] = x[i];

tot = 1; x[sa[1]] = 1;
15
16
17
18
19
20
                                       for (int i = 2; i \le n; ++i) {
                                                      if (\max(sa[i], sa[i-1]) + k > n || y[sa[i]] != y[sa[i-1]] || y[sa[i] +
21
                    \rightarrow k] != y[sa[i - 1] + k]) ++tot;
                                                    x[sa[i]] = tot;
23
24
                                       if (tot == n) break; else m = tot;
25
26
27
           void calc_height(int n) {
                         for (int i = 1; i \le n; ++i) rank[sa[i]] = i;
28
                         for (int i = 1; i <= n; ++i) {
29
30
                                       height[rank[i]] = max(0, height[rank[i-1]] - 1);
31
                                       if (rank[i] == 1) continue;
32
                                       int j = sa[rank[i] - 1];
                                      while (\max(i, j) + \text{height}[\text{rank}[i]] \le n \&\& a[i + \text{height}[\text{rank}[i]]] == a[j + n \&\& a[i + \text{height}[\text{ra
33

    height[rank[i]]]) ++height[rank[i]];
34
                       }
35
```

4.2 后缀自动机

```
static const int MAXL = MAXN * 2; // MAXN is original length static const int alphabet = 26; // sometimes need changing
 3 int l, last, cnt, trans[MAXL][alphabet], par[MAXL], sum[MAXL], seq[MAXL], mxl[MAXL],
       → size[MAXL]; // mxl is maxlength, size is the size of right
    char str[MAXL]:
    inline void init() {
         l = strlen(str + 1); cnt = last = 1;
         for (int i = 0; i \le l * 2; ++i) memset(trans[i], 0, sizeof(trans[i]));
        memset(par, 0, sizeof(*par) * (l * 2 + 1));
memset(mxl, 0, sizeof(*mxl) * (l * 2 + 1));
memset(size, 0, sizeof(*size) * (l * 2 + 1));
11 }
12
    inline void extend(int pos, int c) {
         int p = last, np = last = ++cnt;
mxl[np] = mxl[p] + 1; size[np] = 1;
13
14
15
         for (; p && !trans[p][c]; p = par[p]) trans[p][c] = np;
         if (!p) par[np] = 1;
16
17
         else {
18
              int q = trans[p][c];
19
              if (mxl[p] + 1 == mxl[q]) par[np] = q;
20
              else {
21
22
                   mxl[nq] = mxl[p] + 1;
23
                   memcpy(trans[nq], trans[q], sizeof(trans[nq]));
                   par[nq] = par[q];
par[np] = par[q] = nq;
24
25
26
                    for (; trans[p][c] == q; p = par[p]) trans[p][c] = nq;
27
28
29
         }
30
    inline void buildsam() {
         for (int i = 1; i <= l; ++i) extend(i, str[i] - 'a');
31
        memset(sum, 0, sizeof(*sum) * (l * 2 + 1));
for (int i = 1; i <= cnt; ++i) sum[mxl[i]]++;
for (int i = 1; i <= l; ++i) sum[i] += sum[i - 1];
32
33
34
35
         for (int i = cnt; i; --i) seq[sum[mxl[i]]--] = i;
36
         for (int i = cnt; i; --i) size[par[seq[i]]] += size[seq[i]];
```

4.3 回文自动机

```
int nT, nStr, last, c[MAXT][26], fail[MAXT], r[MAXN], l[MAXN], s[MAXN]; int allocate(int len) {
       l[nT] = len;
       r[nT] = 0;
       fail[nT] = 0;
       memset(c[nT], 0, sizeof(c[nT]));
       return nT++;
8
   void init() {
       nT = nStr = 0;
10
       int newE = allocate(0);
11
       int new0 = allocate(-1):
13
       last = newE;
       fail[newE] = newO;
14
15
       fail[new0] = newE:
16
       s[0] = -1;
17 }
18
   void add(int x) {
19
       s[++nStr] = x:
20
       int now = last;
21
       while (s[nStr - l[now] - 1] != s[nStr]) now = fail[now];
22
23
       if (!c[now][x]) {
            int newnode = allocate(l[now] + 2), &newfail = fail[newnode];
24
           newfail = fail[now];
25
            while (s[nStr - l[newfail] - 1] != s[nStr]) newfail = fail[newfail];
26
           newfail = c[newfail][x];
27
            c[now][x] = newnode;
28
       last = c[now][x];
```

