



Vidyavardhini's College of Engineering and Technology

Department of Artificial Intelligence & Data Science

Experiment No. 9
Travelling Salesperson Problem using Dynamic Approach
Date of Performance:
Date of Submission:

Experiment No. 9

Title: Travelling Salesman Problem

Aim: To study and implement Travelling Salesman Problem.



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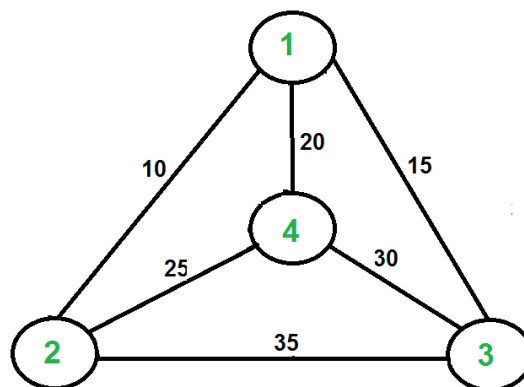
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Objective: To introduce Dynamic Programming approach

Theory:

The **Traveling Salesman Problem (TSP)** is a classic optimization problem in which a salesperson needs to visit a set of cities exactly once and return to the starting city while minimizing the total distance traveled.

Given a set of cities and the distance between every pair of cities, find the **shortest possible route** that visits every city exactly once and returns to the starting point.



For example, consider the graph shown in the figure on the right side. A TSP tour in the graph is 1-2-4-3-1. The cost of the tour is $10+25+30+15$ which is 80. The problem is a famous NP-hard problem. There is no polynomial-time known solution for this problem. The following are different solutions for the traveling salesman problem.

Naive Solution:

- 1) Consider city 1 as the starting and ending point.
- 2) Generate all $(n-1)!$ Permutations of cities.
- 3) Calculate the cost of every permutation and keep track of the minimum cost permutation.
- 4) Return the permutation with minimum cost.

Time Complexity: $O(n!)$

Dynamic Programming:

Let the given set of vertices be $\{1, 2, 3, 4, \dots, n\}$. Let us consider 1 as starting and ending point of output. For every other vertex i (other than 1), we find the minimum cost path with 1 as the starting point, i as the ending point, and all vertices appearing exactly once. Let the cost of this path be $cost(i)$, and the cost of the corresponding Cycle would be $cost(i) + dist(i, 1)$ where $dist(i, 1)$ is the distance from i to 1. Finally, we return the minimum of all $[cost(i) + dist(i, 1)]$ values. This looks simple so far.

Now the question is how to get $cost(i)$? To calculate the $cost(i)$ using Dynamic Programming, we need to have some recursive relation in terms of sub-problems.

Let us define a term $C(S, i)$ be the cost of the minimum cost path visiting each vertex in set S

exactly once, starting at 1 and ending at i. We start with all subsets of size 2 and calculate



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$C(S, i)$ for all subsets where S is the subset, then we calculate $C(S, i)$ for all subsets S of size 3 and so on. Note that 1 must be present in every subset.

If size of S is 2, then S must be $\{1, i\}$,

$$C(S, i) = \text{dist}(1, i)$$

Else if size of S is greater than 2.

$$C(S, i) = \min \{ C(S - \{i\}, j) + \text{dis}(j, i) \} \text{ where } j \text{ belongs to } S, j \neq i \text{ and } j \neq 1.$$

Code :

```
#include <stdio.h>
```

```
#include <stdlib.h>
```

```
#include <limits.h>
```

```
#define V 4 // Number of vertices in the graph
```

```
// Function to find the minimum of two numbers
```

```
int min(int a, int b) {
```

```
    return (a < b) ? a : b;
```

```
}
```

```
// Function to solve TSP using dynamic programming
```

```
int tsp(int graph[][V], int mask, int pos, int n, int dp[][V]) {
```

```
    // If all vertices have been visited
```

```
    if (mask == (1 << n) - 1) {
```

```
        return graph[pos][0]; // Return cost of going back to the starting city
```

```
    }
```

```
    // If this subproblem has already been computed
```

```
    if (dp[mask][pos] != -1) {
```

```
        return dp[mask][pos];
```

```

    }

    int ans = INT_MAX;

    // Try to go to an unvisited city
    for (int city = 0; city < n; city++) {
        if ((mask & (1 << city)) == 0) { // If city has not been visited
            int newAns = graph[pos][city] + tsp(graph, mask | (1 << city), city, n, dp);
            ans = min(ans, newAns);
        }
    }

    return dp[mask][pos] = ans;
}

int main() {
    int graph[V][V] = {{0, 10, 15, 20},
                       {10, 0, 35, 25},
                       {15, 35, 0, 30},
                       {20, 25, 30, 0}};

    int dp[1 << V][V]; // Dynamic programming table to store results of
    subproblems

    // Initialize dp table with -1
    for (int i = 0; i < (1 << V); i++) {
        for (int j = 0; j < V; j++) {
            dp[i][j] = -1;
        }
    }

    int minCost = tsp(graph, 1, 0, V, dp); // Starting from city 0

    printf("Minimum cost of TSP: %d\n", minCost);
}

```

```
    return 0;  
}
```

Output :

```
PS E:\Testing_Lang> cd "e:\Testing_Lang\" ; if ($?) { gcc test.c -o test } ; if ($?) { .\test }  
Minimum cost of TSP: 80  
PS E:\Testing_Lang> 
```

Conclusion:

Travelling Salesman Problem has been successfully implemented. This program demonstrates solving the TSP using dynamic programming. It initializes a 2D array `dp` to store the results of subproblems. The `tsp` function recursively explores all possible permutations of cities and computes the minimum cost of the tour. It uses bitmasking to keep track of visited cities efficiently.