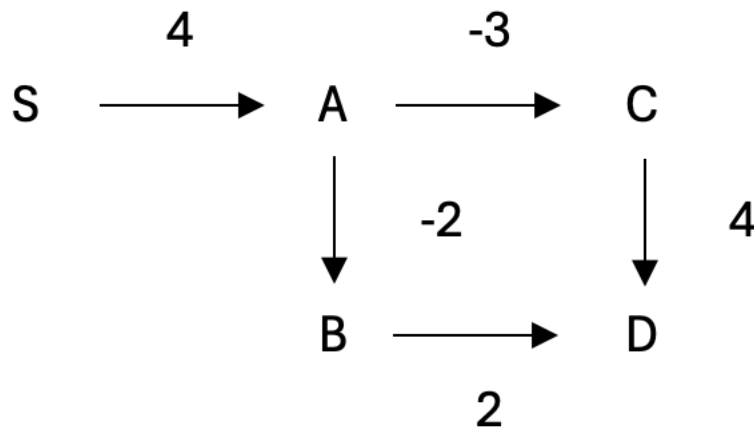


**Question 1** (15 points). Please provide an example of a weighted, directed graph  $G = (V; E)$  which has some edge weights negative, and Dijkstra's algorithm correctly finds the shortest paths from a source  $s$  in the graph. You can assume the graph has no negative cycles. Please clearly indicate the source  $s$  in your graph and which edge weights are negative.

Dijkstra's algorithm is not designed to work correctly with graphs that contain negative edge weights. However, there is an edge case where Dijkstra's algorithm can still function correctly:

- If there are negative edge weights, but **no negative cycles**, Dijkstra's algorithm can work correctly if the negative edges only appear in situations where they don't affect the optimal path. I give an example about no negative cycle as follows:



Iter	PQ	d(S)	d(A)	d(B)	d(C)	d(D)
1	[S]	0	$\infty$	$\infty$	$\infty$	$\infty$
2	[A]	0	4	$\infty$	$\infty$	$\infty$
3	[C,B]	0	4	2	1	$\infty$
4	[B,D]	0	4	2	1	5
5	[D]	0	4	2	1	4
6	[]	0	4	2	1	4

**Dijkstra's algorithm** can still find the shortest path in this graph because the negative edges are not part of the **shortest path** from the source to the other vertices.

**Question 2** (20 points). *There is a network of roads  $G = (V; E)$  connecting a set of cities  $V$ . Each road in  $E$  has an associated length  $l_e$ . There is a proposal to add **one** new road to this network, and there is a list  $E'$  of pairs of cities between which the new road can be built. Each such potential road  $e' \in E'$  has an associated length. As a designer for the public works department you are asked to determine the road  $e' \in E'$  whose addition to the existing network  $G$  would result in the maximum decrease in the driving distance between two fixed cities  $s$  and  $t$  in the network.*

(i) (15 points) *Describe an efficient algorithm by using Dijkstra's algorithm in English for solving this problem.*

(ii) (5 points) *Please explain the running time of your algorithm.*

1.

i. **Initially, calculate  $G = (V; E)$  to use Dijkstra's algorithm twice:**

First Run: Find the shortest path from city  $s$  to all other cities; Second Run: Find the shortest path from city  $t$  to all other cities.

ii. **Evaluate Each Potential Road:**

for  $u, v, l$  in `potential_cities`:

`max_dist = min(dist_s[u] + l + dist_t[v], dist_t[v] + dist_s[u] + l)`

iii. **Determine the Maximum Decrease in Distance:**

`decrease = dist_s[t] - max_dist`

if `decrease > max_deceased`:

`max_deceased = decrease`

iv. **Output the Best Road:**

`def find_max_decrease(self, s, t, potential_cities):`

`dist_s = self.dijkstra(s)`

`dist_t = self.dijkstra(t)`

`max_deceased = 0`

`best_road = None`

for  $u, v, l$  in `potential_cities`:

`max_dist = min(dist_s[u] + l + dist_t[v], dist_t[v] + dist_s[u] + l)`

`decrease = dist_s[t] - max_dist`

if `decrease > max_deceased`:

`max_deceased = decrease`

`best_road = (u, v, l)`

`return best_road, max_deceased`

2.

i. **Dijkstra's Algorithm:** The time complexity of Dijkstra's algorithm is  $O((V+E)\log V)$ . Running algorithm twice will multiply 2.

ii. Comparing the new distances for each road takes  $O(|E'|)$ .

iii. Total time complexity is  $O((V+E)\log V) + O(|E'|)$ .

**Question 3** (25 points). Suppose you are given an infinite supply of coins whose values are one of 1¢, 5¢, 10¢ and a dollar value  $N$ , which is a positive integer.

(i) (10 points) Please describe an efficient greedy algorithm in English to make change for  $N$ ¢ using the three denominations of coins.

