

# Hashing, Hash Function and Hash Table

# What is Hashing

- Hashing is a technique that is used to uniquely identify a specific object from a group of similar objects.
- Examples of how hashing is used in our lives:
  - In universities, each student is assigned a unique roll number that can be used to retrieve information about them.
  - In libraries, each book is assigned a unique number that can be used to determine information about the book
- **Hashing** is the solution that can be used in almost all such situations and performs extremely well compared to above data structures like Array, Linked List, BST in practice

# What is Hashing (Cond...)

- With hashing we get  $O(1)$  search time on average and  $O(n)$  in worst case.
- The idea is to use ***hash function*** that converts a given input number or any other key to a smaller number and uses the small number as an index in a table called ***hash table***.
- Hash Function
  - Hash function maps a big number or string to a small integer that can be used as an index in the hash table
- Hash Table
  - An array that stores pointers to records corresponding to a given input number





Mia	M	i	a	279	4
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				Mia						
0	1	2	3	4	5	6	7	8	9	10

Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1

	Tim				Mia									
0	1	2	3	4	5	6	7	8	9	10				

Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5
Jan	J	74	a	97	n	110	281	6
Ada	A	65	d	100	a	97	262	9
Leo	L	76	e	101	o	111	288	2

Bea	Tim	Leo		Mia	Zoe	Jan			Ada	
0	1	2	3	4	5	6	7	8	9	10



Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5
Jan	J	74	a	97	n	110	281	6
Ada	A	65	d	100	a	97	262	9
Leo	L	76	e	101	o	111	288	2
Sam	S	83	a	97	m	109	289	3
Lou	L	76	o	111	u	117	304	7
Max	M	77	a	97	x	120	294	8
Ted	T	84	e	101	d	100	285	10

Bea	Tim	Leo	Sam	Mia	Zoe	Jan	Lou	Max	Ada	Ted
0	1	2	3	4	5	6	7	8	9	10

Index number = sum ASCII codes Mod size of array

Bea	Tim	Leo	Sam	Mia	Zoe	Jan	Lou	Max	Ada	Ted
0	1	2	3	4	5	6	7	8	9	10

Find Ada

$$\text{Ada} = (65 + 100 + 97) = 262$$

$$262 \bmod 11 = 9$$

myData = Array(9)

Bea	Tim	Leo	Sam	Mia	Zoe	Jan	Lou	Max	Ada	Ted
0	1	2	3	4	5	6	7	8	9	10

Bea	Tim	Leo	Sam	Mia	Zoe	Jan	Lou	Max	Ada	Ted
27/01/1941	08/06/1955	31/12/1945	27/04/1791	20/02/1906	19/06/1978	13/02/1956	27/12/1922	29/04/1898	10/12/1815	17/06/1937
English	English	American	American	Russian	American	Polish	French	German	English	American
Astronomer	Inventor	Mathematician	Inventor	Spice taster	Actress	Linguist	Biologist	Physicist	Mathematician	Philosopher

0    1    2    3    4    5    6    7    8    9    10

# Hashing Algorithm

- Calculation applied to a key to transform it into an address
- For numeric keys, divide the key by the number of available addresses and take the remainder
- $$\text{address} = \text{key} \text{ Mod } n$$
- For alphanumeric keys, divide the sum of ASCII codes in a key by the number of available addresses,  $n$ , and take the remainder

- Folding method divides key into equal parts then adds the parts together
  - The telephone number 01452 8345654, becomes  $01 + 45 + 28 + 34 + 56 + 54 = 212$
  - Depending on size of table, may then divide by some constant and take remainder

# Understand Collision

Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5

Bea	Tim				Mia	Zoe						
0	1	2	3	4	5	6	7	8	9	10		

# Understand Collision

Mia	M	77	i	105	a	97	279
Tim	T	84	i	105	m	109	298
Bea	B	66	e	101	a	97	264
Zoe	Z	90	o	111	e	101	302
Sue	S	83	u	117	e	101	301