```
30  | r[last]++;

31  | }

void count() {

    for (int i = nT - 1; i >= 0; i--) {

        r[fail[i]] += r[i];

    35  | }
```

Chapter 5 数据结构

5.1 KD-Tree

```
long long norm(const long long &x) {
2
       // For manhattan distance
       return std::abs(x);
       // For euclid distance
5
       return x * x;
6
8
   struct Point {
       int x, y, id;
10
       const int& operator [] (int index) const {
  if (index == 0) {
11
12
13
                return x;
14
15
           } else {
                return y;
16
17
18
19
       friend long long dist(const Point &a, const Point &b) {
20
            long long result = 0;
21
22
            for (int i = 0; i < 2; ++i) {
                result += norm(a[i] - b[i]);
23
24
            return result;
25
26
27
28
   } point[N];
   struct Rectangle {
29
30
       int min[2], max[2];
31
       Rectangle() {
32
33
            min[0] = min[1] = INT_MAX; // sometimes int is not enough
            \max[0] = \max[1] = INTMIN;
34
35
36
       void add(const Point &p) {
            for (int i = 0; i < 2; ++i) {
37
                min[i] = std::min(min[i], p[i]);
38
39
                max[i] = std::max(max[i], p[i]);
40
       }
41
42
43
       long long dist(const Point &p) {
44
            long long result = 0;
45
            for (int i = 0; i < 2; ++i) {
46
                // For minimum distance
47
                result += norm(std::min(std::max(p[i], min[i]), max[i]) - p[i]);
48
                // For maximum distance
49
                result += std::max(norm(max[i] - p[i]), norm(min[i] - p[i]));
50
51
            return result;
52
53
54
55
   struct Node {
56
57
       Point seperator;
       Rectangle rectangle;
58
59
       int child[2];
       void reset(const Point &p) {
60
```

CHAPTER 5. 数据结构 6

```
seperator = p;
61
62
             rectangle = Rectangle();
63
             rectangle.add(p);
64
            child[0] = child[1] = 0;
65
66
   } tree[N << 1]:</pre>
67
    int size, pivot;
69
70
    bool compare(const Point &a, const Point &b) {
71
        if (a[pivot] != b[pivot]) {
72
            return a[pivot] < b[pivot];</pre>
73
74
        return a.id < b.id;</pre>
75
76 }
77
    // 左閉右開: build(1, n + 1)
    int build(int l, int r, int type = 1) {
78
79
        pivot = type;
80
        if (l >= r) {
81
            return 0;
82
83
        int x = ++size:
84
        int mid = 1 + r >> 1:
        std::nth_element(point + l, point + mid, point + r, compare);
85
86
        tree[x].reset(point[mid]):
87
        for (int i = l: i < r: ++i)
88
            tree[x].rectangle.add(point[i]);
89
90
        tree[x].child[0] = build(l, mid, type ^ 1);
91
92
        tree[x].child[1] = build(mid + 1, r, type ^1);
93 }
95
96
    int insert(int x, const Point &p, int type = 1) {
97
        if (x == 0)
98
            tree[++size].reset(p);
99
            return size;
100
101
        tree[x].rectangle.add(p);
102
        if (compare(p, tree[x] seperator)) {
103
            tree[x].child[0] = insert(tree[x].child[0], p, type ^ 1);
104
105
            tree[x].child[1] = insert(tree[x].child[1], p, type ^ 1);
106
107
        return x;
108
109
110
    // For minimum distance
111 // For maximum: 下面递归 query 时 0, 1 换顺序;< and >;min and max
112
    void query(int x, const Point &p, std::pair<long long, int> &answer, int type = 1) {
113
        pivot = type;
114
        if (x == 0 || tree[x].rectangle.dist(p) > answer.first) {
115
            return;
116
117
        answer = std::min(answer,
                  std::make_pair(dist(tree[x].seperator, p), tree[x].seperator.id));
118
119
        if (compare(p, tree[x].seperator)) {
            query(tree[x].child[0], p, answer, type ^ 1);
query(tree[x].child[1], p, answer, type ^ 1);
120
121
122
123
            query(tree[x].child[1], p, answer, type ^ 1);
            query(tree[x].child[0], p, answer, type ^ 1);
124
125
126
128
    std::priority_queue<std::pair<long long, int> > answer;
129
130
    void query(int x, const Point &p, int k, int type = 1) {
131
        pivot = type;
        if (x == 0 \mid | (int)answer.size() == k && tree[x].rectangle.dist(p) >
       → answer.top().first) {
```