(ii) (2 points) What is the running time of your algorithm?

(iii) (5 points) Please describe the order of coins and their denominations your algorithm will use when  $N = 13$ ¢?

(iv) (8 points) Suppose you are not given 5¢ but instead given 6¢ denomination. Please explain if your algorithm will still work correctly. Why?

- i. Firstly, Start with the largest denomination (10¢). Next, use the 5¢ denomination. Finally, use the 1¢ denomination.
- ii. The running time is  **$O(1)$** .
- iii.

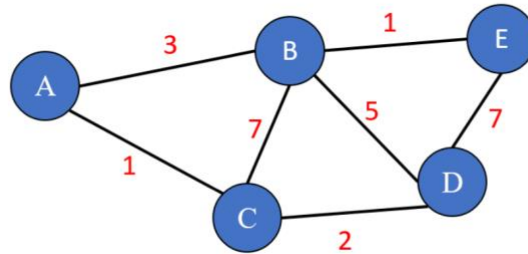
Step	N	Order of Coins	Number of Coins
1	13	10	1
2	3	1	1
3	2	1	1
4	1	1	1

iv.

If the 5¢ coin is replaced with a 6¢ coin, **the greedy algorithm will still work correctly** because:

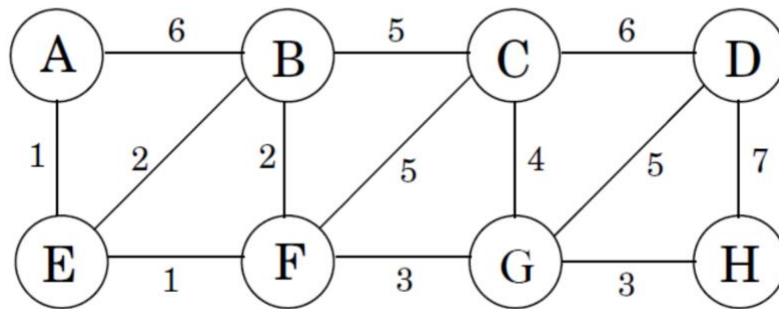
- The greedy approach always tries to use the largest denomination first.
- The problem with a 5¢ coin is that sometimes it would lead to suboptimal results.
- With a 6¢ coin, the algorithm will still select the largest possible denomination, and since 6 is a divisor of 12, it will correctly make change for numbers like 12 or 13 without suboptimal behavior.

**Question 4** (20 points). Please execute Dijkstra's algorithm from vertex A in the following graph and fill the table below.  $d(\cdot)$  denotes the shortest distance from A to the vertex. The column "PQ" should only list the vertices (in order) in the Priority Queue whose distances  $\neq \infty$ . Break all ties lexicographically (i.e, according to alphabetical order).



Iter	PQ	d(A)	d(B)	d(C)	d(D)	d(E)
1	[A]	0	$\infty$	$\infty$	$\infty$	$\infty$
2	[C, B]	0	3	1	$\infty$	$\infty$
3	[B, D]	0	3	1	3	$\infty$
4	[D, E]	0	3	1	3	4
5	[E]	0	3	1	3	4
6	[]	0	3	1	3	4

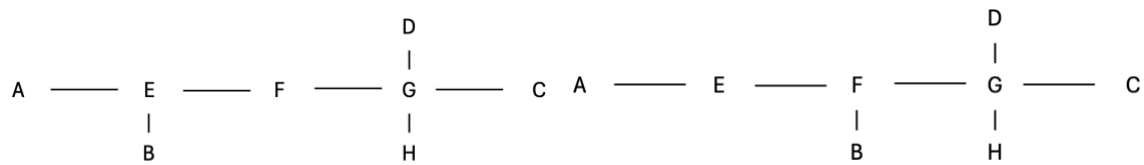
**Question 5** (20 points). Consider the following graph.



1.

Edge	Weight	T/F	Weight
AE	1	T	1
EF	1	T	1
EB	2	T	2
BF	2	F	
GH	3	T	3
FG	3	T	3
CG	4	T	4
BC	5	F	
CF	5	F	
DG	5	T	5
AB	6	F	
DH	7	F	
the cost of MST			19

2.



There are 2 MST in this graph.

3.

**Step1: Sort the list of edges**

Weight	Edge
1	AE,EF
2	EB,BF
3	GH,FG
4	CG
5	BC,CF,DG
6	AB
7	DH

**Step2: Calculate the MST**

Weight	Edge	Find_Vertex	Total_Dist
1	<b>AE,EF</b>	A,E,F	2
2	<b>EB,BF</b>	A,E,F,B	4
3	<b>GH,FG</b>	A,E,F,B,G,H	10
4	<b>CG</b>	A,E,F,B,G,H,C	14
5	<b>BC,CF,DG</b>	A,E,F,B,G,H,C,D	19
6	AB	-	19
7	DH	-	19