Bea	Tim				Mia	Zoe						
0	1	2	3	4	5	6	7	8	9	10		

# Collision Handling

- Irrespective of how good a hash function is, collisions are bound to occur.
- Therefore, to maintain the performance of a hash table, it is important to manage collisions through various collision resolution techniques.
  - Open addressing
    - Linear Probing
    - Quadratic Probing
    - Double Hashing
  - Closed addressing



# Open Addressing (Linear)

Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5
Sue	S	83	u	117	e	101	301	4



Bea	Tim				Mia	Zoe	Sue						
0	1	2	3	4	5	6	7	8	9	10			

Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5
Sue	S	83	u	117	e	101	301	4
Len	L	76	e	101	n	110	287	1



Bea	Tim	Len		Mia	Zoe	Sue				
0	1	2	3	4	5	6	7	8	9	10

Mia	M	77	i	105	a	97	279	4
Tim	T	84	i	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5
Sue	S	83	u	117	e	101	301	4
Len	L	76	e	101	n	110	287	1
Moe	M	77	o	111	e	101	289	3
Lou	L	76	o	111	u	117	304	7
Rae	R	82	a	97	e	101	280	5
Max	M	77	a	97	x	120	294	8
Tod	T	84	o	111	d	100	295	9

Bea	Tim	Len	Moe	Mia	Zoe	Sue	Lou	Rae	Max	Tod
0	1	2	3	4	5	6	7	8	9	10

Find Rae

$$\text{Rae} = (82 + 97 + 101) = 280$$

$$280 \bmod 11 = 5$$

myData = Array(5)

Bea	Tim	Len	Moe	Mia	Zoe	Sue	Lou	Rae	Max	Tod
0	1	2	3	4	5	6	7	8	9	10

Find Rae

$$\text{Rae} = (82 + 97 + 101) = 280$$

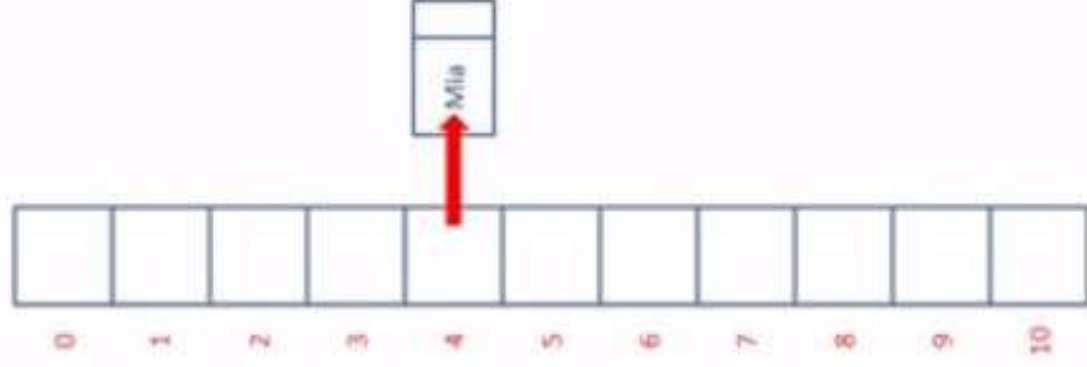
$$280 \bmod 11 = 5$$

myData = Array(5)