```
133
             return:
134
135
        answer.push(std::make_pair(dist(tree[x].seperator, p), tree[x].seperator.id));
136
        if ((int)answer size(\overline{)} > k) {
137
            answer.pop():
138
139
        if (compare(p, tree[x].seperator)) {
140
             query(tree[x].child[0], p, k, type ^ 1);
141
             query(tree[x].child[1], p, k, type ^ 1);
142
143
             query(tree[x].child[1], p, k, type ^ 1);
             query(tree[x] child[0], p, k, type ^ 1);
144
145
146 }
```

5.2 Treap

```
struct Node{
        int mn, key, size, tag;
        bool rev;
        Node* ch[2]:
        Node(int mn, int key, int size): mn(mn), key(key), size(size), rev(0), tag(0){}
        void downtag();
        Node* update(){
8
            mn = min(ch[0] \rightarrow mn, min(key, ch[1] \rightarrow mn));
            size = ch[0] -> size + 1 + ch[1] -> size;
10
            return this:
11
12 | };
13
   typedef pair<Node*, Node*> Pair;
14 | Node *null, *root;
15
   void Node::downtag(){
16
        if(rev){
17
            for(int i = 0; i < 2; i++)
   if(ch[i] != null){</pre>
18
                      ch[i] -> rev ^= 1;
19
                      swap(ch[i] \rightarrow ch[0], ch[i] \rightarrow ch[1]);
20
21
22
            rev = 0:
23
24
25
        if(tag){
            for(int i = 0; i < 2; i++)
26
                 if(ch[i] != null){
27
                      ch[i] \rightarrow key += tag;
28
                      ch[i] \rightarrow mn += tag;
29
                      ch[i] -> tag += tag;
30
31
            tag = 0;
32
33
34
   int r(){
35
        static int s = 3023192386;
36
        return (s += (s << 3) + 1) & (\sim0u >> 1);
37
38
   bool random(int x, int y){
39
        return r() % (x + y) < x;
40
41
   Node* merge(Node *p, Node *q){
42
        if(p == null) return q;
43
        if(g == null) return p;
44
        p -> downtag();
45
        q -> downtag();
46
        if(random(p -> size, q -> size)){
47
            p \rightarrow ch[1] = merge(p \rightarrow ch[1], q);
48
             return p -> update();
49
        }else{
50
            q \rightarrow ch[0] = merge(p, q \rightarrow ch[0]);
51
            return q -> update();
52
        }
53 | }
```

```
54 | Pair split(Node *x, int n) {
55
        if(x == null) return make_pair(null, null);
56
        x -> downtag();
57
        if(n \ll x \rightarrow ch[0] \rightarrow size)
             Pair ret = split(x -> ch[0], n);
x -> ch[0] = ret.second;
58
59
60
              return make pair(ret.first, x -> update());
61
62
        Pair ret = split(x \rightarrow ch[1], n - x \rightarrow ch[0] \rightarrow size - 1);
63
        x \rightarrow ch[1] = ret.first;
64
        return make_pair(x -> update(), ret.second);
65
66
   pair<Node*, Pair> get_segment(int l, int r){
67
        Pair ret = split(root, l - 1);
68
         return make_pair(ret.first, split(ret.second, r - l + 1));
69
70
71
72
73
        null = new Node(INF, INF, 0);
null -> ch[0] = null -> ch[1] = null;
        root = null;
```

5.3 Link/cut Tree

