Vidyavardhini's College of Engineering and Technology Department of Artificial Intelligence & Data Science

_				N I	
$\mathbf{F}\mathbf{V}$	peri	\mathbf{m}	nt	INI (۱/ r
ᆫᄼ			-I I C	1 11 (J.T

Implementation of Bidirectional search for problem solving.

Date of Performance:

Date of Submission:

CSL502: Artificial Intelligence

Lab



Vidyavardhini's College of Engineering and Technology Department of Artificial Intelligence & Data Science

Aim: Implementation of Bidirectional search for problem solving.

Objective: To study the Bidirectional searching techniques and its implementation for problem solving.

Theory:

Bidirectional search is a graph search algorithm which find smallest path from source to goal vertex. It runs two simultaneous search –

- 1. Forward search from source/initial vertex toward goal vertex
- 2. Backward search from goal/target vertex toward source vertex

Bidirectional search replaces single search graph(which is likely to grow exponentially) with two smaller sub

graphs – one starting from initial vertex and other starting from goal vertex. The search terminates when two

graphs intersect.

Algorithm:

Steps for Bidirectional Search Algorithm

1. Initialization:

- o Create two frontiers:
 - One for the forward search starting from the initial node (start).
 - One for the backward search starting from the goal node (goal).
- o Create two sets to keep track of visited nodes for each search direction:
 - visited_start for the forward search.
 - visited_goal for the backward search.
- o Initialize the frontiers by adding the start node to frontier_start and the goal node to frontier_goal.
- o Initialize the visited_start and visited_goal sets with their respective starting nodes.

2. Search Expansion:

o Repeat the following steps until a meeting point is found or one of the frontiers is empty:

• Expand Forward Search:

- Remove the current node from frontier_start.
- Expand all neighboring nodes of the current node.



Vidyavardhini's College of Engineering and Technology

Department of Artificial Intelligence & Data Science

- For each neighbor:
 - If the neighbor is not in visited_start:
 - Add the neighbor to frontier_start.
 - Mark the neighbor as visited in visited_start.
 - Check if the neighbor is already in visited_goal:
 - If yes, a meeting point is found. Proceed to reconstruct the

path.

- Expand Backward Search:
 - Remove the current node from frontier_goal.
 - Expand all neighboring nodes of the current node.
 - For each neighbor:
 - If the neighbor is not in visited_goal:
 - Add the neighbor to frontier_goal.
 - Mark the neighbor as visited in visited_goal.
 - Check if the neighbor is already in visited_start:
 - If yes, a meeting point is found. Proceed to reconstruct the

path.

3. **Meeting Point**:

o The search stops when a node from frontier_start is found in visited_goal or a node from

frontier_goal is found in visited_start. This node is the meeting point.

4. Path Reconstruction:

o Reconstruct the path from the start node to the goal node by combining the paths from both



Vidyavardhini's College of Engineering and Technology Department of Artificial Intelligence & Data Science

searches at the meeting point.

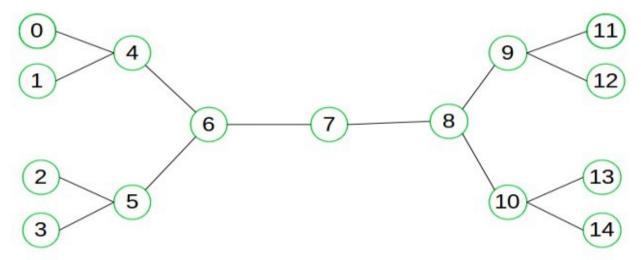
- o Start from the meeting point:
 - Trace back to the start node using the information in visited_start.
 - Trace back to the goal node using the information in visited_goal.
- o Concatenate the two paths to form the complete path from start to goal.

5. **Termination**:

o If one of the frontiers is empty and no meeting point is found, it means there is no path from the

start node to the goal node.

Example:



Suppose we want to find if there exists a path from vertex 0 to vertex 14. Here we can execute two searches, one



Vidyavardhini's College of Engineering and Technology

Department of Artificial Intelligence & Data Science

from vertex 0 and other from vertex 14. When both forward and backward search meet at vertex 7, we know that

we have found a path from node 0 to 14 and search can be terminated now. We can clearly see that we have

successfully avoided unnecessary exploration.

When to use bidirectional approach?

We can consider bidirectional approach when-

- 1. Both initial and goal states are unique and completely defined.
- 2. The branching factor is exactly the same in both directions.

Performance measures

Completeness: Bidirectional Search is complete if we use BFS in both searches.

Time Complexity: Time complexity of bidirectional search using BFS is **O(b**^d**)**.

Space Complexity: Space complexity of bidirectional search is **O**(**b**^d).

Optimal: Bidirectional search is Optimal.

Advantages:

- o Bidirectional search is fast.
- o Bidirectional search requires less memory

Disadvantages:

- o Implementation of the bidirectional search tree is difcult.
- o In bidirectional search, one should know the goal state in advance.