## **Question 1**

# (a) (5 points) Show that $f(x) = x^2 + 4x$ is $O(x^2)$ .

$$0 x^2 + 4x \le c |x^2|$$

$$0 x^2 + 4x \le x^2 + 4x^2 \forall x > 1$$

$$0 x^2 + 4x \le 5x^2 \forall x > 1$$

o 
$$x_0 = 1 \text{ and } c = 5$$

Therefore,  $f(x)=O(x^2)$ .

# (b) (5 points) Show that $f(x) = x^2$ is NOT $O(\sqrt{x})$ .

If 
$$x^2 \le c\sqrt{x}$$
;

$$x^{3/2} \le c$$
;

As  $x \to \infty$ ,  $x^{3/2} \to \infty$ , but c is a constant.

Therefore,  $f(x) = x^2$  is NOT  $O(\sqrt{x})$ .

#### Note:

$$x^2 \le x^2 \quad \forall x > 1$$

$$x^2 \le x^2 \quad \forall x > 1$$

$$x_0 = 1$$
 and  $c = 1$ 

# (c) (5 points) Show that f(x) = x is $\Omega(\log x)$ .

$$x => c \log x$$

$$2^{x} => cx$$

$$x_0 = 1, c = 2$$

Therefore, f(x) = x is  $\Omega(\log x)$ .

(d) (10 points) Show that  $f(x) = (2x^2 - 3)/((3x^4 + x^3 - 2x^2 - 1))$  is  $\Theta(x^{-2})$ .

$$f(x) = (2x^2 - 3)/((3x^4 + x^3 - 2x^2 - 1)$$

 $2x^2 - 3$  is nearly equal to  $2x^2$ ;

 $3x^4 + x^3 - 2x^2 - 1$  is nearly equal to  $3x^4$ ;

 $f(x)\approx 2/3x^2$ ;

#### Big-O:

$$f(x) \le cg(x), x > x_0$$

$$2/3x^2 \le c x^{-2}$$

Therefore, c = 1 and  $x_0 = 1$ ,  $O(x^{-2})$ ;

## Big-Ω:

$$f(x) => cg(x), x > x_0$$

$$2/3x^2 => c x^{-2}$$

Therefore, c = 1/3 and  $x_0 = 1$ ,  $\Omega(x^{-2})$ ;

Thus, that  $f(x) = (2x^2 - 3)/((3x^4 + x^3 - 2x^2 - 1))$  is  $\Theta(x^{-2})$ .

Question 2 (15 points). Please rank the following functions based on their  $O(\cdot)$  complexity of running time. The function that has the least complexity should be ranked 1. Please explain your answer to get full credit.

$$f_1(x) = x \log_2 x$$

$$f_2(x) = 3^x$$

$$f_3(x) = \sqrt{x}$$

$$f_4(x) = x!$$

$$f_5(x) = 2^x$$

#### Rank:

$$f_3(x) \rightarrow f_1(x) \rightarrow f_2(x) \rightarrow f_4(x)$$

Explain my answer(Big-O):

$$f_3(x) = \sqrt{x} <= c\sqrt{x} \quad \forall x > 1 \Rightarrow c = 2; x_0 = 1$$

Big-O Complexity is  $O(\sqrt{x})$ , so it is slower than x.

$$f_1(x) = x \log_2 x \le cx^2 \ \forall x > 1 \Rightarrow c = 1; x_0 = 1$$

Big-O Complexity is  $O(xlog_2x)$ , so it is faster than x.

 $f_5(x)$  and  $f_2(x)$  belong to  $n^x$  type.

Therefore, Big-  $\Omega$  for  $2^x$  is  $x^2$ .

Big-O Complexity is  $O(n^x)$ , so  $3^x$  is faster than  $2^x$ 

$$f_4(x) = x! => c3^x \forall x > 4 \Rightarrow c = 1; x_0 = 7$$

Therefore, Big- $\Omega$  for x! is  $3^x$ .

And then,

## Big-O

$$f_4(x) = x! \le cx! \ \forall x > 1 \Rightarrow c = 2; x_0 = 1$$

#### Result of Rank:

$$f_3(x) \rightarrow f_1(x) \rightarrow f_5(x) \rightarrow f_2(x) \rightarrow f_4(x)$$

#### Question 3:

(a) (15 points) Please describe an efficient algorithm in English using a data structure such as array / linked list / stack / queue to solve this problem.

#### 1. Create a new Stack:

• It can store opening parenthesis ( {, (, [) when they are encountered.

## 2. Traverse the String

- Use "for" to loop through each character into string.
- If it's an opening parenthesis, "push" it into "stack".
- If it's a closing parenthesis:

First, check if the stack is empty.

Secondly, "pop" the top element and check if it matches the closing parenthesis.

Finally, check if the stack is empty.

# (b) (5 points) What is the asymptotic upper bound of complexity of running time for your algorithm?

```
time
Stack<Character> queue = new Stack<>();
                                                                                    1
for(char character : string.toCharArray()){
 if (character == '{' \parallel character == '(' \parallel character == '[') {
                                                                                           Number of opening brackets
   queue.push(character);
                                                                                           Number of opening brackets
 else if(character == '}' \parallel character == ')' \parallel character == ']'){
                                                                                           Number of closing brackets
   if (queue.isEmpty()) {
                                                                                           Number of closing brackets
    return false;
   char left = queue.pop();
                                                                                           Number of closing brackets
                                                                                           Number of closing brackets
   if (!isMatch(left, character)) {
    return false;
                                                                                    1
return queue.isEmpty();
                                                                                    1
```

## Big-O:

```
f(x) = n + 6k + 4(k < n) \approx n + 6n + 4 <= 7n + 4n \forall n > 1 \Rightarrow c = 11; x_0 = 1
f(n) \ grows \ as \ O(n)
```

(c) (40 points) Please write a program in either Python / Java / C++ that realizes your algorithm in (a). To receive full credit, please structure your code, write comments and show the output for the above two examples.

## Example1: I { love [ the {rains}()]}

xinggchen@Xinggs-MacBook-Pro HW1 % cd /Users/xi
ualMachines/jdk-21.jdk/Contents/Home/bin/java -X
orkspaceStorage/5b186955a9a9c5c41e227983c8196a36
The output is true

Example2: "I{love[the{rains]()"; xinggchen@xinggs=MacBook=Pro HW.ualMachines/jdk=21.jdk/Contents, orkspaceStorage/5b186955a9a9c5c4The output is false