```
inline void reverse(int x) {
        tr[x].rev ^= 1; swap(tr[x].c[0], tr[x].c[1]);
   inline void rotate(int x, int k) {
        int y = tr[x] fa, z = tr[y] fa;
        tr[x].fa = z; tr[z].c[tr[z].c[1] == y] = x;
        tr[tr[x].c[k ^ 1]].fa = y; tr[y].c[k] = tr[x].c[k ^ 1];
        tr[x].c[k ^ 1] = y; tr[y].fa = x;
10 }
12
   inline void splay(int x, int w) {
13
        int z = x; pushdown(x);
14
        while (tr[x].fa != w) {
15
            int y = tr[x].fa; z = tr[y].fa;
16
            if (z == w)
                pushdown(z = y); pushdown(x);
rotate(x, tr[y].c[1] == x);
update(y); update(x);
17
18
19
20
21
                 pushdown(z); pushdown(y); pushdown(x);
22
                 int t1 = tr[y] \cdot c[1] == x, t2 = tr[z] \cdot c[1] == y;
23
                 if (t1 == t2) rotate(y, t2), rotate(x, t1);
24
25
26
27
                else rotate(x, t1), rotate(x, t2);
                 update(z); update(x); update(x);
28
29
30
31
32
        update(x);
        if (x != z) par[x] = par[z], par[z] = 0;
    inline void access(int x) {
33
        for (int y = 0; x; y = x, x = par[x]) {
34
            splay(x, 0);
35
            if (tr[x].c[1]) par[tr[x].c[1]] = x, tr[tr[x].c[1]].fa = 0;
            tr[x].c[1] = y; par[y] = 0; tr[y].fa = x; update(x);
37
38
39
40
   inline void makeroot(int x) {
41
        access(x); splay(x, 0); reverse(x);
42
43
44
   inline void link(int x, int y) {
45
        makeroot(x); par[x] = y;
46 | }
48 | inline void cut(int x, int y) {
```

Chapter 6 图论

6.1 基础

```
struct Graph { // Remember to call .init()!
        int e, nxt[M], v[M], adj[N], n;
        bool base;
4
        __inline void init(bool _base, int _n = 0) {
5
            assert(n < N);</pre>
            n = _n; base = _base;
            e = \overline{0}; memset(a\overline{d}j + base, -1, sizeof(*adj) * n);
        __inline int new_node() {
10
            adi[n + base] = -1:
11
            assert(n + base + 1 < N);
12
            return n++ + base;
13
14
        __inline void ins(int u0, int v0) { // directional assert(u0 < n + base \&\& v0 < n + base);
15
16
            v[e] = v0; nxt[e] = adj[u0]; adj[u0] = e++;
17
            assert(e < M);
18
19
        __inline void bi_ins(int u0, int v0) { // bi-directional
20
            ins(u0, v0); ins(v0, u0);
21
22 };
```

6.2 KM

```
struct KM {
        // Truly 0(n^3)
        // 邻接矩阵,不能连的边设为 -INF, 求最小权匹配时边权取负, 但不能连的还是 -INF,
      → 使用时先对 1 -> n 调用 hungary() ,再 get_ans() 求值
        int w[N][N];
5
        int lx[N], ly[N], match[N], way[N], slack[N];
        bool used[N];
        void init() {
8
            for (int i = 1; i <= n; i++) {
    match[i] = 0;</pre>
                lx[i] = 0;

ly[i] = 0;
11
12
                way[i] = 0;
13
14
15
        void hungary(int x) {
16
            match[0] = x;
17
            int j0 = 0;
            for (int j = 0; j <= n; j++) {
    slack[j] = INF;</pre>
18
19
20
                used[j] = false;
21
22
23
24
            do {
                 used[i0] = true;
25
                 int i0 = match[j0], delta = INF, j1 = 0;
                for (int j = 1; j <= n; j++) {
    if (used[j] == false) {
26
27
                          int cur = -w[i0][j] - lx[i0] - ly[j];
28
29
                          if (cur < slack[i]) {</pre>
30
                              slack[j] = cur;
```

CHAPTER 6. 图论