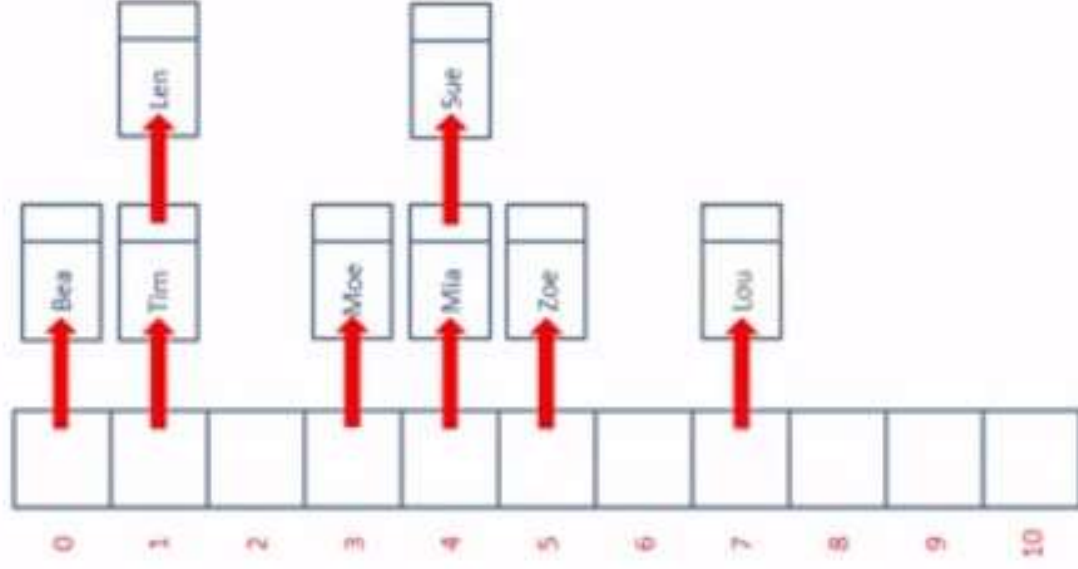


Bea	Tim	Len	Moe	Mia	Zoe	Sue	Lou	Rae	Max	Tod
0	1	2	3	4	5	6	7	8	9	10

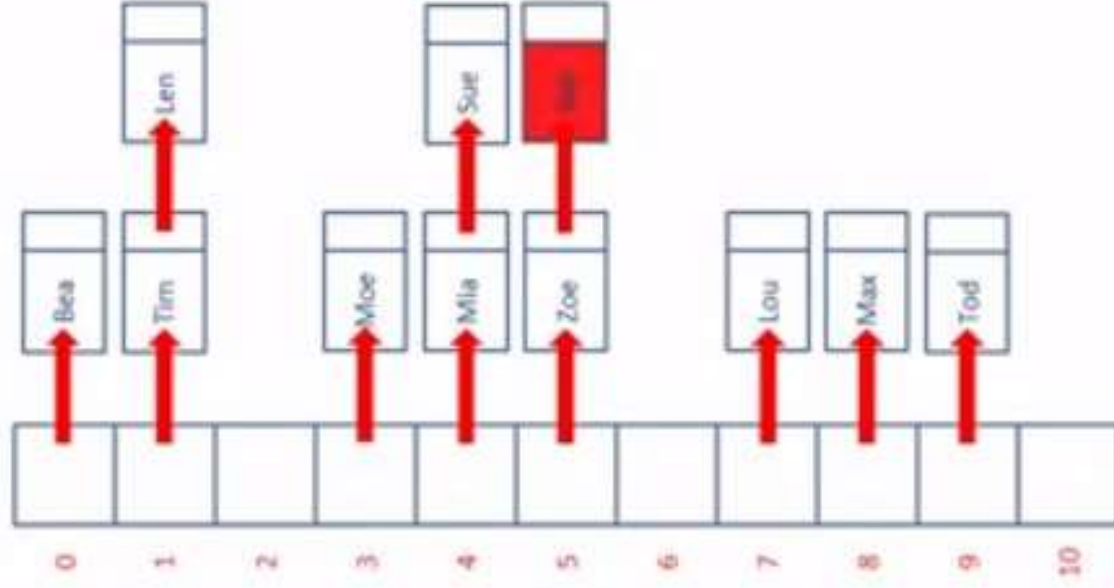
# Closed Addressing (Non-Linear)



Mia M 77 I 105 a 97 279 4



Mia	M	77	l	105	a	97	279	4
Tim	T	84	l	105	m	109	298	1
Bea	B	66	e	101	a	97	264	0
Zoe	Z	90	o	111	e	101	302	5
Sue	S	83	u	117	e	101	301	4
Len	L	76	e	101	n	110	287	1
Moe	M	77	o	111	e	101	289	3
Lou	L	76	o	111	u	117	304	7



Find Rae     $280 \text{ Mod } 11 = 5$

myData = Array(5)



# Objectives of Hash Function

- Minimize collisions
- Uniform distribution of hash values
- Easy to calculate
- Resolve any collisions