

1. No. From the definition $f(n)=O(1)$ means for real number C and N , there exist $f(n) \leq C \cdot 1$ ($n > N$) so we get $f(n) \leq C$. and $f(n)=O(n/100)$ means there exist $f(n) \leq C \cdot n/100$ ($n > N$), we know that when $n > N$, there exist real number which satisfied $C < C \cdot n/100$, so $f(n)=O(1) < O(n/100)$.
2. No. From the definition $g(n)=O(n)$ means for real number C and N , there exist $g(n) \leq C \cdot n$ ($n > N$) so we get $g(n) \leq Cn$, and $g(n)=O(100)$ means there exist $g(n) \leq C \cdot 100$ ($n > N$), we know that when $n > N$, Cn will bigger than $C \cdot 100$ so we cannot claim $g(n) = O(100)$.
3. Yes. $A(n)=f(n)+g(n)=O(1)+O(n)$, so we need to proof that $O(1)+O(n)=O(n)$.
 $O(1)+O(n)=O(n) \Rightarrow$
 $O(1)+O(n) < Cn$;
 $A+Bn < Cn$;
Because $A < A \cdot n$ in this domain.
So we just get $A+B \cdot n < A \cdot n+B \cdot n < C \cdot n \Rightarrow (K) \cdot n < C \cdot n$
So $A(n)=O(n)$.
4. Yes. $B(n)=f(n)g(n)=O(1) \cdot O(n)$, so we need to proof that $O(1) \cdot O(n)=O(n)$
 $O(1) \cdot O(n)=O(n) \Rightarrow$
 $O(1) \cdot O(n)=C \cdot n$;
 $A \cdot B \cdot n=C \cdot n; \Rightarrow K \cdot n=C \cdot n$
So $B(n)=O(n)$;

Bonus

1. $C(n)=O(1)$ means $f(O(n))=O(1)$. We know $g(n)=O(n)$ means $g(n)$ can be written by $A \cdot n+B$, that means the complexity of $g(n)$ is $A \cdot n+B$. if we set $A \cdot n+B$ into code $f(n)$, because the complexity of $f(n)=O(1)$, so the complexity of $f(O(n))=M(A \cdot n+B)=O(n)=O(1)$.
2. $D(n)=O(1)$ means $g(O(1))=O(1)$. We know $f(n)=O(1)$ means the complexity of $f(n)$ is a constant C , then we input that constant to the function $g(n)$, because C is constant, the complexity of $g(C)$ must be constant so $D(n)=O(1)$.