```
31
                               way[j] = j0;
32
33
                           if (slack[j] < delta) {</pre>
34
                               delta = slack[j];
35
                               j1 = j;
36
37
38
                 for (int j = 0; j <= n; j++) {
   if (used[j]) {</pre>
39
40
                           lx[match[j]] += delta;
41
42
                           ly[j] -= delta;
43
44
                      else slack[i] -= delta;
45
46
                 j0 = j1;
47
             } while (match[j0] != 0);
48
49
50
                  int j1 = way[j0];
                 match[j0] = match[j1];
51
52
53
54
55
                  j0 = j1;
             } while (j0);
56
57
58
        int get_ans() {
             int sum = 0;
             for(int i = 1; i <= n; i++) {
59
                  if (w[match[i]][i] == -INF); // 无解
60
                 if (match[i] > 0) sum += w[match[i]][i];
61
62
             return sum;
63
64 | } km;
```

6.3 点双连通分量

bcc.forest is a set of connected tree whose vertices are chequered with cut-vertex and BCC.

```
const bool BCC_VERTEX = 0, BCC_EDGE = 1;
struct BCC {   // N = N0 + M0. Remember to call init(&raw_graph).
    Graph *g, forest; // g is raw graph ptr.
    int dfn[N], DFN, low[N];
    int draw[N]
         int stack[N], top;
         int expand to[N]:
                                         // Where edge i is expanded to in expaned graph.
         // Vertex \overline{i} expaned to i.
         int compress_to[N]; // Where vertex i is compressed to.
bool vertex_type[N], cut[N], compress_cut[N], branch[M];
//std::vector<int> BCC_component[N]; // Cut vertex belongs to none.
10
11
         __inline void init(Graph *raw_graph) {
12
               g = raw_graph;
13
14
         void DFS(int u, int pe)
15
               dfn[u] = low[u] = ++DFN; cut[u] = false;
16
               if (!\sim g->adj[u]) {
17
                    cut[u] = 1;
18
                    compress to[u] = forest.new node();
                    compress_cut[compress_to[u]] = 1;
19
20
21
               for (int e = g->adj[u]; \sim e; e = g->nxt[e]) {
22
                     int v = g->v[e];
23
                     if ((e^pe) > 1 \&\& dfn[v] > 0 \&\& dfn[v] < dfn[u]) {
24
25
26
27
                          stack[top++] = e;
                          low[u] = std::min(low[u], dfn[v]);
                    else if (!dfn[v]) {
28
                          stack[top++] = e; branch[e] = 1;
29
                          DFS(v, e);
30
                          low[u] = std::min(low[v], low[u]);
                          if (low[v] >= dfn[u]) {
31
32
                               if (!cut[u]) {
```

```
33
                              cut[u] = 1:
                              compress_to[u] = forest.new_node();
34
                              compress_cut[compress_to[u]] = 1;
35
36
37
                         int cc = forest.new node();
38
                          forest bi ins(compress to[u], cc):
39
                          compress_cut[cc] = 0;
                          //BCC_component[cc].clear();
40
41
                         do {
42
                              int cur_e = stack[--top];
                              compress_to[expand_to[cur_e]] = cc;
43
44
                              compress_to[expand_to[cur_e^1]] = cc;
45
                              if (branch[cur_e]) {
46
                                  int v = q - v[cur e];
47
                                  if (cut[v])
48
                                       forest.bi_ins(cc, compress_to[v]);
49
                                  else {
50
                                       //BCC_component[cc].push_back(v);
51
                                       compress_to[v] = cc;
52
53
                         } while (stack[top] != e);
54
55
56
                }
57
            }
58
        void solve() {
59
60
            forest.init(q->base);
61
            int n = g -> n;
            for (int i = 0; i < q -> e; i ++)
62
63
                 expand_to[i] = g->new_node();
64
65
            memset(branch, 0, sizeof(*branch) * g->e);
66
            memset(dfn + g->base, 0, sizeof(*dfn) * n); DFN = 0; for (int i = 0; i < n; i++)
67
68
                 if (!dfn[i + g->base]) {
69
                     top = 0;
70
                     DFS(i + g \rightarrow base, -1);
71
72
73
74
   } bcc;
75 bcc.init(&raw_graph);
76 bcc.solve():
   // Do something with bcc.forest ...
```

6.4 边双连通分量

