



Vidyavardhini's College of Engineering and Technology

Department of Artificial Intelligence & Data Science

Experiment No. 7
Kruskal's Algorithm
Date of Performance:
Date of Submission:

Experiment No. 7

Title: Kruskal's Algorithm.

Aim: To study and implement Kruskal's Minimum Cost Spanning Tree Algorithm.

Objective: To introduce Greedy based algorithms.

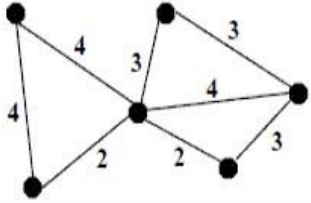

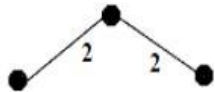
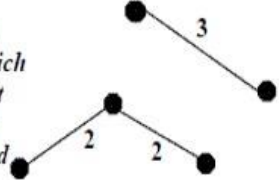
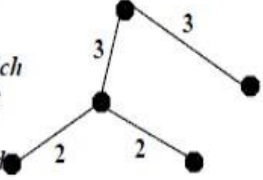
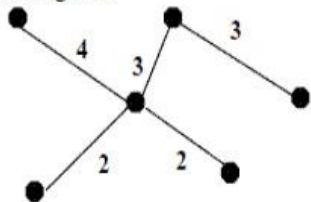
Theory:

Kruskal's algorithm finds a minimum spanning forest of an undirected edge-weighted graph. If the graph is connected, it finds a minimum spanning tree. (A minimum spanning tree of a connected graph is a subset of the edges that forms a tree that includes every vertex, where the sum of the weights of all the edges in the tree is minimized. For a disconnected graph, a minimum spanning forest is composed of a minimum spanning tree for each connected component.) It is a greedy algorithm in graph theory as in each step it adds the next lowest-weight edge that will not form a cycle to the minimum spanning forest.

Example:



Kruskal's Algorithm

<p>1 Given a network.....</p> 	<p>2 Choose the shortest edge (if there is more than one, choose any of the shortest).....</p> 	<p>3 Choose the next shortest edge and add it.....</p> 
<p>4 Choose the next shortest edge which wouldn't create a cycle and add it.</p> 	<p>5 Choose the next shortest edge which wouldn't create a cycle and add it.</p> 	<p>6 Repeat until you have a minimal spanning tree.</p> 

Algorithm and Complexity:



Vidyavardhini's College of Engineering and Technology

Department of Artificial Intelligence & Data Science

```
1  Algorithm Kruskal( $E, cost, n, t$ )
2  //  $E$  is the set of edges in  $G$ .  $G$  has  $n$  vertices.  $cost[u, v]$  is the
3  // cost of edge  $(u, v)$ .  $t$  is the set of edges in the minimum-cost
4  // spanning tree. The final cost is returned.
5  {
6      Construct a heap out of the edge costs using Heapify;
7      for  $i := 1$  to  $n$  do  $parent[i] := -1$ ;
8      // Each vertex is in a different set.
9       $i := 0$ ;  $mincost := 0.0$ ;
10     while  $((i < n - 1)$  and  $(heap \text{ not empty}))$  do
11     {
12         Delete a minimum cost edge  $(u, v)$  from the heap
13         and reheapify using Adjust;
14          $j := Find(u)$ ;  $k := Find(v)$ ;
15         if  $(j \neq k)$  then
16         {
17              $i := i + 1$ ;
18              $t[i, 1] := u$ ;  $t[i, 2] := v$ ;
19              $mincost := mincost + cost[u, v]$ ;
20             Union $(j, k)$ ;
21         }
22     }
23     if  $(i \neq n - 1)$  then write ("No spanning tree");
24     else return  $mincost$ ;
25 }
```

Time Complexity is $O(n \log n)$, Where, n = number of Edges

Code :

```
#include <stdio.h>
#include <stdlib.h>

// Structure to represent an edge in the graph
struct Edge {
    int src, dest, weight;
};

// Structure to represent a subset for union-find
struct Subset {
    int parent;
    int rank;
};
```

```

// Function to find set of an element i
int find(struct Subset subsets[], int i) {
    if (subsets[i].parent != i)
        subsets[i].parent = find(subsets, subsets[i].parent);
    return subsets[i].parent;
}

// Function that does union of two sets of x and y
void Union(struct Subset subsets[], int x, int y) {
    int xroot = find(subsets, x);
    int yroot = find(subsets, y);

    // Attach smaller rank tree under root of high rank tree (Union by Rank)
    if (subsets[xroot].rank < subsets[yroot].rank)
        subsets[xroot].parent = yroot;
    else if (subsets[xroot].rank > subsets[yroot].rank)
        subsets[yroot].parent = xroot;
    else {
        subsets[yroot].parent = xroot;
        subsets[xroot].rank++;
    }
}

// Comparator function to sort edges based on weight
int compare(const void *a, const void *b) {
    struct Edge *a1 = (struct Edge *)a;
    struct Edge *b1 = (struct Edge *)b;
    return a1->weight > b1->weight;
}

// Function to construct and print MST using Kruskal's algorithm
void KruskalMST(struct Edge edges[], int V, int E) {
    struct Edge result[V]; // Store the result MST
    int e = 0; // An index variable used for result[]
    int i = 0; // An index variable used for sorted edges

    // Step 1: Sort all the edges in non-decreasing order of their weight
    qsort(edges, E, sizeof(edges[0]), compare);

```

```

// Allocate memory for creating V subsets
struct Subset *subsets = (struct Subset *)malloc(V * sizeof(struct Subset));

// Create V subsets with single elements
for (int v = 0; v < V; v++) {
    subsets[v].parent = v;
    subsets[v].rank = 0;
}

// Number of edges to be taken is equal to V-1
while (e < V - 1 && i < E) {
    // Step 2: Pick the smallest edge. And increment the index for next iteration
    struct Edge next_edge = edges[i++];

    int x = find(subsets, next_edge.src);
    int y = find(subsets, next_edge.dest);

    // If including this edge doesn't cause cycle, include it in result and increment
    // the index
    // of result for next edge
    if (x != y) {
        result[e++] = next_edge;
        Union(subsets, x, y);
    }
}

// Print the constructed MST
printf("Following are the edges in the constructed MST:\n");
for (i = 0; i < e; ++i)
    printf("%d -- %d == %d\n", result[i].src, result[i].dest, result[i].weight);
}

int main() {
    /* Let us create the following graph
        10
        0-----1
        | \   |
        6|  5\ |15
        |   \|
    */

```

```

    2-----3
    4      */
int V = 4; // Number of vertices in graph
int E = 5; // Number of edges in graph
struct Edge edges[] = {{0, 1, 10}, {0, 2, 6}, {0, 3, 5}, {1, 3, 15}, {2, 3, 4}};

// Function call to construct and print MST
KruskalMST(edges, V, E);

return 0;
}

```

Output :

```

PS E:\Testing_Lang> cd "e:\Testing_Lang\" ; if ($?) { gcc test.c -o test } ; if ($?) { .\test }
Following are the edges in the constructed MST:
2 -- 3 == 4
0 -- 3 == 5
0 -- 1 == 10
PS E:\Testing_Lang> 

```

Conclusion:

This program implements Kruskal's MST algorithm using the Greedy method for a given graph represented using an array of edges. It constructs and prints the Minimum Spanning Tree (MST) of the graph.

Kruskal's algorithm works by sorting all the edges in non-decreasing order of their weight and selecting edges one by one in this sorted order, adding the edge to the MST if it doesn't create a cycle. It continues this process until all vertices are included in the MST.

