



Vidyavardhini's College of Engineering and Technology

Department of Artificial Intelligence & Data Science

Experiment No. 6

Prim's Algorithm

Date of Performance:

Date of Submission:



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Title: Prim's Algorithm.

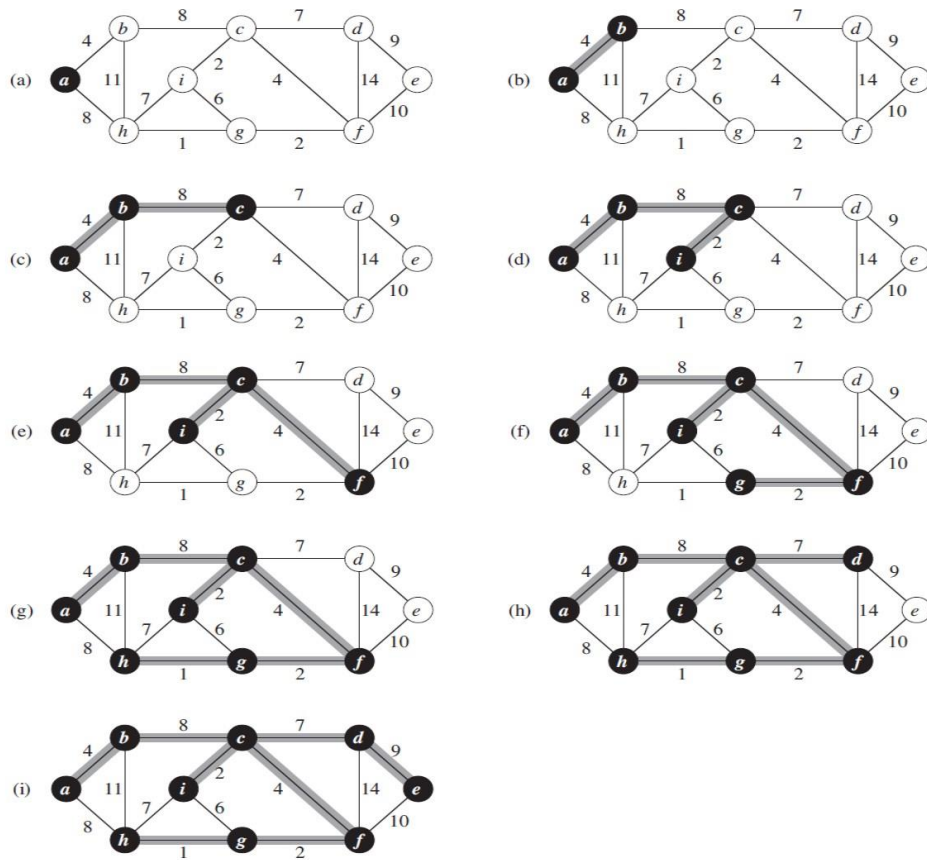
Aim: To study and implement Prim's Minimum Cost Spanning Tree Algorithm.

Objective: To introduce Greedy based algorithms

Theory:

Prim's algorithm is a greedy algorithm that finds a minimum spanning tree for a weighted undirected graph. This means it finds a subset of the edges that forms a tree that includes every vertex, where the total weight of all the edges in the tree is minimized. The algorithm operates by building this tree one vertex at a time, from an arbitrary starting vertex, at each step adding the cheapest possible connection from the tree to another vertex.

Example:



Algorithm and Complexity:



```
1  Algorithm Prim( $E, cost, n, t$ )
2  //  $E$  is the set of edges in  $G$ .  $cost[1 : n, 1 : n]$  is the cost
3  // adjacency matrix of an  $n$  vertex graph such that  $cost[i, j]$  is
4  // either a positive real number or  $\infty$  if no edge  $(i, j)$  exists.
5  // A minimum spanning tree is computed and stored as a set of
6  // edges in the array  $t[1 : n - 1, 1 : 2]$ .  $(t[i, 1], t[i, 2])$  is an edge in
7  // the minimum-cost spanning tree. The final cost is returned.
8  {
9      Let  $(k, l)$  be an edge of minimum cost in  $E$ ;
10      $mincost := cost[k, l]$ ;
11      $t[1, 1] := k$ ;  $t[1, 2] := l$ ;
12     for  $i := 1$  to  $n$  do // Initialize near.
13         if  $(cost[i, l] < cost[i, k])$  then  $near[i] := l$ ;
14         else  $near[i] := k$ ;
15      $near[k] := near[l] := 0$ ;
16     for  $i := 2$  to  $n - 1$  do
17     { // Find  $n - 2$  additional edges for  $t$ .
18         Let  $j$  be an index such that  $near[j] \neq 0$  and
19          $cost[j, near[j]]$  is minimum;
20          $t[i, 1] := j$ ;  $t[i, 2] := near[j]$ ;
21          $mincost := mincost + cost[j, near[j]]$ ;
22          $near[j] := 0$ ;
23         for  $k := 1$  to  $n$  do // Update  $near[ ]$ .
24             if  $((near[k] \neq 0) \text{ and } (cost[k, near[k]] > cost[k, j]))$ 
25                 then  $near[k] := j$ ;
26     }
27     return  $mincost$ ;
28 }
```