```
struct BCC {
       Graph *g, forest;
int dfn[N], low[N], stack[N], tot[N], belong[N], vis[N], top, dfs_clock;
        // tot[] is the size of each BCC, belong[] is the BCC that each node belongs to
        pair<int, int > ori[M]; // bridge in raw graph(raw node)
        bool is_bridge[M];
        __inline void init(Graph *raw_graph) {
            g = raw_graph;
            memset(is_bridge, false, sizeof(*is_bridge) * q -> e);
9
10
            memset(vis + g \rightarrow base, 0, sizeof(*vis) * g \rightarrow n);
11
12
        void tarjan(int u, int from) {
13
            dfn[u] = low[u] = ++dfs\_clock; vis[u] = 1; stack[++top] = u;
14
            for (int p = g -> adj[u]; ~p; p = g -> nxt[p]) {
   if ((p ^ 1) == from) continue;
15
16
                 int v = g -> v[p];
if (vis[v]) {
17
18
                     if (vis[v] == 1) low[u] = min(low[u], dfn[v]);
19
                 } else {
20
                     tarjan(v, p);
                      low[u] = min(low[u], low[v]);
21
                      if (low[v] > dfn[u]) is_bridge[p / 2] = true;
22
23
```

```
24
25
26
              if (dfn[u] != low[u]) return;
              tot[forest.new_node()] = 0;
27
28
29
                   belong[stack[top]] = forest.n;
                   vis(stack(top)) = 2:
30
31
32
33
                   tot[forest.n]++;
                    --top;
              } while (stack[top + 1] != u);
34
35
36
         void solve() {
              forest.init(g -> base);
              int n = g \rightarrow n;
37
              for (int i = 0; i < n; ++i)
                   if (!vis[i + g -> base]) {
   top = dfs_clock = 0;
38
39
40
                        tarjan(i + g \rightarrow base, -1);
41
42
              for (int i = 0; i < g -> e / 2; ++i)
    if (is_bridge[i]) {
43
44
                        int e = forest.e;
45
                        forest.bi_ins(belong[g \rightarrow v[i * 2]], belong[g \rightarrow v[i * 2 + 1]], g \rightarrow
       \hookrightarrow w[i * 2]);
                        ori[e] = make_pair(g -> v[i * 2 + 1], g -> v[i * 2]);
46
47
                        ori[e + 1] = make_pair(q \rightarrow v[i * 2], q \rightarrow v[i * 2 + 1]);
48
49
50 } bcc;
```

Chapter 7 其他

7.1 Dancing Links

```
struct Node {
        Node *1, *r, *u, *d, *col;
        int size, line_no;
        Node() {
            size = 0; line_no = -1;
            l = r = u = d = col = NULL;
   } *root;
10
   void cover(Node *c) {
        c->l->r = c->r; c->r->l = c->l;
11
12
        for (Node *u = c->d; u != c; u = u->d)
13
            for (Node *v = u->r; v != u; v = v->r) {
14
                v->d->u = v->u:
15
                v->u->d = v->d;
16
                 -- v->col->size;
17
18 }
19
20
    void uncover(Node *c) {
21
        for (Node *u = c->u; u != c; u = u->u) {
22
23
24
25
            for (Node *v = u->1; v != u; v = v->1) {
                 ++ v->col->size;
                v\rightarrow u\rightarrow d = v;
                v->d->u = v:
26
27
28
        c - > l - > r = c; c - > r - > l = c;
29 }
31
   std::vector<int> answer;
   bool search(int k) {
32
33
        if (root->r == root) return true;
34
35
        Node *r = NULL:
        for (Node *u = root -> r; u != root; u = u -> r)
36
37
            if (r == NULL || u->size < r->size)
38
        if (r == NULL || r->size == 0) return false;
        else {
```