Time Complexity is $O(n^2)$, Where, n = number of vertices **Theory:**

Code :

```
#include <stdio.h>
```

```
#include <limits.h>
```

```
#define V 5 // Number of vertices in the graph
```

```
// Function to find the vertex with the minimum key value
```

```
int minKey(int key[], int mstSet[]) {
```

```
    int min = INT_MAX, min_index;
```

```
    for (int v = 0; v < V; v++) {
```

```
        if (mstSet[v] == 0 && key[v] < min) {
```

```
            min = key[v];
```

```
            min_index = v;
```

```

    }
}
return min_index;
}

```

// Function to print the MST

```

void printMST(int parent[], int graph[V][V]) {
    printf("Edge  Weight\n");
    for (int i = 1; i < V; i++) {
        printf("%d - %d  %d \n", parent[i], i, graph[i][parent[i]]);
    }
}

```

// Function to construct and print MST for a graph represented using adjacency matrix

```

void primMST(int graph[V][V]) {
    int parent[V]; // Array to store constructed MST
    int key[V];    // Key values used to pick minimum weight edge in cut
    int mstSet[V]; // To represent set of vertices not yet included in MST

```

// Initialize all keys as INFINITE

```

for (int i = 0; i < V; i++) {
    key[i] = INT_MAX;
    mstSet[i] = 0;
}

```

// Always include first vertex in MST.

```

key[0] = 0; // Make key 0 so that this vertex is picked as first vertex
parent[0] = -1; // First node is always root of MST

```

// The MST will have V vertices

```

for (int count = 0; count < V - 1; count++) {
    // Pick the minimum key vertex from the set of vertices not yet included in MST
    int u = minKey(key, mstSet);

```

// Add the picked vertex to the MST set

```

mstSet[u] = 1;

```

```
// Update key value and parent index of the adjacent vertices of the picked vertex.
```

```
// Consider only those vertices which are not yet included in MST
```

```
for (int v = 0; v < V; v++) {
```

```
    // graph[u][v] is non zero only for adjacent vertices of m
```

```
    // mstSet[v] is false for vertices not yet included in MST
```

```
    // Update the key only if graph[u][v] is smaller than key[v]
```

```
    if (graph[u][v] && mstSet[v] == 0 && graph[u][v] < key[v]) {
```

```
        parent[v] = u;
```

```
        key[v] = graph[u][v];
```

```
    }
```

```
}
```

```
}
```

```
// Print the constructed MST
```

```
printMST(parent, graph);
```

```
}
```

```
int main() {
```

```
    /* Let us create the following graph
```

```
      2   3
```

```
    (0)--(1)--(2)
```

```
      |  /\  |
```

```
    6| 8/   \5|7
```

```
      | /    \ |
```

```
    (3)----- (4)
```

```
      9       */
```

```
int graph[V][V] = {{0, 2, 0, 6, 0},
```

```
                  {2, 0, 3, 8, 5},
```

```
                  {0, 3, 0, 0, 7},
```

```
                  {6, 8, 0, 0, 9},
```

```
                  {0, 5, 7, 9, 0}};
```

```
// Print the solution
```

```
primMST(graph);
```

```
return 0;
```

```
}
```

Output :

```
PS E:\Testing_Lang> cd "e:\Testing_Lang\" ; if ($?) { gcc test.c -o test } ; if ($?) { .\test }
Edge  Weight
0 - 1   2
1 - 2   3
0 - 3   6
1 - 4   5
PS E:\Testing_Lang> 
```

Conclusion:

This program implements Prim's MST algorithm using the Greedy method for a given graph represented using an adjacency matrix. It constructs and prints the Minimum Spanning Tree (MST) of the graph.

Prim's algorithm works by starting from an arbitrary vertex and repeatedly adding the minimum-weight edge that connects a vertex in the MST to a vertex outside the MST. It continues this process until all vertices are included in the MST.