```
40
             cover(r);
41
             bool succ = false;
42
             for (Node *u = r->d; u != r && !succ; u = u->d) {
43
                  answer.push_back(u->line_no);
44
                  for (Node *v = u \rightarrow r; v != u; v = v \rightarrow r) // Cover row
45
                      cover(v->col);
46
                  succ \mid = search(k + 1);
                  for (Node *v = u->1; v != u; v = v->1)
47
48
                      uncover(v->col);
49
                  if (!succ) answer.pop_back();
50
51
52
             uncover(r);
             return succ;
53
54
55
    bool entry[CR][CC];
57 Node *who[CR][CC];
    int cr, cc;
60
    void construct() {
         root = new Node();
61
         Node *last = root;
63
         for (int i = 0; i < cc; ++ i) {
             Node *u = new Node();
64
65
             last->r = u; u->l = last;
66
             Node *v = u; u->line_no = i;
             last = u:
67
68
             for (int j = 0; j < cr; ++ j)
69
                  if (entry[j][i]) {
70
                      ++ u->size;
                      Node *cur = new Node();
 71
 72
                      who[j][i] = cur;
 73
                      cur->line_no = j;
 74
                      cur->col = u;
 75
                      cur->u = v; v->d = cur;
76
                      v = cur;
77
78
             v->d = u; u->u = v;
79
80
         last->r = root; root->l = last;
81
         for (int j = 0; j < cr; ++ j) {
82
             Node *last = NULL;
83
             for (int i = cc - 1; i >= 0; -- i)
if (entry[j][i]) {
84
85
                      last = who[i][i];
86
                      break;
87
             for (int i = 0; i < cc; ++ i)
    if (init[j][i]) {</pre>
88
89
90
                      last->r = who[j][i];
                      who[j][i]->l = last;
last = who[j][i];
91
92
93
94
        }
95
96
97
    void destruct() {
98
         for (Node *u = root->r; u != root; ) {
             for (Node *v = u - > d; v != u; ) {
100
                 Node *nxt = v->d;
101
                  delete(v);
102
                 v = nxt;
103
             Node *nxt = u->r:
104
105
             delete(u); u = nxt;
106
107
         delete root;
108
```

Chapter 8 技巧

8.1 真正的释放 STL 容器内存空间

```
template <typename T>
__inline void clear(T& container) {
    container.clear(); // 或者删除了一堆元素
    T(container).swap(container);
}
```

8.2 无敌的大整数相乘取模

Time complexity O(1).

```
1 // 需要保证 x 和 y 非负 long long mult(long long x, long long MODN) {
3 long long t = (x * y - (long long)((long double)x / MODN * y + 1e-3) * MODN) %

→ MODN;
return t < 0 ? t + MODN : t;
}
```

8.3 无敌的读入优化

```
1 // getchar() 读入优化 << 关同步 cin << 此优化 2 // 用 isdigit() 会小幅变慢 3 // 返回 false 表示读到文件尾 namespace Reader {
         const int L = (1 << 15) + 5;
char buffer[L], *S, *T;
         __inline bool getchar(char &ch) {
              if (S == T) {
                   T = (S = buffer) + fread(buffer, 1, L, stdin);
                   if (S == T) {
11
                        ch = EOF;
                         return false;
12
13
14
15
              ch = *S++;
16
              return true;
17
         __inline bool getint(int &x) {
```

```
char ch; bool neg = 0;
for (; getchar(ch) && (ch < '0' || ch > '9'); ) neg ^= ch == '-';
if (ch == EOF) return false;
x = ch - '0';
for (; getchar(ch), ch >= '0' && ch <= '9'; )
x = x * 10 + ch - '0';
if (neg) x = -x;
return true;
}
```

8.4 控制 cout 输出实数精度

```
std::cout << std::fixed << std::setprecision(5);</pre>
```

Chapter 9 提示

9.1 线性规划转对偶

```
maximize \mathbf{c}^T \mathbf{x}
subject to \mathbf{A} \mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq 0
\iff
minimize \mathbf{y}^T \mathbf{b}
subject to \mathbf{y}^T \mathbf{A} \geq \mathbf{c}^T, \mathbf{y} \geq 0
```

9.2 32-bit/64-bit 随机素数

32-bit	64-bit
73550053	1249292846855685773
148898719	1701750434419805569
189560747	3605499878424114901
459874703	5648316673387803781
1202316001	6125342570814357977
1431183547	6215155308775851301
1438011109	6294606778040623451
1538762023	6347330550446020547
1557944263	7429632924303725207
1981315913	8524720079480389849

9.3 NTT 素数及其原根

Prime	Primitive root
1053818881	7
1051721729	6
1045430273	3
1012924417	5
1007681537	3